

UNIVERSITY OF KWAZULU-NATAL (HOWARD
COLLEGE)

MASTERS THESIS

Graph Cut Optimisation

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TAPAMO

*A thesis submitted in fulfillment of the requirements
for the degree of Master of Science in Engineering*

in the

Computer Vision Research Group
Department of Electrical, Electronic and Computer Engineering

April 20, 2016

Declaration of Authorship

I, Ryan NAIDOO, declare that this thesis titled, “Graph Cut Optimisation” and the work presented in it are my own. I confirm that:

- This work was done wholly or mainly while in candidature for a research degree at this University.
- Where any part of this thesis has previously been submitted for a degree or any other qualification at this University or any other institution, this has been clearly stated.
- Where I have consulted the published work of others, this is always clearly attributed.
- Where I have quoted from the work of others, the source is always given. With the exception of such quotations, this thesis is entirely my own work.
- I have acknowledged all main sources of help.
- Where the thesis is based on work done by myself jointly with others, I have made clear exactly what was done by others and what I have contributed myself.

Signed:

Date:

“Thanks to my solid academic training, today I can write hundreds of words on virtually any topic without possessing a shred of information, which is how I got a good job in journalism.”

Dave Barry

UNIVERSITY OF KWAZULU-NATAL (HOWARD COLLEGE)

Abstract

Faculty of Engineering
Department of Electrical, Electronic and Computer Engineering

Master of Science in Engineering

Graph Cut Optimisation

by Ryan NAIDOO

The Thesis Abstract is written here (and usually kept to just this page).
The page is kept centered vertically so can expand into the blank space
above the title too...

Acknowledgements

The acknowledgments and the people to thank go here, don't forget to include your project advisor...

Contents

Declaration of Authorship	iii
Abstract	vii
Acknowledgements	ix
1 Introduction	1
1.1 What is Image Segmentation	1
1.2 Energy Minimisation	1
1.3 Labelling Problems	1
1.4 Labelling Problems as Energy Minimisation	2
1.5 Energy Minimisation Algorithms and Special Cases	2
1.6 Following Chapter Outlines	2
1.7 Other things to talk about	2
1.8 Other things to talk about 2	2
1.9 In Closing	2
2 Mathematical Background	5
2.1 Markov Random Fields	5
2.1.1 Markov Random Fields Theory and Concepts	5
2.1.2 Markov Random Fields in Image Modelling	5
2.1.3 MAP-MRF Approxiamtion via Graph Cuts	5
2.2 Graph Theory and Flow Networks	5
2.3 Max-Flow/Min-Cut Problem	10
2.4 Modelling Images as Graphs	11
2.4.1 Sub-modularity Conditions for Discrete Systems	11
2.4.2 Connectivity	11
2.4.3 Distance Metrics	12
Euclidean Distance	12
Riemannian Distance	12
Learned Distance from Seeds	12
2.5 Max-Flow/Min-Cut Algorithms	12
2.5.1 Ford-Fulkerson	13
2.5.2 Dinic/Edmond-Karp	13
2.5.3 Push-Relabel	13
Push-Relabel Speed Optimisation Heuristics	16
2.5.4 Alpha-Beta Swap	18
2.5.5 Alpha-Expansion	18
3 Interactive Segmentation on Single Channel Data	19
3.1 Seeding	19
3.2 Estimating Probability Distributions	19
3.2.1 Expectation Maximisation (Guassian Mixture Modelling)	19

Fixed-Distribution Modelling	19
Automatically Determining the Optimal Number of Mixtures	20
3.2.2 Naive Bayesian Classification	20
3.2.3 Supervised Learning for Multi-Layered Perceptron	20
4 Interactive Segmentation Extention to Multi-Channel Data	21
5 Automatic Graph Cut Segmentation	23
5.1 Determining FG and BG seeds/probability distributions	23
5.2 Determining Optimal Paramters Settings	23
6 Graph Cut Solution to Chan-Vese Segmentation	25
7 Pre-Processing and Post-Processing Techniques	27
7.1 Removal of Artifacts using Connected Components	27
7.2 Anisotropic Diffusion	27
7.2.1 Coherence Enhancing Diffusion	27
7.2.2 Coherence Enhancing Diffusion with Optimised Ro- tational Invariance	27
7.3 Poisson Denoising	28
7.3.1 Total-Variation Denoising	28
7.4 Contrast Enhancement	28
8 Case-Study: Flourescence Microscopy Cell Segmentation	29
8.1 Removal of Artifacts using Connected Components	29
8.2 Anisotropic Diffusion	29
8.2.1 Coherence Enhancing Diffusion	29
8.2.2 Coherence Enhancing Diffusion with Optimised Ro- tational Invariance	29
8.3 Poisson Denoising	30
8.3.1 Total-Variation Denoising	30
8.4 Contrast Enhancement	30
A Proof of Maximum Flow - Minimum Cut Equivalence	31
B Push-Relabel Algorithm Analysis and Proof of Correctness	33

List of Figures

2.1	Undirected weighted graph G . The degree of each node is shown next to the corresponding node. The graph is simple. The red box shows the vertex set, V_G , and edge set, E_G , and their corresponding norm.	6
2.2	Directed weighted graph (Digraph) D . The in-degree and out-degree is shown next to each node. The graph is simple and not balanced. The red box shows the vertex set, V_D , and edge set, E_D , and their corresponding norm.	7
2.3	Undirected weighted graph H is a subgraph of G in Figure XX, $H \subseteq G$. Directed weighted graph I is a subgraph of D in Figure XX, $I \subseteq D$. The degree of each node is shown next to the corresponding node. The red box shows the vertex set, the edge set and their corresponding norms.	7
2.4	Network N with no flow. The in-degree and out-degree for the source, s , and the sink, t , are shown next to the corresponding node.	8
2.5	Network N with flow. The flow out of the source node, s , is equal to the flow into the sink node, t . For all other nodes, the flow-in is equal to the flow-out. This is the conservation of flow principle. This is only part of the network. The remaining part is the residual graph which shows the amount of reverse flow is available on an edge.	8
2.6	Network N with with a valid cut C . The nodes within the red region are reachable from the source and the nodes within the blue region are able to reach the sink. The cut set, C , is show in the orange filled block.	9
2.7	Network N with with a invalid cut C . The cut does not partition source node s and sink node t into distinct sets.	9
2.8	Network N with with a invalid cut C . The cut partition partitions the graph into more than two sets and the cut intersects the edges <i>ab</i> and <i>ad</i> twice.	10
2.9	Network N with maximum flow. There is no way to push more flow out of the source into the sink without breaking the rules for conservation of flow.	10
2.10	Network N with minimal cut C . The sum of the capacity of all the edges in the cut set is the minimum of all possible valid cuts on the network N	11

List of Tables

List of Algorithms

1	Euclid's algorithm	13
2	Push Operation	15
3	Relabel Operation	16
4	Push-Relabel Main-loop	17

List of Abbreviations

LAH List Abbreviations Here
WSF What (it) Stands For

Physical Constants

Speed of Light $c_0 = 2.997\,924\,58 \times 10^8 \text{ m s}^{-1}$ (exact)

List of Symbols

a	distance	m
P	power	W (J s^{-1})
ω	angular frequency	rad

For/Dedicated to/To my...

Chapter 1

Introduction

1.1 What is Image Segmentation

Welcome to this L^AT_EX Thesis Template, a beautiful and easy to use template for writing a thesis using the L^AT_EX typesetting system.

If you are writing a thesis (or will be in the future) and its subject is technical or mathematical (though it doesn't have to be), then creating it in L^AT_EX is highly recommended as a way to make sure you can just get down to the essential writing without having to worry over formatting or wasting time arguing with your word processor.

L^AT_EX is easily able to professionally typeset documents that run to hundreds or thousands of pages long. With simple mark-up commands, it automatically sets out the table of contents, margins, page headers and footers and keeps the formatting consistent and beautiful. One of its main strengths is the way it can easily typeset mathematics, even *heavy* mathematics. Even if those equations are the most horribly twisted and most difficult mathematical problems that can only be solved on a super-computer, you can at least count on L^AT_EX to make them look stunning.

1.2 Energy Minimisation

Brief energy minimisation and what it is. Huge number of problems in vision are inference problems which can be found from energy minimisation. The relation of the concept of energy in physics is not important but useful.

Vision problems are much more complex involving hundreds or even millions of interdependant variables. Some energies precisely model the desired inference problem while some are coarse approximations. Some energies are easy to optimise while others are known to be NP-hard. Once an accurate energy and satisfying algorithm are available, the associated inference problem is essentially solved.

Many of the most important developments in computer vision began with a proposal for a better energy, better algorithm or a combination of both.

1.3 Labelling Problems

What is a labelling problem? Type of labelling problems. What is needed for a labelling problem. What is a discrete labelling problem in vision? What are the type of labels (semantic or related to geometry).

Data driven criteria - Data influences the outcome. These preferences are derived from machine learning.

Regularisation criteria - When we explicitly prefer some kinds of labelling over others, these criteria are called regularisers. The most prominent in computer vision are spatially coherent labellings. The idea of smoothness in preference to noisy. We know from experience that object and medical data corresponds to coherent labels more often than not - varies from application to application but has now become a rule of thumb. Because data in computer vision tends to be highly correlated in space.

1.4 Labelling Problems as Energy Minimisation

General case of expressing data-driven and regularisation criteria as concrete energy decisions.

smoothness, neighbourhood, joint probability
Markov Random Field, MAP-MRF problem

1.5 Energy Minimisation Algorithms and Special Cases

Dynamic Programming
Binary Energies with Coherence

1.6 Following Chapter Outlines

The remainder of the thesis outline.

1.7 Other things to talk about

Gestalt groupings
Properties of a good segmentation and segmentation algorithm

1.8 Other things to talk about 2

Goal of image segmentation: to cluster pixels into salient image regions i.e. regions corresponding to individual surfaces, objects or natural part of objects.

Difference between bottom-up and top-down image segmentation.

1.9 In Closing

You have reached the end of this mini-guide. You can now rename or overwrite this pdf file and begin writing your own **Chapter1.tex** and the rest of your thesis. The easy work of setting up the structure and framework has been taken care of for you. It's now your job to fill it out!

Good luck and have lots of fun!

Guide written by —
Sunil Patel: www.sunilpatel.co.uk
Vel: LaTeXTemplates.com

Chapter 2

Mathematical Background

2.1 Markov Random Fields

In this section we review MRF's as a pure mathematical/statistical tool used specifically for vision. That is, we only cover the necessary concepts related to understanding the problem of modelling images for analysis and inference purposes. In sub-section 1 we cover the basics, in sub-section 2 we cover how to model an image using MRF, and in sub-section 3 we cover how to find the MAP, or a close enough approximation, to the MRF.

2.1.1 Markov Random Fields Theory and Concepts

MRF theory and concepts. Purely mathematical/statistical.

2.1.2 Markov Random Fields in Image Modelling

Modelling the joint probability of of image using MRFs.

2.1.3 MAP-MRF Approxiamtion via Graph Cuts

How to make MAP estimates on the MRF. How does the graph-cut approach ensure an MAP solution.

2.2 Graph Theory and Flow Networks

In this section we cover Graph Theory and specifically Flow Networks, which is a branch of Graph Theory, which is fundamental to the understanding of image segmentation via graph cuts. With it roots in Germany where Euler tried to find the solution to the Konigsberg bridge problem, graph theory has since blossomed into a rich field of Mathematics with seemingly endless amounts of application. Graph Theory is a huge topic in mathematics and can be applied to many other sciences. Graph theory is part of another more encompassing field of Mathematics known as Combinatronics. Graph theory and applications are more useful than the average person would recognise. They're used in Google Maps to find shortest routes to destinations, in Molecular Chemistry to model the structure of atoms, and the list goes on for quite a while. It is no surprise that it is also found to be useful in image segmentation.

Graph A graph G is a pair (V, E) , where V is the set of nodes/vertices and E is the set of edges consisting of pairs (u, v) where $u, v \in V$. The graph is assumed to be finite i.e. $|V| = n$ and $|E| = m$.

In an **undirected graph**, the edge (u, v) and (v, u) are not distinct. That is, they refer to the same edge. However, in a **directed graph**, the two edge are now distinct. In a directed graph with edge (u, v) , u is known as the **tail** and v is known as the **head**. In directed graphs, edges, also known as arcs, are depicted by placing arrowheads at the head of the edge. Given an edge $e = (u, v)$, u and v are said to be **incident** on e . A graph is said to be **simple** if it does not contain any self-loops. A **self-loop** is an edge with its end points being the same vertex.

Degree The degree of a vertex v is the number of edges incident on it. $\deg(v) = |\{(u, v), (v, u) \in E\}|$. A self-loop counts for 2.

If a graph is directed, also known as a **digraph**, then a node v has an **in-degree** $d_{in}(v)$ and an **out-degree** $d_{out}(v)$. A digraph is said to be **balanced** if $d_{in}(v) = d_{out}(v), \forall v \in V$.

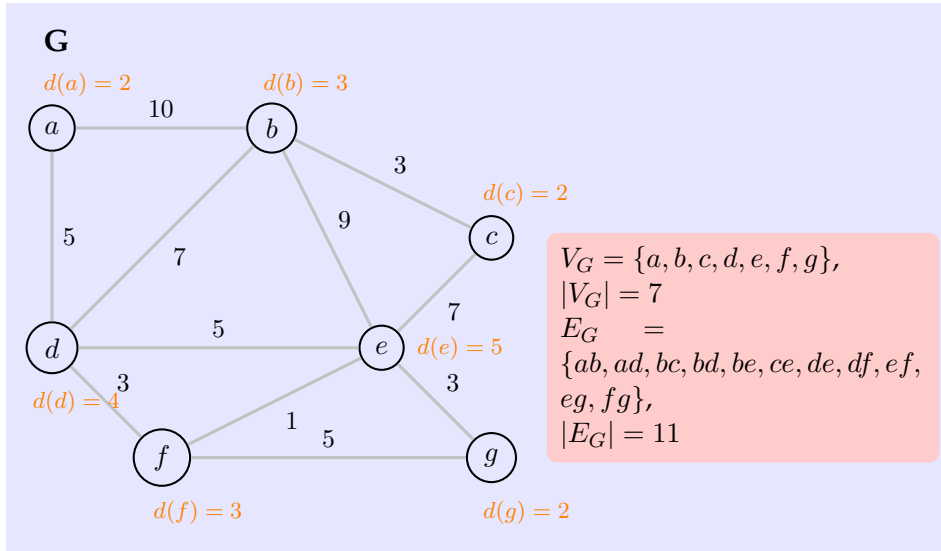


FIGURE 2.1: Undirected weighted graph G . The degree of each node is shown next to the corresponding node. The graph is simple. The red box shows the vertex set, V_G , and edge set, E_G , and their corresponding norm.

Subgraph A graph $G' = (V', E')$ is said to be a sub-graph of $G = (V, E)$, denoted as $G' \subseteq G$, if $V' \subseteq V$ and $E' \subseteq E$.

Clique A clique is a maximal subgraph.

Network A network $N = (V, E)$ is a directed graph with a source node s , a sink node t and a strictly positive capacity on every edge. That is, for each edge $e \in E$, the capacity, $c(\cdot)$, obeys $c(e) \in \mathbb{R}^+$.

The **source node** only has out-going edges, $d_{in}(s) = 0$ and $d_{out}(s) \geq 0$. The **sink node** only has incoming edges, $d_{in} \geq 0$ and $d_{out} = 0$.

Flow A flow $f : V^2 \rightarrow \mathbb{R}^+$ is associated with each edge $e = (u, v)$ such that:

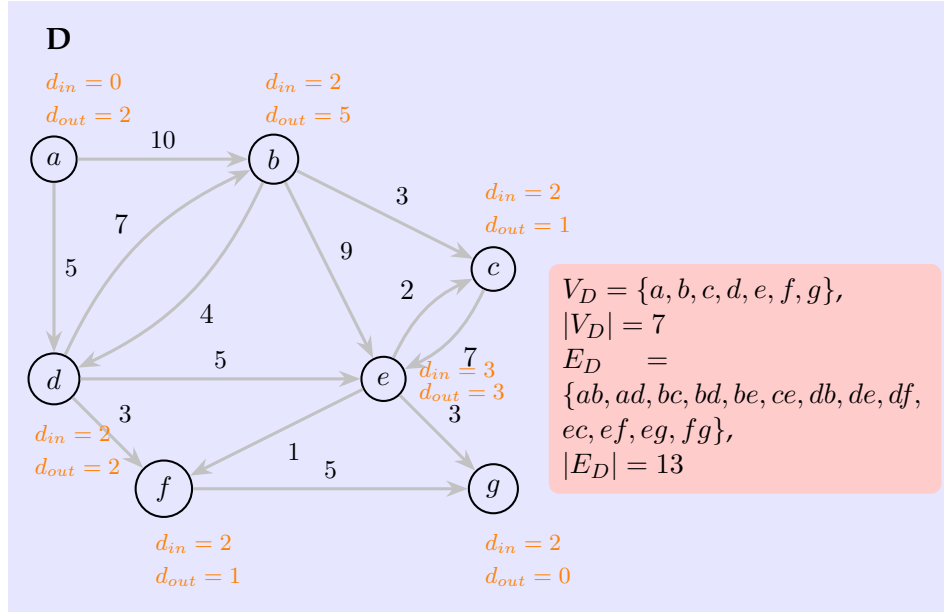


FIGURE 2.2: Directed weighted graph (Digraph) **D**. The in-degree and out-degree is shown next to each node. The graph is simple and not balanced. The red box shows the vertex set, V_D , and edge set, E_D , and their corresponding norm.

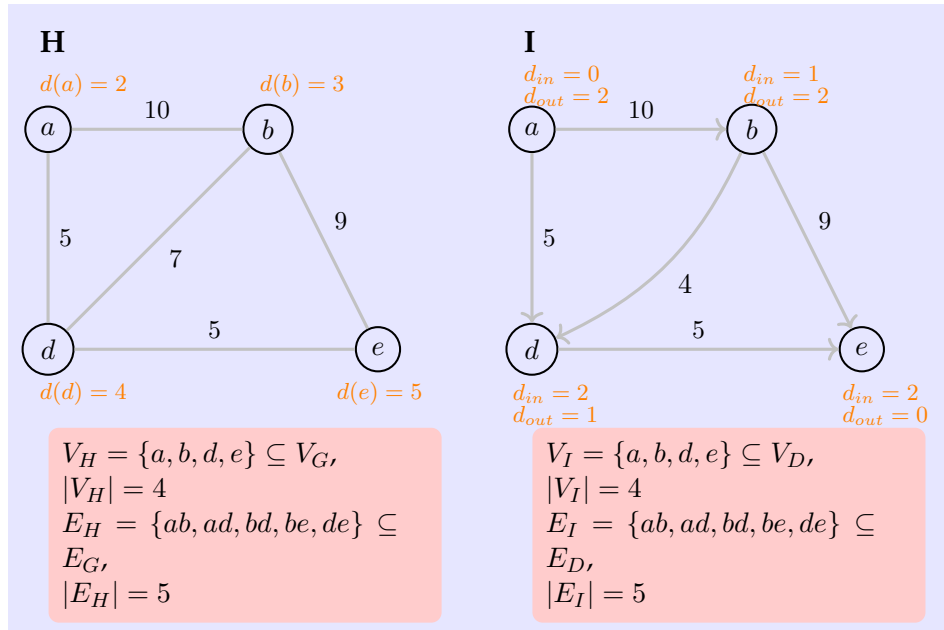


FIGURE 2.3: Undirected weighted graph **H** is a subgraph of **G** in Figure XX, $H \subseteq G$. Directed weighted graph **I** is a subgraph of **D** in Figure XX, $I \subseteq D$. The degree of each node is shown next to the corresponding node. The red box shows the vertex set, the edge set and their corresponding norms.

1. for each edge $e \in E$ we have $0 \leq f(e) \leq c(e)$. That is, the flow is positive and cannot exceed the capacity of the edge.

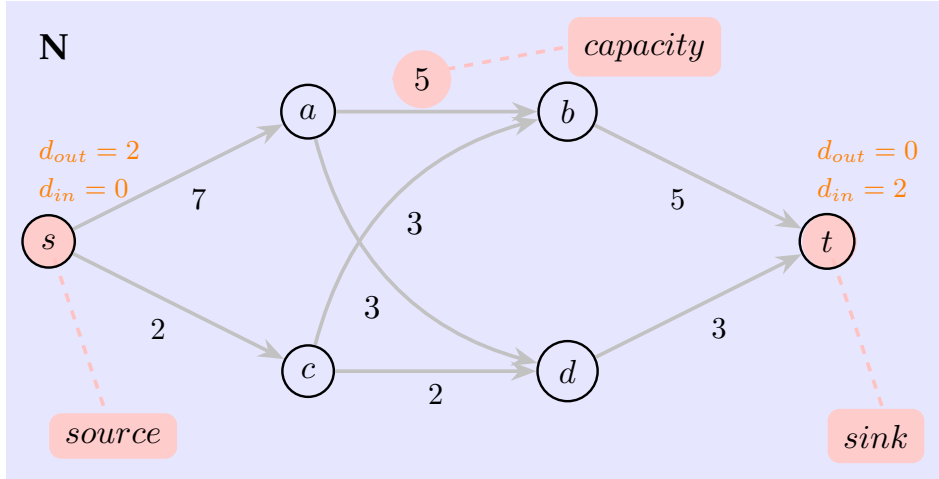


FIGURE 2.4: Network N with no flow. The in-degree and out-degree for the source, s , and the sink, t , are shown next to the corresponding node.

- for each intermediate node $v \in V \setminus \{s, t\}$ the in- and out-flow of that node $\sum_{u \in V^-(v)} f(u, v) = \sum_{u \in V^+(v)} f(v, u)$.

The **total flow** F of a network is then what leave the source s or reaches the sink t :

$$F(N) := \sum_{u \in V} f(s, u) - \sum_{u \in V} f(u, s) = \sum_{u \in V} f(u, t) - \sum_{u \in V} f(t, u) \quad (2.1)$$

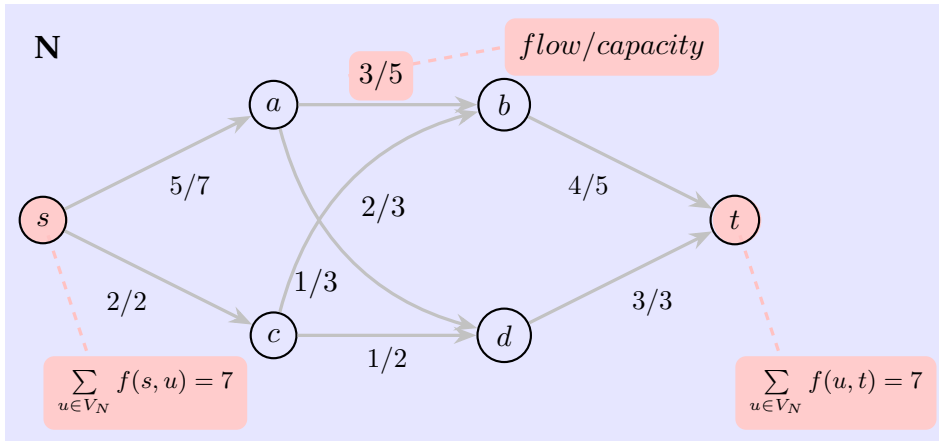


FIGURE 2.5: Network N with flow. The flow out of the source node, s , is equal to the flow into the sink node, t . For all other nodes, the flow-in is equal to the flow-out. This is the conservation of flow principle. This is only part of the network. The remaining part is the residual graph which shows the amount of reverse flow is available on an edge.

Cut A cut of a network $N = (V, E)$ is a partitioning of the vertex set $V = P \cup \bar{P}$ into two disjoint sets P containing the source node s and \bar{P} containing the sink node t . $P \cap \bar{P} = \emptyset$.

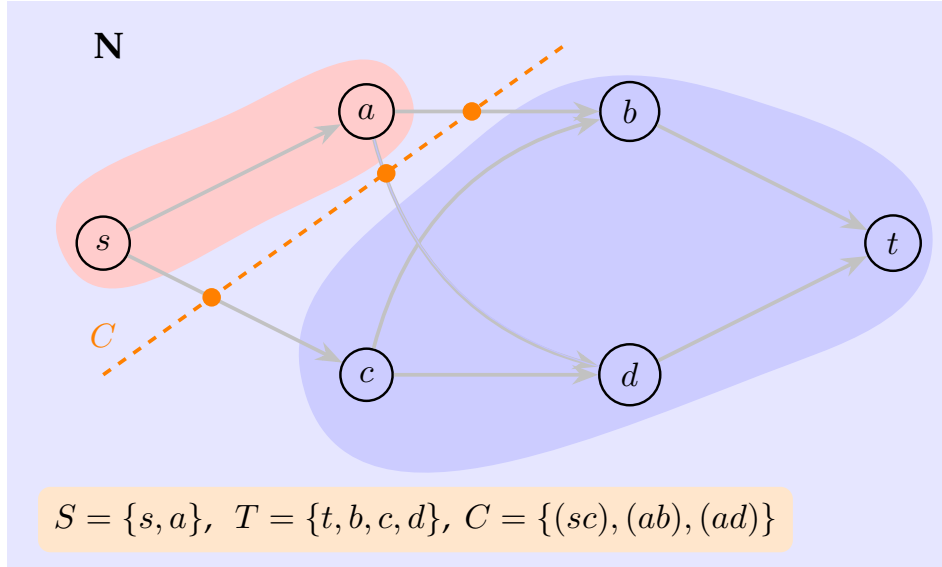


FIGURE 2.6: Network N with with a valid cut C . The nodes within the red region are reachable from the source and the nodes within the blue region are able to reach the sink. The cut set, C , is show in the orange filled block.

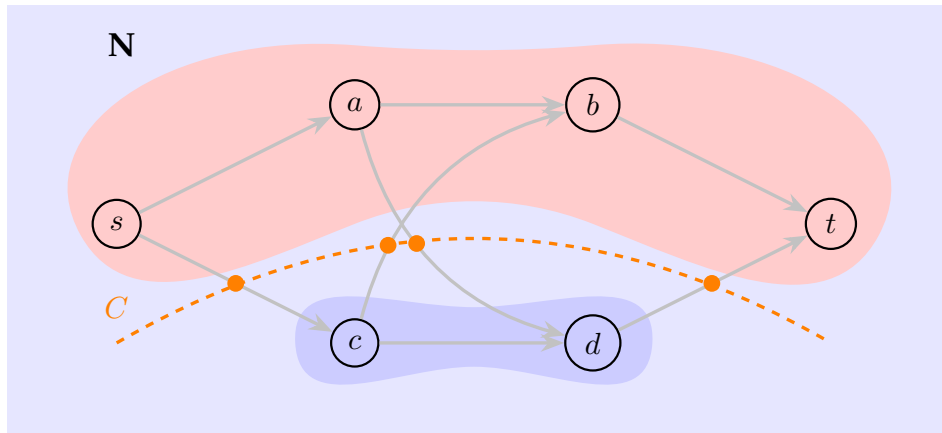


FIGURE 2.7: Network N with with a invalid cut C . The cut does not partition source node s and sink node t into distinct sets.

The **capacity** of a cut is the sum of the edges $(u, v) \in E$ where $u \in P$ and $v \in \bar{P}$:

$$\kappa(P, \bar{P}) = \sum_{u \in P; v \in \bar{P}} c(u, v) \quad (2.2)$$

Maximal Flow The largest amount of flow that can be sent through the source that is able to reach the sink is known as the maximal flow.

Minimal Cut A cut C on a network $N = (V, E)$ is a minimal cut if there exists no other cut C' where $\kappa(C') < \kappa(C)$.

In the next section we show that the Maximal Flow problem and the Minimal Cut problem are duals of each other, commonly known as the Max-Flow/Min-Cut problem.

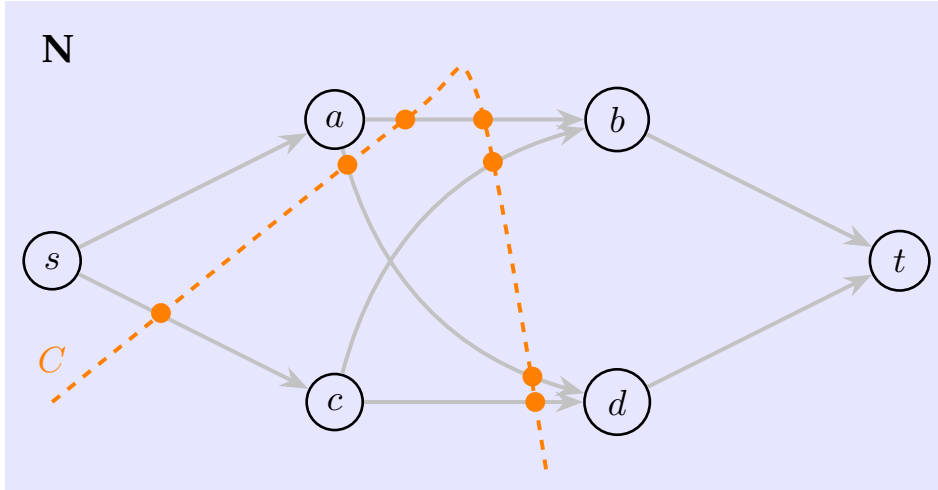


FIGURE 2.8: Network N with with a invalid cut C . The cut partition partitions the graph into more than two sets and the cut intersects the edges ab and ad twice.

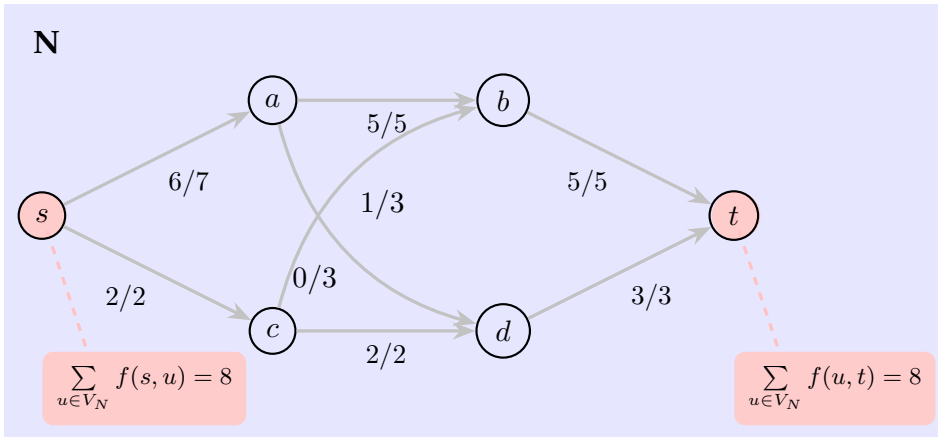


FIGURE 2.9: Network N with maximum flow. There is no way to push more flow out of the source into the sink without breaking the rules for conservation of flow.

2.3 Max-Flow/Min-Cut Problem

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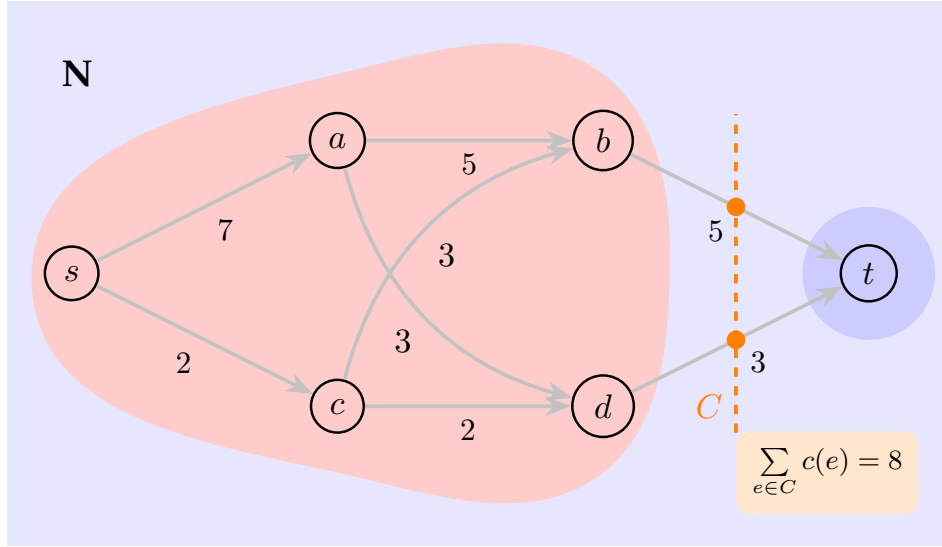


FIGURE 2.10: Network N with minimal cut C . The sum of the capacity of all the edges in the cut set is the minimum of all possible valid cuts on the network N .

2.4 Modelling Images as Graphs

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2.4.1 Sub-modularity Conditions for Discrete Systems

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2.4.2 Connectivity

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2.4.3 Distance Metrics

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Euclidean Distance

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Riemannian Distance

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Learned Distance from Seeds

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2.5 Max-Flow/Min-Cut Algorithms

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2.5.1 Ford-Fulkerson

Nunc posuere quam at lectus tristique eu ultrices augue venenatis. Vestibulum ante ipsum primis in faucibus orci luctus et ultrices posuere cubilia Curae; Aliquam erat volutpat. Vivamus sodales tortor eget quam adipiscing in vulputate ante ullamcorper. Sed eros ante, lacinia et sollicitudin et, aliquam sit amet augue. In hac habitasse platea dictumst.

Algorithm 1 Euclid's algorithm

```

1: procedure EUCLID( $a, b$ )                                ▷ The g.c.d. of  $a$  and  $b$ 
2:    $r \leftarrow a \bmod b$ 
3:   while  $r \neq 0$  do                                     ▷ We have the answer if  $r$  is 0
4:      $a \leftarrow b$ 
5:      $b \leftarrow r$ 
6:      $r \leftarrow a \bmod b$ 
7:   end while
8:   return  $b$                                              ▷ The gcd is  $b$ 
9: end procedure

```

2.5.2 Dinic/Edmond-Karp

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2.5.3 Push-Relabel

Originally developed by Andrew V. Goldberg and Robert E. Tarjan. Previous algorithms, such as Ford-Fulkerson, used the concept of residual networks and augmenting paths to determine max-flow. Push-Relabel used the concept of preflow to determine max-flow instead of augmenting paths. Sometimes referred as the Preflow-Push Algorithm. Preflow is a concept originally developed by A.V. Karzanov.

The algorithm works at converting a preflow, f , into a normal flow and then terminates. This flow also turns out to be the maximum flow. Goldberg and Tarjan defined a generic Push-Relabel algorithm which solves the maximum flow problem.

Preflow A preflow is a real-valued function, f , on vertex pairs. The total flow into a vertex can exceed the flow out of a vertex but not vice versa.

A preflow where all $v \in V - \{s, t\}$ has a flow excess of zero, $e_f(v) = 0$, is a normal flow. The preflow function is also referred to as the **s-t preflow**.

Preflow must satisfy:

1. **Capacity Constraint**
 $\forall u, v \in V, f(u, v) \leq c(u, v)$
2. **Antisymmetry/Skew Symmetry**
 $\forall u, v \in V, f(u, v) = -f(v, u)$
3. **Nonnegative Constrain**
 The flow into $v \in V - \{s\}$ must be greater than or equal to the flow out of v . $\forall u \in V, v \in V - \{s\}, \sum f(u, v) > 0$

Flow Excess Flow excess, $e_f(v)$, is the net flow into v where $v \in V$ for some preflow f .

$$e_f(v) = \begin{cases} \infty & \text{if } v = s \\ \sum_{u \in V} f(u, v) & \text{if } v \in V - \{s\} \end{cases}$$

Active Vertex An active vertex/node is a vertex v which satisfies all of the properties:

1. Not a source or sink, $v \in V - \{s, t\}$
2. Positive flow excess, $e_f(v) > 0$
3. Has a valid label, $d(v) < \infty$

Push-Relabel also uses the concept of a residual graph, $G_f = (V, E_f)$.

Residual Capacity The residual capacity of a preflow is defined as $r_f(v, w) = c(v, w) - f(v, w)$.

Residual Edges The residual edges for a preflow f is defined as the set of edges with positive residual capacity. $E_f = \{(v, w) | r - f(v, w) > 0\}$.

Labelling Push-Relabel also use a valid labelling function, d , to determine which vertex pairs should be selected for the push operation.

A valid labelling, d , is a nonnegative integer function applied to all vertices to denote a label. The labelling is often referred as the height or distance from the sink node, t . This function is sometimes compared to the physical intuition that liquids naturally flow downhill.

A valid labelling for a preflow consists of:

1. For $v \in V, 0 \leq d(v) \leq \infty$
2. $d(s) = |V|$ (source condition)
3. $d(t) = 0$ (sink condition)
4. $d(v) = d(w) + 1$ for every residual edge $(v, w) \in E_f$

A labelling d and a preflow f are said to be compatible if d adheres to the properties above.

The algorithm pushes flow excess starting at the source, s , along all vertices towards the sink, t . The algorithm maintains a compatible vertex labelling function, d , to the preflow, f . The labelling is used to determine where to push the flow excess. The algorithm repeatedly performs either a push or a relabel operation so long as there is an active vertex in G_f .

Push Operation The push operation is used to move flow from one vertex to another. The transfer of excess can be performed across the vertex pair $(v, w) \in E_f$ if:

1. v is an active vertex
2. the edge has positive residual capacity, $r_f(v, w) > 0$
3. the label distance $d(v) = d(w) + 1$

This allows the algorithm to move δ excess flow: $\delta = \min(e_f(v), r_f(v, w))$ from v to w . A push is considered **saturating** if no more flow can be sent over the edge, $\delta = r_f(v, w)$. A push is considered to be **non-saturating** if all the excess from v is pushed over the edge and the edge still has some capacity, $\delta = e_f(v)$.

Algorithm 2 Push Operation

Input: Preflow f , labels d , and (v, w) where $v, w \in V$

Output: Preflow f

Applicable: if $v \in V - \{s, t\}$, $d(v) < \infty$, $e_f(v) > 0$, $r_f(v, w) > 0$ and $d(v) = d(w) + 1$

- 1: **procedure** PUSH(v, w)
 - 2: $\delta := \min(e_f(v), r_f(v, w))$
 - 3: $f(v, w) := f(v, w) + \delta$
 - 4: $f(w, v) := f(w, v) - \delta$
 - 5: $e_f(v) := e_f(v) - \delta$
 - 6: $e_f(w) := e_f(w) + \delta$
 - 7: **return** f
 - 8: **end procedure**
-

Relabel Operation The relabel operation is used to increase the label value of a single active vertex so that excess flow can be pushed out of the active vertex. The relabel operation is performed when all the residual edges of the active vertex have positive residual capacity, $r_f(v, w) > 0$. This implies that v 's label is less than or equal to all vertices, $d(v) \leq d(w)$, meaning that no push operation across the edges is possible given the push condition $d(v) = d(w) + 1$.

The relabel operation for some vertex v selects the smallest label for the vertices with positive residual edges, $r_f(v, w) > 0$. The active vertex is then assigned the smallest label value +1 such that $d(v) := \min_{(v, w) \in E_f} d(w) + 1$. This will allow the vertex v to potentially push its excess flow to at least one of the other vertices during the algorithm's next iteration.

The algorithm initialises the following values in the residual graph before the push and relabel operations in the main loop.

Algorithm 3 Relabel Operation**Input:** Preflow f , labels d , and $v \in V - \{s, t\}$ **Output:** Labels d **Applicable:** if $v \in V - \{s, t\}$, $d(v) < \infty$, $e_f(v) > 0$, and $\forall w \in V, r_f(v, w) > 0$ which implies $d(v) \leq d(w)$

```

1: procedure RELABEL( $v$ )
2:   if  $\{(v, w) \in E_f\} \neq \emptyset$  then
3:      $d(v) := \min(d(w) + 1 | (v, w) \in E_f)$ 
4:   else
5:      $d(v) := \infty$ 
6:   end if
7:   return  $d$ 
8: end procedure

```

1. Initialise the preflow of all edges in the residual graph
2. Initialise the labellings such that:
 - (a) $d(s) = |V|$
 - (b) $d(v) = 0$ for $v \in V - \{s\}$
3. Performs saturation, pushes along all residual edges out of the source $(s, v) \in E_f$ and $v \in V$.

Once complete the algorithm repeatedly performs either a push or a relabel operation against all vertices. The algorithm continues until no operation can be performed. The algorithm terminates when there are no more active vertices.

The analysis and the proof of correctness of the Push-Relabel algorithm can be found in [Appendix B](#).

Push-Relabel Speed Optimisation Heuristics

Discharge Push-Relabel also use a valid labelling function d , to determine which vertex pairs should be selected for the push operation.

FIFO Push-Relabel also use a valid labelling function d , to determine which vertex pairs should be selected for the push operation.

Highest Label First Push-Relabel also use a valid labelling function d , to determine which vertex pairs should be selected for the push operation.

Global Relabel Push-Relabel also use a valid labelling function d , to determine which vertex pairs should be selected for the push operation.

Gap Relabel Push-Relabel also use a valid labelling function d , to determine which vertex pairs should be selected for the push operation.

Algorithm 4 Push-Relabel Main-loop

Input: Network flow graph $G = (V, E)$, s, t and c **Output:** Maximum flow f

```

1: procedure MAIN( $v$ )
2:   for all  $(v, w) \in (V - \{s\})(V - \{s\})$  do
3:      $f(v, w) := 0$ 
4:      $f(w, v) := 0$ 
5:   end for

6:   for all  $v \in V$  do
7:      $f(s, v) := r_f(s, v)$ 
8:      $f(v, s) := -r_f(s, v)$ 
9:   end for

10:   $d(s) = |V|$ 

11:  for all  $v \in V - \{s\}$  do
12:     $d(v) := 0$ 
13:     $e_f(v) := f(s, v)$ 
14:  end for

15:  While there exists an active vertex
16:    while  $\exists v \in V - \{s, t\}$  do  $\triangleright$  with either applicable PUSH() or
      RELABEL() operation
17:      Perform either a PUSH or a RELABEL operation on  $v$ 
18:    end while

19:  return  $f$ 
20: end procedure

```

2.5.4 Alpha-Beta Swap

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2.5.5 Alpha-Expansion

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Chapter 3

Interactive Segmentation on Single Channel Data

3.1 Seeding

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3.2 Estimating Probability Distributions

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3.2.1 Expectation Maximisation (Guassian Mixture Modelling)

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Fixed-Distribution Modelling

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Automatically Determining the Optimal Number of Mixtures

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3.2.2 Naive Bayesian Classification

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3.2.3 Supervised Learning for Multi-Layered Perceptron

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Chapter 4

Interactive Segmentation Extention to Multi-Channel Data

Chapter 5

Automatic Graph Cut Segmentation

5.1 Determining FG and BG seeds/probability distributions

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5.2 Determining Optimal Parameters Settings

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Chapter 6

Graph Cut Solution to Chan-Vese Segmentation

Chapter 7

Pre-Processing and Post-Processing Techniques

7.1 Removal of Artifacts using Connected Components

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7.2 Anisotropic Diffusion

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7.2.1 Coherence Enhancing Diffusion

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7.2.2 Coherence Enhancing Diffusion with Optimised Rotational Invariance

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7.3 Poisson Denoising

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7.3.1 Total-Variation Denoising

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7.4 Contrast Enhancement

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Chapter 8

Case-Study: Flourescence Microscopy Cell Segmentation

8.1 Removal of Artifacts using Connected Components

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8.2 Anisotropic Diffusion

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8.2.1 Coherence Enhancing Diffusion

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8.2.2 Coherence Enhancing Diffusion with Optimised Rotational Invariance

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8.3 Poisson Denoising

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8.3.1 Total-Variation Denoising

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8.4 Contrast Enhancement

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Appendix A

Proof of Maximum Flow - Minimum Cut Equivalence

Write your Appendix content here.

Appendix B

Push-Relabel Algorithm Analysis and Proof of Correctness

Write your Appendix content here.