

THINKING BACK 5 MONTHS



7000 GWEI BASEFEE FLASHBACKS

Enter the L2s



Scaling

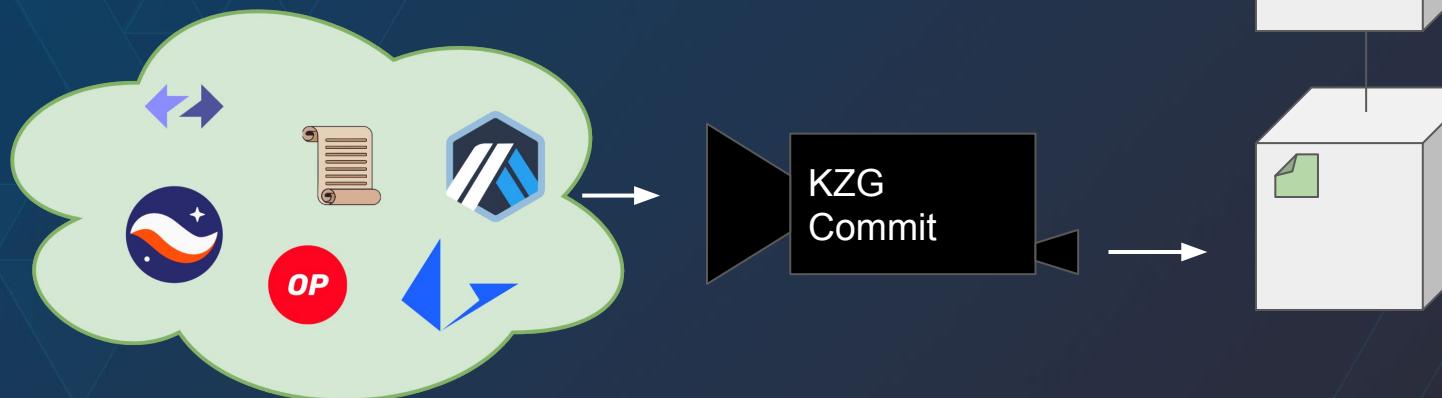
L2s Scale Compute



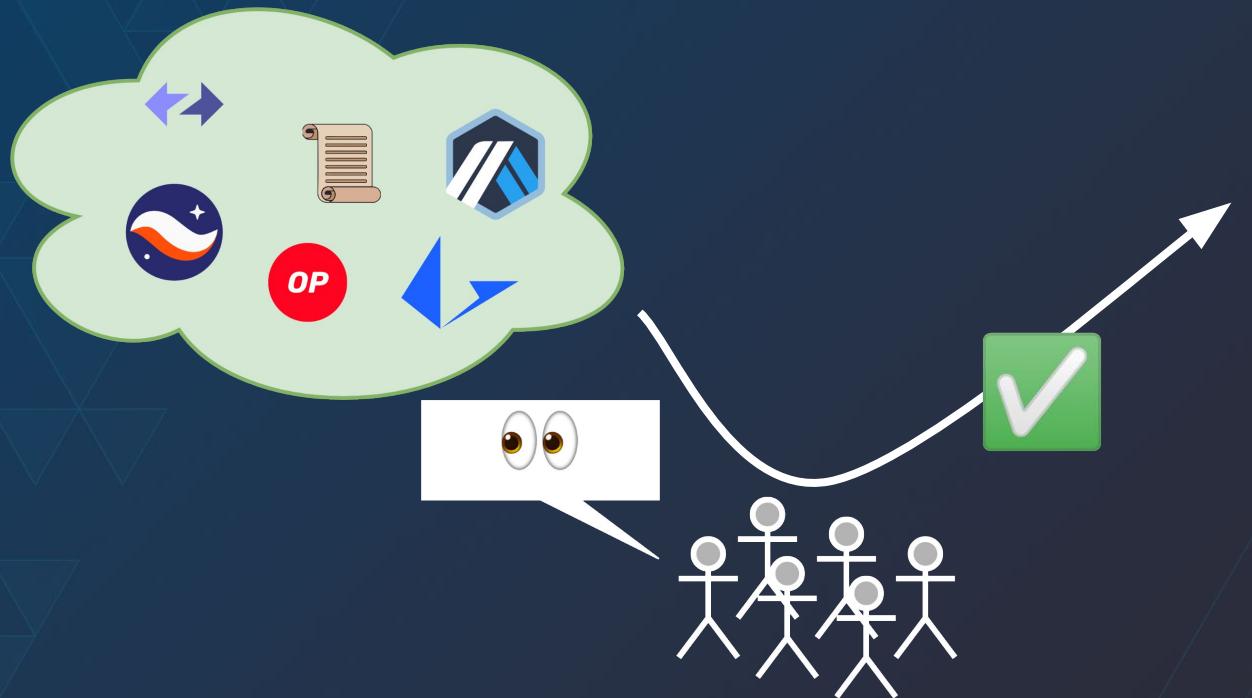
(Proto) Danksharding Scales Data Availability



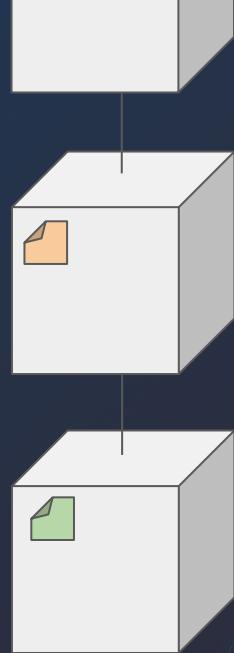
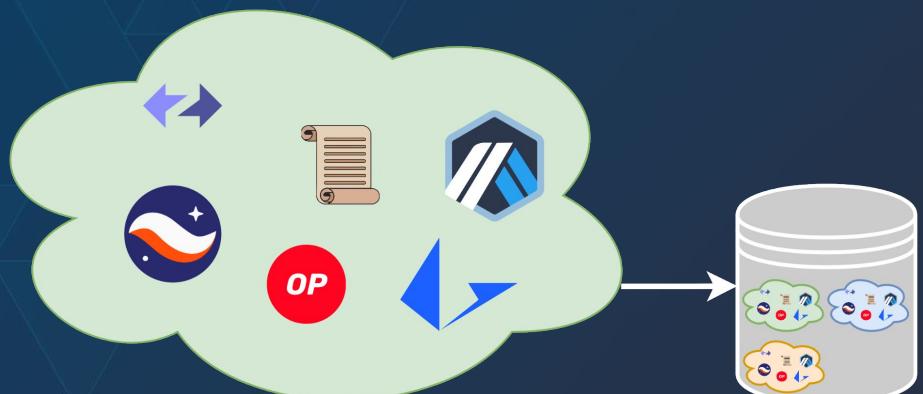
(Proto) Danksharding / EIP-4844



Checking availability

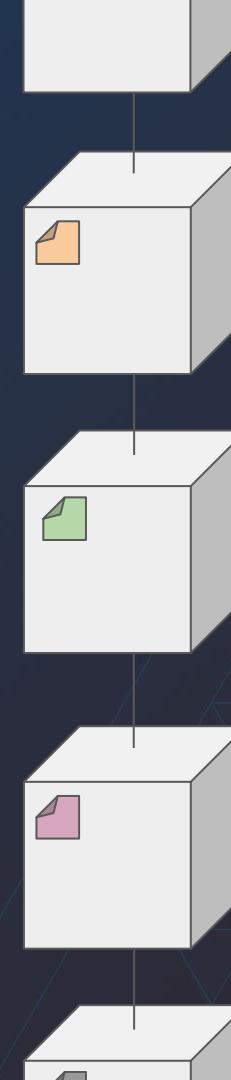
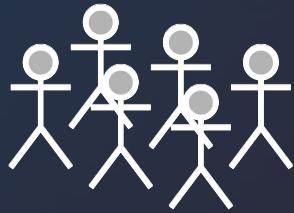


Serving the data



**TWO WEEKS
LATER...**

Throw it all away!



Why KZG?

	Merkle (SHA 256)	Merkle (Arithmetic)	KZG	IPA
Verification inside SNARK	Prohibitive	Doable	Cheap	Cheap
Proof Verification (& low deg. proofs)	Hard	SNARKs needed for easy verify	Easy	$O(N)$ time
Proof generation in $O(n^* \log(n))$	Yes	Yes	Yes	Only with larger proof-size
Commit linearity	No	No	Yes	Yes
Proof linearity	No	No	Yes	No
Security assumption	Very safe	Immature hash functions	Trusted setup + Quantum vuln.	Quantum vuln.

Modified from Vitalik's post:

ethresear.ch/t/13863

Summoning a secret



Only one honest participant required

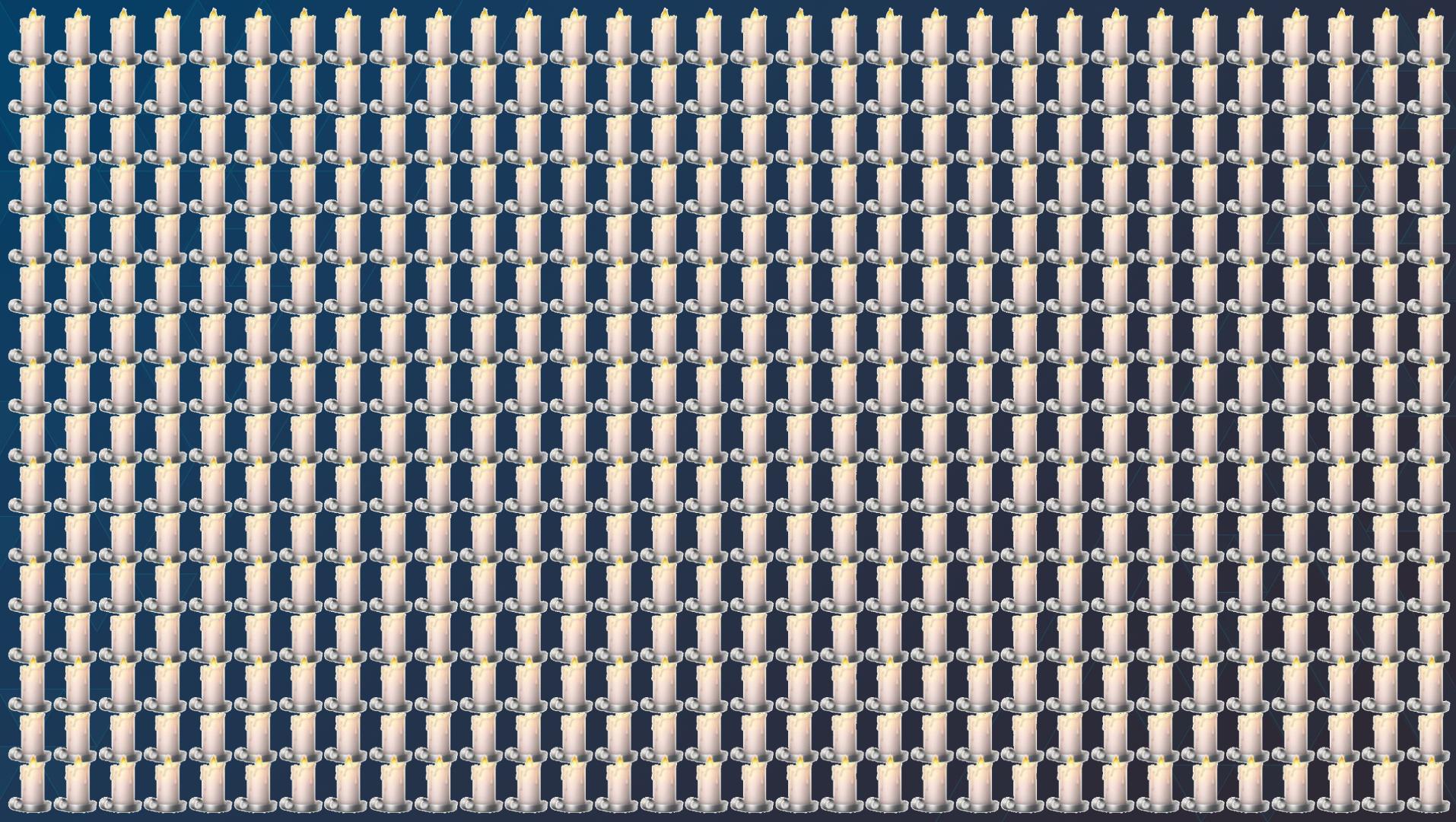


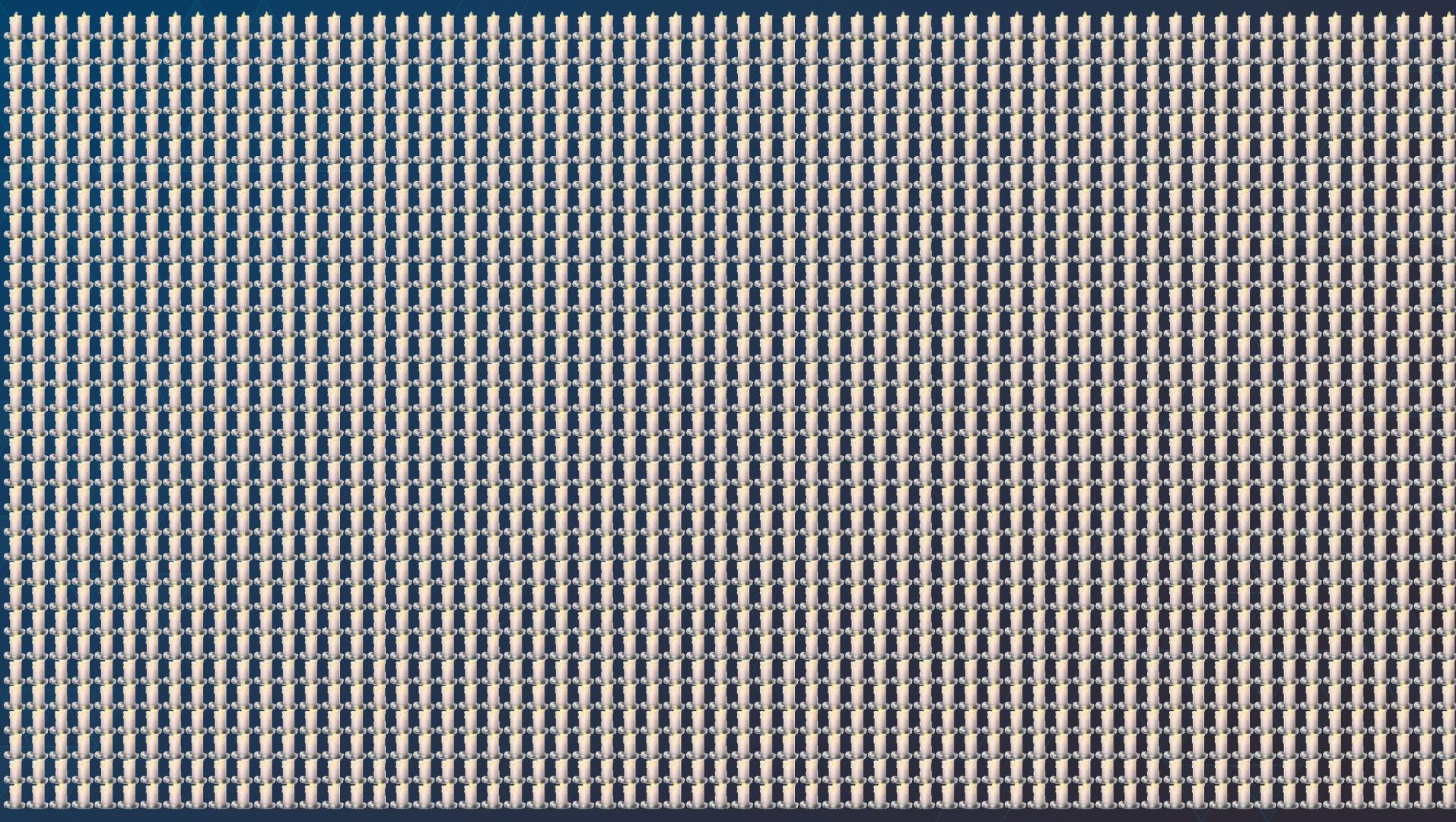
the rebels, the troublemakers,
the round pegs in the square holes.
The ones who see things differently.
They're not fond of rules. And they
have no respect for the status quo.
You can quote them, glorify or vilify
the only thing you can't do is ignore them.
Because they change things.

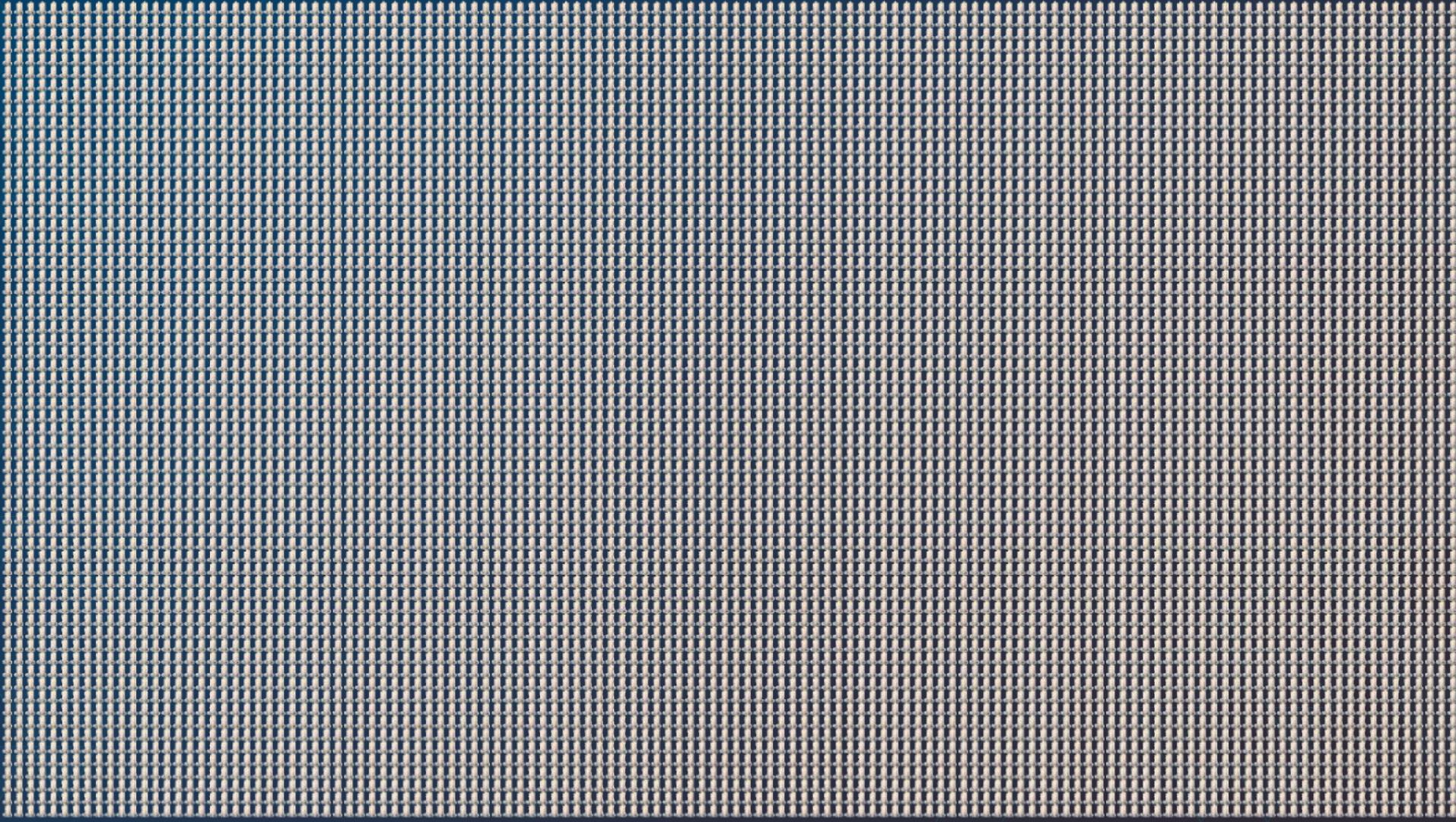
There will be bugs!



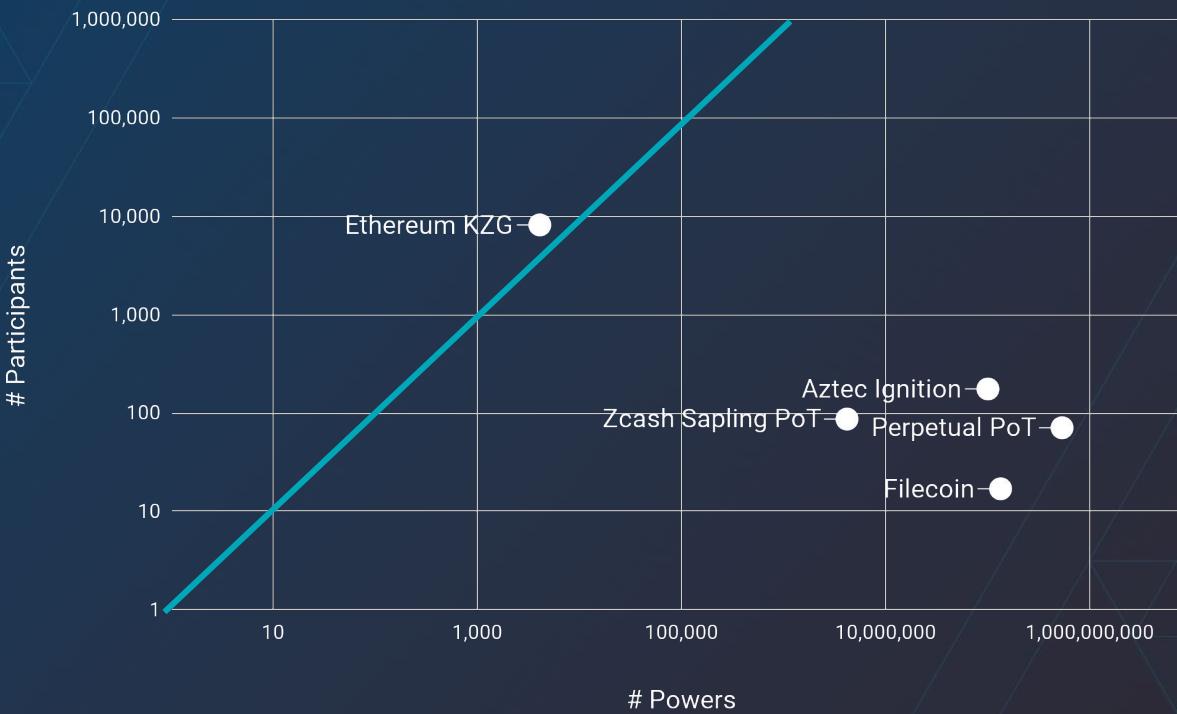




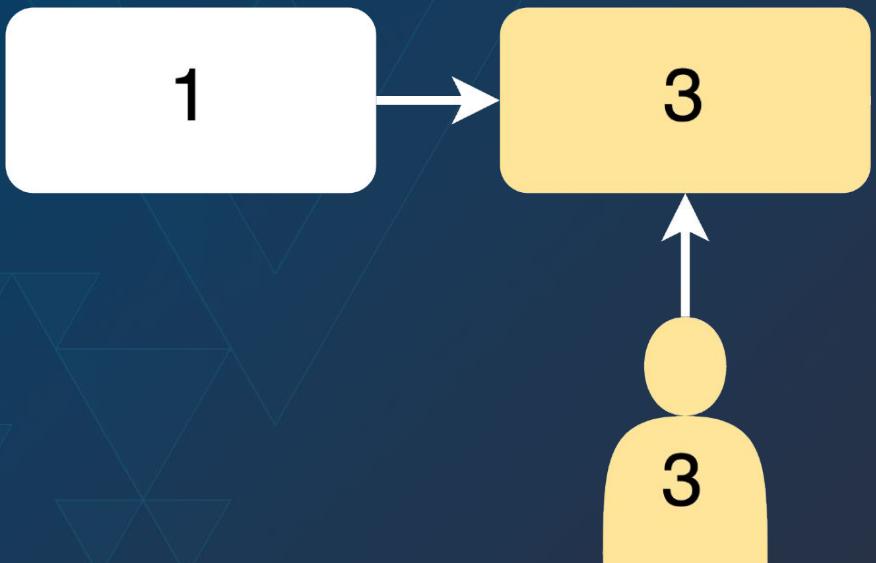




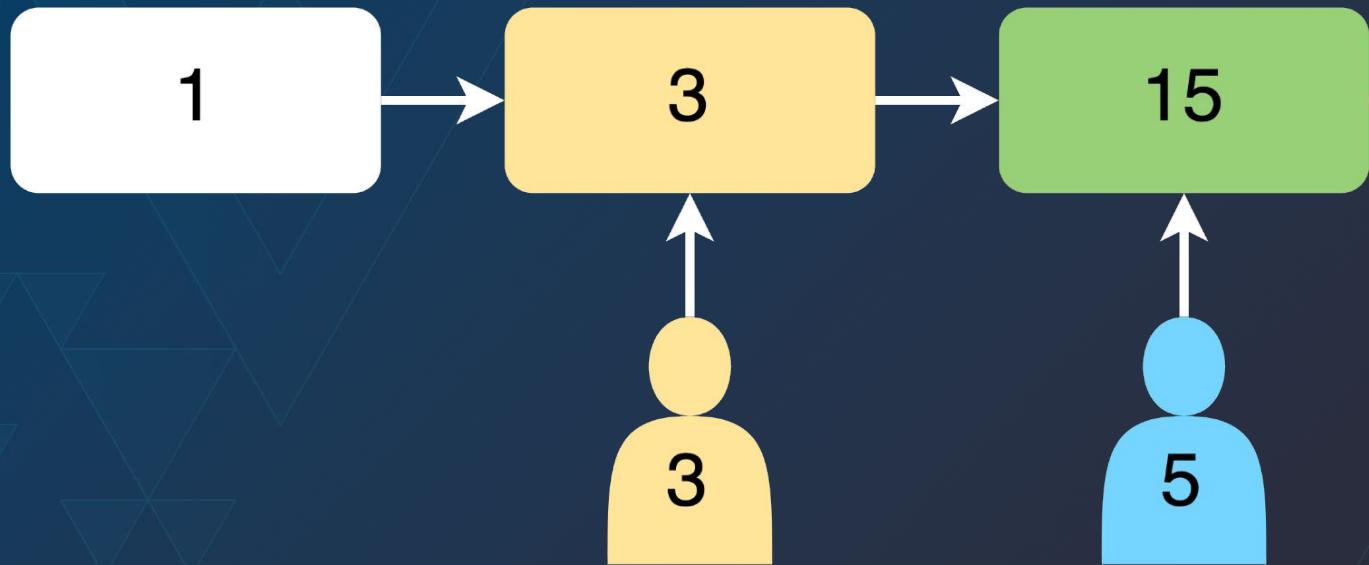
More Participants than Powers!



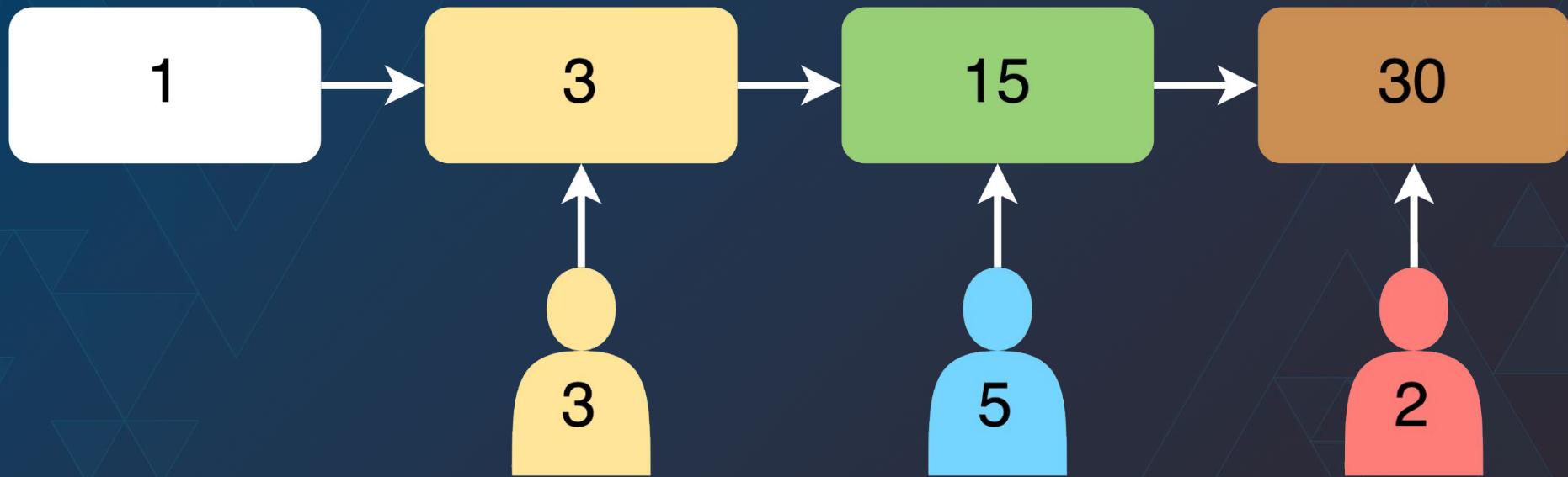
How does the ceremony work



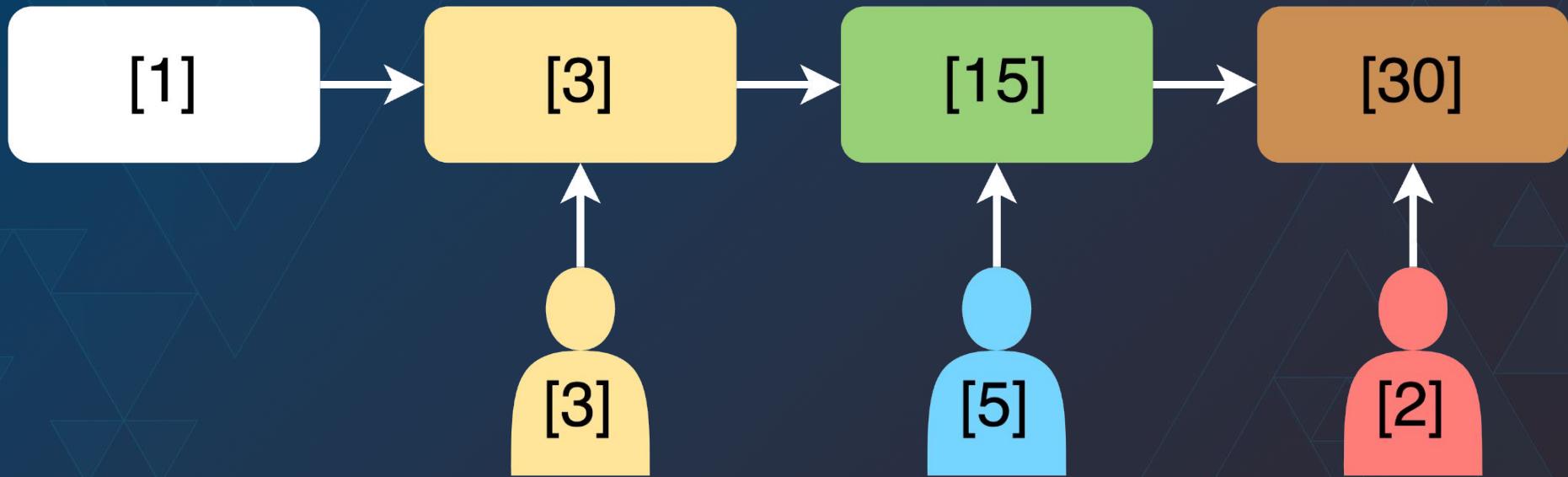
How does the ceremony work



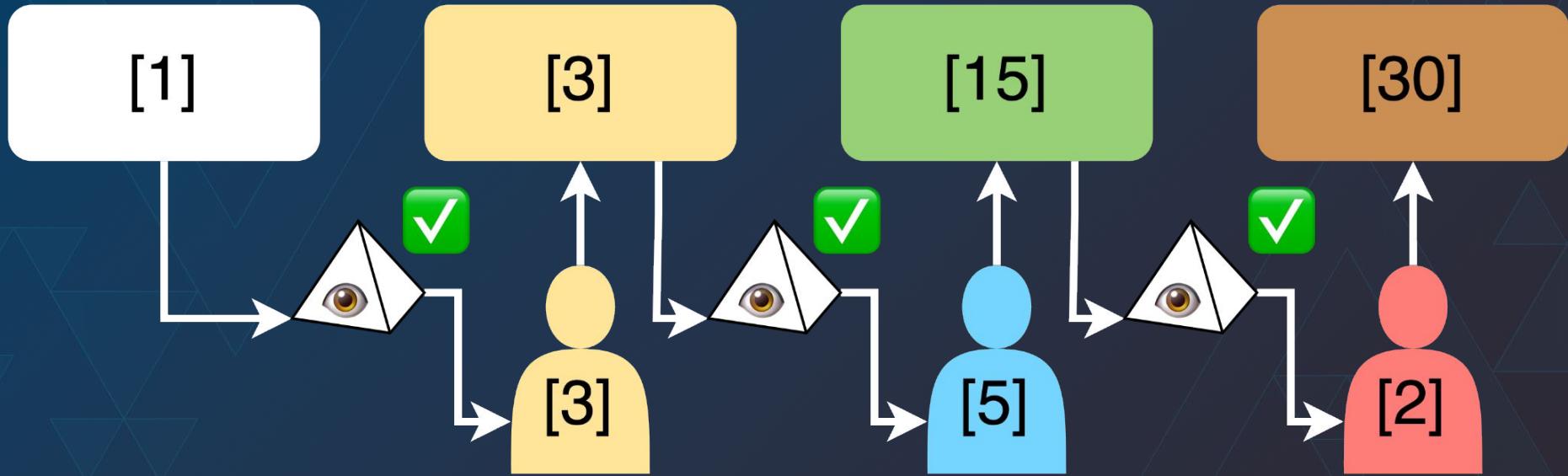
How does the ceremony work

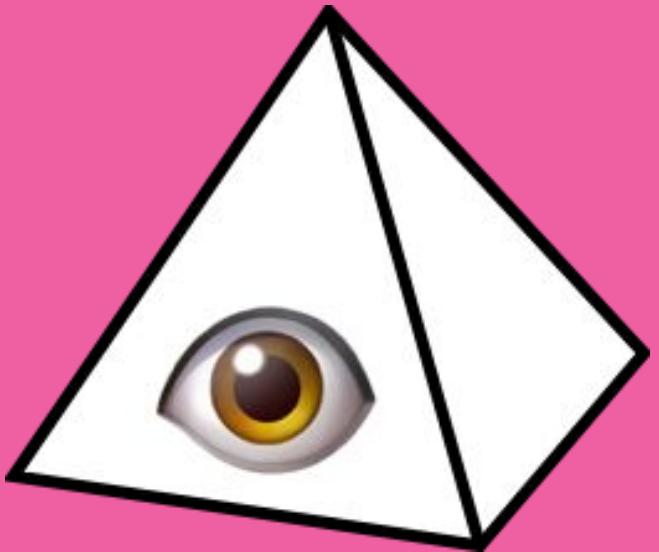


Enter the curves



Sequencer checking things





What could go wrong?

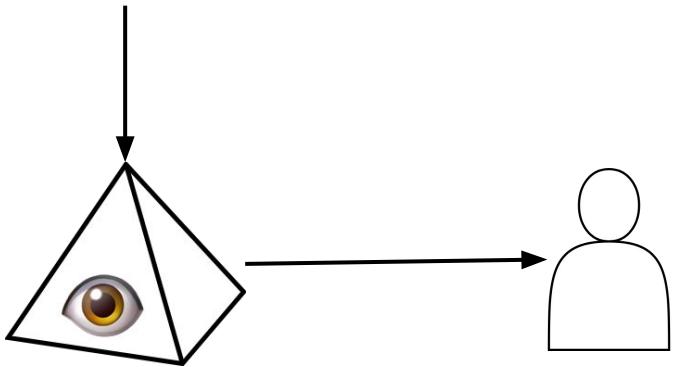
- Censor contributions
 - Lie about verification
 - Not give you a turn
- But, this is attributable!

Powers??

$$[s], [s^2], [s^3], \dots, [s^{2^{12}}]$$

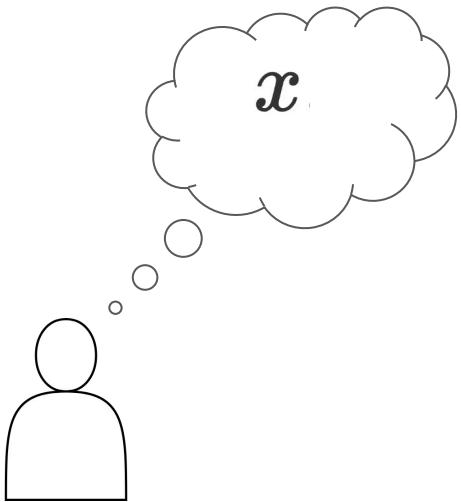
Powers??

$$[s], [s^2], [s^3], \dots, [s^{2^{12}}]$$



Entropy Calculation

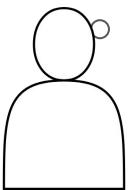
$[s], [s^2], [s^3], \dots, [s^{2^{12}}]$



x

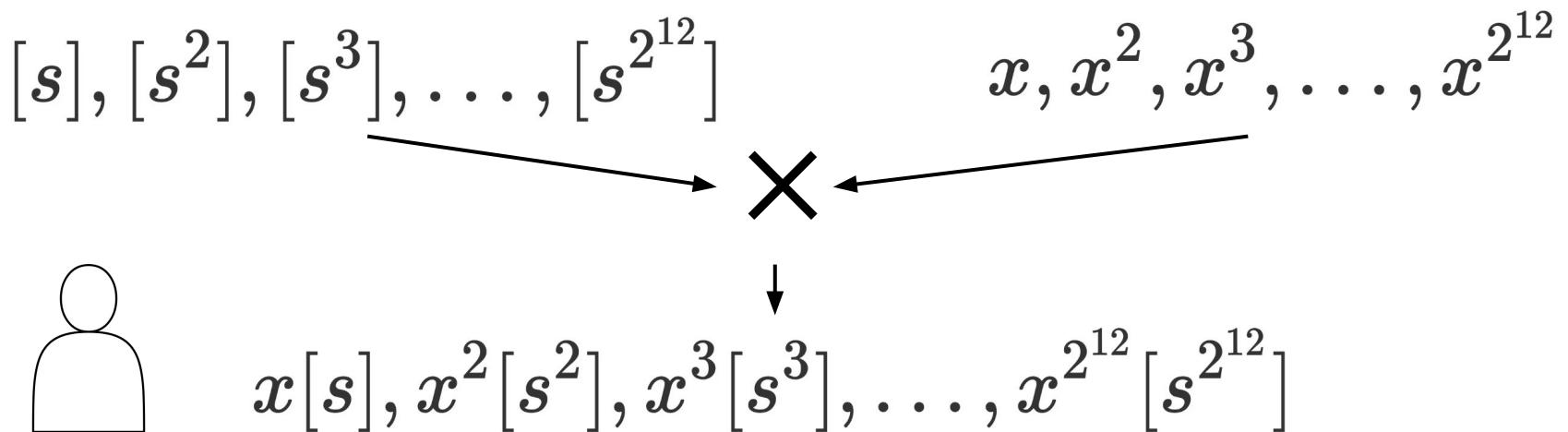
Doing a heckn' big calculate

$[s], [s^2], [s^3], \dots, [s^{2^{12}}]$

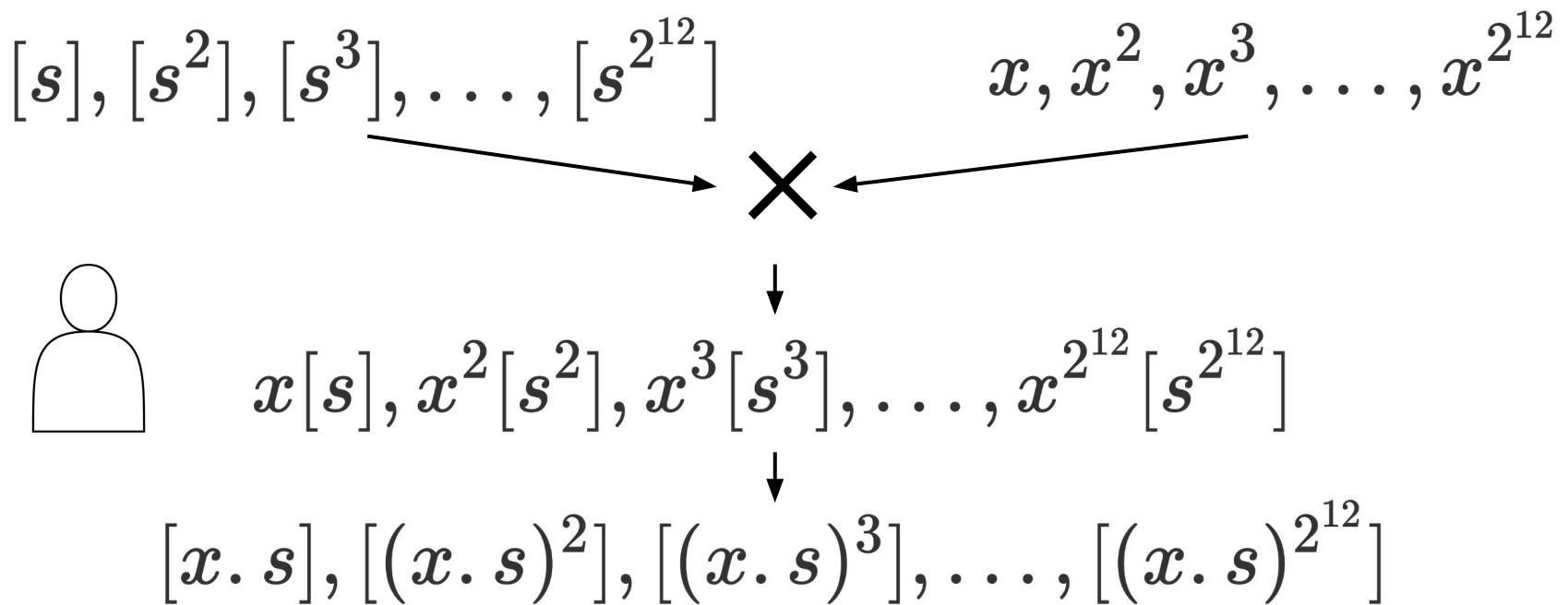


$x, x^2, x^3, \dots, x^{2^{12}}$

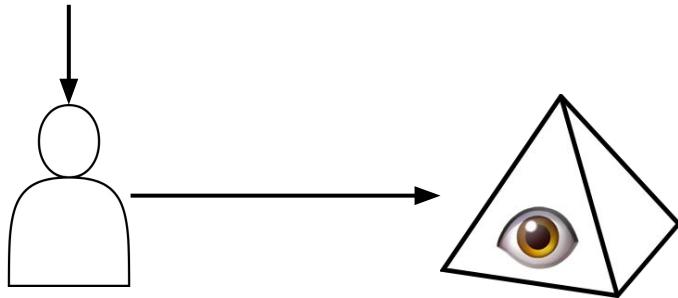
Even more math...



Even more math...



Returning the Powers

$$[x.s], [(x.s)^2], [(x.s)^3], \dots, [(x.s)^{2^{12}}]$$


Grants!

Writing an implementation

-  Roll your own crypto!
-  Implement a client



Crazy randomness generation

-  Vision trip
-  Collect weird entropy
-  Destroy your secret (and PC?!)



Thank you!

Carl Beekhuizen

EF Researcher

carl@ethereum.org



@CarlBeek

Verification

A contribution contains

$$[x \cdot s]_1, [(x \cdot s)^2]_1, [(x \cdot s)^3]_1, \dots, [(x \cdot s)^{2^{12}}]_1$$

$$[s]_2 \quad [x \cdot s]_2$$

Update check

$$e([x \cdot s]_1, g_1) \stackrel{?}{=} e([x]_1, [s]_2)$$

Powers check

$$\forall i \in \mathbb{Z}^+, i \leq 2^{12} : \quad e([x \cdot s]^i_1, [x \cdot s]_2) \stackrel{?}{=} e([(x \cdot s)^{i+1}]_1, g_1)$$

Why KZG?

	Merke (SHA 256)	Merkle (Arithmetic)	KZG	IPA
Verification inside SNARK	Prohibitive	Doable	Cheap	Cheap
Proof Verification (& low deg. proofs)	Hard	SNARKs needed for easy verify	Easy	$O(N)$ time
Proof generation in $O(n^* \log(n))$	Yes	Yes	Yes	Only with larger proof-size
Commit linearity	No	No	Yes	Yes
Proof linearity	No	No	Yes	No
Security assumption	Very safe	Immature hash functions	Trusted setup + Quantum vuln.	Quantum vuln.