# Non-Interactive Hierarchical Id-based Key Agreement in MANETs



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PKG runs following three

Setup: takes security

2) Extract: using msk and

identities of nodes, PKG

generates their secret keys

Shared Key: using params,

its own secret key and peer's

identity, a node computes its

shared key with the peer

parameter as input and

outputs public parameters

(params) and master secret

algorithms:

key (msk)

### Problem Statement

Design a secure and efficient Non-Interactive, Hierarchical, Identity-based Key Agreement scheme for Mobile Ad-hoc Networks which is fully Resilient at each level against arbitrary number of node compromises

## Jargons

- MANET: is an infrastructure-less and wireless network composed of mobile nodes
- Key Agreement Protocol: allows two or more parties to agree on a shared secret key

## Ad-hoc Networks

- MANETs find application in
- > establishing Tactical Networks for Military
- > communication in disaster hit areas
- MANET nodes are constrained in:
  - > Computational capabilities
- > Communicational capabilities
- The nodes are usually mobile and have limited battery supply

### HH-KAS

- Hybrid Hierarchical scheme (HH-KAS) was introduced for key agreement in MANETs by Gennaro et al. in 2008
- HH-KAS scheme comprises:
  - > a linear hierarchical key agreement scheme at non-leaf levels, and
- > SOK key agreement scheme given by Sakai et al. at leaf level.
- HH-KAS is fully resilient at leaf level and resilient upto a threshold at non-leaf levels

#### Tools

- Let  $G_1$ ,  $G_2$ ,  $G_T$  be cyclic prime order groups then, pairing is a efficiently computable map e:  $G_1 \times G_2 \rightarrow G_T$  which satisfies:
- $\triangleright$  Bilinearity:  $\forall a, b \in F_0^*$ ,  $\forall P \in G_1$ ,  $\forall Q \in$  $G_2$ : e(aP, bQ) = e(P,  $\mathring{Q}$ )<sup>ab</sup>
- $\triangleright$  Non-degeneracy:  $e(P, Q) \neq 1$
- Basic Id One way function Scheme (BIOS) is a deterministic key pre-distribution (KPD) scheme introduced by Lee and Stinson
- BIOS achieves perfect resiliency and complete connectivity with fewer keys/node when compared to randomized KPD schemes.

### References

- Gennaro R., Halevi S., Krawczyk H., Rabin T., Reidt S., Wolthusen S.D.: Strongly-Resilient and Non-interactive Hierarchical Key-Agreement in MANETs, ESORICS, 2008
- Lee J., and Stinson D.R., Deterministic Key Predistribution Schemes for Distributed Sensor Networks, SAC, 2004
- Sakai R., Ohgishi K., Kasahara M.: Cryptosystems based
- on pairing. Cryptography and Information Security, 2000 Chatterjee S., Hankerson D., Knapp E., Menezes A.: Comparing two pairing-based aggregate signature schemes. Designs, Codes and Cryptography, 2010

# BIOS-SOK Key Agreement Scheme

PKG

Secret Keys

Shared Key between A and B:

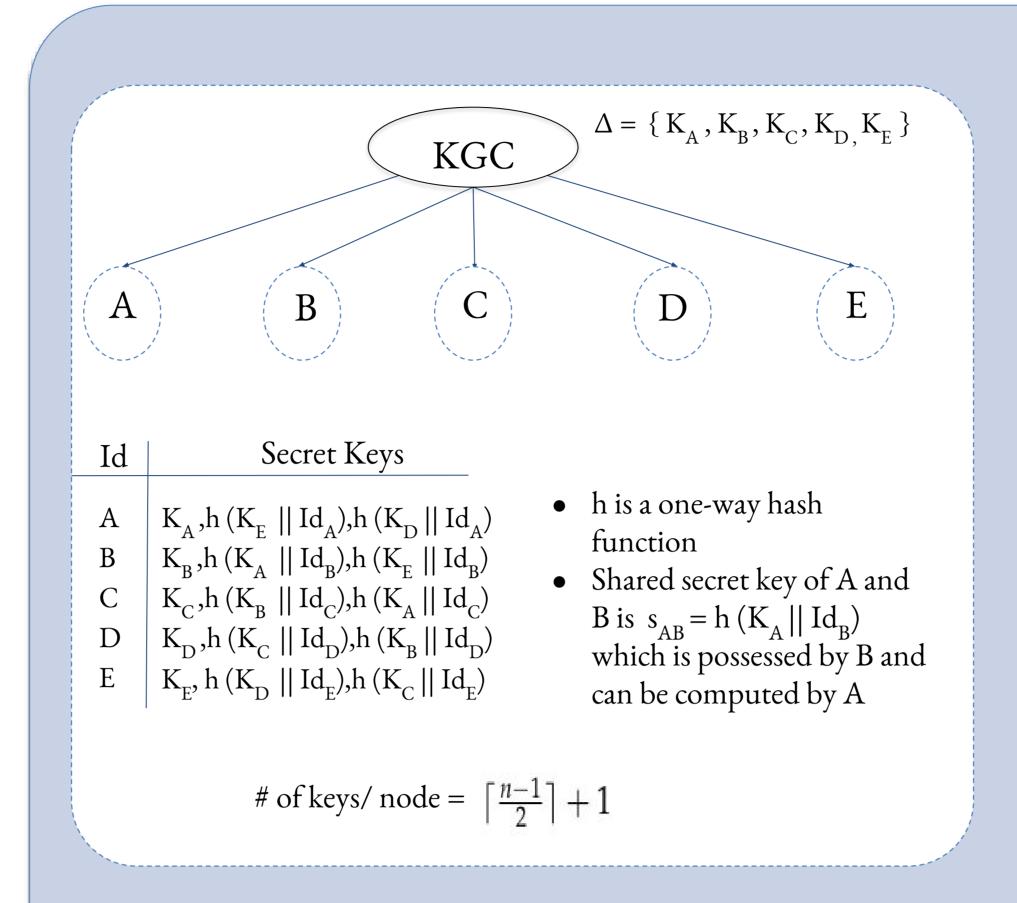
=  $e(H(ID_A), s. H(ID_B))$ 

 $K_{AB} = e(s. H(ID_A), H(ID_B))$ 

 $A \mid s.H.(Id_{\Delta})$ 

 $B \mid s.H.(Id_{R})$ 

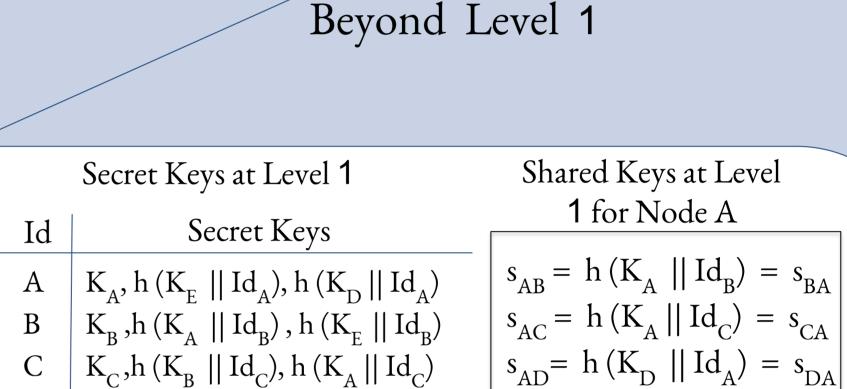
 $C \mid s.H.(Id_C)$ 



# Basic Id One-way Function Scheme

 $\left(A_{1}\right)$ 

#### At Level 1



 $s_{A1A2} = e(s_{AA}.H(A_1),H(A_2)) = e(H(A_1),s_{AA}.H(A_2)) = s_{A2A1}$ 

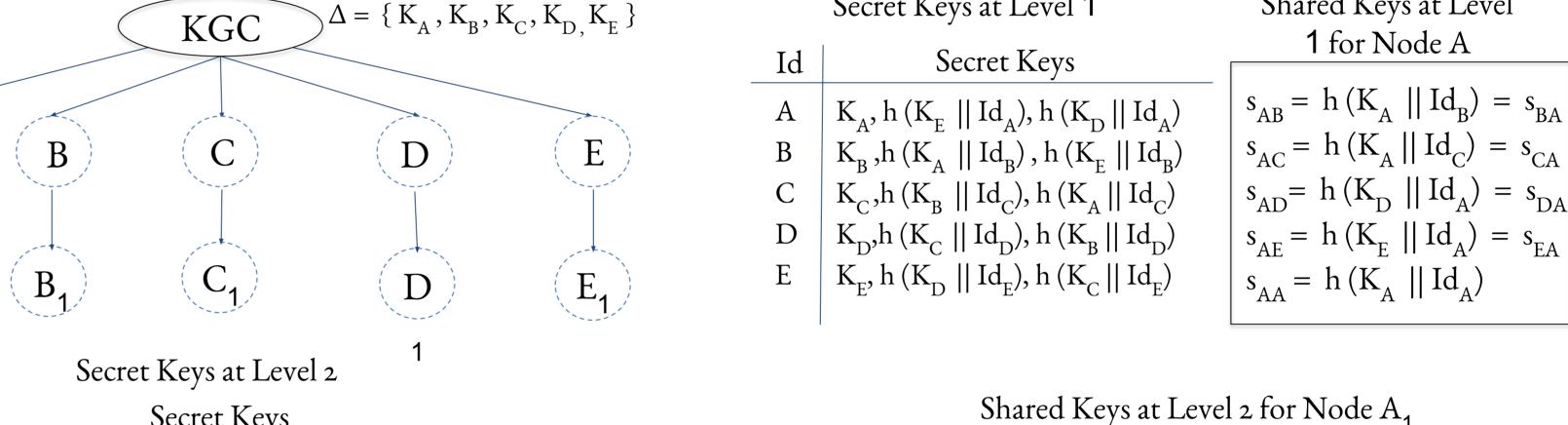
 $s_{A1B1} = e (s_{AB}.H(A_1),H(B_1)) = e(H(A_1),s_{BA}.H(B_1)) = s_{B1A1}$ 

 $s_{A1C1} = e(s_{AC} \cdot H(A_1), H(C_1)) = e(H(A_1), s_{CA} \cdot H(C_1)) = s_{C1A1}$ 

 $s_{A1D1} = e(s_{AD} \cdot H(A_1), H(D_1)) = e(H(A_1), s_{DA} \cdot H(D_1)) = s_{D1A1}$ 

 $s_{A1E1} = e(s_{AE}.H(A_1),H(E_1)) = e(H(A_1),s_{EA}.H(E_1)) = s_{E1A1}$ 

SOK Key Agreement Scheme



Id	Secret Keys
$A_1$	$s_{AA}$ . $H(A_1)$ , $s_{AB}$ . $H(A_1)$ , $s_{AC}$ . $H(A_1)$ , $s_{AD}$ . $H(A_1)$ , $s_{AE}$ . $H(A_1)$
$A_2$	$s_{AA}$ . $H(A_2)$ , $s_{AB}$ . $H(A_2)$ , $s_{AC}$ . $H(A_2)$ , $s_{AD}$ . $H(A_2)$ , $s_{AE}$ . $H(A_2)$
$B_1$	$s_{BA}$ . $H(B_1)$ , $s_{BB}$ . $H(B_1)$ , $s_{BC}$ . $H(B_1)$ , $s_{BD}$ . $H(B_1)$ , $s_{BE}$ . $H(B_1)$
$C_1$	$s_{CA}$ . $H(C_1)$ , $s_{CB}$ . $H(C_1)$ , $s_{CC}$ . $H(C_1)$ , $s_{CD}$ . $H(C_1)$ , $s_{CE}$ . $H(C_1)$
$D_1$	$s_{DA}$ . $H(D_1), s_{DB}$ . $H(D_1), s_{DC}$ . $H(D_1), s_{DD}$ . $H(D_1), s_{DE}$ . $H(D_1)$
E <sub>1</sub>	$s_{EA}$ . $H(E_1)$ , $s_{EB}$ . $H(E_1)$ , $s_{EC}$ . $H(E_1)$ , $s_{ED}$ . $H(E_1)$ , $s_{EE}$ . $H(E_1)$

# BIOS - SOK Key Agreement Scheme

## Comparison of BIOS-SOK and HH-KAS

Scheme :	Polynomial based HH-KAS		Subset based HH-KAS		BIOS-SOK KAS	
Thresholds:	$t_1 = t_2 = 3$	$t_1 = 7, t_2 = 31$	$t_1 = t_2 = 3$	$t_1 = 7, t_2 = 31$	$n_1 = n_2 = 4$	$n_1 = 8, n_2 = 32$
Key-Size (# of group elements)	Root: 100 Leaves: 16	Root : 19008 Leaves : 256	Root : 28768 Leaves : 1800	Root: 8930800 Leaves: 35000	Root: 4 Level 1: 3 Level 2: 4 Leaves: 4	Root: 8 Level 1 : 6 Level 2: 8 Leaves: 32
Shared Key Computation	1 pairing 16 EC mult's	1 pairing 256 EC mult's	1 pairing 450 EC add's 1800 hashing	1 pairing 11000 EC add's 35000 hashing	3 pairings 6 hashing	3 pairings 6 hashing
Shared Key Computation (in terms of field mult's)	39703 m	407623 m	587125 m	11052275 m	47415 m	47415 m

### **Properties**

- Non-Interactive: Any two nodes can compute a shared secret key without any interaction
- Identity-based: to compute the shared secret key, a node only needs its own secret key and peer's identity
- Hierarchical: intermediate nodes in the hierarchy can derive the secret keys for their children
- Resilient: the scheme is fully resilient against compromise of arbitrary number of nodes at each level
- Efficient: Compared to HH-KAS, BIOS-SOK is better in terms of computation time, space requirement and scalability