```
Assignment 2

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Part I.

1ab. Given an array of integers of any size, $n \ge 1$, and an integer m, write an algorithm as a pseudo code (not a program!) that would find all the flipped pairs from the given array. A flipped pair is governed by the equation b = a + m. For instance, given [1, 3, -5, 9, -2, -4, 2, 7, 4, 6], and m=3 the algorithm will return [1, 4] and [3, 6], [-5, -2], and [9, 2]. You must only print pairs where index of a is less than that of b. Find a simple solution with time complexity of O(n2).

```
A: And array of integers n: length of the array A m: integer which defines the relation of a flipped pair (a, b), where b = a + m. S: set of all flipped pair. Initially S < \emptyset function findPairs(A, m): for i = 0 - N do for j = i + 1 - N do if A[i] is equal to A[i] + m do S < S \cup (A[i], A[j]) endif endfor return S
```

1c. Can this task be performed in O(nlogn) or O(logn)? If yes, please provide the pseudo code for your approach. If not, explain why it is not possible.

This task cannot be performed in O(n log n) or O (log n) time. This is because we cannot sort/search the array in logarithmic time since for a flipped pair (a, b), the index of a needs to come before the index of b.

1d. Is a solution in O(n) even a possibility? If yes, please provide the pseudo code for your approach. If not, explain why it is not possible.

This task cannot be performed in O(n) time. This is because for every pair (a, b), the numbers could exist anywhere in the Array and hence we need to traverse the array from index of a to the last index for finding such a pair. This will be repeated for every index of a.

2a. Given a collection of n numbers, write an algorithm, using pseudo code that will output all possible combinations of r numbers that sum up to a prime number. For instance, given $\{1,2,3,4,5\}$, algorithm must output (1,2), (1,4), (2,3), (2,5), and (3,4) for r=2. Write an iterative as well as a recursive solution for this problem.

```
x: a number
function isPrime(x):
       if x <= 1 do
               return false
       for i <- 2 to √x do
               if i \% x == 0 do
                       return false
               endif
       endfor
       return true
A: an array of numbers
r: a positive integer
function findPrimeCombinationsIterative(A, r):
       prevSolns = {}
       primeCombinations = {}
       foreach number in A do
               newSolns = {}
               soln = {number}
               newSolns = newSolns U soln
               foreach prevSoln in prevSolns do
                       if size(prevSoln) == r-1 do
                               sum = 0
                               foreach toAdd in prevSoln do
                                       sum = sum + toAdd
                               endfor
                               sum = sum + number
                               if isPrime(sum) do
                                       soln = prevSoln U number
                                       primeCombinations = primeCombinations U soln
                                       newSolns = newSolns U soln
                               endif
                       else if size(prevSoln) < r-1 do
                               soln = prevSoln ∪ number
                               if soln prevSoln do
                                       newSolns = newSolns U soln
                               endif
                       endif
               endfor
               prevSolns = prevSolns U newSolns
       endfor
       return primeCombinations
```

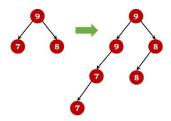
2b. What is the time complexity of your algorithm, in terms of Big-O?

The time complexity is of order $O(n^{r+1})$

2c. What changes will you make to the algorithm to output all permutations instead. The output for the same input will now contain (1,2) as well as (2,1).

While generating new solutions from previous solution, we can store an inverse of the array (i.e. values from end to the start) and that way we can generate all possible permutations.

3. Develop a well-documented pseudo code that inserts a duplicate node as the left child of the node. A sample tree transformation is provided below.



Is it guaranteed to have the resulting tree as a balanced tree (BT) if the input was a BT?

```
T: a tree
```

root: the root node of the tree
function copyNodeToLeft(root):
 if root is null:
 return null
 endif
 copyNodeToLeft(left child of root)
 copyNodeToLeft(right child of root)
 newNode <- makeNode()
 newNode.data <- root.data
 newNode.left <- root.left

root.left <- newNode

return root

Part II.

The Sales.java program defines a Sales class with different attributes corresponding to the information regarding a sale. The Sales class also has multiple accessors and mutators in order to provide flexibility during object creation and modification. There are also methods to convert object to string and to compare different Sales cine objects

The SalesDatabase.java class allows filtering unique records from within each file in a directory tree and writes the filtered records to an output file. This is done by use of many custom methods made with respect to the Sales Database.

InvalidFileException.java implements a custom exception class that is thrown when a file is invalid.

EmptyFolderException.java implements a custom exception class that is thrown when a folder is empty.

DuplicateRecordException.java implements a custom exception class that is thrown when a duplicate record is found.