Java Concurrency Utilities

Based on JavaOne talk given by **David Holmes & Brian Goetz**

Overview

- Rationale and goals for JSR 166
 - Java community process concurrency utilities
- Executors thread pools and scheduling
- Futures
- Concurrent Collections
- Locks, conditions and synchronizers
- Atomic variables

Why Concurrency Utilities

- Java's built-in concurrency primitives wait(), notify(), and synchronized:
 - Hard to use correctly
 - Easy to use incorrectly
 - Too low level for many applications
 - Can lead to poor performance if used incorrectly
 - Leave out lots of useful concurrency constructs

Goals

- Provide efficient, correct & reusable concurrency building blocks
- Enhance scalability, performance, readability, maintainability, and thread-safety of concurrent Java applications

Background

- Found in java.util.concurrent
 - based on Doug Lea's EDU.oswego.cs.dl.util.concurrent package
- APIs take advantage of native JVM constructs & Java Memory Model guarantees specified in JSR 133

Building Blocks

- Executor, ThreadPool, and Future
- Concurrent collections:
 - BlockingQueue, ConcurrentHashMap, CopyOnWriteArray
- Locks and Conditions
- Synchronizers: Semaphores, Barriers, etc.
- Atomic Variables
 - Low-level compare-and-set operation

Executor

- Standardizes asynchronous invocation
- Separates job submission from execution policy
 - anExecutor.execute(aRunnable)
 - not new Thread(aRunnable).start()
- Two code styles supported:
 - Actions: Runnables
 - Functions: Callables
 - Also has lifecycle mgmt: e.g., cancellation, shutdown
- Executor usually created via **Executors** factory class
 - Configures ThreadPoolExecutor
 - Customizes shutdown methods, before/after hooks, saturation policies, queuing

Executor & Executor Service

ExecutorService adds lifecycle management to Executor

```
public interface Executor {
  void execute(Runnable command);
}
public interface ExecutorService extends Executor {
  void shutdown();
  List<Runnable> shutdownNow();
  boolean isShutdown();
  boolean isTerminated();
  boolean awaitTermination( long timeout, TimeUnit unit);
  // other convenience methods for submitting tasks
}
```

Creating Executors

• Executors factory methods

```
public class Executors {
    static ExecutorService newSingleThreadedExecutor();
    static ExecutorService newFixedThreadPool(int n);
    static ExecutorService newCachedThreadPool(int n);
    static ScheduledExecutorService newScheduledThreadPool(int n);
    // additional versions & utility methods
```

(Not) Executor Example

• Thread per message Web Server (no limit on thread creation)

```
class WebServer {
  public static void main( String [] args) {
    ServerSocket socket = new ServerSocket ( 80 );
  while ( true ) {
    final Socket connection = socket.accept();
    Runnable r = new Runnable () {
      public void run () {handleRequest(connection);}
    };
    new Thread (r).start();
  }
}
```

Executor Example

- Thread pool web server better resource management
- class WebServer {
 Executor pool = Executors.newFixedThreadPool(7);
 public static void main(String[] args) {
 ServerSocket socket = new ServerSocket(80);
 while (true) {
 final Socket connection = socket.accept();
 Runnable r = new Runnable() {
 public void run() {handleRequest(connection);}
 }
 }
 pool.execute(r);
 }
 }

Future and Callable

 Callable is functional analog of Runnable interface Callable<V> {

```
interface Callable<V> {
    V call() throws Exception;
}
```

}

Future represents result of asynchronous computation

```
interface Future<V> {
   V get() throws InterruptedException, ExecutionException;
   V get(long timeout, TimeUnit unit);
   boolean cancel(boolean mayInterrupt);
   boolean isCancelled();
   boolean isDone();
```

Using Futures

- Client initiates asynchronous computation via oneway message
- Client receives a "handle" to the result: a Future
- Client does other work while waiting for result
- When ready Client requests result from Future, blocking if necessary until result is available
- Client uses result

FutureTask

- A cancellable asynchronous computation
- A base implementation of Future
- Can wrap a Callable or Runnable
 - Allows FutureTask to be submitted to an Executor

Future Example

- See: FutureStringReverser.java
- See: FutureTaskStringReverser.java

Another Future Example

• Implementing a cache with Future

ScheduledExecutorService

- For deferred and recurring tasks, can schedule
 - Callable or Runnable to run once with a fixed delay after submission
 - Schedule a Runnable to run periodically at a fixed rate
 - Schedule a Runnable to run periodically with a fixed delay between executions
- Submission returns a ScheduledFutureTask handle which can be used to cancel the task
- Like **Timer**, but supports pooling and is more robust

Concurrent Collections

- Pre-1.5 Java class libraries had few concurrent (vs. synchronized) classes
 - Synchronized collections:
 - · Hashtable, Vector, and Collections.synchronized*
 - Often required locking during iteration
 - Locking becomes is a source of contention
- Java 1.5 concurrent collections:
 - Allow multiple operations to overlap
 - Some differences in semantics

Queues

Queue interface added to java.util

```
interface Queue<E> extends Collection<E> {
   boolean offer(E x); // try to insert
   E poll(); //retrieve and remove. Return null if empty
   E remove() throws NoSuchElementException; //retrieve and remove
   E peek(); // retrieve, don't remove. Return null if empty
   E element() throws NoSuchElementException; //retrieve, don't remove
}
```

- Thread-safe and non-thread safe implementations
 - Non-thread-safe LinkedList
 - Non-thread-safe PriorityQueue
 - Thread-safe non-blocking ConcurrentLinkedQueue

Blocking Queues

- Extends Queue to provide blocking operations
 - Retrieval: wait for queue to become nonempty
 - Insertion: wait for capacity to be available
- Common in producer-consumer designs
- Can be bounded or unbounded
- Implementations provided:
 - LinkedBlockingQueue (FIFO, may be bounded)
 - PriorityBlockingQueue (priority, unbounded)
 - ArrayBlockingQueue (FIFO, bounded)
 - SynchronousQueue (rendezvous channel)
- See API for details

Producer-Consumer Examples

- See:
 - ProducerConsumerPrimitive.java (wait/notify)
 - ProducerConsumerConcUtil.java (BlockingQueue)

Concurrent Collections

- ConcurrentHashMap Concurrent (scalable) alternative to Hashtable or Collections.synchronizedMap
 - Multiple reads can overlap each other
 - Reads can overlap writes
 - Retrieval operations reflect the results of the most recently completed update operations holding at onset of operation
 - Up to 16 writes can overlap
 - Iterators do not throw **ConcurrentModificationException**
- CopyOnWriteArrayList
 - Optimized for case where iteration is much more frequent than insertion or removal. E.g., event listeners

Performance Comparison

- ConcurrentHashMap vs. Collections.synchronizedMap
- See HashMapPerfTest.java
- Note: incrementCount() is not safe

Locks and Lock Support

- High-level locking interface
- Adds non-blocking lock acquisition

```
interface Lock {
  void lock();
  void lockInterruptibly() throws IE;
  boolean tryLock();
  boolean tryLock(long time,TimeUnit unit) throws IE;
  void unlock();
  Condition newCondition() throws UnsupportedOperationException;
}
```

ReentrantLock

- Flexible, high-performance lock implementation
- Implements a reentrant mutual exclusion lock (like Java intrinsic locks) but with extra features
 - Can interrupt a thread waiting to acquire a lock
 - Can specify a timeout while waiting for a lock
 - Can poll for lock availability
 - Can have multiple wait-sets per lock via the Condition interface
- Outperforms built-in monitor locks in most cases, but slightly less convenient to use (requires finally block to release lock)

Lock Example

- Locks not automatically released
 - Must release lock in **finally** block

```
Lock lock = new ReentrantLock();
...
lock.lock();
try {
    // perform operations protected by lock
} catch (Exception ex) {
    // restore invariants
} finally {
    lock.unlock();
}
```

ReadWrite Locks

- ReadWriteLock interface defines a pair of locks;
 - one for readers; one for writers

```
interface ReadWriteLock {
  Lock readLock();
  Lock writeLock();
}
```

- ReentrantReadWriteLock class
 - Multiple readers, single writer
 - Allows writer to acquire read lock
 - Allows writer to downgrade to read lock
 - Supports "fair" and "non-fair" (default) acquisition

Read/Write Lock Example

```
class RWDictionaryRWL {
    private final Map<String, Data> m = new TreeMap<String, Data>();
    private final ReentrantReadWriteLock rwl = new ReentrantReadWriteLock();
    private final Lock r = rwl.readLock();
    private final Lock w = rwl.writeLock();
    public Data get(String key) {
        r.lock();
        try { return m.get(key); }
        finally { r.unlock(); }
    }
    public Data put(String key, Data value)
        w.lock();
        try { return m.put(key, value); }
        finally { w.unlock(); }
    }
}
```

Read/Write Lock Example

- See
 - RWDictionary .java & RWDictionary RWL.java

Condition

• Condition lets you wait for a condition to hold (like wait), but adds several features

```
interface Condition {
  void await() throws IE;
  boolean await( long time, TimeUnit unit ) throws IE;
  long awaitNanos( long nanosTimeout) throws IE;
  void awaitUninterruptibly()
  boolean awaitUntil( Date deadline) throws IE;
  void signal();
  void signalAll();
}
```

Condition (cont.)

- Many improvements over wait()/notify()
 - Multiple conditions per lock
 - Absolute and relative time-outs
 - Timed waits tell you why you returned
 - Convenient uninterruptible wait

Condition Example

```
class BoundedBufferCond {
    Lock lock = new ReentrantLock ();
    Condition notFull = lock.newCondition();
    Condition notEmpty = lock.newCondition();
    Object[] items = new Object[100];
    int putptr, takeptr, count;

public void put( Object x) throws IE {
    lock.lock();
    try {
      while (count == items.length) notFull.await();
      items[putptr] = x;
      if (++putptr == items.length) putptr = 0;
      ++count;
      notEmpty.signal();
    } finally { lock.unlock(); }
}
```

Condition Example (cont.)

```
public Object take() throws IE {
    lock.lock();
    try {
        while (count == 0) notEmpty.await();
        Object x = items[takeptr];
        if (++takeptr == items.length) takeptr = 0;
        --count;
        notFull.signal();
        return x;
        } finally { lock.unlock(); }
    }
}
```

Condition Example (cont.)

- Previous example in BoundedBufferCond.java
- See also: BoundedBufferPrim.java

Synchronizers

- Utilities for coordinating access and control
- **CountDownLatch** allows one or more threads to wait for a set of threads to complete an action
- CyclicBarrier allows a set of threads to wait until they all reach a specified barrier point
- **Semaphore** Dijkstra counting semaphore, managing some number of permits
- Exchanger allows two threads to rendezvous and exchange data, such as exchanging an empty buffer for a full one

CountDownLatch

- Latching variables are conditions that once set never change
- Often used to start several threads, but have them wait for a signal before continuing
- See: CountDownLatchTest.java

CyclicBarrier

- Allows threads to wait at a common barrier point
- Useful when a fixed-sized party of threads must occasionally wait for each other
- Cyclic Barriers can be re-used after threads released
- Can execute a Runnable once per barrier point
 - After the last thread arrives, but before any are released
 - Useful for updating shared-state before threads continue
- See: CyclicBarrierEx1.java & CyclicBarrierEx2.java

Semaphore

- Semaphore maintain a logical set of permits
- acquire() blocks until a permit is free, then takes it
- release() adds a permit, releasing a blocking acquirer
- Often used to restrict the number of threads that can access some resource
 - But can be used to implement many sync disciplines
- See: SemaphoreTunnel.java & SemaphoreBuffer.java

Exchanger

- Synch. point where two threads exchange objects
- A bidirectional SynchronizedQueue
- Each thread presents some object on entry to the exchange() method, and receives the object presented by the other thread on return
- See ExchangerTest.java

Atomic Variables

- · Holder classes for scalars, references and fields
- Supports atomic operations
 - Compare-and-set (CAS)
 - Get and set and arithmetic (where applicable)
- Ten main classes: { int, long, ref } X { value, field, array }
 - E.g. **AtomicInteger** useful for counters, sequences, statistics
- Essential for writing efficient code on MPs
 - Nonblocking data structures & optimistic algorithms
 - Reduce overhead/contention updating "hot" fields
- JVM uses best construct available on platform
 - CAS, load-linked/store-conditional, locks

Atomic Variables

• See: CounterTest.java