

# TD - Semantics and Verification ${ m VI-~LTL}$ Friday 4th March 2022

TA: Ralph Sarkis ralph.sarkis@ens-lyon.fr

Recall that:

- $\llbracket \Diamond \varphi \rrbracket_{\rho} = \{ \sigma \in (2^{AP})^{\omega} \mid \exists i, \sigma \upharpoonright i \in \llbracket \varphi \rrbracket_{\rho} \}$
- $\llbracket \Box \varphi \rrbracket_{\rho} = \{ \sigma \in (2^{AP})^{\omega} \mid \forall i, \sigma \upharpoonright i \in \llbracket \varphi \rrbracket_{\rho} \}$
- $\llbracket \varphi \cup \psi \rrbracket_{\rho} = \{ \sigma \in (2^{AP})^{\omega} \mid \exists i, \sigma \upharpoonright i \in \llbracket \psi \rrbracket_{\rho}, \forall j < i, \sigma \upharpoonright j \in \llbracket \varphi \rrbracket_{\rho} \}$
- $\varphi \mathsf{W} \psi := \neg(\neg \psi \mathsf{U} \neg (\varphi \lor \psi))$

Moreover, a state s of a transitition system satisfies a LTL formula  $\varphi$  with parameter  $\rho$  if and only if  $\operatorname{Traces}(s) \subseteq \llbracket \varphi \rrbracket_{\rho}$ .

#### Exercise 1.

Let L be a complete lattice and let  $f: L \to L$  be a monotone function. Show that  $\mu(f) = \bigwedge \{a \in L \mid A \in L \}$  $L \mid f(a) \leq a$  (resp.  $\nu(f) = \bigvee \{a \in L \mid a \leq f(a)\}\)$  is the least fixpoint (resp. greatest fixpoint) of f.

# Exercise 2.

Show that:

- 1.  $\neg(\varphi \mathsf{W} \psi) \equiv \neg \psi \mathsf{U} (\neg \varphi \land \neg \psi)$
- 2.  $\neg(\varphi \cup \psi) \equiv \neg \psi \cup (\neg \varphi \wedge \neg \psi)$
- 3.  $\bigcirc(\varphi \cup \psi) \equiv \bigcirc\varphi \cup \bigcirc\psi$
- 4.  $\Diamond \phi \equiv \neg \Box \neg \phi$
- 5.  $\Diamond \phi \equiv \phi \lor \bigcirc \Diamond \phi$

#### Exercise 3.

Let  $\phi$  be a formula with parameters  $\rho$ . We define  $\phi_{\square}(X) = \phi \wedge \bigcirc X$ . Show that for every valuation  $\rho$ ,  $\llbracket \Box \phi \rrbracket_{\rho}$  is the greatest fixpoint of  $\llbracket \phi_{\Box} \rrbracket_{\rho}$ .

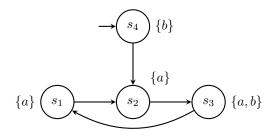
# Exercise 4.

Show that:

- 1.  $\llbracket \top \mathsf{U} \varphi \rrbracket_{\rho} = \{ \sigma \in (2^{AP})^{\omega} \mid \exists i, \sigma \upharpoonright i \in \llbracket \varphi \rrbracket_{\rho} \}$ 2.  $\llbracket \varphi \mathsf{W} \perp \rrbracket_{\rho} = \{ \sigma \in (2^{AP})^{\omega} \mid \forall i, \sigma \upharpoonright i \in \llbracket \varphi \rrbracket_{\rho} \}$

## Exercise 5.

Consider the following transition system over the set of atomic propositions  $\{a, b\}$ :



Indicate for each of the following LTL formulae the set of states for which these formulae are fulfilled:

A.  $\bigcirc a$  C.  $\Box b$  E.  $\Box (b \cup a)$  B.  $\bigcirc \bigcirc \bigcirc b$  D.  $\Box \lozenge a$  F.  $\lozenge (a \cup b)$ 

## Exercise 6.

A formula  $\phi$  is said to be in positive normal form if the negations is only applied to atoms or variables (for instance, with  $a, b \in AP$ ,  $a \land \neg b$  is in positive normal form but  $\neg (a \lor b)$  is not).

- 1. Show that every formula is equivalent to a formula in positive normal form.
- 2. We replace the operator W by the release operator R satisfying  $\neg(\phi \ R \ \psi) \equiv (\neg \phi) \ U \ (\neg \psi)$ . Show that, with this operator, every formula is still equivalent to a formula in positive normal form.
- $3. \,$  Compare the two methods regarding the size of the formulae.