

# 자료구조 스터디 1차시(성능이란?)

자료구조를 공부하는 목적: 더 좋은 성능의 프로그램을 설계하고 개발하기 위함

∴ 자료구조를 공부하기에 앞서 프로그램의 성능을 분석하고 측정하는 방법에 대해 알아야한다.

※ 실제로 이론 시험의 대부분이 여기서 배운 내용을 토대로 나옴

## performance analysis vs performance measurement

- performance measurement
  - machine dependent: running times
  - Measure actual time on an actual computer
  - Use system clock
  - Measure running time in C

```
#include <time.h>

int main(void){
    clock_t start = clock();
    //do something
    double duration = ((double)(clock() - start)) / CLOCKS_PER_SEC;
}
```

```
#include <time.h>

int main(void){
    time_t start = time(NULL);
    //do something
    double duration = (double) difftime(time(NULL),start);
}
```

- Performance Analysis
  - Analyzing time and space complexity independent of the machine for the program
    - Time complexity: amount of computation time that a program needs to run to completion
    - Space complexity: amount of memory that a program needs to run to completion

## Space Complexity

- $S(P) = c + S_p(n)$ 
  - $S(P)$ : space requirement of program P
  - c: constant (for fixed space requirements)
    - Fixed part: space independent of the program's input/output size
      - Instruction space
      - space for simple variables
      - fixed-size structured variables

- constants
  - $n$  : Instance characteristics (ex: I/O size, number)
  - $S_p$ : space requirement of algorithm  $p$

- ```
float abc(float a, float b, float c){
    return a+b+b*c+(a+b-c)/(a+b)+4.00;
}
```

- $S_{abc}(n) = 0$

- ```
float sum(float list[],int n){
    float tempsum = 0;
    int i;
    for(i=0;i<n;i++)
        tempsum += list[i];
    return tempsum;
}
```

- $S_{sum}(n) = 4(n+3)$

- ```
float rsum(float list[],int n){
    if(n) return rsum(list,n-1)+list[n-1];
    return 0;
}
```

- $S_{rsum}(n) = 12(n+1)$

## Time complexity

- $T(P) = c + T_p(n)$ 
  - $T(P)$  : time taken by program  $P$
  - $c$  : compile time
  - $n$  : instance characteristics
  - $T_p$  : run time
- Divide the program into distinct steps
  - Assignment
  - comparison
  - return

- ```
float sum (float list[], int n){
    float tempsum = 0;           //1
    int i;                       //0
    for (i = 0; i < n; i++)       //n+1
        tempsum += list[i];      //n
    return tempsum;              //1
}
```

//total: 2n+3

- ```
float rsum(float list[], int n){
    if (n)                       //n+1
        return rsum(list, n-1) + list[n-1]; //n
    return list[0];              //1
}
```

//total: 2n+2

- ```

void add(int a[][MAX_SIZE],int b[][MAX_SIZE],int c[][MAX_SIZE]int row, int col){
    int i,j;                                //0
    for(i=0;i<rows;i++)                      //rows+1
        for(j=0;j<cols;j++)                 //rows*(cols+1)
            c[i][j] = a[i][j]+b[i][j];      //rows*cols
}

//total: 2rows*cols + 2rows + 1

```

- ```

count;

float sum (float list[], int n){
    float tempsum = 0; count++; /*for assignment*/
    int i;
    for (i = 0; i < n; i++){
        count++;                /*for the for loop*/
        tempsum += list[i]; count++; /*for assignment*/
    }
    count++; /*for last execution of for*/
    count++; /*for return*/
    return tempsum;
}

float rsum(float list[], int n){
    count++; /* for if conditional*/
    if (n){
        count++; /*for return*/
        return rsum(list, n-1) + list[n-1];
    }
    count ++; /*for return*/
    return list[0];
}

void add(int a[][MAX_SIZE],int b[][MAX_SIZE],int c[][MAX_SIZE]int row, int col){
    int i,j;
    for(i=0;i<rows;i++){
        count++; /*for i for loop*/
        for(j=0;j<cols;j++){
            count++; /*for j for loop*/
            c[i][j] = a[i][j]+b[i][j]; count++; /*for assignment*/
        }
        count++; /*for last execution of j for*/
    }
    count++; /*for last execution of i for*/
}

int main(void){
    float l[MAX_SIZE];
    int n,row,col,a[MAX_SIZE][MAX_SIZE],b[MAX_SIZE][MAX_SIZE],c[MAX_SIZE][MAX_SIZE];

    count = 0;
    sum(l,n);
    printf("%d\n",count);

    count = 0;
    rsum(l,n);
    printf("%d\n",count);

    count = 0;

```

```

add(a,b,c,row,col);
printf("%d\n",count);
}

```

## Asymptotic notation (O, Ω, Θ)

- To make meaningful statements about the complexity of the program's time and space
- $1(\text{constant}) < \log n < n(\text{linear}) < n \log n < n^2(\text{quadratic}) < n^3(\text{cubic}) < \dots < 2^n(\text{exponential})$
- big-O notation
  - $f(n) = O(g(n))$  (read as " $f$  of  $n$  is **big oh** of  $g$  of  $n$ ") iff  $\exists c, n_0 > 0, s. t. f(n) \leq cg(n) \forall n, n \geq n_0$
  - $\forall n, n \geq n_0, g(n)$  is upper bound on  $f(n)$
  - e.g:
    - $3n + 3 = O(n)$  as  $3n + 3 \leq 4n$  for  $n \geq 3$
    - $3n + 3 = O(n^2)$  as  $3n + 3 \leq 3n^2$  for  $n \geq 2$
    - $10n^2 + 4n + 2 =$
    - $6 * 2^n + n^2 =$
- Omega notation
  - $f(n) = \Omega(g(n))$  (read as " $f$  of  $n$  is **omega** of  $g$  of  $n$ ") iff  $\exists c, n_0 > 0, s. t. f(n) \geq cg(n) \forall n, n \geq n_0$
  - $\forall n, n \geq n_0, g(n)$  is lower bound on  $f(n)$
  - e.g:
    - $3n + 2 = \Omega(n)$  as  $3n + 2 \geq 3n$  for  $n \geq 1$
    - $10n^2 + 4n + 2 = \Omega(n^2)$  as  $10n^2 + 4n + 2 \geq n^2$  for  $n \geq 1$
    - $10n^2 + 4n + 2 = \Omega(n)$  as  $10n^2 + 4n + 2 \geq n$  for  $n \geq 1$
    - $6 * 2^n + n^2 =$
- Theta notation
  - $f(n) = \Theta(g(n))$  (read as " $f$  of  $n$  is **theta** of  $g$  of  $n$ ") iff  $\exists c_1, c_2, n_0 > 0, s. t. c_1g(n) \leq f(n) \leq c_2g(n) \forall n, n \geq n_0$
  - $\forall n, n \geq n_0, g(n)$  is both upper and lower bound on  $f(n)$
  - $f(n) = \Theta(g(n))$  means  $f(x) = O(g(x))$  and  $f(x) = \Omega(g(x))$
  - e.g:
    - $3n + 2 = \Theta(n)$  as  $3n \leq 3n + 2 \leq 4n$  for  $n \geq 2$
    - $10n^2 + 4n + 2 =$
    - $6 * 2^n + n^2 =$
    - $2 * n * m + 2 * n + 1 =$

## Recurrence equation

- Recurrence equation
  - $T(n) = T(n-1) + 1$
  - $T(0) = 1$
  - Solve  $T(n)$ 

$$\begin{aligned}
 T(n) &= T(n-1) + 1 \\
 &= (T(n-2) + 1) + 1 \\
 &= ((T(n-3) + 1) + 1) + 1 \\
 &\dots \\
 &= (\dots((T(0) + 1) + 1)\dots) + 1 = n + 1 = O(n)
 \end{aligned}$$

- Recurrence equation (assume that  $n = 2k$ )
  - $T(n) = T(n/2) + 1$
  - $T(1) = 1$
  - Solve  $T(n)$ , and compute big-oh function
 
$$T(n) = T(n/2) + 1$$

$$= (T(n/2^2) + 1) + 1$$

$$= ((T(n/2^3) + 1) + 1) + 1$$

$$\dots$$

$$= (\dots((T(n/2^k) + 1) + 1) \dots) + 1 = 1+k = 1+\log_2 n$$

$$= O(\log_2 n)$$

## 자료구조란?

- 데이터를 저장하고 관리하는 방식
- 예시: 전화번호를 저장하는 경우
  1. 전화번호를 알게된 순서대로 작성하기
  2. 사전순으로 작성하기
  3.  $\sim$ 까지 각 공간을 미리 크게 할당해두고, 그 공간에 저장하기
  4. 힙 구조로 저장하기
  5. 해시 구조로 저장하기
- 어떤 식으로 저장하냐에 따라 저장속도와 검색 속도, 알고리즘 실행속도, 저장 공간 등에 영향을 미침.

## 과제

### 과제1

score.txt로부터 이름, 국어,영어,수학 성적을 입력받아 이름과 평균을 result.txt파일과 화면에 출력하시오

- score.txt

```
박지원 100 100 100
박지투 85 90 80
이성희 60 70 80
이우진 98 73 85
순기 75 95 84
```

### 과제2

표준입력으로 정수  $n$ 을 입력받아  $n*n$  행렬을 2개 A,B 만들고 각 행렬을 0~20 사이의 랜덤한 수로 채우시오.  
행렬 A와 B를 출력하고 두 행렬의 합 C를 출력하시오.

### 과제3

$n$ 의 사이즈에 따른 selectionSort 실행 시간을 측정하시오.

[복잡도 참고 사항](#)

