

## Linear Data Structure

### Definition

A **data structure** is said to be **linear** if its elements form a sequence or a **linear** list.

### Examples:

- Array
- Linked List
- Stacks
- Queues

### Operations on linear Data Structures

- *Traversal* : Visit every part of the **data structure**
- *Search* : Traversal through the data structure for a given element
- *Insertion* : Adding new elements to the **data structure**
- *Deletion* : Removing an element from the **data structure**.
- *Sorting* : Rearranging the elements in some type of order(e.g. Increasing or Decreasing)
- *Merging* : Combining two similar **data** structures into one

## STACK – a pile of elements

### Introduction:

1. Stack is basically a data object
2. The operational semantic (meaning) of stack is LIFO i.e. last in first out

**Definition:** It is an ordered list of elements  $n$ , such that  $n > 0$  in which all insertions and deletions are made at one end called the top.

### Primary operations defined on a stack:

1. PUSH : add an element at the top of the list.
2. POP : remove at the top of the list.
3. Also "IsEmpty()" and IsFull()" function, which tests whether a stack is empty or full respectively.

### Example:

1. **Practical daily life** : a pile of heavy books kept in a vertical box, dishes kept one on top of another

2. **In computer world:** In processing of subroutine calls and returns ; there is an explicit use of stack of return addresses.

Also in evaluation of [arithmetic expressions](#), stack is used.

Large number of stacks can be expressed using a single one dimensional stack only. Such an array is called a [multiple stack array](#).

## PROBLEM – 1

### Tower of Hanoi

Tower of Hanoi is a historical problem, which can be easily expressed using recursion. There are N disks of decreasing size stacked on one needle, and two other empty needles. It is required to stack all the disks onto a second needle in the decreasing order of size. The third needle can be used as a temporary storage. The movement of the disks must confirm to the following rules -

1. Only one disk may be moved at a time
2. A disk can be moved from any needle to any other.
3. The larger disk should not rest upon a smaller one.

**Question:** write a c program to implement tower of Hanoi using stack?

**Solution:** /\* Program of towers of Hanoi. \*/

```
#include <stdio.h>
#include <conio.h>

void move ( int, char, char, char ) ;

void main( )
{
    int n = 3 ;
    clrscr( ) ;
    move ( n, 'A', 'B', 'C' ) ;
    getch( ) ;
}

void move ( int n, char sp, char ap, char ep )
{
    if ( n == 1 )
        printf ( "\nMove from %c to %c ", sp, ep ) ;
    else
    {
        move ( n - 1, sp, ep, ap ) ;
```

```

        move ( 1, sp, ' ', ep ) ;
        move ( n - 1, ap, sp, ep ) ;
    }
}

```

## PROBLEM – 2

### Function Calls and Stack

A stack is used by programming languages for implementing function calls. Write a program to check how function calls are made using stack.

**Solution:** /\* To show the use of stack in function calls \*/

```

#include <stdio.h>
#include <conio.h>
#include <stdlib.h>
#include <dos.h>

unsigned int far *ptr ;
void ( *p )( void ) ;

void f1( ) ;
void f2( ) ;

void main( )
{
    f1( ) ;
    f2( ) ;
    printf ( "\nback to main..." ) ;
    exit ( 1 ) ;
}

void f1( )
{
    ptr = ( unsigned int far * ) MK_FP ( _SS, _SP + 2 ) ;
    printf ( "\n%d", *ptr ) ;
    p = ( void ( * )( ) ) MK_FP ( _CS, *ptr ) ;
    ( *p )( ) ;
    printf ( "\nI am f1( ) function " ) ;
}

void f2( )

```

```
{  
    printf ( "\nI am f2( ) function" );  
}
```

## **PUSH & POP**

### **Algorithms**

#### ***Push (item,array , n, top)***

```
{  
    If ( n>= top)  
        Then print "Stack is full" ;  
    Else  
        {  
            top = top + 1;  
            array[top] = item ;  
        }  
}
```

#### ***Pop (item,array,top)***

```
{  
    if ( top<= 0)  
        Then print " stack is empty".  
    Else  
        {  
            item = array[top];  
            top = top - 1;  
        }  
}
```

## PUSH OPERATION

### Working of a Stack

11

Stack is implemented here as a one dimensional array of size 7

### Addition of element to Stack

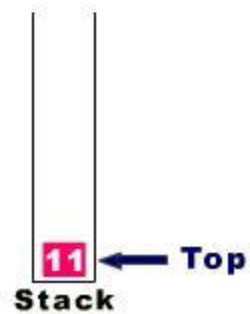


push 11  
on Stack

Step - 1

### Working of a Stack

### Addition of element to Stack

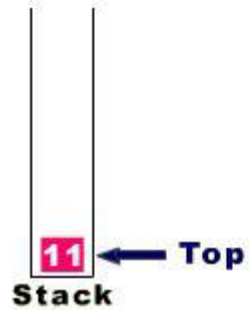


Step – 2

## Working of a Stack

23

**Addition of element  
to Stack**



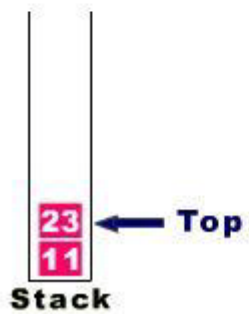
push 23  
on Stack

Step – 3

## Working of a Stack

-8

**Addition of element  
to Stack**



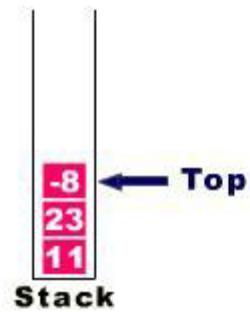
push -8  
on Stack

Step – 4

### Working of a Stack

16

Addition of element  
to Stack



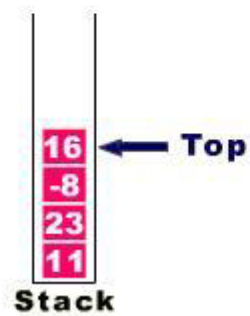
push 16  
on Stack

Step – 5

### Working of a Stack

27

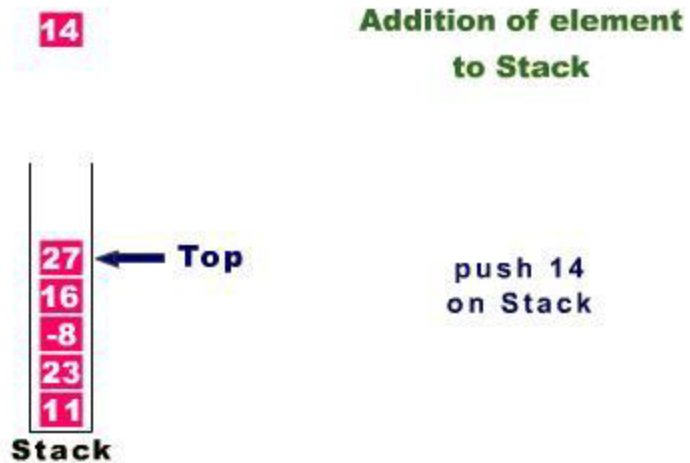
Addition of element  
to Stack



Push 27  
on Stack

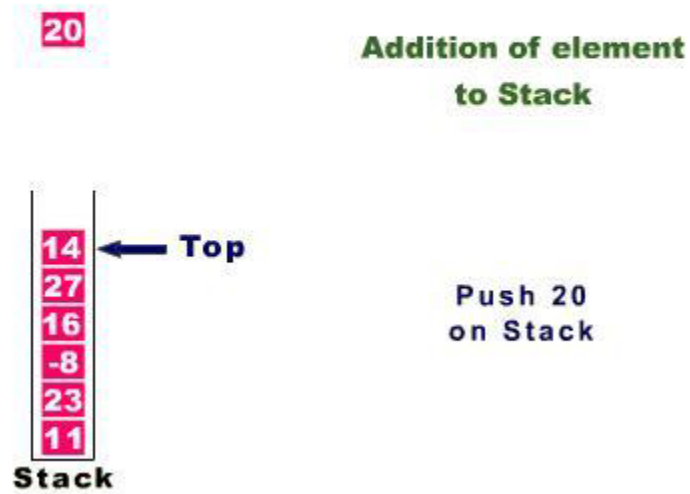
Step – 6

## Working of a Stack



Step – 7

## Working of a Stack



Step – 8



## Working of a Stack

**Addition of element  
to Stack**

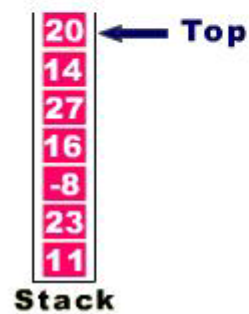


Step – 9

## POP OPERATION

## Working of a Stack

**Deletion of element  
by calling pop**

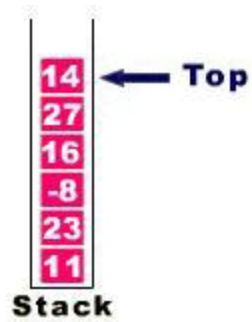


Step – 10

## Working of a Stack

20

Deletion of element  
by calling pop



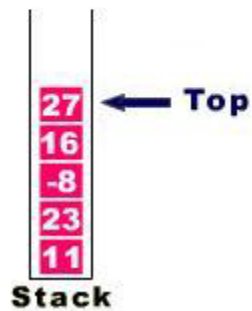
pop 20  
from Stack

Step – 11

## Working of a Stack

14

Deletion of element  
by calling pop



pop 14  
from Stack

Step – 12

## Arithmetic Expressions

**Arithmetic expressions are expressed as combinations of:**

1. Operands
2. Operators (arithmetic, Boolean, relational operators)

Various rules have been formulated to specify the order of evaluation of combination of operators in any expression.

**The arithmetic expressions are expressed in 3 different notations:**

### 1. *Infix*:

- In this if the operator is binary; the operator is between the 2 operands.
- And if the operator is unary, it precedes the operand.

### 2. *Prefix*:

- In this notation for the case of binary operators, the operator precedes both the operands.
- Simple algorithm using stack can be used to evaluate the final answer.

### 3. *Postfix*:

- In this notation for the case of binary operators, the operator is after both the corresponding operands.
- Simple algorithm using stack can be used to evaluate the final answer.

*Always remember that the order of appearance of operands does not change in any Notation. What changes is the position of operators working on those operands.*

## **RULES FOR EVALUATION OF ANY EXPRESSION:**

An expression can be interpreted in many different ways if parentheses are not mentioned in the expression.

- For example the below given expression can be interpreted in many different ways:
- Hence we specify some basic rules for evaluation of any expression :

**A priority table is specified for the various type of operators being used:**

PRIORITY LEVEL	OPERATORS
6	** ; unary - ; unary +
5	* ; /
4	+ ; -
3	< ; > ; <= ; >= ; != ; !< ; !>
2	Logical and operation
1	Logical or operation

### Algorithm for evaluation of an expression E which is in prefix notation:

- We assume that the given prefix notation starts with IsEmpty() .
- If number of symbols = n in any infix expression then number of operations performed = some constant times n.
- Here *next token* function gives us the next occurring element in the expression in a left to right scan.
- The *PUSH* function adds element x to stack Q which is of maximum length n

Evaluate (E)

```

{
    Top = 0;
    While (1)
    {
        x = next token (E)
        If (x == infinity)
        {
            Print value of stack [top]
            as the output of the expression
        }
    }
}

Else
{
    If (x == operand)
        PUSH (Q, top, n, x);
    If (x == operator)
    {
        Pop correct number
        of operands according to the
        the operator (unary/binary)
        and then perform the operation
        and store result onto
        the stack
    }
}

```

### Algorithm for evaluation of an expression E, which is in postfix notation:

- We assume that the postfix notation specifies end of expression by appending NULL at the end of expression.
- here *next token* function gives us the next occurring element in the expression in a left to right scan.

- The *PUSH* function adds element  $x$  to stack  $Q$  which is of maximum length  $n$

Evaluate (E)

```

{
    Top = 0;
    While (1)
    {
        x = next token (E)
        If (x == infinity)
        {
            Return (stack [top]);
        }
    }
}

Else
{
    If (x == operand)
        PUSH (Q, top, n, x);
    If (x == operator)
    {
        Pop correct number of
        operands according to the ope-
        rands according to the operat-
        or (unary/binary) and then
        perform the operation and
        store result onto the stack
    }
}

```