

Prelim: March 17 7³⁰ - 9³⁰

- Double sided A4 cheatsheet
- Material up to today

HW3 posted → March 19

► Conditional MI

Recap: $I(X;Y) := D_{KL}(P_{XY} \parallel P_X \otimes P_Y)$

relation b/t MI and entropy

$$\begin{aligned} I(X;Y) &= H(X) + H(Y) - H(X,Y) \\ &= H(X) - H(X|Y) \\ &= H(Y) - H(Y|X) \end{aligned}$$

Let $P_{XYZ} \in \mathcal{P}(X \times Y \times Z)$ and consider $P_{Y|XZ}(\cdot|z)$,
 $z \in \mathbb{Z}$, then $P_{XY|Z}(\cdot|x) \in \mathcal{P}(X \times Y)$.

We define

$$I(X;Y|Z=z) := D_{KL}\left(P_{XY|Z}(\cdot|z) \parallel P_{X|Z}(\cdot|z) \otimes P_{Y|Z}(\cdot|z)\right)$$

Definition (Mutual Information):

For $X, Y, Z \sim P_{XYZ} \in \mathcal{P}(X \times Y \times Z)$, the conditional MI between X and Y given Z is

$$I(X; Y|Z) = D_{KL}(P_{XY|Z} \| P_{X|Z} \otimes P_{Y|Z} | P_Z)$$

$$= \mathbb{E}_{z \sim P_Z} [I(X; Y|Z=z)]$$

Comments:

- ① $I(X; Y|Z)$ is a functional of P_{XYZ} (not of e.g. $P_{XY|Z}$)
- ② It can be shown that

$$\begin{aligned} I(X; Y|Z) &= H(X|Z) + H(Y|Z) - H(XY|Z) \\ &= H(X|Z) - H(X|YZ) \\ &= H(Y|Z) - H(Y|XZ) \end{aligned}$$

Definition (Markov Chain)

Let $X, Y, Z \sim P_{XYZ} \in \mathcal{P}(X \times Y \times Z)$. We say that $X \rightarrow Y \rightarrow Z$ forms a Markov chain if

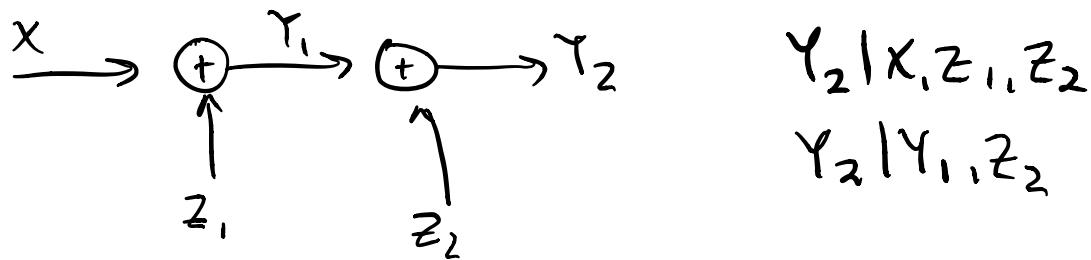
$$P_{XYZ} = P_X P_{Y|X} \underbrace{P_{Z|XY}}_{\text{independent of past}} = P_X P_{Y|X} \underbrace{P_{Z|Y}}_{\text{independent of past}}$$

Example

$$\left. \begin{array}{l} X \sim N(0, \sigma^2) \\ Z_1 \sim N(0, \eta_1^2) \\ Z_2 \sim N(0, \eta_2^2) \end{array} \right\} \text{mutually independent}$$

Define $Y_1 = X + Z_1$,

$$Y_2 = X + Z_1 + Z_2 = Y_1 + Z_2$$



Proposition (Equivalent Conditions)

The following are equivalent

$$X \rightarrow Y \rightarrow Z \Leftrightarrow P_{XYZ} = P_X P_{Y|X} P_{Z|Y}$$

$$\Leftrightarrow P_{XZ|Y} = P_{X|Y} P_{Z|Y}$$

$$\Leftrightarrow P_{Z|XY} = P_{Z|Y}$$

$$\Leftrightarrow X \perp\!\!\!\perp Z | Y \quad (X, Z \text{ conditionally independent } | Y)$$

$$\Leftrightarrow Z \rightarrow Y \rightarrow X$$

...

Proposition (More properties of MI)

Let $(X, Y, Z) \sim P_{XYZ}$. Then

① $I(X; Y|Z) \geq 0$ w/ equality iff $X \rightarrow Z \rightarrow Y$

② Chain Rule

$$\begin{aligned} \text{(i) Small } I(X, Y; Z) &= I(X; Z) + I(Y; Z | X) \\ &= I(Y; Z) + I(X; Z | Y) \end{aligned}$$

$$\text{(ii) Full: } I(X_1, \dots, X_n; Z) = I(X_1; Y) + \sum_{i=2}^n I(X_i; Y | X_1, \dots, X_{i-1})$$

③ Data Processing Inequality

If $X \rightarrow Y \rightarrow Z$, then

$$I(X; Y) \geq I(X; Z)$$

w/ equality iff $X \rightarrow Z \rightarrow Y$.

④ If f is a bijection then

$$I(X; Y) = I(X; f(Y))$$

⑤ Concavity / Convexity

For $(X, Y) \sim P_{XY}$ denote $I(X; Y)$ as $I(P_X; P_{Y|X})$.

Then

(i) For fixed $P_{Y|X}$, $P_X \mapsto I(P_X; P_{Y|X})$ is concave

(ii) For fixed P_X , $P_{Y|X} \mapsto I(P_X; P_{Y|X})$ is convex

Proof of proposition

① By def \rightarrow complete details

② "Trivial"

$X \rightarrow Y \rightarrow Z$ MC

$$③ I(X; Y, Z) = I(X; Y) + I(X; Z|Y) \geq I(X; Y) + 0$$

$$= I(X; Z) + \underbrace{I(X; Y|Z)}_{\text{?/0}}$$

w/ equality

$$\Rightarrow I(X; Z) \leq I(X; Y)$$

condition pops out!

④ For any f deterministic

$$X \rightarrow Y \rightarrow f(Y)$$

$$\xrightarrow{\text{DPI}} I(X; Y) \geq I(X; f(Y))$$

But when f a bijection $X \rightarrow f(Y) \rightarrow Y$

$$I(X; Y) \leq I(X; f(Y))$$

$$\Rightarrow I(X; Y) = I(X; f(Y))$$

IV Letter Typical Sequences

Overview for Binary Alphabets

- Let $\mathcal{X} = \{0, 1\}$ and consider its n -fold extension \mathcal{X}^n . This is the set of all binary sequences of length n . Elements of \mathcal{X}^n are denoted as

$$x^n := (x_1, \dots, x_n) \in \mathcal{X}^n$$

Q1: How many sequences are in \mathcal{X}^n ?

A1: $|\mathcal{X}^n| = |\mathcal{X}|^n = 2^n$

Now let $P \in \mathcal{P}(\mathcal{X})$, i.e. $P = \text{Ber}(\alpha)$, $\alpha \in (0, 1)$

non-inclusive otherwise
not interesting

Let $X \sim P$ and X_1, X_2, X_3, \dots be iid copies of X . In other words: For $n \in \mathbb{N}$, $(X_1, \dots, X_n) \sim P^{\otimes n}$

Any time we draw $x^n := (X_1, \dots, X_n)$ we get a value in \mathcal{X}^n .

Observation For any $x^n \in \mathcal{X}^n$, we have $P^{\otimes n}(x^n) > 0$.

Indeed if x^n has $k \leq n$ 1's, then

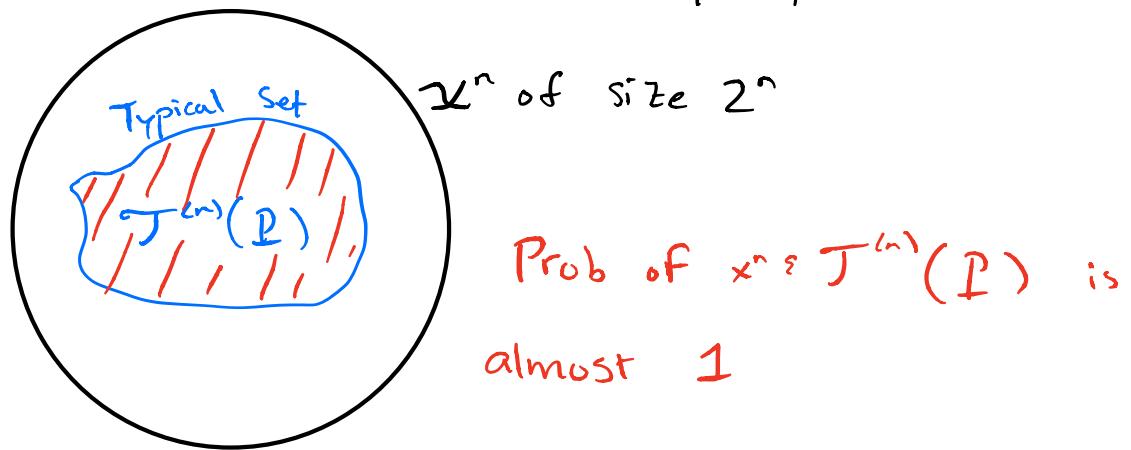
$$P^{\otimes n} = \alpha^k (1-\alpha)^{n-k} > 0 \quad \text{when } \alpha \in (0, 1)$$

Q2: Despite the above, what is the likely portion of 1's that we'll get by drawing x^n ?

A2: Expect to see roughly $n\alpha$ 1's, $n(1-\alpha)$ 0's.

Goal: By exploiting the above observation, we will define for a subset of \mathcal{X}^n ($\mathcal{T}^{(n)}(P)$), such that

$$(i) |\mathcal{T}^{(n)}(P)| \ll |\mathcal{X}^n| \Rightarrow \frac{|\mathcal{T}^{(n)}(P)|}{|\mathcal{X}^n|} \xrightarrow{n \rightarrow \infty} 0$$



(ii) Set "absorbs almost all the probability" in the sense that

$$\lim_{n \rightarrow \infty} P^{\otimes n}(\mathcal{T}^{(n)}(P)) = 1$$

$\xrightarrow{\quad}$

$$P^{\otimes n}(\{x^n \in \mathcal{X}^n \mid x^n \in \mathcal{T}^{(n)}(P)\})$$

To formally set this up we need a few definitions.

Definition (Empirical Frequency):

Let X be discrete and $x^n \in X^n$. Define

$$N_{x^n}(a) := \sum_{i=1}^n \mathbb{1}_{\{x_i = a\}} \quad a \in X$$

deterministic sequence

and then set $\tilde{v}_{x^n}(a) := \frac{1}{n} N_{x^n}(a)$, $a \in X$, and

observe that $\tilde{v}_{x^n}(a)$ is a valid pmf on X .

The pmf $\tilde{v}_{x^n}(a)$ is called the empirical frequency of $x^n \in X^n$.