

HONORS REAL ANALYSIS LECTURE 21

ROHAN KARAMEL

ABSTRACT. This lecture reviews material to prepare for the upcoming second midterm.

Theorem. *Let G be a connected d -regular non-Eulerian graph. If \bar{G} is connected, then \bar{G} is Eulerian.*

Proof. We know that Eulerian is equivalent to being connected and having all vertices of even degree. Since G is connected and d -regular, we only need to show that d must be even. We also know that we cannot have a odd-regular graph on odd vertices, so n must be even. Thus \bar{G} has $n - 1 - d$ edges, therefore d must be even as the number of vertices cannot be odd. Thus, \bar{G} is Eulerian. \square

Theorem. *If $G - |S|$ has more than $|S|$ components, then G is not hamiltonian.*

Proof. If for all non-adjacent vertices u, v in G , $\deg(u) + \deg(v) \geq n$. Then, $\delta(G) \geq \frac{n}{2} \implies G$ is hamiltonian. \square

Theorem. *There exists a graph with $\kappa(G) \geq n$ but G is not hamiltonian.*

Proof. Consider $K_{n,n+1}$, $\kappa(G) = n$ but G is not hamiltonian. \square

Theorem. *There exists a graph with $\kappa(G) \geq n$ but G does not contain a hamiltonian path.*

Proof. Consider $K_{n,n+2}$ where $\kappa(G) = n$ but G does not contain a hamiltonian path. \square

Theorem. *Every tree has at most 1 perfect matching.*

Proof. Recall that a perfect matching is a matching that uses every vertex. Suppose for the sake of contradiction that there exists at least 2 perfect matchings. Then, there must exist a vertex that is matched to 2 different vertices in the 2 perfect matchings. \square

Theorem. *The number of spanning trees is equal to*

$$\frac{1}{n} \lambda_1 \lambda_2 \dots \lambda_{n-1}$$

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Example. Consider the adjacency matrix below:

$$\begin{bmatrix} 2 & -1 & 0 & -1 \\ -1 & 2 & -1 & 0 \\ 0 & -1 & 2 & -1 \\ -1 & 0 & -1 & 2 \end{bmatrix}$$

Theorem. We know that $\kappa(G) \leq \lambda(G) \leq \delta(G) \leq \lfloor \frac{2m}{n} \rfloor$. If G is a tree, then all terms are equal.

Proof. If G is a tree, then $\kappa(G) = \lambda(G) = \delta(G) = n - 1$. \square