

Aim:- Implement Linear search to find an item in the list.

Theory:-

### Linear Search

Linear search is one of the simplest searching algorithm in which targeted item is sequentially matched with each item in the list.

It is worst searching algorithm with worst case time complexity. It is a linear approach. On the other hand in case of an ordered list, instead of searching the list in sequence, A binary search is used which will start by examining the middle term.

Linear search is a technique to compare each and every element with the key element to be found, if both of them matches, the algorithm returns that element found and its position is also found.

### Unsorted array :-

Algorithm :-  
 Step 1: Create an empty set and assign it to a variable.

Step 2: Accept the total no. of elements to be inserted into the list from the user. say 'n'

Defn: we for loop for adding the elements into the list. print the new list.

Step 3: Accept an element from the user that to be searched in the list.

Step 4: we for loop in a range from '0' to the total no. of elements to search the element from the list.

Step 5: we if loop that the element in the list is equal to the element accepted from user.

Step 6: If the element is found then print the statement that the element is found along with the element's position.

### Unsorted array :-

```
def linear(arry, x):
    for i in range(len(arry)):
        if arry[i] == x:
            return i
```

inp = input("Enter elements in array : ").split()

arry = [int(i) for i in inp]
print("Enter Element in array noo : ", arry)

x1 = int(input("Enter the statements to be searched : "))
x2 = linear(arry, x1)

if x2 == x1:
 print("Element found at location", x2)
else:
 print("Element not found")

### Output :-

Enter elements in array 3 2 4 5 1  
 Element in array noo : [3 2 4 5 1]

Element to be searched 4

Element found at location 2

so if we enter it loop to point that the element is not found. if the element which is accepted from user is not found in the list.

Step 9: Draw the output of given algorithm.

### Sorted array:-

Q) Sorted linear search:-  
 Sorting means to arrange the elements in increasing or decreasing order.

### Algorithm:-

Step 1: Create empty list and assign it to a variable.

Step 2: Accept total no. of elements to be inserted into the list from user, say 'n'.

Step 3: use for loop for using append() method to add the element in the list.

Step 4: use sort() method to sort the accepted element and assign in increasing order the list then print the list.

Step 5: use if statement to give the range in which element is found in given range then display "Element not found".

Step 6: Then use else statement, if element is not found in range then satisfy the given condition.

```
def linear_search(x):
    for i in range (len(arr)):
        if arr[i] == x:
            return i
    return -1

inp = input("Enter elements in array :").split()
array = []
for ind in inp:
    array.append(int(ind))

print("Elements in array are : ", array)
array.sort()

x1 = int(input("Enter a element to be searched :"))
x2 = linear_search(x1)

if x2 == -1:
    print ("Element found at position ", x2)
else:
    print ("Element not found")
```

### Output:-

Enter elements in array: 1 2 3 4 5

Enter elements in array as : [1, 2, 3, 4, 5]

Enter element to be searched: 2

The element is found at position 1.

step:- we use for loop to range from 0 to the total no. of elements to be searched before accepting this accept or search no from user using input statement.

step:- we use if loop that the element is in the list is equal to the element accepted from user.

step:- If the element is found then print the statement that the element is found along with the element position.

step:- we another if loop to print that the element is not found if the element which is accepted from user is not there in the list.

step:- Atlast the input and output of above algorithm.

## Practical :-

Q8

Aim:- Implement binary search to find a search no in the list

```
Coding :-  
a = [ ]  
n = int(input("Enter element:"))  
for b in range(0,n):  
    b = int(input("Enter a element:"))  
    a.append(b)
```

### Theory :- Binary Search

Binary Search is also known as half interval search, logarithmic search or binary chop is a search algorithm that finds the position of a target value within a sorted array. If you are looking for the number which is at the end of the list then you need to search entire list in linear search, which is time consuming. This can be avoided by using binary search.

### Algorithm:-

- Step 1:- Create empty list and assign it to a variable
  - Step 2:- Using input method, accept the range of given list
  - Step 3:- Use for loop, add elements in list using append() method
- Step 4:- We sort() method to sort the accepted element and assign it in increasing ordered list print the list after sorting.

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```
else:  
    if (s < a[m]):  
        i = m + 1  
    else:  
        f = m + 1  
  
else:  
    if (s < a[m]):  
        i = m + 1  
    else:  
        f = m + 1  
  
print("Element not found")  
Point("Element not found")  
Point("Element is", a)  
for i in range(f, n):  
    m = int((f + i) / 2)  
    Point(m)  
    if (s == a[m]):  
        print("Found element at", m)  
        break
```

steps we if loop to give the range in which element is found in given range then display a message "Element not found".

If element is not found in range then satisfy the below condition.

report accept an argument & say if the element that element has to be searched.

Suppose initialize first to 0 and last to last element of the list. as array is starting from 0. hence it is initialized 1 less than the total count.

Suppose for loop & assign the given range.

~~if point statement in list and still the element to be searched not found then find the middle element (mid).~~

Suppose if the item to be searched is still less than the middle term then

initialize  $last = mid + 1$

Else  
initialize  $first = mid - 1$

Suppose repeat till you found the element. Stick two inputs & output of above algorithm.

## Practical :-

### Coding:-

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Aim:- Implementation of Bubble Sort program on given list.

Theory:- bubble sort is based on the idea of repeatedly comparing pairs of adjacent if they exist in element and then swapping their position if they exist in the wrong order. This is the simplest form of sorting available. In this, we sort the given element in ascending or descending order by comparing two adjacent elements at a time.

### Algorithm:-

Step 1:- Bubble Sort algorithm Starts by comparing the first two elements of an array and swapping if necessary.

Step 2:- If we want to sort the elements of array in descending order then first element is greater than second then; we need to swap the element.

Step 3:- If the first element is smaller than record then we do not swap the the element.

Print ("Enter before sort ", a)

```
n = len(a)
for i in range(0, n):
    for j in range(n-1):
        if a[i] < a[j]:
            temp = a[i]
            a[i] = a[j]
            a[j] = temp
```

Print ("Enter after sorting ", a)

Output:-  
=> Enter the Element : 5 10 11 2 1  
Enter before sort [5 10 11 2 1]  
Enter after sorting [1 2 5 10 11]

Step:- again second and third is smaller than element are compared and swapped if it is necessary and this process go on until last and second last element is compared and swapped.

Step:- If there are n elements to be sorted then the process mentioned above should be repeated n-1 times to get the required result.

Step:- Since the output and input of done algorithm of bubble sort stepwise.

## Practical 4

Point ("Rahul Yadav")

class Stack:

    global tos

    def \_\_init\_\_(self):

        self.l = [0, 0, 0, 0, 0, 0]

        self.tos = -1

    def push(self, data):

        n = len(self.l)

        if self.tos == n - 1:

            print("Stack is full")

        else:

            self.tos = self.tos + 1

            self.l[self.tos] = data

    def pop(self):

        if self.tos < 0:

            print("Stack empty")

        else:

            k = self.l[self.tos]

            print("data =", k)

        self.tos = self.tos - 1

    X = stack()

Step 1:- Define methods push & pop under the class stack.

Step 2:- use if statement to give the condition that if length of given list then print stack is full.

Aim: Implementation of stack using Python list.

Theory:- A Stack is a linear data structure that can be represented in the real world in the form of a physical stack of a the element in the Stack are added or removed only from one position i.e. the topmost position. Thus, the stack works on the LIFO (Last in first out) principle.

Algorithm:-

Step 1:- Create a class Stack with instance variable items.

Step 2:- Define the init method with self argument & initialize the initial value & initialized to an empty list.

Step 3:- Define methods push & pop under the class stack.

Output:-

Rahul yadav

```
>>> S.push ('ra')
>>> S.push ('u')
>>> S.push ('o')
>>> S.push ('v')
```

```
>>> S.push ('c')
>>> S.push ('e')
>>> S.push ('s')
>>> S.push ('t')
```

```
>>> S.push ('g')
>>> S.push ('a')
>>> S.push ('d')
>>> S.push ('f')
```

Stack is full

```
>>> S.pop()
data=80
```

```
>>> S.pop()
data=20
```

```
>>> S.pop()
data=60
```

```
>>> S.pop()
data=50
```

```
>>> S.pop()
data=40
```

```
>>> S.pop()
data=30
```

```
>>> S.pop()
data=20
```

```
>>> S.pop()
data=10
```

Stack empty

Step 8:- Or else print statement as insert the element into the stack and initialize the value.

Step 6:- Push method used to insert the element but pop method used to delete the element from the stack.

Step 9:- If in pop method, value is less than 1 then return the stack is empty or else delete the element from stack at topmost position.

Step 10:- Assign the element values in push method & point the given value in popped out.

Step 11:- First condition check whether the no. of elements are 2000 while the second condition is whether it is designed any value. If it has not assign any value, then they can be seen that stack is empty.

Step 12:- Attach the input & output of above algorithm.

Algorithm

## Practical 15

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```
Print ("Quick Sort")
def partition (arr, low, high):
    i = low - 1
```

```
Pivot = arr[high]
for j in range (low, high):
    if arr [j] <= pivot:
        i = i + 1
        arr[i], arr[j] = arr[j], arr[i]
arr[i+1], arr[high] = arr[high], arr[i+1]
return i + 1
```

Algorithm:-

Step1:- Quick sort first selects a value, which is called pivot value. first value element. Since all our first pivot value since we know that first will eventually end up as part in that list.

Step2:- The partition process will happen next. It will find the split point and at the same time move other items to the appropriate side of the last either less than or greater than pivot. The last either less than or greater than pivot will be partitioned by looking two position marker. Let calls them leftvalue and rightmost at the beginning and end of remaining item in the list.

```
def quicksort (arr, low, high):
    if low < high:
        pi = Partition (arr, low, high)
        quicksort (arr, low, pi-1)
        quicksort (arr, pi+1, high)
```

```
arr = input ("Enter Elements in the list").split()
list1 = []
quicksort (arr, 0, len(arr) - 1)
print ("Elements in the list are : " , list1)
```

Step3:- We began by increasing leftmark until we locate a value that is greater than the Pv, we increase rightmark until we find value that is less than the Pv. At the point where rightmark becomes less than leftmark we stop. The position of rightmark is now the split point.

Output:-

Quick sort

Initial elements in the list are 21, 22, 20, 30, 24, 56

Elements in list are: [21, 22, 20, 30, 24, 56]

Elements after quick sort are [20, 21, 24, 30, 56]

~~Step 9:~~ The p.v can be exchanged with the content of split point and p.v is now in place.

~~Step 10:~~ In addition all the items to left of split point and all less than p.v and all the items to left to the right split point are greater than p.v. The list can now be divided at split point and quick sort can be invoked recursively on the two halves.

~~Step 11:~~ The quick sort function invokes a recursive function quick sort helps.

~~Step 12:~~ Quick sort helps, begins with same base as the merge sort.

~~Step 13:~~ If length of the list is less than 0 or equal one it is already sorted.

~~Step 14:~~ If it is greater than it can be partitioned and recursive function.

~~Step 15:~~ The partition function implement the process described earlier.

~~Step 16:~~ Display and write the coding and output of above algorithm.

## Practical 1

class

Queue:

Aim:- Implementing a Queue using Python list.

Theory:- Queue is a linear data structure which has 2 references front & rear implementing a queue using Python list in the simplest as the Python list provides inbuilt function.

Queue () : creates a new empty Queue.

Enqueue ( ) : Insert an element at the rear of the queue & similar to that of insertion of linked using tail

Dequeue ( ) : Return the element which uses of the front, the front is moved to the successive element. A dequeue operation cannot remove element of the queue is empty.

Algorithm:-

Step 1:- Defines a class Queue & assign global variables then defines init() method with self argument in init() assign or initialize the initial value with the help of self argument.

Step 2:- Define a empty list & define enqueue() with 2 argument assign the length of empty list.

q = Queue()

```

class Queue:
    global f
    global r
    def __init__(self):
        self.t = 0
        self.t = 0
        self.l = [0, 0, 0]
    def enqueue(self, data):
        n = len(self.l)
        if self.r == n:
            self.l[self.r] = data
            self.r = self.r + 1
        print("Element inserted....", data)
    else:
        print("Queue is full!")
    def dequeue(self):
        n = len(self.l)
        if self.f == n:
            print(self.l[self.f])
            self.f = self.f + 1
        else:
            print("Element deleted....")
            self.f = self.f + 1
    else:
        print("Queue is empty")

```

```

Output:
>>> Q.add(10)
element insert 10
>>> Q.add(20)
element insert 20
>>> Q.add(3)
element insert 3
>>> Q.add(4)
Queue is full!
>>> Q.remove()

```

Step 4 : Define dequeue() with self argument, under this, we if statement that front is empty or else, give that front is at 2000 & using that delete the element from front side & increment it by 1.

~~Step 5 :~~ Now call the dequeue() function & give the element that has to be added in the list after adding & some for deleting & display the list after deleting the element from the list.

Aim:- Evaluation of given string using stack  
i.e Postfix.

Theory:- The postfix expression is free from parentheses further we took care of the priorities of operation in program. A given postfix expression is always from left to right.

### Algorithm:-

- Step 1:- Define evaluate as function then create a empty stack in python.
- Step 2:- convert the string to a list by using the string method split.
- Step 3:- calculate the length of string & print.
- Step 4:- for loop to assign the range of string then given condition using if statement.
- Step 5:- Scan the token list for left to right. If token is opened convert it from a string to integer & push the value onto the p.

```

def evaluate(s):
    k = s.split()
    n = len(k)
    stack = []
    for i in range(n):
        if (k[i].isdigit() == True):
            stack.append(int(k[i]))
        elif k[i] == '+':
            a = stack.pop()
            b = stack.pop()
            stack.append(int(a) + int(b))
        elif k[i] == '-':
            a = stack.pop()
            b = stack.pop()
            stack.append(int(b) - int(a))
        elif k[i] == '*':
            a = stack.pop()
            b = stack.pop()
            stack.append(int(b) * int(a))
        return stack.pop()

s = "8 6 9 * +"
a = evaluate(s)
print("The evaluated value is:", a)

```

Output :-  
The evaluated value is : 62

Step 1:- If the token is an operator it will  
need two operands pop the top twice. The  
first pop is second operand & the second  
pop is the first operand.

Step 2:- Perform the arithmetic operation push  
the result back on the stack.

Step 3:- When the input expression has been completely  
processed the result is on the stack. Pop  
the stack & return the value.

Step 4:- Point the result & string after the  
evaluation of postfix.

Step 5:- Write output & input of above algorithm.

## Practical no:-8

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Aim:- Implementation of single linked list by adding the nodes from the last position.

Theory:- A linked list is a linear data structure which stores the element in a node. The individual element of linked list called a Node. Node consists of 2 parts  
 ① Data ② Next . Data stores all the information about the element. For example roll no, name, address etc whereas next refers to the next node.

Algorithm:-

Step 1:- Transfer storing of a linked list means listing all nodes in linked list in order to perform some operation on them.

Step 2:- The entire linked list can be accessed with first node of the linked list.

Step 3:- Thus, the entire linked list can be traversed using the node which is referred by the head pointer of the linked list.

Step 4:- Now that we know that we can traverse the entire linked list using:

Start = linkedlist()

Point (head.data)

```
class node:
    global data
    global next
    def __init__(self, item):
        self.data = item
        self.next = None

class linkedlist():
    global a
    def __init__(self):
        self.s = None

    def add(self, item):
        newnode = node(item)
        if self.s == None:
            self.s = newnode
        else:
            head = self.s
            while head.next != None:
                head = head.next
            head.next = newnode
            head = head.next

    def addB(self, item):
        newnode = node(item)
        if self.s == None:
            self.s = newnode
        else:
            newnode.next = self.s
            self.s = newnode

    def display(self):
        head = self.s
        while head.next != None:
            print(head.data)
            head = head.next
        print(head.data)
```

Output :-  
 >>> start.add(50)  
 >>> start.add(60)  
 >>> start.add(70)  
 >>> start.add(80)  
 >>> start.add(90)  
 >>> start.add(100)  
 >>> start.add(110)  
 >>> start.add(120)

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Step 10: Similarly, we can traverse rest of nodes in the linked list using same method by while loop.

Step 11: Our concept now is to find terminating condition for the while loop.

Step 12: The last node in the linked list is referred by tail of linked list. Since, the last node of linked list does not have any next node, the value in the next field of last node is null.

Step 13: So we can refer here last node of linked list self = None.

Step 14: - we have to now see how to start traversing the linked list & how to identify whether we have reached the last node or linked list or not.

Step 15: - write the coding & input, output of above algorithm.

Code :-

```
def sort(arr, l, r):
    n1 = m - l + 1
    n2 = r - m
    L = [0] * (n1)
    R = [0] * (n2)
    for i in range(0, n1):
        L[i] = arr[l+i]
    for j in range(0, n2):
        R[j] = arr[m+1+j]
    i = 0
    j = 0
    k = l
    while i < n1 and j < n2:
        if L[i] <= R[j]:
            arr[k] = L[i]
            i = i + 1
        else:
            arr[k] = R[j]
            j = j + 1
        k = k + 1
    while i < n1:
        arr[k] = L[i]
        i = i + 1
        k = k + 1
    while j < n2:
        arr[k] = R[j]
        j = j + 1
        k = k + 1
    mergeSort(arr, l, r)
```

Algorithm :-  
Implementation of mergesort by using python

Theory :- Merge sort is a divide and conquer algorithm. It divides input array into two halves which itself for the two halves & the merge the two sorted halves. The merge function is used for merging two halves.

Algorithm :-

Step 1:- The list is divided into left & right in each recursive call until no adjacent elements are obtained.

Step 2:- Now begins the sorting process. Two i & j iterators transverse the two halves in each call. The k iterators however traverse the whole list & makes changes along the way.

Step 3:- If the value at i is smaller than the value of j ( $L[i]$ ) is assigned to the  $arr[i+1]$ . No & k is incremented if not then  $R[j]$  is chosen.

def mergeSort(arr, l, r):
 if l < r:
 m = int((l+(r-1))/2)
 mergeSort(arr, l, m)
 mergeSort(arr, m+1, r)
 mergeSort(arr, l, m)

 if l < r:
 m = int((l+(r-1))/2)
 mergeSort(arr, l, m)
 mergeSort(arr, m+1, r)
 mergeSort(arr, l, m)

Step 4:- This way, the values being assigned through  $arr[i:j]$  are all sorted.

steps:- At the end of this loop, one of the  
values may not have been transposed  
completely remaining slots in the list.

Steps:- Thus, the range sort has been  
implemented.

Output:-

arr = [12, 23, 34, 56, 78, 45, 86, 98, 42]

point (arr)

n = len (arr)

mergeSort (arr, 0, n-1)

point (arr)

[12, 23, 34, 56, 78, 45, 86, 98, 42]

[12, 23, 34, 56, 78, 45, 86, 98]

Code:-

```

set1 = set()
set2 = set()
for i in range(8, 15):
    set1.add(i)
for i in range(1, 12):
    set2.add(i)
print("set1:", set1)
print("set2:", set2)
print("\n")
set3 = set1 | set2
print("union of set1 and set2; set3", set3)
set4 = set1 & set2
print("intersection of set1 & set2; set4", set4)
print("\n")
if set3 > set4: print("set3 is superset of set4")
elif set3 < set4: print("set3 is subset of set4")
else: print("set2 is same as set4")
if set4 < set3:
    print("set4 is subset of set3")
    print("\n")
set5 = set3 - set4
print("elements in set3 and not in set4; set5")
print("\n")

```

Algorithm:-

Step 1:- Define two empty set as set1 and set2 now use for statement providing the range of above 2 sets.

Step 2:- Now add() method is required for addition the element according to given range then point the sets for addition.

Step 3:- Find the union and intersection of above 2 sets by using |, & method. point the sets of union & intersection of set3.

Step 4:- Use if statement to find out the subset and superset of set3 and set4. Display the above set.

Step 5:- Display that element in set3 is not in set4 using mathematical operation.

Step 6:- use clear() to remove or delete the sets and point the set after clearing the element present in the set.

If set u. is disjoint(sets):

Point("set4 and sets are mutually exclusive")

set4.clear()

Point("After applying clear, set5 is empty set;")

Point("set5:", set5)

Output:-

set1: {8, 9, 10, 11, 12, 13, 14}

set2: {1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 13}

union set4 & set2: set3

{1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13, 14}

Intersection of set1 & set2: {8, 9, 10, 11}

set3 is superset of set4  
elements in set3 and not in set4 : set5

{1, 2, 3, 4, 5, 6, 7, 12, 13, 14}

set4 & set5 are mutually exclusive after  
applying clear, set5 is empty set

set5 = set()

## #Node

class node:

def \_\_init\_\_(self, value):

self.left = None

self.val = value

self.right = None

class BST:

def \_\_init\_\_(self):

self.root = None

def add(self, value):

p = node(value)

if self.root == None:

self.root = p

print("Root is added successfully")

p.val

else:

h = self.root

if p.val &lt; h.val:

if h.left == None:

h.left = p

print(p.val, "None is added successfully")

break;

else:

h = h.left

else:

if h.right == None:

h.right = p

print(p.val, "None is added to right side successfully")

break;

else:

h = h.right

Theory :- Binary tree is a tree which supports maximum of 2 children for any node within the tree. Thus any particular node can have either 0 or 1 or 2 children. There is another identity of binary tree that is ordered such that one child is identified as left child and other as right child.

Inorder : 1) Traverse the left subtree. The left subtree inform right have left and right subtrees.

Preorder : i) visit the root node. traverse the left subtree and right subtree. traverse the right subtree.

Postorder : 1) Traverse the left subtree. The left subtree inform right have left & right subtrees.

Program :- Program based on binary search tree by implementation in-order, pre-order & post-order traversed.

Algorithm:-

Step 1:- Define class node & define init() with 2 arguments. Initialize the value in this method.

Step 2:- Again, define a class BST that is binary search tree with init() with self argument & assign the root is none.

Step 3:- Define add() for adding the node.  
Define a variable p that  
p = node(value)

Step 4:- use if statement for checking the condition the root is none then we else statement for if node is less than the main node then put on orange than in left side.

Step 5:- use while loop for checking the node is less than or greater than the main node & break the loop if it is not satisfying.

```

def inorder (root):
    if root==None:
        return
    else:
        inoder (root.left)
        print (root.value)
        inoder (root.right)

def preoder (root):
    if root==None:
        return
    else:
        print (root.value)
        preoder (root.left)
        preoder (root.right)

def postoder (root):
    if root==None:
        return
    else:
        postoder (root.left)
        postoder (root.right)
        print (root.value)

t= BST()
#Output:-
'''t.add(1)
root is added successfully
'''t.add(2)
node is added to rightside successfully

```

```

>>> t.add(4)
4 node is added to right side successfully
>>> t.add(5)
5 node is added to right side successfully
>>> t.add(3)
3 node is added to right left side successfully
>>> print("In gnode:", gnode(t.root))
gnode:

```

```

1
2
3
4
5

```

gnode: None

```

>>> print("In Preorder:", Preorder(gnode,t.root))
Preorder:

```

```

1
2
3
4
5

```

Preorder: None

```

>>> print("In Postorder:", Postorder(gnode,t.root))
Postorder:

```

```

3
4
5
1
2

```

Postorder: None

Step 9:- After this, left side tree & right subtree repeat this method to arrange the node according to binary search tree.

Step 8:- Define Inorder() Preorder() & Postorder() with root as argument & use it statement that root is none & return that all.

Step 9:- Inorder, else statement used for giving that condition if first left, root & then right node.

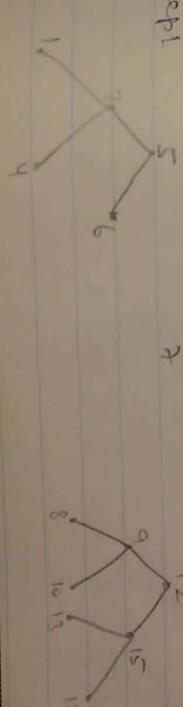
Step 10:- For Preorder, we have to give condition in else that first root, left & the right node.

Step 11:- for postorder in else part assign left & right & root.

Step 11:- Display the output & input.

- INORDER (LVR)

Step 1



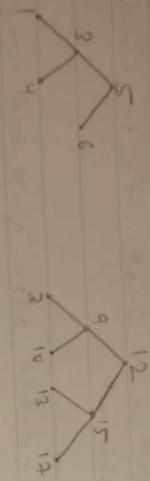
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Step 2: 3 5 6 7 9 12 15

Step 3: 1 3 4 5 6 7 8 9 10 11 12 13 15 17

### • PREORDER (NLR)

Step 1:

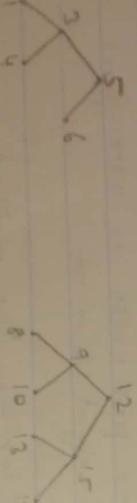


Step 2: 7 5 3 6 12 9 15

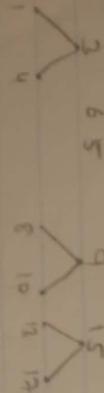
Step 3: 7 5 3 1 4 6 12 9 8 16 15 13 17

### • POSTORDER (LRN)

Step 1:

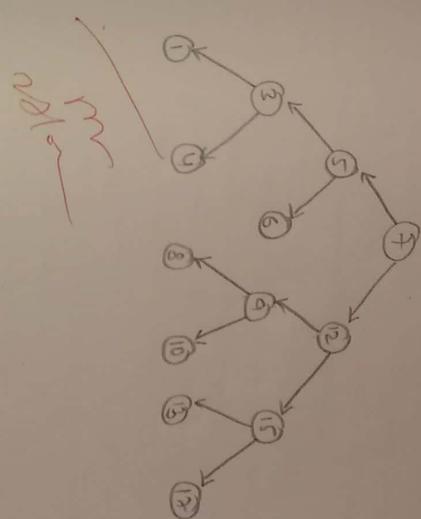


Step 2: 3 6 5



Step 3: 1 4 3 6 5 8 10 9 13 14 15 7

\* Binary search tree:



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