

Agenda

1. K^{th} Smallest Element
2. Morris Inorder Traversal
3. LCA in Binary Tree
4. LCA in BST
5. In time and out time

Binary Tree

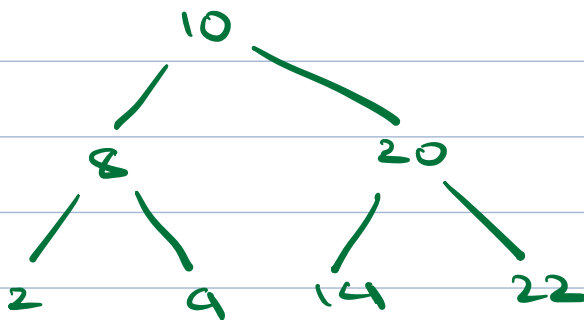
- Hierarchical data structure, composed of tree nodes
- Can have at max 2 children: left and right

Binary Search Tree (BST)

$O(N)$

- For every node x
All nodes in LST $\leq x$
All nodes in RST $> x$

1. Given a BST and a positive integer K , find K^{th} smallest element in BST.



K	ans
3	9
5	14

0 1 2 3 4 5 6
2 8 9 10 14 20 22

Do inorder traversal of BST

Approach
1 :

↓
store in arr[]

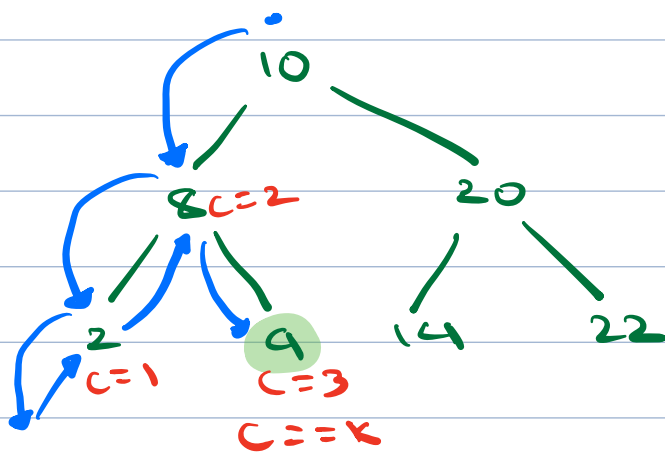
↓
return arr[K-1]

TC: $O(N)$

SC: $O(H + N)$

↓
Recursive stack

Approach 2: keep a global count variable while doing inorder traversal



$cnt = 0 \neq 3$ $k = 3$

ans = 9

ans

cnt = 0

void inorder (Node root) {

if (root == NULL)
return

inorder (root.left)

cnt++

if (cnt == k) {

ans = root.data
return

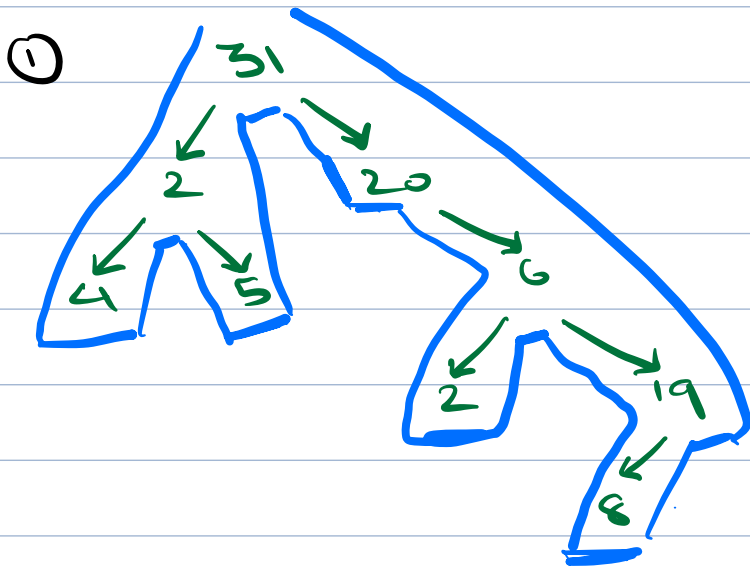
if (cnt < k) {

inorder (root.right)

TC: $O(N)$

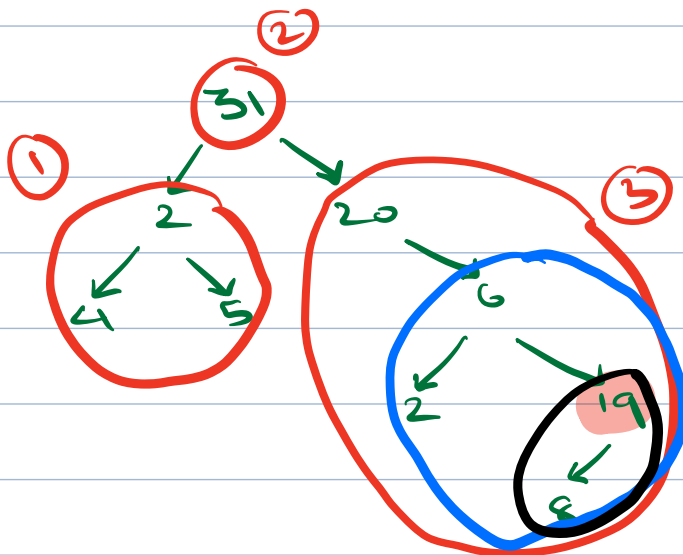
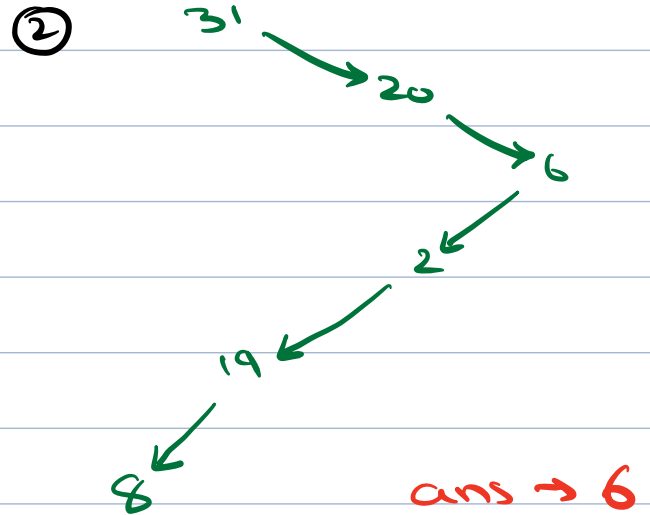
SC: $O(H)$

Q. Which is the last node printed in inorder traversal ?



4 2 5 31 20 2 6 8 19

ans → 19

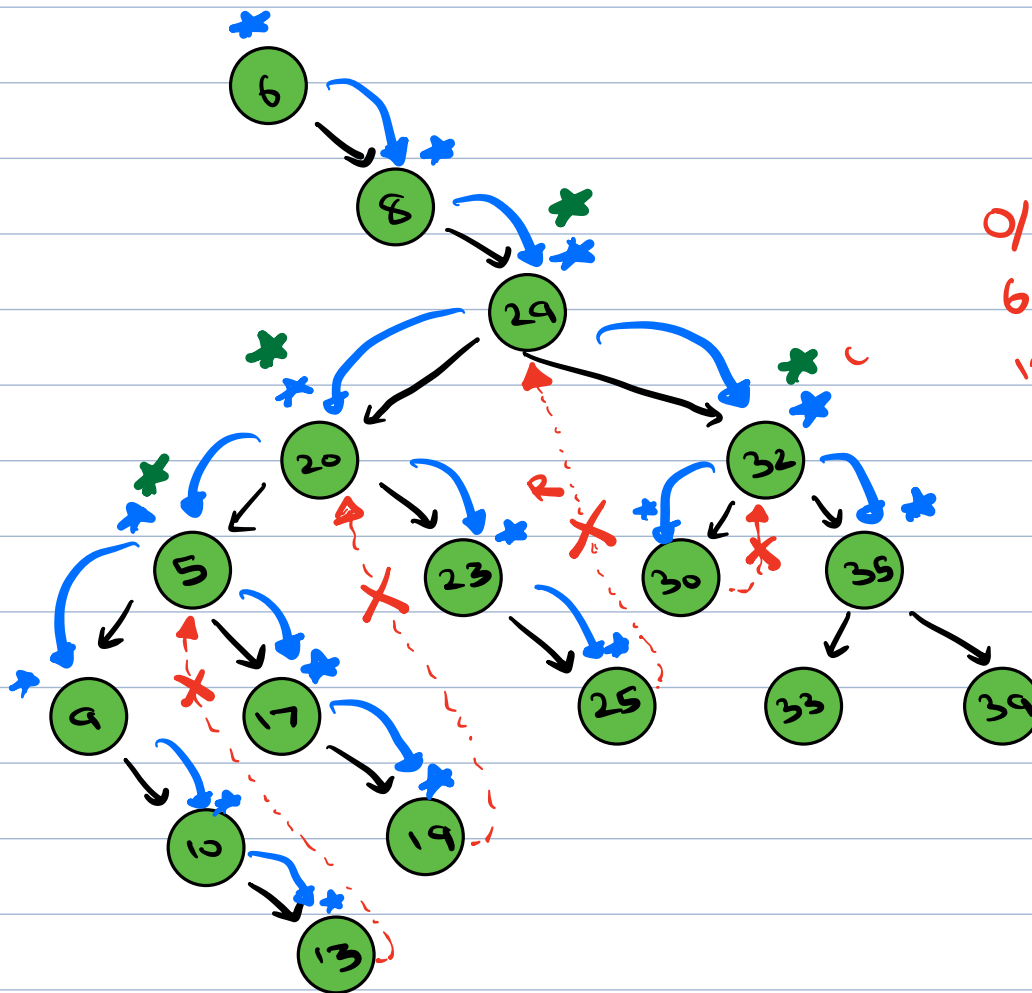


LNR

Rightmost node of tree → last node in inorder

2. Morris Inorder Traversal

SC $\rightarrow O(1)$



O/P

6 8 9 10
13 5 17
19 20 23
25 29 30
32

25, right = 29

Last node of LST of 29



Rightmost node of LST

SC

Recursive $\rightarrow O(H)$

Iterative $\rightarrow O(H)$



$O(1)$

```
void inorder (Node root)
```

```
Node cur = root
```

```
while (cur != NULL) {
```

```
    if (cur.left == NULL) {
```

```
        print (cur.data)
```

```
        cur = cur.right
```

```
    } else {
```

```
        Node temp = cur.left
```

```
        while (temp.right != NULL &&  
            temp.right != cur) {
```

```
            temp = temp.right
```

```
        if (temp.right == NULL) {
```

```
            temp.right = cur
```

```
            cur = cur.left
```

```
        } else if (temp.right == cur) {
```

```
            temp.right = NULL
```

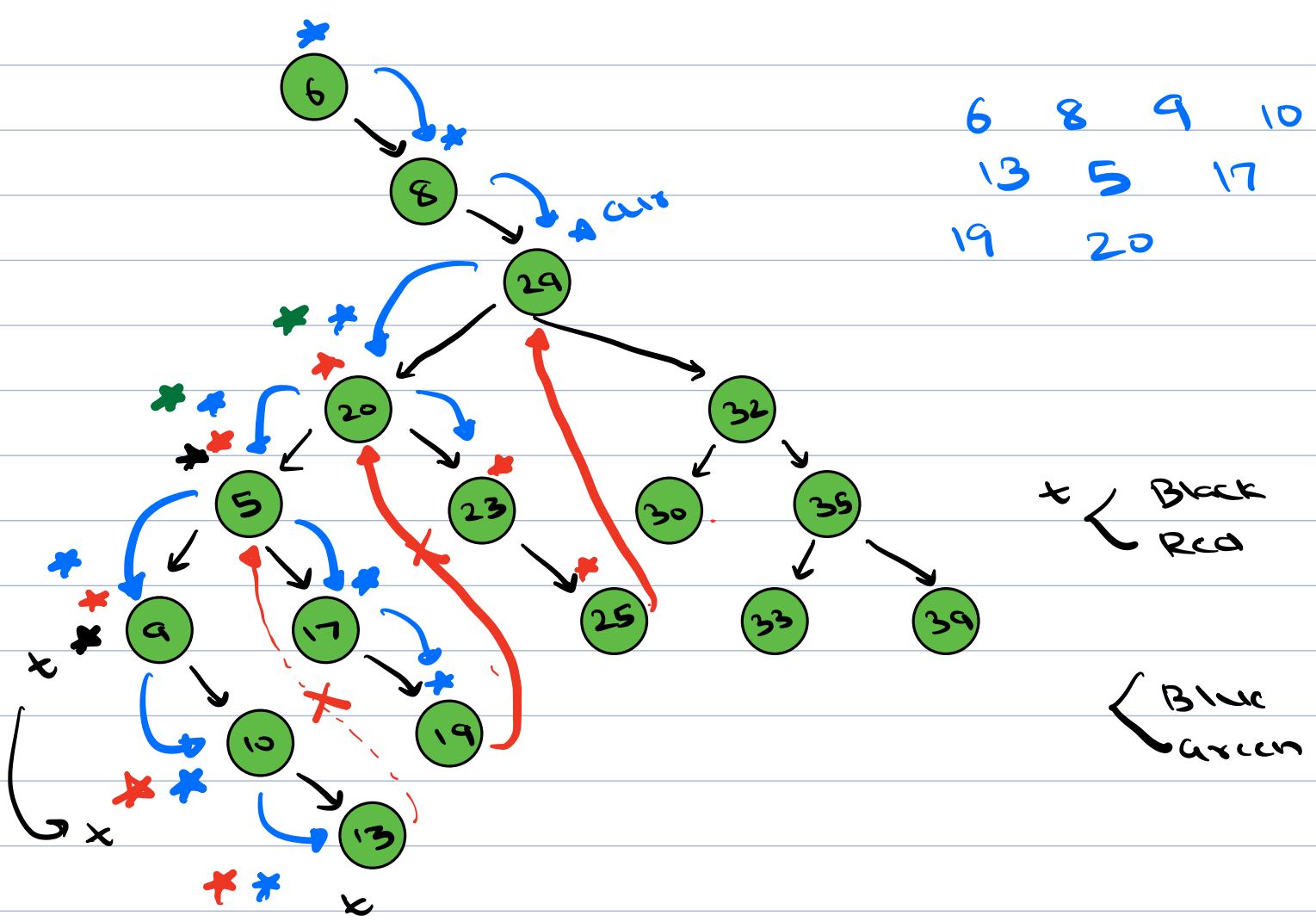
```
            print (cur.data)
```

```
            cur = cur.right
```

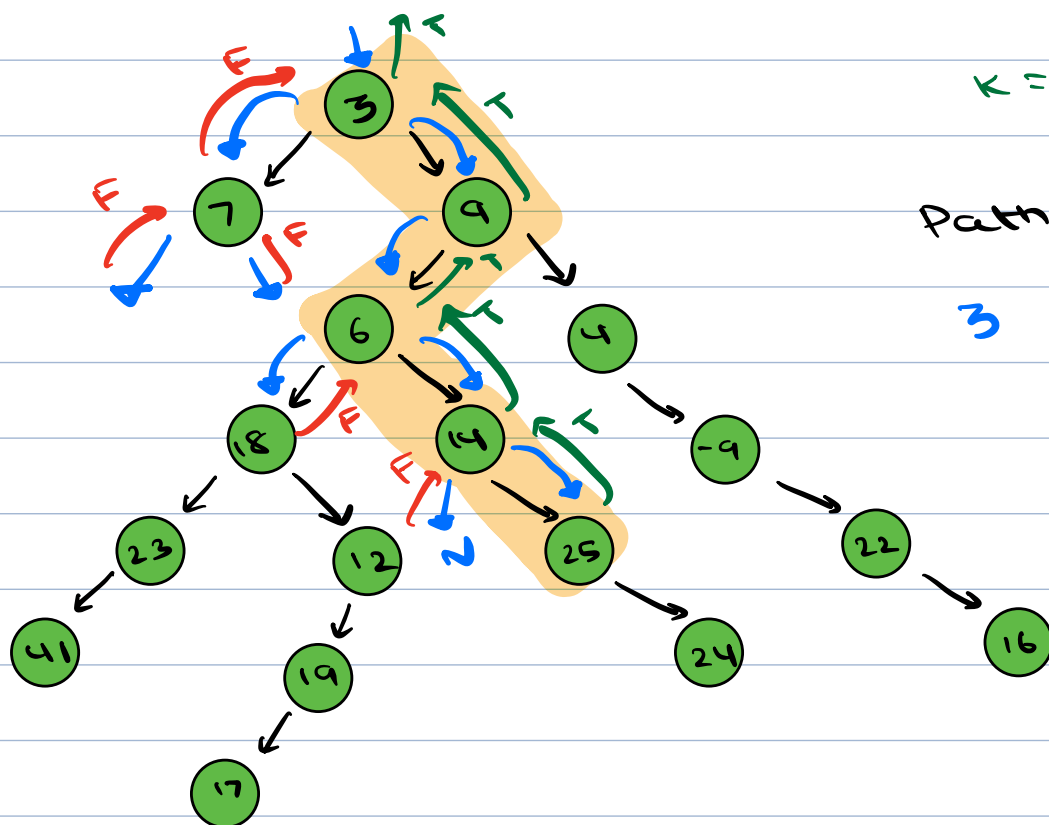
```
    }
```

```
}
```

TC: O(N)
SC: O(1)



3. Search K in Binary Tree



$K = 25$

Path (3 → 25)

3 9 6 14 25

// Given root, find K in tree

bool search (Node root, int K) {

if (root == NULL)

return false

if (root->data == K)

return true

return (search (root->left, K) ||
search (root->right, K))

TC: $O(N)$

SC: $O(H)$

✓
 $\log_2 N \rightarrow N$

4. Path from root node to K

list <int> path

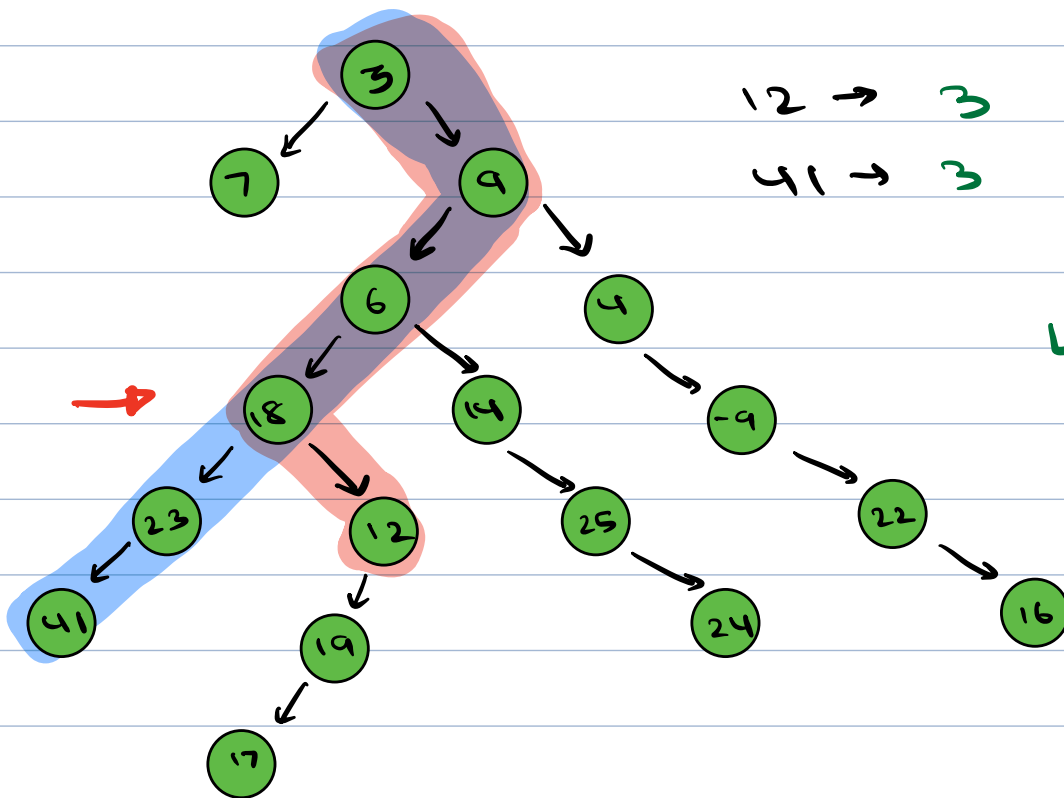
// Given root, find k in tree

```
bool search (Node root, int k) {  
    if (root == NULL)  
        return false  
    if (root.data == k) {  
        path.add (root.data)  
        return true  
    }  
    if (search (root.left, k) == true ||  
        search (root.right, k) == true) {  
        path.add (root.data)  
        return true  
    }  
    else  
        return false  
}
```

TC: $O(N)$
SC: $O(H)$
✓
 $\log_2 N \rightarrow N$

* Your path is in reverse order
(k \rightarrow root)
↓
Reverse it

Lowest Common Ancestor in Binary Tree

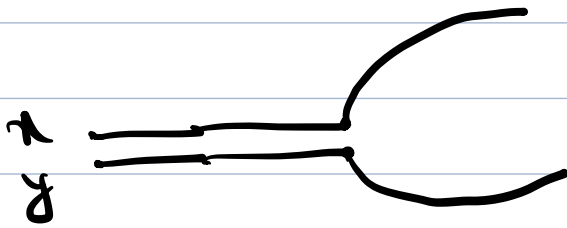


12 → 3 9 6 18 12
41 → 3 9 6 18 23 41

$LCA(41, 12) = 18$

ans = 18

$LCA(18, 41) = 18$



18 : 3 9 6 18
41 : 3 9 6 18 23 41

Last node where path diverge → LCA

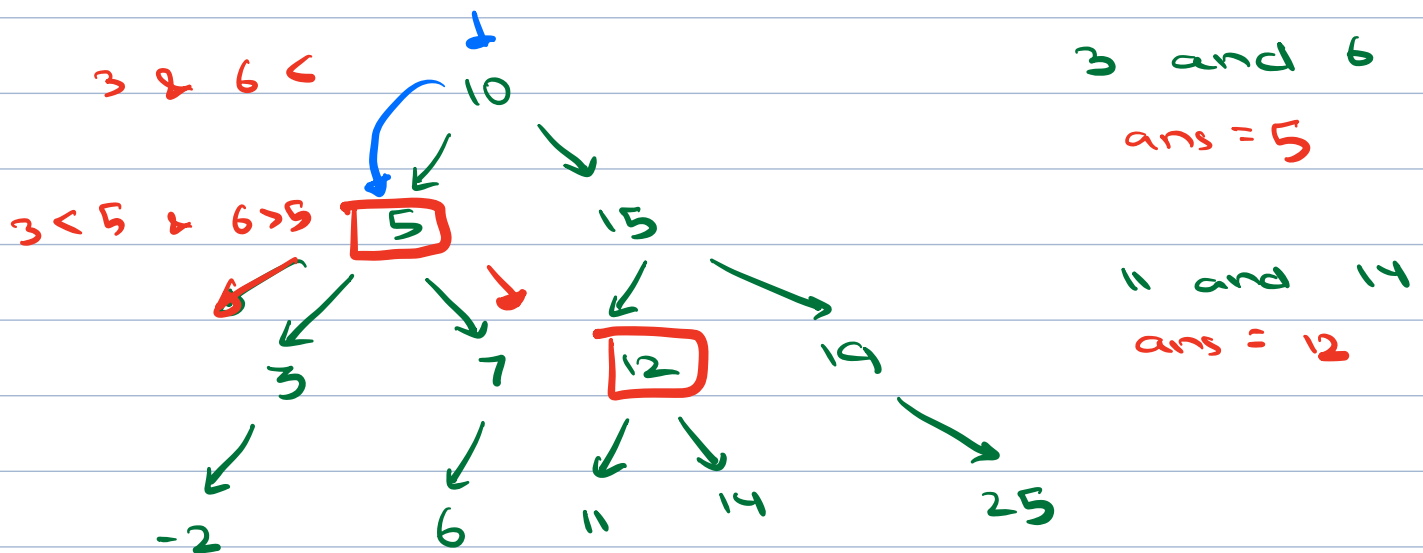
$LCA(x, y)$

- ① Path (root → x)
- ② Path (root → y)
- ③ Traversing both paths simultaneously until first non-common node → break

TC : $O(N)$

SC : $O(H)$

LCA of 2 nodes in BST



Node cur = root

while (cur != NULL) <

if (x < cur.data & y < cur.data)

cur = cur.left

else if (x > cur.data & y > cur.data)

cur = cur.right

else

return cur

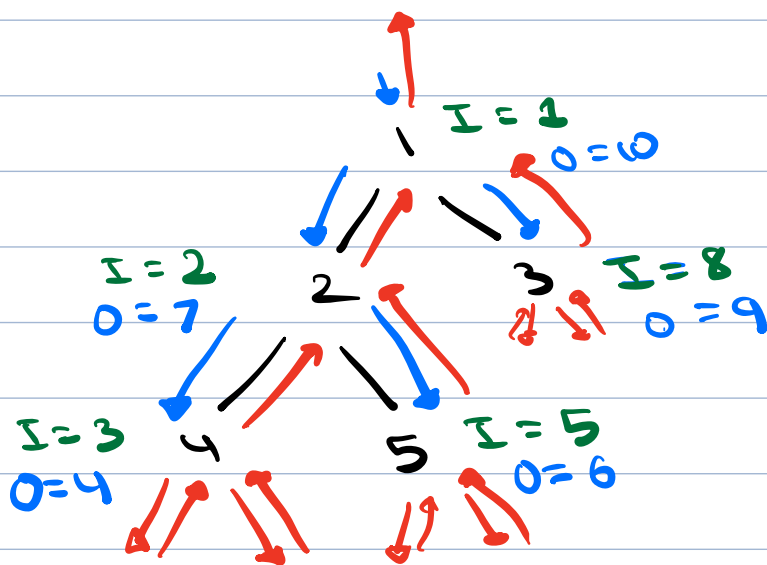
>

LCA(x, y)

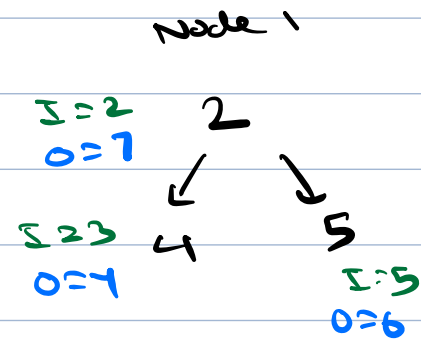
Tc: O(H)

\downarrow
 $\log_2 N \rightarrow N$

In time and out time



T = x z y x b b 7 8 9
~~11~~



Node 1

Node 2

In 1

<

In 2

Out 1

>

Out 2

Node 2 is in subtree of Node 1
 Hashmap < Node, int > intime, outtime

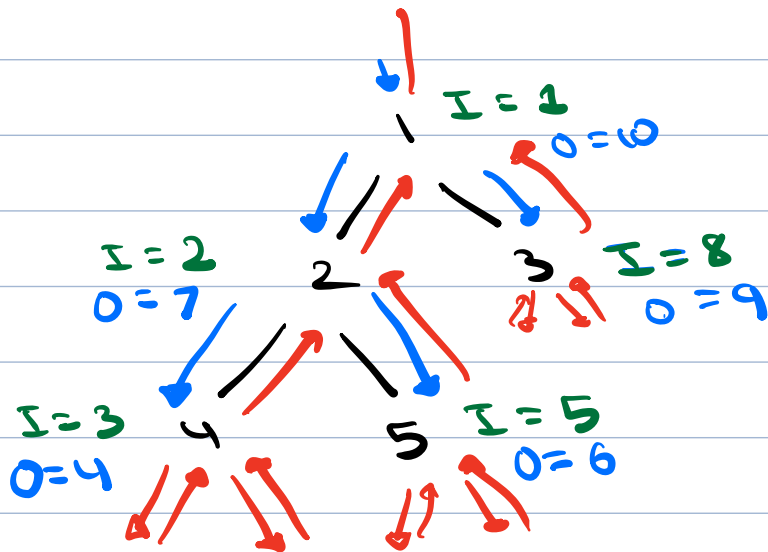
T = 1

```
void dfs (Node root) {
    if (root == NULL)
        return;

    intime [root] = T
    T++
    dfs (root.left)
    dfs (root.right)
    outtime [root] = T
    T++
}
```

Find LCA

- Calculate in time, outtime
- Store parent of each node
- Store depth of each node

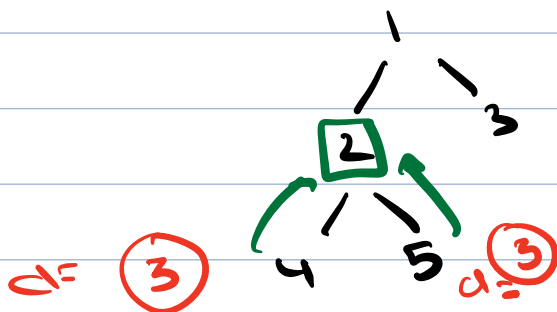


$$LCA(2, 5) = 2$$

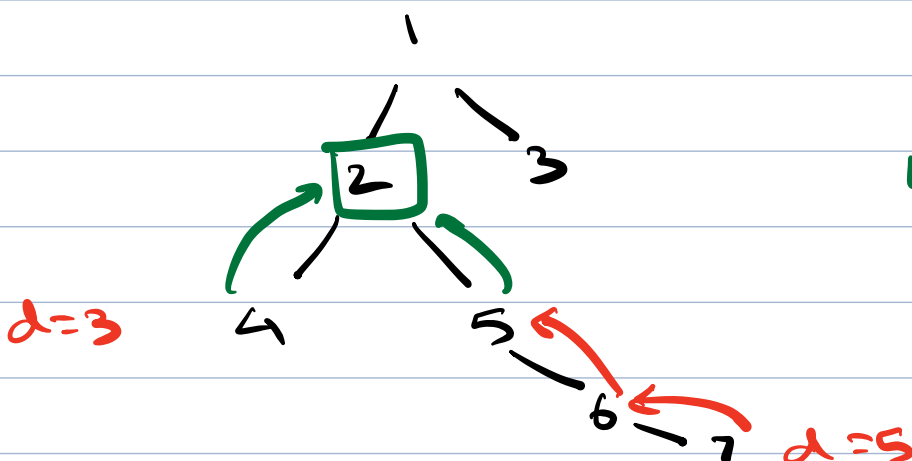
5 is in subtree of 2
 ↓ Node 2 ↓ Node 1

$$Out\ 2 < Out\ 1$$

$$In\ 2 > In\ 1$$



$$LCA(4, 5) = 2$$



$$LCA(4, 7) = 2$$

Node find LCA (x, y)

if (inTime[x] <= inTime[y] &&
outTime[x] >= outTime[y]) <

if y is
subtree
of x

return x

else if (inTime[x] >= inTime[y] &&
outTime[x] <= outTime[y]) <

return y

else <

while (depth[x] > depth[y])
x = parent[x]

while (depth[x] < depth[y])
y = parent[y]

while (x != y) <
x = parent[x]
y = parent[y]

return x

Precomputation \rightarrow TC : $O(N)$ SC : $O(N)$

Query for LCA \rightarrow TC : $O(1)$ SC : $O(1)$
