

## Class Assignment - 6

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### Part 1: Implementation of LRU Page Replacement in xv6

- **Step 1 – Modifications in `proc.h`**

File Name: `proc.h`

Code Added: New `frameinfo` structure

```
struct frameinfo {  
    uint va;      // Virtual address of page  
    pte_t *pte;   // Page table entry  
    int ref;      // Reference (optional for CLOCK)  
    uint last_used; // Time of last access (for LRU)  
};
```

Added fields in `struct proc` :

```
struct frameinfo frames[16];  
  
int framecount;
```

- **Step 2 – Modifications in `vm.c`**

File Name: `vm.c`

Code Added:

```
extern uint ticks; // Global tick counter from trap.c  
  
void update_lru_access(struct proc *p, uint va) {  
    for (int i = 0; i < p->framecount; i++) {  
        if (p->frames[i].va == va) {  
            p->frames[i].last_used = ticks; // update on access  
            break;  
        }  
    }  
}
```

```

    }
}

```

## Replacing existing `allocvm()` from new function :

```

int allocvm(pde_t *pgdir, uint oldsz, uint newsz) {
    char *mem;
    uint a;

    if (newsz >= KERNBASE) return 0;
    if (newsz < oldsz) return oldsz;

    a = PGROUNDUP(oldsz);
    for (; a < newsz; a += PGSIZE) {
        mem = kalloc();
        if (mem == 0) {
            cprintf("allocvm out of memory\n");
            deallocvm(pgdir, newsz, oldsz);
            return 0;
        }
        memset(mem, 0, PGSIZE);
        if (mappages(pgdir, (char*)a, PGSIZE, V2P(mem), PTE_W | PTE_U)
< 0) {
            cprintf("allocvm out of memory (2)\n");
            deallocvm(pgdir, newsz, oldsz);
            kfree(mem);
            return 0;
        }

        // ----- LRU Update -----
        struct proc *curproc = myproc();
        if (curproc) {
            if (curproc->framecount < 16) {
                curproc->frames[curproc->framecount].va = a;
                curproc->frames[curproc->framecount].pte = walkpgdir(pgdir,
(void*)a, 0);
                curproc->frames[curproc->framecount].last_used = ticks;
                curproc->framecount++;
            } else {
                // Select victim using LRU
                int victim_index = 0;

```

```

        for (int i = 1; i < curproc->framecount; i++) {
            if (curproc->frames[i].last_used <
                curproc->frames[victim_index].last_used) {
                victim_index = i;
            }
        }

        // Free victim
        pte_t *vppte = curproc->frames[victim_index].pte;
        uint vpa = PTE_ADDR(*vppte);
        kfree(P2V(vpa));
        *vppte = 0;

        // Replace with new
        curproc->frames[victim_index].va = a;
        curproc->frames[victim_index].pte = walkpgdir(pgdir, (void*)a,
0);
        curproc->frames[victim_index].last_used = ticks;
    }
}
}
return newsz;
}

```

- **Step 3 : User Program mytest.c**

**File Name:** mytest.c

**Code Added:**

```

#include "types.h"

#include "user.h"

#define MAX_PAGES 5

#define TOTAL_ACCESSES 15

int main(int argc, char *argv[]) {

```

```

int lru[MAX_PAGES];

int last_used[MAX_PAGES];

int i, j, page, hit = 0, miss = 0, time = 0;

for (i = 0; i < MAX_PAGES; i++) {

    lru[i] = -1;

    last_used[i] = -1;

}

printf(1, "Starting LRU page replacement simulation...\n");

int accesses[TOTAL_ACCESSES] = {0, 1, 2, 3, 4, 1, 5, 0, 6, 1, 2, 7, 3,
8, 4};

for (i = 0; i < TOTAL_ACCESSES; i++) {

    page = accesses[i];

    time++;

    int found = 0;

    for (j = 0; j < MAX_PAGES; j++) {

        if (lru[j] == page) {

            found = 1;

            last_used[j] = time;

            break;

        }

    }

    if (found) {

        hit++;

        printf(1, "Access page %d: HIT\n", page);

    } else {

        miss++;

```

```

        int replaced = -1;
        for (j = 0; j < MAX_PAGES; j++) {
            if (lru[j] == -1) {
                lru[j] = page;
                last_used[j] = time;
                replaced = j;
                break;
            }
        }

        if (replaced == -1) {
            int lru_index = 0, min_time = last_used[0];
            for (j = 1; j < MAX_PAGES; j++) {
                if (last_used[j] < min_time) {
                    min_time = last_used[j];
                    lru_index = j;
                }
            }

            printf(1, "Access page %d: MISS, replacing page %d\n",
                page, lru[lru_index]);
            lru[lru_index] = page;
            last_used[lru_index] = time;
        } else {
            printf(1, "Access page %d: MISS, placed in free frame\n", page);
        }
    }
}

```

```

}

printf(1, "LRU simulation completed.\n");

printf(1, "Total hits: %d\n", hit);

printf(1, "Total misses: %d\n", miss);

    exit();

}

```

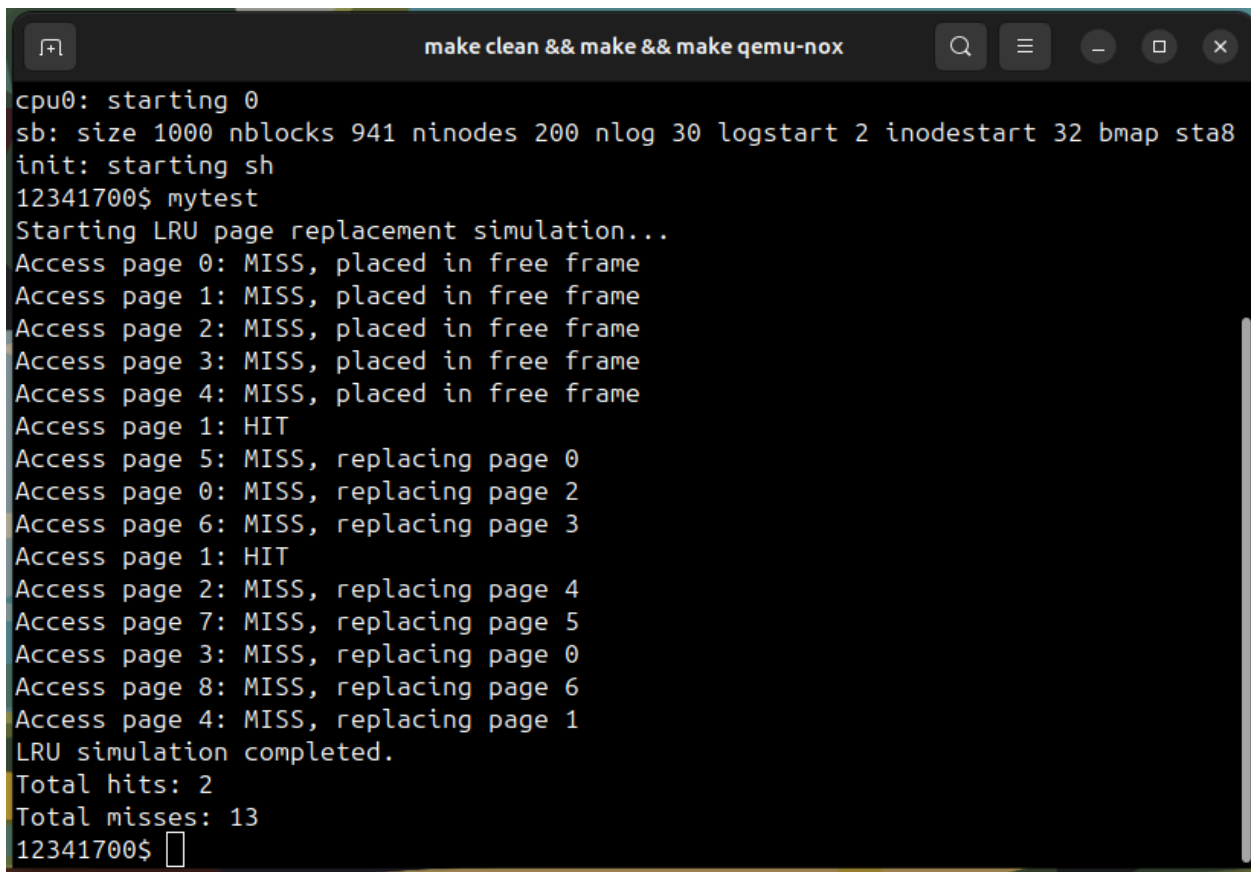
## • Step 4 : Makefile Update

File Name: **Makefile**

Code Added:

```
_mytest\
```

## • OUTPUT :



```

make clean && make && make qemu-nox
cpu0: starting 0
sb: size 1000 nblocks 941 ninodes 200 nlog 30 logstart 2 inodestart 32 bmap sta8
init: starting sh
12341700$ mytest
Starting LRU page replacement simulation...
Access page 0: MISS, placed in free frame
Access page 1: MISS, placed in free frame
Access page 2: MISS, placed in free frame
Access page 3: MISS, placed in free frame
Access page 4: MISS, placed in free frame
Access page 1: HIT
Access page 5: MISS, replacing page 0
Access page 0: MISS, replacing page 2
Access page 6: MISS, replacing page 3
Access page 1: HIT
Access page 2: MISS, replacing page 4
Access page 7: MISS, replacing page 5
Access page 3: MISS, replacing page 0
Access page 8: MISS, replacing page 6
Access page 4: MISS, replacing page 1
LRU simulation completed.
Total hits: 2
Total misses: 13
12341700$ 

```

**Explanation :** The output shows the working of the LRU (Least Recently Used) page replacement algorithm.

- At first, the free frames are filled with pages 0–4, resulting in misses.
- After the frames are full, every new page request may cause either a hit (if the page is already in memory) or a miss (if it is not).
- On a miss, the least recently used page is replaced.
- In this simulation, there were 2 hits and 13 misses in total.

This demonstrates how LRU tries to keep the most recently accessed pages in memory while replacing the least used ones when new pages arrive.

## Part 2: FIFO Page Replacement Practical in xv6-public

- **Step 1 – Modifications in `proc.h`**

**File Name:** `proc.h`

**Code Added:**

```
#define MAX_PAGES 20

struct proc {

    int pages[MAX_PAGES];

    int page_count;

};
```

**Explanation :**

- **Step 2 – Modifications in `vm.c`**

**File Name:** `vm.c`

**Code Added:**

```
int allocvm(pde_t *pgdir, uint oldsz, uint newsz) {
    char *mem;
    uint a;

    if(newsz >= KERNBASE)
        return 0;
    if(newsz < oldsz)
        return oldsz;

    a = PGROUNDUP(oldsz);
    for(; a < newsz; a += PGSIZE) {
        mem = kalloc();
        if(mem == 0) {
            cprintf("allocvm out of memory\n");
            deallocvm(pgdir, newsz, oldsz);
            return 0;
        }
        memset(mem, 0, PGSIZE);
    }
```



```

if(mappages(pgdir, (char*)a, PGSIZE, V2P(mem), PTE_W|PTE_U) < 0)
{
    kfree(mem);
    deallocvm(pgdir, newsz, oldsz);
    return 0;
}

struct proc *cur = myproc();

// Case 1: Still under MAX_PAGES
if(cur->page_count < MAX_PAGES) {
    cur->pages[cur->page_count++] = (int)a;
    cprintf("Allocated page %d at VA 0x%x\n", cur->page_count-1, a);
}
// Case 2: Need FIFO eviction
else {
    int evict = cur->pages[0];
    char *victim = (char*)P2V(PTE_ADDR(*walkpgdir(pgdir, (void*)evict,
0)));
    kfree(victim);

    // Shift FIFO queue
    for(int i = 1; i < MAX_PAGES; i++)
        cur->pages[i-1] = cur->pages[i];

    cur->pages[MAX_PAGES-1] = (int)a;
    cprintf("Evicted page at VA 0x%x, allocated new page at VA 0x%x\n",
evict, a);
}
}
return newsz;
}

```

- **Step 3 : User Program mytest.c**

**File Name: mytest.c**

**Code Added:**

```

#include "types.h"

#include "user.h"

#define FRAME_SIZE 5    // Number of frames in memory
#define TOTAL_ACCESSES 15 // Total number of page accesses

int main(int argc, char *argv[]) {
    int fifo[FRAME_SIZE];    // Stores pages in memory
    int next_to_replace = 0; // FIFO pointer
    int hit = 0, miss = 0;

    // Page access sequence
    int accesses[TOTAL_ACCESSES] = {0,1,2,3,4,1,5,0,6,1,2,7,3,8,4};

    // Initialize memory slots
    for(int i = 0; i < FRAME_SIZE; i++)
        fifo[i] = -1;

    printf(1, "Starting FIFO Page Replacement Simulation...\n");

    for(int i = 0; i < TOTAL_ACCESSES; i++) {
        int page = accesses[i];
        int found = 0;

        // Check if page is already in memory (HIT)
        for(int j = 0; j < FRAME_SIZE; j++) {
            if(fifo[j] == page) {
                found = 1;
            }
        }
    }
}

```

```

        break;
    }
}

if(found) {
    hit++;

    printf(1, "Access page %d: HIT\n", page);
} else {
    miss++;

    printf(1, "Access page %d: MISS, replacing page %d\n",
           page, fifo[next_to_replace]);

    fifo[next_to_replace] = page;

    next_to_replace = (next_to_replace + 1) % FRAME_SIZE;
}
}

printf(1, "\nFIFO Simulation Completed.\n");
printf(1, "Total Hits: %d\n", hit);
printf(1, "Total Misses: %d\n", miss);

exit();
}

```

- **Step 4 : Makefile Update**

**File Name:** **Makefile**

**Code Added:**

```
_mytest\
```

**OUTPUT :** FIFO Page Replacement (Access sequence = {0, 1, 2, 3, 4, 1, 5, 0, 6, 1, 2, 7, 3, 8, 4})

**1. For FRAME SIZE = 3**

```
make clean && make && make qemu-nox

Booting from Hard Disk..xv6...
cpu0: starting 0
sb: size 1000 nblocks 941 ninodes 200 nlog 30 logstart 2 inodestart 32 bmap sta8
Allocated page 0 at VA 0x0
Allocated page 1 at VA 0x1000
Allocated page 2 at VA 0x2000
init: starting sh
Allocated page 0 at VA 0x0
Allocated page 1 at VA 0x1000
Allocated page 2 at VA 0x2000
Allocated page 3 at VA 0x3000
12341700$ mytest
Allocated page 0 at VA 0x0
Allocated page 1 at VA 0x1000
Allocated page 2 at VA 0x2000
Starting FIFO page replacement simulation...
Access page 0: MISS, replacing page -1
Access page 1: MISS, replacing page -1
Access page 2: MISS, replacing page -1
Access page 3: MISS, replacing page 0
Access page 4: MISS, replacing page 1
Access page 1: MISS, replacing page 2
Access page 5: MISS, replacing page 3
Access page 0: MISS, replacing page 4
Access page 6: MISS, replacing page 1
Access page 1: MISS, replacing page 5
Access page 2: MISS, replacing page 0
Access page 7: MISS, replacing page 6
Access page 3: MISS, replacing page 1
Access page 8: MISS, replacing page 2
Access page 4: MISS, replacing page 7
FIFO simulation completed.
Total hits: 0
Total misses: 15
12341700$
```

**Explanation :** Frames = 3

The first three accesses (0, 1, 2) fill the frames. When page 3 is accessed, page 0 (the oldest) is evicted. Next, page 4 evicts page 1. Then, page 1 evicts page 2. When page 5 arrives, page 3 is evicted. The process continues this way, always removing the oldest page.

Drawback: even pages that are still useful can be evicted just because they were loaded earlier (classic FIFO anomaly).

## 2. For FRAME SIZE = 4

```
make clean && make && make qemu-nox

SeaBIOS (version 1.16.3-debian-1.16.3-2)

iPXE (https://ipxe.org) 00:03.0 CA00 PCI2.10 PnP PMM+1EFCAF60+1EF0AF60 CA00

Booting from Hard Disk..xv6...
cpu0: starting 0
sb: size 1000 nblocks 941 ninodes 200 nlog 30 logstart 2 inodestart 32 bmap sta8
Allocated page 0 at VA 0x0
Allocated page 1 at VA 0x1000
Allocated page 2 at VA 0x2000
init: starting sh
Allocated page 0 at VA 0x0
Allocated page 1 at VA 0x1000
Allocated page 2 at VA 0x2000
Allocated page 3 at VA 0x3000
12341700$ mytest
Allocated page 0 at VA 0x0
Allocated page 1 at VA 0x1000
Allocated page 2 at VA 0x2000
Starting FIFO page replacement simulation...
Access page 0: MISS, replacing page -1
Access page 1: MISS, replacing page -1
Access page 2: MISS, replacing page -1
Access page 3: MISS, replacing page -1
Access page 4: MISS, replacing page 0
Access page 1: HIT
Access page 5: MISS, replacing page 1
Access page 0: MISS, replacing page 2
Access page 6: MISS, replacing page 3
Access page 1: MISS, replacing page 4
Access page 2: MISS, replacing page 5
Access page 7: MISS, replacing page 0
Access page 3: MISS, replacing page 6
Access page 8: MISS, replacing page 1
Access page 4: MISS, replacing page 2
FIFO simulation completed.
Total hits: 1
Total misses: 14
12341700$ █
```

### Explanation : Frames = 4

The first four accesses (0, 1, 2, 3) fill the frames. When page 4 is accessed, page 0 is evicted. The next access to page 1 is a hit because it is still in memory. When page 5 arrives, page 1 is evicted. When page 0 is accessed again, page 2 is evicted. The process continues in the same manner.

With four frames, there are fewer evictions compared to three frames. However, FIFO may sometimes show Belady's anomaly, where having more frames can unexpectedly increase the number of misses.

### 3. For FRAME SIZE = 5

```
make clean && make && make qemu-nox

Booting from Hard Disk..xv6...
cpu0: starting 0
sb: size 1000 nblocks 941 ninodes 200 nlog 30 logstart 2 inodestart 32 bmap sta8
Allocated page 0 at VA 0x0
Allocated page 1 at VA 0x1000
Allocated page 2 at VA 0x2000
init: starting sh
Allocated page 0 at VA 0x0
Allocated page 1 at VA 0x1000
Allocated page 2 at VA 0x2000
Allocated page 3 at VA 0x3000
12341700$ mytest
Allocated page 0 at VA 0x0
Allocated page 1 at VA 0x1000
Allocated page 2 at VA 0x2000
Starting FIFO page replacement simulation...
Access page 0: MISS, replacing page -1
Access page 1: MISS, replacing page -1
Access page 2: MISS, replacing page -1
Access page 3: MISS, replacing page -1
Access page 4: MISS, replacing page -1
Access page 1: HIT
Access page 5: MISS, replacing page 0
Access page 0: MISS, replacing page 1
Access page 6: MISS, replacing page 2
Access page 1: MISS, replacing page 3
Access page 2: MISS, replacing page 4
Access page 7: MISS, replacing page 5
Access page 3: MISS, replacing page 0
Access page 8: MISS, replacing page 6
Access page 4: MISS, replacing page 1
FIFO simulation completed.
Total hits: 1
Total misses: 14
12341700$
```

#### Explanation : Frames = 5

The first five accesses (0, 1, 2, 3, 4) fill the frames. The next access to page 1 is a hit. When page 5 arrives, page 0 is evicted. Then, page 0 comes back and evicts page 1. Later, page 6 arrives and evicts page 2. The process continues.

With five frames, the number of misses is much lower compared to three or four frames, because more pages can stay resident in memory.