AFFINE PAVING OF PARTIAL FLAG QUIVER VARIETY

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ABSTRACT. In this article, we establish an affine paving for partial flag quiver varieties when the quiver is of Dynkin type. By copying results in [3, section 6] word by word, the same problem for affine quiver reduced to the case where the representation is regular quasi-simple. The idea of the proof mainly comes from [3], and the result is a natural continuation of [4].

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1. Introduction

Let Q be a quiver of Dynkin or affine type(without loops), $X \in \text{Rep}(Q)$ be an quiver representation.¹ We are interested in three objects related to $X \in \text{Rep}(Q)$:

quiver Grassmannian
$$\operatorname{Gr}^{KQ}(X) := \{M_1 \mid M_1 \subseteq X\}$$

partial flag variety $d \geqslant 1$ $\operatorname{Flag_d}(X) := \{0 \subseteq M_1 \subseteq \cdots M_d \subseteq X\}$
strict partial flag variety $d \geqslant 2$ $\operatorname{Flag_{d,str}}(X) := \{0 \subseteq M_1 \subseteq \cdots M_d \subseteq X \mid x.M_{i+1} \subseteq M_i\}^2$

It's easy to see that $\operatorname{Flag}_1(X) = \operatorname{Gr}^{KQ}(X)$. These geometrical objects can be divided into different pieces according to the dimension vectors of M_1, \ldots, M_d , and each piece have its own natural (complex/Zarisky) topology. It was proved in [3] that $\operatorname{Gr}^{KQ}(X)$ have an

¹We fix the base field $K = \mathbb{C}$ for convinience.

²for any $x \in Q_1, i \in \{2, ..., d\}$.

	$\operatorname{Gr}^{KQ}(X)$	$\operatorname{Flag_d}(X)$	$\operatorname{Flag}_{\operatorname{d,str}}(X)$	
$A \over D$	[3, Section 5]	[4, Theorem 2.20]	Corollary 4.2	
E		Section 4		
$ ilde{ ilde{A}} ilde{ ilde{D}}$	[3, Section 6]	Section 5		
\tilde{E}		reduced to the regular quasi-finite case.		

Table 1. Until now, except the \tilde{E} case we've proved the affine paving properties for these varieties.

affine paving, and in [4] that $\operatorname{Flag}_{\operatorname{d}}(X)$ have the same property when Q is Dynkin quiver of type A/E. Here we go one step further, the results are concluded in the Table 1.

The idea of proof is very simple: first, we view the partial flag quiver variety as the quiver Grassmannian of the more complicated quiver; then we establish the decomposition so that one may solve the problem by induction; finally we set a special way of decomposition for each indecomposable module so that we can avoid meeting the bad decomposition. These contents are in Section 2,3,4, accordingly.

Before the end of this section, let us see in one example how partial flag variety is viewed as quiver Grassmannian.

Example 1.1. Let $Q: x \longrightarrow y \longleftarrow z \longrightarrow w$ be a quiver, and let $X: X_x \longrightarrow X_y \longleftarrow X_z \longrightarrow X_w$ be a quiver representation, then the varieties $\operatorname{Flag}_3(X), \operatorname{Flag}_{3,str}(X)$ can be viewed as quiver Grassmannian in Figure 1:

Figure 1

Acknowledgement

First, I would like to thank my supervisor, Professor Jens Niklas Eberhardt, for introducing me this specific problem of partial flag quiver variety and giving some advises. I would also like to thank Hans Franzen for answering some questions I had for understanding the article [3].

2. Preliminary Facts

In this section, we will collect all the knowledges required in the later sections, and fix the notations.

2.1. **Extended quiver.** We follow [4, 2.2,2.3] in this subsection, but we will have some small variations and different notations. We want to view partial flag variety as the quiver Grassmannian. Intuitively, the partial flag variety contains more information than the quiver Grassmannian. So we need to use bigger quiver, and encode these informations in the extra arrows.

Definition 2.1 (extended quiver). Fix the quiver Q and the integer $d \ge 1$, the **extended** quiver Q_d is defined as follows:

• The vertex set of Q_d is defined as the Cartesian product of the vertex set of Q and $\{1, \ldots, d\}$, i.e.

$$v(Q_d) = v(Q) \times \{1, \dots, d\}$$

• We have two types of arrows: for each $(i,r) \in v(Q) \times \{1,\ldots,d-1\}$, there is one arrow from (i,r) to (i,r+1); for every arrow $i \longrightarrow j$ in quiver Q and $r \in \{1,\ldots,d\}$, there is one arrow from (i,r) to (j,r).

The extended quiver Q_d is exactly the same quiver as $\hat{\Gamma}_d$ in [4, Definition 2.2]. The next definition is a small variation of it:

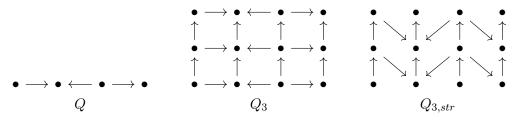
Definition 2.2 (strict extended quiver). Fix the quiver Q and the integer $d \ge 2$, the strict extended quiver $Q_{d,str}$ is defined as follows:

• The vertex set of Q_d is defined as the Cartesian product of the vertex set of Q and $\{1, \ldots, d\}$, i.e.

$$v(Q_{d,str}) = v(Q) \times \{1, \dots, d\}$$

• We have two types of arrows: for each $(i,r) \in v(Q) \times \{1,\ldots,d-1\}$, there is one arrow from (i,r) to (i,r+1); for every arrow $i \longrightarrow j$ in quiver Q and $r \in \{2,\ldots,d\}$, there is one arrow from (i,r) to (j,r-1).

Example 2.3. Here is the picture of new quiver:



Now we define the special bound quiver algebras for later use.

Definition 2.4 (algebra of an extended quiver). For an extended quiver Q_d , let KQ_d be the corresponding path algebra, and I be the ideal of KQ_d identifying all the paths with same sources and targets. The algebra of the extended quiver Q_d is defined as

$$R_d := KQ_d/I$$
.

We also have the "strict" version.

Definition 2.5 (algebra of a strict extended quiver). For an extended quiver $Q_{d,str}$, let $KQ_{d,str}$ be the corresponding path algebra, and I be the ideal of $KQ_{d,str}$ identifying all the paths with same sources and targets. The algebra of the strict extended quiver $Q_{d,str}$ is defined as

$$R_{d,str} := KQ_{d,str}/I.$$

By an aesthentically desirable abuse of notation, we abbreviate the notations R_d and $R_{d str}$ as R.

2.2. **canonical functor** Φ **.** We still follow [4, 2.3] in this subsection with a few variations. The functor Φ realizes a partial flag as a quiver subrepresentation.

Definition 2.6 (canonical functor Φ). The canonical functor $\Phi : \operatorname{Rep}(Q) \longrightarrow \operatorname{Mod}(R)$ is defined as follows:

- $(\Phi(X))_{(i,r)} := X_i;$
- $(\Phi(X))_{(i,r)\to(i,r+1)} := \mathrm{Id}_{X_i};$
- $(\Phi(X))_{(i,r)\to(i,-)} := X_{i\to j}$

For the module $T \in \text{Mod}(R)$, we define $Gr(T) := \{T' \subseteq T \text{ as the submodule}\}.$

Proposition 2.7. Fix the representation $X \in \text{Rep}(Q)$, we have the isomorphism

$$\operatorname{Flag}(X) \cong \operatorname{Gr}(\Phi(X)).$$

Proof. The two side give us same amount of informations. Or, you can easily construct the bijection from $\operatorname{Flag}(X)$ to $\operatorname{Gr}(\Phi(X))$.

2.3. Ext-vanishing properties. We would like to show some higher rank Extension group to be 0, which would be a key ingredient in the proof of the next section.

Lemma 2.8. Let $M, N, X, S \in \text{Rep}(Q), V, W, T \in \text{Mod}(R)$.

- (1) gl. dim $R \leq 2$;
- (2) The functor $\Phi : \operatorname{Rep}(Q) \longrightarrow \operatorname{Mod}(R)$ is exact;
- (3) Φ maps projective module to projective module, and maps injective module to injective module:
- (4) $\operatorname{Hom}_{KQ}(M, N) \cong \operatorname{Hom}_{R}(\Phi(M), \Phi(N)), \quad \operatorname{Ext}_{KQ}^{i}(M, N) \cong \operatorname{Ext}_{R}^{i}(\Phi(M), \Phi(N));$
- (5) proj. dim $\Phi(M) \leq 1$, inj. dim $\Phi(M) \leq 1$;
- (6) Suppose $V \subseteq \Phi(X)$, $W \subseteq \Phi(S)$, then $\operatorname{Ext}_R^2(W,T) = 0$, $\operatorname{Ext}_R^2(T,\Phi(X)/V) = 0$.

Proof.

For (1), we just need to check minimal projective resolution of S(i) in Mod(R).

For (2), if we have the short exact sequence in Rep(Q):

$$0 \longrightarrow X \longrightarrow Y \longrightarrow S \longrightarrow 0$$

then for each vector $i \in v(Q)$, we have the short exact sequence

$$0 \longrightarrow X_i \longrightarrow Y_i \longrightarrow S_i \longrightarrow 0$$
,

so the complex

$$0 \longrightarrow \Phi(X) \longrightarrow \Phi(Y) \longrightarrow \Phi(S) \longrightarrow 0$$

is exact.

For (3), we reduced to the case of indecomposable projective modules, and observe that

$$\Phi(P(i)) = P((i,1)), \qquad \Phi(I(i)) = I((i,d)).$$

For (4), the surjection of map

$$\operatorname{Hom}_{KO}(M,N) \cong \operatorname{Hom}_{R}(\Phi(M),\Phi(N))$$

follows by the commutative diagram:

$$M_i \xrightarrow{\phi_{(i,r+1)}} N_i$$

$$Id \uparrow \qquad \qquad Id \uparrow$$

$$M_i \xrightarrow{\phi_{(i,r)}} N_i$$

and the isomorphism

$$\operatorname{Ext}^i_{KQ}(M,N) \cong \operatorname{Ext}^i_R(\Phi(M),\Phi(N))$$

follows by the projective resolution of M.

For (5), Notice that the minimal projective resolution of M is of length 1, and $\Phi(-)$ sends projective resolution of M to projective resolution of $\Phi(M)$ by (3), thus we get proj. dim $\Phi(M) \leq 1$. The injective dimension of $\Phi(M)$ is computed in the similar way.

For (6), the short exact sequence

$$0 \longrightarrow W \longrightarrow \Phi(S) \longrightarrow \Phi(S)/W \longrightarrow 0$$

induces the long exact sequence

$$\cdots \longrightarrow \operatorname{Ext}_R^2(\Phi(S),T) \longrightarrow \operatorname{Ext}_R^2(W,T) \longrightarrow \operatorname{Ext}_R^3(\Phi(S)/W,T) \longrightarrow \cdots$$

By (1) and (5), $\operatorname{Ext}_R^3(\Phi(S)/W,T)$ and $\operatorname{Ext}_R^2(\Phi(S),T)$ are both 0, so $\operatorname{Ext}_R^2(W,T)=0$. Similarly, from the short exact sequence

$$0 \longrightarrow V \longrightarrow \Phi(X) \longrightarrow \Phi(X)/V \longrightarrow 0$$

we get the induced long exact sequence

$$\cdots \longrightarrow \operatorname{Ext}_R^2(T, \Phi(X)) \longrightarrow \operatorname{Ext}_R^2(T, \Phi(X)/V) \longrightarrow \operatorname{Ext}_R^3(T, V) \longrightarrow \cdots$$
 so $\operatorname{Ext}_R^2(T, \Phi(X)/V) = 0$.

2.4. How much do we understand the quiver representation? To understand the category Rep(Q), one should understand indecomposable modules (as well as their relations). This has almost been done in the Auslander-Reiten theory. For example, when the quiver Q is of Dynkin type, then there are only finite indecomposable representations (up to isomorphism) and each indecomposable representation corresponds to the positive root of Dynkin diagram. One can compute the Auslander-Reiten quiver by knitting algorithm and get the structure of indecomposable representations. Moreover, one can directly get Hom space between M and N by looking at nontrivial paths from M to N^3 .

We will use the Auslander-Reiten quiver to find "good monomorphisms" in Section 4,5. For more informations about Auslander-Reiten theory, one can see [2].

³These paths may be linear dependent, so it's not too easy.

3. Main Theorem

The following notations will be used throughout this article:

• f, g are dimension vectors of quiver Q_d or $Q_{d,str}$. $Gr_f(T) := \{T' \subseteq T \mid \underline{\dim} T' = f\}$ is defined as the set of subrepresentations with dimension vector f.

We use

$$[M, N]^{i} := \dim_{K} \operatorname{Ext}_{KQ}^{i}(M, N), \qquad [M, N] := \dim_{K} \operatorname{Hom}_{KQ}(M, N)$$

$$[T, T']^{i} := \dim_{K} \operatorname{Ext}_{R}^{i}(T, T'), \qquad [T, T'] := \dim_{K} \operatorname{Hom}_{R}(T, T')$$

$$\langle T, T' \rangle_{R} := \sum_{i=0}^{\infty} (-1)^{i} [T, T']^{i} = [T, T'] - [T, T']^{1} + [T, T']^{2}$$

$$\langle f, g \rangle_{R} := \sum_{i \in v(R)} f_{i}g_{i} - \sum_{b \in a(R)} f_{s(b)}g_{t(b)} + \sum_{c \in va(R)} f_{s(c)}g_{t(c)}$$

for the shorthand notation, where

 $v(R) := \{ \text{vertices in quiver } KQ_d \text{ or } KQ_{d,str} \}$

 $a(R) := \{ arrows in quiver KQ_d \text{ or } KQ_{d,str} \}$

 $va(R) := \{\text{"virtual arrows" in quiver } KQ_d \text{ or } KQ_{d,str} \}$

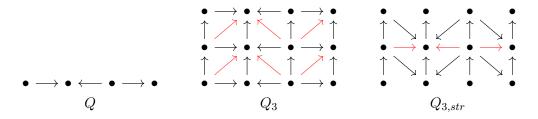


FIGURE 2. virtual arrow: can be thought as the "face" of the quiver

Let $\eta: 0 \longrightarrow X \stackrel{\iota}{\longrightarrow} Y \stackrel{\pi}{\longrightarrow} S \longrightarrow 0$ be a short exact sequence in Rep(Q). Consider the canonical **non-continuous** map

$$\Psi: \operatorname{Gr}(\Phi(Y)) \longrightarrow \operatorname{Gr}(\Phi(X)) \times \operatorname{Gr}(\Phi(S)) \qquad U \longmapsto ([\Phi(\iota)]^{-1}(U), [\Phi(\pi)](U))$$

and $\Psi_{f,g}$ is the map Ψ restricted to the preimage of $\mathrm{Gr}_f(\Phi(X)) \times \mathrm{Gr}_g(\Phi(S))$.

The goal of this section is to prove the following theorems:

Theorem 3.1. When η splits, Ψ is surjective. Moreover, $\Psi_{f,g}$ is a Zarisky-locally trivial affine bundle of rank $\langle g, \underline{\dim} \Phi(X) - f \rangle_R$.

Theorem 3.2 (follows [3, Theorem 32]). When η does not split and $[S, X]^1 = 1$,

$$\operatorname{Im} \Psi_{f,g} = \left(\operatorname{Gr}_f(\Phi(X)) \times \operatorname{Gr}_g(\Phi(S))\right) \setminus \left(\operatorname{Gr}_f(\Phi(X_S)) \times \operatorname{Gr}_{g-\underline{\operatorname{\mathbf{dim}}}} \Phi(S^X) \left(\Phi(S/S^X)\right)\right)$$

where

$$X_S := \max \left\{ M \subseteq X \mid [S, X/M]^1 = 1 \right\} \subseteq X$$

$$S^X := \max \left\{ M \subseteq S \mid [M, X]^1 = 1 \right\} \subseteq S$$

Moreover, $\Psi_{f,g}$ is a Zarisky-locally trivial affine bundle of rank $\langle g, \underline{\dim} \Phi(X) - f \rangle_R$ over $\operatorname{Im} \Psi_{f,g}$.

We will spend the rest of the section proving these theorems. We decide the image as well as the fiber of Ψ respectively.

Lemma 3.3 (follows [3, Lemma 21]). The element $(V, W) \in Gr(\Phi(X)) \times Gr(\Phi(S))$ lies in the image of Ψ if and only if the canonical map $\operatorname{Ext}^1(\Phi(S), \Phi(X)) \longrightarrow \operatorname{Ext}^1(W, \Phi(X)/V)$ maps η to 0.

Proof. The canonical map is defined as follows:

so $\bar{\eta} = 0$ if and only if the last short exact sequence splits, that means, there exist a submodule $U \subseteq \Phi(Y)$, such that $\pi(U) = W$ and $U \cap \Phi(X) = V$.

Corollary 3.4. Resume the notations of Lemma 3.3 When η splits, then Ψ is surjective.

Lemma 3.5. the canonical map $\operatorname{Ext}^1(\Phi(S), \Phi(X)) \longrightarrow \operatorname{Ext}^1(W, \Phi(X)/V)$ is surjective.

Proof. By using the long exact sequence of Extension groups and the Lemma 2.8(6), the maps

$$\operatorname{Ext}^1(\Phi(S), \Phi(X)) \longrightarrow \operatorname{Ext}^1(W, \Phi(X)) \qquad \operatorname{Ext}^1(W, \Phi(X)) \longrightarrow \operatorname{Ext}^1(W, \Phi(X)/V)$$
 are both surjective. Thus the composition is also surjective.

Corollary 3.6. Let S, X be KQ-modules, and $W \subseteq \Phi(S), V \subseteq \Phi(X)$ be R-modules. When $[S, X]^1 = 1$, then $[W, \Phi(X)/V]^1 = 0$ or 1.

Suppose η generates $\operatorname{Ext}^1(S,X)$, then

$$(V,W)\in\operatorname{Im}\Psi\iff [W,\Phi(X)/V]^1=0.$$

In the case where η generates $\operatorname{Ext}^1(S,X)$, we want to describe $\operatorname{Im}\Psi$ more precisely. For this reason we need to introduce two new R-modules:

$$\widetilde{X_S} := \max \left\{ V \subseteq \Phi(X) \mid [\Phi(S), \Phi(X)/V]^1 = 1 \right\} \subseteq \Phi(X)$$

$$\widetilde{S^X} := \max \left\{ W \subseteq \Phi(S) \mid [W, \Phi(X)]^1 = 1 \right\} \subseteq \Phi(S)$$

 $\widetilde{X_S}$ and $\widetilde{S^X}$ are well-defined because of the following lemma:

Lemma 3.7 (follows [3, Lemma 27]).

(i) Let $V, V' \subset \Phi(X)$ such that $[\Phi(S), \Phi(X)/V]^1 = [\Phi(S), \Phi(X)/V']^1 = 1$. Then $[\Phi(S), \Phi(X)/V + V']^1 = 1$.

(ii) Let
$$W, W' \subset \Phi(S)$$
 such that $[W, \Phi(X)]^1 = [W', \Phi(X)]^1 = 1$. Then $[W \cap W', \Phi(X)]^1 = 1$.

Proof. We only prove (i). (ii) is similar.

From the short exact sequence

$$0 \longrightarrow \Phi(X)/V \cap V' \longrightarrow \Phi(X)/V \oplus \Phi(X)/V' \longrightarrow \Phi(X)/V + V' \longrightarrow 0,$$

we get the long exact sequence

$$\cdots \to \operatorname{Ext}^1\!\!\left(\Phi(S), \frac{\Phi(X)}{V \cap V'}\right) \to \operatorname{Ext}^1\!\!\left(\Phi(S), \frac{\Phi(X)}{V}\right) \bigoplus \operatorname{Ext}^1\!\!\left(\Phi(S), \frac{\Phi(X)}{V'}\right) \to \operatorname{Ext}^1\!\!\left(\Phi(S), \frac{\Phi(X)}{V + V'}\right) \to \cdots$$

By Corollary 3.6, $[\Phi(S), \Phi(X)/V \cap V']^1 \leq 1$, $[\Phi(S), \Phi(X)/V + V']^1 \leq 1$, and this forces that $[\Phi(S), \Phi(X)/V + V']^1 = 1$.

Lemma 3.8 (follows [3, Lemma 31(1)(2)], and the proof is same). Let $f: X \longrightarrow \tau S$ be a non-zero morphism, then $X_S = \ker(f)$; also, $\Phi(f): \Phi(X) \longrightarrow \Phi(\tau S)$ is a non-zero morphism, $\widetilde{X_S} = \ker(\Phi(f))$.

Corollary 3.9.
$$\widetilde{X_S} = \Phi(X_S)$$
. $(since \ \widetilde{X_S} = \ker(\Phi(f)) = \Phi(\ker(f)) = \Phi(X_S))$

By the similar argument, one can show that $\widetilde{S}^X = \Phi(S^X)$.

Lemma 3.10 (follows [3, Lemma 31(6)]). Given $V \subseteq \Phi(X)$ and $W \subseteq \Phi(S)$, we have

$$[W, \Phi(X)/V]^1 = 0 \iff V \nsubseteq \Phi(X_S) \text{ or } W \not\supseteq \Phi(S^X).$$

Proof. \Leftarrow : without lose of generation suppose $V \nsubseteq \Phi(X_S)$, then

$$V \nsubseteq \Phi(X_S) \iff [\Phi(S), \Phi(X)/V]^1 = 0 \Rightarrow [W, \Phi(X)/V]^1 = 0.$$

 \Rightarrow : If not, then $V \subseteq \Phi(X_S)$ and $W \supseteq \Phi(S^X)$, and⁴

$$[W, \Phi(X)/V]^1 \geqslant [\Phi(S^X), \Phi(X)/\Phi(X_S)]^1 = [S^X, X/X_S]^1 = 1.$$

We get the contradiction!

Corollary 3.11. When η generates $\operatorname{Ext}^1(S,X)$, we have

$$\operatorname{Im} \Psi_{f,g} = \left(\operatorname{Gr}_f(\Phi(X)) \times \operatorname{Gr}_g(\Phi(S))\right) \setminus \left(\operatorname{Gr}_f(\Phi(X_S)) \times \operatorname{Gr}_{g-\underline{\operatorname{\mathbf{dim}}} \Phi(S^X)} \left(\Phi(S/S^X)\right)\right)$$

Lemma 3.12. For $(V, W) \in \text{Im } \Psi$, the preimage of (V, W) is a torsor of $\text{Hom}_R(W, \Phi(X)/V)$. Or we could say, there is one non-canonical isomorphism

$$\Psi^{-1}((V,W)) \cong \operatorname{Hom}_R(W,\Phi(X)/V).$$

 $^{{}^{4}[}S^{X}, X/X_{S}]^{1} = 1$ follows from [3, Lemma 31(5)]

Proof. Recall the commutative diagram

$$\eta \in \operatorname{Ext}^{1}(\Phi(S), \Phi(X)) \qquad 0 \longrightarrow \Phi(X) \longrightarrow \Phi(Y) \stackrel{\pi}{\longrightarrow} \Phi(S) \longrightarrow 0$$

$$\downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \uparrow \qquad \uparrow \qquad \uparrow$$

$$\operatorname{Ext}^{1}(W, \Phi(X)) \qquad 0 \longrightarrow \Phi(X) \longrightarrow \pi^{-1}(W) \longrightarrow W \longrightarrow 0$$

$$\downarrow \qquad \qquad \downarrow \qquad \qquad \downarrow \qquad \qquad \downarrow$$

$$\bar{\eta} \in \operatorname{Ext}^{1}(W, \Phi(X)/V) \qquad 0 \longrightarrow \Phi(X)/V \stackrel{\iota}{\longrightarrow} \pi^{-1}(W)/V \stackrel{\pi'}{\longrightarrow} W \longrightarrow 0$$

When $(V, W) \in \text{Im } \Psi$, $\bar{\eta}$ is split, and each split morphism θ give us an element in $\Psi^{-1}((V, W))$. If we fix one split morphism θ_0 , then the other split morphisms are all of the form $\theta_0 + \iota \circ f$ where $f \in \text{Hom}_R(W, \Phi(X)/V)$ (and this form is unique). So

$$\Psi^{-1}((V, W)) \cong \{\theta : \text{ split morphism}\} \cong \operatorname{Hom}_R(W, \Phi(X)/V).$$

Lemma 3.13. For two R-modules T, T', we have

$$\langle T, T' \rangle_R \cong \langle \underline{\mathbf{dim}} T, \underline{\mathbf{dim}} T' \rangle_R$$

Proof. We construct one canonical resolution of T,

$$0 \longrightarrow_{\substack{c \in Q_{2} \\ c \in D_{1} \\ c \in D_{2}}} Re_{t(c)} \otimes_{k} e_{s(c)} T \longrightarrow_{\substack{b \in Q_{1} \\ b \in Q_{2}}} Re_{t(b)} \otimes_{k} e_{s(b)} T \longrightarrow_{\substack{c \in Q_{2} \\ c \in D_{3} \\ c \in D_{2}}} Re_{i} \otimes_{k} e_{i} T \longrightarrow_{\substack{c \in Q_{2} \\ c \in D_{3} \\ c \in D_{3}}} r \otimes x \longrightarrow_{\substack{r \in C_{1} \otimes x \\ -r \in C_{1} \otimes x - r \otimes b_{2} x}} r \otimes x \longrightarrow_{\substack{r \in C_{1} \otimes x \\ -r \in C_{2} \otimes x - r \otimes b_{3} x}} r \otimes x \longrightarrow_{\substack{r \in C_{1} \otimes x \\ -r \in C_{2} \otimes x - r \otimes b_{3} x}} r \otimes x \longrightarrow_{\substack{r \in C_{1} \otimes x \\ -r \in C_{2} \otimes x - r \otimes b_{3} x}} r \otimes x \longrightarrow_{\substack{r \in C_{1} \otimes x \\ -r \in C_{2} \otimes x - r \otimes b_{3} x}} r \otimes_{\substack{r \in C_{1} \otimes x \\ -r \in C_{2} \otimes x - r \otimes b_{3} x}} r \otimes_{\substack{r \in C_{1} \otimes x \\ -r \in C_{2} \otimes x - r \otimes b_{3} x}} r \otimes_{\substack{r \in C_{1} \otimes x \\ -r \in C_{2} \otimes x - r \otimes b_{3} x}} r \otimes_{\substack{r \in C_{1} \otimes x \\ -r \otimes x - r \otimes b_{3} x}} r \otimes_{\substack{r \in C_{1} \otimes x \\ -r \otimes x - r \otimes b_{3} x}} r \otimes_{\substack{r \in C_{1} \otimes x \\ -r \otimes x - r \otimes b_{3} x}} r \otimes_{\substack{r \in C_{1} \otimes x \\ -r \otimes x - r \otimes b_{3} x}} r \otimes_{\substack{r \in C_{1} \otimes x \\ -r \otimes x - r \otimes b_{3} x}} r \otimes_{\substack{r \in C_{1} \otimes x \\ -r \otimes x - r \otimes b_{3} x}} r \otimes_{\substack{r \in C_{1} \otimes x \\ -r \otimes x - r \otimes b_{3} x}} r \otimes_{\substack{r \in C_{1} \otimes x \\ -r \otimes x - r \otimes b_{3} x}} r \otimes_{\substack{r \in C_{1} \otimes x \\ -r \otimes x - r \otimes b_{3} x}} r \otimes_{\substack{r \in C_{1} \otimes x \\ -r \otimes x - r \otimes b_{3} x}} r \otimes_{\substack{r \in C_{1} \otimes x \\ -r \otimes x - r \otimes b_{3} x}} r \otimes_{\substack{r \in C_{1} \otimes x \\ -r \otimes x - r \otimes b_{3} x}} r \otimes_{\substack{r \in C_{1} \otimes x \\ -r \otimes x - r \otimes b_{3} x}} r \otimes_{\substack{r \in C_{1} \otimes x \\ -r \otimes x - r \otimes b_{3} x}} r \otimes_{\substack{r \in C_{1} \otimes x \\ -r \otimes x - r \otimes b_{3} x}} r \otimes_{\substack{r \in C_{1} \otimes x \\ -r \otimes x - r \otimes x - r \otimes b_{3} x}} r \otimes_{\substack{r \in C_{1} \otimes x \\ -r \otimes x - r \otimes x - r \otimes x}} r \otimes_{\substack{r \in C_{1} \otimes x \\ -r \otimes x - r \otimes x}} r \otimes_{\substack{r \in C_{1} \otimes x \\ -r \otimes x - r \otimes x}} r \otimes_{\substack{r \in C_{1} \otimes x \\ -r \otimes x - r \otimes x}} r \otimes_{\substack{r \in C_{1} \otimes x \\ -r \otimes x - r \otimes x}} r \otimes_{\substack{r \in C_{1} \otimes x \\ -r \otimes x - r \otimes x}} r \otimes_{\substack{r \in C_{1} \otimes x \\ -r \otimes x - r \otimes x}} r \otimes_{\substack{r \in C_{1} \otimes x \\ -r \otimes x - r \otimes x}} r \otimes_{\substack{r \in C_{1} \otimes x \\ -r \otimes x - r \otimes x}} r \otimes_{\substack{r \in C_{1} \otimes x \\ -r \otimes x - r \otimes x}} r \otimes_{\substack{r \in C_{1} \otimes x \\ -r \otimes x - r \otimes x}} r \otimes_{\substack{r \in C_{1} \otimes x \\ -r \otimes x - r \otimes x}} r \otimes_{\substack{r \in C_{1} \otimes x \\ -r \otimes x - r \otimes x}} r \otimes_{\substack{r \in C_{1} \otimes x \\ -r \otimes x - r \otimes x}} r \otimes_{\substack{r \in C_{1} \otimes x \\ -r \otimes x - r \otimes x}} r \otimes_{\substack{r \in$$

by applying the functor $\operatorname{Hom}_R(-,T')$ to this resolution, we get

$$\langle T, T' \rangle_R \cong \langle \underline{\dim} T, \underline{\dim} T' \rangle_R$$

Proof of Theorem 3.1 and 3.2. We have already computed Im Ψ in Corollary 3.4 and 3.11. For the rank of the affine bundle, we have

$$(V, W) \in \operatorname{Im} \Psi_{f,g} \Longrightarrow [W, \Phi(X)/V]^1 = 0$$

 $\Longrightarrow [W, \Phi(X)/V] = \langle W, \Phi(X)/V \rangle_R = \langle f, \underline{\dim} \Phi(X) - g \rangle_R$

4. Application: Dynkin Case

Before discussing the affine paving property, let me introduce some new numerical concepts, which can be seen as the measure of the "complexity" of the representation.

Fix an **indecomposable** quiver representation $M \in \text{Rep}(\mathbb{Q})$, we define the order of M by

$$\operatorname{ord}(M) := \max_{i \in v(Q)} \dim_K M_i.$$

When the quiver Q is of type E, we denote by $e \in v(Q)$ the unique vertex which is connected to three other vertices, and the number

$$\operatorname{ord}_e(M) := \dim_K M_e = [P(e), M]$$

is equal to $\operatorname{ord}(M)$ unless $\operatorname{ord}_e(M) = 0$.

By Theorem 3.1, we just need to focus on the case of indecomposable modules. The next lemma tells us, for the "easy representation", we can prove the affine paving property easily.

Lemma 4.1 (follows [4, Lemma 2.22]). For the representation $M \in \text{Rep}(Q)$ satisfying $\text{ord}(M) \leq 2$ and the dimension vector f, the variety $\text{Gr}_f(\Phi(M))$ is either empty or is a singleton or is a direct product of some copies of \mathbb{P}^1 .

Proof. For each vertex i of the bigger quiver, the dimension of V_i is 0,1,2; then the variety $\operatorname{Gr}_f(\Phi(M))$ is naturally included to a direct product of \mathbb{P}^1 , and the information of arrows just reduce the number of \mathbb{P}^1 (maybe to singleton or empty).

Corollary 4.2 ([4, Theorem 2.20]). Assume that Q is a Dynkin quiver of type A or D, $M \in \text{Rep}(Q)$ is the representation. then the Grassmannian $\text{Gr}(\Phi(M))$ has an affine paving.

Corollary 4.3. Assume that Q is a affine quiver of type A or D, $M \in \text{Rep}(Q)$ is the regular quasi-simple representation. then the Grassmannian $\text{Gr}(\Phi(M))$ has an affine paving.

In the rest of this section we focus on the "difficult" indecomposable representation of E_6, E_7, E_8 . The idea is to design the special route for each case, and use Theorem 3.2 in the process. Notice that even though the Auslander-Reiten quivers look quite different for different quiver (with same type), they can have the same form is we use the number $\operatorname{ord}_e(M)$ to represent the representation M:

Lemma 4.4. For every indecomposable representation Y of type E with $\operatorname{ord}(Y) > 2$, there is a minimal section mono $f: X \longrightarrow Y$.

Proof. Just observe the Auslander-Reiten sequence. The chosen minimal section monos are represented in Figure 4. Notice that for the most time $\operatorname{ord}_e(-)$ is enough to guarantee the map to be a mono.

⁵Some representations M are hidden when $\operatorname{ord}_e(M)=0$. In [1] the Figure 3 is called the starting functions.

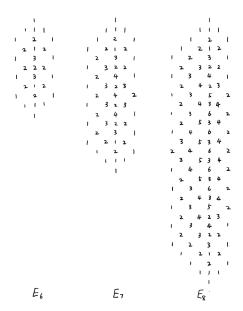


FIGURE 3. central imformation of Auslander-Reiten quiver

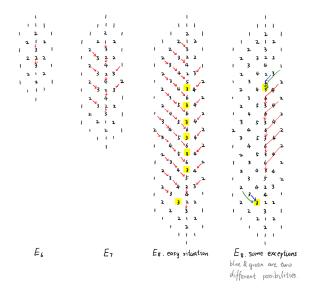


FIGURE 4. minimal section monos

Remark 4.5. The condition ord(Y) > 2 in the lemma can not be removed.

Lemma 4.6. Let $X \hookrightarrow Y$ be an minimal section mono, and S := Y/X be the quotient. Then we have the short exact sequence

$$\eta: 0 \longrightarrow X \longrightarrow Y \longrightarrow S \longrightarrow 0$$

and the following values:

[W'N], [W'N]/N	X	Y	2
X	- 0	-0	00
Y	00	- 0	1 0
S	0 -	00	I 0

In particular, S is indecomposable and rigid; $[S, X]^1 = 1$, so X_S and S^X are well-defined.

Proof. We know that [X,X]=[Y,Y]=1 and $[X,X]^1=[Y,Y]^1=0$. By the definition of minimal section mono, we get [X,Y]=1,[Y,X]=0 and $[X,Y]^1=[Y,X]^1=0$. By applying the functors [Y,-],[-,S],[X,-],[-,X],[-Y] to the short exact sequence η we get the results.

In the following two lemmas we will describe the representations S^X and X_S more clearly.

Lemma 4.7. Take the same notations as in Lemma 4.6. Then $S^X = S$.

Proof. Let $\iota: N \longrightarrow S$ be a proper non-zero subrepresentation of S, we need to prove that $\iota^*\eta: 0 \longrightarrow X \longrightarrow Y' \longrightarrow N \longrightarrow 0$ splits.

$$\iota^*\eta: \qquad \qquad 0 \longrightarrow X \longleftarrow Y' \longrightarrow N \longrightarrow 0$$

$$\downarrow \qquad \qquad \downarrow^{\eta} \qquad \downarrow^{\iota}$$

$$\eta: \qquad \qquad 0 \longrightarrow X \longrightarrow Y \longrightarrow S \longrightarrow 0$$

We decompose $Y' = \bigoplus_i Y_i'$ as the direct sum of indecomposable representations. Since the map $X \longrightarrow Y$ is the minimal section mono, we get $Y_i' = X$ or $Y_i' = Y$ or $X \stackrel{0}{\longrightarrow} Y_i'$ for all i. If there exists i such that $Y_i' = X$, then ι^* splits; if there exists i such that $Y_i' = Y$, then η is isomorphism, we get ι is isomorphism; if for every i the map $X \longrightarrow Y_i'$ is 0, then the map $X \longrightarrow Y'$ is 0, we also get the contradiction.

Lemma 4.8 (follows [3, Lemma 36], proof is exactly the same). Let $E \longrightarrow X$ be the minimal right almost split morphism ending in X, then we can decompose E as $E = E' \oplus \tau X_1$. When Y is not projective, X_S is isomorphic to $\ker(E \longrightarrow \tau Y) \cong E' \oplus \ker(\tau X_1 \longrightarrow \tau Y)$; when Y is projective, $X_S \cong E$.

Corollary 4.9. When $X \longrightarrow Y$ is irreducible monomorphism, the representation X_S is either 0 or an indecomposable representation with property that $X_S \longrightarrow X$ is also an irreducible monomorphism.

Remark 4.10. We can not copy everything in [3, Lemma 56], sometimes it would happen that $X_S = F \oplus T$ with F and T indecomposable, $F \hookrightarrow X$ is irreducible but $T \longrightarrow X/F$ is not a good mono.

$$\bullet \longrightarrow \bullet \longrightarrow \bullet \longrightarrow \bullet \longleftarrow \bullet \longleftarrow \bullet$$

For example, take the quiver of type E_7 : $\bullet \longrightarrow \bullet \longrightarrow \bullet \longrightarrow \bullet \longleftarrow \bullet \longleftarrow \bullet$ take $Y = \frac{1}{12321}$, $X = \frac{1}{112321}$, then $X_S = \frac{1}{111210} \oplus \frac{0}{000111} = F \oplus T$, $X/F = \frac{0}{001111}$, the map $T \longrightarrow X/F$ is not a good mono.

Luckily, we can avoid this bad situation by carefully choosing the minimal section mono $X \longrightarrow Y$. The minimal section monos I chose are presented in Figure 4. In appendix we will write down the induction process in detail for some examples.

5. Application: Affine Case

For the affine case, we just need to follow [3, Section 6], and change everything from Gr(-) to $Gr(\Phi(-))$. There is no difference except the Proposition 48, in which the authors proved the affine paving properties of quasi-simple regular representations. So we reduced the question to the case of quasi-simple regular representation. Combined with Corollary 4.3, we've proved the affine paving properties for A, D cases.

For an regular quasi-simple representation Y of type E, it's possible that there's no short exact sequence

$$\eta: 0 \longrightarrow X \longrightarrow Y \longrightarrow S \longrightarrow 0$$

such that $[S,X]^1 \leq 1$. Then we can no longer use Theorem 3.1 or 3.2. We leave this small tail for interested readers.

Appendix A

In this appendix we solve every case in Figure 4.

When the minimal section mono $X \longrightarrow Y$ is irreducible, we use Theorem 3.2 to get morphism

$$Gr(\Phi(Y)) \longrightarrow Gr(\Phi(X)) \times Gr(\Phi(S))$$
 or $Gr(\Phi(X)) \setminus Gr(\Phi(X_S))$

By observation of Figure 4, $\operatorname{ord}_e(S) = \operatorname{ord}_e(Y) - \operatorname{ord}_e(X)$ is smaller or equal to 2, so by Lemma 4.1 ord_e(S) has the affine paving property. Let $Y_1 := X$, $X_1 := X_S$, $S_1 := Y_1/X_1$, we again use Theorem 3.2 to get morphism

$$\operatorname{Gr}(\Phi(X)) \longrightarrow \operatorname{Gr}(\Phi(X_1)) \times \operatorname{Gr}(\Phi(S_1)) \text{ or } \operatorname{Gr}(\Phi(X_1)) \setminus \operatorname{Gr}(\Phi(X_{1S_1}))$$

 $\operatorname{Gr}(\Phi(X)) \setminus \operatorname{Gr}(\Phi(X_S)) \longrightarrow \operatorname{Gr}(\Phi(X_1)) \times \operatorname{Gr}(\Phi(S_1))$

Luckily $\operatorname{ord}_{e}(S_{1})$ is still smaller or equal to 2. We can continue this process until the order of representations are small enough.

In the exception cases the game is similar, but we need to discuss a little more complicated. Let us look at some examples. (We simplify the notations: Gr(M) as $Gr_f(\Phi(M))$, U(M,N) as $\operatorname{Gr}_f(\Phi(M)) \setminus \operatorname{Gr}_f(\Phi(M))$, and we also ignore the dimension vectors.)

Example A.1. In the case of Figure 5(a), if $X_1 \longrightarrow Y$ is injective, then

$$\operatorname{Gr}(Y) \longrightarrow \operatorname{Gr}(X_1) \times \operatorname{Gr}(Y/X_1) \text{ or } U(X_1, X)$$

$$\operatorname{Gr}(X_1) \longrightarrow \operatorname{Gr}(X) \times \operatorname{Gr}(X_1/X) \text{ or } U(X, X_S)$$

$$U(X_1, X) \longrightarrow \operatorname{Gr}(X) \times \operatorname{Gr}(X_1/X)$$

$$U(X, X_S) \longrightarrow \operatorname{Gr}(X_S) \times \operatorname{Gr}(X/X_S).$$

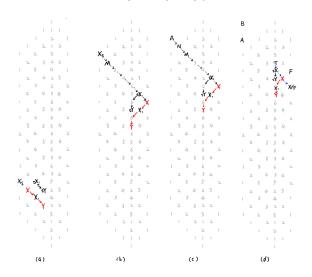


Figure 5. special cases

When $X_1 \longrightarrow Y$ is not injective, we get

$$Gr(Y) \longrightarrow Gr(X) \times Gr(Y/X)$$
 or $U(X, X_S)$.

Since the map $\tau X_1 \longrightarrow \tau Y$ is injective, from Lemma 4.8 we get $X_S \longrightarrow X$ is irreducible monomorphism. Thus

$$U(X, X_S) \longrightarrow Gr(X_S) \times Gr(X/X_S).$$

These maps give the variety Gr(Y) an affine paving from bottom to top.

Example A.2. In Figure 5(b), we would like to prove that Gr(Y) has the affine paving property. We have

$$Gr(Y) \longrightarrow Gr(X) \times Gr(Y/X)$$
 or $U(X, X_S)$.

When the map $M \longrightarrow X$ is not monomorphism, we get

$$U(X, X_S) \longrightarrow Gr(X_S) \times Gr(X/X_S);$$

when the map $M \longrightarrow X$ is monomorphism, we get

$$U(X, X_S) = U(X, M) \bigsqcup U(M, X_S)$$
$$U(X, M) \longrightarrow Gr(M) \times Gr(X/M)$$
$$U(M, X_S) \longrightarrow Gr(X_S) \times Gr(M/X_S).$$

Since the order of X, Y/X, X_S , X/X_S , M, X/M, M/X_S are small or equal to 2, the induction process stops, we get Gr(Y) has the affine paving property.

Example A.3. In the case of Figure 5(c), we have

$$Gr(Y) \longrightarrow Gr(X) \times Gr(Y/X)$$
 or $U(X, X_S)$

where $X_S = \ker(\tau X_1 \longrightarrow \tau Y)$. When $X_S = 0$ we're done; if not, then $A \neq 0$ and $X_S = A$, we decompose $X_S \longrightarrow Y$ as compositions of minimal section monos:

Case 1: $M \longrightarrow X$ is not injective, then

$$U(X, X_S) = U(X, N) \bigsqcup U(N, X_S)$$
$$U(X, N) \longrightarrow Gr(N) \times Gr(X/N)$$
$$U(N, X_S) \longrightarrow Gr(X_S) \times Gr(N/X_S).$$

Case 2: $M \longrightarrow X$ is injective, then

$$U(X, X_S) = U(X, M) \bigsqcup U(M, N) \bigsqcup U(N, X_S)$$

$$U(X, M) \longrightarrow Gr(M) \times Gr(X/M)$$

$$U(M, N) \longrightarrow Gr(N) \times Gr(M/N)$$

$$U(N, X_S) \longrightarrow Gr(X_S) \times Gr(N/X_S).$$

Since Gr(X), Gr(Y/X), Gr(N), ... have affine paving property, we conclude that Gr(Y) has also the affine paving property.

Example A.4. Finally we begin to handle the most difficult case(Figure 5(d)). When $X \longrightarrow Y$ is not injective, we get

$$Gr(Y) \longrightarrow Gr(F) \times Gr(Y/F)$$
 or $U(F,?)$

then we get the result ⁶.

When $X \longrightarrow Y$ is injective, we have

$$Gr(Y) \longrightarrow Gr(X) \times Gr(Y/X)$$
 or $U(X, X_S)$

where $X_S = F \oplus \ker(\tau X_1 \longrightarrow \tau Y) = F \oplus T$ by Lemma 4.8. Since $X \longrightarrow Y$ is injective, we get A = 0, thus B = 0 also, and then the sectional map $T \longrightarrow X/F$ in injective. We thus get two short exact sequence satisfying the conditions in 3.2:

$$\eta:$$
 $0 \longrightarrow F \longrightarrow X \xrightarrow{\pi} X/F \longrightarrow 0$
 $\xi:$ $0 \longrightarrow T \longrightarrow X/F \xrightarrow{\pi'} X/X_S \longrightarrow 0$

Let $N \in Gr(X)$ be a subrepresentation, it's obvious that $N \in Gr(X_S) \iff \pi' \circ \pi(N) = 0$, so

$$N \in U(X, X_S) \iff \pi' \circ \pi(N) \neq 0$$

 $\iff \pi(N) \notin Gr(T)$
 $\iff \pi(N) \in U(X/F, T)$
 $\iff \Psi_n(N) \in Gr(F) \times U(X/F, T)$

Thus the Zarisky-locally trivial affine bundle map

$$U(X,F) \longrightarrow Gr(F) \times Gr(X/F)$$

restricted to the Zarisky-locally trivial affine bundle map

$$U(X, X_S) \longrightarrow Gr(F) \times U(X/F, T).$$

Finally, by applying the short exact sequence ξ to Theorem 3.2 we get the map

$$U(X/F,T) \longrightarrow Gr(X/F) \times Gr(T)$$
.

 $^{{}^{6}\}mathrm{Gr}(F)$ is empty or a singleton, so is U(F,?).

Since all the Grassmannians Gr(X), Gr(Y/X), Gr(F), Gr(X/F), Gr(T) have the affine paving property, we conclude that Gr(Y) has the affine paving property.

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