RAC performance tuning

By

Riyaj Shamsudeen



These are receive metrics

RAC CR Wait Events

GC cr block 2-way/3-way

GC cr block busy

GC cr block congested/

GC cr grants congested.

GC cr multiblock request

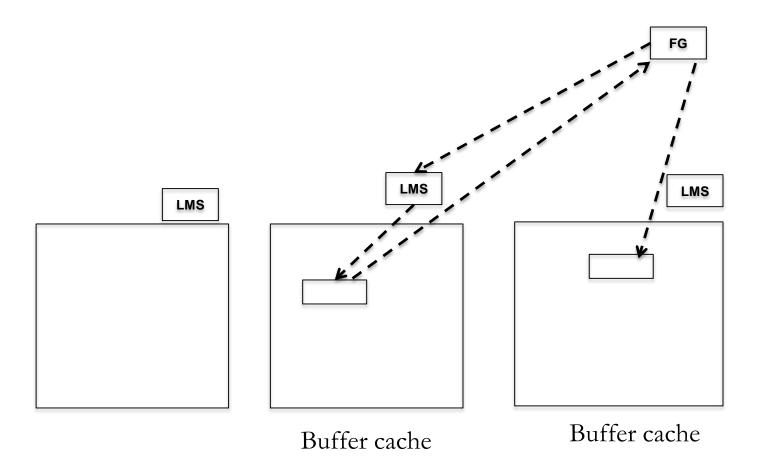
GC cr grants 2-way

CR Wait events

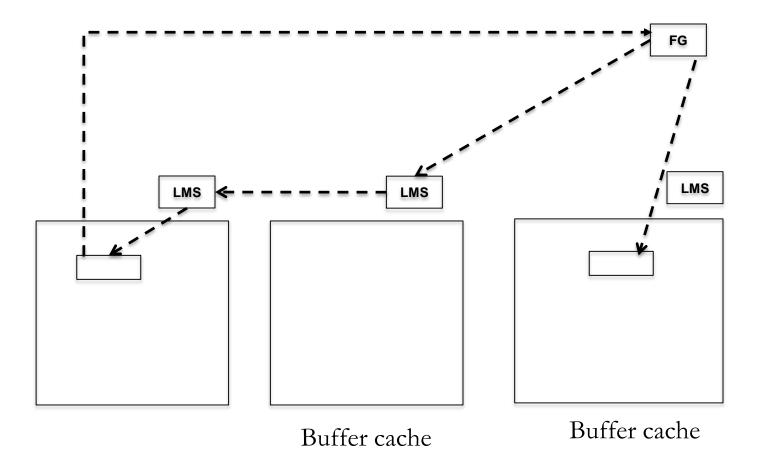
Following are the top wait events associated with CR mode transfers:

gc cr block 2-way gc cr block 3-way	Transfers without congestion or concurrency.
gc cr multi block request	Multi block read
gc cr block busy gc buffer busy (acquire/release)	Concurrency related
gc cr grant 2-way	Grants
gc cr grant congested gc cr block congested	Congestion related

2-way



3-way



Event: GC CR block 2/3-way

• 'gc cr block 2-way': block owner and master is the same instance.

■ 3-way: The owner and master instance are different.

No additional overhead incurred.

Dba_objects.dba_object_id or data_object_id

Analysis

gc cr block 2-way/3-way are baseline wait events.

If these events have higher impact, then differentiate.

Elapsed time per event wait is high.

Numerous waits for these events.

Concurrency or congestion issues are not included in to these events.

Diagnostics – Longer individual events (1)

 Review the histogram for this event using event_histogram.sql script.

41% of waits took between 2-4ms in this example below. WAIT_TIME_MILLI WAIT_COUNT INST_ID EVENT PER 1 gc cr block 2-way 1 3720856 1.11 148016413 1 gc cr block 2-way 44.25 1 gc cr block 2-way 140006974 41.86 1 gc cr block 2-way 40140870 12 .74 1 gc cr block 2-way 16 2491886 1 gc cr block 2-way 32 43253 .01 1 gc cr block 2-way 8192 0 1 gc cr block 2-way 16384 0 24

Demo: event_histogram.sql

Possible reasons

- High CPU usage in the nodes.
- Network performance or Network configuration issue.
- Platform issues as SMP scaling or NUMA related.

Good baseline indicators for cache fusion performance.

Diagnostics – Excessive waits (2)

 Use ASH or AWR to understand which object is inducing numerous waits for this wait event.

@ash_gcwait_to_obj.sql Enter value for event_name: gc cr block 2-way INST_ID OWNER OBJECT_NAME OBJECT_TYPE CNT 1 APPLSYS 118 **TABLE** FND_CONCURRENT_PROCESSES 144 1 INV **TABLE** MTL_SERIAL_NUMBERS 1 INV MTL_TRANSACTIONS_INTERFACE_N1 INDEX 176 184 1 APPLSYS FND_CONCURRENT_REQUESTS **TABLE** 1 INV 211 **TABLE** MTL_MATERIAL_TRANSACTIONS 1 INV 216 MTL_TRANSACTIONS_INTERFACE TABLE Undo Header/Undo block? 18483

For undo header blocks/undo blocks, current_obj# is set to 0 and for undo blocks, curent_obj# is set to -1.

Recommendations

Consider application affinity.

Increase SGA size.

RAC CR Wait Events

GC cr block 2-way/3-way

GC cr block busy

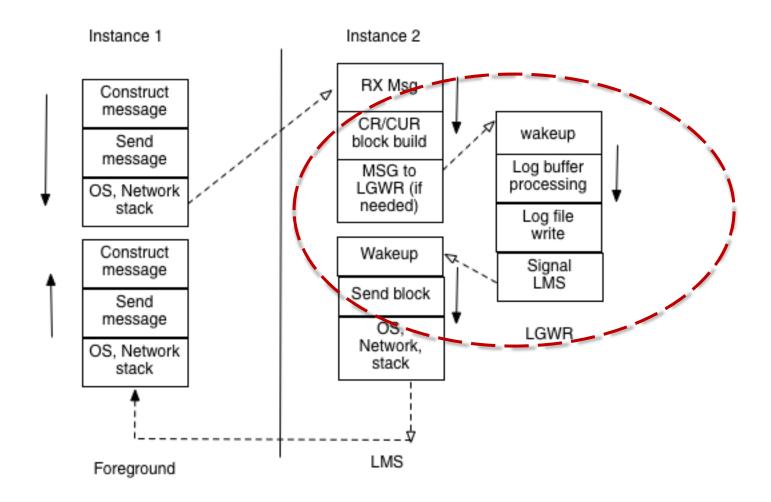
GC cr block congested/

GC cr grants congested.

GC cr multiblock request

GC cr grants 2-way

gc cr block busy (one possibility)



Event: GC CR block busy

- Implies that building of the block had to wait for an event, such as 'log flush sync'.
- For example, if the block is undergoing changes, then LMS process need to create a consistent version of the block applying undo records.
- But, LMS process must wait for LGWR to do a log flush sync event before sending the block. Log flush sync event is similar to log file sync event during commit processing.

RAC CR Wait Events

GC cr block 2-way/3-way

GC cr block busy

GC cr block congested/

GC cr grants congested.

GC cr multiblock request

GC cr grants 2-way

Event: gc cr block congested/gc cr grants congested

- These wait events indicate that there were <u>CPU resource</u> <u>starvation or higher global cache workload</u> issues.
- For example, sudden spikes in global cache workload,
 CPU/memory starvation etc.
- Reducing CPU usage by tuning costly SQL statement.
- As with any wait events, how much is the impact?

Event: gc cr multi block request

Full table scan.

■ In 11gR2, just one request is sent to read a range of blocks.

nam='gc cr multi block request' ela= 19329 file#=8 block#=3557 class#=1 obj#=77779

nam='db file scattered read' ela= 2483 file#=8 block#=3550 blocks=8 obj#=77779

Recommendations

- Reducing full table scan will reduce the time spent on this event. But, do you need to?
- Reduce Dynamic Sampling in RAC.
- Direct path reads and Adaptive direct path read features reads the block directly in to the PGA of the process, even for non-parallel processes.
- So, direct path reads do not suffer from GC wait events.

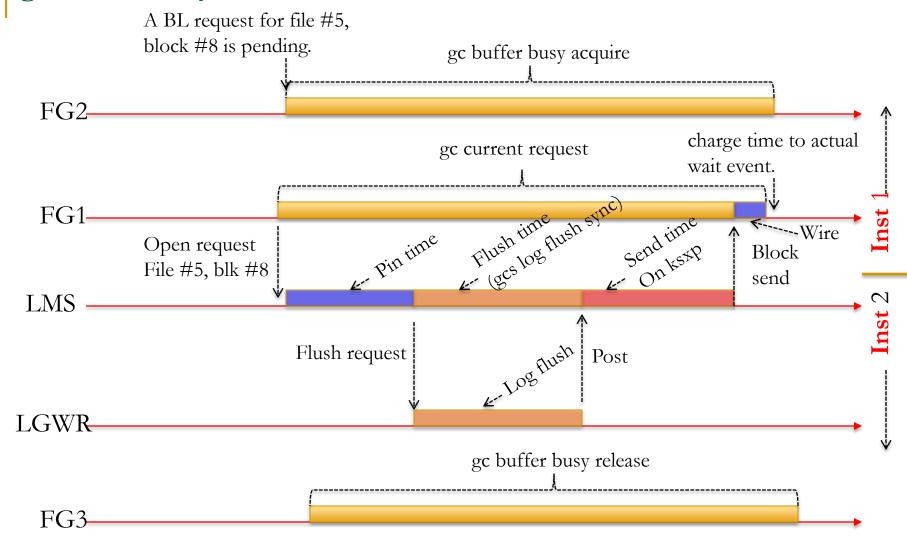
nam='direct path read' ela= 5024 file number=8 first dba=3648 block cnt=32 obj#=77779 nam='direct path read' ela= 479 file number=8 first dba=3712 block cnt=32 obj#=77779

CURRENT Wait events

Following are the top wait events associated with CURRENT mode transfers:

gc current block 2-way	Transfers without
gc current block 3-way	congestion or concurrency.
gc current multi block request	
gc current block busy	
gc buffer busy	Concurrency related
gc current grant 2-way	Grants
gc current grant congested	Congestion related
gc current block congested	O

gc buffer busy



Event: GC Buffer busy waits

- GC Buffer Busy waits are usually symptoms.
- GC buffer busy waits indicates that buffer is busy waiting for some sort of Global event.
 - →■ Another session is working on that buffer and that session is waiting for a global cache event.
 - → We need to understand why that session 2 is waiting for global cache event.
- 11g differentiates this event in to two events: 'gc buffer busy acquire' and 'gc buffer busy release'

Example

Here is one example, where one session stuck waiting for a block.

```
WAIT #65: nam='gc current request' ela= 1220754 file#=1017 block#=1198237 id#=33554433 obj#=1598669 wait #65: nam='gc current request' ela= 1220767 file#=1017 block#=1198237 id#=33554433 obj#=1598669 wait #65: nam='gc current request' ela= 1220810 file#=1017 block#=1198237 id#=33554433 obj#=1598669 wait #65: nam='gc current request' ela= 1220756 file#=1017 block#=1198237 id#=33554433 obj#=1598669 wait #65: nam='gc current request' ela= 1220763 file#=1017 block#=1198237 id#=33554433 obj#=1598669 wait #65: nam='gc current request' ela= 1220763 file#=1017 block#=1198237 id#=33554433 obj#=1598669 id#=33554433 obj
```

■ Gv\$session shows that <u>other</u> sessions are waiting for 'gc buffer busy' waits.

	P1	P2	Р3	COUNT(*)
• • •	1017	1198215	65537	35
	995	707545	65537	37
	382	123951	65537	60
	96	373926	66336	65
	1017	1198237	65537	235

Diagnostics

Is the problem with one or few objects or is it widespread?

Few objects

Widespread

Probably, an application coding or concurrency issue.

Insert statement or SELECT?

Possible environmental issue.
Review performance metrics
at server level.

For Insert statements, use hash partitioning or other methods to improve concurrency

Tune SQL statement or improve concurrency in SQL Execution

Example (AWR)

Many 'gc buffer busy' waits will result in high receive time for both CR and CUR traffic.

Global Cache and Enqueue Services - Workload Characteristics

Avg global enqueue get time (ms):	7.4
Avg global cache cr block receive time (ms):	222.0
Avg global cache current block receive time (ms):	27.5
Avg global cache cr block build time (ms):	0.0
Avg global cache cr block send time (ms):	0.1
Global cache log flushes for cr blocks served %:	2.7
Avg global cache cr block flush time (ms):	15879.9

Example-log file sync (AWR)

High gc buffer busy waits are due to 'log file sync' waits as below.

Top 5 Timed Foreground Events Avg % DB wait time Wait Class Waits Time(s) (ms) Event 45.8 Commit log file sync 2,054 23,720 11548 gc buffer busy acquire 19,505 10,382 532 20.0 Cluster gc cr block busy 4,655 9.0 Cluster 5,407 861 6.6 Configurat enq: SQ - contention 140 3,432 24514 db file sequential read 38,062 1,305 34 2.5 User I/O

Example - Back ground waits (AWR)

LMS process also will wait for 'gcs log flush sync'.

				Avg			
		%Time	Total Wait	wait	Waits	% bg	
Event	Waits	-outs	Time (s)	(ms)	/txn	time	
gcs log flush sync	80,695	51	1,862	23	34.7	32.9	
log file parallel write	44,129	0	880	20	19.0	15.6	
Log archive I/O	1,607	0	876	545	0.7	15.5	
gc cr block busy	729	71	752	1031	0.3	13.3	
db file parallel write	25,752	0	434	17	11.1	7.7	
enq: CF - contention	166	64	307	1850	0.1	5.4	

Example 2 – busy waits (AWR)

These waits are tagged to busy events as these waits for 'gcs log flush sync event'.

CR	Avg Time	(ms)		Current	Avg Time	(ms)			
Inst	Block								
No	Class	All	Immed	Busy (Congst	All	Immed	Busy	Congst
2	data blo	50.5	1.2	9532.1	23.6	11.0	0.9	3550.4	5.0
3	data blo	300.8	2.3	2151.1	6.4	81.9	1.3	6474.9	27.7
3	others	746.7	1.0	1729.1	N/A	50.2	0.8	175.5	29.0
3	undo blo	0.9	0.8	16.0	N/A	N/A	N/A	N/A	N/A
3	undo hea	1986.1	0.8	1.3E+04	N/A	1.1	0.	7 5.	3 N/
2	undo hea	701.5	0.8	1.4E+04	N/A	3974.6	0.9	9 5.2E	+04
2	others	0.8	0.8	N/A	N/A	1.0	0.9	1.3	N/A
2	undo blo	3.0	0.9	10.7	N/A	N/A	N/A	N/A	N/A

Gc buffer busy – Objects related

- ASH or trace files can be used to identify the blocks suffering from excessive 'gc buffer busy' waits.
- Only use this method if there is no other issue affecting background processes.
- To understand gc buffer busy waits:
 - Identify the object and object type
 - Identify the block type
- Ash_gcwait_to_obj.sql can be used to identify the object name and type.

Gc buffer busy – Identify object

@ash_gcwait_to_obj.sql

INST_ID EVENT	OWNER	OBJECT_NAME	OBJECT_TYPE	CNT
1 gc buffer busy acquire	RS	T_GEN_INS_01_N1	INDEX	3
2 gc buffer busy acquire	RS	T_GEN_INS_01_N1	INDEX	10
2 gc buffer busy release	RS	T_GEN_INS_01	TABLE	4

Gc buffer busy – Identify block

<pre>@ash_gcwait_to_block.sql</pre>			
INST_ID EVENT	CURRENT_FILE# CURRENT	_BLOCK#	CNT
2 gc buffer busy acquire	4	103582	4
2 gc buffer busy acquire	4	103607	3
2 gc buffer busy acquire	4	103603	3
1 gc buffer busy acquire	4	103582	2
2 gc buffer busy release	4	103615	2
1 gc buffer busy acquire	4	103583	1
2 gc buffer busy release	4	103586	1
2 gc buffer busy release	4	103613	1
		1	

In this example, many different blocks are involved in 'gc buffer busy' waits. A symptom of heavy insert load.

Demo: ash_gcwait_to_block.sql

Object type	Block type	Possible issue(s)
Table	Segment header	Freelists, freelist groups, deletes from one node, ASSM bugs etc
	Segment blocks	Heavy insert workload, heavy update to few blocks, SQL performance issues scanning few objects aggressively.
Index	Leaf block	Numerous inserts on indexed columns, sequence generated keys
	Root block/branch block	Insert + numerous index scans, concurrent inserts and deletes etc
Undo	Undo header block	Numerous short transactions
	Undo block	Long pending transaction and CR block generation
Dictionary	SEQ\$ (object)	Sequences with nocache or order set and excessive access to sequence.

Measuring performance correctly

- It is essential to **review metrics from all nodes**. Common mistake I see is that DBAs review the metrics from just one node.
- Performance problem from one node usually affects other nodes. Funnily, unhealthy node itself might not see the slowness.
- Most Global cache wait events are symptoms, not necessarily problems.
- For example, Log file write slowness in one node can induce massive 'gc buffer busy' waits in other nodes.

Reviewing all nodes

- It is easy to create AWR reports from all nodes using my script: Refer awrrpt_all_gen.sql.
 - [Don't forget that access to AWR report needs license]
- Or use my script gc_traffic_processing.sql to review Global cache performance.

Default collection period is 60 seconds.... Please wait for at least 60 seconds...

blk TX 	CUR	R fls tm	UR pin tm 	JR b1k TX	CR snd tm	CR fls tm	CR bld 	CR blk Tx	Inst
.23	1	.2	1.62	34909	.23	.88	.08	67061	2
.26		.08	.61	28303	.26	2.19	.17	38207	3
.19		.24	1.76	40578	.2	1.76	.06	72820	4
.25		.44	2.69	30717	.23	2.42	.09	84355	5
		.24	1.76	40578	. 2	1.76	.06	72820	3 4 5

Demo: swing, awrrpt_all_gen.sql awrrpt_all_range_gen.sql gc_traffic_processing.sql awrgrpt.sql

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33

Caution

- Don't use gv\$views directly to find averages as that can be misleading. Use AWR reports or custom scripts.
- gv\$views are aggregated data and persistent from the instance restart.
- For example this query output can be <u>misleading</u>
- (As the output of this query is aggregated over the life of the instances):

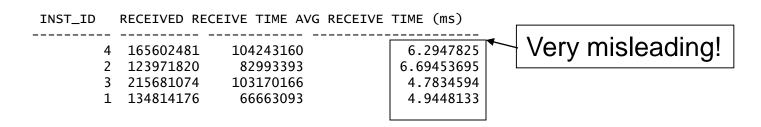
```
select b1.inst_id, b2.value "RECEIVED",
b1.value "RECEIVE TIME",
  ((b1.value / b2.value) * 10) "AVG RECEIVE TIME (ms)"
from gv$sysstat b1, gv$sysstat b2
where b1.name = 'gc cr block receive time' and
b2.name = 'gc cr blocks received' and b1.inst_id = b2.inst_id
```

gc_traffic_print.sql

You can use my script to print global cache performance data for the past minute.

 Inst	CR blocks Rx	 CR time 	CUR blocks Rx	 CUR time	CR blocks Tx	 CUR blocks Tx 	 Tot blocks
1	40999		!	4.82	25070	17855	- ·
3	12471 28795			5.28 3.97	31269 28946	9772 4248	
4	33105	4.54 	12136	4.68	29517	13645	88403

During the same time frame, output of the script from prior slide:



Demo: gc_traffic_print.sql

GC Send time

- Global cache send time is important metrics.
- Global cache Send time in one node can affect Receive time in all other nodes.
- For that matter, GC receive performance of unhealthy node will not be worse, but other nodes will suffer from GC receive performance
- For example, if node 2 is suffering from GC send latencies, other nodes will suffer from GC receive latencies, but node 2 might not suffer from GC receive latencies.

AWR –Send metrics

These are Send metrics

Send metrics are broken down to various metrics. This section shows efficiency of LMS processes:

```
Avg global cache cr block build time (ms): 0.1

Avg global cache cr block send time (ms): 0.3

Global cache log flushes for cr blocks served %: 18.7

Avg global cache cr block flush time (ms): 4.9
```

■ LMS processing delay ~=

Time to build block +

Time to wait for Log flush sync +

Time to send the block over the network.

AWR: GC Efficiency section

 Global cache efficiency section shows the local vs remote buffer cache access.

Access to remote cache
Should be less than 10%

Disk access should be less than 10% for OLTP.

GC CR latency

■ GC CR latency ~=

Time spent in sending message to LMS +

LMS processing (building blocks etc) +

LGWR latency (if any) +

LMS send time +

Wire latency

Processing in the remote nodes

Averages can be misleading. Always review both total time and average to understand the issue.

Example – high undo – high GC timings

Three instances are suffering from CR latency, except instance
 2.

Wait time	Node 1	Node 2	Node 3	Node 4
Avg. CR block receive time	18.2	6.7	20.0	17.3
Avg CUR block receive time	14.6	5.0	11.6	17.3

■ In RAC, node suffering from chronic issues causes GC performance issues in other nodes. With that logic in mind, node 2 should be suffering from chronic issues.

Example – high undo – latency breakdown

Sum of flush time is higher, but it is comparable across the cluster.

But, notice the build time in node 2.

Statistics	Node 1	Node 2	Node 3	Node 4	Total
		j	·		
gc cr block build time	11,392	148,666	5,267	6,632	171,957
Gc cr block flush time	56,634	75,751	34,406	53,031	219,822
Gc cr block send time	9,153	7,779	4,018	7,905	28,855

Example – high undo – CR fabrication

■ For CR blocks, time is spent in building blocks, which indicates consistent block generation.

Very high value compared to other nodes.

Statistics	Node 1	Node 2/	Node 3	Node 4
data blocks consistent Reads – undo records applied	2,493,242	86,988,512	3,090,308	7,208,575
db block changes	6,276,149	43,898,418	20,698,189	14,259,340

GC CUR latency

■ GC CUR latency ~=

Time spent in sending message to LMS +

LMS processing (pin blocks etc, defer) +

LGWR latency (if any) +

LMS send time +

Wire latency

Statistics

gc current block flush time gc current block pin time gc current block send time

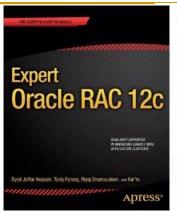
Gv\$instance_cache_transfer

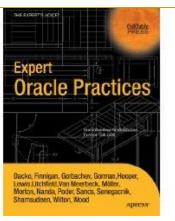
- This is an ultra important view to see the traffic of packets received.
- 11g introduces many columns adding the time component to this view.

Gv\$instance_cache_transfer					
Name	Null?	Туре			
INST_ID		NUMBER Current instance			
INSTANCE		NUMBERInstance received from			
CLASS		VARCHAR2(18)			
LOST		NUMBER Block class: data block, undo header, undo			
LOST_TIME		NUMBER block, 1st level bmb etc			
CR_BLOCK		NUMBER			
CR_BLOCK_TIME		NUMBER			
CR_2HOP		NUMBER			
CR_2HOP_TIME		NUMBER — Time tracking column introduced in 11g.			
CR_BUSY		NUMBER			
CR_BUSY_TIME		NUMBER			
CR_CONGESTED		NUMBER			

THANK YOU









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