Support Recovery for Orthogonal Matching Pursuit

Upper and Lower bounds

Raghav Somani, Chirag Gupta, Prateek Jain and Praneeth Netrapalli

September 24, 2018

Sparse Regression

$$\bar{\mathbf{x}} = \underset{\|\mathbf{x}\|_0 \le s^*}{\min} f(\mathbf{x}) \tag{1.1}$$

 $\mathbf{x} \in \mathbb{R}^d$ and $s^* << d$.

 ℓ_0 norm counts the number of non-zero elements.

Applications

- Resource constrained Machine Learning
- High dimensional Statistics
- Bioinformatics



Sparse Regression

$$\bar{\mathbf{x}} = \underset{\|\mathbf{x}\|_{0} \le s^{*}}{\min} f(\mathbf{x}) \tag{1.1}$$

 $\mathbf{x} \in \mathbb{R}^d$ and $s^* << d$.

 ℓ_0 norm counts the number of non-zero elements.

Applications

- Resource constrained Machine Learning
- High dimensional Statistics
- Bioinformatics



Sparse Linear Regression (SLR)

- Sparse Linear Regression is a representative problem. Results typically extend easily to general case.
- With $f(\mathbf{x}) = \|\mathbf{A}\mathbf{x} \mathbf{y}\|_2^2$, SLR's objective is to find

$$\overline{\mathbf{x}} = \underset{\|\mathbf{x}\|_0 \le s^*}{\arg\min} \|\mathbf{A}\mathbf{x} - \mathbf{y}\|_2^2$$
 (2.1)

where $\mathbf{A} \in \mathbb{R}^{n \times d}, \mathbf{x} \in \mathbb{R}^d$ and $\mathbf{y} = \mathbb{R}^n$.

Unconditionally, it is NP hard (reduction to 3 set cover problem).



Sparse Linear Regression (SLR)

- Sparse *Linear* Regression is a representative problem. Results typically extend easily to general case.
- With $f(\mathbf{x}) = \|\mathbf{A}\mathbf{x} \mathbf{y}\|_2^2$, SLR's objective is to find

$$\bar{\mathbf{x}} = \underset{\|\mathbf{x}\|_0 \le s^*}{\arg\min} \|\mathbf{A}\mathbf{x} - \mathbf{y}\|_2^2$$
 (2.1)

where $\mathbf{A} \in \mathbb{R}^{n \times d}, \mathbf{x} \in \mathbb{R}^d$ and $\mathbf{y} = \mathbb{R}^n$.

Unconditionally, it is NP hard (reduction to 3 set cover problem).



Sparse Linear Regression (SLR)

- Sparse Linear Regression is a representative problem. Results typically extend easily to general case.
- With $f(\mathbf{x}) = \|\mathbf{A}\mathbf{x} \mathbf{y}\|_2^2$, SLR's objective is to find

$$\bar{\mathbf{x}} = \underset{\|\mathbf{x}\|_{0} \leq s^{*}}{\operatorname{arg \, min}} \|\mathbf{A}\mathbf{x} - \mathbf{y}\|_{2}^{2} \tag{2.1}$$

where $\mathbf{A} \in \mathbb{R}^{n \times d}, \mathbf{x} \in \mathbb{R}^d$ and $\mathbf{y} = \mathbb{R}^n$.

Unconditionally, it is NP hard (reduction to 3 set cover problem).

Despite being NP hard, SLR is tractable under certain assumptions.

Despite being NP hard, SLR is tractable under certain assumptions.

Incoherence -

If
$$\Sigma = \mathbf{A}^T \mathbf{A}$$
, then $\max_{i \neq j} |\Sigma_{ij}| \leq M$

If $M \leq \frac{1}{2s^*-1}$ and $\mathbf{y} = \mathbf{A}\bar{\mathbf{x}} \implies \bar{\mathbf{x}}$ is unique sparsest solution, and OMP can recover $\bar{\mathbf{x}}$ in s^* steps.

Restricted Isometry Property (RIP) -

$$\begin{aligned} & \left\| \mathbf{A}_{\mathbf{S}}^{T} \mathbf{A}_{\mathbf{S}} - \mathbf{I} \right\|_{2} \le \delta_{|\mathbf{S}|} & (\delta_{s} \le M(s-1) & \forall \ s \ge 2) \\ \Rightarrow & (1 - \delta_{s}) \left\| \mathbf{v} \right\|_{2}^{2} \le \left\| \mathbf{A} \mathbf{v} \right\|_{2}^{2} \le (1 + \delta_{s}) \left\| \mathbf{v} \right\|_{2}^{2} & \forall \ \mathbf{v} \ \text{s.t.} \ \left\| \mathbf{v} \right\|_{0} \le s \end{aligned}$$

Null space property -

$$orall \mathbf{S} \in [d] ext{ s.t. } |\mathbf{S}| \leq s ext{, if } \mathbf{v} \in \mathsf{Null}(\mathbf{A}) \setminus \{\mathbf{0}\} ext{, then } \|\mathbf{v_S}\|_1 \leq \|\mathbf{v_{S^c}}\|_1 \\ \Longrightarrow \left\{ \mathbf{v} \in \mathbb{R}^d \mid \mathbf{A}\mathbf{v} = \mathbf{0} \right\} \cap \left\{ \mathbf{v} \in \mathbb{R}^d \mid \|\mathbf{v_{S^c}}\|_1 \leq \|\mathbf{v_S}\|_1 \right\} = \mathbf{0}$$

Restricted Strong Convexity (RSC) -

$$\|\mathbf{A}\mathbf{x} - \mathbf{A}\mathbf{z}\|_{2}^{2} \ge \rho_{s}^{-} \|\mathbf{x} - \mathbf{z}\|_{2}^{2} \ \forall \ \mathbf{x}, \mathbf{z} \in \mathbb{R}^{d} \ \text{s.t.} \ \|\mathbf{x} - \mathbf{z}\|_{0} \le s$$



Despite being NP hard, SLR is tractable under certain assumptions.

Incoherence -

If
$$\Sigma = \mathbf{A}^T \mathbf{A}$$
, then $\max_{i \neq j} |\Sigma_{ij}| \leq M$

If $M \leq \frac{1}{2s^*-1}$ and $\mathbf{y} = \mathbf{A}\bar{\mathbf{x}} \implies \bar{\mathbf{x}}$ is unique sparsest solution, and OMP can recover $\bar{\mathbf{x}}$ in s^* steps.

Restricted Isometry Property (RIP) -

$$\begin{aligned} & \left\| \mathbf{A}_{\mathbf{S}}^{T} \mathbf{A}_{\mathbf{S}} - \mathbf{I} \right\|_{2} \leq \delta_{|\mathbf{S}|} & (\delta_{s} \leq M(s-1) \quad \forall \ s \geq 2) \\ \Longrightarrow & (1 - \delta_{s}) \left\| \mathbf{v} \right\|_{2}^{2} \leq \left\| \mathbf{A} \mathbf{v} \right\|_{2}^{2} \leq (1 + \delta_{s}) \left\| \mathbf{v} \right\|_{2}^{2} \quad \forall \ \mathbf{v} \ \text{s.t.} \ \left\| \mathbf{v} \right\|_{0} \leq s. \end{aligned}$$

Null space property -

$$orall \mathbf{S} \in [d] ext{ s.t. } |\mathbf{S}| \leq s ext{, if } \mathbf{v} \in \mathsf{Null}(\mathbf{A}) \setminus \{\mathbf{0}\} ext{, then } \|\mathbf{v_S}\|_1 \leq \|\mathbf{v_{S^c}}\|_1 \ \Rightarrow \ \left\{\mathbf{v} \in \mathbb{R}^d \mid \mathbf{A}\mathbf{v} = \mathbf{0}\right\} \cap \left\{\mathbf{v} \in \mathbb{R}^d \mid \|\mathbf{v_{S^c}}\|_1 \leq \|\mathbf{v_S}\|_1 \right\} = \mathbf{0}$$

Restricted Strong Convexity (RSC) -

$$\|\mathbf{A}\mathbf{x} - \mathbf{A}\mathbf{z}\|_{2}^{2} \ge \rho_{s}^{-} \|\mathbf{x} - \mathbf{z}\|_{2}^{2} \ \forall \ \mathbf{x}, \mathbf{z} \in \mathbb{R}^{d} \ \text{s.t.} \ \|\mathbf{x} - \mathbf{z}\|_{0} \le s$$



Despite being NP hard, SLR is tractable under certain assumptions.

Incoherence -

If
$$\Sigma = \mathbf{A}^T \mathbf{A}$$
, then $\max_{i \neq j} |\Sigma_{ij}| \leq M$

If $M \leq \frac{1}{2s^*-1}$ and $\mathbf{y} = \mathbf{A}\bar{\mathbf{x}} \implies \bar{\mathbf{x}}$ is unique sparsest solution, and OMP can recover $\bar{\mathbf{x}}$ in s^* steps.

Restricted Isometry Property (RIP) -

$$\begin{aligned} & \left\| \mathbf{A}_{\mathbf{S}}^{T} \mathbf{A}_{\mathbf{S}} - \mathbf{I} \right\|_{2} \leq \delta_{|\mathbf{S}|} & (\delta_{s} \leq M(s-1) \quad \forall \ s \geq 2) \\ \Longrightarrow & (1 - \delta_{s}) \left\| \mathbf{v} \right\|_{2}^{2} \leq \left\| \mathbf{A} \mathbf{v} \right\|_{2}^{2} \leq (1 + \delta_{s}) \left\| \mathbf{v} \right\|_{2}^{2} \quad \forall \ \mathbf{v} \ \text{s.t.} \ \left\| \mathbf{v} \right\|_{0} \leq s. \end{aligned}$$

Null space property -

$$\forall \ \mathbf{S} \in [d] \ \text{s.t.} \ |\mathbf{S}| \leq s, \text{ if } \mathbf{v} \in \mathsf{Null}(\mathbf{A}) \setminus \{\mathbf{0}\}, \text{ then } \|\mathbf{v}_{\mathbf{S}}\|_1 \leq \|\mathbf{v}_{\mathbf{S}^c}\|_1 \\ \Longrightarrow \left\{ \mathbf{v} \in \mathbb{R}^d \mid \mathbf{A}\mathbf{v} = \mathbf{0} \right\} \cap \left\{ \mathbf{v} \in \mathbb{R}^d \mid \|\mathbf{v}_{\mathbf{S}^c}\|_1 \leq \|\mathbf{v}_{\mathbf{S}}\|_1 \right\} = \mathbf{0}$$

Restricted Strong Convexity (RSC) -

$$\|\mathbf{A}\mathbf{x} - \mathbf{A}\mathbf{z}\|_2^2 \ge \rho_s^- \|\mathbf{x} - \mathbf{z}\|_2^2 \ \forall \ \mathbf{x}, \mathbf{z} \in \mathbb{R}^d \ \text{s.t.} \ \|\mathbf{x} - \mathbf{z}\|_0 \le s$$

Despite being NP hard, SLR is tractable under certain assumptions.

Incoherence -

If
$$\Sigma = \mathbf{A}^T \mathbf{A}$$
, then $\max_{i \neq j} |\Sigma_{ij}| \leq M$

If $M \leq \frac{1}{2s^*-1}$ and $\mathbf{y} = \mathbf{A}\bar{\mathbf{x}} \implies \bar{\mathbf{x}}$ is unique sparsest solution, and OMP can recover $\bar{\mathbf{x}}$ in s^* steps.

Restricted Isometry Property (RIP) -

$$\begin{aligned} & \left\| \mathbf{A}_{\mathbf{S}}^{T} \mathbf{A}_{\mathbf{S}} - \mathbf{I} \right\|_{2} \leq \delta_{|\mathbf{S}|} & (\delta_{s} \leq M(s-1) \quad \forall \ s \geq 2) \\ \Longrightarrow & (1 - \delta_{s}) \left\| \mathbf{v} \right\|_{2}^{2} \leq \left\| \mathbf{A} \mathbf{v} \right\|_{2}^{2} \leq (1 + \delta_{s}) \left\| \mathbf{v} \right\|_{2}^{2} \quad \forall \ \mathbf{v} \ \text{s.t.} \ \left\| \mathbf{v} \right\|_{0} \leq s. \end{aligned}$$

Null space property -

$$\forall \ \mathbf{S} \in [d] \ \text{s.t.} \ |\mathbf{S}| \leq s, \ \text{if} \ \mathbf{v} \in \mathsf{Null}(\mathbf{A}) \setminus \{\mathbf{0}\}, \ \text{then} \ \|\mathbf{v_S}\|_1 \leq \|\mathbf{v_{S^c}}\|_1 \\ \Longrightarrow \left\{ \mathbf{v} \in \mathbb{R}^d \mid \mathbf{A}\mathbf{v} = \mathbf{0} \right\} \cap \left\{ \mathbf{v} \in \mathbb{R}^d \mid \|\mathbf{v_{S^c}}\|_1 \leq \|\mathbf{v_S}\|_1 \right\} = \mathbf{0}$$

Restricted Strong Convexity (RSC) -

$$\|\mathbf{A}\mathbf{x} - \mathbf{A}\mathbf{z}\|_{2}^{2} \ge \rho_{s}^{-} \|\mathbf{x} - \mathbf{z}\|_{2}^{2} \ \forall \ \mathbf{x}, \mathbf{z} \in \mathbb{R}^{d} \ \text{s.t.} \ \|\mathbf{x} - \mathbf{z}\|_{0} \le s$$



Despite being NP hard, SLR is tractable under certain assumptions.

Incoherence -

If
$$\Sigma = \mathbf{A}^T \mathbf{A}$$
, then $\max_{i \neq j} |\Sigma_{ij}| \leq M$

If $M \leq \frac{1}{2s^*-1}$ and $\mathbf{y} = \mathbf{A}\bar{\mathbf{x}} \implies \bar{\mathbf{x}}$ is unique sparsest solution, and OMP can recover $\bar{\mathbf{x}}$ in s^* steps.

Restricted Isometry Property (RIP) -

$$\begin{aligned} & \left\| \mathbf{A}_{\mathbf{S}}^{T} \mathbf{A}_{\mathbf{S}} - \mathbf{I} \right\|_{2} \leq \delta_{|\mathbf{S}|} & (\delta_{s} \leq M(s-1) \quad \forall \ s \geq 2) \\ \Longrightarrow & (1 - \delta_{s}) \left\| \mathbf{v} \right\|_{2}^{2} \leq \left\| \mathbf{A} \mathbf{v} \right\|_{2}^{2} \leq (1 + \delta_{s}) \left\| \mathbf{v} \right\|_{2}^{2} \quad \forall \ \mathbf{v} \ \text{s.t.} \ \left\| \mathbf{v} \right\|_{0} \leq s. \end{aligned}$$

Null space property -

$$\forall \mathbf{S} \in [d] \text{ s.t. } |\mathbf{S}| \leq s, \text{ if } \mathbf{v} \in \text{Null}(\mathbf{A}) \setminus \{\mathbf{0}\}, \text{ then } \|\mathbf{v}_{\mathbf{S}}\|_{1} \leq \|\mathbf{v}_{\mathbf{S}^{c}}\|_{1} \\ \implies \left\{\mathbf{v} \in \mathbb{R}^{d} \mid \mathbf{A}\mathbf{v} = \mathbf{0}\right\} \cap \left\{\mathbf{v} \in \mathbb{R}^{d} \mid \|\mathbf{v}_{\mathbf{S}^{c}}\|_{1} \leq \|\mathbf{v}_{\mathbf{S}}\|_{1}\right\} = \mathbf{0}$$

Restricted Strong Convexity (RSC) -

$$\|\mathbf{A}\mathbf{x} - \mathbf{A}\mathbf{z}\|_{2}^{2} \ge \rho_{s}^{-} \|\mathbf{x} - \mathbf{z}\|_{2}^{2} \ \forall \ \mathbf{x}, \mathbf{z} \in \mathbb{R}^{d} \ \text{s.t.} \ \|\mathbf{x} - \mathbf{z}\|_{0} \le s$$

Incoherence \implies RIP \implies Null space property \implies RSC RSC is the weakest and the most popular assumption.



SLR can be modelled as

$$\mathbf{y} = \mathbf{A}\bar{\mathbf{x}} + \boldsymbol{\eta} \tag{2.2}$$

where $\eta \sim \mathcal{N}(\mathbf{0}, \sigma^2 \mathbf{I}_{n \times n})$, $\operatorname{supp}(\bar{\mathbf{x}}) = \mathbf{S}^*$ and $|\mathbf{S}^*| = s^*$.

$$\Rightarrow \mathbf{y} = \mathbf{A}_{\mathbf{S}^*} \bar{\mathbf{x}}_{\mathbf{S}^*} + \boldsymbol{\eta} \tag{2.3}$$

Model with deterministic conditions on η can also be analyzed.

- **Bounding Generalization error** Upper bound $G(\mathbf{x}) := \frac{1}{n} \|\mathbf{A}(\mathbf{x} \bar{\mathbf{x}})\|_2^2$ where the rows of \mathbf{A} are i.i.d.
- **Support Recovery** Recover the true features of A, i.e., find a $S \supseteq S^*$.

SLR can be modelled as

$$y = A\bar{x} + \eta \tag{2.2}$$

where $\eta \sim \mathcal{N}(\mathbf{0}, \sigma^2 \mathbf{I}_{n \times n})$, $\operatorname{supp}(\bar{\mathbf{x}}) = \mathbf{S}^*$ and $|\mathbf{S}^*| = s^*$.

$$\implies \mathbf{y} = \mathbf{A}_{\mathbf{S}^*} \bar{\mathbf{x}}_{\mathbf{S}^*} + \boldsymbol{\eta} \tag{2.3}$$

Model with deterministic conditions on η can also be analyzed.

- **Bounding Generalization error** Upper bound $G(\mathbf{x}) := \frac{1}{n} \|\mathbf{A}(\mathbf{x} \bar{\mathbf{x}})\|_2$ where the rows of \mathbf{A} are i.i.d.
- **Outport Recovery** Recover the true features of A. i.e., find a $S \supset S^*$.

SLR can be modelled as

$$y = A\bar{x} + \eta \tag{2.2}$$

where $\eta \sim \mathcal{N}(\mathbf{0}, \sigma^2 \mathbf{I}_{n \times n})$, $\operatorname{supp}(\bar{\mathbf{x}}) = \mathbf{S}^*$ and $|\mathbf{S}^*| = s^*$.

$$\implies \mathbf{y} = \mathbf{A}_{\mathbf{S}^*} \bar{\mathbf{x}}_{\mathbf{S}^*} + \boldsymbol{\eta} \tag{2.3}$$

Model with deterministic conditions on η can also be analyzed.

- **Output** Bounding Generalization error Upper bound $G(\mathbf{x}) := \frac{1}{n} \|\mathbf{A}(\mathbf{x} \bar{\mathbf{x}})\|_2^2$ where the rows of \mathbf{A} are i.i.d.
- ② Support Recovery Recover the true features of A, i.e., find a $S \supseteq S^*$.

SLR can be modelled as

$$y = A\bar{x} + \eta \tag{2.2}$$

where $\eta \sim \mathcal{N}(\mathbf{0}, \sigma^2 \mathbf{I}_{n \times n})$, $\operatorname{supp}(\bar{\mathbf{x}}) = \mathbf{S}^*$ and $|\mathbf{S}^*| = s^*$.

$$\implies \mathbf{y} = \mathbf{A}_{\mathbf{S}^*} \bar{\mathbf{x}}_{\mathbf{S}^*} + \boldsymbol{\eta} \tag{2.3}$$

Model with deterministic conditions on η can also be analyzed.

- **Output** Bounding Generalization error Upper bound $G(\mathbf{x}) := \frac{1}{n} \|\mathbf{A}(\mathbf{x} \bar{\mathbf{x}})\|_2^2$ where the rows of \mathbf{A} are i.i.d.
- **② Support Recovery** Recover the true features of A, i.e., find a $S \supseteq S^*$.

Algorithms to solve SLR

The literature mainly studies 3 classes of algorithms

Existing SLR algorithms

- ullet ℓ_1 minimization based (LASSO based). E.g. Dantzig selector
- Non-convex penalty based. E.g. IHT, SCAD penalty, Log-sum penalty
- Greedy methods. E.g. Orthogonal Matching Pursuit (OMP)

We study SLR under RSC assumption for OMP algorithm.

Algorithms to solve SLR

The literature mainly studies 3 classes of algorithms

Existing SLR algorithms

- ullet ℓ_1 minimization based (LASSO based). E.g. Dantzig selector
- Non-convex penalty based. E.g. IHT, SCAD penalty, Log-sum penalty
- Greedy methods. E.g. Orthogonal Matching Pursuit (OMP)

We study SLR under RSC assumption for OMP algorithm.

Set initial support set $S_0 = \phi \& x_0 = 0$.

- ${f r}_0 : {f r}_0 = {f y} {f A}{f x}_0 = {f y}_0$.
- At k^{th} iteration ($k \ge 1$)
 - From the left-over columns of A (in $A_{S_{k-1}^c}$), find the column with maximum absolute inner product with \mathbf{r}_{k-1} .

$$[((A_{i_1 i_1 r_{k-1}}) | ((A_{i_1 i_1 r_{k-1}}) | \dots | ((A_{i_1 i_1 r_{k-1}}) | \dots | ((A_{i_{k-k+1} i_1 r_{k-1}}) |)]]$$

- Include i_i into the set: $\mathbf{S}_k = \mathbf{S}_{k-1} \cup \{i_i\}$.
- ullet Fully optimize on \mathbf{S}_k : $\mathbf{x}_k = rg \min_{\mathsf{supp}(\mathbf{x}) \subseteq \mathbf{S}_k} \|\mathbf{y} \mathbf{A}\mathbf{x}\|_2^2$ (Simple least squares)
- Update residual: $\mathbf{r}_k = \mathbf{y} \mathbf{A}\mathbf{x}_k$.

Set initial support set $S_0 = \phi \& x_0 = 0$.

$$\cdot$$
 residual $\mathbf{r}_0 = \mathbf{y} - \mathbf{A}\mathbf{x}_0 = \mathbf{y}$.

At k^{th} iteration $(k \ge 1)$

- From the left-over columns of **A** (in $A_{\mathbf{S}_{k-1}^*}$), find the column with maximum absolute inner product with \mathbf{r}_{k-1} .
 - $\left[\left|\left(\Delta_{11},r_{k-1}\right)\right|,\left|\left(\Delta_{21},r_{k-1}\right)\right|,\dots,\left|\left(\Delta_{1j},r_{k-1}\right)\right|,\dots,\left|\left(\Delta_{1j-k+1},r_{k-1}\right)\right|\right]\right]$
- Include i_i into the set: $\mathbf{S}_k = \mathbf{S}_{k-1} \cup \{i_i\}$.
- ullet Fully optimize on \mathbf{S}_k : $\mathbf{x}_k = rg \min_{\mathsf{supp}(\mathbf{x}) \subseteq \mathbf{S}_k} \|\mathbf{y} \mathbf{A}\mathbf{x}\|_2^2$ (Simple least squares)
- Update residual: $\mathbf{r}_k = \mathbf{y} \mathbf{A}\mathbf{x}_k$.

Set initial support set $S_0 = \phi \& x_0 = 0$. \therefore residual $\mathbf{r}_0 = \mathbf{y} - \mathbf{A}\mathbf{x}_0 = \mathbf{y}$. At k^{th} iteration (k > 1)

$$\left[\left| \left\langle \mathbf{A}_{i_1}, \mathbf{r}_{k-1} \right\rangle \right| \, \left| \left\langle \mathbf{A}_{i_2}, \mathbf{r}_{k-1} \right\rangle \right| \, \dots \, \left| \left\langle \mathbf{A}_{i_j}, \mathbf{r}_{k-1} \right\rangle \right| \, \dots \, \left| \left\langle \mathbf{A}_{i_{d-k+1}}, \mathbf{r}_{k-1} \right\rangle \right| \right]$$

- Include i_i into the set: $S_k = S_{k-1} \cup \{i_i\}$.
- Fully optimize on \mathbf{S}_k : $\mathbf{x}_k = \underset{\mathsf{supp}(\mathbf{x}) \subseteq \mathbf{S}_k}{\arg\min} \|\mathbf{y} \mathbf{A}\mathbf{x}\|_2^2$ (Simple least squares).
- Update residual: $\mathbf{r}_k = \mathbf{y} \mathbf{A}\mathbf{x}_k$.

Set initial support set $S_0 = \phi \& x_0 = 0$.

$$\therefore$$
 residual $\mathbf{r}_0 = \mathbf{y} - \mathbf{A}\mathbf{x}_0 = \mathbf{y}$.

At k^{th} iteration $(k \ge 1)$

$$\left[\left| \left\langle \mathbf{A}_{i_1}, \mathbf{r}_{k-1} \right\rangle \right| \, \left| \left\langle \mathbf{A}_{i_2}, \mathbf{r}_{k-1} \right\rangle \right| \, \dots \, \left| \left\langle \mathbf{A}_{i_j}, \mathbf{r}_{k-1} \right\rangle \right| \, \dots \, \left| \left\langle \mathbf{A}_{i_{d-k+1}}, \mathbf{r}_{k-1} \right\rangle \right| \, \right]$$

- Include i_i into the set: $\mathbf{S}_k = \mathbf{S}_{k-1} \cup \{i_i\}$.
- Fully optimize on \mathbf{S}_k : $\mathbf{x}_k = \underset{\mathsf{supp}(\mathbf{x}) \subseteq \mathbf{S}_k}{\arg\min} \|\mathbf{y} \mathbf{A}\mathbf{x}\|_2^2$ (Simple least squares).
- Update residual: $\mathbf{r}_k = \mathbf{y} \mathbf{A}\mathbf{x}_k$.

Set initial support set $S_0 = \phi \& x_0 = 0$.

$$\therefore$$
 residual $\mathbf{r}_0 = \mathbf{y} - \mathbf{A}\mathbf{x}_0 = \mathbf{y}$.

At k^{th} iteration $(k \ge 1)$

$$\left[\left| \langle \mathbf{A}_{i_1}, \mathbf{r}_{k-1} \rangle \right| \ \left| \langle \mathbf{A}_{i_2}, \mathbf{r}_{k-1} \rangle \right| \ \dots \ \left| \langle \mathbf{A}_{i_j}, \mathbf{r}_{k-1} \rangle \right| \ \dots \ \left| \langle \mathbf{A}_{i_{d-k+1}}, \mathbf{r}_{k-1} \rangle \right| \right] \right]$$

- Include i_i into the set: $\mathbf{S}_k = \mathbf{S}_{k-1} \cup \{i_i\}$.
- Fully optimize on \mathbf{S}_k : $\mathbf{x}_k = \underset{\mathsf{supp}(\mathbf{x}) \subseteq \mathbf{S}_k}{\arg\min} \|\mathbf{y} \mathbf{A}\mathbf{x}\|_2^2$ (Simple least squares).
- Update residual: $\mathbf{r}_k = \mathbf{y} \mathbf{A}\mathbf{x}_k$.



Set initial support set $S_0 = \phi \& x_0 = 0$. \therefore residual $\mathbf{r}_0 = \mathbf{y} - \mathbf{A}\mathbf{x}_0 = \mathbf{y}$.

At k^{th} iteration $(k \ge 1)$

$$\left[\left| \langle \mathbf{A}_{i_1}, \mathbf{r}_{k-1} \rangle \right| \, \left| \langle \mathbf{A}_{i_2}, \mathbf{r}_{k-1} \rangle \right| \, \dots \, \left| \left\langle \mathbf{A}_{i_j}, \mathbf{r}_{k-1} \right\rangle \right| \, \dots \, \left| \left\langle \mathbf{A}_{i_{d-k+1}}, \mathbf{r}_{k-1} \right\rangle \right| \right]$$

- Include i_i into the set: $S_k = S_{k-1} \cup \{i_i\}$.
- Fully optimize on \mathbf{S}_k : $\mathbf{x}_k = \underset{\mathsf{supp}(\mathbf{x}) \subseteq \mathbf{S}_k}{\arg\min} \|\mathbf{y} \mathbf{A}\mathbf{x}\|_2^2$ (Simple least squares).
- Update residual: $\mathbf{r}_k = \mathbf{y} \mathbf{A}\mathbf{x}_k$.



Set initial support set $S_0 = \phi \& x_0 = 0$. \therefore residual $\mathbf{r}_0 = \mathbf{y} - \mathbf{A}\mathbf{x}_0 = \mathbf{y}$.

At k^{th} iteration $(k \ge 1)$

• From the left-over columns of A (in $A_{S_{k-1}^c}$), find the column with maximum absolute inner product with \mathbf{r}_{k-1} .

$$\left[\left| \langle \mathbf{A}_{i_1}, \mathbf{r}_{k-1} \rangle \right| \, \left| \langle \mathbf{A}_{i_2}, \mathbf{r}_{k-1} \rangle \right| \, \dots \, \left| \langle \mathbf{A}_{i_j}, \mathbf{r}_{k-1} \rangle \right| \, \dots \, \left| \langle \mathbf{A}_{i_{d-k+1}}, \mathbf{r}_{k-1} \rangle \right| \right]$$

- Include i_i into the set: $S_k = S_{k-1} \cup \{i_i\}$.
- Fully optimize on \mathbf{S}_k : $\mathbf{x}_k = \underset{\mathsf{supp}(\mathbf{x}) \subseteq \mathbf{S}_k}{\arg\min} \|\mathbf{y} \mathbf{A}\mathbf{x}\|_2^2$ (Simple least squares).
- Update residual: $\mathbf{r}_k = \mathbf{y} \mathbf{A}\mathbf{x}_k$.



Set initial support set $S_0 = \phi \& x_0 = 0$. \therefore residual $\mathbf{r}_0 = \mathbf{y} - \mathbf{A}\mathbf{x}_0 = \mathbf{y}$. At k^{th} iteration $(k \ge 1)$

• From the left-over columns of A (in $A_{S_{k-1}^c}$), find the column with maximum absolute inner product with \mathbf{r}_{k-1} .

$$\left[\left| \langle \mathbf{A}_{i_1}, \mathbf{r}_{k-1} \rangle \right| \, \left| \langle \mathbf{A}_{i_2}, \mathbf{r}_{k-1} \rangle \right| \, \dots \, \left| \left\langle \mathbf{A}_{i_j}, \mathbf{r}_{k-1} \right\rangle \right| \, \dots \, \left| \left\langle \mathbf{A}_{i_{d-k+1}}, \mathbf{r}_{k-1} \right\rangle \right| \right]$$

- Include i_i into the set: $S_k = S_{k-1} \cup \{i_i\}$.
- Fully optimize on \mathbf{S}_k : $\mathbf{x}_k = \underset{\mathsf{supp}(\mathbf{x}) \subseteq \mathbf{S}_k}{\arg\min} \|\mathbf{y} \mathbf{A}\mathbf{x}\|_2^2$ (Simple least squares).
- Update residual: $\mathbf{r}_k = \mathbf{y} \mathbf{A}\mathbf{x}_k$.



Set initial support set $S_0 = \phi \& x_0 = 0$. \therefore residual $\mathbf{r}_0 = \mathbf{y} - \mathbf{A}\mathbf{x}_0 = \mathbf{y}$. At k^{th} iteration $(k \ge 1)$

• From the left-over columns of A (in $A_{S_{k-1}^c}$), find the column with maximum absolute inner product with \mathbf{r}_{k-1} .

$$\left[\left| \langle \mathbf{A}_{i_1}, \mathbf{r}_{k-1} \rangle \right| \, \left| \langle \mathbf{A}_{i_2}, \mathbf{r}_{k-1} \rangle \right| \, \dots \, \left| \left\langle \mathbf{A}_{i_j}, \mathbf{r}_{k-1} \right\rangle \right| \, \dots \, \left| \left\langle \mathbf{A}_{i_{d-k+1}}, \mathbf{r}_{k-1} \right\rangle \right| \right]$$

- Include i_i into the set: $\mathbf{S}_k = \mathbf{S}_{k-1} \cup \{i_i\}$.
- Fully optimize on \mathbf{S}_k : $\mathbf{x}_k = \underset{\mathsf{supp}(\mathbf{x}) \subseteq \mathbf{S}_k}{\arg\min} \|\mathbf{y} \mathbf{A}\mathbf{x}\|_2^2$ (Simple least squares).
- Update residual: $\mathbf{r}_k = \mathbf{y} \mathbf{A}\mathbf{x}_k$.



Set initial support set $\mathbf{S}_0 = \phi \& \mathbf{x}_0 = \mathbf{0}$. \therefore residual $\mathbf{r}_0 = \mathbf{y} - \mathbf{A}\mathbf{x}_0 = \mathbf{y}$. At k^{th} iteration $(k \ge 1)$

• From the left-over columns of A (in $A_{S_{k-1}^c}$), find the column with maximum absolute inner product with \mathbf{r}_{k-1} .

$$\left[\left| \langle \mathbf{A}_{i_1}, \mathbf{r}_{k-1} \rangle \right| \ \left| \langle \mathbf{A}_{i_2}, \mathbf{r}_{k-1} \rangle \right| \ \dots \ \left| \left\langle \mathbf{A}_{i_j}, \mathbf{r}_{k-1} \right\rangle \right| \ \dots \ \left| \left\langle \mathbf{A}_{i_{d-k+1}}, \mathbf{r}_{k-1} \right\rangle \right| \right]$$

- Include i_j into the set: $\mathbf{S}_k = \mathbf{S}_{k-1} \cup \{i_j\}$.
- Fully optimize on \mathbf{S}_k : $\mathbf{x}_k = \underset{\mathsf{supp}(\mathbf{x}) \subseteq \mathbf{S}_k}{\arg\min} \|\mathbf{y} \mathbf{A}\mathbf{x}\|_2^2$ (Simple least squares).
- Update residual: $\mathbf{r}_k = \mathbf{y} \mathbf{A}\mathbf{x}_k$.



Set initial support set $S_0 = \phi \& x_0 = 0$. \therefore residual $\mathbf{r}_0 = \mathbf{y} - \mathbf{A}\mathbf{x}_0 = \mathbf{y}$. At k^{th} iteration (k > 1)

$$\left[\left| \langle \mathbf{A}_{i_1}, \mathbf{r}_{k-1} \rangle \right| \ \left| \langle \mathbf{A}_{i_2}, \mathbf{r}_{k-1} \rangle \right| \ \dots \ \left| \left\langle \mathbf{A}_{i_j}, \mathbf{r}_{k-1} \right\rangle \right| \ \dots \ \left| \left\langle \mathbf{A}_{i_{d-k+1}}, \mathbf{r}_{k-1} \right\rangle \right| \right]$$

- Include i_i into the set: $\mathbf{S}_k = \mathbf{S}_{k-1} \cup \{i_i\}$.
- Fully optimize on \mathbf{S}_k : $\mathbf{x}_k = \underset{\mathsf{supp}(\mathbf{x}) \subseteq \mathbf{S}_k}{\arg\min} \|\mathbf{y} \mathbf{A}\mathbf{x}\|_2^2$ (Simple least squares).
- Update residual: $\mathbf{r}_k = \mathbf{y} \mathbf{A}\mathbf{x}_k$.



```
 \begin{aligned} & \textbf{Result: OMP sparse estimate } \hat{\mathbf{x}}_s^{\text{OMP}} = \mathbf{x}_s \\ & \mathbf{S}_0 = \phi, \, \mathbf{x}_0 = \mathbf{0}, \, \mathbf{r}_0 = \mathbf{y} \\ & \textbf{for } k = 1, 2, \dots, s \, \textbf{do} \\ & | j \leftarrow \arg\max_{i \not \in \mathbf{S}_{k-1}} | \mathbf{A}_i^T \mathbf{r}_{k-1} | \quad \text{(Greedy selection)} \\ & \mathbf{S}_k \leftarrow \mathbf{S}_{k-1} \cup \{j\} \\ & \mathbf{x}_k \leftarrow \arg\min_{\substack{\mathbf{supp}(\mathbf{x}) \subseteq \mathbf{S}_k \\ \mathbf{r}_k \leftarrow \mathbf{y} - \mathbf{A} \mathbf{x}_k} \end{aligned}
```

Algorithm 1: Orthogonal Matching Pursuit (OMP) for SLR

Note that $\mathbf{A}_i^T \mathbf{r}_{k-1} \propto [\nabla f(\mathbf{x}_{k-1})]_i$ for $f(\mathbf{x}) = \|\mathbf{A}\mathbf{x} - \mathbf{y}\|^2$

end

$$\begin{aligned} & \textbf{Result: OMP sparse estimate } \hat{\mathbf{x}}_s^{\text{OMP}} = \mathbf{x}_s \\ & \mathbf{S}_0 = \phi, \, \mathbf{x}_0 = \mathbf{0}, \, \mathbf{r}_0 = \mathbf{y} \\ & \textbf{for } k = 1, 2, \dots, s \, \textbf{do} \\ & j \leftarrow \underset{i \not\in \mathbf{S}_{k-1}}{\operatorname{max}} |\mathbf{A}_i^T \mathbf{r}_{k-1}| \quad \text{(Greedy selection)} \\ & \mathbf{S}_k \leftarrow \mathbf{S}_{k-1} \cup \{j\} \\ & \mathbf{x}_k \leftarrow \underset{\underset{\text{supp}(\mathbf{x}) \subseteq \mathbf{S}_k}{\operatorname{arg min}} \quad \|\mathbf{A}\mathbf{x} - \mathbf{y}\|_2^2 \\ & \mathbf{r}_k \leftarrow \mathbf{y} - \mathbf{A}\mathbf{x}_k \end{aligned}$$

Algorithm 1: Orthogonal Matching Pursuit (OMP) for SLR

Note that
$$\mathbf{A}_i^T \mathbf{r}_{k-1} \propto \left[\nabla f(\mathbf{x}_{k-1}) \right]_i$$
 for $f(\mathbf{x}) = \left\| \mathbf{A} \mathbf{x} - \mathbf{y} \right\|_2^2$

end

Orthogonal Matching Pursuit for general $f(\mathbf{x})$

```
\begin{aligned} & \textbf{Result: OMP sparse estimate } \hat{\mathbf{x}}_s^{\text{OMP}} = \mathbf{x}_s \\ & \mathbf{S}_0 = \phi, \, \mathbf{x}_0 = \mathbf{0} \\ & \textbf{for } k = 1, 2, \dots, s \textbf{ do} \\ & j := \underset{i \not\in \mathbf{S}_{k-1}}{\operatorname{arg max}} \left| \left[ \nabla f(\mathbf{x}_{k-1}) \right]_i \right| \\ & \mathbf{S}_k := \mathbf{S}_{k-1} \cup \{j\} \\ & \mathbf{x}_k := \underset{\text{supp}(\mathbf{x}) \subseteq \mathbf{S}_k}{\operatorname{arg min}} f(\mathbf{x}) \end{aligned}
```

Algorithm 2: OMP for a general function $f(\mathbf{x})$

Key quantities

• Restricted Smoothness (ρ^+) & Restricted Strong Convexity (ρ^-)

$$\rho_s^- \|\mathbf{x} - \mathbf{z}\|_2^2 \le \|\mathbf{A}\mathbf{x} - \mathbf{A}\mathbf{z}\|_2^2 \le \rho_s^+ \|\mathbf{x} - \mathbf{z}\|_2^2$$
(4.1)

 $\forall \mathbf{x}, \mathbf{z} \in \mathbb{R}^d \text{ s.t. } \|\mathbf{x} - \mathbf{z}\|_0 \leq s.$

• Restricted condition number $(\widetilde{\kappa}_s)$

$$\widetilde{\kappa}_s = \frac{\rho_1^+}{\rho_s^-} \tag{4.2}$$

We also define

$$\kappa_s = \frac{\rho_s^+}{\rho_s^-} \ge \widetilde{\kappa}_s$$
(4.3)

Key quantities

• Restricted Smoothness (ρ^+) & Restricted Strong Convexity (ρ^-)

$$\rho_s^- \|\mathbf{x} - \mathbf{z}\|_2^2 \le \|\mathbf{A}\mathbf{x} - \mathbf{A}\mathbf{z}\|_2^2 \le \rho_s^+ \|\mathbf{x} - \mathbf{z}\|_2^2$$
(4.1)

 $\forall \mathbf{x}, \mathbf{z} \in \mathbb{R}^d \text{ s.t. } \|\mathbf{x} - \mathbf{z}\|_0 \leq s.$

• Restricted condition number $(\widetilde{\kappa}_s)$

$$\widetilde{\kappa}_s = \frac{\rho_1^+}{\rho_s^-} \tag{4.2}$$

We also define

$$\kappa_s = \frac{\rho_s^+}{\rho_s^-} \ge \widetilde{\kappa}_s \tag{4.3}$$

Key quantities

• Restricted Smoothness (ρ^+) & Restricted Strong Convexity (ρ^-)

$$\rho_s^- \|\mathbf{x} - \mathbf{z}\|_2^2 \le \|\mathbf{A}\mathbf{x} - \mathbf{A}\mathbf{z}\|_2^2 \le \rho_s^+ \|\mathbf{x} - \mathbf{z}\|_2^2$$
(4.1)

 $\forall \mathbf{x}, \mathbf{z} \in \mathbb{R}^d \text{ s.t. } \|\mathbf{x} - \mathbf{z}\|_0 \leq s.$

• Restricted condition number $(\widetilde{\kappa}_s)$

$$\widetilde{\kappa}_s = \frac{\rho_1^+}{\rho_s^-} \tag{4.2}$$

We also define

$$\kappa_s = \frac{\rho_s^+}{\rho_s^-} \ge \widetilde{\kappa}_s \tag{4.3}$$

Lower bounds for Fast rates

• If $\hat{\mathbf{x}}_{\ell_0}$ is the best ℓ_0 estimate in the set of s^* -sparse vectors then one can show

$$\sup_{\|\bar{\mathbf{x}}\|_{0} \leq s^{*}} \frac{1}{n} \mathbb{E}\left[\|\mathbf{A}(\hat{\mathbf{x}}_{\ell_{0}} - \bar{\mathbf{x}})\|_{2}^{2} \right] \lesssim \frac{\sigma^{2} s^{*}}{n}$$
(4.4)

- Non-tractable since computing $\hat{\mathbf{x}}_{\ell_0}$ involves searching all $\binom{d}{s^*}$ subsets.
- (Y. Zhang, Wainwright & Jordan'15) \exists $\mathbf{A} \in \mathbb{R}^{n \times d}$ s.t. any poly-time algorithm satisfies

$$\sup_{\|\bar{\mathbf{x}}\|_{0} \leq s^{*}} \frac{1}{n} \mathbb{E}\left[\left\| \mathbf{A} (\hat{\mathbf{x}}_{\mathsf{poly}} - \bar{\mathbf{x}}) \right\|_{2}^{2} \right] \gtrsim \frac{\sigma^{2} s^{*1 - \delta} \tilde{\kappa}_{s^{*}}}{n} \qquad \forall \ \delta > 0$$
 (4.5)

Consequence - Any estimator $\hat{\mathbf{x}}$ achieving fast rate must either **not be poly-time** or must return $\hat{\mathbf{x}}_{\text{poly}}$ that is **not** s^* -sparse.

Lower bounds for Fast rates

• If $\hat{\mathbf{x}}_{\ell_0}$ is the best ℓ_0 estimate in the set of s^* -sparse vectors then one can show

$$\sup_{\|\bar{\mathbf{x}}\|_{0} \leq s^{*}} \frac{1}{n} \mathbb{E}\left[\|\mathbf{A}(\hat{\mathbf{x}}_{\ell_{0}} - \bar{\mathbf{x}})\|_{2}^{2} \right] \lesssim \frac{\sigma^{2} s^{*}}{n}$$
(4.4)

- Non-tractable since computing $\hat{\mathbf{x}}_{\ell_0}$ involves searching all $\binom{d}{s^*}$ subsets.
- (Y. Zhang, Wainwright & Jordan'15) \exists $\mathbf{A} \in \mathbb{R}^{n \times d}$ s.t. any poly-time algorithm satisfies

$$\sup_{\|\bar{\mathbf{x}}\|_{0} \leq s^{*}} \frac{1}{n} \mathbb{E}\left[\left\| \mathbf{A} (\hat{\mathbf{x}}_{\mathsf{poly}} - \bar{\mathbf{x}}) \right\|_{2}^{2} \right] \gtrsim \frac{\sigma^{2} s^{*1 - \delta} \tilde{\kappa}_{s^{*}}}{n} \qquad \forall \ \delta > 0$$
 (4.5)

Consequence - Any estimator $\hat{\mathbf{x}}$ achieving fast rate must either **not be poly-time** or must return $\hat{\mathbf{x}}_{\text{poly}}$ that is **not** s^* -sparse.

Lower bounds for Fast rates

• If $\hat{\mathbf{x}}_{\ell_0}$ is the best ℓ_0 estimate in the set of s^* -sparse vectors then one can show

$$\sup_{\|\bar{\mathbf{x}}\|_{0} \leq s^{*}} \frac{1}{n} \mathbb{E}\left[\|\mathbf{A}(\hat{\mathbf{x}}_{\ell_{0}} - \bar{\mathbf{x}})\|_{2}^{2} \right] \lesssim \frac{\sigma^{2} s^{*}}{n}$$
(4.4)

- Non-tractable since computing $\hat{\mathbf{x}}_{\ell_0}$ involves searching all $\binom{d}{s^*}$ subsets.
- (Y. Zhang, Wainwright & Jordan'15) \exists $\mathbf{A} \in \mathbb{R}^{n \times d}$ s.t. any poly-time algorithm satisfies

$$\sup_{\|\bar{\mathbf{x}}\|_{0} \le s^{*}} \frac{1}{n} \mathbb{E}\left[\left\| \mathbf{A}(\hat{\mathbf{x}}_{\mathsf{poly}} - \bar{\mathbf{x}}) \right\|_{2}^{2} \right] \gtrsim \frac{\sigma^{2} s^{*1 - \delta} \widetilde{\kappa}_{s^{*}}}{n} \qquad \forall \ \delta > 0 \tag{4.5}$$

Consequence - Any estimator $\hat{\mathbf{x}}$ achieving fast rate must either **not be poly-time** or must return $\hat{\mathbf{x}}_{\text{poly}}$ that is **not** s^* -sparse.

Lower bounds for Fast rates

• If $\hat{\mathbf{x}}_{\ell_0}$ is the best ℓ_0 estimate in the set of s^* -sparse vectors then one can show

$$\sup_{\|\bar{\mathbf{x}}\|_{0} \leq s^{*}} \frac{1}{n} \mathbb{E}\left[\|\mathbf{A}(\hat{\mathbf{x}}_{\ell_{0}} - \bar{\mathbf{x}})\|_{2}^{2} \right] \lesssim \frac{\sigma^{2} s^{*}}{n}$$
(4.4)

- Non-tractable since computing $\hat{\mathbf{x}}_{\ell_0}$ involves searching all $\binom{d}{s^*}$ subsets.
- (Y. Zhang, Wainwright & Jordan'15) \exists $\mathbf{A} \in \mathbb{R}^{n \times d}$ s.t. any poly-time algorithm satisfies

$$\sup_{\|\bar{\mathbf{x}}\|_{0} \le s^{*}} \frac{1}{n} \mathbb{E}\left[\left\| \mathbf{A}(\hat{\mathbf{x}}_{\mathsf{poly}} - \bar{\mathbf{x}}) \right\|_{2}^{2} \right] \gtrsim \frac{\sigma^{2} s^{*1 - \delta} \widetilde{\kappa}_{s^{*}}}{n} \qquad \forall \ \delta > 0 \tag{4.5}$$

Consequence - Any estimator $\hat{\mathbf{x}}$ achieving fast rate must either **not be poly-time** or must return $\hat{\mathbf{x}}_{\text{poly}}$ that is **not** s^* -sparse.

- Tightest known upper bounds for poly-time algorithms like IHT, OMP and Lasso were at least $\widetilde{\kappa}$ times worse than known lower bounds (Jain'14, T. Zhang'10, Y. Zhang'17).
- (T. Zhang'10) If $\hat{\mathbf{x}}_s$ be the output of OMP after $s \gtrsim s^* \widetilde{\kappa}_{s+s^*} \log \kappa_{s+s^*}$ iterations, then with high probability

$$\frac{1}{n} \|\mathbf{A}(\mathbf{x}_s - \bar{\mathbf{x}})\|_2^2 \lesssim \frac{1}{n} \sigma^2 s^* \widetilde{\kappa}_{s+s^*}^2 \log \kappa_{s+s^*}. \tag{4.6}$$

With slight modification to T. Zhang's analysis we get

Generalization error for OMP

If $\hat{\mathbf{x}}_s$ is the output of OMP after $s \gtrsim s^* \widetilde{\kappa}_{s+s^*} \log \kappa_{s+s^*}$ iterations, then with high probability

$$\frac{1}{n} \|\mathbf{A}(\mathbf{x}_s - \bar{\mathbf{x}})\|_2^2 \lesssim \frac{1}{n} \sigma^2 s^* \widetilde{\kappa}_{s+s^*} \log \kappa_{s+s^*}. \tag{4.7}$$

Matches the fast rate lower bound up to log factors



- Tightest known upper bounds for poly-time algorithms like IHT, OMP and Lasso were at least $\tilde{\kappa}$ times worse than known lower bounds (Jain'14, T. Zhang'10, Y. Zhang'17).
- (T. Zhang'10) If $\hat{\mathbf{x}}_s$ be the output of OMP after $s \gtrsim s^* \widetilde{\kappa}_{s+s^*} \log \kappa_{s+s^*}$ iterations, then with high probability

$$\frac{1}{n} \|\mathbf{A}(\mathbf{x}_s - \bar{\mathbf{x}})\|_2^2 \lesssim \frac{1}{n} \sigma^2 s^* \tilde{\kappa}_{s+s^*}^2 \log \kappa_{s+s^*}. \tag{4.6}$$

With slight modification to T. Zhang's analysis we get

Generalization error for OMF

If $\hat{\mathbf{x}}_s$ is the output of OMP after $s \gtrsim s^* \tilde{\kappa}_{s+s^*} \log \kappa_{s+s^*}$ iterations, then with high probability

$$\frac{1}{n} \|\mathbf{A}(\mathbf{x}_s - \bar{\mathbf{x}})\|_2^2 \lesssim \frac{1}{n} \sigma^2 s^* \widetilde{\kappa}_{s+s^*} \log \kappa_{s+s^*}. \tag{4.7}$$

Matches the fast rate lower bound up to log factors

- Tightest known upper bounds for poly-time algorithms like IHT, OMP and Lasso were at least $\tilde{\kappa}$ times worse than known lower bounds (Jain'14, T. Zhang'10, Y. Zhang'17).
- (T. Zhang'10) If $\hat{\mathbf{x}}_s$ be the output of OMP after $s \gtrsim s^* \widetilde{\kappa}_{s+s^*} \log \kappa_{s+s^*}$ iterations, then with high probability

$$\frac{1}{n} \|\mathbf{A}(\mathbf{x}_s - \bar{\mathbf{x}})\|_2^2 \lesssim \frac{1}{n} \sigma^2 s^* \widetilde{\kappa}_{s+s^*}^2 \log \kappa_{s+s^*}. \tag{4.6}$$

With slight modification to T. Zhang's analysis we get

Generalization error for OMP

If $\hat{\mathbf{x}}_s$ is the output of OMP after $s \gtrsim s^* \widetilde{\kappa}_{s+s^*} \log \kappa_{s+s^*}$ iterations, then with high probability

$$\frac{1}{n} \|\mathbf{A}(\mathbf{x}_s - \bar{\mathbf{x}})\|_2^2 \lesssim \frac{1}{n} \sigma^2 s^* \widetilde{\kappa}_{s+s^*} \log \kappa_{s+s^*}. \tag{4.7}$$

Matches the fast rate lower bound up to log factors

- Tightest known upper bounds for poly-time algorithms like IHT, OMP and Lasso were at least $\tilde{\kappa}$ times worse than known lower bounds (Jain'14, T. Zhang'10, Y. Zhang'17).
- (T. Zhang'10) If $\hat{\mathbf{x}}_s$ be the output of OMP after $s \gtrsim s^* \widetilde{\kappa}_{s+s^*} \log \kappa_{s+s^*}$ iterations, then with high probability

$$\frac{1}{n} \|\mathbf{A}(\mathbf{x}_s - \bar{\mathbf{x}})\|_2^2 \lesssim \frac{1}{n} \sigma^2 s^* \widetilde{\kappa}_{s+s^*}^2 \log \kappa_{s+s^*}. \tag{4.6}$$

With slight modification to T. Zhang's analysis we get

Generalization error for OMP

If $\hat{\mathbf{x}}_s$ is the output of OMP after $s \gtrsim s^* \widetilde{\kappa}_{s+s^*} \log \kappa_{s+s^*}$ iterations, then with high probability

$$\frac{1}{n} \|\mathbf{A}(\mathbf{x}_s - \bar{\mathbf{x}})\|_2^2 \lesssim \frac{1}{n} \sigma^2 s^* \widetilde{\kappa}_{s+s^*} \log \kappa_{s+s^*}. \tag{4.7}$$

Matches the fast rate lower bound up to log factors.

- Support recovery results are known for SCAD/MCP penalty based methods under bounded incoherence (Loh'14).
- For greedy algorithms like HTP and PHT, known support recovery results require poor dependence of $\widetilde{\kappa}$ on $|\overline{x}_{\min}|$ (Shen'17).
- If S is the support set of the s^{th} OMP iterate $\hat{\mathbf{x}}_s$, and if $\mathbf{S}^* \setminus \mathbf{S} \neq \phi$, then there is a **large additive decrease** in objective if $|\bar{x}_{\min}|$ is larger than the appropriate noise level.

Large decrease in objective

If $\hat{\mathbf{x}}_s$ is the output of OMP after $s \gtrsim s^* \widetilde{\kappa}_{s+s^*} \log \kappa_{s+s^*}$ iterations s.t. $\mathbf{S}^* \setminus \mathbf{S} \neq \phi$ and $|\bar{x}_{\min}| \gtrsim \frac{\sigma \gamma \sqrt{\rho_1^+}}{\rho_{s+s^*}^-}$ then with high probability

$$\|\mathbf{A}\hat{\mathbf{x}}_{s} - \mathbf{y}\|_{2}^{2} - \|\mathbf{A}\hat{\mathbf{x}}_{s+1} - \mathbf{y}\|_{2}^{2} \gtrsim \sigma^{2}$$
 (4.8)

where $\left\|\mathbf{A}_{\mathbf{S}^* \setminus \mathbf{S}}^T \mathbf{A}_{\mathbf{S}} \left(\mathbf{A}_{\mathbf{S}}^T \mathbf{A}_{\mathbf{S}}\right)^{-1}\right\|_{\infty} \leq \gamma$ and $\mathbf{S} = \operatorname{supp}(\hat{\mathbf{x}}_s)$.

y is similar to standard incoherence condition.



- Support recovery results are known for SCAD/MCP penalty based methods under bounded incoherence (Loh'14).
- For greedy algorithms like HTP and PHT, known support recovery results require poor dependence of $\widetilde{\kappa}$ on $|\bar{x}_{\min}|$ (Shen'17).
- If S is the support set of the s^{th} OMP iterate $\hat{\mathbf{x}}_s$, and if $\mathbf{S}^* \setminus \mathbf{S} \neq \phi$, then there is a **large additive decrease** in objective if $|\bar{x}_{\min}|$ is larger than the appropriate noise level.

Large decrease in objective

If $\hat{\mathbf{x}}_s$ is the output of OMP after $s \gtrsim s^* \widetilde{\kappa}_{s+s^*} \log \kappa_{s+s^*}$ iterations s.t. $\mathbf{S}^* \setminus \mathbf{S} \neq \phi$ and $|\bar{x}_{\min}| \gtrsim \frac{\sigma \gamma \sqrt{\rho_1^+}}{\rho_{s+s^*}^-}$ then with high probability

$$\|\mathbf{A}\hat{\mathbf{x}}_s - \mathbf{y}\|_2^2 - \|\mathbf{A}\hat{\mathbf{x}}_{s+1} - \mathbf{y}\|_2^2 \gtrsim \sigma^2$$

$$\tag{4.8}$$

where $\left\|\mathbf{A}_{\mathbf{S}^* \setminus \mathbf{S}}^T \mathbf{A}_{\mathbf{S}} \left(\mathbf{A}_{\mathbf{S}}^T \mathbf{A}_{\mathbf{S}}\right)^{-1}\right\|_{\infty} \leq \gamma$ and $\mathbf{S} = \operatorname{supp}(\hat{\mathbf{x}}_s)$.

is similar to standard incoherence condition.



- Support recovery results are known for SCAD/MCP penalty based methods under bounded incoherence (Loh'14).
- For greedy algorithms like HTP and PHT, known support recovery results require poor dependence of $\widetilde{\kappa}$ on $|\bar{x}_{\min}|$ (Shen'17).
- If S is the support set of the s^{th} OMP iterate $\hat{\mathbf{x}}_s$, and if $\mathbf{S}^* \setminus \mathbf{S} \neq \phi$, then there is a **large additive decrease** in objective if $|\bar{x}_{\min}|$ is larger than the appropriate noise level.

Large decrease in objective

If $\hat{\mathbf{x}}_s$ is the output of OMP after $s \gtrsim s^* \widetilde{\kappa}_{s+s^*} \log \kappa_{s+s^*}$ iterations s.t. $\mathbf{S}^* \setminus \mathbf{S} \neq \phi$ and $|\bar{x}_{\min}| \gtrsim \frac{\sigma \gamma \sqrt{\rho_1^+}}{\rho_{s+s^*}^-}$ then with high probability

$$\|\mathbf{A}\hat{\mathbf{x}}_s - \mathbf{y}\|_2^2 - \|\mathbf{A}\hat{\mathbf{x}}_{s+1} - \mathbf{y}\|_2^2 \gtrsim \sigma^2$$
(4.8)

where $\left\|\mathbf{A}_{\mathbf{S}^* \backslash \mathbf{S}}^T \mathbf{A}_{\mathbf{S}} \left(\mathbf{A}_{\mathbf{S}}^T \mathbf{A}_{\mathbf{S}}\right)^{-1}\right\|_{\infty} \leq \gamma \text{ and } \mathbf{S} = \text{supp}(\hat{\mathbf{x}}_s).$

y is similar to standard incoherence condition.



- Support recovery results are known for SCAD/MCP penalty based methods under bounded incoherence (Loh'14).
- For greedy algorithms like HTP and PHT, known support recovery results require poor dependence of $\widetilde{\kappa}$ on $|\bar{x}_{\min}|$ (Shen'17).
- If S is the support set of the s^{th} OMP iterate $\hat{\mathbf{x}}_s$, and if $\mathbf{S}^* \setminus \mathbf{S} \neq \phi$, then there is a **large additive decrease** in objective if $|\bar{x}_{\min}|$ is larger than the appropriate noise level.

Large decrease in objective

If $\hat{\mathbf{x}}_s$ is the output of OMP after $s \gtrsim s^* \widetilde{\kappa}_{s+s^*} \log \kappa_{s+s^*}$ iterations s.t. $\mathbf{S}^* \setminus \mathbf{S} \neq \phi$ and $|\bar{x}_{\min}| \gtrsim \frac{\sigma \gamma \sqrt{\rho_1^+}}{\rho_{s+s^*}^-}$ then with high probability

$$\|\mathbf{A}\hat{\mathbf{x}}_{s} - \mathbf{y}\|_{2}^{2} - \|\mathbf{A}\hat{\mathbf{x}}_{s+1} - \mathbf{y}\|_{2}^{2} \gtrsim \sigma^{2}$$
 (4.8)

where $\left\|\mathbf{A}_{\mathbf{S}^* \setminus \mathbf{S}}^T \mathbf{A}_{\mathbf{S}} \left(\mathbf{A}_{\mathbf{S}}^T \mathbf{A}_{\mathbf{S}}\right)^{-1}\right\|_{\infty} \leq \gamma \text{ and } \mathbf{S} = \operatorname{supp}(\hat{\mathbf{x}}_s).$

y is similar to standard incoherence condition.

- Support recovery results are known for SCAD/MCP penalty based methods under bounded incoherence (Loh'14).
- For greedy algorithms like HTP and PHT, known support recovery results require poor dependence of $\tilde{\kappa}$ on $|\bar{x}_{\min}|$ (Shen'17).
- If S is the support set of the s^{th} OMP iterate $\hat{\mathbf{x}}_s$, and if $\mathbf{S}^* \setminus \mathbf{S} \neq \phi$, then there is a large additive decrease in objective if $|\bar{x}_{\min}|$ is larger than the appropriate noise level.

Large decrease in objective

If $\hat{\mathbf{x}}_s$ is the output of OMP after $s \gtrsim s^* \tilde{\kappa}_{s+s^*} \log \kappa_{s+s^*}$ iterations s.t. $\mathbf{S}^* \setminus \mathbf{S} \neq \phi$ and $|\bar{x}_{\min}| \gtrsim \frac{\sigma \gamma \sqrt{\rho_1^+}}{\rho_{-+}^+}$ then with high probability

$$\|\mathbf{A}\hat{\mathbf{x}}_{s} - \mathbf{y}\|_{2}^{2} - \|\mathbf{A}\hat{\mathbf{x}}_{s+1} - \mathbf{y}\|_{2}^{2} \gtrsim \sigma^{2}$$
 (4.8)

where $\left\|\mathbf{A}_{\mathbf{S}^* \setminus \mathbf{S}}^T \mathbf{A}_{\mathbf{S}} \left(\mathbf{A}_{\mathbf{S}}^T \mathbf{A}_{\mathbf{S}}\right)^{-1}\right\| \leq \gamma \text{ and } \mathbf{S} = \operatorname{supp}(\hat{\mathbf{x}}_s).$

 γ is similar to standard incoherence condition.

Since $\|\mathbf{A}\mathbf{x} - \mathbf{y}\|_2^2 \ge 0 \ \forall \ \mathbf{x} \in \mathbb{R}^d$, the number of extra iterations cannot be too large.

Support recovery and infinity norm bound

If $\hat{\mathbf{x}}_s$ is the output of OMP after $s \gtrsim s^* \widetilde{\kappa}_{s+s^*} \log \kappa_{s+s^*}$ iterations, s.t

 $|\bar{x}_{\min}| \gtrsim rac{\sigma \gamma \sqrt{
ho_1^+}}{
ho_{-s+s}^-}$, then with high probability

- \mathbf{O} $\mathbf{S}^* \subseteq \operatorname{supp}(\hat{\mathbf{x}}_s)$

where
$$\left\|\mathbf{A}_{\mathbf{S}^* \setminus \mathbf{S}}^T \mathbf{A}_{\mathbf{S}} \left(\mathbf{A}_{\mathbf{S}}^T \mathbf{A}_{\mathbf{S}}\right)^{-1}\right\|_{\infty} \leq \gamma$$
 and $\mathbf{S} = \operatorname{supp}(\hat{\mathbf{x}}_s)$.

- Condition on $|\bar{x}_{\min}|$ scales as $\frac{1}{\sqrt{n}}$ since both $\rho_{s+s^*}^-$ and ρ_1^+ have a factor of n. Also it is better by at-least $\sqrt{\kappa}$ than that in other recent works.
- \bullet γ is allowed to be very large.



Since $\|\mathbf{A}\mathbf{x} - \mathbf{y}\|_2^2 \ge 0 \ \forall \ \mathbf{x} \in \mathbb{R}^d$, the number of extra iterations cannot be too large.

Support recovery and infinity norm bound

If $\hat{\mathbf{x}}_s$ is the output of OMP after $s \gtrsim s^* \tilde{\kappa}_{s+s^*} \log \kappa_{s+s^*}$ iterations, s.t.

$$|ar{x}_{\min}|\gtrsim rac{\sigma\gamma\sqrt{
ho_1^+}}{
ho_{s+s^*}^-}$$
 , then with high probability

where
$$\left\|\mathbf{A}_{\mathbf{S}^* \setminus \mathbf{S}}^T \mathbf{A}_{\mathbf{S}} \left(\mathbf{A}_{\mathbf{S}}^T \mathbf{A}_{\mathbf{S}}\right)^{-1}\right\|_{\infty} \leq \gamma \text{ and } \mathbf{S} = \operatorname{supp}(\hat{\mathbf{x}}_s).$$

• Condition on $|\bar{x}_{\min}|$ scales as $\frac{1}{\sqrt{n}}$ since both $\rho_{s+s^*}^-$ and ρ_1^+ have a factor of n. Also it is better by at-least \sqrt{k} than that in other recent works.

Since $\|\mathbf{A}\mathbf{x} - \mathbf{y}\|_2^2 \ge 0 \ \forall \ \mathbf{x} \in \mathbb{R}^d$, the number of extra iterations cannot be too large.

Support recovery and infinity norm bound

If $\hat{\mathbf{x}}_s$ is the output of OMP after $s \gtrsim s^* \tilde{\kappa}_{s+s^*} \log \kappa_{s+s^*}$ iterations, s.t.

 $|\bar{x}_{\min}| \gtrsim rac{\sigma \gamma \sqrt{
ho_1^+}}{
ho_{\circ+\circ}^-}$, then with high probability

- $\|\hat{\mathbf{x}}_s \bar{\mathbf{x}}\|_{\infty} \lesssim \sigma \sqrt{\frac{\log s}{\rho_s^-}}$

where $\left\|\mathbf{A}_{\mathbf{S}^* \setminus \mathbf{S}}^T \mathbf{A}_{\mathbf{S}} \left(\mathbf{A}_{\mathbf{S}}^T \mathbf{A}_{\mathbf{S}}\right)^{-1}\right\|_{\infty} \leq \gamma \text{ and } \mathbf{S} = \text{supp}(\hat{\mathbf{x}}_s).$

- Condition on $|\bar{x}_{\min}|$ scales as $\frac{1}{\sqrt{n}}$ since both $\rho_{s+s^*}^-$ and ρ_1^+ have a factor of n. Also it is better by at-least $\sqrt{\tilde{\kappa}}$ than that in other recent works.
- \bullet γ is allowed to be very large.



Since $\|\mathbf{A}\mathbf{x} - \mathbf{y}\|_2^2 \ge 0 \ \forall \ \mathbf{x} \in \mathbb{R}^d$, the number of extra iterations cannot be too large.

Support recovery and infinity norm bound

If $\hat{\mathbf{x}}_s$ is the output of OMP after $s \gtrsim s^* \tilde{\kappa}_{s+s^*} \log \kappa_{s+s^*}$ iterations, s.t.

 $|ar{x}_{\min}|\gtrsim rac{\sigma\gamma\sqrt{
ho_1^+}}{
ho_{=+*}^-}$, then with high probability

- $\|\hat{\mathbf{x}}_s \bar{\mathbf{x}}\|_{\infty} \lesssim \sigma \sqrt{\frac{\log s}{\rho_s^-}}$

where $\left\|\mathbf{A}_{\mathbf{S}^* \setminus \mathbf{S}}^T \mathbf{A}_{\mathbf{S}} \left(\mathbf{A}_{\mathbf{S}}^T \mathbf{A}_{\mathbf{S}}\right)^{-1}\right\|_{\infty} \leq \gamma \text{ and } \mathbf{S} = \text{supp}(\hat{\mathbf{x}}_s).$

- Condition on $|\bar{x}_{\min}|$ scales as $\frac{1}{\sqrt{n}}$ since both $\rho_{s+s^*}^-$ and ρ_1^+ have a factor of n. Also it is better by at-least $\sqrt{\tilde{\kappa}}$ than that in other recent works.
- γ is allowed to be very large.



- (Y. Zhang'15)'s lower bounds were for algorithms that output s^* -sparse solutions which does not apply for OMP when it is run for more than s^* iterations.
- We provide matching lower bounds for Support recovery as well as
- Fool OMP into picking incorrect indexes. Large support size ⇒ large
- Construct an evenly distributed x̄

$$\bar{x}_i = \begin{cases} \frac{1}{\sqrt{s^*}} & \text{if } 1 \leq i \leq s^* \\ 0 & \text{if } i > s^* \end{cases} \implies \operatorname{supp}(\bar{\mathbf{x}}) = \{1, 2, \dots, s^*\}$$

- Construct $\mathbf{M}^{(\epsilon)} \in \mathbb{R}^{n \times d}$ parameterized by ϵ
 - $\bullet \ \ \mathbf{M}_{1:s^*}^{(\epsilon)} \ \text{are random } s^* \text{ orthogonal column vectors s.t. } \left\| \mathbf{M}_i^{(\epsilon)} \right\|_2^2 = n \ \forall \ i \in [s^*].$
 - $\bullet \ \ \mathbf{M}_i^{(\epsilon)} = \sqrt{1-\epsilon} \left[\tfrac{1}{\sqrt{s^*}} \sum_{j=1}^{s^*} \mathbf{M}_j^{(\epsilon)} \right] + \sqrt{\epsilon} \mathbf{g}_i \ \forall \ i \not \in [s^*] \ \text{where} \ \mathbf{g}_i \text{'s are}$ orthogonal to each other and $\mathbf{M}_{1:s^*}^{(\epsilon)}$ with $\|\mathbf{g}_i\|_2^2 = n$.

- (Y. Zhang'15)'s lower bounds were for algorithms that output s^* -sparse solutions which does not apply for OMP when it is run for more than s^* iterations.
- We provide matching lower bounds for Support recovery as well as generalization error for OMP.
- \bullet Fool OMP into picking incorrect indexes. Large support size \implies large generalization error.
- ullet Construct an evenly distributed $ar{x}$

$$\bar{x}_i = \begin{cases} \frac{1}{\sqrt{s^*}} & \text{if } 1 \leq i \leq s^* \\ 0 & \text{if } i > s^* \end{cases} \implies \mathrm{supp}(\bar{\mathbf{x}}) = \{1, 2, \dots, s^*\}$$

- Construct $\mathbf{M}^{(\epsilon)} \in \mathbb{R}^{n \times d}$ parameterized by ϵ
 - $\bullet \ \mathbf{M}_{1:s^*}^{(\epsilon)} \text{ are random } s^* \text{ orthogonal column vectors s.t. } \left\| \mathbf{M}_i^{(\epsilon)} \right\|_2^2 = n \ \forall \ i \in [s^*].$
 - $\mathbf{M}_i^{(\epsilon)} = \sqrt{1-\epsilon} \left[\frac{1}{\sqrt{s^*}} \sum_{j=1}^{s^*} \mathbf{M}_j^{(\epsilon)} \right] + \sqrt{\epsilon} \mathbf{g}_i \ \forall \ i \not\in [s^*] \ \text{where } \mathbf{g}_i \text{'s are}$ orthogonal to each other and $\mathbf{M}_{1:s^*}^{(\epsilon)}$ with $\|\mathbf{g}_i\|_2^2 = n$.

- (Y. Zhang'15)'s lower bounds were for algorithms that output s^* -sparse solutions which does not apply for OMP when it is run for more than s^* iterations.
- We provide matching lower bounds for Support recovery as well as generalization error for OMP.
- \bullet Fool OMP into picking incorrect indexes. Large support size \implies large generalization error.
- Construct an evenly distributed \bar{x}

$$\bar{x}_i = \begin{cases} \frac{1}{\sqrt{s^*}} & \text{if } 1 \leq i \leq s^* \\ 0 & \text{if } i > s^* \end{cases} \implies \mathrm{supp}(\bar{\mathbf{x}}) = \{1, 2, \dots, s^*\}$$

- Construct $\mathbf{M}^{(\epsilon)} \in \mathbb{R}^{n \times d}$ parameterized by ϵ
 - $\bullet \ \mathbf{M}_{1:s^*}^{(\epsilon)} \text{ are random } s^* \text{ orthogonal column vectors s.t. } \left\| \mathbf{M}_i^{(\epsilon)} \right\|_2^2 = n \ \forall \ i \in [s^*].$
 - $\mathbf{M}_i^{(\epsilon)} = \sqrt{1-\epsilon} \left[\frac{1}{\sqrt{s^*}} \sum_{j=1}^{s^*} \mathbf{M}_j^{(\epsilon)} \right] + \sqrt{\epsilon} \mathbf{g}_i \ \forall \ i \not\in [s^*] \ \text{where } \mathbf{g}_i \text{'s are}$ orthogonal to each other and $\mathbf{M}_{1:s^*}^{(\epsilon)}$ with $\|\mathbf{g}_i\|_2^2 = n$.

- (Y. Zhang'15)'s lower bounds were for algorithms that output s^* -sparse solutions which does not apply for OMP when it is run for more than s^* iterations.
- We provide matching lower bounds for Support recovery as well as generalization error for OMP.
- Construct an evenly distributed x̄

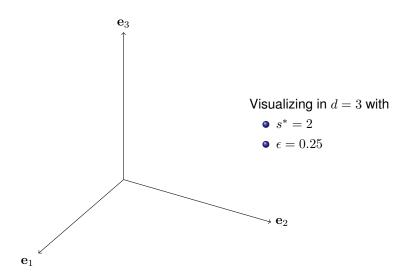
$$\bar{x}_i = \begin{cases} \frac{1}{\sqrt{s^*}} & \text{if } 1 \leq i \leq s^* \\ 0 & \text{if } i > s^* \end{cases} \implies \text{supp}(\bar{\mathbf{x}}) = \{1, 2, \dots, s^*\}$$

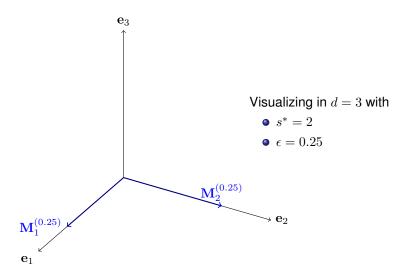
- Construct $\mathbf{M}^{(\epsilon)} \in \mathbb{R}^{n \times d}$ parameterized by ϵ
 - $\bullet \ \mathbf{M}_{1:s^*}^{(\epsilon)} \text{ are random } s^* \text{ orthogonal column vectors s.t. } \left\| \mathbf{M}_i^{(\epsilon)} \right\|_2^2 = n \ \forall \ i \in [s^*].$
 - $\mathbf{M}_i^{(\epsilon)} = \sqrt{1-\epsilon} \left[\frac{1}{\sqrt{s^*}} \sum_{j=1}^{s^*} \mathbf{M}_j^{(\epsilon)} \right] + \sqrt{\epsilon} \mathbf{g}_i \ \forall \ i \not\in [s^*] \ \text{where } \mathbf{g}_i \text{'s are}$ orthogonal to each other and $\mathbf{M}_{1:s^*}^{(\epsilon)}$ with $\|\mathbf{g}_i\|_2^2 = n$.

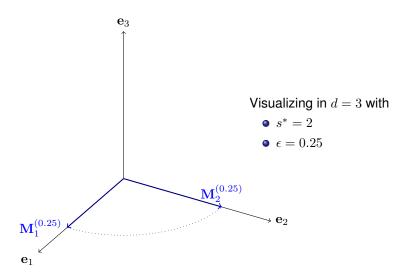
- (Y. Zhang'15)'s lower bounds were for algorithms that output s^* -sparse solutions which does not apply for OMP when it is run for more than s^* iterations.
- We provide matching lower bounds for Support recovery as well as generalization error for OMP.
- Construct an evenly distributed x̄

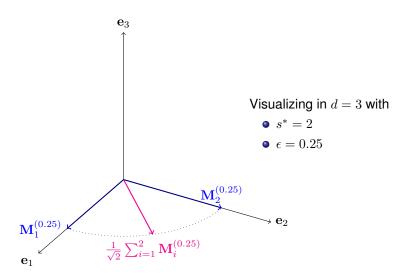
$$\bar{x}_i = \begin{cases} \frac{1}{\sqrt{s^*}} & \text{if } 1 \leq i \leq s^* \\ 0 & \text{if } i > s^* \end{cases} \implies \text{supp}(\bar{\mathbf{x}}) = \{1, 2, \dots, s^*\}$$

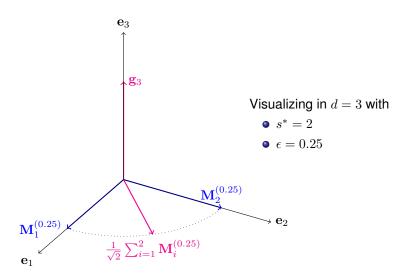
- ullet Construct $\mathbf{M}^{(\epsilon)} \in \mathbb{R}^{n \times d}$ parameterized by ϵ
 - $\bullet \ \ \mathbf{M}_{1:s*}^{(\epsilon)} \ \text{ are random } s^* \text{ orthogonal column vectors s.t. } \left\| \mathbf{M}_i^{(\epsilon)} \right\|_2^2 = n \ \forall \ i \in [s^*].$
 - $\mathbf{M}_i^{(\epsilon)} = \sqrt{1 \epsilon} \left[\frac{1}{\sqrt{s^*}} \sum_{j=1}^{s^*} \mathbf{M}_j^{(\epsilon)} \right] + \sqrt{\epsilon} \mathbf{g}_i \ \forall \ i \not\in [s^*] \ \text{where } \mathbf{g}_i \text{'s are orthogonal to each other and } \mathbf{M}_{1:s^*}^{(\epsilon)} \ \text{with } \|\mathbf{g}_i\|_2^2 = n.$

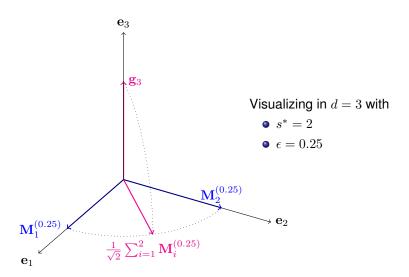


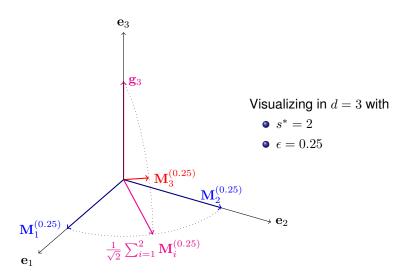


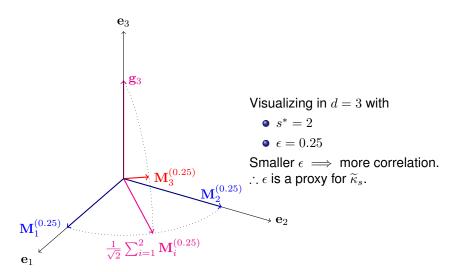














Lower bounds

Noiseless case

For $s^* \le d \le n$, $\exists \ \epsilon > 0$ s.t. when OMP is executed on SLR problem with $y = \mathbf{M}^{(\epsilon)} \mathbf{\bar{x}}$ for $s \le d - s^*$ iterations

- ullet $\widetilde{\kappa}_s(\mathbf{M}^{(\epsilon)})\lesssim rac{s}{s^*}$ and $\gamma\leq \sqrt{rac{2}{3}}$
- $\bullet \ \mathbf{S}^* \cap \mathsf{supp}(\mathbf{\hat{x}}_s) = \phi$

Noisy case

For $s^* \leq s \leq d^{1-\alpha}$ where $\alpha \in (0,1)$, $\exists \ \epsilon > 0$ s.t. when OMP is executed on SLR problem with $y = \mathbf{M}^{(\epsilon)} \bar{\mathbf{x}} + \boldsymbol{\eta}$ where $\boldsymbol{\eta} \sim \mathcal{N}\left(\mathbf{0}, \sigma^2 \mathbf{I}_{n \times n}\right)$, then

- ullet $\widetilde{\kappa}_s(\mathbf{M}^{(\epsilon)})\lesssim rac{s}{s^*}$ and $\gamma\leq rac{1}{2}$
- with high probability $\frac{1}{n} \| \mathbf{A} \hat{\mathbf{x}}_s \mathbf{A} \bar{\mathbf{x}} \|_2^2 \gtrsim \frac{1}{n} \sigma^2 \widetilde{\kappa}_{s+s^*} s^*$
- with high probability $S^* \cap \operatorname{supp}(\hat{\mathbf{x}}_s) = \phi$

 $\implies s \gtrsim \widetilde{\kappa}_s s^*$ iterations are indeed necessary. Addition of noise can only help.



Lower bounds

Noiseless case

For $s^* \le d \le n$, $\exists \ \epsilon > 0$ s.t. when OMP is executed on SLR problem with $y = \mathbf{M}^{(\epsilon)} \bar{\mathbf{x}}$ for $s \le d - s^*$ iterations

- ullet $\widetilde{\kappa}_s(\mathbf{M}^{(\epsilon)})\lesssim rac{s}{s^*}$ and $\gamma\leq \sqrt{rac{2}{3}}$
- $\bullet \ \mathbf{S}^* \cap \mathsf{supp}(\mathbf{\hat{x}}_s) = \phi$

Noisy case

For $s^* \leq s \leq d^{1-\alpha}$ where $\alpha \in (0,1)$, $\exists \ \epsilon > 0$ s.t. when OMP is executed on SLR problem with $y = \mathbf{M}^{(\epsilon)} \bar{\mathbf{x}} + \boldsymbol{\eta}$ where $\boldsymbol{\eta} \sim \mathcal{N}\left(\mathbf{0}, \sigma^2 \mathbf{I}_{n \times n}\right)$, then

- ullet $\widetilde{\kappa}_s(\mathbf{M}^{(\epsilon)})\lesssim rac{s}{s^*}$ and $\gamma\leq rac{1}{2}$
- with high probability $\frac{1}{n} \| \mathbf{A} \hat{\mathbf{x}}_s \mathbf{A} \bar{\mathbf{x}} \|_2^2 \gtrsim \frac{1}{n} \sigma^2 \widetilde{\kappa}_{s+s^*} s^*$
- with high probability $S^* \cap \text{supp}(\hat{\mathbf{x}}_s) = \phi$

 $\implies s \gtrsim \widetilde{\kappa}_s s^*$ iterations are indeed necessary. Addition of noise can only help.



Lower bounds

Noiseless case

For $s^* \le d \le n$, $\exists \ \epsilon > 0$ s.t. when OMP is executed on SLR problem with $y = \mathbf{M}^{(\epsilon)} \bar{\mathbf{x}}$ for $s \le d - s^*$ iterations

- ullet $\widetilde{\kappa}_s(\mathbf{M}^{(\epsilon)})\lesssim rac{s}{s^*}$ and $\gamma\leq \sqrt{rac{2}{3}}$
- $\bullet \ \mathbf{S}^* \cap \mathsf{supp}(\mathbf{\hat{x}}_s) = \phi$

Noisy case

For $s^* \leq s \leq d^{1-\alpha}$ where $\alpha \in (0,1)$, $\exists \ \epsilon > 0$ s.t. when OMP is executed on SLR problem with $y = \mathbf{M}^{(\epsilon)} \bar{\mathbf{x}} + \boldsymbol{\eta}$ where $\boldsymbol{\eta} \sim \mathcal{N}\left(\mathbf{0}, \sigma^2 \mathbf{I}_{n \times n}\right)$, then

- ullet $\widetilde{\kappa}_s(\mathbf{M}^{(\epsilon)})\lesssim rac{s}{s^*}$ and $\gamma\leq rac{1}{2}$
- with high probability $\frac{1}{n} \| \mathbf{A} \hat{\mathbf{x}}_s \mathbf{A} \bar{\mathbf{x}} \|_2^2 \gtrsim \frac{1}{n} \sigma^2 \widetilde{\kappa}_{s+s^*} s^*$
- ullet with high probability $\mathbf{S}^* \cap \operatorname{supp}(\mathbf{\hat{x}}_s) = \phi$

 $\implies s \gtrsim \widetilde{\kappa}_s s^*$ iterations are indeed necessary. Addition of noise can only help.

<ロ > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < 回 る の ○ □ < □ る の ○ □ < □ る の ○ □ < □ る の ○ □ < □ る の ○ □ < □ る の ○ □ る の ○ □ < □ る の ○ □ < □ る の ○ □ < □ る の ○ □ < □ る の ○ □ < □ る の ○ □ < □ る の ○ □ < □ る の ○ □ < □ る の ○ □ < □ る の ○ □ < □ る の ○ □ < □ る の ○ □ < □ る の ○ □ < □ る の ○ □ < □ る の ○ □ る の ○ □ < □ る の ○ □ < □ る の ○ □ < □ る の ○ □ < □ る の ○ □ < □ る の ○ □ < □ る の ○ □

Simulations

We perform simulations on the lower bound instance class.

 $\mathbf{M}^{(\epsilon)} \in \mathbb{R}^{1000 \times 100}$ and $s^* = 10$.

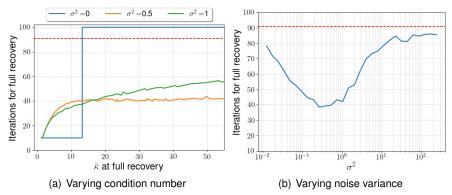


Figure: Number of iterations required for recovering the full support of $\bar{\mathbf{x}}$ with respect to the restricted condition number $(\tilde{\kappa}_{s+s^*})$ of design matrix and the variance of noise (σ^2) .