EXP NO:8 DATE:

GENERATE THREE ADDRESS CODE FOR A SIMPLE PROGRAM USING LEX AND YACC

AIM:

To design and implement a LEX and YACC program that generates three-address code (TAC) for a simple arithmetic expression or program. The program will:

Recognize expressions like addition, subtraction, multiplication, and division.

Generate three-address code that represents the operations in a way that could be directly translated into assembly code or intermediate code for a compiler.

ALGORITHM:

1. Lexical Analysis (LEX) Phase:

Input: A string containing an arithmetic expression (e.g., a = b + c * d;). Output: A stream of tokens such as identifiers (variables), numbers (constants), operators, and special characters (like =, ;, (), etc.).

1. Define the Token Patterns:

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o ID: Identifiers (variables) are strings starting with a letter and followed by letters or digits (e.g., a, b, result).
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o NUMBER: Constants (e.g., 1, 5, 100).

o OPERATOR: Arithmetic operators (+, -, *, /).

o ASSIGNMENT: Assignment operator (=).

o PARENTHESIS: Parentheses for grouping ((and)).

o WHITESPACE: Spaces, tabs, and newline characters (which should be ignored).

2. Write Regular Expressions for the Tokens:

```
o ID -> [a-zA-Z_][a-zA-Z0-9_]*
```

o NUMBER -> [0-9]+

o OPERATOR -> [\+\-*/]

o ASSIGN -> "="

o PAREN -> [\(\)]

o WHITESPACE -> [\t\n]+ (skip whitespace)

3. Action on Tokens:

o When a token is matched, pass it to YACC using yylval to store the token values.

2. Syntax Analysis and TAC Generation (YACC) Phase:

Input: Tokens provided by the LEX lexical analyzer.

Output: Three-address code for the given arithmetic expression.

1. Define Grammar Rules:

o Assignment:

bash

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statement: ID '=' expr

This means an expression is assigned to a variable.

o Expressions:

bash

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expr: expr OPERATOR expr

An expression can be another expression with an operator (+, -, *, /).

bash

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```
expr: NUMBER
expr: ID
expr: '(' expr ')'
2. Three-Address Code Generation:
o For every arithmetic operation, generate a temporary variable (e.g., t1, t2, etc.)
to hold intermediate results.
o For a = b + c, generate:
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t1 = b + c
a = t1
o For a = b * c + d, generate:
ini
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t1 = b * c
t2 = t1 + d
```

3. Temporary Variable Management:

o Keep a counter (temp_count) for generating unique temporary variable names (t0, t1, t2, ...).

o Each time a new operation is encountered, increment the temp_count to generate a new temporary variable.

4. Rule Actions:

a = t2

o When a rule is matched (e.g., expr OPERATOR expr), generate the TAC and assign temporary variables for intermediate results.

Detailed Algorithm:

- 1. Initialize Lexical Analyzer:
- o Define the token patterns for ID, NUMBER, OPERATOR, ASSIGN, PAREN, and WHITESPACE.
- 2. Define the Syntax Grammar:
- o Define grammar rules for:

Assignments: ID = expr

Expressions: expr -> expr OPERATOR expr, expr -> NUMBER, expr

-> ID, expr -> (expr)

3. Token Matching:

o LEX: Match input characters against the defined regular expressions for tokens.

o YACC: Use the tokens to parse and apply grammar rules.

4. TAC Generation:

o For Assignment:

Upon parsing ID = expr, generate a temporary variable for the result of expr and assign it to the variable ID.

o For Arithmetic Operations:

For each operator (e.g., +, -, *, /), generate temporary variables for intermediate calculations.

5. Output TAC:

o Print the generated three-address code, with each expression and its intermediate results represented by temporary variables.

PROGRAM:

```
LEX file (expr.l)
#include "y.tab.h"
%}
%%
[0-9]+ { yylval.str = strdup(yytext); return NUMBER; }
[a-zA-Z_][a-zA-Z0-9_]* { yylval.str = strdup(yytext); return ID; }
[+\-*/=()] { return yytext[0]; }
[\t\n] { /* Ignore whitespace */ }
. { printf("Unexpected character: %s\n", yytext); }
YACC Program expr.y
%{
#include <stdio.h>
#include <stdlib.h>
#include <string.h>
int temp_count = 0;
char* new_temp() {
char* temp = (char*)malloc(8);
sprintf(temp, "t%d", temp_count++);
return temp;
void emit(char* result, char* op1, char op, char* op2) {
printf("%s = %s %c %s\n", result, op1, op, op2);
void emit_assign(char* id, char* expr) {
printf("%s = %s\n", id, expr);
%}
%union {
char* str;
}
%token <str> ID NUMBER
%type <str> expr term factor
%left '+' '-'
%left '*' '/'
%%
statement : ID '=' expr { emit_assign($1, $3); }
expr: expr'+' term { $$ = new_temp(); emit($$, $1, '+', $3); }
| expr '-' term { $$ = new_temp(); emit($$, $1, '-', $3); }
| term { $$ = $1; }
term : term '*' factor { $$ = new_temp(); emit($$, $1, '*', $3); }
| term '/' factor { $$ = new_temp(); emit($$, $1, '/', $3); }
| factor { $$ = $1; }
factor: '(' expr ')' { $$ = $2; }
| NUMBER { $$ = $1; }
| ID { $$ = $1; }
%%
int main() {
yyparse();
return 0;
}
void yyerror(const char* s) {
```

```
fprintf(stderr, "Error: %s\n", s);
}
```

OUTPUT:

```
$ ./a.out

a = b + c * d

t0 = c * d

t1 = b + t0

a = t1

$ ./a.out

x = (a + b) * (c - d)

t0 = a + b

t1 = c - d

t2 = t0 * t1

x = t2

$ ./a.out

y = m / n + p

t0 = m / n

t1 = t0 + p

y = t1
```

RESULT:

Thus the process effectively tokenizes the input, parses it according to defined grammar rules, and generates the corresponding Three-Address Code, facilitating further compilation or interpretation stages.