

# Affine pavings of partial flag varieties

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# Process

- 1 Setting and Statement
- 2 Case study
- 3 Auslander–Reiten theory
- 4 Sketch of proof

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# Affine paving

## Setting

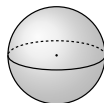
$K = \mathbb{C}$ ,  $X$ : algebraic variety over  $K$ .

## Definition

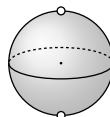
An **affine paving** of  $X$  is a filtration

$$0 = X_0 \subset X_1 \subset \cdots \subset X_d = X$$

with  $X_i$  closed and  $X_{i+1} \setminus X_i \cong \mathbb{A}_K^k$ .



$$\mathbb{P}^1 = \{\infty\} \sqcup \mathbb{A}^1$$



$\mathbb{P}^1 \setminus \{0, \infty\}$  has no affine paving

# Quiver and quiver representation

Quiver is a graph. It has some vertices & arrows.  
In this talk, all the quivers are finite and connected.

# Quiver and quiver representation

We focus on the Dynkin quiver.

That means, the graph of the Dynkin diagrams in the ADE series.

# Partial flag variety

## Definition

Fix a quiver  $Q$  and  $M \in \text{rep}(Q)$ ,

$$\text{Flag}_d(M) := \{F: 0 \subseteq N_1 \subseteq \cdots \subseteq N_d \subseteq M\}$$

$$\text{Flag}_{\underline{\mathbf{f}}}(M) := \{F: 0 \subseteq N_1 \subseteq \cdots \subseteq N_d \subseteq M \mid \underline{\dim} M_i = \underline{\mathbf{f}}_i\}$$

## Example

$$Q = \bullet, M = \mathbb{C}^n, \underline{\mathbf{f}} := \begin{pmatrix} n \\ \vdots \\ 1 \end{pmatrix}$$

$$\text{Flag}_d(\mathbb{C}^n) = \{F: 0 \subseteq N_1 \subseteq \cdots \subseteq N_d \subseteq \mathbb{C}^n\}$$

$$\text{Flag}_1(\mathbb{C}^n) = \{F: 0 \subseteq N_1 \subseteq \mathbb{C}^n\} = \sqcup_{k=0}^n \text{Gr}(n, k)$$

$$\text{Flag}_{\underline{\mathbf{f}}}(\mathbb{C}^n) = \text{complete flags of } \mathbb{C}^n$$

$$\text{Flag}_{(k)}(\mathbb{C}^n) = \text{Gr}(n, k)$$

# Statement

## Theorem

*For a Dynkin quiver  $Q$  and  $M \in \text{rep}(Q)$ ,*

*$\text{Flag}_d(M)$  has an affine paving.*



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Task 1.  $Q = \bullet$ ,  $M = \mathbb{C}^n$ 

In this case,

$$\begin{array}{ccc} \mathrm{GL}_n(\mathbb{C}) \curvearrowright \mathbb{C}^n & \rightsquigarrow & \mathrm{GL}_n(\mathbb{C}) \curvearrowright \mathrm{Flag}_d(\mathbb{C}^n) \\ & \rightsquigarrow & B \curvearrowright \mathrm{Flag}_d(\mathbb{C}^n) \end{array}$$

$\mathrm{Flag}_d(\mathbb{C}^n)$  has an affine paving given by Schubert cells (i.e.,  $B$ -orbits).

### Note

When  $Q = \bullet \longrightarrow \bullet$ ,  $\mathrm{Flag}_{\mathbf{f}}(M)$  have no natural group actions.

# Idea of affine pavings

Find a nice short exact sequence

$$0 \longrightarrow X \xrightarrow{\iota} M \xrightarrow{\pi} S \longrightarrow 0$$

which induces a nice morphism

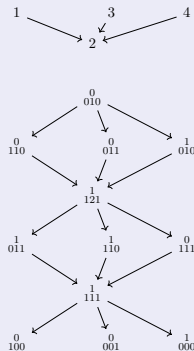
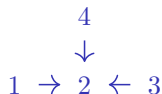
$$\begin{aligned} \Psi : \operatorname{Flag}_d(M) &\longrightarrow \operatorname{Flag}_d(X) \times \operatorname{Flag}_d(S) \\ F &\longmapsto (\iota^{-1}(F), \pi(F)) \end{aligned}$$

We construct the affine paving of  $\operatorname{Flag}_d(M)$  from the affine paving of  $\operatorname{Flag}_d(X)$  and  $\operatorname{Flag}_d(S)$ . Then, we use mathematical induction.

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Another example:  $D_4$



For other examples, see [here](#).

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