

# CSE 34 HW#1 Solutions

1- a) Let  $f(n) = 5n^2 - 7n + 9$  and  
 $g(n) = n^2 - 4n + 1$

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To show  $f(n) = O(g(n))$  we should find constants  $c, n_0 > 0$  such that

$$f(n) \leq c g(n) \text{ for all } n \geq n_0.$$

for  $c=6$  we have

$$5n^2 - 7n + 9 \leq 6(n^2 - 4n + 1)$$

$$0 \leq n^2 - 17n - 3$$

roots of  $n^2 - 17n - 3$  is  $-0.15$  and  $17.15$

that is

	$-0.15$	$17.15$
	+	-

so ~~that~~  $c=6$  and  $n_0=18$

satisfy the definition.

1-b) Let  $2n^3 - 6n^2 + n - 3 = f(n)$

To show  $f(n) = \Theta(n^3)$

we should find constants  $c_1, c_2, n_0 > 0$

such that  $0 \leq c_1 n^3 \leq f(n) \leq c_2 n^3$

for  $c_1 = 1$  and  $c_2 = 2$  we have

$$n^3 \leq 2n^3 - 6n^2 + n - 3 \leq 2n^3$$



$$0 \leq n^3 - 6n^2 + n - 3$$

In this part we will use induction to prove correctness.

for  $n=6$  the above inequality is satisfied.

$$0 \leq 6n^2 - n + 3$$

$$\Delta = 1 - 4 \cdot 6 \cdot 3 < 0$$

so there is no root and because  $a=6$  in  $ax^2 + bx + c$  is positive the above inequality is positive for all  $n$ .

So we will now show that ~~for any k~~ for any  $k$  the above inequality is true for  $k+1$  whenever the equality is true for  $k > 6$ .

assume  $0 \leq k^3 - 6k^2 + k - 3$  is true for  $k > 6$ .

we will investigate if  $0 \leq (k+1)^3 - 6(k+1)^2 + (k+1) - 3$

$\Rightarrow$  next page

1-6) continue. ...

when we subtract the first inequality from the first one we get

$$3k^2 - 3k + 1 \quad \text{whose roots are}$$

$$\Delta = 9 - 4 \cdot 3 \cdot 1 = -3 \quad \text{which is always positive.}$$

so we get if

$$0 \leq k^3 - 6k^2 + k - 3 \quad \text{is true for } k \geq 6$$

it is true for ~~all~~  $k+1$ ,  $(k+1)$  also.

$$\text{so } H(n) = O(n^3) \quad c_1 = 1 \quad c_2 = 2$$

$n_0 = 6$  satisfies the definition.

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1-c) Let's first prove that  $(\log n)^3$  is  $O(n)$ . We will prove this by using limits. We will use L'Hopital 3 times.

$$\lim_{n \rightarrow \infty} \frac{n}{(\log n)^3} = \lim_{n \rightarrow \infty} \frac{n}{3(\log n)^2} = \lim_{n \rightarrow \infty} \frac{n}{6 \log n} = \lim_{n \rightarrow \infty} \frac{n}{6} = \frac{n}{6} = \infty$$

so  $(\log n)^3$  is  $O(n)$ .

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$$7n^2 + 5n(\log n)^3 \leq 12n^2 \quad \text{because}$$

$$(\log n)^3 < n \quad \text{for all } n > 4$$

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so for  $c=12$  and  $n_0=2$

$$7n^2 + 5n(\log n)^3 \text{ is } O(n^2).$$

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$$2-a) \sum_{i=0}^{n-2} \sum_{j=i+1}^{n-1} 1 = \sum_{i=0}^{n-2} n-i+1 = 1+2+3+\dots+n-1$$

$$= \frac{n^2 - n}{2}$$

$$f(n) = \frac{n^2 - n}{2}$$

$$2-b) \sum_{i=0}^N 1 + \sum_{i=0}^{N/2} 1 + \sum_{i=0}^{N/4} 1 + \dots + \sum_{i=0}^1 1$$

$$= N + N/2 + N/4 + \dots + 1 = 2N - 1$$

$$f(N) = 2N - 1$$

3-a) Solution can be found on the internet.

3-b) To show that  $(n+a)^b = \Theta(n^b)$

we should find constants  $c_1, c_2, n_0 > 0$  such that

$$0 \leq c_1 n^b \leq (n+a)^b \leq c_2 n^b \text{ for all } n \geq n_0.$$

Note that:

$$n+a \leq n+|a| \leq 2n \text{ when } |a| \leq n.$$

and

$$n+a \geq n-|a| \geq \frac{1}{2}n \text{ when } |a| \leq \frac{1}{2}n$$

Thus, when  $n \geq 2|a|$ ,

$$0 \leq \frac{1}{2}n \leq n+a \leq 2n.$$

Since  $b > 0$ , the inequality still holds when all parts are raised to the power  $b$ .

$$0 \leq \left(\frac{1}{2}n\right)^b \leq (n+a)^b \leq (2n)^b$$

$$0 \leq \left(\frac{1}{2}\right)^b n^b \leq (n+a)^b \leq 2^b n^b$$

Thus  $c_1 = (1/2)^b$ ,  $c_2 = 2^b$  and  $n_0 = 2|a|$

satisfy the definition.

4-)

$$a) \lim_{n \rightarrow \infty} \frac{(n^2 + 1)^{10}}{n^{20}} = \lim_{n \rightarrow \infty} \left(1 + \frac{1}{n^2}\right)^{10} = 1$$

$$\text{so } (n^2 + 1)^{10} = \Theta(n^{20})$$

$$b) \lim_{n \rightarrow \infty} \frac{\sqrt{10n^2 + 7n + 3}}{n} = \sqrt{10}$$

$$\text{so } \sqrt{10n^2 + 7n + 3} = \Theta(n)$$

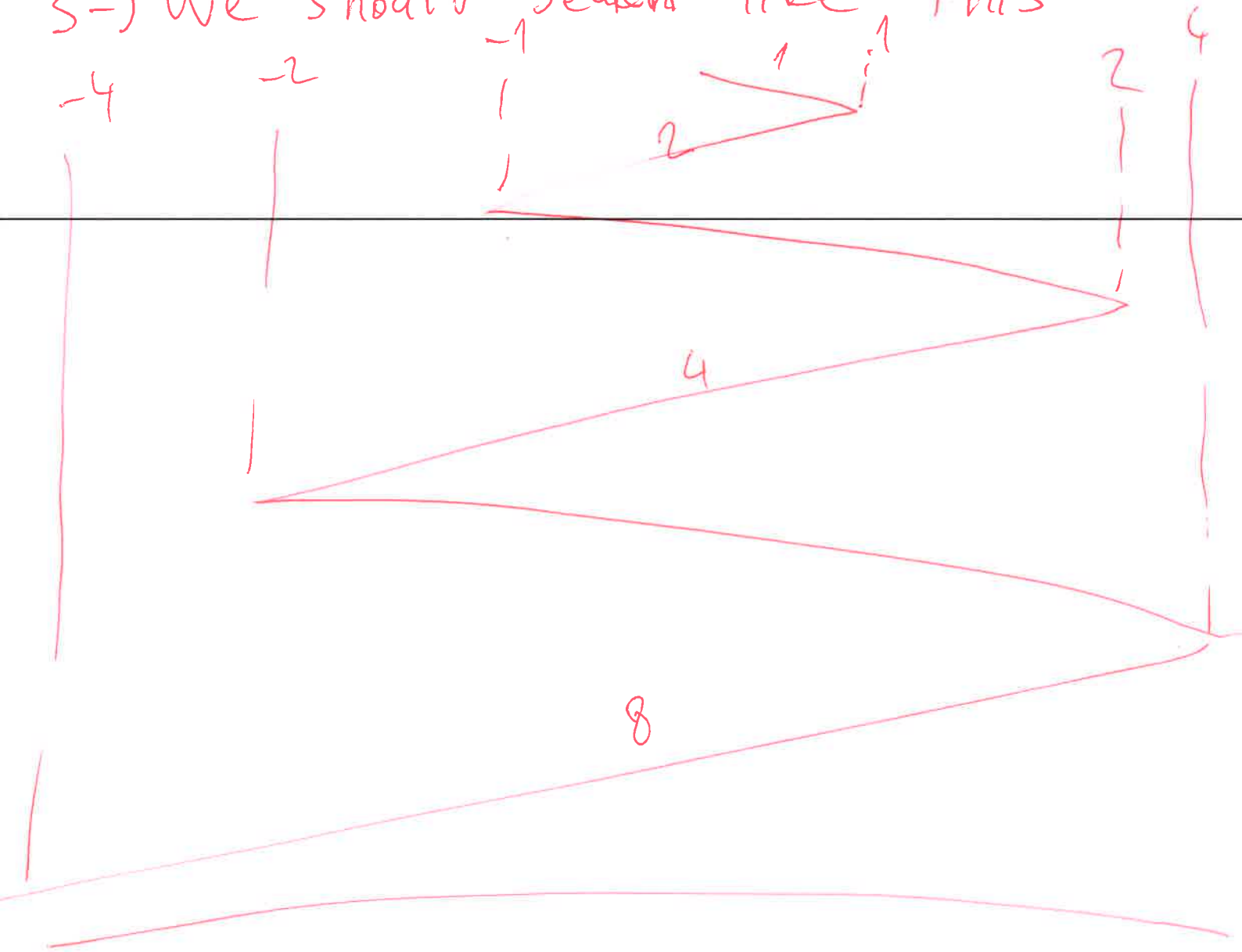
$$c) 2n^2 \lg(n+2)^2 + (n+2)^2 \lg \frac{n}{2} = 2n^2 \lg(n+2)^2 + (n+2)^2 (\lg n - 1) \\ \in \Theta(n \lg n) + \Theta(n^2 \lg n) = \Theta(n^2 \lg n)$$

$$d) 2^{n+1} + 3^{n-1} = 2^{\frac{1}{2}} 2^n + 3^n \cdot \frac{1}{3} \in \Theta(2^n) + \Theta(3^n) = \Theta(3^n)$$

$$e) \log_2 n - 1 \leq \lfloor \log_2 n \rfloor \leq \log_2 n$$

$$\frac{1}{2} \log_2 n \leq \log_2 n - 1 \leq \log_2 n \\ \log_2 n - 1 \in \Theta(\log_2 n) \quad \text{so } \lfloor \log_2 n \rfloor \in \Theta(\log_2 n) \\ \log_2 n \in \Theta(\log_2 n)$$

5-) We should search like this



So total steps ~~is~~ is less than

$$2(2 + 4 + 8 + \dots + n/2 + n)$$

which is equal to  $4 \cdot (1+2+\dots+n/2)$

$$= 4 \cdot (n-1) = 4n - 4$$

total steps is less than

$4n - 4$  which is  $O(n)$ .



6-)

$$a) \sum_{i=0}^{n-1} (i^2+1)^2 = \sum_{i=0}^{n-1} (i^4 + 2i^2 + 1) = \sum_{i=0}^{n-1} i^4 + 2 \sum_{i=0}^{n-1} i^2 + \sum_{i=0}^{n-1} 1$$

$$\in \Theta(n^5) + \Theta(n^3) + \Theta(n) = \Theta(n^5)$$

$$b) \sum_{i=2}^{n-1} \log_2 i^2 = \sum_{i=2}^{n-1} 2 \log_2 i = 2 \sum_{i=2}^{n-1} \log_2 i = 2 \sum_{i=1}^n \log_2 i - 2 \log_2 n$$

$$\left( 2 \sum_{i=1}^n \log_2 i = 2 (\log 1 + \log 2 + \log 3 + \dots + \log n) < 2 \cdot n \cdot \log n, \right. \\ \left. \frac{1}{2} n \log n < 2 \sum_{i=1}^n \log_2 i \right)$$

$$\text{so } 2 \sum_{i=1}^n \log_2 i \in \Theta(n \log n) \text{ and } 2 \log_2 n \in \Theta(\log n)$$

$$\text{hence } \sum_{i=2}^{n-1} \log_2 i^2 \in \Theta(n \log n)$$

$$c) \sum_{i=1}^n (i+1) \cdot 2^{i-1} = \sum_{i=1}^n i \cdot 2^{i-1} + \sum_{i=1}^n 2^{i-1} = \frac{1}{2} \sum_{i=1}^n i \cdot 2^i + \sum_{j=0}^{n-1} 2^j$$

$$\in \Theta(n 2^n) + \Theta(2^n) = \Theta(n 2^n)$$

$$d) \sum_{i=0}^{n-1} \sum_{j=0}^{i-1} (i+j) = \sum_{i=0}^{n-1} \left( \sum_{j=0}^{i-1} i + \sum_{j=0}^{i-1} j \right) = \sum_{i=0}^{n-1} \left( i^2 + \frac{(i-1)i}{2} \right)$$

$$= \sum_{i=0}^{n-1} \left( \frac{3}{2} i^2 - \frac{1}{2} i \right) = \frac{3}{2} \sum_{i=0}^{n-1} i^2 - \frac{1}{2} \sum_{i=0}^{n-1} i \in \Theta(n^3) - \Theta(n^2) = \Theta(n^3)$$