

Today's Content:

Modular Arithmetics

Basic And/or principle

Intro to Subsets vs Subsequences

Subsequence with given sum

Modular Arithmetics

$$\left. \begin{array}{l} \text{any} \\ \text{number} \end{array} \right\} \% M = \begin{array}{cc} \min & \max \\ \{0, & M-1\} \end{array}$$

$$(a + b) \% m = (a \% m + b \% m) \% m$$

$$(a \times b) \% m = (a \% m \times b \% m) \% m$$

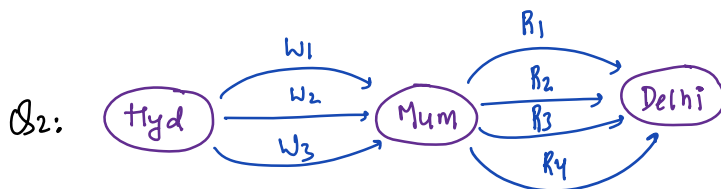
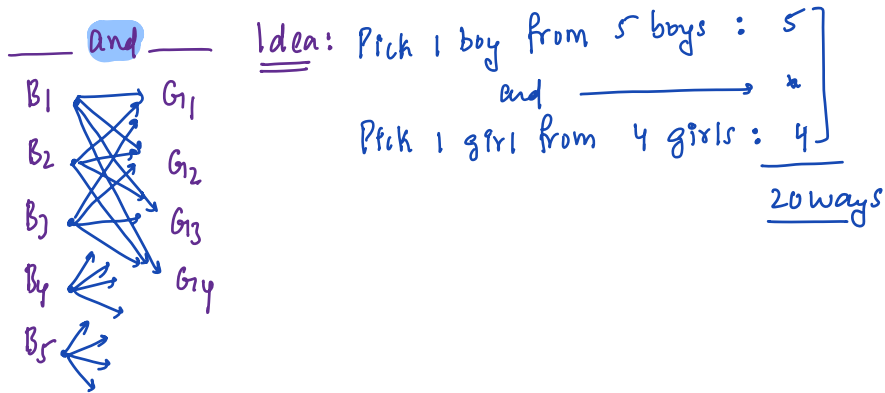
$$\underline{a \quad b \quad m} \quad (a - b) \% m = (a \% m - b \% m + m) \% m$$

$$14 \quad 7 \quad 4 \quad \frac{(14-7) \% 4}{3} = \frac{(14 \% 4 - 7 \% 4 + 4) \% 4}{(2 - 3 + 4) \% 4} = 3 \% 4 =$$

$$11 \quad 7 \quad 3 \quad \frac{(11-7) \% 3}{1} = \frac{(11 \% 3 - 7 \% 3 + 3) \% 3}{(2 - 1 + 3) \% 3} = 4 \% 3 = 1$$

And/OR Principle

Q1. 5 girls & 4 boys, how many pairs : $5 \times 4 = 20$ pairs



ways to go from Hyd → Delhi via Mumbai

1. Ways Hyd → Mum and ways Mum → Delhi

$$3 \times 4 = 12 \text{ ways}$$

Q3: 3 T/F, how many ways we can answer all of them.

Q ₁	Q ₂	Q ₃	
T	T	T	# ways to Q ₁ : 2 ways
F	F	F	and *
T	T	T	# ways to Q ₂ : 2 ways
T	T	F	and *
T	F	T	# ways to Q ₃ : 2 ways
T	F	F	
F	T	T	
F	T	F	
F	F	T	
F	F	F	

obs

1. and = X

2. OR = +

8 ways

Subarrays vs Subsequence vs Subsets

Subarray: Continuous part of an array.

Subsequence: Sequence obtained by deleting none or more ele from arr[]

1. Data should be arranged based on inc order of index
2. Empty Subsequence is also valid.

Ex: arr[6] = { 0 1 2 3 4 5 }
 { 3 2 1 9 6 8 }

{ 1 6 8 } Subseq → { * * 1 * 6 8 }

→ { 3 4 0 } Order miss → { 3 * * 9 6 * }

{ 2 9 8 } Subseq → { * 2 * 9 * 8 }

{ 1 9 6 8 } Subseq → { * * 1 9 6 8 }

{ 2 1 9 6 } Subseq → { * 2 1 9 6 * }

{ } Subseq → { * * * * * }

{ 3 2 1 9 6 8 } Sub → { 3 2 1 9 6 8 }

Subarr[] vs Subsequence

arr[6] = { 0 1 2 3 4 5 }
 { 3 2 1 9 6 8 }

Subarray: Subsequence

{ 3 2 1 } Yes

{ 1 9 6 8 } Yes

{ 3 1 6 8 } No

{ 2 1 9 } Yes

Yes obs:

1. All subarrays are Subsequences Yes

2. All subsequences are Subarrays No

Yes

Yes

$arr[3] = \begin{matrix} 0 & 1 & 2 \\ 3 & 1 & 8 \end{matrix} \xrightarrow{\text{sort[] in inc}} arr[] = \begin{matrix} 0 & 1 & 2 \\ 1 & 3 & 8 \end{matrix}$

<u>man</u>	<u>sum</u>	<u>All Subsequences</u>		<u>All Subsequences</u>	<u>sum</u>	<u>man</u>
—	0	{ }	←	{ }	0	—
3	3	{ 3 }	←	{ 3 }	3	3
1	1	{ 1 }	←	{ 1 }	1	1
8	8	{ 8 }	←	{ 8 }	8	8
3	4	{ 3 1 }	← Order diff but data same	{ 1 3 }	4	3
8	11	{ 3 8 }	←	{ 3 8 }	11	8
8	9	{ 1 8 }	←	{ 1 8 }	9	8
8	12	{ 3 1 8 }	← Order diff but data same	{ 1 3 8 }	12	8

obs: If we sort arr[]

1. Order before & after sorting change
But data remains same.

2. Sum of Subsequences }
 Max of Subsequences } Are not effected by
 Min of Subsequence } sorting

Subsets: Subsets obtained by deleting none or more ele from $arr[]$

1. Empty Subset is also valid.

$arr[3] = \{3 \ 1 \ 8\} \xrightarrow{\text{sort}[]} arr[] = \{1 \ 3 \ 8\}$

<u>All Subsets</u>		<u>All Subsets</u>
$\{ \}$	\longleftrightarrow	$\{ \}$
$\{3\} \{1\} \{8\}$	\longleftrightarrow	$\{1\} \{3\} \{8\}$
$\{3 \ 1\}$	$\xleftarrow{\text{Order: Not Matter Data: Same}}$	$\{1 \ 3\}$
$\{1 \ 8\}$	\longleftrightarrow	$\{1 \ 8\}$
$\{3 \ 8\}$	\longleftrightarrow	$\{3 \ 8\}$
$\{3 \ 1 \ 8\}$	$\xleftarrow{\text{Order: Not Matter Data: Same}}$	$\{1 \ 3 \ 8\}$

Count of Subsets & Subsequences

$arr[3] = \begin{matrix} 0 & 1 & 2 \\ \{6 & 7 & 4\} \end{matrix}$

 Subseq
 Subset = $\{2 \ 4 \ 2 \ 4 \ 2\}$

make choice at index 0 = 2
 and
 make choice at index 1 = 2
 and
 make choice at index 2 = 2
 $\left. \begin{matrix} * \\ * \\ * \end{matrix} \right\} = 8$

$\{ \} \{6\} \{7\} \{4\} \{6 \ 7\} \{7 \ 4\} \{6 \ 4\} \{6 \ 7 \ 4\}$

$arr[4] = \{7 \ 2 \ 9 \ 5\} =$

 Seq/sub = $\{2 * 2 * 2 * 2\} = 2^4 = 16$

$arr[n] = \{a_1 \ a_2 \ a_3 \ \dots \ a_n\}$

Seq/sub = $\{2 * 2 * 2 \dots 2\} = 2^n$

Q1) Given an $arr[N]$, check if There exists a subsequence with $sum = k$

Constraints

$$1 \leq N \leq 20$$

$$1 \leq arr[i] \leq 10^6$$

Ex: $arr[5] = \{ \overset{0}{2} \ \overset{1}{5} \ \overset{2}{3} \ \overset{3}{11} \ \overset{4}{7} \}$

$k=20 : \{2 \ 11 \ 7\}$ True

$k=8 : \{5 \ 3\}$ True

$k=9 : \{2 \ 7\}$ True

$k=19 : \{5 \ 3 \ 11\}$ True

$k=50 : \text{False}$

Idea: For all subsequence get sum & compare $= k$.

If There exists such sequence: Return True

If No such sequence: False

Generate all Subseq:

$\begin{matrix} 0 & 1 & 2 \\ \rightarrow & \rightarrow & \rightarrow \end{matrix}$

 $A = [7 \ 5 \ 3]$

 $\rightarrow \text{delete an ele} : 0$

 $\rightarrow \text{select an ele} : 1$

decimal	Bi		All Sub
0	<u>0 0 0</u>	\rightarrow	$\{\}$
1	0 0 1	$=$	$\{7\}$
2	0 1 0	$=$	$\{5\}$
3	<u>0 1 1</u>	\rightarrow	$\{7, 5\}$
4	1 0 0	$=$	$\{3\}$
5	1 0 1	$=$	$\{7, 3\}$
6	<u>1 1 0</u>	\rightarrow	$\{5, 3\}$
7	1 1 1	$=$	$\{7, 5, 3\}$

$$A = \begin{bmatrix} 2 & 6 & 4 & 5 \end{bmatrix}$$

<u>decimal</u>	3	2	1	0	<u>subseq</u>
0	0	0	0	0	{ }
1	0	0	0	1	{ 2 }
2	0	0	1	0	{ 6 }
3	0	0	1	1	{ 2 6 }
4	0	1	0	0	{ 4 }
5	0	1	0	1	{ 2 4 }
6	0	1	1	0	{ 6 4 }
7	0	1	1	1	{ 2 6 4 }
8	1	0	0	0	{ 5 }
9	1	0	0	1	{ 2 5 }
10	1	0	1	0	{ 6 5 }
11	1	0	1	1	{ 2 6 5 }
12	1	1	0	0	{ 4 5 }
13	1	1	0	1	{ 2 4 5 }
14	1	1	1	0	{ 6 4 5 }
15	1	1	1	1	{ 2 6 4 5 }

obs:

1) $N=3$: we have 8 sub = $[0, 7]$, check 3 bits

2) $N=4$: we have 16 sub = $[0, 15]$, check 4 bits

$$N: \text{Subseq} = 2^N = [0, 2^N - 1] \quad \text{check n bits}$$

boolean subseq(int arr[], int N, int k) { TC: $O(2^N \times N) = O(N \times 2^N)$ SC: $O(1)$

for (num = 0; num < (1 < N); num++) { // (1 < N) = 2^N

int sum = 0;

// For a num get required subsequence.

for (i = 0; i < N; i++) { // i is indicating bit pos

if ((num >> i) & 1 == True) { // ith bit pos in num is 1 or 0.

sum = sum + arr[i]

if (sum == k) { return True }

return False;

A = [7 5 3] ; num = { 0, $2^3 - 1$ } = { 0, 7 }, i = [0, 2]

num :	S	2	1	0 ... i
0	S = 0	0	0	0
1	S = 7	0	0	1 {7}
2	S = 5	0	1 {5}	0
3	S = 12	0	1 {5}	1 {7}
4	S = 3	1 {3}	0	0
5	S = 10	1 {3}	0	1 {7}
6	S = 8	1 {3}	1 {5}	0
7	S = 15	1 {3}	1 {5}	1 {7}

max N = 20.

$$20 \times 2^{20} < 10^8 = 4 \text{es}$$