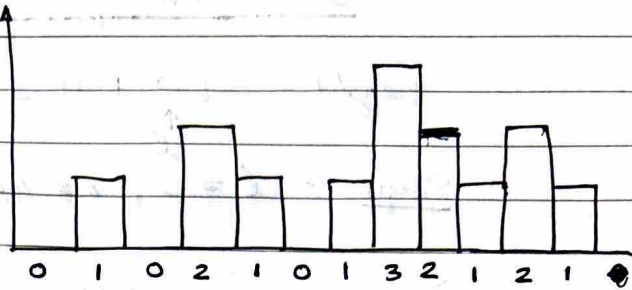
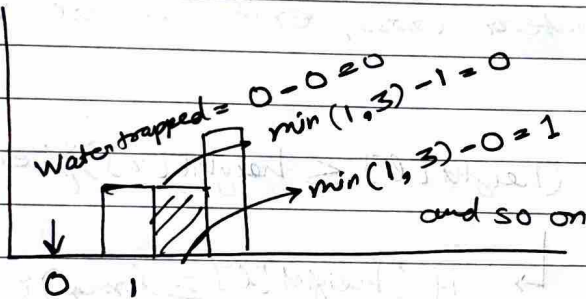


Trapping Rain WaterInput

height = [0, 1, 0, 2, 1, 0, 1, 3, 2, 1, 2, 1]

Output = 6Intuition

For at any index, water trapped = $\min(l_{\max}, r_{\max}) - \text{current height}$



Better solution Using prefix Sum & suffix Sum which gives height max (left) & height max (right) respectively

Step 1 → for loop from $i=0$ to height.lengthStep 2 → prefix = [0, 1, 1, 2, 2, 2, 2, 3, 3, 3, 3, 3]Step 3 → suffix = [3, 3, 3, 3, 3, 3, 3, 3, 2, 2, 2, 1]Step 4 → for loop again from $i=0$ to height.length
int trapped = 0;for ($i=0 \rightarrow \text{height.length}$)trapped += $\min(\text{prefix}[i], \text{suffix}[i]) - \text{height}[i]$ Step 5 → Return trapped

Optimized Approach

Using Two pointers

height = [0, 1, 0, 2, 1, 0, 1, 3, 2, 1, 2, 1]

Step 1 → $l = 0$, $r = \text{height.length} - 1$ ← two pointers

Step 2 → $l_{\max} = 0$, $r_{\max} = 0$ ← to store max in both directions

Step 3 → trapped = 0

Step 4 - In all two pointers ($l \leq r$) ← once pointers cross, exit so use while ($l \leq r$)

Step 5

if ($\text{height}[l] \leq \text{height}[r]$) { // ensure right has some big building

 ↳ if ($\text{height}[l] \geq l_{\max}$) {

$l_{\max} = \text{height}[l]$; // store

 } else { // left max till pointer

 trapped += $l_{\max} - \text{height}[l]$;

 }

$l++$;

}

Step 6 else {

 if ($\text{height}[r] \geq r_{\max}$) { // left has some big building

$r_{\max} = \text{height}[r]$

 } else {

 trapped += $r_{\max} - \text{height}[r]$

 }

classmate $r--$;

Step 7 Return trapped;