

# LDPC Code

- The LDPC code, namely low density parity check code, is essentially a linear block code.
- LDPC codes provide a means to control errors in data transmissions over unreliable or noisy communication channels.
- LDPC can be use in both BSC and BEC channel.

# LDPC Code-Encoding

- LDPC code encode  $k$  message bit into  $n$  bit where  $n > k$ .
- $n - k$  is the size of Syndrome vector which is used for error correction and detection.
- LDPC code transfer  $k$ -bit into  $n$ -bit by Generator matrix( $G$ ).
- The dimension of Generator matrix is  $n \times k$ . This  $n \times k$  matrix multiply with

$k \times 1$  vector of message which will give  $n \times 1$  vector. That  $n \times 1$  vector is our codeword which will transfer through the channel.

$$C(n \times 1) =$$

$C$ =Codeword Matrix

$G$ =Generator Matrix

$M$ =Information Matrix

	$C_0$	$C_1$	$C_2$	$C_3$	$C_4$	$C_5$	$C_6$	$C_7$	$C_8$	$C_9$	$C_{10}$	$C_{11}$	$C_{12}$	$C_{13}$	$C_{14}$
0	1	0	0	0	0	0	0	1	0	0	1	1	1	0	1
1	0	1	0	0	0	0	0	1	1	0	0	1	1	1	0
2	0	0	1	0	0	0	0	0	1	1	1	0	0	0	1
3	0	0	0	1	0	0	0	1	0	1	1	1	0	0	0
4	0	0	0	0	1	0	0	0	1	0	1	1	1	0	0
5	0	0	0	0	0	1	0	0	0	1	0	1	1	1	0
6	0	0	0	0	0	0	1	0	0	0	1	0	1	1	1

$\underbrace{\hspace{10em}}_{\mathbf{I}} \quad \underbrace{\hspace{10em}}_{\mathbf{X}}$

# LDPC-Decoding

- LDPC code decode  $n$ -bit to check whether it have an error or not, It uses Parity-check matrix( $H$ ) for decoding.
- The dimension of  $H$  matrix is  $(n-k) \times k$ .
- It multiply the  $H$  matrix to the codeword which we have received over channel. it will the  $(n-k) \times 1$  matrix which is called Syndrome Matrix, if that matrix is zero matrix then we can say that there is no error in the codeword but if it is non-zero then there is error in the codeword.

$$S(n-k \times 1) = H(n-k \times n) * C(n \times 1)$$

$C$ =Codeword Matrix

$H$ =Parity check Matrix

$S$ =Syndrome Matrix

$$\mathbf{H} = \begin{bmatrix} 0 & 1 & 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 & 0 & 1 \end{bmatrix}$$

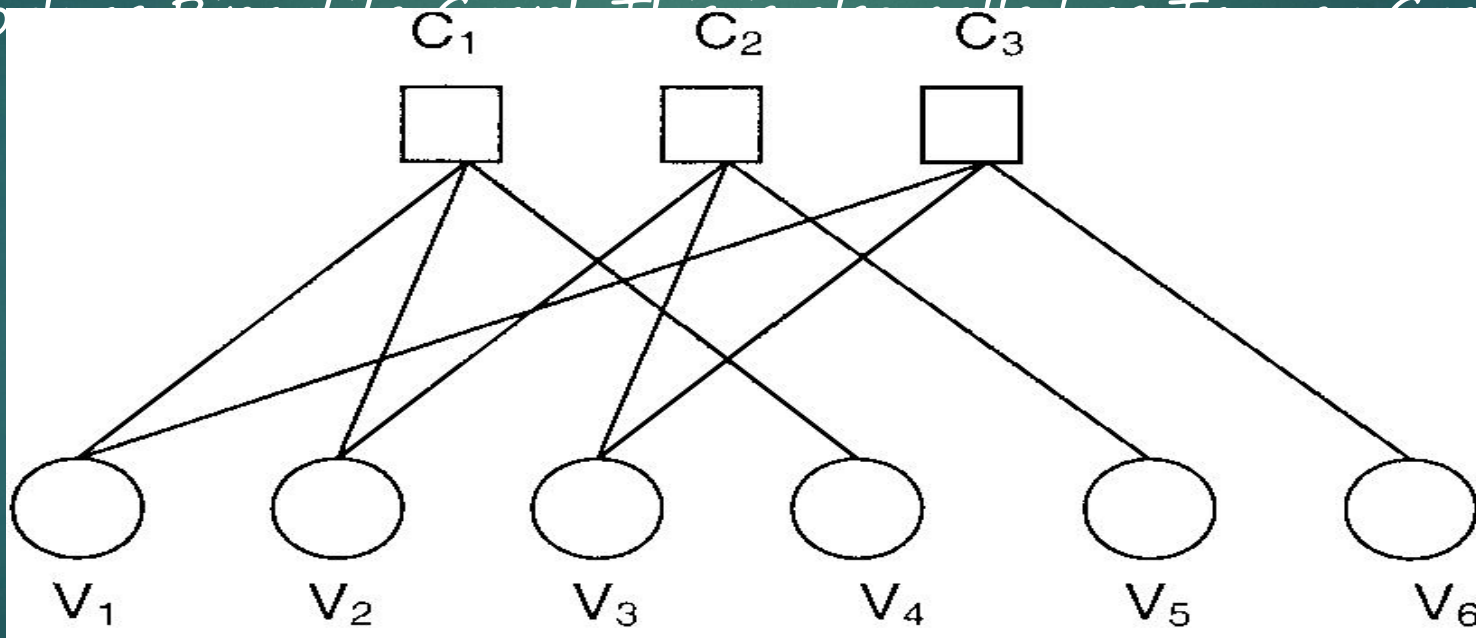


# Product Code

- Product Codes are a special case of LDPC code.
- Product code are restricted as the last rows and last columns of  $H$  matrix is for parity checking,
- If  $k$  i.e. #message bits increases then the size of parity check matrix increases squarely, but the parity check bits in  $n$ -bit message increases linearly.
- Hence the power of product code decreases as  $k$  increases and hence on a noisy channel if data is transmitted with large  $k$ .

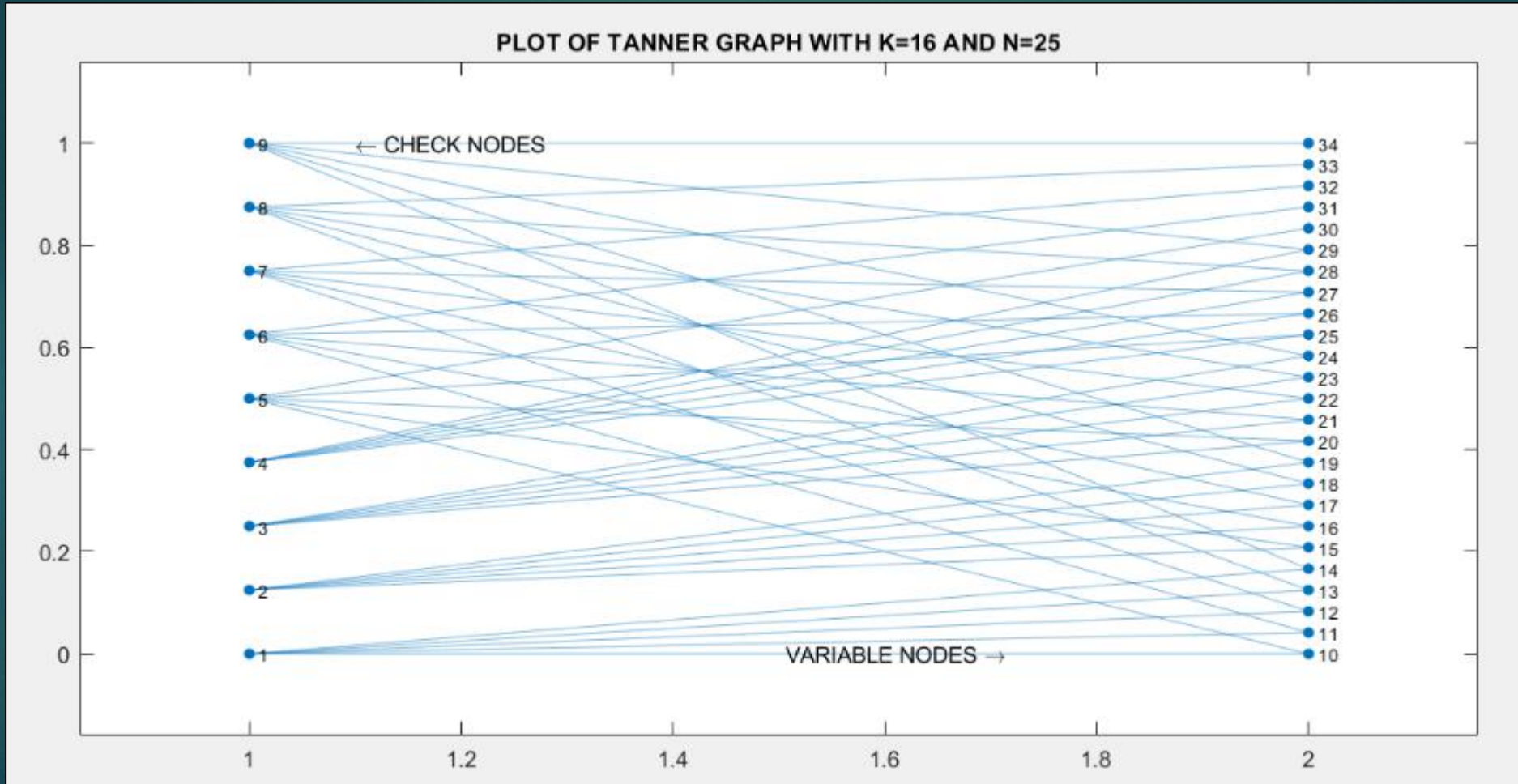
# Message Passing

- Message Passing is for correction the error. Here, All the bit of codeword is called Variable nodes(VN) and all the bit Syndrome is called check nodes(CN).
- Message can be passed from VN to CN and CN to VN. But message can not passed from VN to VN or CN to CN.
- It will work as Bipartite Graph. This is called as Tanner Graph.



# Tanner Graph

For  $n=25$  and  $k=16$

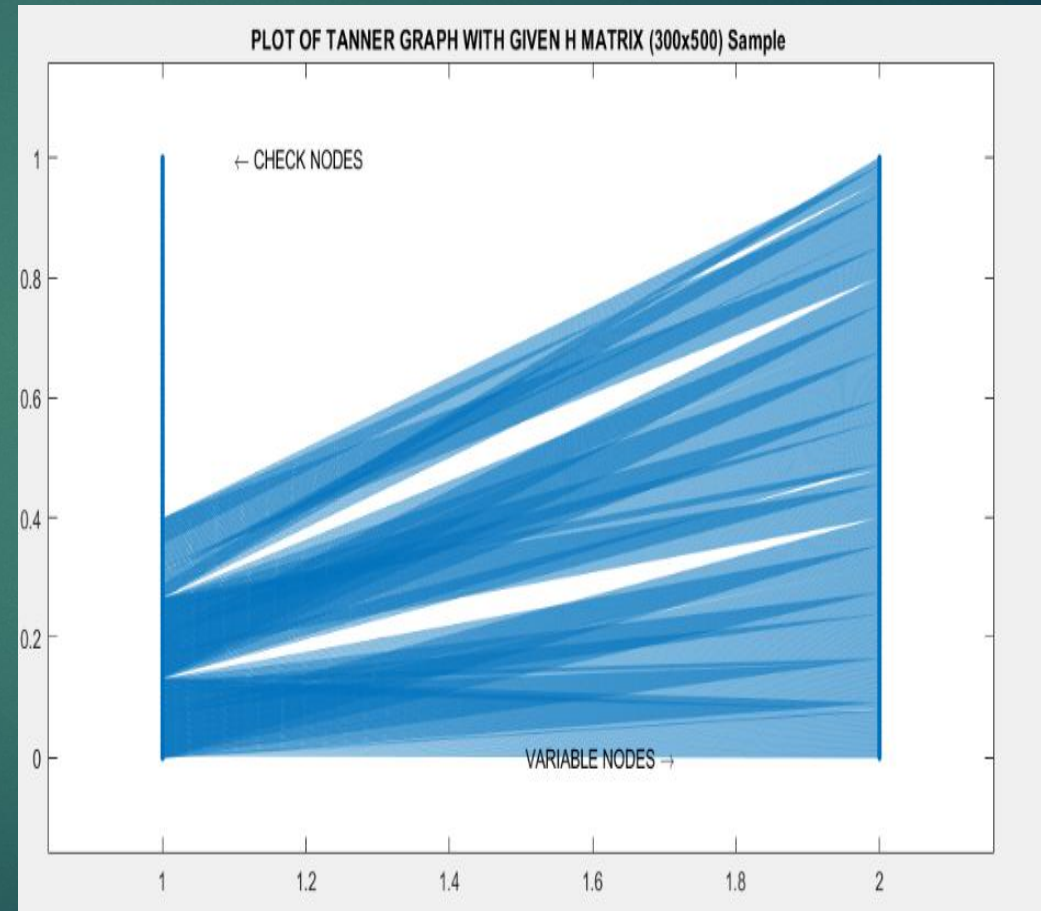
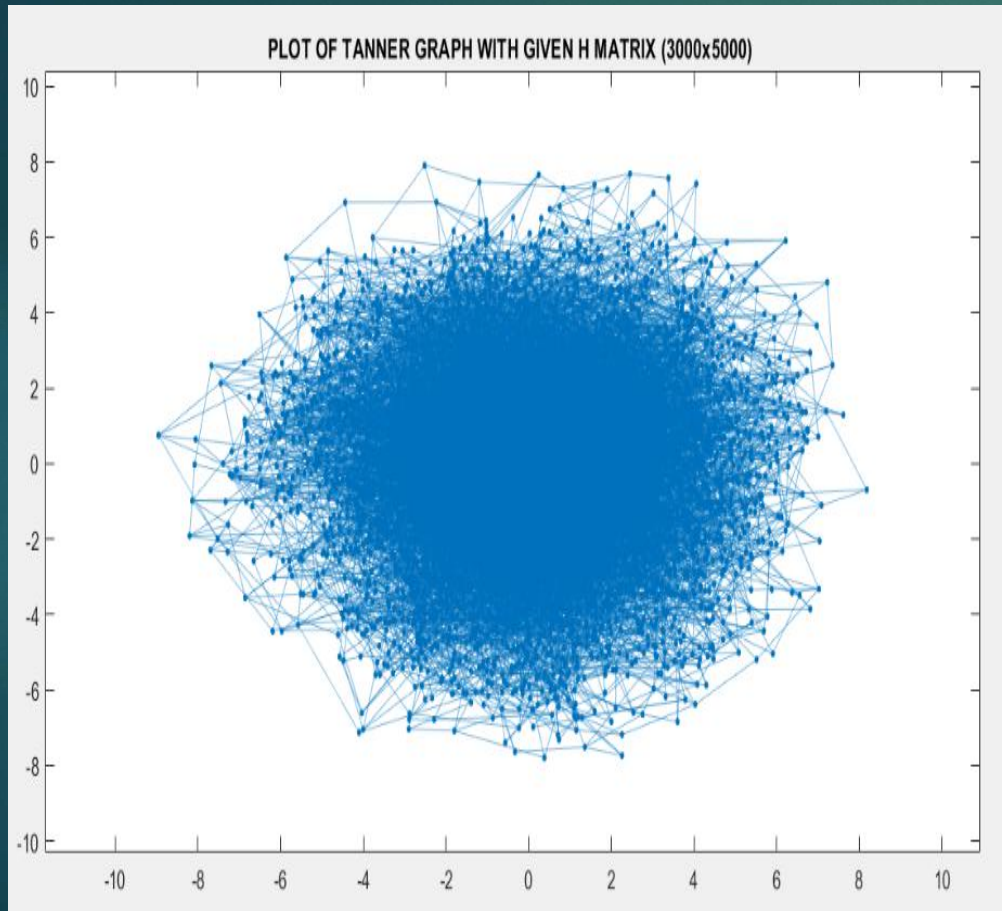




# Tanner Graph

For  $n=5000$  and  $k=2000$

for  $H(1:300,1:500)$



# Message Passing Decoding on Product Code(9,4) over BEC

```
Command Window
>> codeword_reshaped

codeword_reshaped =

     1     1     0     0     1     1     1     0     1

>> codeword_noisy

codeword_noisy =

     1    -1     0     0    -1    -1     1     0     1

>> decoded_codeword

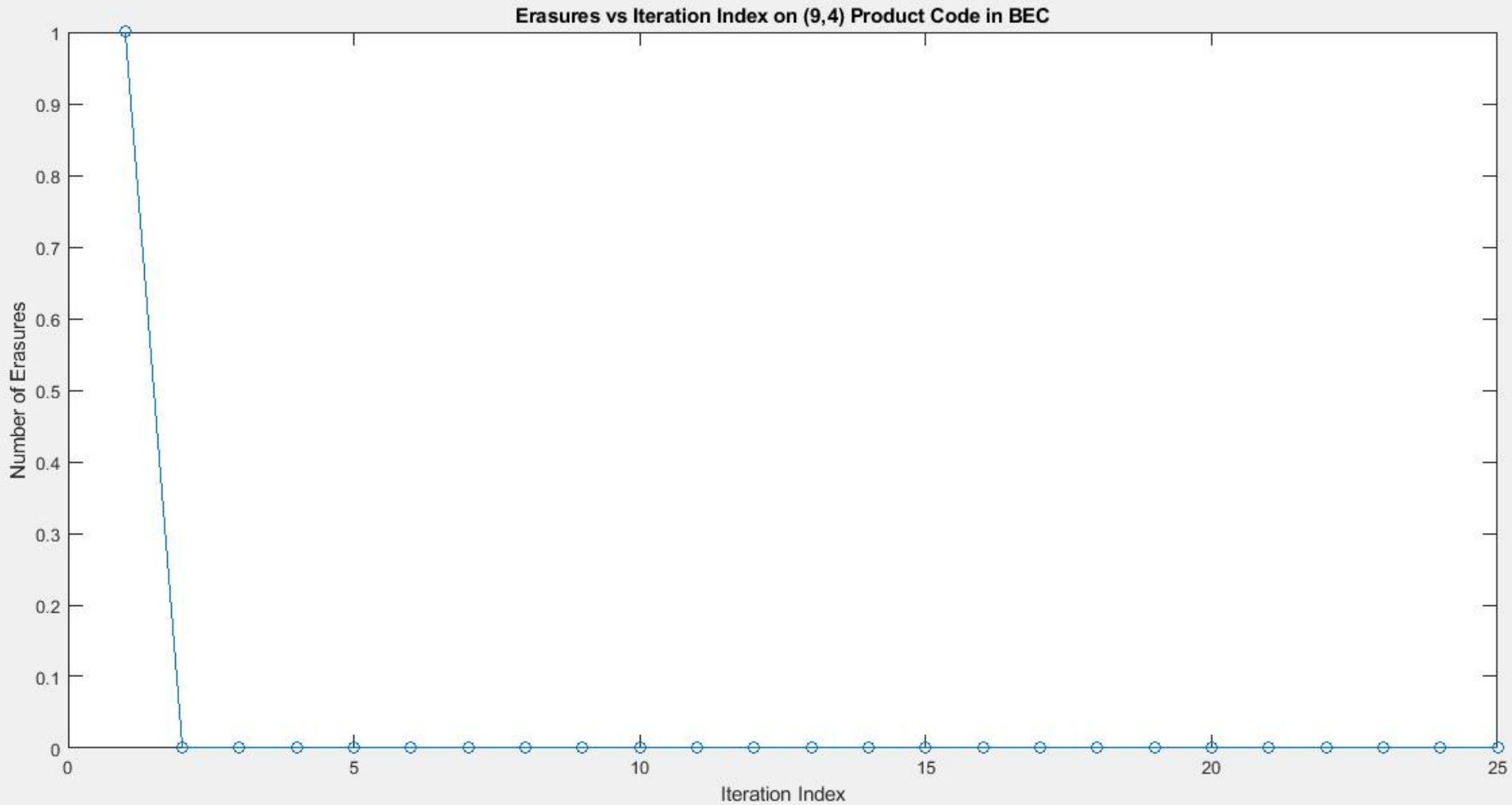
decoded_codeword =

     1     1     0     0     1     1     1     0     1

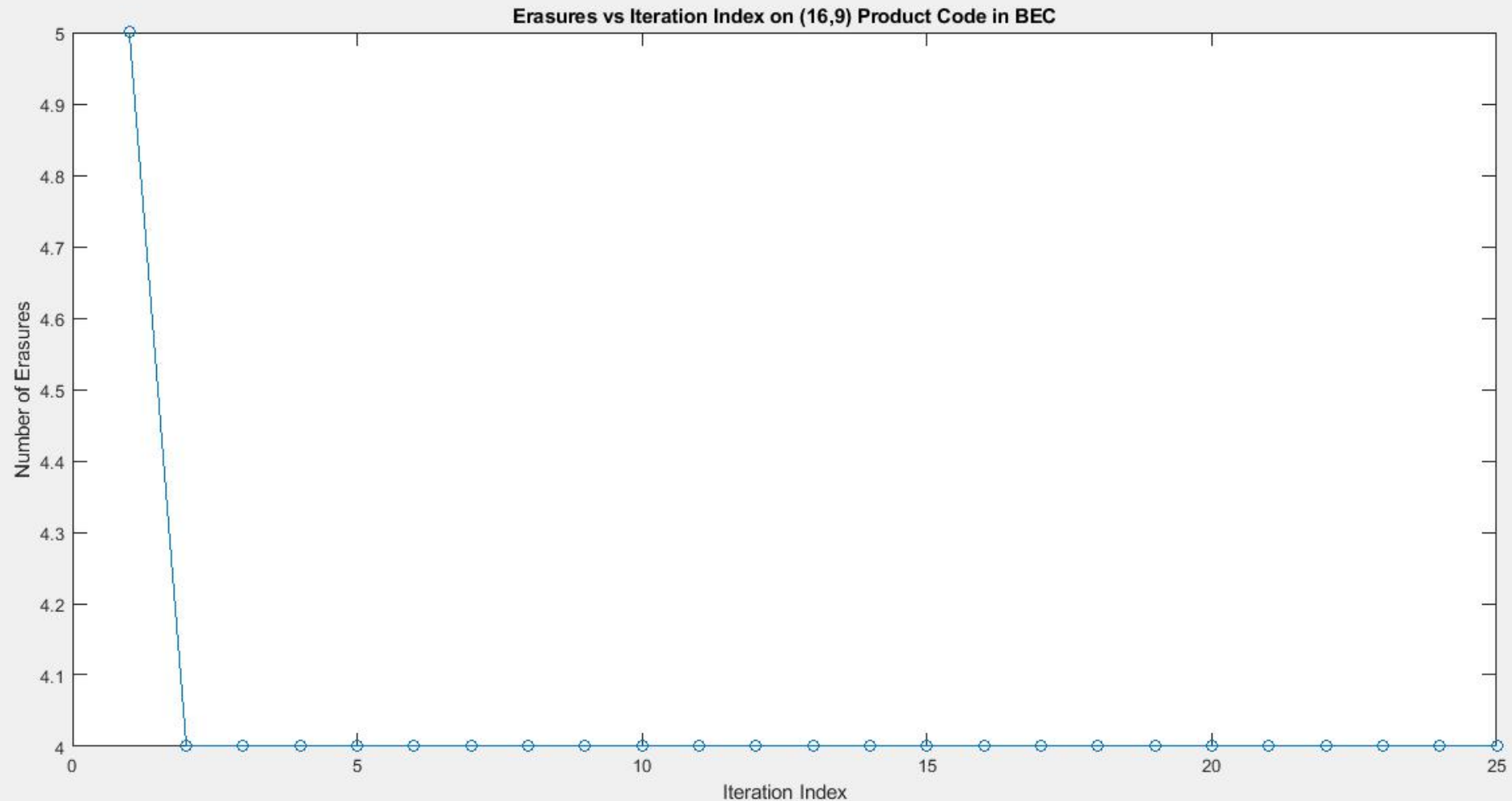
fx >>
```



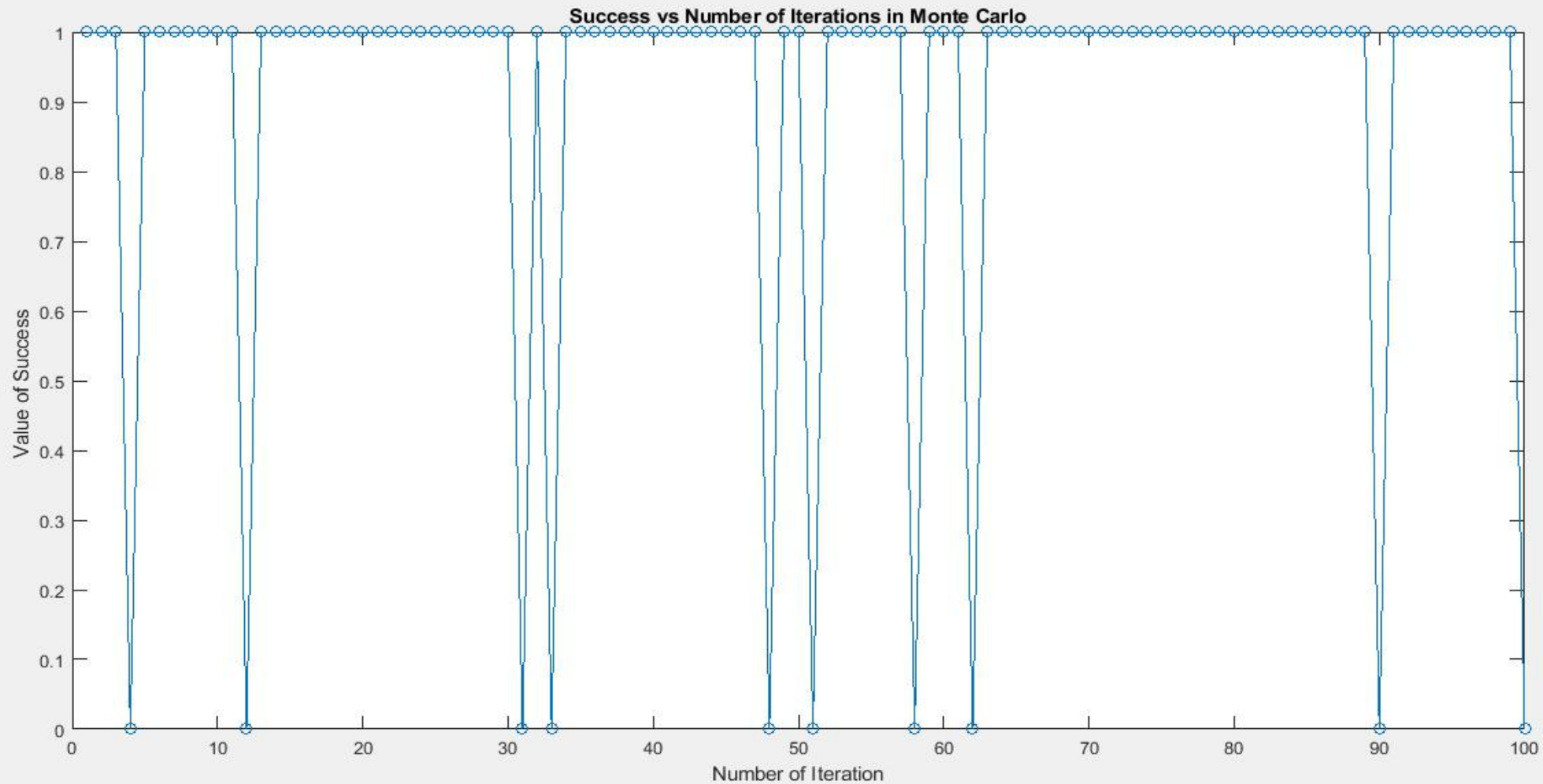
# BEC Simulation - For (9,4) Product Code



# BEC Simulation - For (16,9) Product Code

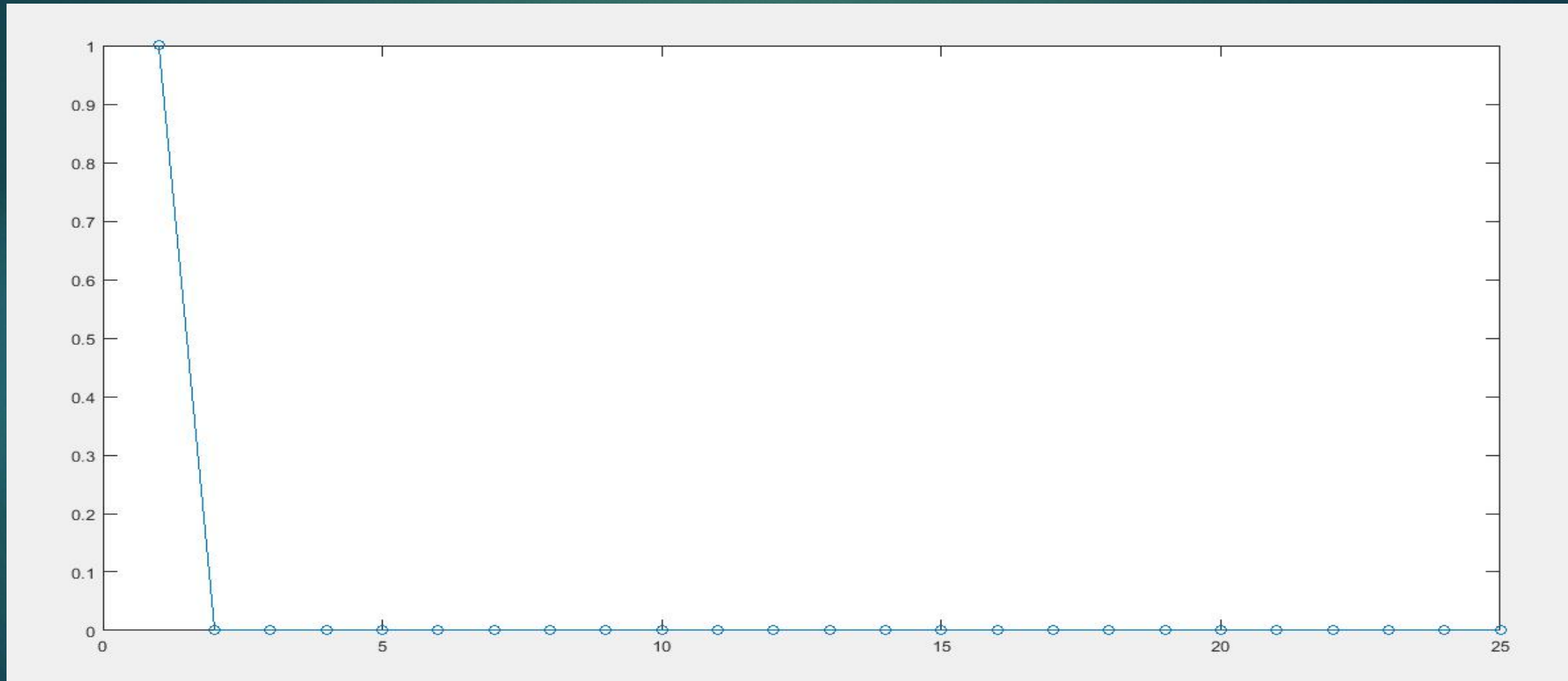


# Monte Carlo Simulation - BEC Product Code(9,4)





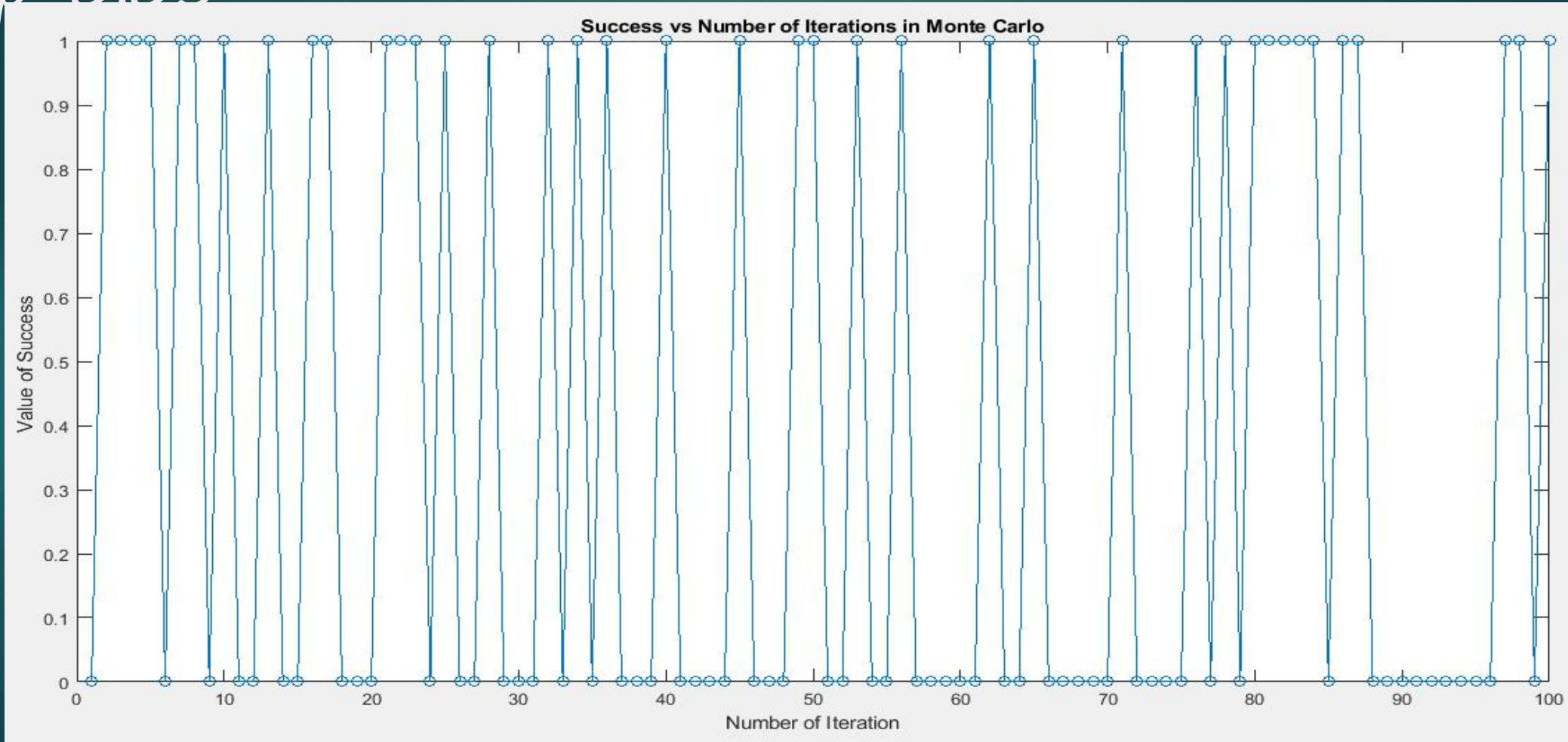
# BSC Simulation- For (25,16) Product Code (Bit Errors vs Iteration Index)



# Monte Carlo Simulation - BSC Product

Code(25,16)

$p=0.05$







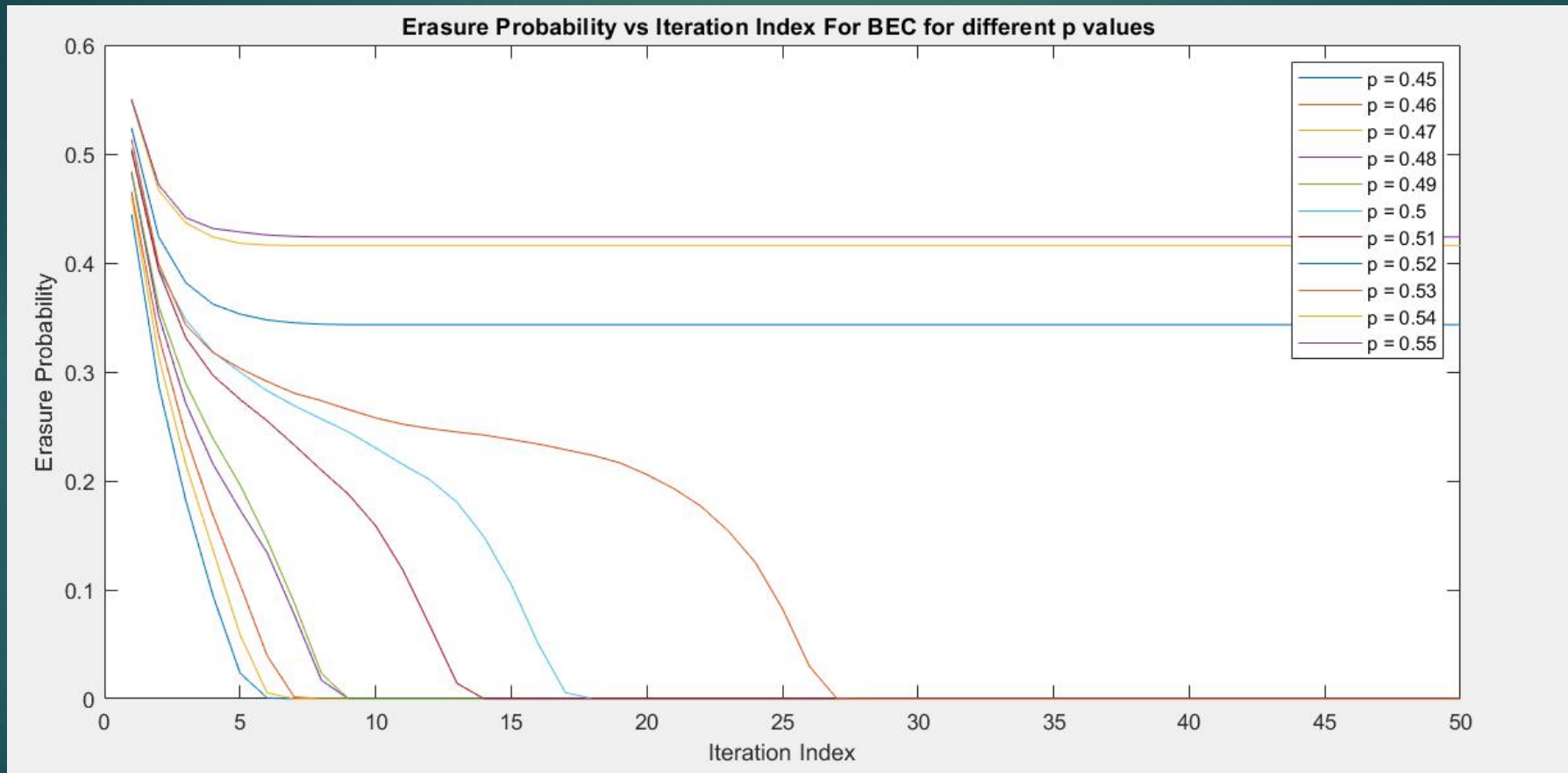




Erasure Probability vs Iteration Index For BEC for different  $p$  values(0.45-0.55)

$H(3000 \times 5000)$  50 Iterations this time.

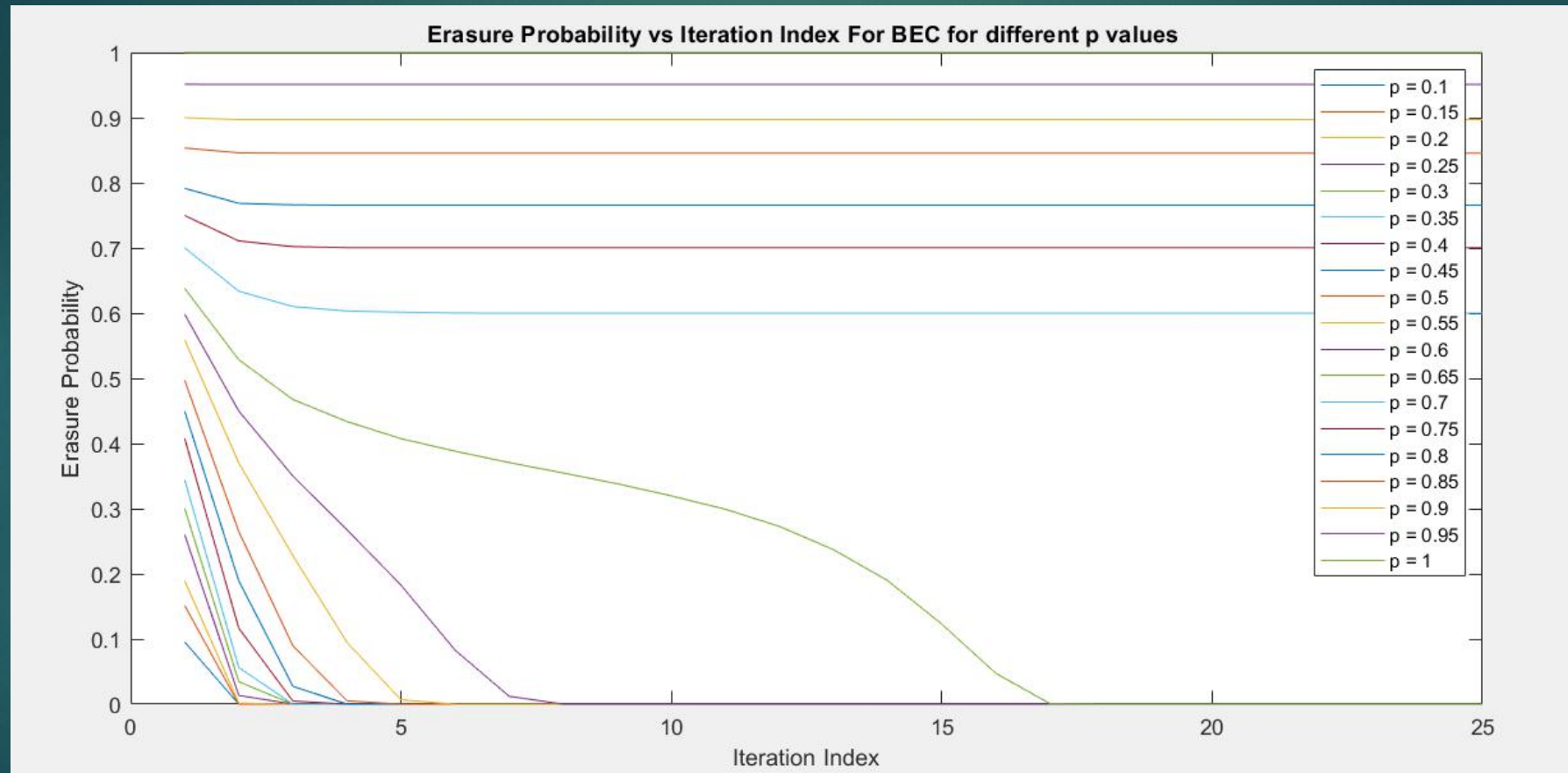
Didn't get decoded for  $p=0.55, 0.54$  and  $0.52$





# Erasure Probability vs Iteration Index For BEC for different $p$ values(0.1-1)

$H(3792 \times 5056)$





THANK YOU...