

COMP3121 ASSIGNMENT1 QUESTION1

1. You are given an array A of n distinct integers.

- (a) You have to determine if there exists a number (not necessarily in A) which can be written as a sum of squares of two distinct numbers from A in two different ways (note: $m^2 + n^2$ and $n^2 + m^2$ counts as a single way) and which runs in time $n^2 \log n$ in the **worst case** performance. Note that the brute force algorithm would examine all quadruples of elements in A and there are $\binom{n}{4} = O(n^4)$ such quadruples. (10 points)
- (b) Solve the same problem but with an algorithm which runs in the **expected time** of $O(n^2)$. (10 points)

Answer:

- (A) First, we create and initialize array B and size of $\frac{n(n-1)}{2}$ which to store the all result for pair. $O(\frac{n(n-1)}{2}) = O(n^2)$

Then each element in A , we need go through all $\frac{n(n-1)}{2}$ pairs $(A[K], A[M])$, $K < M$, of **distinct** integers in A . And for every pair we need compute $(O(1) \cdot n^2)$ and store the result $(A[K]^2 + A[M]^2)$ in array $B(O(1) \cdot n^2)$. It will take $O(n^2) + O(n^2) + O(n^2) = O(n^2)$ time.

Then we sort array B . because the size of b is $\frac{1}{2}(n^2 - n) < n^2$ we use n^2 for further prove for convenient. We use merge sort can sort in $O(n \log n)$ in worst case. because the size of n should be n^2 in our case. It become $O(n^2 \log n^2) = O(2n^2 \log n) = O(n^2 \log n)$

Then for each element in B , we use binary search that whether same element in B . If find it should return true if not return false. In the worst case it will go through every element in B and do binary search every time. It will take $O(n^2 \log n)$.

So worst case should take $O(n^2 \log n) + O(n^2 \log n) + O(n^2) = O(n^2 \log n)$

- (B) First, we create and initialize array B and size of $\frac{n(n-1)}{2}$ which to store the all result for pair. $O(\frac{n(n-1)}{2}) = O(n^2)$

Then each element in A , we need go through all $\frac{n(n-1)}{2}$ pairs $(A[K], A[M])$, $K < M$, of distinct integers in A . And for every pair we need compute $(O(1) \cdot n^2)$ and store the result $(A[K]^2 + A[M]^2)$ in array $B(O(1) \cdot n^2)$. It will take $O(n^2) + O(n^2) + O(n^2) = O(n^2)$ time.

Then we do not sort and use hash table to check if a same element exist in the array: each insertion and lookup takes $O(1)$ expected time.

Then we hash all elements of B and then go through elements of B again, this time for each element in B checking in $O(1)$ time at whether appear multiple in the corresponding slot of the hash table. If appear multiple it is true, not are false. This will be done in expected time of $O(n^2)$.