Influential coalitions for Boolean Functions

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Abstract

For $f: \{0,1\}^n \to \{0,1\}$ and $S \subset \{1,2,\ldots,n\}$, let $J_S^+(f)$ be the probability that, for x uniform from $\{0,1\}^n$, there is some $y \in \{0,1\}^n$ with f(y) = 1 and $x \equiv y$ off S. We are interested in estimating, for given $\mathbb{E}(f)$ and m, the least possible value of $\max\{J_S^+(f): |S| = m\}$.

A theorem of Kahn, Kalai, and Linial (KKL) gave some understanding of this issue and led to several stronger conjectures. Here we improve the positive consequences of the KKL Theorem and disprove a pair of conjectures from the late 80s, as follows.

- (1) The KKL Theorem implies that there is a fixed $\alpha>0$ so that if $\mathbb{E} f \approx 1/2$, and c>0, then there is a set S of size at most αcn with $J_S^+(f) \geq 1 n^{-c}$. We show that for every $\delta>0$ there is an f with $\mathbb{E} f \approx 1/2$ and $J_S^+(f) \leq 1 n^{-C}$ for every S of size $(1/2 \delta)n$, where $C = C_\delta$. This disproves a conjecture of Benny Chor from 1989.
- (2) The KKL Theorem also implies that there are fixed, positive c and δ such that for any f with $\mathbb{E}f \geq n^{-c}$ there is some S of size $(1/2 \delta)n$ with $J_S^+(f) > 0.9$. We improve this, showing that for every C > 0 there is some $\delta = \delta(C) > 0$ such that if $\mathbb{E}f \geq n^{-C}$ then there is a set S of size $(1/2 \delta)n$, with $J_S^+(f) > 0.9$.
- (3) We also show that for fixed $\delta > 0$ there are $c, \alpha > 0$ and Boolean functions f such that $\mathbb{E}f > \exp[-n^{1-c}]$ and $J_S^+(f) \le \exp[-n^{\alpha}]$ for each S of size $(1/2 \delta)n$. This disproves a conjecture of the third author from the late 80s.

1 Introduction

For a set T we use $\Omega(T)$ for the discrete cube $\{0,1\}^T$ and μ_T for the uniform probability measure on $\Omega(T)$. In this paper f will always be a

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Boolean function on $\Omega([n])$ (that is, $f:\Omega([n]) \to \{0,1\}$, where, as usual, $[n] = \{1,\ldots,n\}$). We write μ for $\mu_{[n]}$ and usually use \mathbb{E} for expectation with respect to μ (e.g. $\mathbb{E}f = \mathbb{E}f = \mu\{x: f(x) = 1\}$). We reserve x,y for elements of $\Omega([n])$ and set $|x| = \sum x_i$.

Following Ben-Or and Linial [2] we define, for a given f and $S \subset [n]$, the influence of S toward one to be

$$I_S^+(f) = \mu_{[n] \setminus S}(\{u \in \Omega([n] \setminus S) : \exists v \in \Omega(S), f(u, v) = 1\}) - \mathbb{E}f.$$
 (1)

Similarly, the influence of S toward zero is

$$I_S^-(f) = \mu_{[n] \setminus S}(\{u \in \Omega([n] \setminus S) : \exists v \in \Omega(S), f(u, v) = 0\}) - (1 - \mathbb{E}f) \tag{2}$$

and the (total) influence of S is

$$I_S(f) = I_S^+(f) + I_S^-(f).$$

Suppose $\mathbb{E}f=1/2$. It then follows from a theorem of Kahn, Kalai and Linial [10] (Theorem 2.2 below, henceforth "KKL") that for every $a\in(0,1)$ there is an $S\subset[n]$ of size an with $I_S^+(f)\geq 1/2-n^{-c}$, where c>0 depends on a. (See Theorem 2.3.) Benny Chor conjectured in 1989 that one can in fact achieve $I_S^+(f)\geq 1/2-c^n$ (where, again, c<1 depends on a). The conjecture has been "in the air" since that time, though as far as we know it has appeared in print only in [11, 13].

In this paper we disprove Chor's conjecture and another, similar conjecture from the same period. On the other hand, we improve the preceding consequence of KKL.

For our purposes the subtracted terms in (1) and (2) are mostly a distraction, and it sometimes seems clearer to speak of $J_S^+(f) := I_S^+(f) + \mathbb{E}f$ and $J_S^-(f) := I_S^-(f) + (1 - \mathbb{E}f)$. Thus, for example, $J_S^+(f)$ is the probability that a uniform setting of the variables in $[n] \setminus S$ doesn't force f = 0, and Chor's conjecture predicts an S with $J_S^+(f) \ge 1 - c^n$. The following statement shows that this need not be the case.

Theorem 1.1. For any fixed $\alpha, \delta \in (0,1)$ there are a C and an f with $\mathbb{E}f = \alpha$ and $J_S^+(f) < 1 - n^{-C}$ for every $S \subseteq [n]$ of size $(1/2 - \delta)n$.

We should note that one cannot expect to go much beyond $|S| = (1/2 - \delta)n$; for example if $\mathbb{E}f = 1/2$, then it follows from the "Sauer-Shelah Theorem" (Theorem 2.5) that there is an S of size n/2 with $J_S^+(f) = 1$.

Another consequence of KKL (see Theorem 2.4 below) is that there is a $\beta > 0$ such that for any f with $\mathbb{E}f > n^{-\beta}$ there is an S of size (say) 0.1n with

influence 1 - o(1). A conjecture of the second author, again from the late 80s, asserts that the same conclusion holds even assuming only $\mathbb{E}f > (1-\varepsilon)^n$ for sufficiently small ε . This conjecture turns out to be false as well:

Theorem 1.2. For any fixed $\varepsilon, \delta > 0$ there are an $\alpha > 0$ and Boolean functions f such that $\mathbb{E}f > (1 - \varepsilon)^n$ and no set of size $(1/2 - \delta)n$ has influence to 1 more than $\exp[-n^{\alpha}]$.

This can be strengthened a bit to require $\mathbb{E}f > \exp[-n^{1-c}]$ for some fixed $c = c_{\delta} > 0$.

While the preceding, rather optimistic conjectures turn out to be false, we do show that the first of the aforementioned consequences of KKL can be improved:

Theorem 1.3. For each C > 0, there is a $\delta > 0$ such that for any f with $\mathbb{E}f > n^{-C}$, there is an $S \subset [n]$ of size at most $(1/2 - \delta)n$ with $|\pi_I(A)| > .9$.

(Of course, as elsewhere in this discussion, ".9" could be any preset $\rho < 1$.)

Note that the gap between Theorems 1.2 and 1.3 is substantial and our modest progress is likely not the final word on the problem. For example, could it be that there is some fixed β such that there are f's with $\mathbb{E}f > n^{-\beta}$ for which no S of size 0.1n has influence $\Omega(1)$? We will discuss this question further in the next section.

The examples proving Theorems 1.1 and 1.2 are given in Section 3. Each of these is of the form $f = \wedge_{i=1}^m C_i$, where the C_i 's are random \vee 's of k literals using k distinct variables (henceforth "k-clauses"). These f's, which may be thought of as variants of the "tribes" construction of Ben-Or and Linial (see below), were inspired by a paper of Ajtai and Linial [1] and share with it the following curious feature. It's easy to see that any f can be converted to a monotone (i.e. increasing) f' with $\mathbb{E}(f') = \mathbb{E}f$ and each influence (I_S^+) and so on) for f' no larger than the corresponding influence for f; thus it's natural to look for f's with small influences among the increasing functions. But the present random examples, like those of [1], do not do this, and it's not easy to see what one gets by monotonizing them.

The proof of Theorem 1.3 is given in Section 4. The argument goes roughly as follows. We employ two strategies, both variants of the analysis in [10]. The first (described in Section 4.1) uses the total influence, assumed sufficiently large. If at some point this total influence becomes "small," we switch to a different procedure (Sections 4.2 and 4.3) that combines the incremental argument from [10] with the Sauer-Shelah lemma. Perhaps the main novelty is in combining harmonic analysis in the spirit of [10] with more purely combinatorial ingredients.

2 Background and perspective

Influence

We write $I_{\ell}(f)$ for $I_{\{\ell\}}(f)$. A form of the classic edge isoperimetric inequality for Boolean functions is

Theorem 2.1. For any (Boolean function) f with $\mathbb{E}f = t$,

$$I(f) := \sum_{k=1}^{n} I_{\ell}(f) \ge 2t \log_2(1/t). \tag{3}$$

(This convenient version is easily derived from the precise statement, due to Hart [7]; see also [9, Sec. 7] for a simple inductive proof.)

While (3) is exact or close to exact (depending on t), it typically gives only a weak lower bound on the maximum of the $I_{\ell}(f)$'s, namely

$$\max_{\ell} I_{\ell}(f) \ge 2t \log_2(1/t)/n. \tag{4}$$

For t not too close to 0 or 1, the following statement from [10] gives better information.

Theorem 2.2 (KKL). There is a fixed c > 0 such that for any f with $\mathbb{E}f = t$, there is an $\ell \in [n]$ with

$$I_{\ell}(f) \ge ct(1-t)\log n/n. \tag{5}$$

Recall that $J_{\ell}(f) = \mathbb{E}f + I_{\ell}(f)$. Repeated application of Theorem 2.2 gives the following two corollaries.

Theorem 2.3. For all $a, t \in (0,1)$ there is a c such that for any f with $\mathbb{E}f = t$ there is an $S \subseteq [n]$ with $|S| \le an$ and

$$J_S^+(f) \ge 1 - n^{-c}$$

(that is,
$$I_S^+(f) \ge (1-t) - n^{-c}$$
).

Similarly (either by the same argument or by applying Theorem 2.3 to the function 1 - f(x)) there is a small S' with $J_{S'}(f) \ge 1 - n^{-c}$ (i.e. $I_{S'}(f) \ge t - n^{-c}$), and combining these observations we find that there is in fact a small S'' (e.g. $S \cup S'$) with $I_{S''}(f) \ge 1 - n^{-c}$.

Theorem 2.4. For every $\delta, \epsilon > 0$, there is an $\alpha > 0$ such that for large enough n and any f with $\mathbb{E}f \geq n^{-\alpha}$, there is an $S \subseteq [n]$ with $|S| = \delta n$ and

$$J_S^+(f) \ge 1 - \epsilon$$
.

The conjecture of Chor stated in Section 1 asserts that the n^{-c} in Theorem 2.3 can be replaced by something exponential in n, and the conjecture stated before Theorem 1.2 proposes a similar weakening of the $n^{-\alpha}$ lower bound on $\mathbb{E}f$ in Theorem 2.4. As already noted, we will show below that these conjectures are incorrect.

Tribes

The original "tribes" examples of Ben-Or and Linial [2] are Boolean functions of the form $f = \bigvee_{i=1}^{m} C_i$, where the "tribes" C_i are \land 's of k (distinct) variables and each variable belongs to exactly one tribe. The dual of such an f (so "dual tribes") is $g = \bigwedge_{i=1}^{m} D_i$, where D_i is the \lor of the variables in C_i (so again, each variable belongs to exactly one D_i).

When $k = \log n - \log \log n - \log \ln(1/t)$, we have $1 - \mathbb{E}f = \mathbb{E}(g) \approx t$ (where $\log = \log_2$ and f, g are as above). For fixed $t \in (0, 1)$ both constructions show that Theorem 2.2 is sharp (up to the value of c).

On the other hand, when $t = O(n^{-c})$ for a fixed c > 0, f shows that (4) is tight up to a multiplicative constant, depending on c; for example, $k = 2\log n - \log\log n$ gives $\mathbb{E}f \approx 1/(2n)$ and $I_{\ell}(f) \approx 2\log n/n^2 = \Theta(\mathbb{E}f\log(1/\mathbb{E}f)/n)$ for each ℓ . (In contrast, for $\mathbb{E}(g) \approx 1/n$, we should take $k = \log n - 2\log\log n - 1$, in which case $I_{\ell}(g) = \Theta(\log^2 n/n^2)$ and (4) is off by a log.)

For f (again, as above) with $\mathbb{E}f \in (\Omega(1), 1 - \Omega(1))$, there are sets of size $\log n$ with large influence towards 1, while no set of size $o(n/\log n)$ has influence $\Omega(1)$ towards 0. (The corresponding statement with the roles of 0 and 1 reversed holds for g.) The Ajtai-Linial construction mentioned in the introduction shows that there are Boolean functions h with $\mathbb{E}(h) \approx 1/2$ and $I_S(h) < o(1)$ for every S of size $o(n/\log^2 n)$.

Trace

We now briefly consider influences from a different point of view. For a set X let $2^X = \{S : S \subset X\}$, $\binom{X}{k} = \{S \subset X : |S| = k\}$, $\binom{X}{< k} = \{S \subset X : |S| < k\}$ and $\binom{n}{< k} = \sum_{i=0}^{k-1} \binom{n}{i}$. For $\mathcal{F} \subset 2^X$ and $Y \subset X$, the trace of \mathcal{F} on Y is

$$\mathcal{F}_{|Y} = \{ S \cap Y : S \in \mathcal{F} \}.$$

Let X = [n]. The following "Sauer-Shelah Theorem" determines, for every n and m, the minimum T such that for each $\mathcal{F} \subseteq 2^X$ of size T there is some $Y \in \binom{X}{m}$ on which the trace of \mathcal{F} is complete, meaning $\mathcal{F}_{|Y} = 2^Y$. Such a Y is said to be shattered by \mathcal{F} .

Theorem 2.5 (The Sauer-Shelah Theorem). If $\mathcal{F} \subset 2^X$ and $|\mathcal{F}| > \binom{n}{< r}$, then \mathcal{F} shatters some $Y \in \binom{X}{r}$.

That this is sharp is shown by $\mathcal{F} = {X \choose < r}$, the *Hamming ball* of radius r-1 about \emptyset with respect to the usual Hamming metric on $2^X \equiv \Omega(X)$.

Theorem 2.5 was proved around the same time by Sauer [14], Shelah and Perles [15], and Vapnik and Chervonenkis [16]. It has many connections, applications and extensions in combinatorics, probability theory, model theory, analysis, statistics and other areas.

We identify $\Omega([n])$ and $2^{[n]}$ in the usual way. The connection between traces and influences is as follows. Let f be a Boolean function on $\Omega([n])$ and $\mathcal{F} = f^{-1}(1)$. It is easy to see that for $S \subseteq [n]$ and $T = [n] \setminus S$,

$$J_S^+(f) = 2^{-|T|} |\mathcal{F}_{|T}|.$$

Thus, in the language of traces, we are interested in the effect of relaxing " \mathcal{F} shatters Y" to require only that $\mathcal{F}_{|Y}$ contain a large fraction of 2^Y .

The following arrow notation (e.g. [4, 6]) is convenient. Write $(N, n) \to (M, r)$ if every $\mathcal{F} \subseteq 2^{[n]}$ of size N has a trace of size at least M on some $S \in \binom{[n]}{r}$; for example the Sauer-Shelah Theorem says $\binom{n}{r}+1, n \to (2^r, r)$.

One might hope that Hamming balls would again give the best examples in our relaxed setting, which would say, for example, that for $m \leq n$,

$$\left(\binom{n}{< r} + 1, n \right) \to \left(\binom{m}{< r} + 1, m \right). \tag{6}$$

But (6), which was first considered by Bollobás and Radcliffe [3] and would have implied both of the conjectures disproved here, was shown in [3] to be false for fixed r and (large) m = n/2. (For r = n/2 and m = n - 1, it fails for the original tribes example discussed above.)

A consequence of (6) is that for fixed $\delta, \epsilon > 0$ and large r,

$$(\binom{n}{<(1+\epsilon)r/2},n) \to ((1-\delta)2^r,r),$$

which would imply our second conjecture from the introduction. Here a counterexample with $n \gg r$ was given by Kalai and Shelah [12], but this seems not very relevant to present concerns, for which the regime of interest has n a little smaller than 2r.

Two problems

Question 2.6. For fixed $\alpha, \delta > 0$, what is the largest $t \in (0, 1/2)$ for which one can find Boolean functions f with $\mathbb{E}f = t$ and $I_S^+(f) < \alpha$ for every $S \subseteq [n]$ of size $(1/2 - \delta)n$?

As far as we know $t > n^{-\beta}$ (with β depending on α, δ) is possible. The influence of sets of half the variables is of special interest:

Question 2.7. Given $\mathbb{E}f \ll 1$ what can be said about the maximum of $J_S^+(f)$ for |S| = n/2? What is the smallest t such that for each f with $\mathbb{E}f = t$ there is some S of size n/2 with $J_S^+(f) \geq 1/2$?

3 Boolean functions without influential coalitions

In each construction we consider, for suitable k and m, $f = \bigwedge_{i=1}^{m} C_i$, where the C_i 's are random \vee 's of k literals using k distinct variables (henceforth "k-clauses") and show that f is likely to have the desired properties. Every C_i can be regarded as a list of specifications for the values of k variables. We use g_i for the specification associated with C_i , and write $C_i \sim x$ if some entry of x agrees with g_i . We say C_i misses $S \subseteq [n]$ if the indices of all variables in C_i lie in $[n] \setminus S$.

Let $s = (1/2 - \delta)n$. We will always use S for an s-subset of [n] and (for such an S) set $m_S = |\{i : C_i \text{ misses } S\}|$. (Following common practice we omit irrelevant floor and ceiling symbols, pretending all large numbers are integers. As in the case of k, m and s, parameters not declared to be constants are assumed to be functions of n.) We use log for \log_2 .

Both constructions will make use of the next two observations, with Theorem 1.1 following immediately from these and Theorem 1.2 requiring a little more work.

Lemma 3.1. If
$$k = o(\sqrt{n})$$
 and $(1/2 + \delta)^k m = \omega(n)$ then w.h.p.

$$m_S \sim (1/2 + \delta)^k m \ \forall S \in {[n] \choose s}.$$
 (7)

(where, as usual, $a_n \sim b_n$ means $a_n/b_n \to 1$ and with high probability (w.h.p.) means with probability tending to 1, both as $n \to \infty$).

Proof. For a given S, m_S has the binomial distribution B(m,p), with $p = \binom{n-s}{k}/\binom{n}{k} \sim (1/2+\delta)^k$ (using $k = o(\sqrt{n})$ for the " \sim "). Thus $\mathbb{E}m_S = mp$ and, by "Chernoff's Inequality" (e.g. [8, Theorem 2.1]),

$$\Pr(m_S \notin ((1-\zeta)mp, (1+\zeta)mp)) < \exp[-\Omega(\zeta^2 mp)],$$

for $\zeta \in (0,1)$. Applying this with a ζ which is both $\omega(\sqrt{n/(mp)})$ and o(1) gives $\Pr(m_S \not\sim mp) < 2^{-\omega(n)}$, and the union bound then gives (7).

The next lemma is stated to cover both applications, though nothing so precise is needed for Theorem 1.1.

Lemma 3.2. If there is a ξ for which

$$\exp[-\xi^2 n] = o((1 - 2^{-k})^m) \tag{8}$$

and

$$[(1+2\xi)/4]^k = o(1/m), \tag{9}$$

then w.h.p.

$$\mathbb{E}f \sim (1 - 2^{-k})^m. \tag{10}$$

Proof. This is a simple second moment method calculation (similar to what's done in [1], though described differently there).

Recalling that x, y always denote elements of $\{0, 1\}^n$, write A_x for the event $\{f(x) = 1\}$ and $\mathbf{1}_x$ for its indicator, and set $X = \sum \mathbf{1}_x = 2^n \mathbb{E} f$. Then $\Pr(A_x) = (1 - 2^{-k})^m$ and $\mathbb{E} X = (1 - 2^{-k})^m 2^n$; so we just need to show $\mathbb{E} X^2 \sim \mathbb{E}^2 X$ (equivalently, $\mathbb{E} X^2 < (1 + o(1))\mathbb{E}^2 X$), since Chebyshev's Inequality then gives $\Pr(|X - \mathbb{E} X| > \zeta \mathbb{E} X) = o(1)$ for any fixed $\zeta > 0$.

We have

$$\mathbb{E}X^2 = \sum_x \sum_y \mathbb{E}\mathbf{1}_x \mathbf{1}_y = \sum_x \Pr(A_x) \sum_y \Pr(A_y | A_x),$$

so will be done if we show that for a fixed x,

$$\sum_{y} \Pr(A_y | A_x) < (1 + o(1))(1 - 2^{-k})^m 2^n.$$

Since the sum is the same for all x, it's enough to prove this when $x = \underline{0}$. Set $Z = \{y : |y| < (1/2 - \xi)n\}$ and recall that by Chernoff's Inequality, $|Z| < \exp[-2\xi^2 n]2^n$. It is thus enough to show that (for $x = \underline{0}$)

$$y \notin Z \Rightarrow \Pr(A_y|A_x) < (1+o(1))(1-2^{-k})^m,$$
 (11)

since then, using (8), we have

$$\sum_{y} \Pr(A_y|A_x) < |Z| + \sum_{y \notin Z} \Pr(A_y|A_x) < (1 + o(1))(1 - 2^{-k})^m 2^n.$$

Now since $x = \underline{0}$, we have $A_x = \{g_i \neq \underline{1} \ \forall i\}$; so if, for a given $y \notin Z$, we set $\beta = \beta_y = \Pr(C_i \sim y | g_i \neq \underline{1})$ (a function of |y|), then $\Pr(A_y | A_x) = \beta^m$. Aiming for a bound on β , we have

$$1 - 2^{-k} = \Pr(C_i \sim y)$$

$$= \Pr(g_i = \underline{1}) \Pr(C_i \sim y | g_i = \underline{1}) + \Pr(g_i \neq \underline{1}) \Pr(C_i \sim y | g_i \neq \underline{1})$$

$$= 2^{-k} \Pr(C_i \sim y | g_i = \underline{1}) + (1 - 2^{-k})\beta$$

and

$$\Pr(C_i \sim y | g_i = \underline{1}) > 1 - (1 - |y|/n)^k \ge 1 - (1/2 + \xi)^k =: 1 - \nu$$

(using the fact that if $g_i = \underline{1}$, then $g_i \nsim y$ iff all indices of variables in C_i belong to $\{j : y_j = 0\}$). Combining, we have

$$\beta < (1 - 2^{-k})^{-1} [1 - 2^{-k} - 2^{-k} (1 - \nu)] = (1 - 2^{-k}) \left[1 + \frac{2^{-k} (\nu - 2^{-k})}{(1 - 2^{-k})^2} \right],$$

which with (9) gives $\beta^m < (1 + o(1))(1 - 2^{-k})^m$ (which is (11)).

Proof of Theorem 1.1. Notice that it's enough to prove this with $\mathbb{E}f \sim \alpha$ (rather than "= α "); for then, since $f^{-1}(1) \subseteq g^{-1}(1)$ trivially implies $J_S^+(f) \leq J_S^+(g)$ for all S, we can choose $\beta \in (\alpha, 1)$ and a g with $\mathbb{E}(g) \sim \beta$ possessing the desired small influences, and shrink $g^{-1}(1)$ to produce f.

Let $k = C \log n$, with $C = C_{\delta}$ chosen so that $(1 + 2\delta)^k = \omega(n)$ (e.g. $C = 1/\delta$ does this), and $m = 2^k \ln(1/\alpha) = n^C \ln(1/\alpha)$. Here all we use from Lemma 3.1 (whose hypotheses are satisfied for our choice of k and m) is the fact that w.h.p. $m_S \neq 0$ for all S, whence each $J_S^+(f)$ is at most $1 - 2^{-k} = 1 - n^{-C}$. On the other hand, by Lemma 3.2 (with, for example, $\xi = 0.1$), we have $\mathbb{E}f \sim \alpha$ w.h.p. So w.h.p. f meets our requirements.

Proof of Theorem 1.2. Here, intending to recycle n, m and f, we rename these quantities \mathbf{n} , \mathbf{m} and \mathbf{f} . We may of course assume δ is fairly small. Let (for example) $\xi = \delta/3$, fix ε with $0 < \varepsilon < \xi^2$, and set $k = (1 + \delta) \log \mathbf{n}$ and $\mathbf{m} = \varepsilon 2^k \mathbf{n}$. These values are easily seen to give the hypotheses of Lemmas 3.1 and 3.2. In particular, we can say that w.h.p. the supports of the C_i 's are chosen so (7) holds (note this says nothing about the values specified by the g_i 's) and

$$\mathbb{E}(\mathbf{f}) \sim (1 - 2^{-k})^{\mathbf{m}} \sim e^{-\varepsilon \mathbf{n}}.$$
 (12)

Set $n = (1/2 + \delta)\mathbf{n}$. Fix $S \in {[\mathbf{n}] \choose s}$, set $m = m_S$, and let $f = f_S$ be the \wedge of the m C_i 's—w.l.o.g. C_1, \ldots, C_m —that miss S. Thus f is the \wedge of $m \sim (1/2 + \delta)^k \mathbf{m} = \varepsilon (1 + 2\delta)^k \mathbf{n}$ random k-clauses from a universe of n variables. Theorem 1.2 (with $\alpha = \delta$) thus follows from

Claim A.
$$\Pr(\mathbb{E}f > \exp[-n^{\delta}]) < o(2^{-n})$$

(since then w.h.p. we have $\mathbb{E}(f_S) \leq \exp[-n^{\delta}]$ for every S).

Remarks. The actual bound in Claim A will be $\exp[-\Omega(m)]$, so much smaller than 2^{-n} . Note that here it doesn't matter whether we take μ to be our original measure (i.e. μ uniform on $\{0,1\}^n$) or uniform measure on $Q:=\{0,1\}^n$; but it's now more natural to think of the latter—and we will do so in what follows—since our original universe plays no further role in this discussion. It may also be worth noting that, unlike in the proof of Lemma 3.2, the second moment method is not strong enough to give the exponential bound in Claim A.

Claim B. If $X \subseteq Q$, $\mu(X) = \beta > \exp[-o(n/\log^2 n)]$ and $\zeta = o(2^{-k})$, then for a random k-clause C,

$$\Pr(\mu(C \land X) > (1 - \zeta)\mu(X)) < 1/2$$

(where
$$C \wedge X = \{x \in X : C \sim x\}$$
).

Remark. This is probably true for β greater than something like $\exp[-n/k]$. The bound in the claim is just what the proof gives, and is more than enough for us since we're really interested in much larger β .

To see that Claim B implies Claim A, set $f_j = \wedge_{i=1}^j C_i$ and notice that $\mathbb{E} f \geq \beta$ implies (for example)

$$|\{i : \mathbb{E}(f_i) < (1 - \frac{5\ln(1/\beta)}{m})\mathbb{E}(f_{i-1})\}| < m/5$$
 (13)

(and, of course, $\mathbb{E}(f_i) \geq \beta$ for all i). But if we take $\beta = \exp[-n^{\delta}]$ then our choice of parameters gives

$$\zeta := 5m^{-1}\ln(1/\beta) = o(2^{-k})$$

(using $m/\ln(1/\beta)=\Theta(n^{(1+\delta)\log(1+2\delta)+1-\delta})$ and $2^k=n^{1+\delta}$), so Claim B bounds the probability of (13) by

$$\binom{m}{m/5} 2^{-4m/5} = o(2^{-\mathbf{n}}).$$

Proof of Claim B. Let G be the bipartite graph on $Q \cup W$, where W is the set of (n-k)-dimensional subcubes of Q and, for $(x,D) \in Q \times W$, we take $x \sim D$ if $x \in D$. (So we've gone to complements: for a clause C the corresponding subcube is $D = \{y : C \not\sim y\}$, so $C \wedge X = X \setminus D$.)

Assuming Claim B fails at X, fix $T \subseteq W$ with |T| = |W|/2 and

$$D \in T \implies \mu(D \cap X) < \zeta \mu(X),$$

and set $R = W \setminus T$.

Consider the experiment: (i) choose x uniformly from X; (ii) choose D uniformly from the members of W containing x; (iii) choose y uniformly from D.

Claim C.
$$\Pr(y \in X) > (2 - o(1))\beta$$
.

Proof. Since each triple (x, D, y) with $x \in X$ and $x, y \in D$ is produced by (i)-(iii) with probability $|X|^{-1} \cdot 2^k |W|^{-1} \cdot 2^{k-n}$, we just need to show that the number of such triples with $y \in X$ is at least

$$(2 - o(1))\beta |X||W|2^n 2^{-2k} = (2 - o(1))|X|^2 |W|2^{-2k}.$$

Writing d for degree in G, we have

$$\sum_{x \in X} d_T(x) = \sum_{D \in T} d_X(D) < |T|\zeta|X|,$$

implying

$$\sum_{D \in R} d_X(D) = \sum_{x \in X} (d(x) - d_T(x))$$

$$> |X||W|2^{-k} - \zeta|T||X| = (1 - o(1))|X||W|2^{-k}.$$

The number of (x, D, y)'s as above is thus

$$\sum_{D \in W} d_X^2(D) \ge \sum_{D \in R} d_X^2(D) \ge (\sum_{D \in R} d_X(D))^2 / |R|$$

$$> (1 - o(1))|X|^2 |W|^2 |R|^{-1} 2^{-2k} = (2 - o(1))|X|^2 |W| 2^{-2k}.$$

Let T(x) be the random element of Q gotten from x by choosing K uniformly from $\binom{[n]}{k}$ and randomly (uniformly, independently) revising the x_i 's with $i \notin K$. Then y gotten from x by (ii) and (iii) above is just T(x),

so the next assertion contradicts Claim C, completing the proof of Claim B (and Theorem 1.2).

Claim D. If $\mu(X) > \exp[-o(n/\log^2 n)]$ and x is uniform from X, then $\Pr(T(x) \in X) < (1 + o(1))\mu(X)$.

Remark. If X is a subcube of codimension n/k, say $X = \{x : x \equiv 0 \text{ on } L\}$ with |L| = n/k, then for any $x \in X$,

$$\Pr(T(x) \in X) = \sum_{t} \Pr(|K \cap L| = t) 2^{-(|L| - t)} = \mu(X) \sum_{t} \Pr(|K \cap L| = t) 2^{t},$$

and, since $|K\cap L|$ is essentially Poisson with mean 1, the sum is approximately $e^{-1}\sum_t 2^t/t! = e$. So Claim D fails for $\mu(X) = 2^{-n/k}$ and, as earlier, it's natural to guess that it holds if $\mu(X)$ is much bigger than this.

Proof of Claim D. Let $Q_r = \{y \in Q : |y| \le r\}$. The assumption on $\mu(X)$ implies that $\mu(Q_{r-1}) < \mu(X) \le \mu(Q_r)$ for some r > (1/2 - o(1/k))n, so Claim D follows from

Claim E. If $\varphi = o(1/k)$ and $r > (1/2 - \varphi)n$, then for any $x \in Q$ and $X \subseteq Q$ with

$$\mu(X) \le \mu(Q_r),\tag{14}$$

we have

$$\Pr(T(x) \in X) < (1 + o(1))\mu(Q_r).$$
 (15)

Proof. We may assume $x = \underline{0}$, so that $\Pr(T(x) = y)$ is a decreasing function of |y|. We thus maximize $\Pr(T(x) \in X)$ subject to (14) by taking $X = Q_r$, and (15) is then a routine calculation using

$$\mu(X) = \Pr(\operatorname{Bin}(n, 1/2) \le r)$$

and

$$\Pr(T(x) \in X) = \Pr(\operatorname{Bin}(n-k, 1/2) \le r)$$

(where $Bin(\cdot, \cdot)$ denotes a binomially distributed r.v.).

Influential coalitions for Boolean functions 4

4.1 Use of the total influence

Let

$$f: \{0,1\}^n \to \{0,1\} \text{ with } \mathbb{E}[f] > n^{-C_0} (C_0 > 1).$$

Write $f = \sum \hat{f}(S)w_S$ and $I_j = ||f|_{\varepsilon_j=1} - f|_{\varepsilon_j=0}||_1$.

Let $C^{(1)} = C^{(1)}(C_0)$ be a sufficiently large constant. Assume

$$\sum I_j = \sum |S| |\hat{f}(S)|^2 > C^{(1)} \log n. \mathbb{E}[f].$$
 (16)

Then

$$I_j > C^{(1)} \frac{\log n}{n} \mathbb{E}[f]$$
 for some j .

Replace f by $f_1 = \pi_{\hat{j}}(f)$ obtained by projection on $\{1, \ldots, n\} \setminus \{j\}$. Hence

$$\mathbb{E}[f_1] = \mu \left[[f|_{\varepsilon_j = 0} = 1] \cup [f|_{\varepsilon_j = 1} = 1] \right] > \mathbb{E}[f] + \frac{1}{2} I_j$$
$$> \left(1 + \frac{C^{(1)} \log n}{2n} \right) \mathbb{E}[f].$$

If f_1 again satisfies (16), repeat the construction.

Either one obtains $\tilde{f} = \pi_{\hat{B}_0}(f), \mathbb{E}[\tilde{f}] > \frac{9}{10}$ after at most

$$|B_0| < \frac{3C_0}{C^{(1)}}n < \frac{n}{10}$$

steps, or (16) fails after $k_1 < \frac{3C_0}{C^{(1)}}n$, before achieving this.

We distinguish 2 cases Case 1. $k_1 \leq \frac{10^{-10}}{C^{(1)}}n$.

We then switch to a different strategy for further amplification of $\pi_{\hat{B}_0}(f)$ that will be described in Sections 2 and 3. Case 2. $k_1 > \frac{10^{-10}}{C^{(1)}}n$.

Note that

$$\mathbb{E}[\pi_{\hat{B}_0}(f)] > \left(1 + \frac{C^{(1)} \log n}{2n}\right)^{k_1} \mathbb{E}[f] \gtrsim e^{\frac{1}{2 \cdot 10^{10}} \log n} \mathbb{E}[f]$$

$$\geq n^{\frac{1}{2 \cdot 10^{10}}} \mathbb{E}[f].$$

Replace $C^{(1)}$ by $C^{(2)} = \frac{C^{(1)}}{10^{50}C_0}$ and repeat the preceding.

Hence $k_2 < \frac{3C_0}{C^{(2)}}n$. If $k_2 < \frac{10^{-10}}{C^{(2)}}n$ we switch to the §2, §3 procedure. If $k_2 \geq \frac{10^{-10}}{C^{(2)}} n$, we gained another factor $\frac{1}{n^{2.10^{10}}}$.

Hence, after at most $r = O(C_1)$ steps, we obtain $\mathbb{E}[\pi_{\hat{B}_0}(f)] > \frac{9}{10}$ for some $B_0 \subset \{1, \ldots, n\}$ satisfying

$$|B_0| < 3C_0n\left(\frac{1}{C^{(1)}} + \frac{1}{C^{(2)}} + \dots + \frac{1}{C^{(r)}}\right) < \frac{n}{4}$$

(where $C^{(1)}$ is chosen to ensure $C^{(r)} > 100C_0$, hence $\log C^{(1)} \sim C_0 \log C_0$) unless at some earlier stage, we switched to the amplification strategy from Sections 2 and 3 applied to $\pi_{\hat{B}_0}(f)$, where $B_0 \subset \{1, \ldots, n\}$ satisfies now

$$|B_0| < C_0 n \left(\frac{3}{C^{(1)}} + \dots + \frac{3}{C^{(\rho-1)}} \right) + \frac{10^{-10}}{C^{(\rho)}} n < \frac{2.10^{-10}}{C^{(\rho)}} n.$$
 (17)

Denoting again $f = \pi_{\hat{B}_0}(f)$, it satisfies

$$\sum |S| |\hat{f}(S)|^2 \le C^{(\rho)}(\log n) \mathbb{E}[f]$$
(18)

and proceed with the amplification using a different method described next. Set $C_2 = C^{(\rho)}$.

4.2 Second strategy: Preparations

Assume

$$\sum |S| |\hat{f}(S)|^2 \le C_2 \log n \, \mathbb{E}[f]. \tag{19}$$

Set $\delta = 10^{-8}C_2^{-1}$ and let $A \subset \{1, \dots, n\} \setminus B_0$ be a random set of size δn . Let C_3 be another parameter and write

$$\sum_{|S\cap A|>10^{-6}\log n} |\hat{f}(S)|^2 \le \frac{1}{C_3 \log n} \sum_{|S|>C_3 \log n} |S| \, |\hat{f}(S)|^2 + \sum_{\substack{|S|\le C_3 \log n \\ |S\cap A|>10^{-6}\log n}} |\hat{f}(S)|^2$$

$$\stackrel{(2.1)}{<} \frac{C_2}{C_3} \mathbb{E}[f] + \frac{10^6}{\log n} \sum_{|S|\le C_3 \log n} |S\cap A|.|\hat{f}(S)|^2.$$

Taking expectation in A, we get an estimate

$$\left(\frac{C_2}{C_3} + 10^6 \delta C_3\right) \mathbb{E}[f] < 10^3 \sqrt{\delta C_2} \,\mathbb{E}[f] < \frac{1}{10} \mathbb{E}[f] \tag{20}$$

for appropriate choice of C_3 .

Hence

$$\sum_{|S \cap A| \le 10^{-10} \log n} |\hat{f}(S)|^2 > \frac{9}{10} \mathbb{E}[f].$$

Write

$$\{1,\ldots,n\}\setminus B_0 = A \cup A', \varepsilon = (x,x') \in \{0,1\}^A \times \{0,1\}^{A'}.$$

Define

$$g = \sum_{|S \cap A| \le 10^{-6} \log n} \hat{f}(S) w_S = \sum_{\substack{T \subset A \\ |T| \le 10^{-6} \log n}} \left[\sum_{S \cap A = T} \hat{f}(S) w_{S \cap A'} \right] w_T$$

and

$$\Omega = \left\{ x' \in \{0, 1\}^{A'}; \|g_{x'}\|_2^2 > \frac{1}{2} \|f_{x'}\|_2^2 \right\}.$$

Then

$$||g1_{\Omega^C}||_2^2 \le \frac{1}{2}\mathbb{E}_{x'}[||f_{x'}||_2^2] = \frac{1}{2}\mathbb{E}[f]$$

and

$$||f1_{\Omega}||_2^2 \ge ||g1_{\Omega}||_2^2 > \frac{9}{20} \mathbb{E}[f].$$

Fix $x' \in \Omega$ and write $f_{x'}(x) = \sum_{T} \widehat{f_{x'}}(T) w_T(x)$. Then

$$||f_{x'}||_{2}^{2} = ||f_{x'}||_{\frac{3}{2}}^{\frac{3}{2}} \ge \left(\sum_{T} |\widehat{f}_{x'}(T)|^{2} 2^{-|T|}\right)^{\frac{3}{4}}$$

$$> 2^{-\frac{3}{4}10^{-6}\log n} ||g_{x'}||_{2}^{\frac{3}{2}}$$

$$> \frac{1}{2}n^{-\frac{3}{4}10^{-6}} ||f_{x'}||_{2}^{\frac{3}{2}}$$

so that

$$\mathbb{E}[f_{x'}] > \frac{1}{16} n^{-310^{-6}} \text{ for } x' \in \Omega.$$
 (21)

Also

$$||f|\Omega||_1 > \frac{9}{20}\mathbb{E}[f] > \frac{9}{20}n^{-C_0}.$$
 (22)

We fix A, Ω and replace f by $f|\Omega$. Hence

$$||f_{x'}||_1 \gtrsim n^{-3.10^{-6}} \text{ of } f_{x'} \neq 0$$
 (23)

nd

$$\frac{|\Omega|}{\mathbb{E}[f]} < n^{3.10^{-6}}.\tag{24}$$

Second strategy: Iteration

Write f redefined above, we start another iterative construction with selection of coordinates from A. Fix $x' \in \Omega$ and set $F = f_{x'} = \sum_{T \subset A} \hat{F}(T) w_T$. Assume $||F||_1 < \frac{9}{10}$. We distinguish 2 cases.

Case I.

$$\sum_{j \in A} I_j(F) = \sum_{T \subset A} |\hat{F}(T)|^2 |T| > 10^{-3} (\log n) ||F||_2^2.$$

Case II.

$$\sum_{j \in A} I_j(F) \le 10^{-3} \log n ||F||_2^2.$$

In case II, write using hypercontractivity

$$\sum_{j \in A} I_{j}(F) \sim \sum_{j \in A} \|F|_{\varepsilon_{j}=1} - F|_{\varepsilon_{j}=0}\|_{\frac{3}{2}}^{\frac{3}{2}}$$

$$\geq \sum_{j \in A} \left(\sum_{\substack{T \subset A \\ j \in T}} |\hat{F}(T)|^{2} 2^{-|T|}\right)^{\frac{3}{4}}$$

$$\geq 2^{-\frac{\log n}{50}} \left(\sum_{\substack{T \subset A \\ 0 < |T| < \frac{\log n}{50}}} |\hat{F}(T)|^{2}\right) \cdot \frac{1}{\max_{j \in A} (\sum_{\substack{T \subset A \\ j \in T}} |\hat{F}(T)|^{2})^{\frac{1}{4}}}$$

$$\gtrsim n^{-\frac{1}{50}} \left(\frac{19}{20} \mathbb{E}[F] - \mathbb{E}[F]^{2}\right) \frac{1}{\max_{j \in A} I_{j}(F)^{\frac{1}{4}}}$$

$$\gtrsim \frac{1}{10} n^{-\frac{1}{50}} \frac{10^{3}}{\log n} \left(\sum_{j \in T} I_{j}(F)\right) \frac{1}{\max_{j \in A} I_{j}(F)^{\frac{1}{4}}}.$$

Therefore, there is some $j = j_{x'} \in A$ such that

$$I_j(f_{x'}) \gtrsim \frac{n^{-\frac{4}{50}}}{(\log n)^4} > n^{-\frac{1}{12}}.$$
 (25)

Partition

$$\Omega' \equiv \left\{ x' \in \Omega; \|f_{x'}\|_1 < \frac{9}{10} \right\} = \Omega_I \cup \Omega_{II}$$

according to x' is Case I, II.

Hence

$$\sum_{j \in A} I_j(f|\Omega_I) > 10^{-3} \log n ||f|\Omega_I||_2^2$$

and we choose $j \in A$ such that

$$I_j(f|\Omega_I) > \frac{10^{-3} \log n}{|A|} ||f|\Omega_I||_2^2.$$
 (26)

For $x' \in \Omega_{II}$, choose $j = j_{x'} \in A$ for which (25) holds. Hence

$$\mathbb{E}[\pi_{\widehat{j_{x'}}}(f)1_{\Omega_{II}}] = \mathbb{E}_{x'}[\mathbb{E}_x[\pi_{\widehat{j_{x'}}}(f_{x'})]1_{\Omega_{II}}] > \left(1 + \frac{1}{2n^{-\frac{1}{12}}}\right) ||f1_{\Omega_{II}}||_2.$$
 (27)

Set $j_{x'} = j$ for $x' \in \Omega_I$, so that also by (26)

$$\mathbb{E}[\pi_{\widehat{j}_{x'}}(f)1_{\Omega_I}] > \left(1 + \frac{10^{-3}\log n}{|A|}\right) ||f1_{\Omega_I}||_2^2.$$
 (28)

¿From (27), (28)

$$\mathbb{E}[\pi_{\widehat{j_{x'}}}(f)1_{\Omega'}] > \left(1 + \frac{10^{-3}\log n}{|A|}\right) ||f1_{\Omega'}||_2^2.$$
 (29)

Replace f by f_1 defined by

$$\begin{cases} (f_1)_{x'} = f_{x'} \text{ if } \mathbb{E}[f_{x'}] \ge \frac{9}{10} \\ (f_1)_{x'} = \pi_{\widehat{j}_{x'}}(f_{x'}) \text{ if } \mathbb{E}[f_{x'}] < \frac{9}{10}. \end{cases}$$

By (29)

$$|\Omega| \ge \mathbb{E}[f_1] > \mathbb{E}[f] + \frac{10^{-3} \log n}{|A|} \mathbb{E}\left[f 1_{\|f_{x'}\|_1 < \frac{9}{10}}\right]$$
 (30)

and repeat the process described in Section 3 to f_1 .

We terminate when

$$\mathbb{E}\left[f1_{\|f_{x'}\|_{1}<\frac{9}{10}}\right]<\frac{1}{2}\|f\|_{1}.\tag{31}$$

By (30), this will happen after at most k steps, with

$$\left(1 + \frac{10^{-3}\log n}{2|A|}\right)^k \mathbb{E}[f] \le |\Omega|.$$

Therefore, by (24)

$$k < 3.10^{-6} \frac{2|A|}{10^{-3}} < 6.10^{-3}|A|.$$
 (32)

The Boolean function \tilde{f} obtained satisfies by (31)

$$\mathrm{mes}_{x'} \Big[\|\tilde{f}_{x'}\|_1 \geq \frac{9}{10} \Big] \geq \mathbb{E} \big[\tilde{f} \, \mathbf{1}_{[\|\tilde{f}_{x'}\|_1 \geq \frac{9}{10}]} \big] \geq \frac{1}{2} \|\tilde{f}\|_1 \geq \frac{1}{2} \|f\| > \frac{1}{5} n^{-C_0}$$

by (22).

Also, $\tilde{f}_{x'} = \pi_{\hat{A}_{x'}}(f_{x'})$ where $A_{x'} \subset A$ is obtained as

$$A_{x'} = A^{(I)} \cup A_{x'}^{(II)}.$$

Here $A^{(I)}$ consists of the coordinates introduced in **Case I** and $A^{(II)}_{x'}$ in **Case II** alternative.

Note that by (25), a **Case II** coordinate corresponds to a measure increment $\sim n^{-\frac{1}{12}}$ in the x'-section, implying that

$$|A_{x'}^{(II)}| \lesssim n^{\frac{1}{12}}.$$
 (33)

Also

$$|A^{(I)}| < 6.10^{-3}|A|. (34)$$

Let $\ell \sim n^{\frac{1}{12}}$ satisfy $|A_{x'}^{(II)}| \leq \ell$ for all x'. Partition

$$\Omega = \bigcup_{\substack{V \subset A \\ |V| \le \ell}} \Omega_V$$

with

$$\Omega_V = \{ x' \in \Omega; A_{x'}^{(II)} = V \}.$$

One can then specify some V such that

$$\operatorname{mes}\left[\Omega_{V} \cap \left[\|\tilde{f}_{x'}\|_{1} > \frac{9}{10}\right]\right] > \frac{\frac{1}{5}n^{-C_{0}}}{\sum_{j \leq \ell} \binom{|A|}{j}} > e^{-n^{\frac{1}{11}}}.$$

At this point, invoke the Sauer-Shelah lemma to produce a subset. $A'' \subset A'$ satisfying

$$\pi_{A''} \left[\Omega_V \cap \left[\|\tilde{f}_{x'}\|_1 > \frac{9}{10} \right] \right] = \{0, 1\}^{A''}$$
 (35)

and

$$|A''| > \frac{|A'|}{2} - O(n^{\frac{1}{2} + \frac{1}{22}}). \tag{36}$$

Define

$$B_1 = (A' \backslash A'') \cup V \cup A^{(I)} \tag{37}$$

and

$$B = B_0 \cup B_1$$
 with B_0 the set in (17).

Hence

$$|B_1| < \frac{|A'|}{2} + O(n^{\frac{1}{2} + \frac{1}{22}}) + O(n^{\frac{1}{12}}) + 6.10^{-3}|A|$$

$$|B| < \frac{n}{2} - \left(\frac{1}{2} - 6.10^{-3}\right)|A| + O(n^{\frac{1}{2} + \frac{1}{22}}) + \frac{2.10^{-10}}{C_2}n$$

(by (17), since C_2 in (19) is $C^{(\rho)}$ from (17)).

$$<\frac{n}{2} + O(n^{\frac{1}{2} + \frac{1}{22}}) - \left(\frac{1}{2} - 6.10^{-3} - 2.10^{-2}\right)|A|$$

(since $|A| = 10^{-8}C_2^{-1}n = \delta n$)

$$< \left(\frac{1}{2} - \frac{\delta}{3}\right)n\tag{38}$$

where $\delta = \delta(C_0) > 0$. Next, for $\varepsilon = (x, x')$

$$\tilde{f}(\varepsilon) \equiv \max_{\pi_{B^c}(\gamma) = \pi_{B^c}(\varepsilon)} f(\gamma) = \max\{\pi_{\hat{B}_0}(f)(y', y); \pi_{A''}(y') = \pi_{A''}(x') \text{ and } \pi_{(A \setminus (V \cup A^I))}(y) = \pi_{(A \setminus (V \cup A^I))}(x)\}$$
(39)

$$\begin{array}{ll}
(3.11) & \max_{x \in \mathcal{A}(V \cup A^I)} F(y) \\
& = \pi_{(A \setminus (V \cup A^I))}(y) = \pi_{A \setminus (V \cup A^I)}(x)
\end{array} \tag{40}$$

where $F = [\pi_{\hat{B}_0}(f)]_{y'}$, for some $y' \in \Omega_V, \|\tilde{F}\|_1 > \frac{9}{10}$.

Since $\tilde{F}(x) \leq (39)$, it follows that $\|(39)\|_{L^1_x} > \frac{9}{10}$ and therefore $\|\tilde{f}\|_1 > \frac{9}{10}$. This completes the proof.

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