

Conversion from NFA to DFA

In this section, we will discuss the method of converting NFA to its equivalent DFA. In NFA, when a specific input is given to the current state, the machine goes to multiple states. It can have zero, one or more than one move on a given input symbol. On the other hand, in DFA, when a specific input is given to the current state, the machine goes to only one state. DFA has only one move on a given input symbol.

Let, $M = (Q, \Sigma, \delta, q_0, F)$ is an NFA which accepts the language $L(M)$. There should be equivalent DFA denoted by $M' = (Q', \Sigma', q_0', \delta', F')$ such that $L(M) = L(M')$.

Steps for converting NFA to DFA:

Step 1: Initially $Q' = \phi$

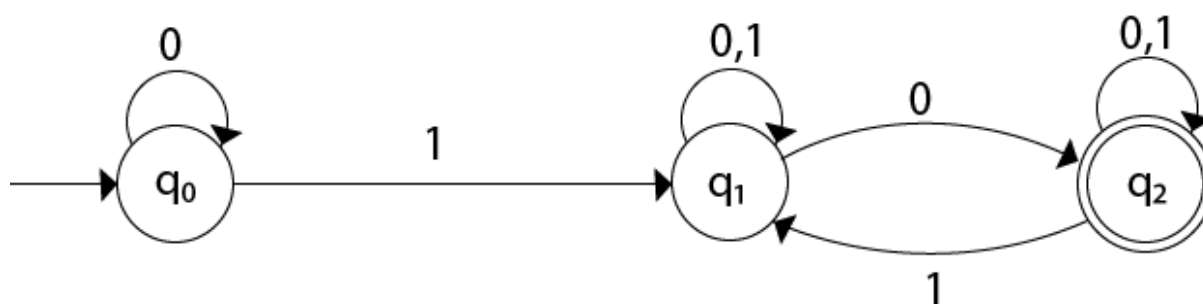
Step 2: Add q_0 of NFA to Q' . Then find the transitions from this start state.

Step 3: In Q' , find the possible set of states for each input symbol. If this set of states is not in Q' , then add it to Q' .

Step 4: In DFA, the final state will be all the states which contain F (final states of NFA)

Example 1:

Convert the given NFA to DFA.



Solution: For the given transition diagram we will first construct the transition table.

State	0	1
$\rightarrow q_0$	q_0	q_1

q1	{q1, q2}	q1
*q2	q2	{q1, q2}

Now we will obtain δ' transition for state q0.

1. $\delta'([q0], 0) = [q0]$
2. $\delta'([q0], 1) = [q1]$

The δ' transition for state q1 is obtained as:

1. $\delta'([q1], 0) = [q1, q2]$ (**new** state generated)
2. $\delta'([q1], 1) = [q1]$

The δ' transition for state q2 is obtained as:

1. $\delta'([q2], 0) = [q2]$
2. $\delta'([q2], 1) = [q1, q2]$

Now we will obtain δ' transition on [q1, q2].

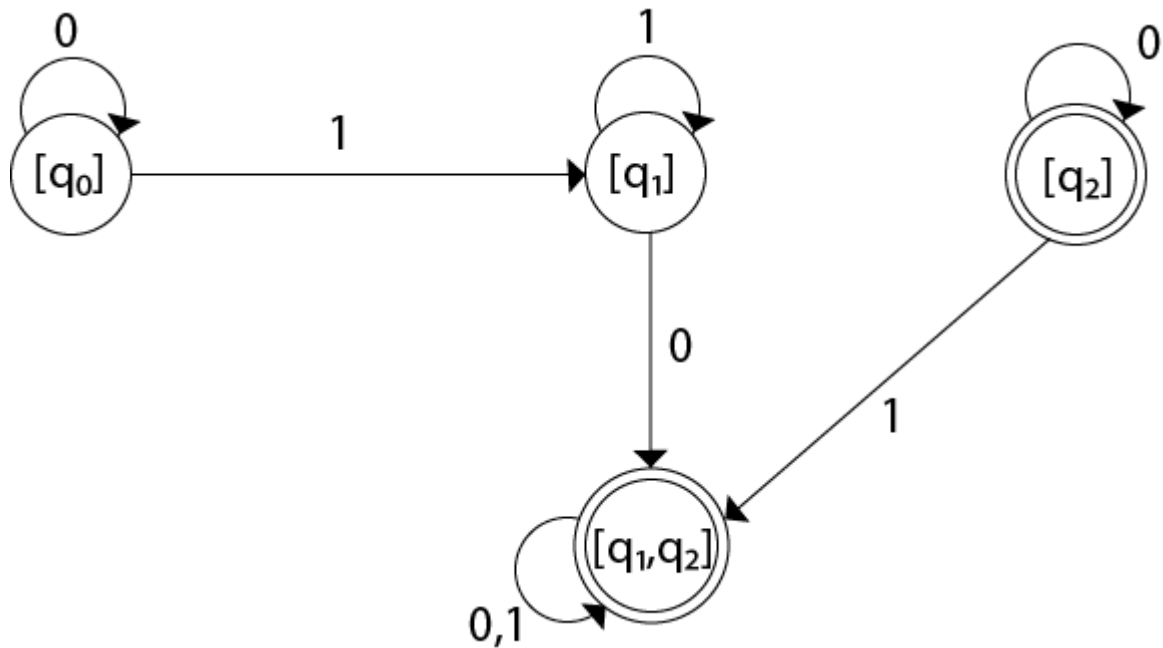
1. $\delta'([q1, q2], 0) = \delta(q1, 0) \cup \delta(q2, 0)$
2. $= \{q1, q2\} \cup \{q2\}$
3. $= [q1, q2]$
4. $\delta'([q1, q2], 1) = \delta(q1, 1) \cup \delta(q2, 1)$
5. $= \{q1\} \cup \{q1, q2\}$
6. $= \{q1, q2\}$
7. $= [q1, q2]$

The state [q1, q2] is the final state as well because it contains a final state q2. The transition table for the constructed DFA will be:

State	0	1
$\rightarrow[q0]$	[q0]	[q1]
[q1]	[q1, q2]	[q1]

*[q2]	[q2]	[q1, q2]
*[q1, q2]	[q1, q2]	[q1, q2]

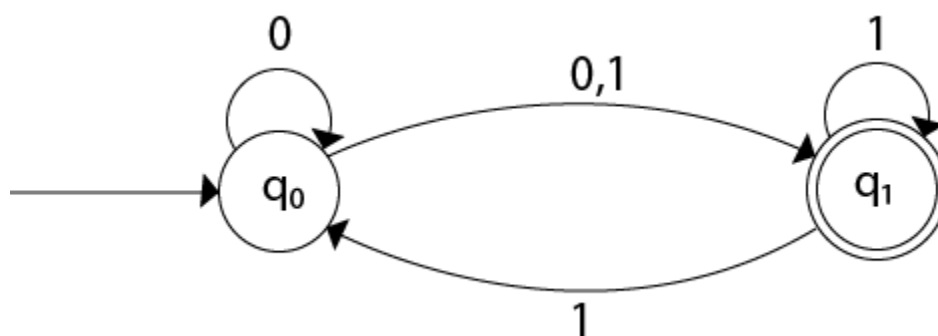
The Transition diagram will be:



The state q_2 can be eliminated because q_2 is an unreachable state.

Example 2:

Convert the given NFA to DFA.



Solution: For the given transition diagram we will first construct the transition table.

State	0	1
$\rightarrow q_0$	$\{q_0, q_1\}$	$\{q_1\}$
$*q_1$	ϕ	$\{q_0, q_1\}$

Now we will obtain δ' transition for state q_0 .

1. $\delta'([q_0], 0) = \{q_0, q_1\}$
2. $\quad = [q_0, q_1]$ (**new** state generated)
3. $\delta'([q_0], 1) = \{q_1\} = [q_1]$

The δ' transition for state q_1 is obtained as:

1. $\delta'([q_1], 0) = \phi$
2. $\delta'([q_1], 1) = [q_0, q_1]$

Now we will obtain δ' transition on $[q_0, q_1]$.

1. $\delta'([q_0, q_1], 0) = \delta(q_0, 0) \cup \delta(q_1, 0)$
2. $\quad = \{q_0, q_1\} \cup \phi$
3. $\quad = \{q_0, q_1\}$
4. $\quad = [q_0, q_1]$

Similarly,

1. $\delta'([q_0, q_1], 1) = \delta(q_0, 1) \cup \delta(q_1, 1)$
2. $\quad = \{q_1\} \cup \{q_0, q_1\}$
3. $\quad = \{q_0, q_1\}$
4. $\quad = [q_0, q_1]$

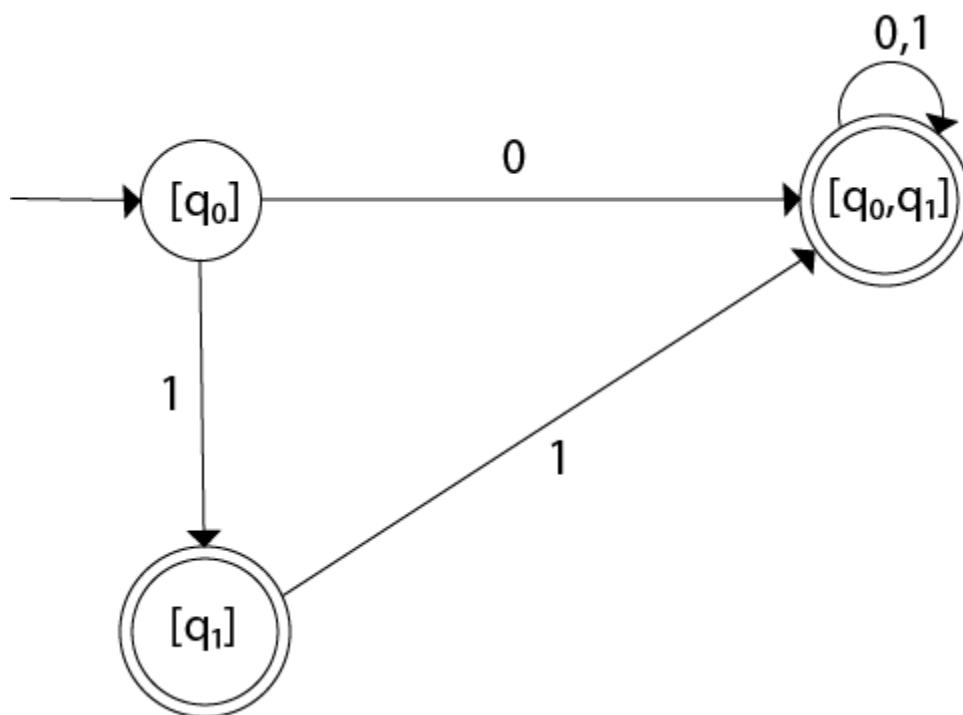
As in the given NFA, q_1 is a final state, then in DFA wherever, q_1 exists that state becomes a final state. Hence in the DFA, final states are $[q_1]$ and $[q_0, q_1]$. Therefore set of final states $F = \{[q_1], [q_0, q_1]\}$.

ADVERTISEMENT

The transition table for the constructed DFA will be:

State	0	1
$\rightarrow[q_0]$	$[q_0, q_1]$	$[q_1]$
$*[q_1]$	ϕ	$[q_0, q_1]$
$*[q_0, q_1]$	$[q_0, q_1]$	$[q_0, q_1]$

The Transition diagram will be:



Even we can change the name of the states of DFA.

Suppose

1. $A = [q_0]$
2. $B = [q_1]$
3. $C = [q_0, q_1]$

With these new names the DFA will be as follows:

