Last update: February 2, 2010

LOCAL SEARCH ALGORITHMS

CMSC 421: Chapter 4, Sections 3-4

Iterative improvement algorithms

In many optimization problems, the **path** to a goal is irrelevant; the goal state itself is the solution

Then state space = a set of goal states find one that satisfies constraints (e.g., no two classes at same time) or, find **optimal** one (e.g., highest possible value, least possible cost)

In such cases, can use *iterative improvement* algorithms; keep a single "current" state, try to improve it

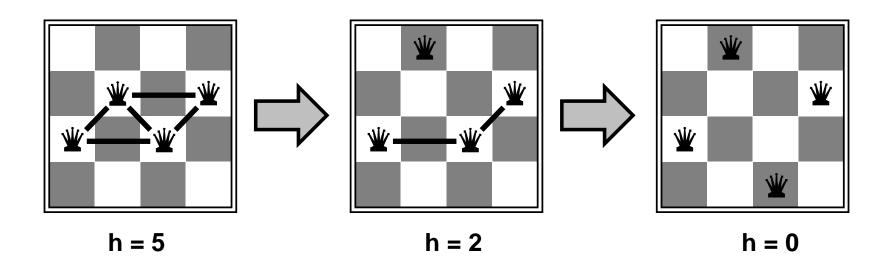
- ♦ Constant space
- ♦ Suitable for online as well as offline search

Example: the n-Queens Problem

- \diamondsuit Put n queens on an $n \times n$ chessboard
- \diamondsuit No two queens on the same row, column, or diagonal

Iterative improvement:

Start with one queen in each column move a queen to reduce number of conflicts



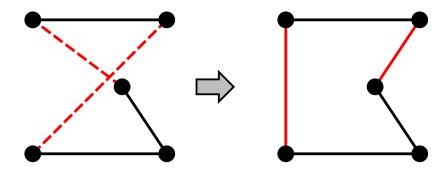
Even for very large n (e.g., n = 1 million), this usually finds a solution almost instantly

Example: Traveling Salesperson Problem

- ♦ Given a complete graph (edges between all pairs of nodes)
- ♦ A tour is a cycle that visits every node exactly once
- \Diamond Find a least-cost *tour* (simple cycle that visits each city exactly once)

Iterative improvement:

Start with any tour, perform pairwise exchanges



Variants of this approach get within 1% of optimal very quickly with thousands of cities

Outline

- ♦ Hill-climbing
- ♦ Simulated annealing
- ♦ Genetic algorithms (briefly)
- ♦ Local search in continuous spaces (very briefly)

Hill-climbing (or gradient ascent/descent)

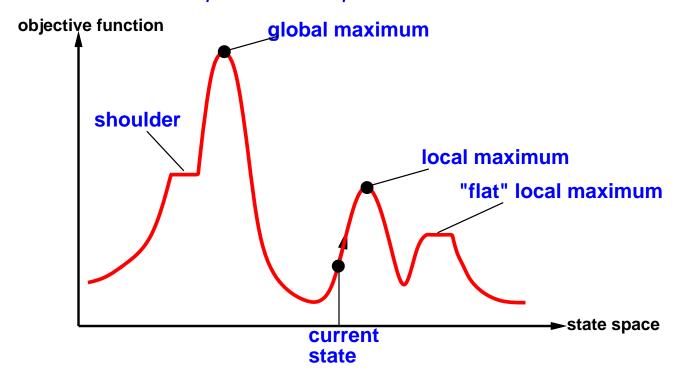
"Like climbing Everest in thick fog with amnesia"

At each step, move to a neighbor of higher value in hopes of getting to a solution having the highest possible value

Can easily modify this for problems where we want to minimize rather than maximize

Hill-climbing, continued

Useful to consider state space landscape



Random-restart hill climbing: repeat with randomly chosen starting points Russell & Norvig say it's trivially complete; they're almost right

If finitely many local maxima, then $\lim_{\text{restarts}\to\infty} P(\text{complete}) = 1$

Simulated annealing

Idea: escape local maxima by allowing some "bad" moves but gradually decrease their size and frequency

```
function SIMULATED-ANNEALING (problem, schedule) returns a solution state
   inputs: problem, a problem
             schedule, a mapping from time to "temperature"
   local variables: current, a node
                        next, a node
                        T, a "temperature" controlling prob. of downward steps
   current \leftarrow Make-Node(Initial-State[problem])
   for i \leftarrow 1 to \infty do
        T \leftarrow schedule[i]
        if T = 0 then return current
        next \leftarrow a randomly selected successor of current
        \Delta E \leftarrow \text{Value}[next] - \text{Value}[current]
        if \Delta E > 0 then current \leftarrow next
        else with probability e^{\Delta E/T}, set current \leftarrow next
```

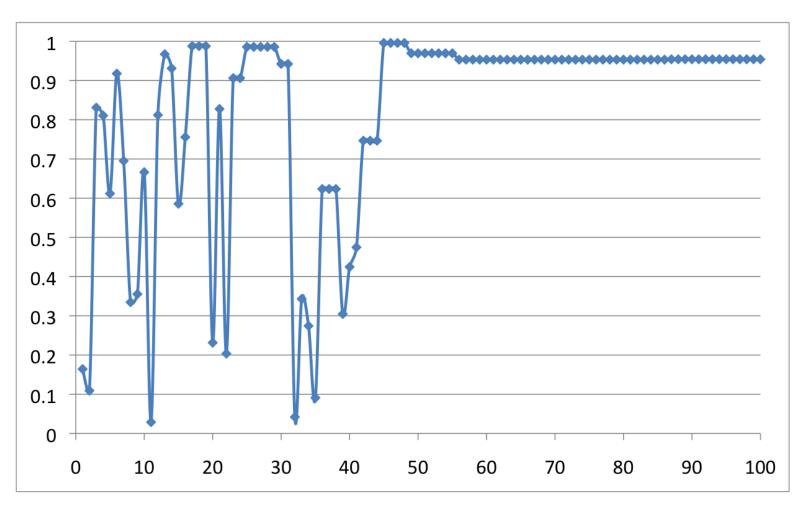
A simple example

Each state is a number $x \in [0,1]$, initial state is 0, all states are neighbors, $Value(x) = x^2$, 100 iterations, $schedule[i] = 10 \times 0.9^i$

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```

Simple example, continued

100 iterations, each state is a number $x \in [0,1]$, initial state is x=0, $Value(x)=x^2$, all states are neighbors, $schedule[i]=10\times 0.9^i$



Properties of simulated annealing

At fixed "temperature" T, probability of being in any given state x reaches Boltzman distribution

$$p(x) = \alpha e^{\frac{E(x)}{kT}}$$

for every state x other than x^* and for small T,

$$p(x^*)/p(x) = e^{\frac{E(x^*)}{kT}}/e^{\frac{E(x)}{kT}} = e^{\frac{E(x^*)-E(x)}{kT}} \gg 1$$

From this it can be shown that if we decrease T slowly enough, $\Pr[\operatorname{reach}\ x^*]$ approaches 1

Devised by Metropolis et al., 1953, for physical process modelling Widely used in VLSI layout, airline scheduling, etc.

Local beam search

```
function BEAM-SEARCH(problem, k) returns a solution state start with k randomly generated states loop generate all successors of all k states if any of them is a solution then return it else select the k best successors
```

Not the same as k parallel searches Searches that find good states will recruit other searches to join them

Problem: often all k states end up on same local hill

Stochastic beam search:

choose k successors randomly, biased towards good ones

Close analogy to natural selection

Genetic algorithms

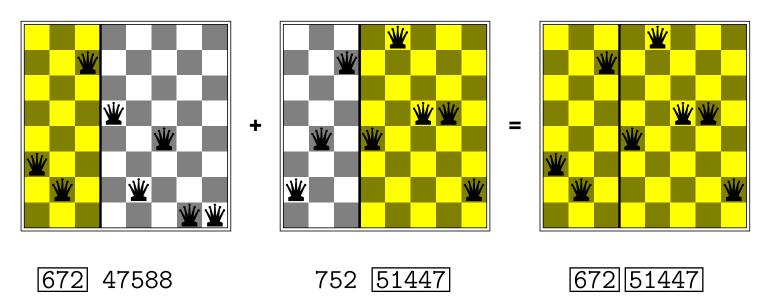
Genetic algorithms

= stochastic local beam search + generate successors from **pairs** of states

Each state should be a string of characters; Substrings should be meaningful components

Example: *n*-queens problem

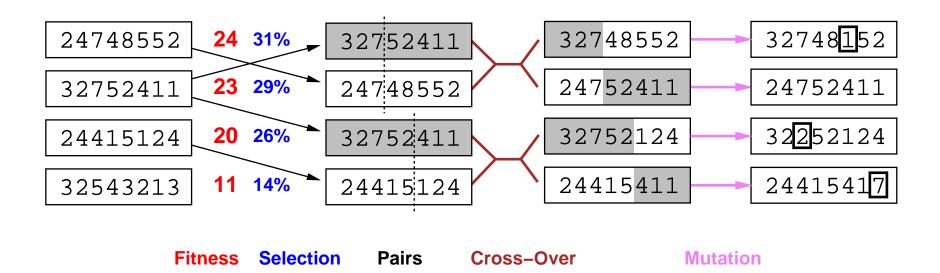
i'th character = row where i'th queen is located



Genetic algorithms

Genetic algorithms

= stochastic local beam search + generate successors from **pairs** of states

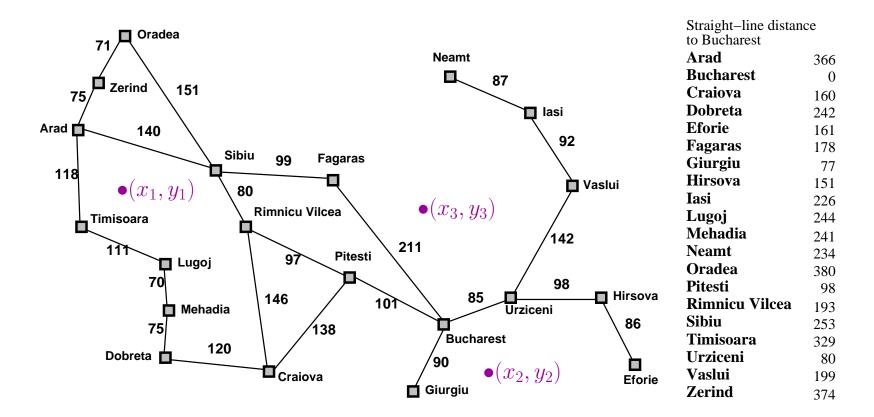


Genetic algorithms \neq biological evolution for example, real genes encode replication machinery

Hill-climbing in continuous state spaces

Suppose we want to put three airports in Romania – what locations?

- \diamondsuit 6-D state space defined by (x_1,y_2) , (x_2,y_2) , (x_3,y_3)
- \diamondsuit Objective function $f(x_1,y_2,x_2,y_2,x_3,y_3)$ measures desirability, e.g., sum of squared distances from each city to nearest airport



Hill-climbing in continuous state spaces

A technique from numerical analysis:

Given a surface z = f(x, y), and a point (x, y), a **gradient** is a vector

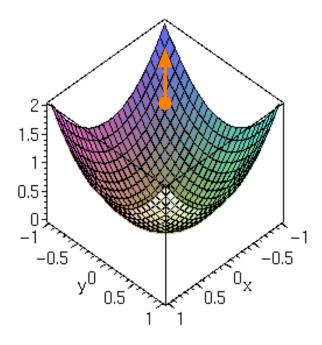
$$\nabla f(x,y) = \left(\frac{\partial f}{\partial x}, \frac{\partial f}{\partial y}\right)$$

The vector points in the direction of the steepest slope, and its length is proportional to the slope.

Gradient methods compute ∇f and use it to increase/reduce f,

e.g., by
$$\mathbf{x} \leftarrow \mathbf{x} - \alpha \nabla f(\mathbf{x})$$

If $\nabla f = 0$ then you've reached a local maximum/minimum

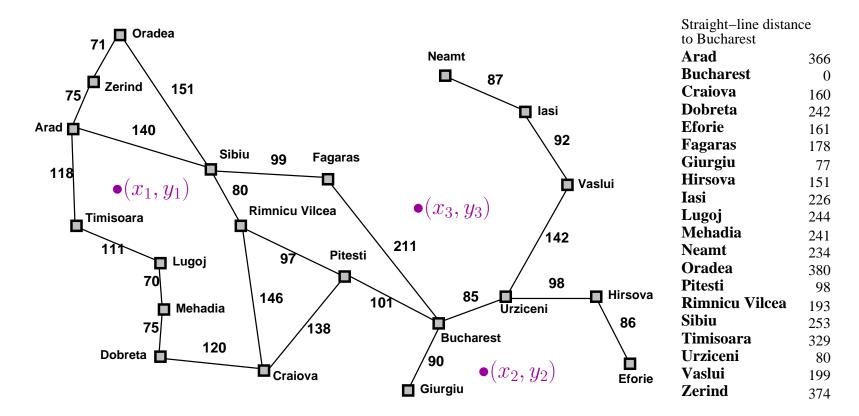


Hill-climbing in continuous state spaces

Suppose we want to put three airports in Romania – what locations?

$$\nabla f = \left(\frac{\partial f}{\partial x_1}, \frac{\partial f}{\partial y_1}, \frac{\partial f}{\partial x_2}, \frac{\partial f}{\partial y_2}, \frac{\partial f}{\partial x_3}, \frac{\partial f}{\partial y_3}\right)$$

Look for $x_1, y_1, x_2, y_2, x_3, y_3$ such that $\nabla f(x_1, y_1, x_2, y_2, x_3, y_3) = 0$



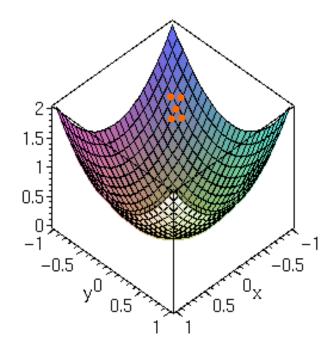
Continuous state spaces, continued

Sometimes can solve for $\nabla f(\mathbf{x}) = 0$ exactly (e.g., with one city)

Newton-Raphson (1664, 1690) iterates $\mathbf{x} \leftarrow \mathbf{x} - \mathbf{H}_f^{-1}(\mathbf{x}) \nabla f(\mathbf{x})$ to solve $\nabla f(\mathbf{x}) = 0$, where $\mathbf{H}_{ij} = \partial^2 f / \partial x_i \partial x_j$

Discretization methods turn continuous space into discrete space

e.g., *empirical gradient* considers $\pm \delta$ change in each coordinate



Homework

```
Problems 4.1,
4.2,
4.9 (but you don't need to suggest a way to calculate it)
4.11,
4.12
```

10 points each, 50 points total

Due in one week