Characters, Strings, and Floats

CS 350: Computer Organization & Assembler Language Programming

A. Why?

- We need to represent textual characters in addition to numbers.
- Floating-point numbers provide a way to separate the magnitude of a number from the number of bits of significance it has.

B. Outcomes

At the end of today, you should:

- Know how textual characters are represented.
- Know why we have floating-point numbers and how they are represented.

C. Characters and Strings

- A character is represented by a bit string.
- ASCII is one scheme for characters; it uses 8 bits per character.
- The ASCII representations of
 - The digits 0-9 are hex 30-39 (decimal 48-57).
 - The letters A-Z are hex 41-5A (decimal 65-90)
 - The letters a–z are hex 61–7A (decimal 97–122)
 - The space character is hex 20 (decimal 32)
- Unicode is a newer scheme that extends ASCII to use 16 bits per character.
 - Includes symbols from other languages and from math, logic, and other fields.
 - Strings are sequences of characters (null-terminated or length-specified).
- $\bullet\,$ ASCII C strings are null-terminated: They include an extra 8 bit of zeros.
 - E.g., "A9z" would be hex 41 39 7A 00.
 - The character representation of 00_{16} (equals $0000\ 0000_2$) is '\0'.

D. Non-Whole and Floating-Point Numbers

- In decimal, non-whole numbers are represented using a decimal point and a fractional part that sums up negative powers of 10.
- In binary, non-whole numbers use a binary point and negative powers of 2.
 - E.g., $1.01101_2 = 1 + \frac{1}{4} + \frac{1}{8} + \frac{1}{32} = 1 \cdot \frac{13}{32} = 1.40625_{10}$
- True real numbers can take an infinite number of digits.
 - We approximate them with a finite number of digits.
 - Tradeoff representing numbers with a large magnitude vs numbers with a lot of precision. **Magnitude**: Distance from zero. **Precision**: Number of significant digits. The **range** of a set of numbers is the distance between the most positive and most negative numbers.
- Floating-point numbers use separate bits for magnitude and precision
 - Based on scientific notation. E.g. $1.011_2 \times 2^{-56}$ or $-1.101_2 \times 2^{12}$.
 - Note 1: The bit to the left of the binary point will always be 1.
 - Note 2: The overall sign of the number and the sign of the exponent can be different.

E. IEEE Floating-Point Number Standard

- IEEE = Institute of Electrical & Electronics Engineers.
- Break up floating-point number into a sign bit S, an exponent field E, and a fraction field F. If the sign S is 1, the number is negative; if it is 0, the number is non-negative.
- Floating-point numbers come in various sizes. For a 32-bit IEEE floating-point number, the exponent is 8 bits and the fraction is 23 bits.
- The IEEE representation differs from scientific notation in two ways: (1) The fraction part omits the leading 1. it's understood to be there but it's not written; (2) The exponent field is 127 plus the scientific notation exponent.

- exponent is 139-127=12. To get the actual fraction, we prepend "1." to F and get 1.011010...0. So N represents -1.01101×2^{12} .
- Example: For the other direction, if we were asked for the IEEE representation of -1.01101×2^{12} , we'd use 1 for the negative sign, drop the 1. and use 01101... 0 for the fraction, and add 127 to the exponent 12. Since $12 + 127 = 139 = 1000 \ 1011$ for the exponent; this gives 1 1000 1011 01101 followed by eighteen 0's.
- In general, the case where $1 \le E < 255$ is the standard case above.
 - In this case, S E F represents $(-1)^S \times 1.F \times 2^X$ where X = E 127.
- There are some other cases (these won't be on the tests):
 - The case where E = 0 is used for floating-point zero and for extremely small numbers (very close to zero).
 - If E = 0 and Fraction = 0 we have +0 or -0.
 - If E=0 and $Fraction \neq 0$ we have $N=(-1)^S\times 0.Fraction_2\times 2^{-126}$ (***)
 - The case where E=255 used for $+\infty$, $-\infty$, and NaN (Not a Number). (***)

Other Data

CS 350: Computer Organization & Assembler Language Programming

A. Why?

- We need to represent textual characters in addition to numbers.
- We use floating-point numbers to represent non-whole numbers (numbers not evenly divisible by 1).

B. Outcomes

After this activity, you should

- Be able to describe the difference between characters and their ASCII representations.
- Be able to represent floating-point numbers in binary and using the IEEE representation.

C. Questions

- 1. Let '0', '1', ..., '9' be the ASCII representations of the digits 0 9. (a) What are the values of '0', '1', ..., '9'? Which (if any) of the following are true? (b) '2'+'3' = '5' (c) '2'+3 = '5' (d) 2+'3' = '5' (e) 2+3 = '5'.
- 2. Let N be an integer ≥ 0 and ≤ 9 . What are the integer values of the following: (a) '0'+N (b) '0'+'N' (c) 'A'+N (d) 'a'+N (e) Which of (a) (d) represent printable characters and what are those characters?
- 3. Let X represent one of the capital letters A Z. What are the integer values of: (a) X 'A' (b) X 'A' + 'a'? (c) Which of (a) and (b) represent printable characters and what are those characters?
- 4. What decimal number does 10.011₂ represent?
- 5. (a) What is the binary representation of 6.4375_{10} ? (b) What is its scientific notation representation? (I.e., 1.something \times 2 raised to some power.)
- 6. What are the scientific notation and IEEE 32-bit representations for 10.011₂?

- 7. (a) If we add 1.1111 + 0.1110 using 5 significant digits, does truncation (of nonzero bits) occur before the addition? After the addition? (b) Repeat, for 1.1001 + .11101. (c) Repeat, for 1.1101 + .01101.
- 8. In general, is floating-point addition associative?

Solution

- 1a. '0' = 48, '1' = 49 ..., '9' = 57.
- 1b. $'2'+'3' = 50+51 = 101 \text{ is } \neq 53 = '5'$
- 1c. '2'+3 = 50+3=53 is = '5'
- 1d. 2+'3' = 2+53=55 is = '5'
- 1e. 2+3 = 5 is $\neq 53 = 5$.
- 2. Given N is an integer, $0 \le N \le 9$, we have
- 2a. '0' + N = 48 + N = the ASCII representation of N
- 2b. '0'+'N' = 48+78 = 126 (which happens not to = '~')¹.
- 2c. 'A' + N is the representation of the N+1st capital letter of the alphabet, (If N=0 then 'A' + N= 'A', if N=2, then 'A' + N= 'B', etc.)
- 2d. 'a'+ N is the representation of the N+1st lowercase letter of the alphabet. (If N=0 then 'a'+ N= 'a', if N=2, then 'a'+ N= 'b', etc.)
- 2e. (See answers above.)
- 3a. If X is the nth letter of the uppercase alphabet $(1 \le n \le 26)$ then X = 'A' + n-1, so X 'A' = n-1. (If X = 'A', then X 'A' = 0; if X = 'B', then X 'A' = 1; ...; if X = 'Z', then X 'A' = 25.)
- 3b. (X 'A') + 'a' = n 1 + 'a'. This sum is the lowercase version of the uppercase X.
- 4. $10.011_2 = 2 + .011_2 = 2 + (11_2/100_2) = 2^3/8$. Also, $2 + .011_2 = 2 + .25 + .125 = 2.375$ (which = $2^3/8$, of course).
- 5a. $6.4375_{10} = 6 + 0.25 + .125 + .0625 = 110.0111_2$.
- 5b. $110.0111_2 = 1.100111_2 \times 2^2$
- 6. $10.011_2 = 1.0011_2 \times 2^1$. For the IEEE representation, the sign S = 0 (since the number is ≥ 0), the exponent $E = 127 + 1 = 128 = 1000\ 0000_2$, and the fraction F is 1.0011 less the leading "1." plus enough zeros to make 23 bits, so $F = 0011\ 000000\ 00000\ 00000\ 00000\ 0000$. Combining, we get the 32-bit

 $^{^1}$ You're not expected to know 'N' = 78 or '~' = 126 for the tests.

- representation: 0 1000 0000 0011 000000 00000 00000 000. (Not asked for: This is hex 40180000.)
- 7a. We don't need to drop any bits to align the binary points of 1.1111 and 0.1110 (so no truncation), but we have to convert the result $10.1101 \rightarrow 10.110$ because we only have 5 significant digits, and we do truncate the rightmost 1.
 - 1.1111 +0.1110 -----10.1101
- 7b. Aligning 1.1001 + .11101 causes truncation of 1 in the right-hand value; we get 1.1000 + 0.1110; addition yields 10.0110, which turns into 10.011 (no truncation of nonzeros)
- 7c. Aligning 1.1101 + .01101 turns it into 1.1101 + .0110 (truncates a 1); addition yields 10.00111; we truncate 11 to get 10.001.
- 8. No, in general floating-point addition (and arithmetic in general) isn't associative.

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