



# Auction Algorithms

- Subham Ghosh
- Manu Viswanadhan



# Assignment Problem

- N persons and M objects and benefits  $a_{ij}$
- Match one-to-one basis
- Maximize total benefit  $\sum_{i=1}^n a_{ij_i}$



# Economic equilibrium Problem

- Match  $n$  persons and objects and maximize the profit
- Person  $i$  would want object  $j_i$ , such that

$$a_{ij_i} - p_{ji} = \max_{j=1,\dots,n} a_{ij} - p_j$$

- Equilibrium if all persons are happy



# Basic auction algorithm

- Choose object  $j_i$  with maximal value

$$j_i \in \arg \max_{j=1,\dots,n} a_{ij} - p_j$$

- Get object, increment its price

$$p_{j_i} = p_{j_i} + \gamma_i$$

$$\gamma_i = v_i - w_i$$

$$v_i = \max_j a_{ij} - p_j,$$

$$w_i = \max_{j \neq j_i} a_{ij} - p_j$$



# Points to note

- $y_i$  is always non decreasing
- The object becomes less desirable as it is bid more
- What if  $v_i$  equals  $w_i$



# Modified Auction Algorithm

- Person is almost happy if

$$a_{ij_i} - p_{ji} = \max_{j=1,\dots,n} a_{ij} - p_j - \epsilon$$

$$\gamma_i = v_i - w_i + \epsilon$$

- almost at equilibrium

# Optimality

- Total benefit within  $n\epsilon$  of the optimal value
- When weights are integers and

$$\epsilon < \frac{1}{n}$$

- optimal assignment



# $\epsilon$ -scaling

- Convergence depends on  $\epsilon$  and maximum object value ( $C$ )
- Low initial prices
  - number of rounds  $\propto C/\epsilon$
- Fasten process with  $\epsilon$ -scaling





# Centralized Algorithm

- Centralized Controller
- Bidding process in synchronized rounds
- Communication between controller and person threads through RMI calls



# Controller

- Asks the current person to bid
- If the person returns with no assignment, terminate the person
- If the person returns with an assignment, assign next round to displaced person, else go on with next person in queue



# Person

- Person finds the object offering the best value
- If this exists, assigns itself to that and returns the assignment
- Otherwise returns terminated



# Points to Note

- Overdependence on controller
- Not really distributed
- Bidding and selection can be parallelized



# Asynchronous-Parallel Algorithm

- Controller starts all person and object threads
- Responsible for termination detection

# Person

- If currently not assigned to an object, sends a bid to the object having max benefit
- If there are no objects with positive benefit, it reports terminated to the controller
- On receiving accept from the object, updates the price of the object and reports the new assignment to controller



# Object

- On receiving bid from object, if it is higher than current price, sends accept, else reject
- If bid is accepted, then informs all other objects of its new price



# Termination Detection - Controller

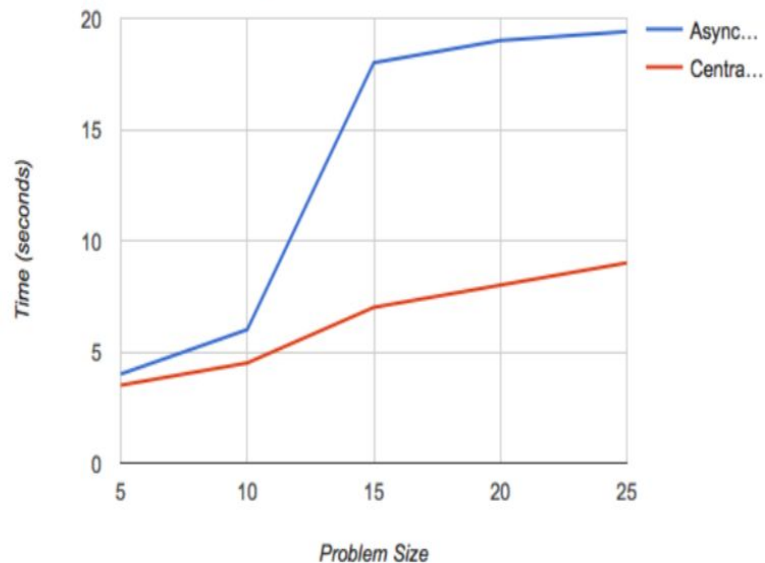
- Keeps track of the persons assigned and their respective objects.
- Process terminated when
  - All persons have an object assigned or are terminated
  - The number of assigned objects are  $\min(N_{\text{person}}, N_{\text{object}})$



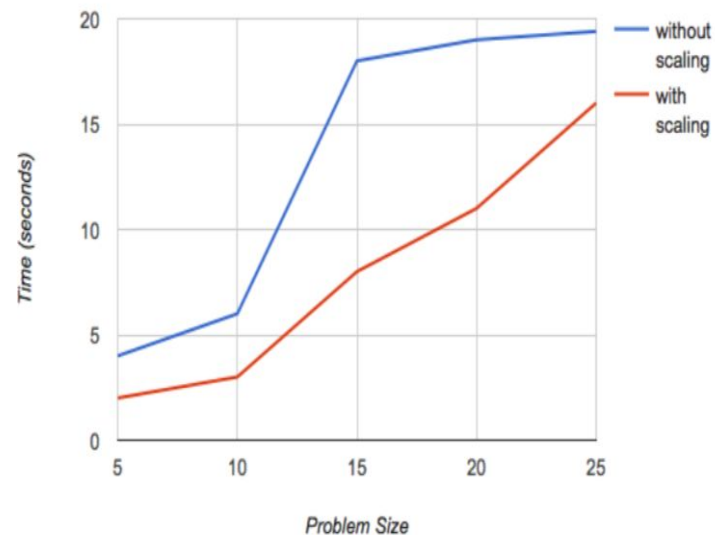


# Experiments

Centralized v Asynchronous-Parallel



Impact of Epsilon Scaling





# Demo