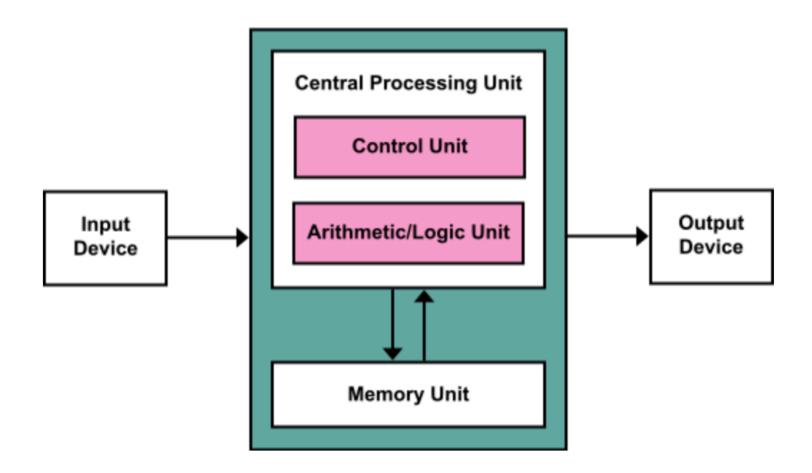
Part 8 Register Machines

von Neumann Machines

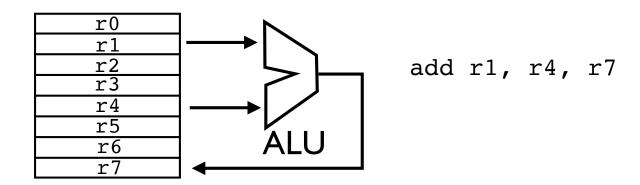


Arithmetic/Logic

CPUs have instructions that perform <u>single</u> arithmetic operations

```
add, sub, mul, div, and, or, xor, not, eq, lt, ...
```

 These operations are applied to values, typically supplied from "registers" on the CPU

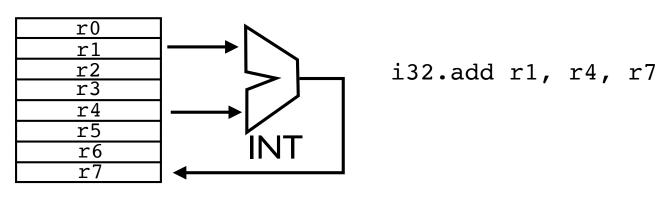


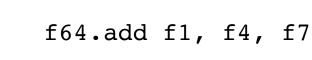
Data Types

There are two main datatypes (of varying sizes)

```
Integers (i8, i16, i32, i64, i128, etc.) Floats (f32, f64, f128, etc.)
```

Often handled by different ALUs & instructions





The Reality

- CPUs have limited resources
- Limited registers
- Limited memory
- Limited instructions
- Code needs to map to this constrained setting

Evaluate and show your work

Evaluate and show your work

Evaluate and show your work

$$2 + 3 * (10 - 2) + 5$$
 $r1 = 2$ $r2 = 3$ $2 + 3 * 8 + 5$ $26 + 5$

Evaluate and show your work

Evaluate and show your work

$$2 + 3 * (10 - 2) + 5$$
 $r1 = 2$ $r2 = 3$ $r3 = 10$ $r4 = 2$ $r4 = 2$ $r4 = 2$

Evaluate and show your work

Evaluate and show your work

```
2 + 3 * (10 - 2) + 5 r1 = 2

r2 = 3

2 + 3 * 8 + 5 r3 = 10

r4 = 2

r3 = sub(r3, r4); r3 = 8

r3 = r2 = mul(r2, r3); r2 = 24

r3 = r2 = mul(r2, r3); r3 = 8
```

Evaluate and show your work

```
2 + 3 * (10 - 2) + 5 r1 = 2

r2 = 3

2 + 3 * 8 + 5 r3 = 10

r4 = 2

r3 = sub(r3, r4); r3 = 8

r3 = sub(r2, r3); r2 = 24

r3 = sub(r3, r4); r3 = 8

r3 = sub(r3, r4); r3 = 8
```

Evaluate and show your work

```
2 + 3 * (10 - 2) + 5 r1 = 2

r2 = 3

2 + 3 * 8 + 5 r3 = 10

r4 = 2

r3 = sub(r3, r4); r3 = 8

r3 = sub(r2, r3); r2 = 24

r3 = sub(r3, r4); r3 = 8

r3 = sub(r3, r4); r3 = 8
```

Evaluate and show your work

```
2 + 3 * (10 - 2) + 5 r1 = 2

r2 = 3

2 + 3 * 8 + 5 r3 = 10

r4 = 2

2 + 24 + 5 r3 = sub(r3, r4); r3 = 8

r2 = mul(r2, r3); r2 = 24

r3 = r1 = add(r1, r2); r1 = 26

r3 = r1 = add(r1, r2); r1 = 31
```

Use registers (4 required)

Evaluate more efficiently

Evaluate more efficiently

$$2 + 3 * (10 - 2) + 5$$
 $r1 = 10$
 $2 + 3 * 8 + 5$
 $2 + 24 + 5$
 $26 + 5$

Evaluate more efficiently

$$2 + 3 * (10 - 2) + 5$$
 $r1 = 10$ $r2 = 2$
 $2 + 3 * 8 + 5$
 $2 + 24 + 5$
 $26 + 5$

Evaluate more efficiently

Evaluate more efficiently

```
2 + 3 * (10 - 2) + 5 r1 = 10 r2 = 2 r1 = sub(r1, r2); r1 = 8 r2 = 3 r3 + 24 + 5 r4 = 10 r5 = 10 r6 = 10 r6 = 10 r6 = 10 r6 = 10 r7 = 10 r8 = 10
```

Evaluate more efficiently

```
2 + 3 * (10 - 2) + 5 r1 = 10 r2 = 2 r1 = sub(r1, r2); r1 = 8 r2 = 3 r1 = mul(r2, r1); r1 = 24 r1 = r1
```

Evaluate more efficiently

```
2 + 3 * (10 - 2) + 5   r1 = 10   r2 = 2   r1 = 30   r2 = 3   r2 = 3   r1 = 30   r2 = 3   r2 = 3   r2 = 3
```

```
r1 = 10

r2 = 2

r1 = sub(r1, r2) ; r1 = 8

r2 = 3

r1 = mul(r2, r1) ; r1 = 24

r2 = 2
```

Evaluate more efficiently

```
2 + 3 * (10 - 2) + 5 r1 = 10 r2 = 2 r1 = sub(r1, r2) ; r1 = 8 r2 = 3 r1 = mul(r2, r1) ; r1 = 24 r2 = 2 r1 = add(r2, r1) ; r1 = 26
```

Evaluate more efficiently

Evaluate more efficiently

```
2 + 3 * (10 - 2) + 5 r1 = 10

r2 = 2

2 + 3 * 8 + 5 r1 = sub(r1, r2) ; r1 = 8

r2 = 3

r1 = mul(r2, r1) ; r1 = 24

r2 = 2

r1 = add(r2, r1) ; r1 = 26

r2 = 5

r1 = add(r1, r2) ; r1 = 31
```

Use registers (2 required!)

Evaluate with fewer instructions

Evaluate with fewer instructions

Evaluate with fewer instructions

$$2 + 3 * (10 - 2) + 5$$
 $r1 = 2$ $r2 = 10$ $2 + 3 * 8 + 5$ $26 + 5$

Evaluate with fewer instructions

Evaluate with fewer instructions

Evaluate with fewer instructions

```
2 + 3 * (10 - 2) + 5 r1 = 2 r2 = 10 r2 = sub(r2, r1) ; r2 = 8 r3 = 3 r2 = mul(r3, r2) ; r2 = 24 r3
```

Evaluate with fewer instructions

```
2 + 3 * (10 - 2) + 5 r1 = 2

r2 = 10

r2 = sub(r2, r1) ; r2 = 8

r3 = 3

r3 = 3
```

Evaluate with fewer instructions

```
2 + 3 * (10 - 2) + 5 r1 = 2

r2 = 10

r2 = sub(r2, r1) ; r2 = 8

r3 = 3

r3 = 3
```

Evaluate with fewer instructions

```
2 + 3 * (10 - 2) + 5 r1 = 2

r2 = 10

2 + 3 * 8 + 5 r2 = sub(r2, r1) ; r2 = 8

r3 = 3

r3
```

Now takes 8 instructions! (but 3 registers)

New Questions!

- What happens if you run out of registers?
- What is the most efficient use of registers?
- The register scheduling problem
- It is non-trivial

Register Spilling

 If registers are exhausted, must "spill" an existing register to memory and read it back later

```
r1 = 10
...
store(r1, temp1) ; spill r1 -> temp
r1 = 123 ; Load a new value into r1
...
r1 = load(temp1) ; load temp -> r1
...
```

Memory is slow. You want to avoid spilling

Register Scheduling

- To efficiently use registers, must perform advanced "analysis" on the code to know how values can be reused
- Example:

$$(x + y)/a + (x + y)/b$$

Maybe could rewrite

$$t = x + y$$

 $t/a + t/b$

"Common subexpression elimination"

Example

```
(x+y)/2 + (x+y)/4
```

```
r1 = load(x)

r2 = load(y)

r3 = add(r1, r2)

r4 = 2

r5 = div(r3, r4)

r6 = load(x)

r7 = load(y)

r8 = add(r6, r7)

r9 = 4

r10 = div(r8,r9)

r11 = add(r5,r10)
```

```
r1 = load(x)
r2 = load(y)
r3 = add(r1, r2)
r4 = 2
r5 = div(r3, r4)
    ; already in r1
    ; already in r2
    ; already in r3
r9 = 4
r10 = div(r3, r4)
r11 = add(r5, r10)
```

Register Lifetimes

```
r1 = load(x)
r2 = load(y)
r3 = add(r1, r2) ; r1, r2 not referenced again
r4 = 2
r5 = div(r3, r4) ; r4 not referenced again
r9 = 4
r10 = div(r3, r4) ; r3, r4 not referenced again
r11 = add(r5, r10) ; r5, r10 not reference again
```

Rewrite with register overwriting

```
r1 = load(x)

r2 = load(y)

r1 = add(r1, r2)

r2 = 2

r2 = div(r1, r2)

r3 = 4

r3 = div(r1, r3)

r1 = add(r2, r3)
```

A Moment of Reflection

All CPUs are similar









- But the low-level details vary. For example, number of registers, variety of instructions, etc.
- Question: Do you write a compiler for a single model of a specific CPU? No.

Abstract Machines

- Compilers often target an "abstract machine"
- A generic "CPU"
- With a standard set of basic "instructions"

Intermediate Representation

- The abstract machine is programmed using "intermediate representation" or IR Code
- It's like a generic machine code
 - Mimics architecture of actual CPUs
 - "Easy" to translate to actual machine code

Big Picture

Source

```
print 2 + 3 * 4;
```

AST



IR



```
('i32.const', 2, 'r1'),

('i32.const', 3, 'r2'),

('i32.const', 4, 'r3'),

('i32.mul', 'r2', 'r3', 'r4')

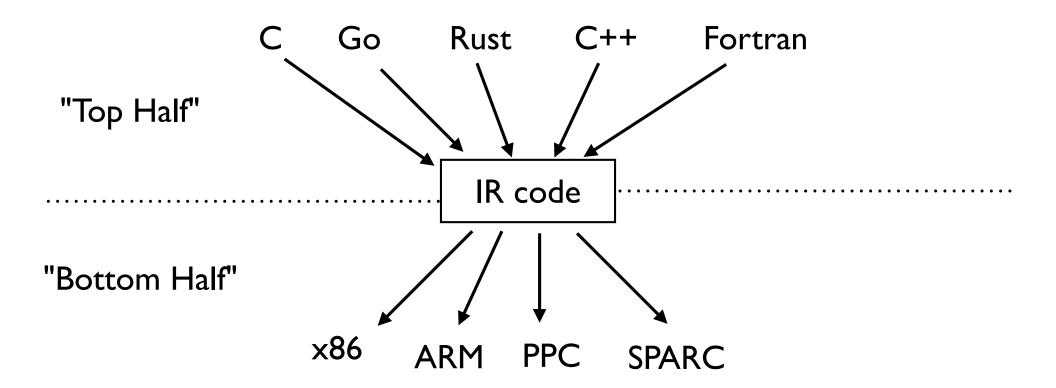
('i32.add', 'r1', 'r4', 'r5')
```





Metal

Compiler Design



IR Code

- What does IR code look like?
- It embodies a few essential concepts
 - A Model of Computation
 - Control Flow
 - Memory
 - Modules
- Examples: LLVM, WebAssembly

Project

- Compile Wabbit to IR
 - See wabbit/llvm.py
 - See wabbit/wasm.py