

nous we know that h is a random no. with a complex gaussian distribu this random no. can be represented as n= rejo : where v, o are random values now the received symbo can be represented as y= rejo x(a+oj) +n. as n, of o are random no.s this received symbol can lie any where on the complex plane if we take r=1, $0=\pi/2$ of n=0 tuen y will become $y = 1e^{j\pi/2}(a) + n = 0 + aj$ which is infact the symbol S2 so basically the symbols after getting transmitted through the channel end up somewhere, we don't have any control where they will land. This was more of an intuitive proof, lets prove the above with numbers.

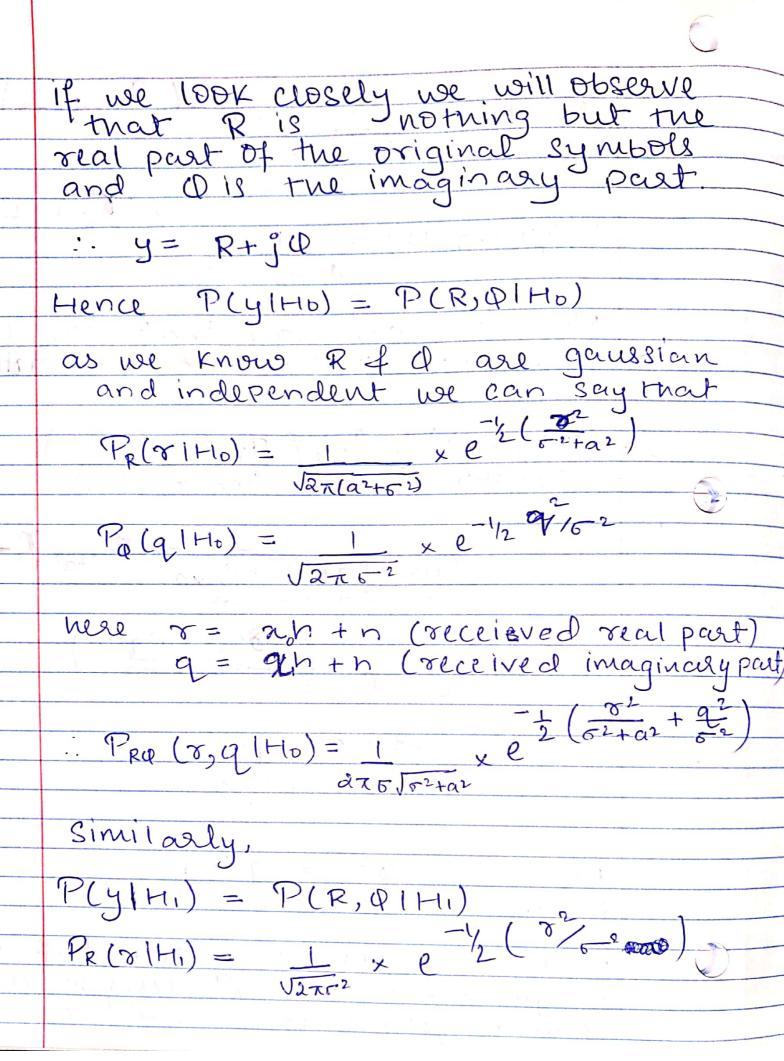
Let us apply our basic ML rule Ho => 40= (a+0) h + n H1 => 41= (0+aj)h+n use can clearly see that, as ha No (0,1) of na No (0,0-2) both y, & yo ~ N. (0,0) + N. (0,02) if both y, f yo will have the same distribution we will never be able to form a decision boundary. The above statement gives us a idea, How can use make the distributions' of y.f yo dafferent? There is one way,

If our symbols would have been like in the case of on-off Keying,

we indeed will be able to adrieve thal Now our ML rube becomes: Ho= yo= (a+ oj)h+n; f $H_1 \Rightarrow y_1 = Oxh + n$

Hence now, yo ~ No(0,02) + No(0,52) 41 N N(0, 52) Now we can do something to detect different. But the problem is we need to use the symbols given to us in the problem statement; To solve that problem we can use a combined method. Silo Si és a+oj or (a,0) in complex S2 is 0+aj or (0,a) in complex plane. * we can interpret Si as an 'a' followed by a 'O' in the on off Keying example. whereas, * we can interpret S1 as a 'O' followed by an 'a' in the on off Keying example. $S_1 \Longrightarrow [a, o]$ $S_2 \Longrightarrow [o, a]$ So to transmit si we will first transmit 'a' through the channel followed by a 'o'.

similarly, we will transmit a 0' through the channel followed by an (a'. if we want to transmit Sz. after transmitting two symbols separately, we will make a desision rule by taking pairs of received symbols. let P, I denote the first symbol received and the second symbol received, respectively. Tuen, $R = ah + n =) R \sim \mathcal{N}_{c}(0,a^{2}) + \mathcal{N}_{c}(0,e^{2})$ $Q = oh + n =) Q \sim \mathcal{N}_{c}(0,e^{2})$ Ho=) $R = Oh + n \Rightarrow R N N_c(O_5 = 2)$ $Q = Qh + n \Rightarrow Q N N_c(O_5 = 2) + N_c(O_5 = 2)$ $H_1 \Rightarrow$ we know $N_c(0,a^2) + N_c(0,\sigma^2) = N_c(0,a^2+\sigma^2)$ tuerefore on simplification R~ Nc (0, a2+52) Ho =O~ No (0,52) $H_1 \Rightarrow 0 \sim \mathcal{N}_c(0, \alpha^2 + \sigma^2)$ R~ N((0,52)



Po(q|H1) = 1 × e
$$\frac{1}{2}$$
 ($\frac{9^{2}}{6^{2}+\alpha^{2}}$)

 $\frac{1}{2}$ ($\frac{8^{2}}{6^{2}}$ + $\frac{9^{2}}{6^{2}+\alpha^{2}}$)

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Code quide:	
There are 3 scripts	
a) Main.m - for demo	
b) commSystem.m - simulated	system
 There are 3 scripts a) Main.m — for demo b) commSystem.m — simulated c) bit-Exekate.m — calculating	bit eggor rate
to get a demo of the pri	oject
 run Main.m, please ensus	e all
 the three files are in the	2 same
 to get a demo of the protein main. In please ensus the three files are in the folder and are added to po	eth.
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