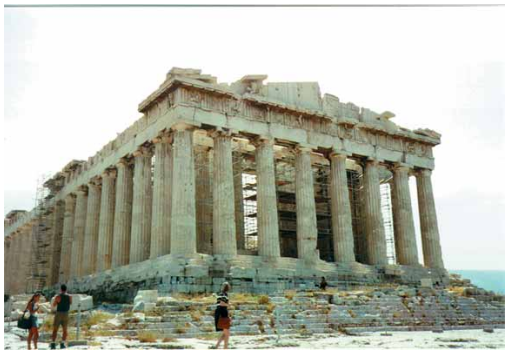
The difference between a good and a poor architect is that the poor architect succumbs to every temptation and the good one resists it.
-- Ludwig Wittgenstein

Chapter 2

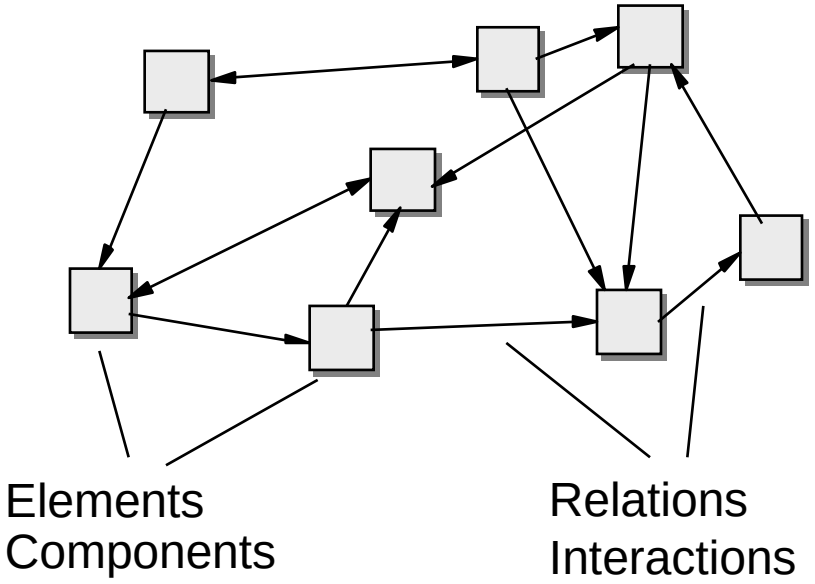
Architecture



2.1 Coarse structure

A system (e.g. operating system, software system) consists of

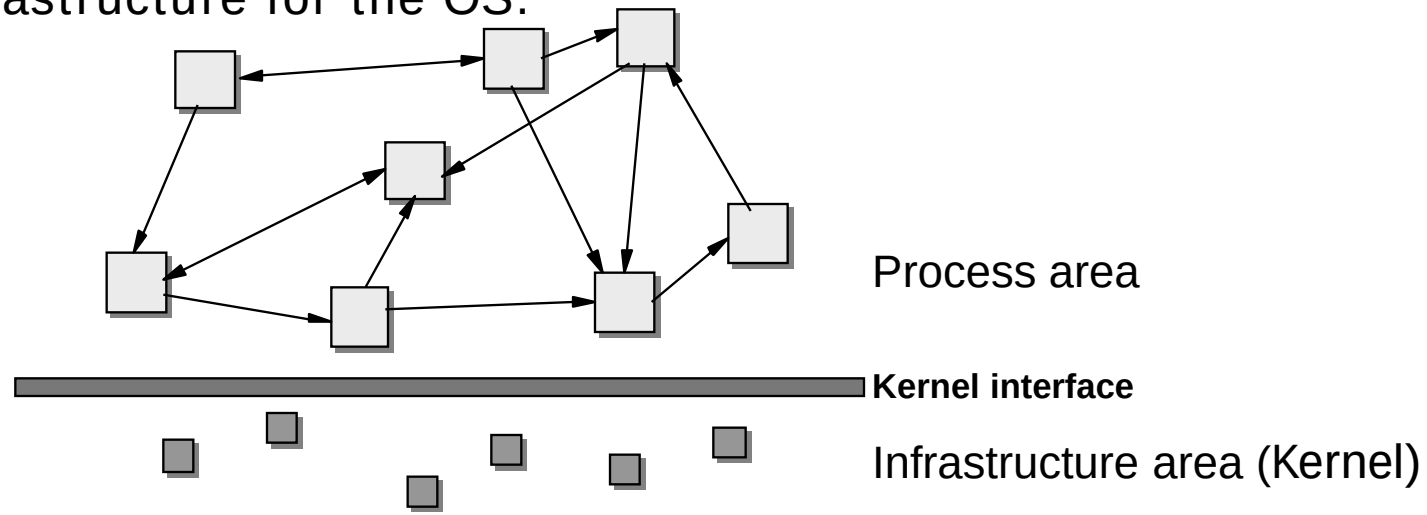
- Elements
- Relations between elements



The elements are functional units with interactions of different kinds in between (data flow, request flow, synchronization, call, communication, ...).

Coarse division of an OS

- In operating systems the elements/components are the *processes*.
- An operating system consists of a set of interacting processes.
- Since processes are not offered by the hardware, there must be something that provides the concept of a process and the interaction between processes.
- This „something“ is called kernel of the operating system. It provides the basic infrastructure for the OS.



- In a first coarse structure, we therefore distinguish two areas:
  - Process area, where the essential OS functionality is located.
  - Kernel area that provides the fundamental infrastructure for these processes.

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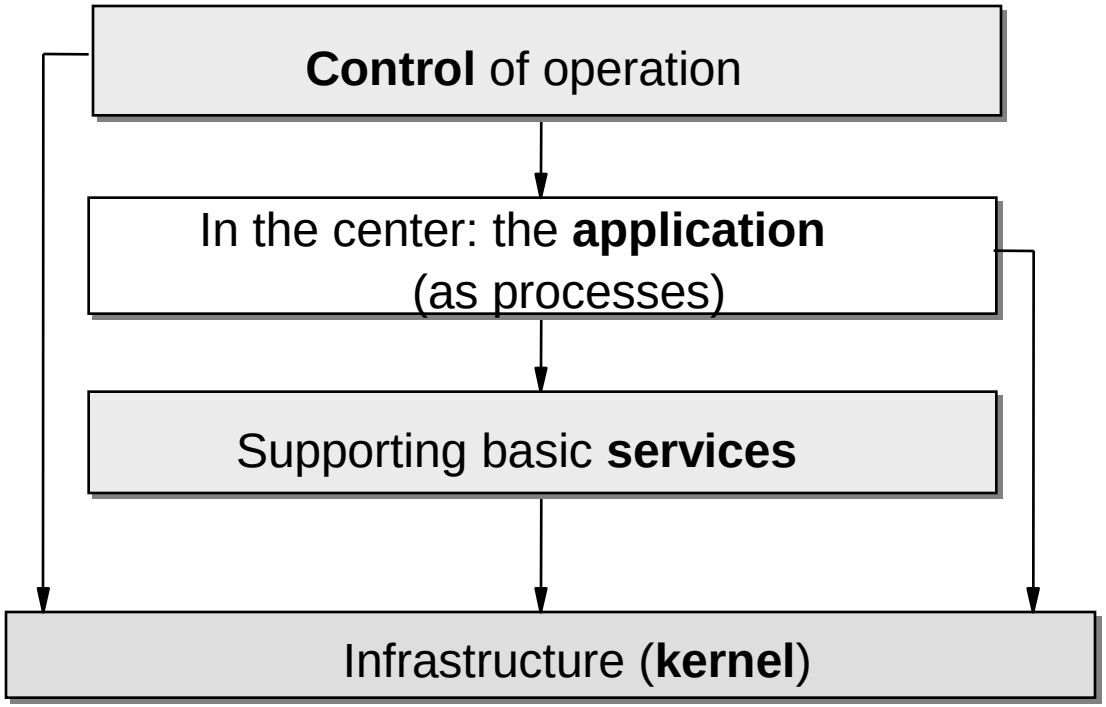
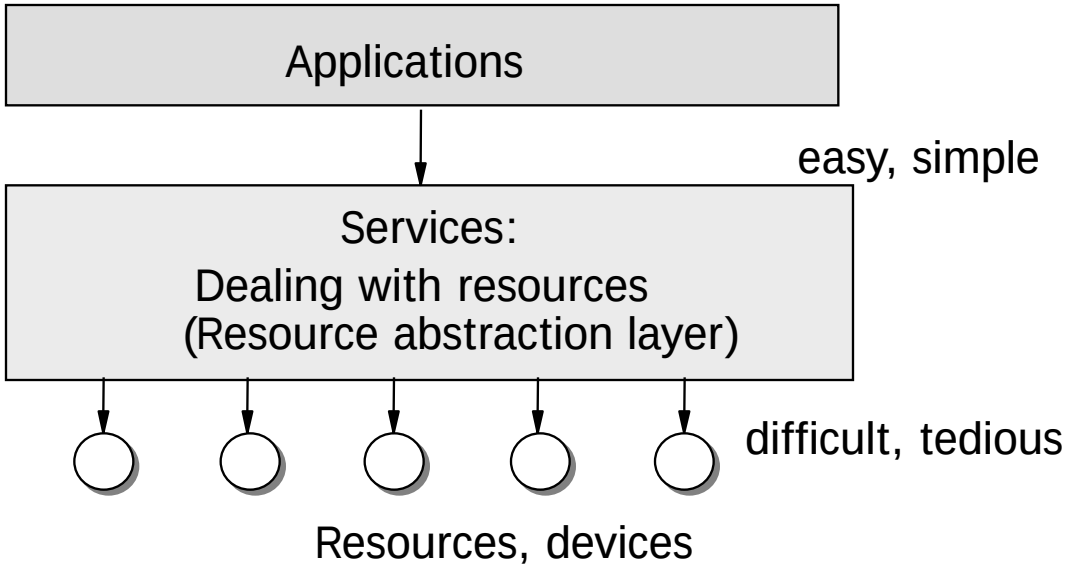
Finer resolution: Process area



Services

Finer resolution: Services



Distinction

Logical resource:	„thought up“ for organizational reasons realized by real resources Examples: file, window
Real (physical) resource:	real existence, to be touched Examples: disc, keyboard, ...

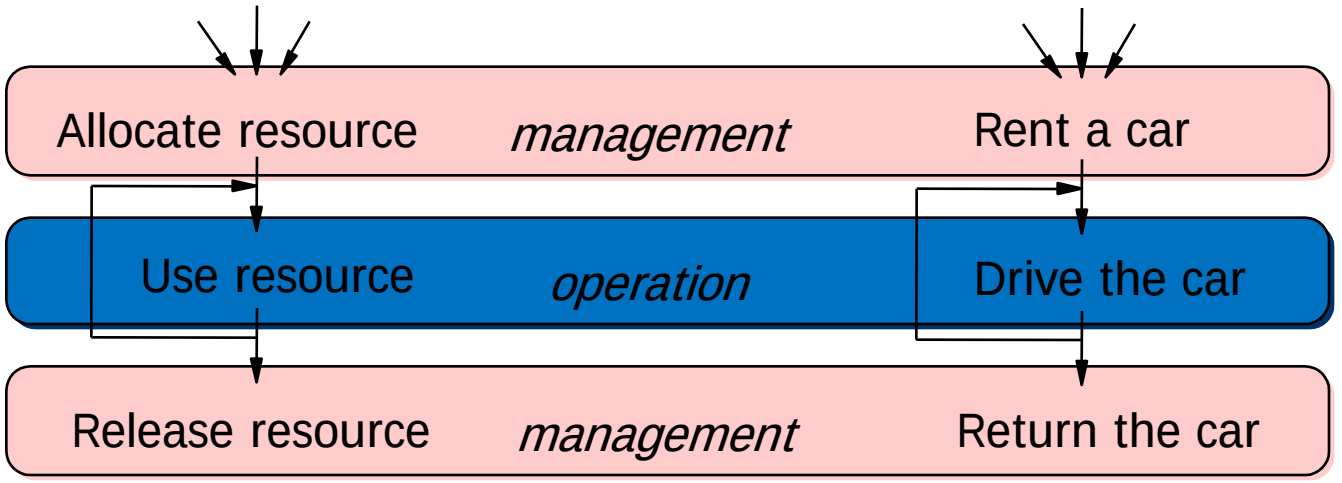
How to deal with resources

Two aspects:

Resource management: only in case of competition for resources: who may accesses what and when?

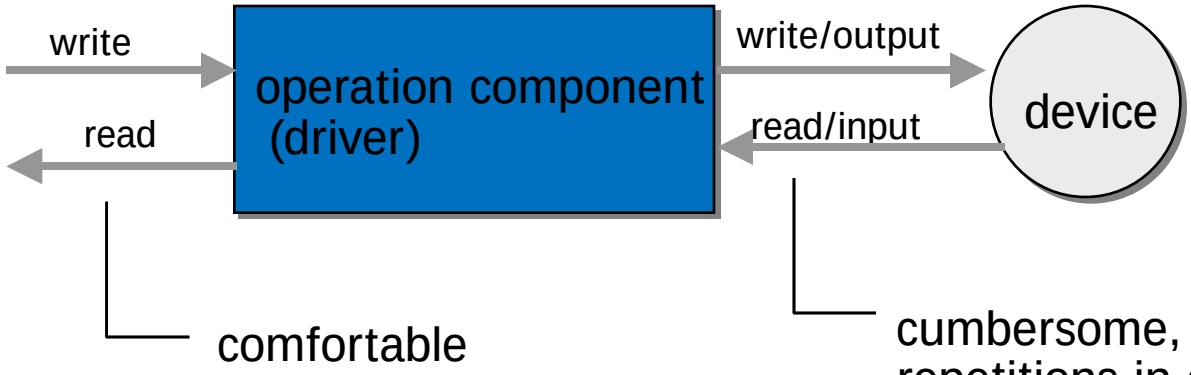
Resource operation: real usage, e.g. data transport

Example:

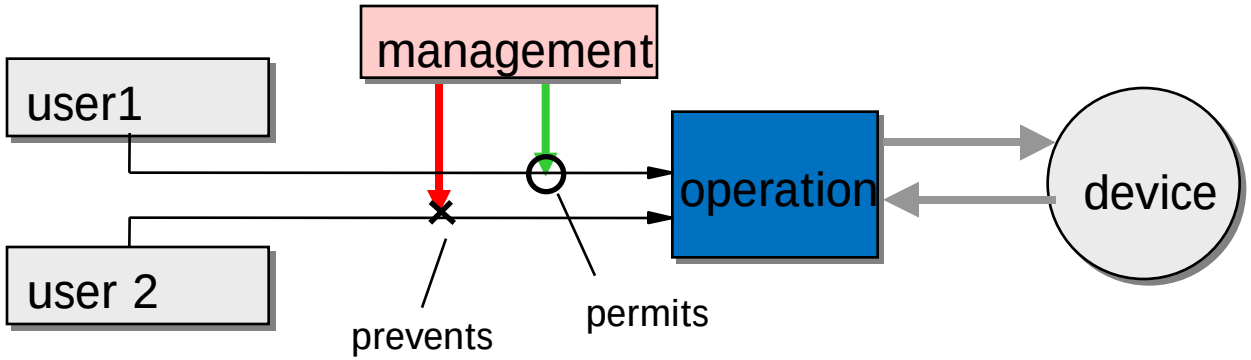


Dealing with resources

The term "operation"

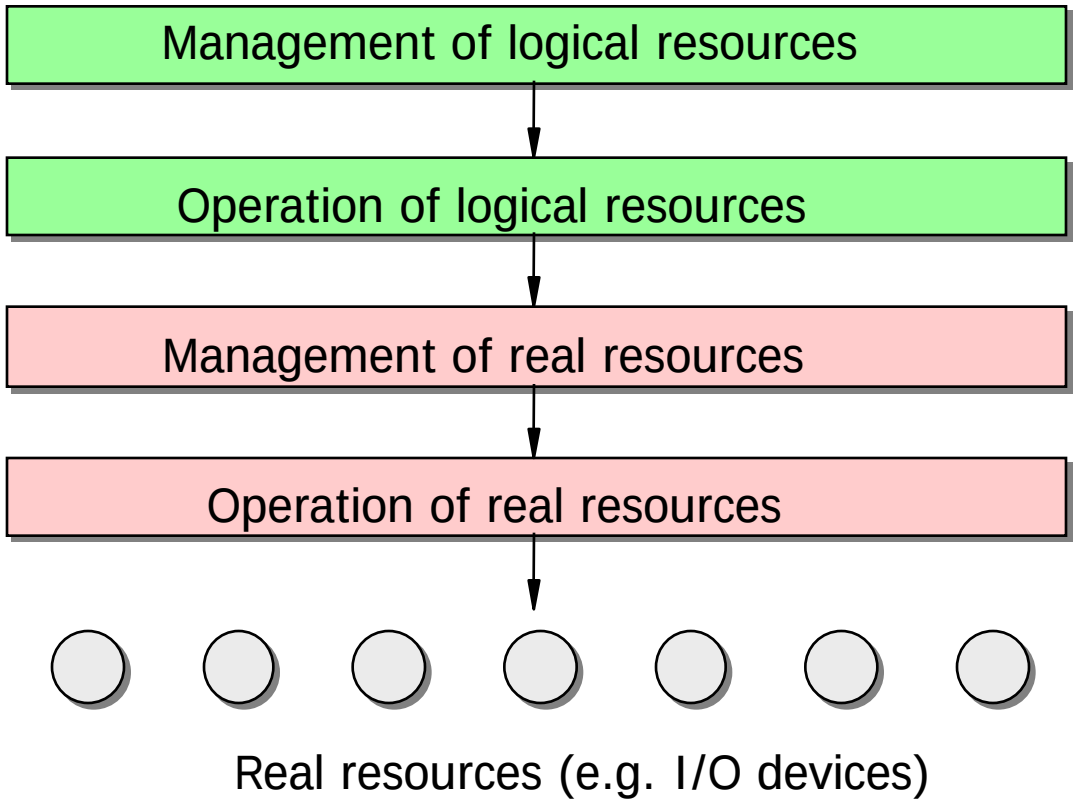


The term "management"



Operating system services

- Structure of service layer

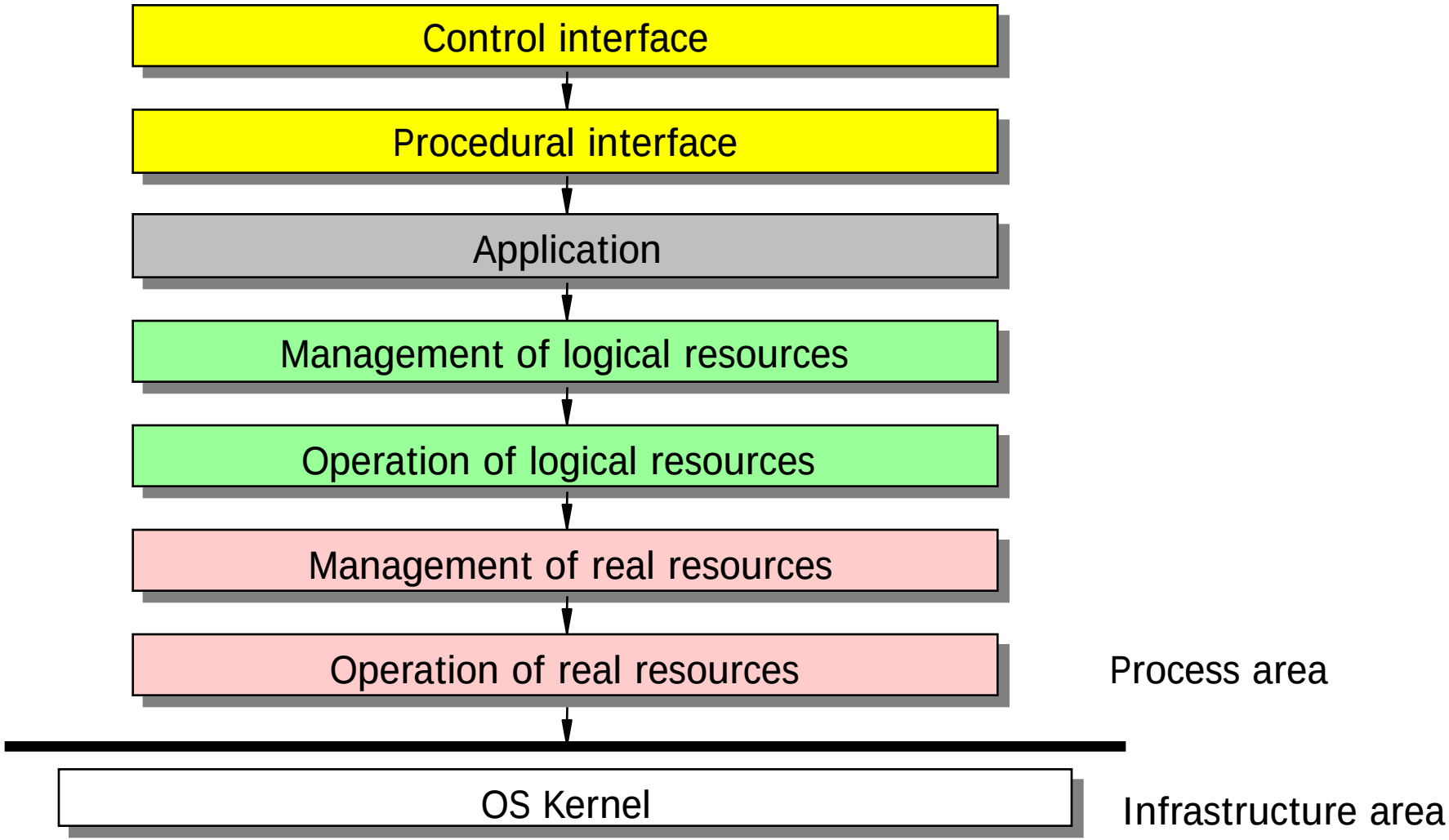


Control

We distinguish

- Control interface
  - Interaction between human and machine
    - OS-commands
    - User interface (textual, graphical, touch, acoustic, …)
- Procedural interface
  - Possibility to define complex requests to the OS
  - Programming language notation with embedded OS commands to define and control complex tasks

Overview

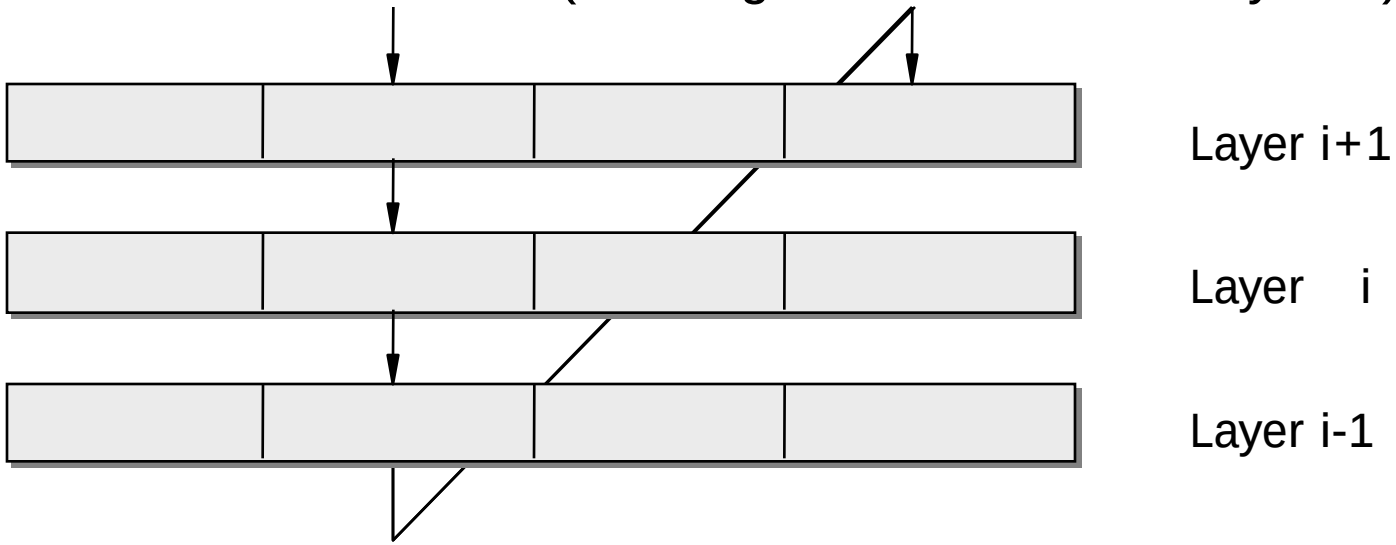

Control interface
Procedural interface
Application
Management of logical resources
Operation of logical resources
Management of real resources
Operation of real resources
Process area
OS Kernel
Infrastructure area

Remarks

Each layer may be partitioned.

Operation device A	Operation device B	Operation device C	Operation device D

Upward calls are allowed (as long as there are no cycles).



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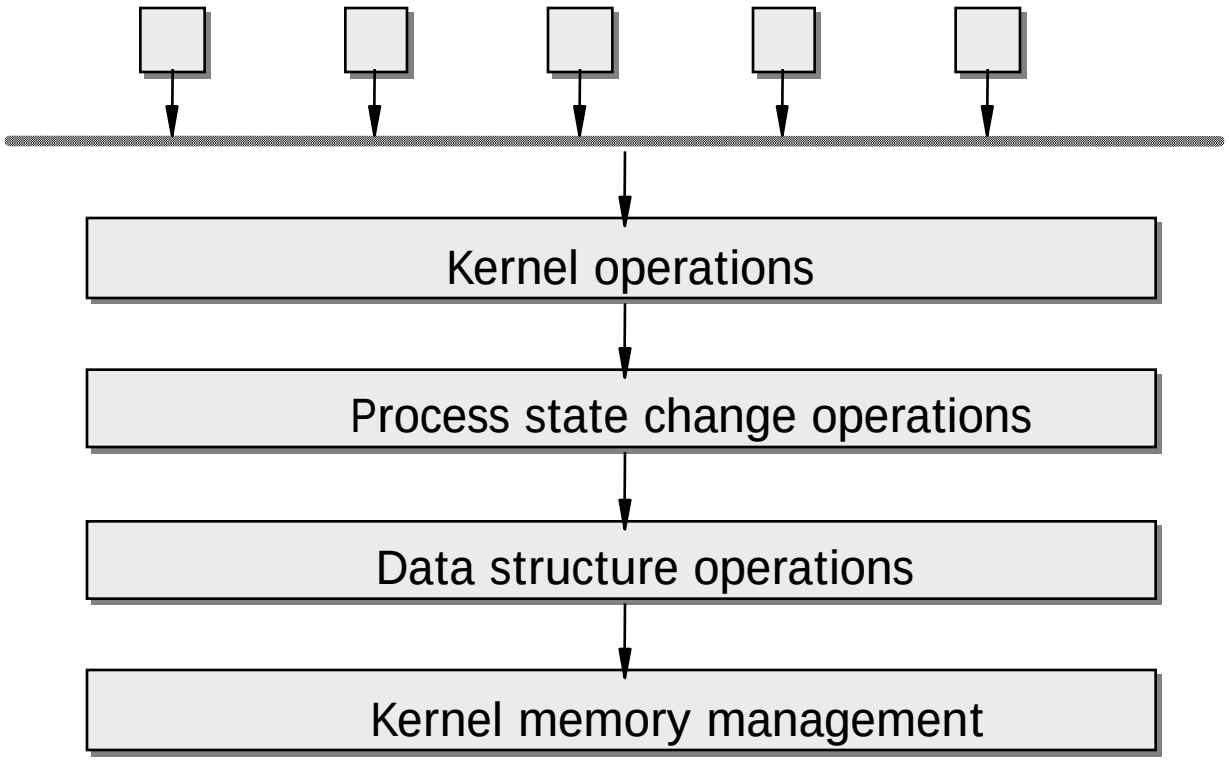
Finer resolution: Kernel



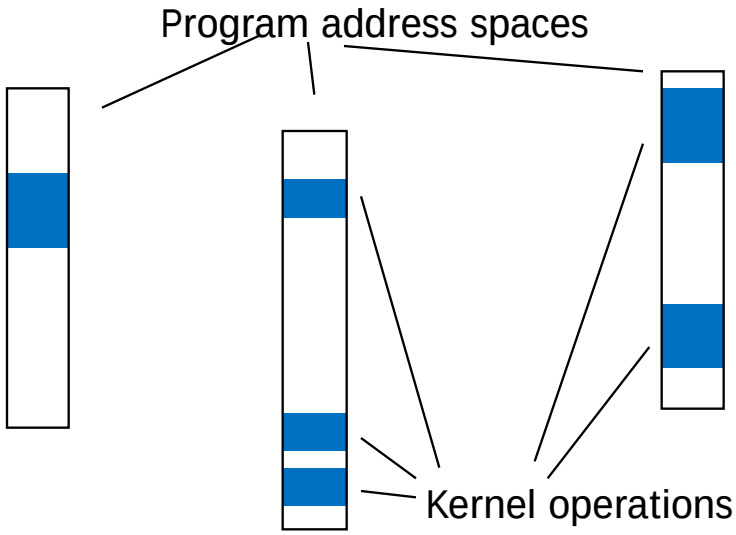
Processes

Kernel interface

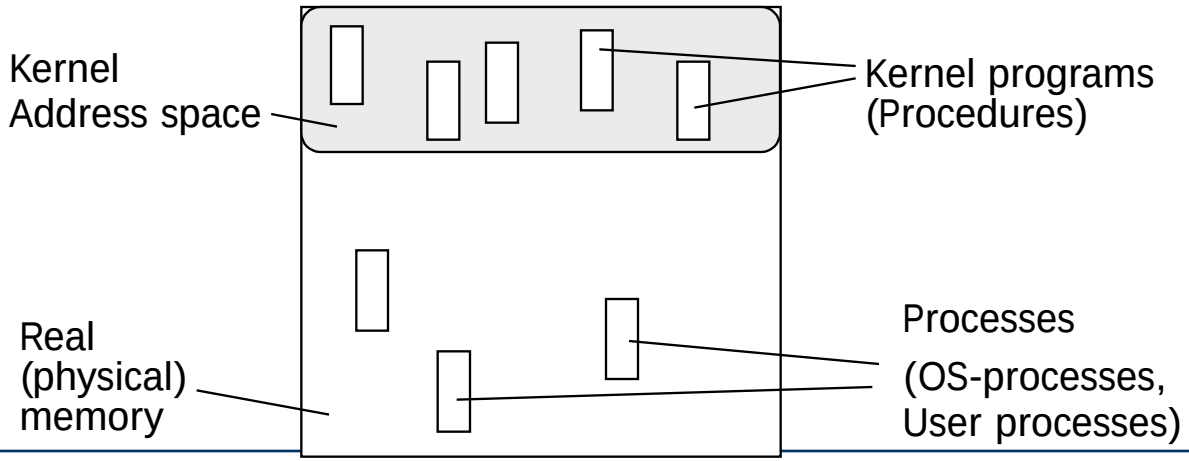
Dynamic systems only

Types of realization of the kernel

- Scattered across programs



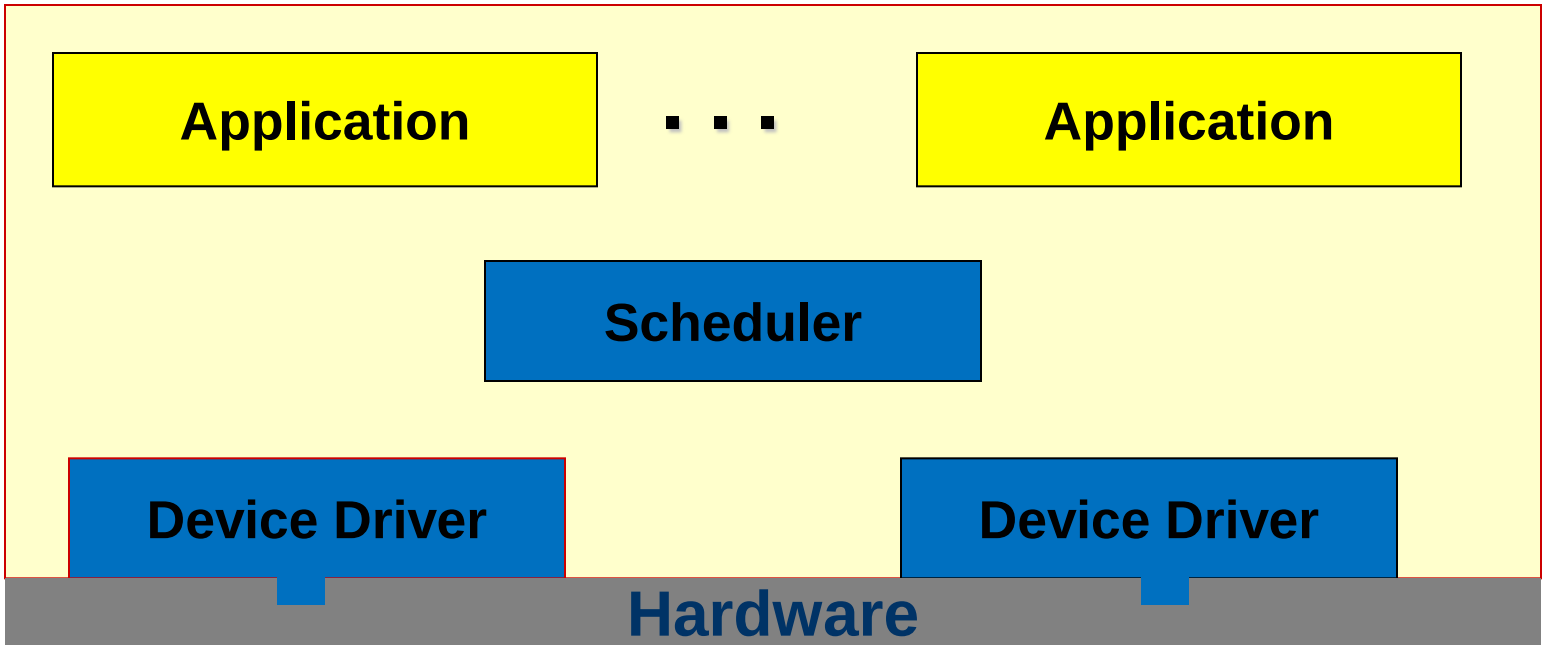
- Resident, compact, sealed



Microkernel architecture

- The architecture just described is called microkernel architecture in the literature.
- Historically, people simply wanted to make a distinction to UNIX, in which the whole resource management (e.g. file system) is part of the kernel (macrokernel).
- Therefore the sizes of kernels differ widely.
- *Mach* (OSF-1) uses several MByte memory footprint, while *Cosy* (own development) only needs some 100KB.
- Such deviations in size sometimes lead to names like *Nanokernel* or *Picokernel*.
- There is no general agreement what should be in the kernel.
- However, process management and process communication are essential.

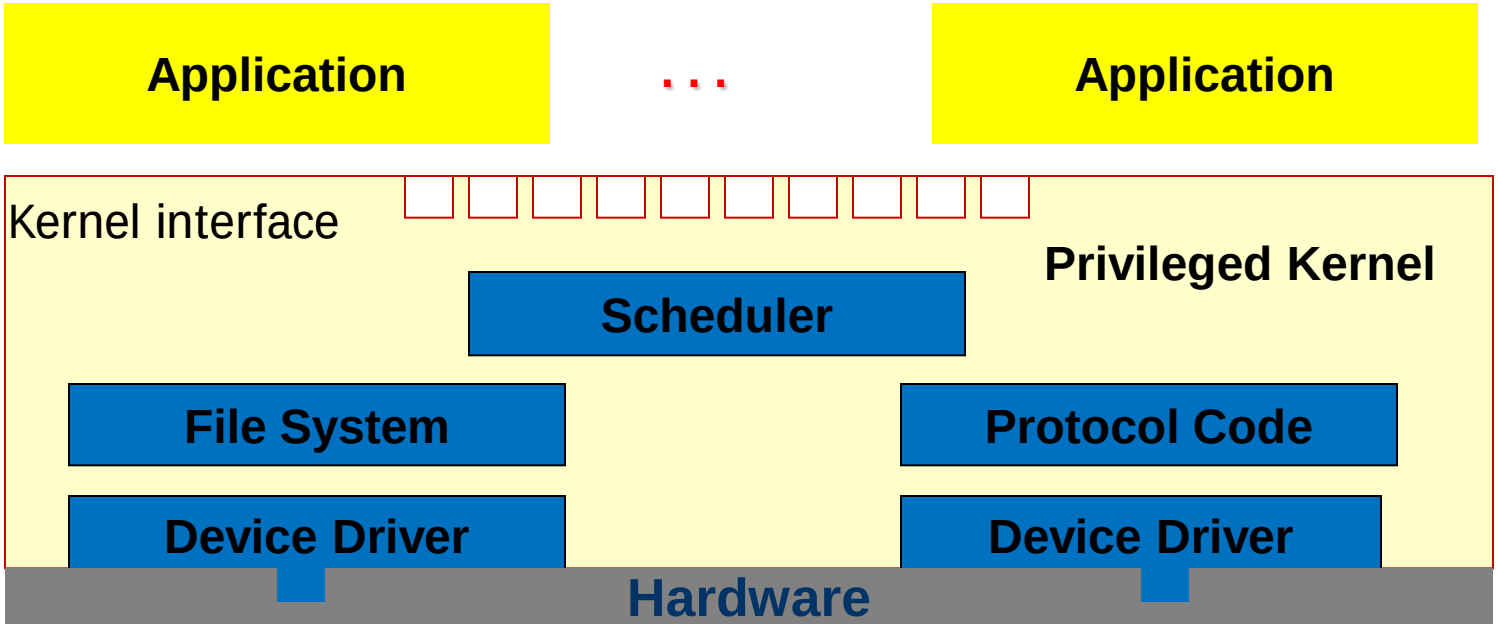
Monolithical Systems



Besides microkernel OS, there are other approaches available:

In **monolithical systems** there is no strict separation between application and operating system.
Appropriate for small static OS, e.g. in embedded systems

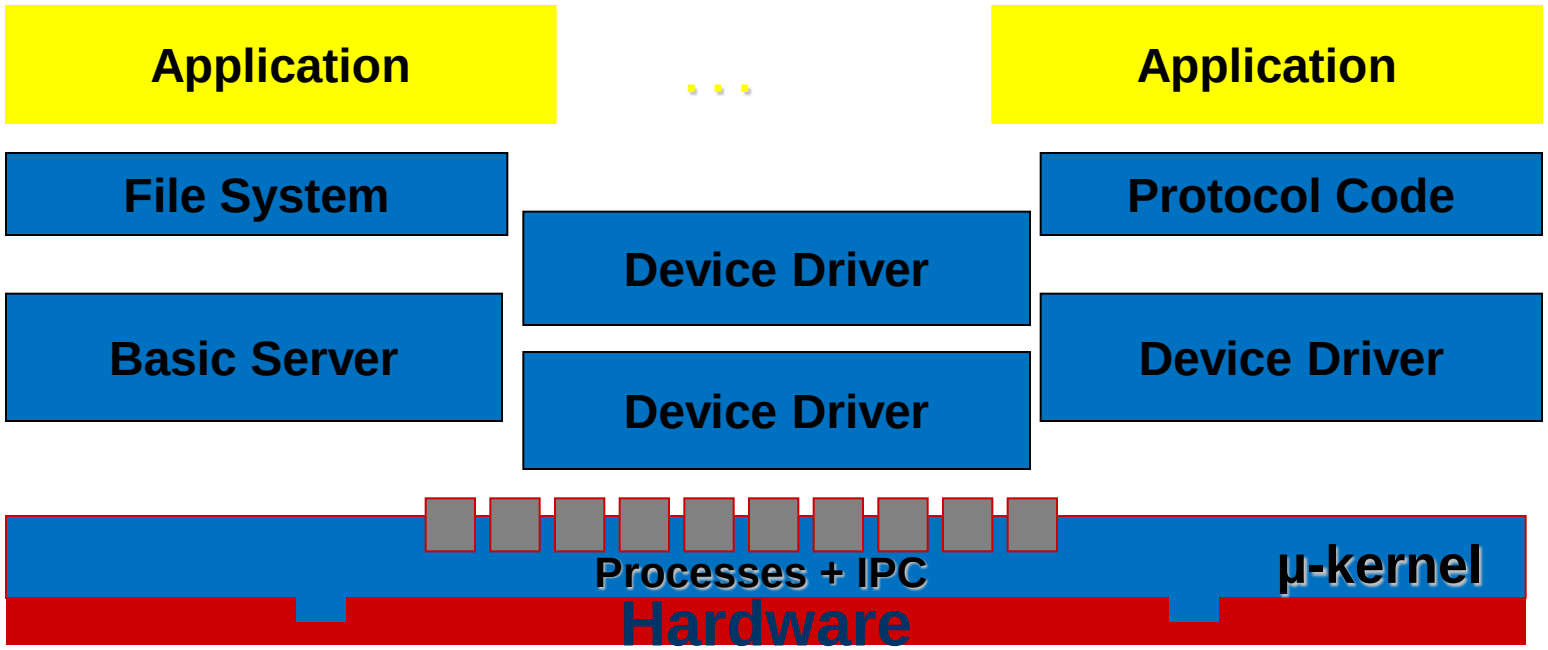
Monolithical OS-Kernel



Monolithical OS-**kernels** do have a separation (protection) between application and OS, but no separation among OS components.

The whole kernel needs to be trustworthy.

Microkernel OS



The kernel comprises process management (initializing, dispatching) and interprocess communication only.

Advantages of a Microkernel architecture

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- Clear kernel interface supports modular structure.
- Realization of services outside the kernel:
  - leads to more security and stability since the kernel will not be affected by faulty services,
  - improves flexibility and extendibility since services can be added and removed arbitrarily, even during operation.
- The safety-critical part of the system (kernel) is relatively small and can be verified easier.
- Usually, only the kernel needs to run in privileged mode.
- Microkernel architecture allows the coexistence of several alternative interfaces between OS and applications.

Drawbacks of Microkernel architecture

- Usually worse performance

Why?

- **Interplay of components outside the kernel requires more interprocess communication and therefore more kernel calls.**

2.4 General Design Principles

KISS (keep it small and simple; keep it simple, stupid)

- Occam's Razor: "Plurality should not be assumed without necessity."
  (*William of Ockham*, *ca. 1285-1349*)

- "There are two ways of constructing a software design: One way is to make it so simple that there are obviously no deficiencies, and the other way is to make it so complicated that there are no obvious deficiencies. The first method is far more difficult."
  (*C.A.R. Hoare, 1934-*)

- "Everything should be made as simple as possible, but not simpler." (*Albert Einstein, 1879-1955*)

- "Perfection (in design) is achieved not when there is nothing more to add, but rather when there is nothing more to take away." (*Antoine de Saint-Exupéry, 1900-1944)*

Modularization

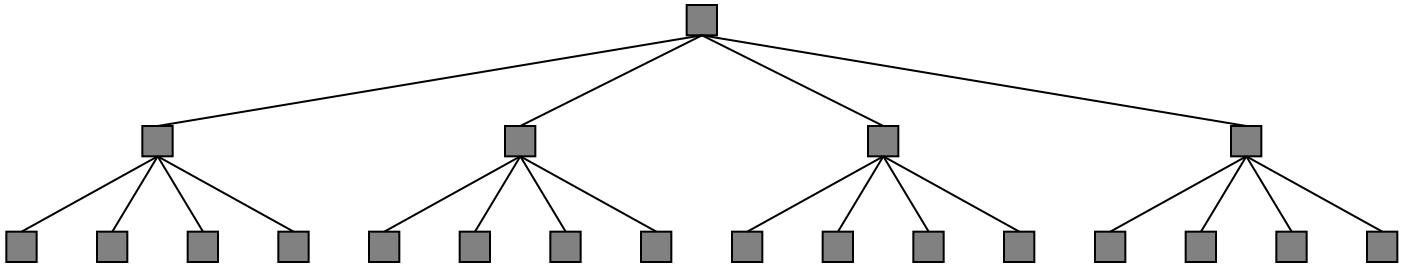
- The system is decomposed into a set of modules such that
  - the interaction (information and control flows) within the module is high,
  - the interaction between modules is low,
  - the interfaces between the modules are simple,
  - the modules are easily understandable due to their limited size and complexity.

- The principle can be applied hierarchically.

Hierarchization

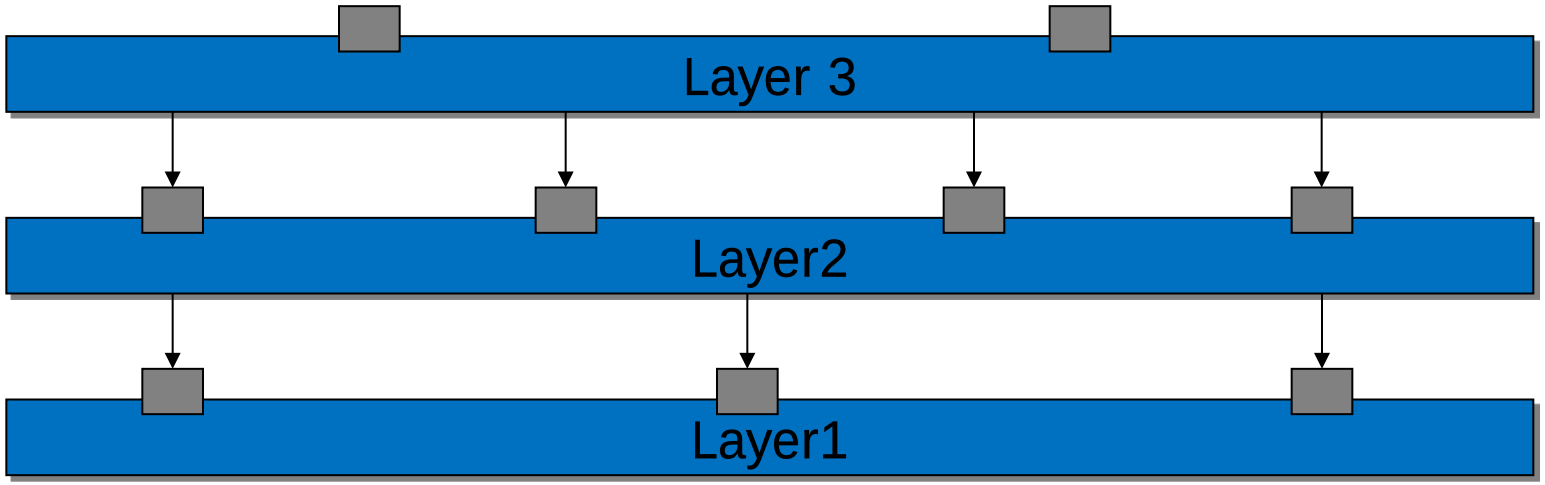
- Tree-like organization of homogeneous elements

- Goals:
  - scalability,
  - mastering complexity

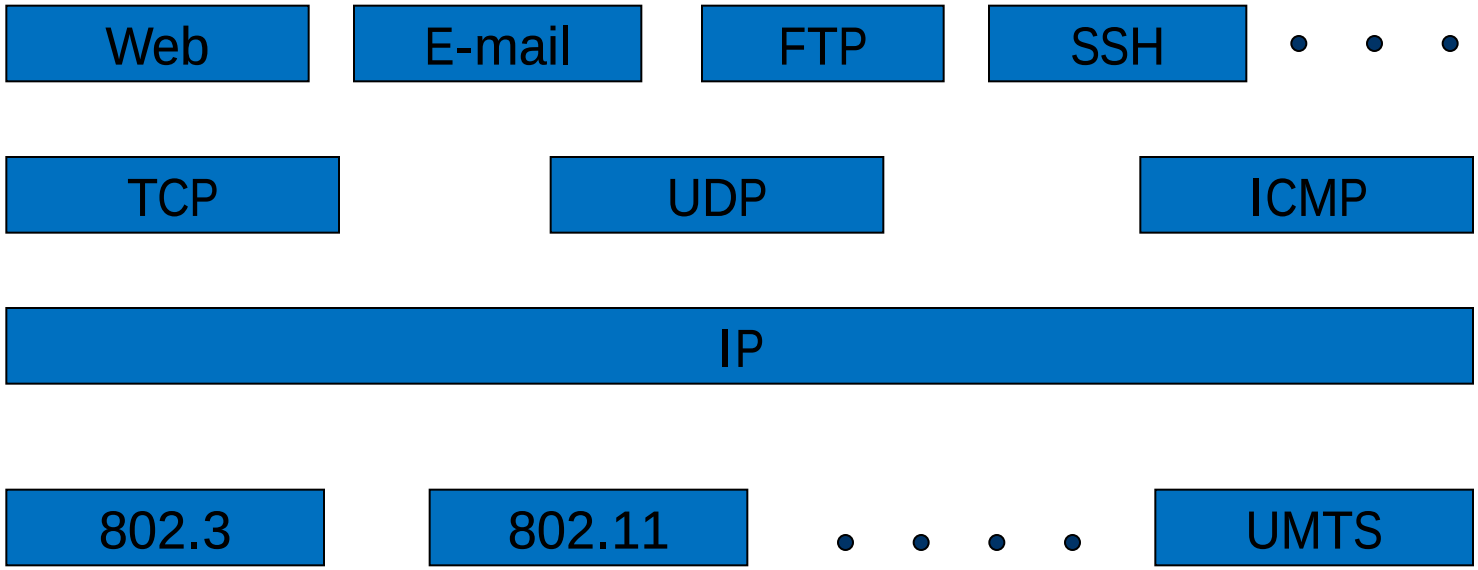


Layering

- Decomposition of a system's functionality into layers :
  - Simple, more universally usable functions more at the bottom
  - More complex and specific functions higher up
- Each layer represents an abstraction of lower layers.
- Each layer provides an interface that can be used by higher layers.



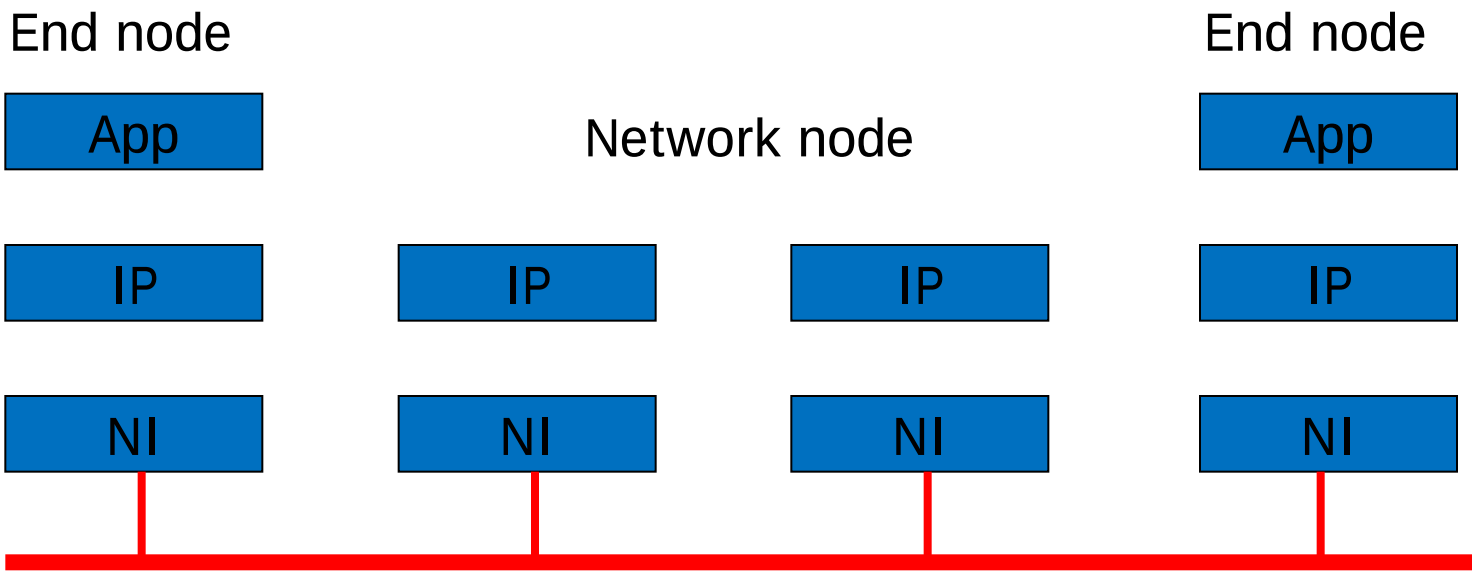
Example for Layering: Internet Protocol Stack



End-to-End Principle (E2E) (Internet)

- End-to-end-Principle:

  „A function of service should be carried out within a general layer only if it is needed by all clients of that layer and if it can be completely implemented in that layer.“

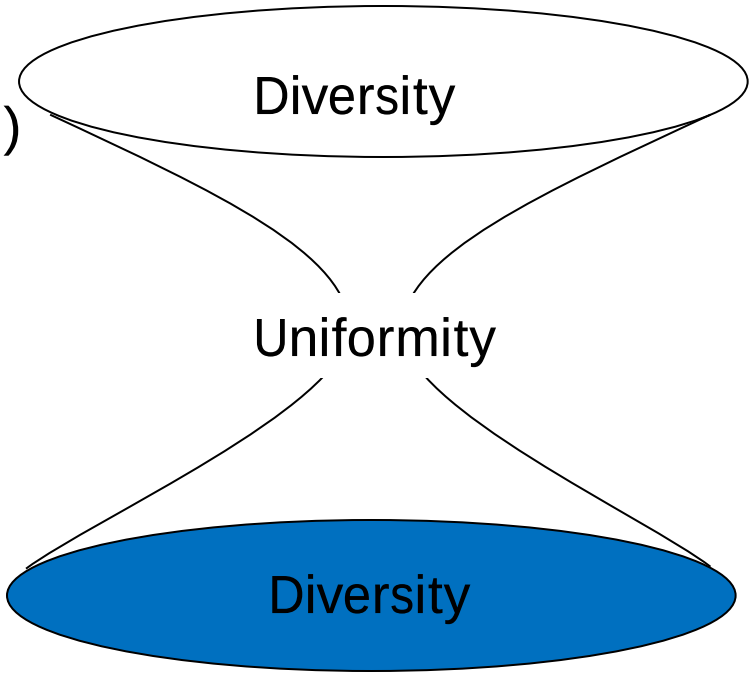


Hourglass-Architecture (Internet)

Various Applications
(file transfer, media streaming, web, email, VoIP)

Internet protocol (IP)

Various networks, media, and signal representations

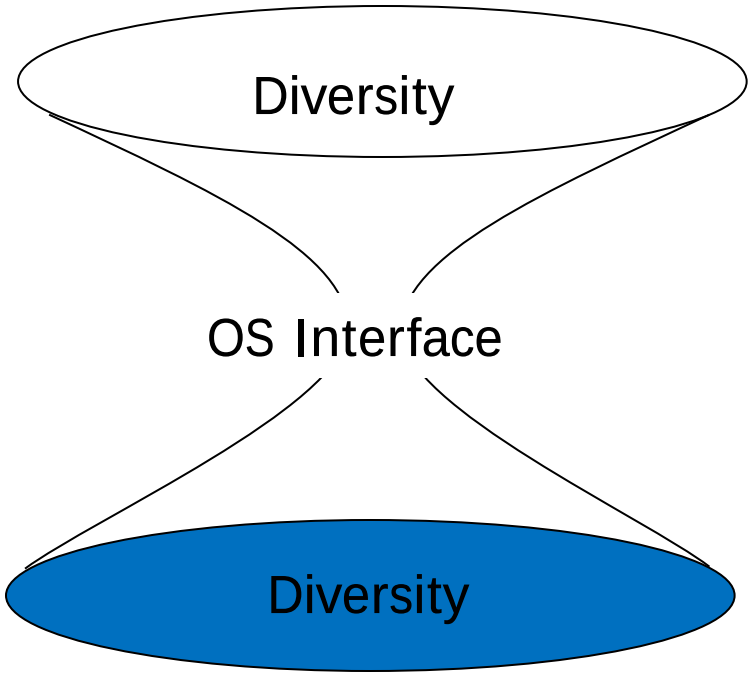


Hourglass-Architecture in OS

Various applications

OS programming interface

Various Hardware



E2E in the OS context: application neutrality

- A stable, universal programming interface should be provided.
- OS for a universal computer should be application neutral.
- Support for specific requirements should be placed in the topmost layer possible.

Separation of policy and mechanism

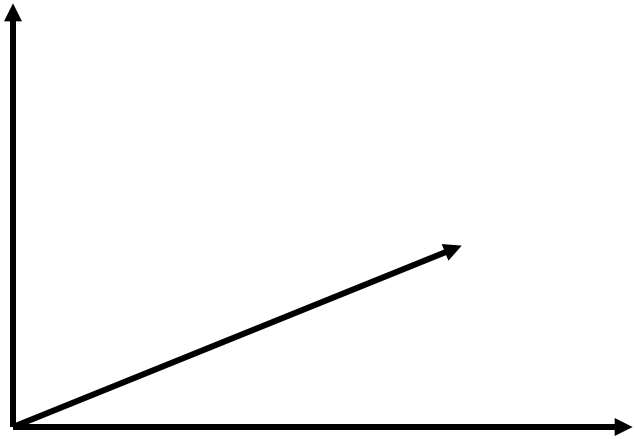
- Application neutrality sometimes requires compromises:

  Lower layers may provide mechanisms that can be parameterized in higher layers to suit specific application requirements.

- Examples:
  - Scheduling
  - Paging
  - Security

Orthogonality

- Functions and concepts of an OS should be independent of each other.
- Each component should exhibit orthogonal design criteria.
- Orthogonality means freedom of combination.



SPOT-Rule (Single Point of Truth)

- No copies or repetitions
- For code: each functionality is implemented exactly once.
- For data: each piece of information that is managed by the system has exactly one representation.

- Usage of SPOT avoids inconsistencies.

Further reading

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