

# Modelling Nondeterminism, Probability and Termination

The story of Ralph's internship with his amazing supervisors Matteo and Valeria.

Ralph Sarkis

ENS de Lyon

April 12th, 2021

Overview of the presentation:

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- ▶ **Set** monads and equational theories modelling

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- ▶ **1Met** monads and quantitative equational theories modelling



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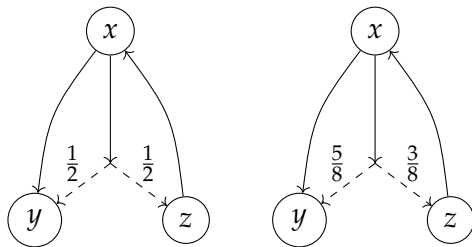
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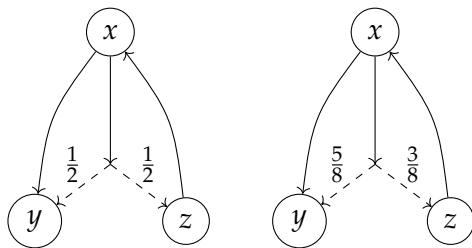
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- ▶ Combining nondeterminism, probability and termination **NEW!**

# Motivation



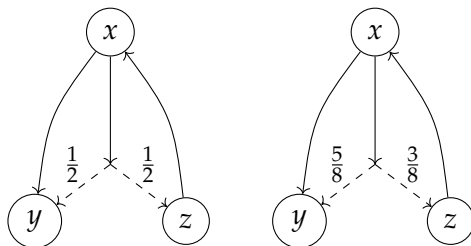
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## Questions

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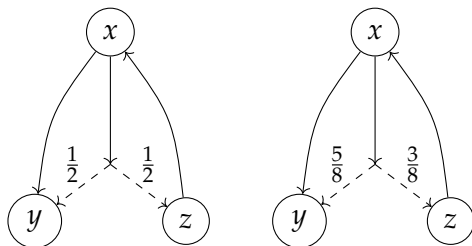


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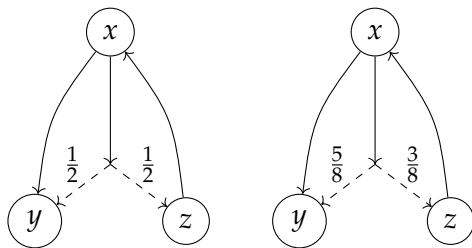
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Categorically

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A link between these two pictures: **algebraic presentations of monads**.

# Nondeterministic Choices (in **Set**)

Categorically



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- Powerset monad  $\mathcal{P}$ :

$$\mathcal{P}(X) = \{\text{non-empty finite subsets of } X\}$$

$$\eta_X = x \mapsto \{x\}$$

$$\mu_X = \mathcal{F} \mapsto \bigcup_{S \in \mathcal{F}} S$$

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- Theory of (sup-)semilattices:  
A binary operation  $\oplus$  satisfying

$$x \oplus x = x \quad I$$

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- Operational semantics

$$\frac{q \xrightarrow{a} r \quad q \xrightarrow{a} s}{q \xrightarrow{a} r \oplus s} \quad \frac{q \xrightarrow{a} r \quad q' \xrightarrow{a} r'}{q \oplus q' \xrightarrow{a} r \oplus r'} \\ \frac{q \downarrow \text{accept}}{q \oplus r \downarrow \text{accept}}$$

# Probabilistic Choices (in **Set**)

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$$\mathcal{D}(X) = \{\text{finitely supported} \\ \text{probability distributions on } X\}$$

$$\eta_X = x \mapsto \delta_x \text{ (Dirac)}$$

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- Operational semantics

$$\frac{\frac{q \xrightarrow{a,p} r \quad q \xrightarrow{a,1-p} s}{q \xrightarrow{a} r +_p s} \quad \frac{q \xrightarrow{a} r \quad q' \xrightarrow{a} r'}{q +_p q' \xrightarrow{a} r +_p r'}}{\frac{q \downarrow o(q) \quad r \downarrow o(r)}{q +_p r \downarrow po(q) + (1-p)o(r)}}$$

# Termination (in **Set**)

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- Maybe monad  $\cdot + \mathbf{1}$ :

$$X + \mathbf{1} = X \sqcup \{\star\}$$

$$\eta_X = \text{inl}^{X+\mathbf{1}}$$

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$$\frac{q \nrightarrow}{q \rightarrow \star}$$

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- Operational semantics: Combine all previous rules.

[BSV19]

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## Equationally

- ▶ Equations become quantitative

$$x =_{\varepsilon} y \quad (\varepsilon \in [0, 1])$$

- ▶ Quantitative equational logic

# Quantitative Equational Logic

$\vdash t =_0 t$	Refl
$\vdash t =_1 s$	1-bounded
$t =_\varepsilon s \vdash s =_\varepsilon t$	Symm
$t =_\varepsilon s, s =_\delta u \vdash t =_{\varepsilon+\delta} u$	Triang
$t =_\varepsilon s \vdash t =_{\varepsilon+\delta} s$	Max (for $\delta > 0$ )
$\forall \varepsilon > \delta, t =_\varepsilon s \vdash t =_\delta s$	Arch (for $\delta \geq 0$ )
$\forall 1 \leq i \leq n, t_i =_\varepsilon s_i \vdash f(t_1, \dots, t_n) =_\varepsilon f(s_1, \dots, s_n)$	Nexp (for $f$ an $n$ -ary)
$\Gamma \vdash t =_\varepsilon s \implies \sigma(\Gamma) \vdash \sigma(t) =_\varepsilon \sigma(s)$	Subst (for $\sigma$ a substitution)
$\phi \in \Gamma \implies \Gamma \vdash \phi$	Assumption

Introduced by Mardare, Panangaden and Plotkin [MPP16].

# Nondeterminism Choices (in **1Met**)

Categorically



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## Categorically

- Hausdorff lifting of  $d : X \times X \rightarrow [0, 1]$  into  $H(d) : \mathcal{P}(X) \times \mathcal{P}(X) \rightarrow [0, 1]$ :

$$H(d)(S, T) = \max \left\{ \sup_{s \in S} \inf_{t \in T} d(s, t), \sup_{t \in T} \inf_{s \in S} d(s, t) \right\}$$

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- ▶ Hausdorff lifting of  $\mathcal{P}$ :  $\widehat{\mathcal{P}}(X, d) = (\mathcal{P}(X), H(d))$ .

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- ▶ Theory of quantitative (sup-)semilattices: binary operation  $\oplus$  satisfying

$$\vdash x \oplus x =_0 x \quad I$$

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$$\vdash (x \oplus y) \oplus z =_0 x \oplus (y \oplus z) \quad A$$

$$x_1 =_{\varepsilon_1} y_1, x_2 =_{\varepsilon_2} y_2 \vdash x_1 \oplus x_2 =_{\max\{\varepsilon_1, \varepsilon_2\}} y_1 \oplus y_2 \quad H$$

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$$K(d)(\varphi, \psi) = \inf_{\omega} \sum_{x_1, x_2 \in X} \omega(x_1, x_2) \cdot d(x_1, x_2),$$

where  $\omega$  ranges over couplings of  $\varphi$  and  $\psi$ .

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## Categorically

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- Theory of quantitative convex algebras

$$\vdash x +_p x =_0 x \quad I_p$$

$$\vdash x +_p y =_0 y +_p x \quad C_p$$

$$\vdash (x +_q y) +_p z =_0 x +_{pq} (y +_{\frac{p\bar{q}}{pq}} z) \quad A_p$$

$$x_1 =_{\varepsilon_1} y_1, x_2 =_{\varepsilon_2} y_2 \vdash x_1 +_p x_2 =_{p\varepsilon_1 + (1-p)\varepsilon_2} y_1 +_p y_2 \quad K_p$$

# Termination (in **1Met**)

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- The maybe monad  $\cdot + \hat{\mathbf{1}}$  puts  $\star$  at distance 1 from every other point:

$$X + \hat{\mathbf{1}} = X \sqcup \{\star\}$$

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- Theory of pointed spaces:

$$\vdash x =_1 \star$$

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## Equationally

- Quantitative theory of convex semilattices: combine all previous axioms plus

$$\vdash (x \oplus y) +_p z =_0 (x +_p z) \oplus (y +_p z) \qquad D_p$$

[MV20]

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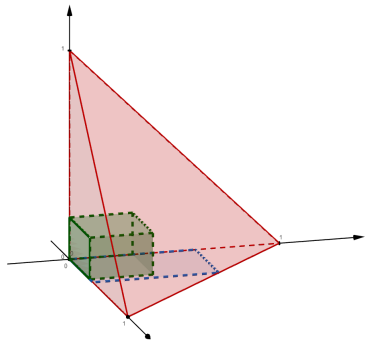


Figure 1: Examples of  $\perp$ -closed sets.

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## On 1Met

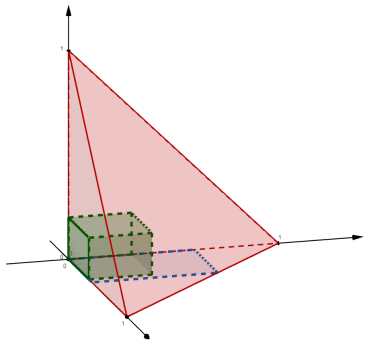


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## On 1Met

- ▶ The theory of quantitative convex semilattices with bottom corresponds to  $\mathcal{C}^\downarrow$  with the Kantorovitch and Hausdorff liftings of the metrics.

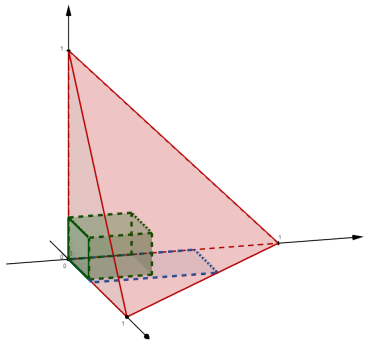


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# Summary

On Set		On 1Met	
$\mathcal{P}$	Semilattices	$\hat{\mathcal{P}}$	Quantitative semilattices
$\mathcal{D}$	Convex algebras	$\hat{\mathcal{D}}$	Quantitative convex algebras
$\cdot + \mathbf{1}$	Pointed sets	$\cdot + \hat{\mathbf{1}}$	Pointed spaces
$\mathcal{C}$	Convex semilattices (CS)	$\hat{\mathcal{C}}$	Quantitative CS
$\mathcal{C}(\cdot + \mathbf{1})$	Pointed CS	$\hat{\mathcal{C}}(\cdot + \hat{\mathbf{1}})$	Quantitative pointed CS
$\mathcal{C} + \mathbf{1}$	CS with $\perp$ and black-hole	Trivial	Quantitative CS with $\perp$ and black-hole
$\mathcal{C}^\downarrow$	CS with $\perp$	$\hat{\mathcal{C}}^\downarrow$	Quantitative CS with $\perp$

See Section VI of preprint.

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- ▶ Can our results be obtained *compositionally*? Close to answering yes with Daniela Petrişan.
- ▶ Where can quantitative algebraic reasoning be applied?

## References

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Merci !