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Branch – CSE(3<sup>rd</sup> year)

## **LP LAB Assignment - 3**

### Implementing CYK Algorithm

Question:-

1. Implement CYK algorithm for deciding membership of a string in CFG in CNF. Assume that terminals and non-terminals are represented by a single alphabet and each grammar rule is given as a string where the first symbol is left side and the remaining portion is the corresponding right side of the production. Modify it to produce the number of non-identical derivation sequences for a given string in the grammar. Also generate the distinct derivation sequences, if possible.

Code:-

```

#include <bits/stdc++.h>
using namespace std;
#define pb push_back
struct node
{
    char c; // to store non-terminal which can drive cerresponding index's
    character in input string int count;
    array<int, 2> A[50]; // to store index of non-terminals (Like index of S A
C)
    array<int, 2> B[50]; // to store the index of cell which contain non-
terminal
    array<int, 2> C[50]; // to store the index of previous row from which we
    derive current row
};
#define vn vector<struct node>
#define vvn vector<vn>
map<string, unordered_set<char>>> P;
vector<vvn> Table;
vector<string> res; // To store Derivation Sequence
void derivation_sequence(vvn &R, int p)
{
    int cnt = 0;
    // checking wheather we reached to 0th row of Table
    // means row which contain original input string
    for (auto T : R[p])
    {
        if (T.count == 0)
        {
            cnt++;
        }
        else
        {
            break;
        }
    }
    if (cnt == R[p].size())
    {
        int z = 0;
        string s;
        for (auto cell : R)
        {
            for (auto T : cell)

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        {
            s.pb(T.c);
        }
        if ((z + 1) != R.size())
        {
            s += "->";
        }
        z++;
    }
    res.pb(s);
    return;
}
for (int i = 0; i < R[p][cnt].count; i++)
{
    vn cell;
    array<int, 2> A = R[p][cnt].A[i];
    array<int, 2> B = R[p][cnt].B[i];
    array<int, 2> C = R[p][cnt].C[i];
    for (int j = 0; j < R[p].size(); j++)
    {
        if (j != cnt)
        {
            cell.pb(R[p][j]);
        }
        else
        {
            cell.pb(Table[C[0]][B[0]][A[0]]);
            if (C[1] != 0)
            {
                cell.pb(Table[C[1]][B[1]][A[1]]);
            }
        }
    }
    R.pb(cell);
    derivation_sequence(R, p + 1);
    R.pop_back();
}
}
void CYK(string &w, char start)
{
    int n = w.size();
    // make 0th row of input string
    vn R;
    for (int i = 0; i < n; i++)
    {

```

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    vn cell;
    struct node T;
    T.c = w[i];
    T.count = 0;
    cell.pb(T);
    R.pb(cell);
}
Table.pb(R);
// We check if 'A -> a' is a production rule and 'a = w[i]'
// and if any then we store 'A'
// so first row of table
R.clear();
for (int i = 0; i < n; i++)
{
    vn cell;
    int j = 0;
    string s;
    s.pb(w[i]);
    // check production from which we can directly derive current input
symbol
    for (auto A : P[s])
    {
        struct node T;
        T.c = A;
        T.count = 1;
        T.A[j] = {0, 0};
        T.B[j] = {i, i}; // as from ith input symbol we derive this row
        T.C[j] = {0, 0}; // as from 0th row(which is input string) we are
        deriving this row
        cell.pb(T);
        j++;
    }
    R.pb(cell);
}
Table.pb(R);
// We will check for rule A -> BC
for (int l = 2; l <= n; l++)
{
    R.clear();
    for (int i = 0; i < n - l + 1; i++)
    {
        vn cell;
        for (int j = 1; j < l; j++)
        {
            for (int k = 0; k < Table[j][i].size(); k++)

```

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{
    for (int r = 0; r < Table[l - j][j + i].size(); r++)
    {
        string s;
        s.pb(Table[j][i][k].c);
        s.pb(Table[l - j][j + i][r].c);
        if (P.find(s) != P.end())
        {
            for (auto A : P[s])
            {
                int idx = -1;
                // check if any cell already contain A if A -> BC
                is a production for (int q = 0; q < cell.size(); q++)
                {
                    if (cell[q].c == A)
                    {
                        idx = q;
                        break;
                    }
                }
                if (idx == -1)
                {
                    struct node T;
                    T.c = A;
                    T.count = 1;
                    T.A[0] = {k, r};
                    T.B[0] = {i, j + i};
                    T.C[0] = {j, l - j};
                    cell.pb(T);
                }
                else
                {
                    int z = cell[idx].count;
                    cell[idx].count++;
                    cell[idx].A[z] = {k, r};
                    cell[idx].B[z] = {i, j + i};
                    cell[idx].C[z] = {j, l - j};
                }
            }
        }
    }
}
R.pb(cell);
}

```

```

    Table.pb(R);
}
// Printing Table
cout << "Table :- \n";
for (int i = Table.size() - 1; i >= 0; i--)
{
    if (i != 0)
    {
        cout << i << " ";
    }
    else
    {
        cout << " ";
    }
    for (int j = 0; j < Table[i].size(); j++)
    {
        if (i != 0)
        {
            cout << "{";
        }
        else
        {
            cout << " ";
        }
        for (int k = 0; k < Table[i][j].size(); k++)
        {
            cout << Table[i][j][k].c;
        }
        if (i != 0)
        {
            cout << "}\t";
        }
        else
        {
            cout << "\t";
        }
    }
    cout << "\n";
}
int idx = -1;
// Checking in Table at n,0 whether start is present or not
// if not than given string doesn't satisfy Given Grammar
for (int i = 0; i < Table[n][0].size(); i++)
{
    if (Table[n][0][i].c == start)

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    {
        idx = i;
    }
}
if (idx != -1)
{
    cout << w << " is Accepted\n";
    R.clear();
    vn cell;
    // We will start with (n,0) cell if that cell contain start symbol
    cell.pb(Table[n][0][idx]);
    R.pb(cell);
    // To find derivation sequence
    derivation_sequence(R, 0);
    cout << "No. of Derivation Sequence : " << res.size() << endl;
    cout << "Derivation Sequences :-\n";
    for (auto s : res)
    {
        cout << s << "\n";
    }
}
else
{
    cout << w << " is not Accepted by Given Grammer\n";
}
}
int main()
{
    int n; // no. of production
    cout << "Enter No. of Production : ";
    cin >> n;
    vector<string> prod(n); // set of production in CNF
    cout << "Enter Production Rule(In CNF) :-\n";
    for (int i = 0; i < n; i++)
    {
        cin >> prod[i]; // e.g. S->AB
    }
    char start; // start symbol
    for (int i = 0; i < n; i++)
    {
        string s = prod[i];
        P[s.substr(3)].insert(s[0]); // e.g. P[AB] = S
        if (i == 0)
        {
            start = s[0];
        }
    }
}

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    }
}
cout << "Input String :- ";
string w; // input string for which membership have to decide
cin >> w;
CYK(w, start);
}

```

```

rajat_goyal@Rajat_laptop:/mnt/c/Users/rajat/Desktop/CYK$ ./cyk
Enter No. of Production : 8
Enter Production Rule(In CNF) :-
S->AB
S->BC
A->BA
A->a
B->CC
B->b
C->AB
C->a
Input String :- baaba
Table :-
5 {SAC}
4 {} {CSA}
3 {} {B} {B}
2 {SA} {B} {CS} {SA}
1 {B} {CA} {CA} {B} {CA}
  b   a   a   b   a
baaba is Accepted
No. of Derivation Sequence : 2
Derivation Sequences :-
S->BC->bC->bAB->baB->baCC->baABC->baaBC->baabC->baaba
S->AB->BAB->bAB->baB->baCC->baABC->baaBC->baabC->baaba
rajat_goyal@Rajat_laptop:/mnt/c/Users/rajat/Desktop/CYK$

```