

# COMP 6651: Assignment 9

Fall 2020

**Submission through Moodle is due by November 22nd at 23:55**

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1. Recall that a language  $L$  belongs to the class  $\mathcal{NP}$  if there exists a polynomial time verification algorithm  $A$  and a constant  $d$  such that for every  $x \in \{0, 1\}^*$ 
  - if  $x \in L$  then there exists  $y$  such that  $|y| \leq |x|^d$  and  $A(x, y) = 1$ ;
  - if  $x \notin L$  then for every  $y$  with  $|y| \leq |x|^d$  we have  $A(x, y) = 0$ .
  - (a) On a test, Alex mistakenly forgot to specify the second item (if  $x \notin L \dots$ ). Alex's definition defines another class of languages,  $Alex - \mathcal{NP}$ . Determine the class  $Alex - \mathcal{NP}$ . Give a simple description of this class. Prove your answer.
  - (b) Another student, Joana, mistakenly forgot to specify the first item (if  $x \in L \dots$ ). Joana's definition defines another class of languages,  $Joana - \mathcal{NP}$ . Determine the class  $Joana - \mathcal{NP}$ . Give a simple description of this class. Prove your answer.
  - (c) Another student, Steve, replaced both items with the following: for every  $x \in \{0, 1\}^*$  there exists  $y$  such that  $x \in L$  if and only if  $|y| \leq |x|^d$  and  $A(x, y) = 1$ . This is again an incorrect definition of  $\mathcal{NP}$ . This definitions gives another class of languages,  $Steve - \mathcal{NP}$ . Determine the class  $Steve - \mathcal{NP}$ . Give a simple description of this class. Prove your answer.
2. Let  $G = (V, E)$  be a simple undirected graph. Recall that a set  $S \subseteq V$  is called a *vertex cover* if every edge has at least one endpoint in  $S$ , i.e.,  $\forall \{u, v\} \in E$  we have  $\{u, v\} \cap S \neq \emptyset$ . The Minimum Vertex Cover problem (MVC) asks to find a vertex cover of minimum size for a given graph  $G$ .
  - (a) State the optimization version, decision version, and search version of MVC.
  - (b) Let  $MVC - DEC$  denote the decision version of MVC. Recall that  $CLIQUE = \{\langle G, k \rangle : G \text{ has a clique of size at least } k\}$  is the decision version of the Maximum Clique problem. Show that  $CLIQUE \leq_p MVC - DEC$ . Give a reduction and prove its correctness and that it runs in polynomial time.
3. Recall that a simple undirected graph  $G = (V, E)$  is  $k$ -colorable if there exists a coloring  $c : V \rightarrow \{1, 2, \dots, k\}$  such that for every edge  $\{u, v\} \in E$  we have  $c(u) \neq c(v)$ . Define  $k - COL = \{\langle G \rangle : G \text{ is } k\text{-colorable}\}$ . Prove that  $3 - COL \leq_p 4 - COL$ . Give a reduction and prove its correctness and that it runs in polynomial time.