

# Machine-Level Programming III: Procedures

adapted for CS367@GMU

# Mechanisms in Procedures

## ■ Passing control

- To beginning of procedure code
- Back to return point

## ■ Passing data

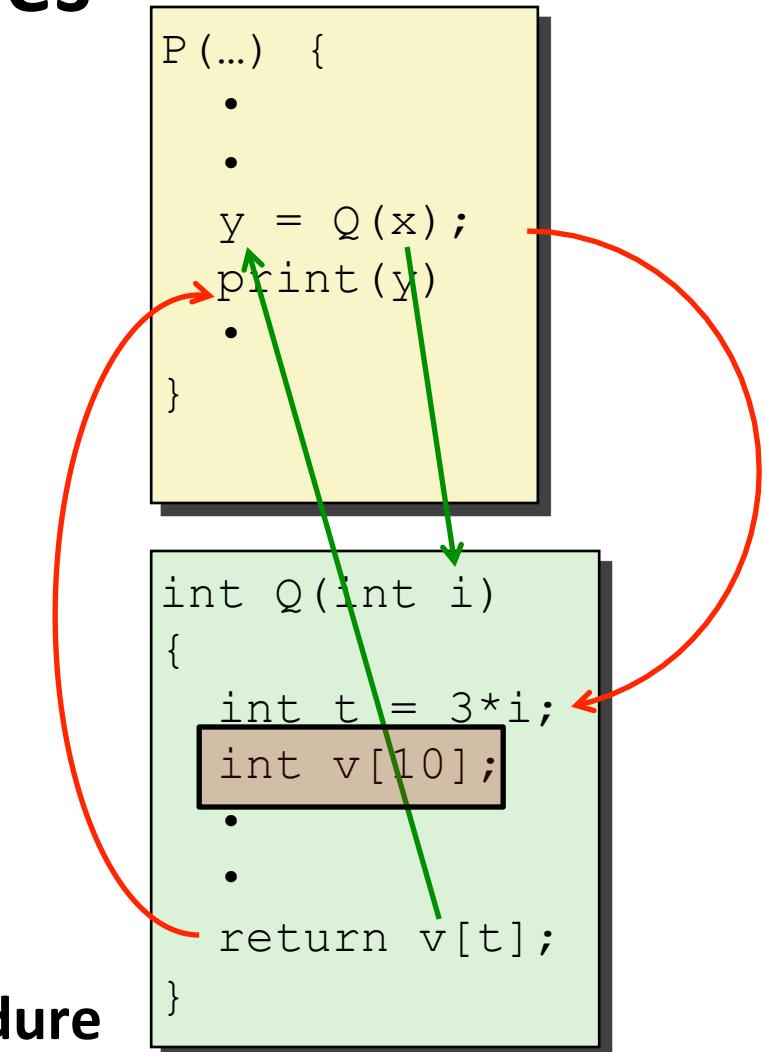
- Procedure arguments
- Return value

## ■ Memory management

- Allocate during procedure execution
- Deallocate upon return

## ■ Mechanisms all implemented with machine instructions

## ■ x86-64 implementation of a procedure uses only those mechanisms required



# Today

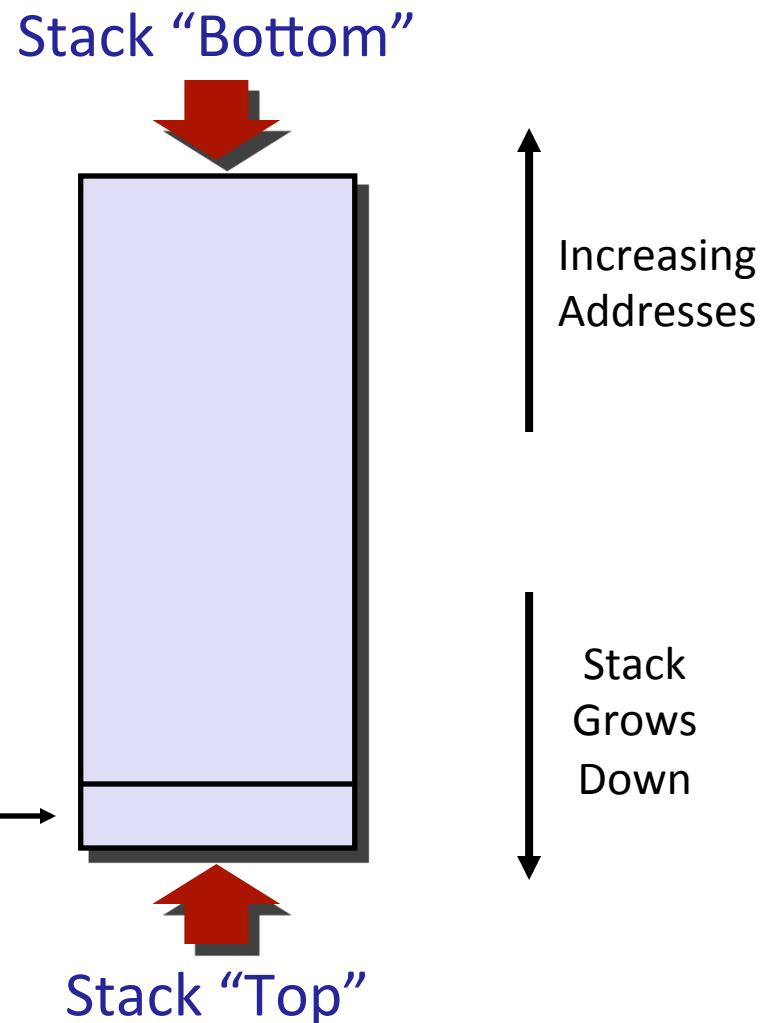
## ■ Procedures

- Stack Structure
- Calling Conventions
  - Passing control
  - Passing data
  - Managing local data
- Illustration of Recursion

# x86-64 Stack

- Region of memory managed with stack discipline
- Grows toward lower addresses
- Register `%rsp` contains lowest stack address
  - address of “top” element

Stack Pointer: `%rsp` →



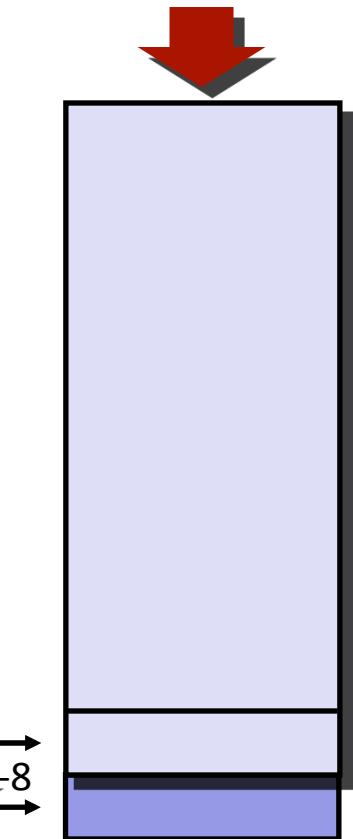
# x86-64 Stack: Push

## ■ **pushq Src**

- Fetch operand at Src
- Decrement  $\%rsp$  by 8
- Write operand at address given by  $\%rsp$

Stack Pointer:  $\%rsp$

Stack “Bottom”

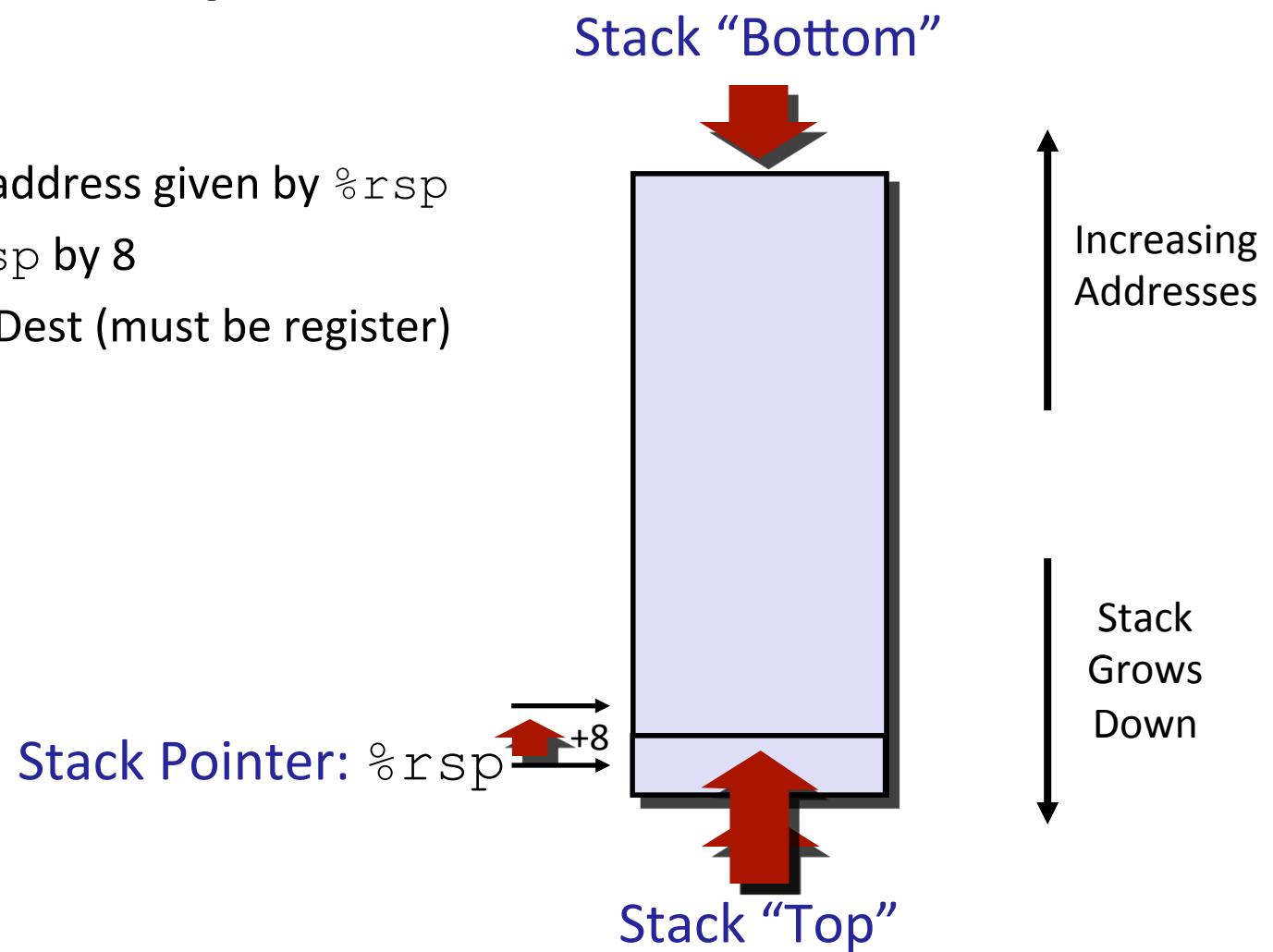


Stack “Top”

# x86-64 Stack: Pop

## ■ **popq Dest**

- Read value at address given by `%rsp`
- Increment `%rsp` by 8
- Store value at Dest (must be register)



# Today

## ■ Procedures

- Stack Structure
- Calling Conventions
  - Passing control
  - Passing data
  - Managing local data
- Illustration of Recursion

# Code Examples

```
void multstore  
    (long x, long y, long *dest)  
{  
    long t = mult2(x, y);  
    *dest = t;  
}
```

```
0000000000400540 <multstore>:  
400540: push    %rbx          # Save %rbx  
400541: mov     %rdx,%rbx    # Save dest  
400544: callq   400550 <mult2>  # mult2(x,y)  
400549: mov     %rax,(%rbx)    # Save at dest  
40054c: pop     %rbx          # Restore %rbx  
40054d: retq               # Return
```

```
long mult2  
    (long a, long b)  
{  
    long s = a * b;  
    return s;  
}
```

```
0000000000400550 <mult2>:  
400550: mov     %rdi,%rax      # a  
400553: imul    %rsi,%rax      # a * b  
400557: retq               # Return
```

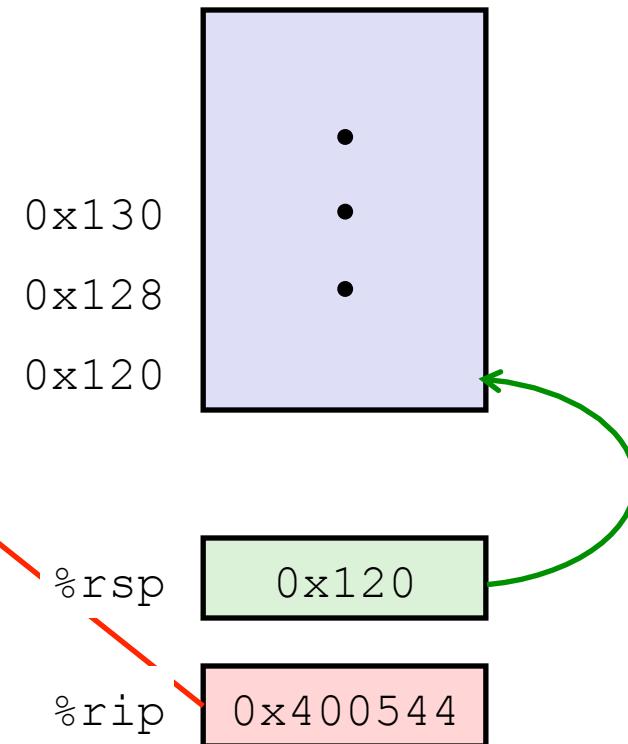
# Procedure Control Flow

- Use stack to support procedure call and return
- **Procedure call: `call` label**
  - Push return address on stack
  - Jump to label
- **Return address:**
  - Address of the next instruction right after call
  - Example from disassembly
- **Procedure return: `ret`**
  - Pop `code` address from stack
  - Jump to address, resume execution there

# Control Flow Example #1

```
000000000400540 <multstore>:  
    •  
    •  
    400544: callq  400550 <mult2>  
    400549: mov     %rax, (%rbx)  
    •  
    •
```

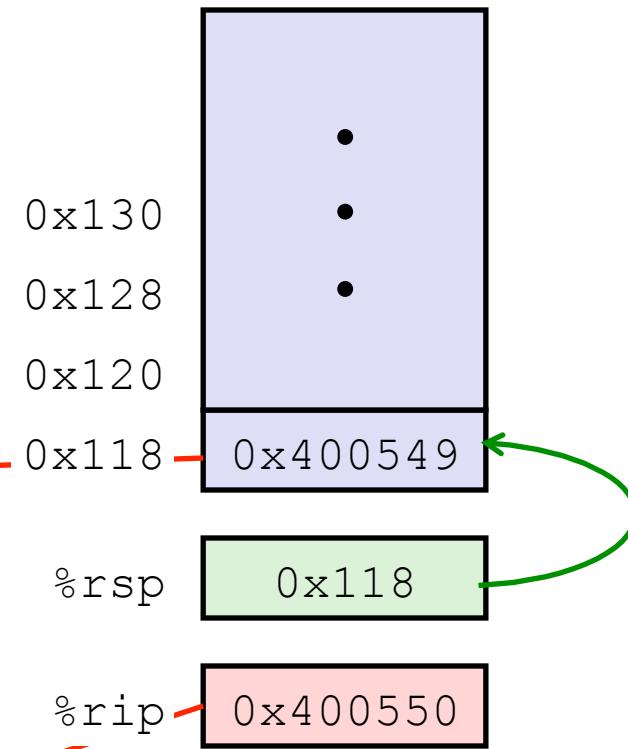
```
000000000400550 <mult2>:  
    400550: mov     %rdi,%rax  
    •  
    •  
    400557: retq
```



# Control Flow Example #2

```
000000000400540 <multstore>:  
•  
•  
400544: callq 400550 <mult2>  
400549: mov    %rax, (%rbx) ←
```

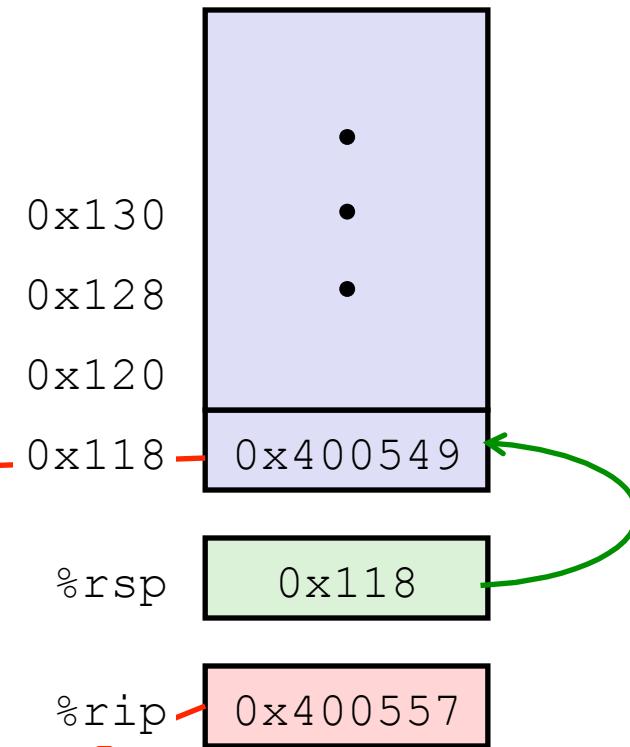
```
000000000400550 <mult2>:  
400550: mov    %rdi,%rax ←  
•  
•  
400557: retq
```



# Control Flow Example #3

```
000000000400540 <multstore>:  
    •  
    •  
    400544: callq  400550 <mult2>  
    400549: mov     %rax, (%rbx) ←
```

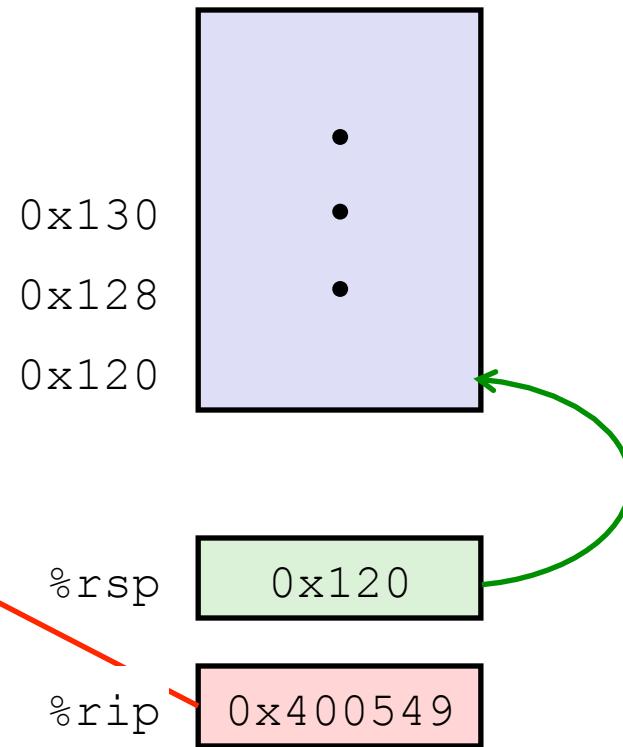
```
000000000400550 <mult2>:  
    400550: mov     %rdi,%rax  
    •  
    •  
    400557: retq ←
```



# Control Flow Example #4

```
000000000400540 <multstore>:  
    •  
    •  
    400544: callq  400550 <mult2>  
    400549: mov     %rax, (%rbx) ←
```

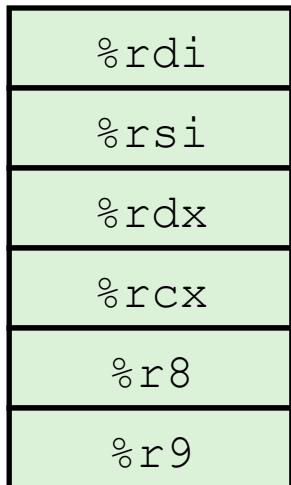
```
000000000400550 <mult2>:  
    400550: mov     %rdi,%rax  
    •  
    •  
    400557: retq
```



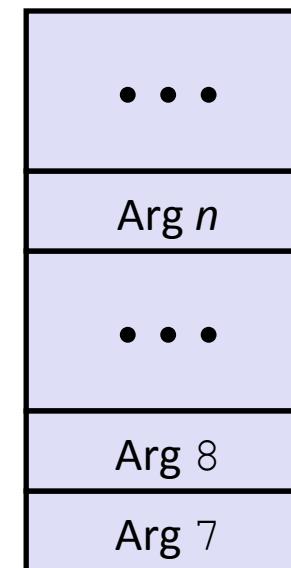
# Procedure Data Flow

## Registers

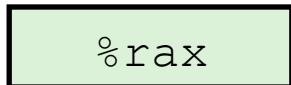
- First 6 arguments



## Stack



- Return value



- Only allocate stack space when needed

# Data Flow Examples

```
void multstore
    (long x, long y, long *dest)
{
    long t = mult2(x, y);
    *dest = t;
}
```

```
0000000000400540 <multstore>:
# x in %rdi, y in %rsi, dest in %rdx
...
400541: mov    %rdx,%rbx          # Save dest
400544: callq  400550 <mult2>    # mult2(x,y)
# t in %rax
400549: mov    %rax,(%rbx)       # Save at dest
...
```

```
long mult2
    (long a, long b)
{
    long s = a * b;
    return s;
}
```

```
0000000000400550 <mult2>:
# a in %rdi, b in %rsi
400550: mov    %rdi,%rax          # a
400553: imul   %rsi,%rax          # a * b
# s in %rax
400557: retq               # Return
```

# Stack-Based Languages

## ■ Languages that support recursion

- e.g., C, Pascal, Java
- Code must be “Reentrant”
  - Multiple simultaneous instantiations of single procedure
- Need some place to store state of each instantiation
  - Arguments
  - Local variables
  - Return pointer

## ■ Stack discipline

- State for given procedure needed for limited time
  - From when it's called to when it returns
- Callee returns before caller does (strictly nested calls)

## ■ Stack allocated in **Frames**

- state for single procedure instantiation

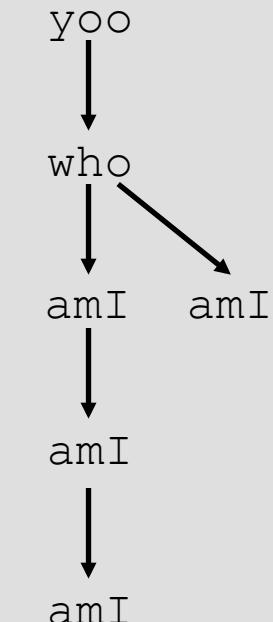
# Call Chain Example

```
yoo (...)  
{  
    •  
    •  
    who () ;  
    •  
    •  
}
```

```
who (...)  
{  
    • • •  
    amI () ;  
    • • •  
    amI () ;  
    • • •  
}
```

```
amI (...)  
{  
    •  
    •  
    amI () ;  
    •  
    •  
}
```

Example  
Call Chain

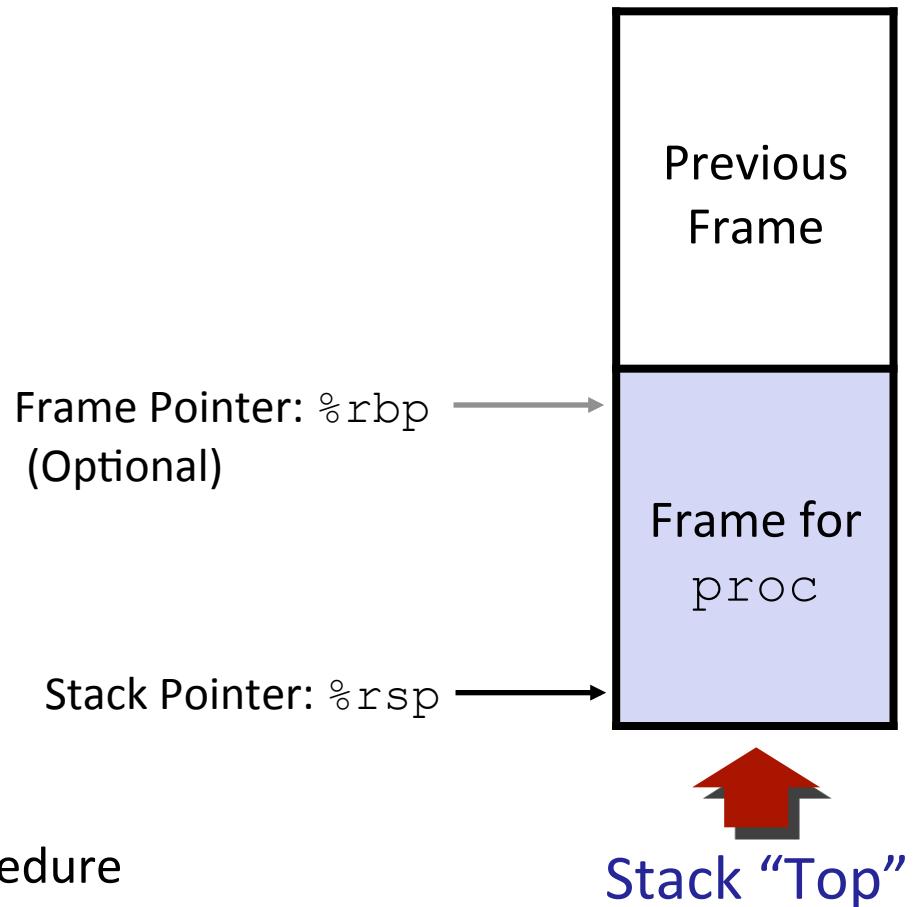


Procedure `amI ()` is recursive

# Stack Frames

## ■ Contents

- Return information
- Local storage (if needed)
- Temporary space (if needed)

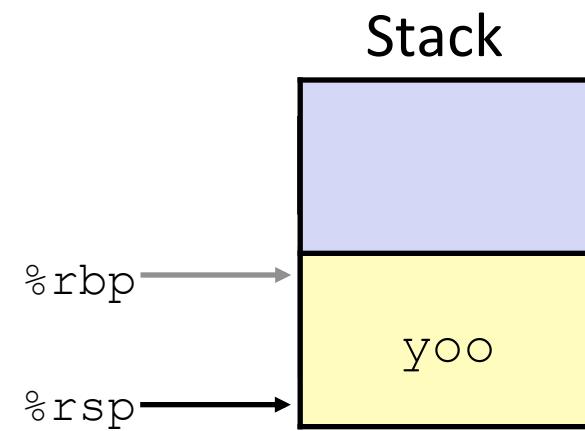
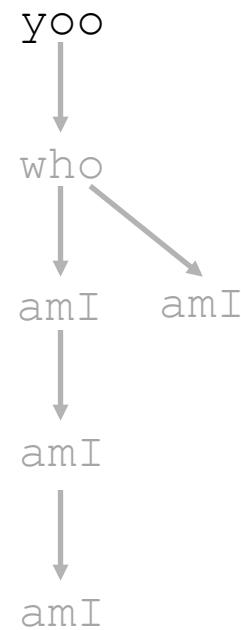


## ■ Management

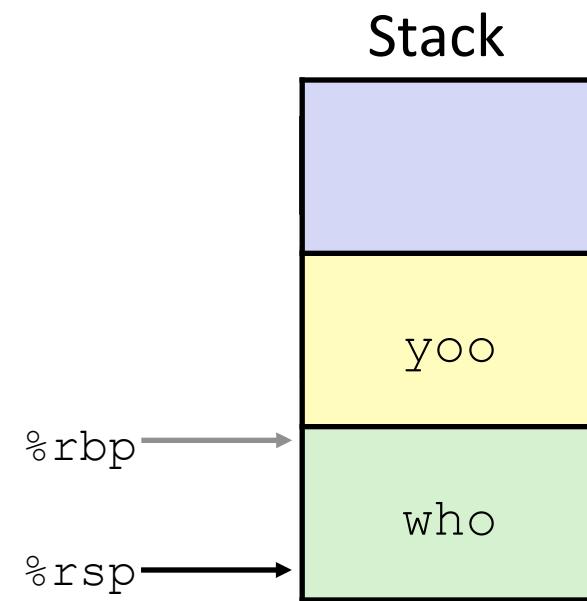
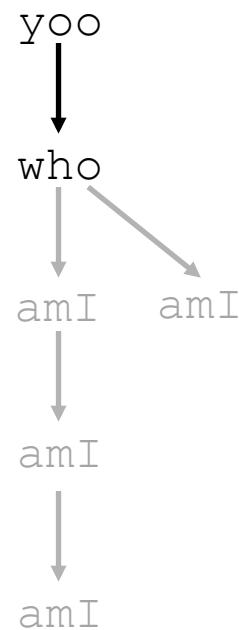
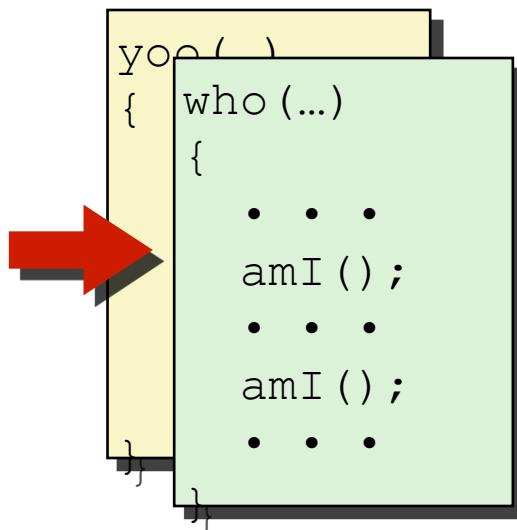
- Space allocated when enter procedure
  - “Set-up” code
  - Includes push by **call** instruction
- Deallocated when return
  - “Finish” code
  - Includes pop by **ret** instruction

# Example

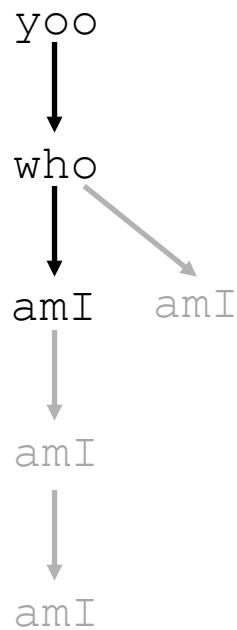
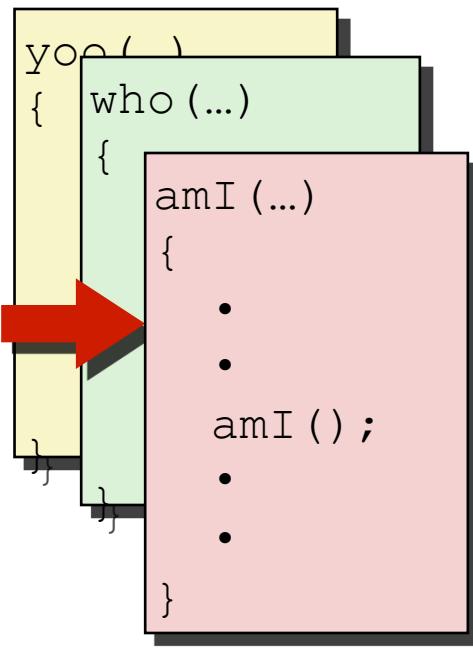
```
    yoo (...)  
    {  
        •  
        •  
        who () ;  
        •  
        •  
    }
```



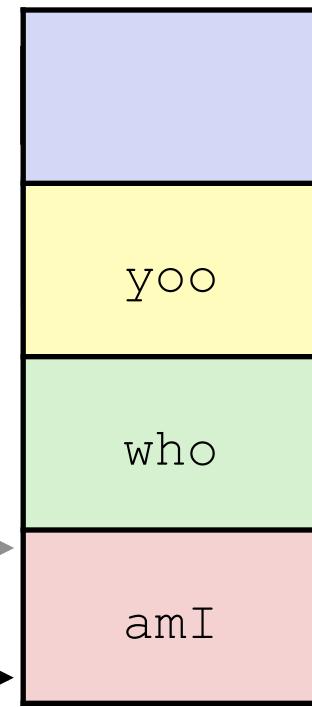
# Example



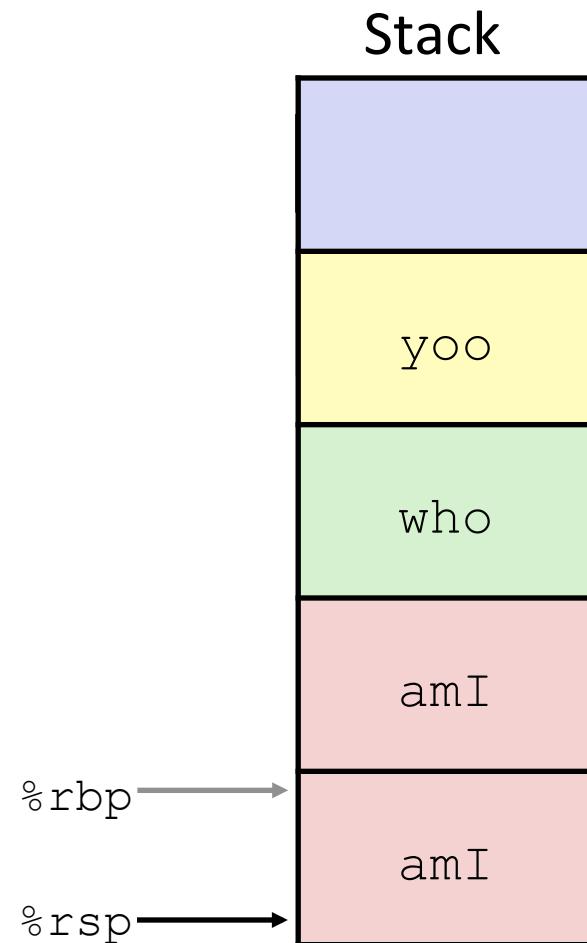
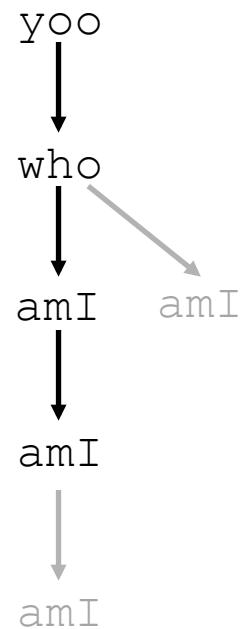
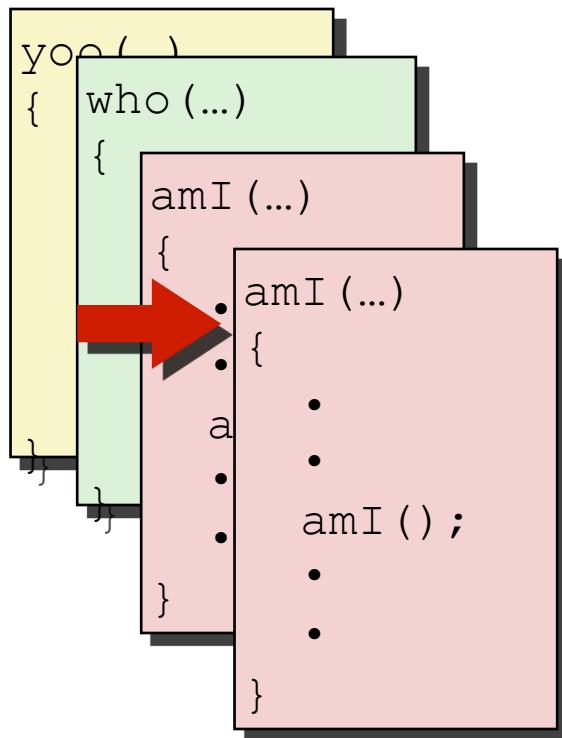
# Example



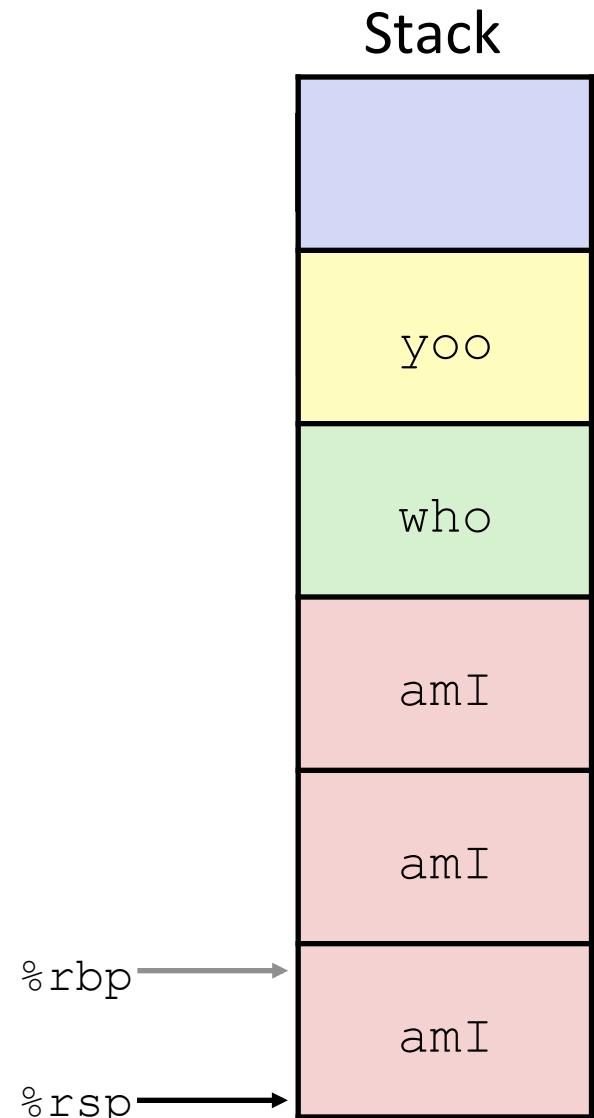
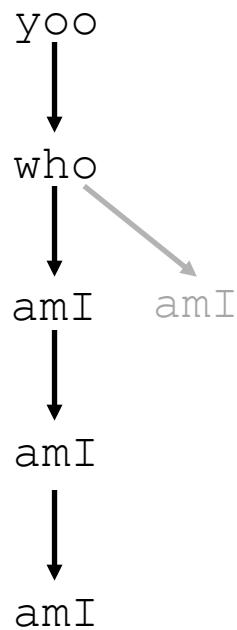
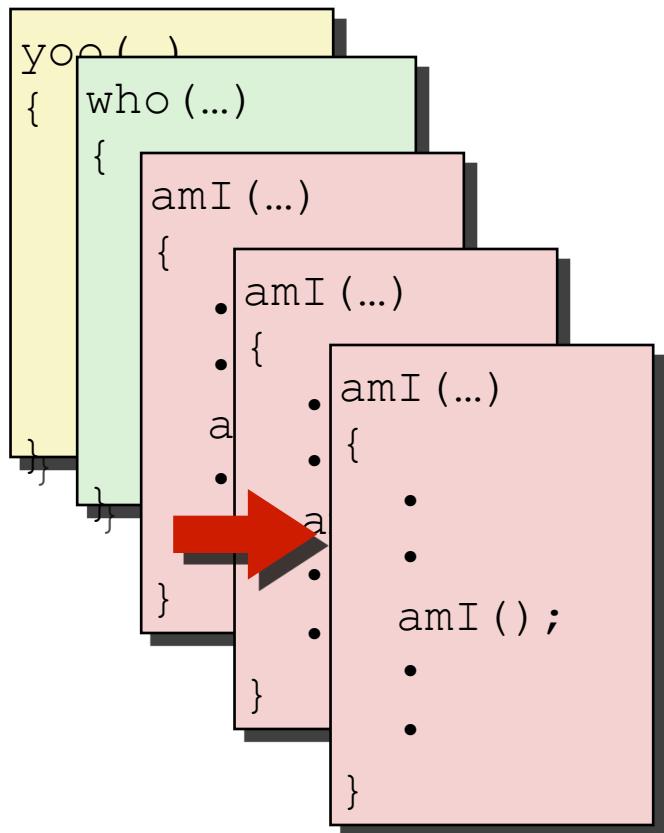
Stack



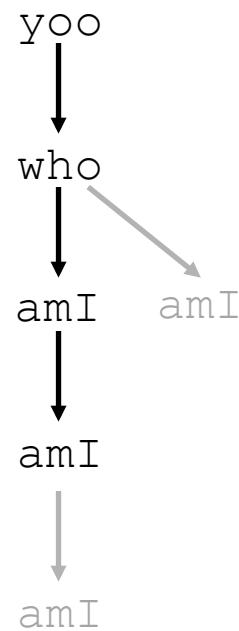
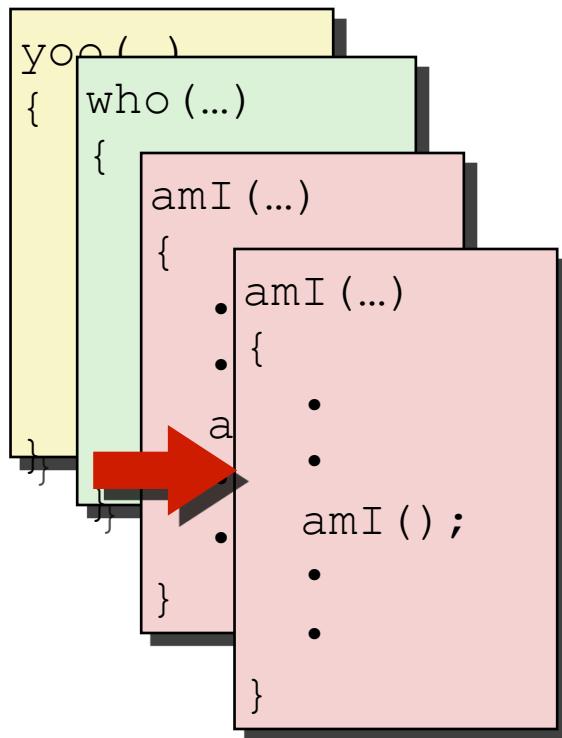
# Example



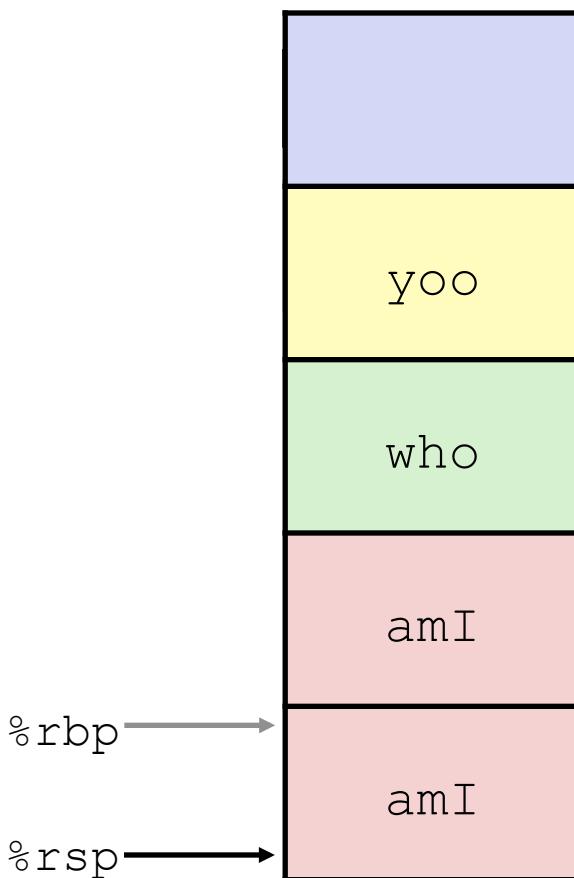
# Example



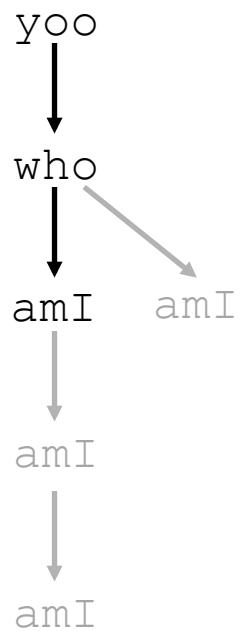
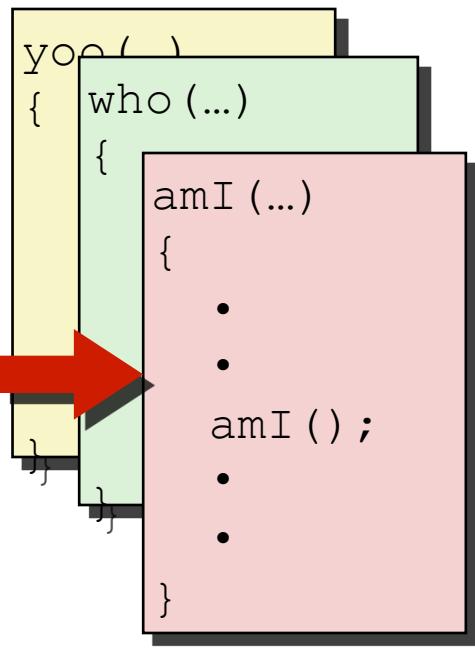
# Example



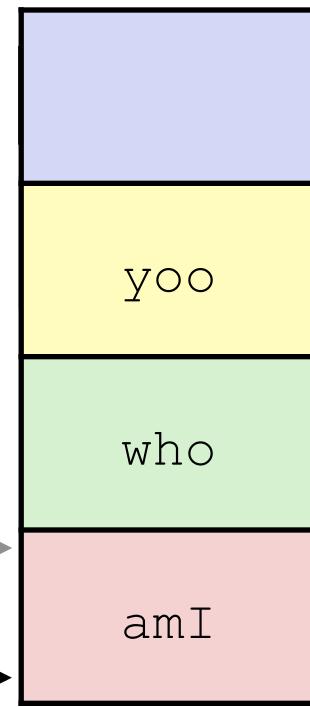
Stack



# Example



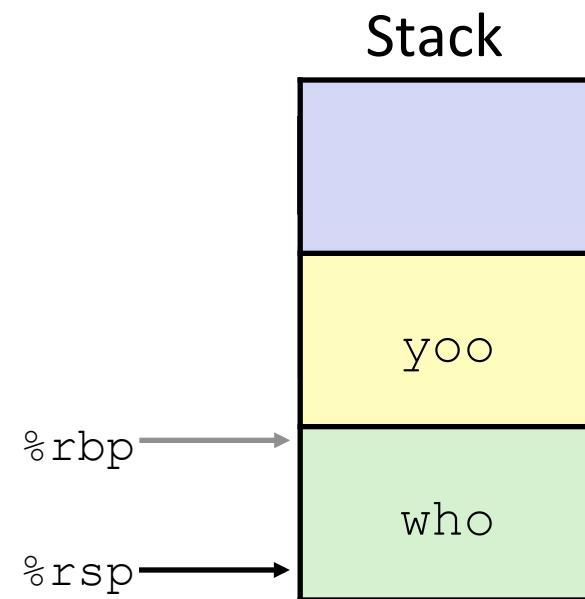
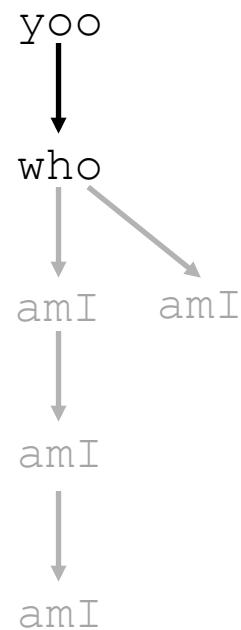
Stack



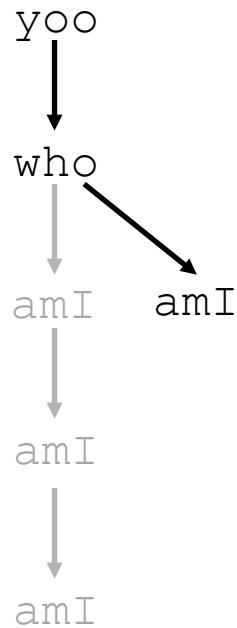
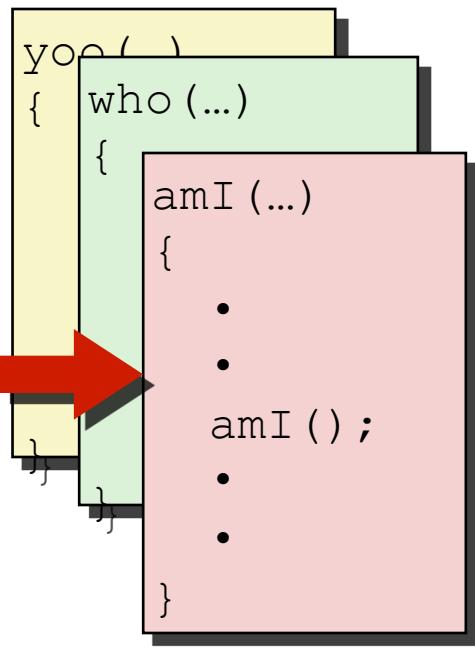
# Example

```
yoo()
{
    who(...)

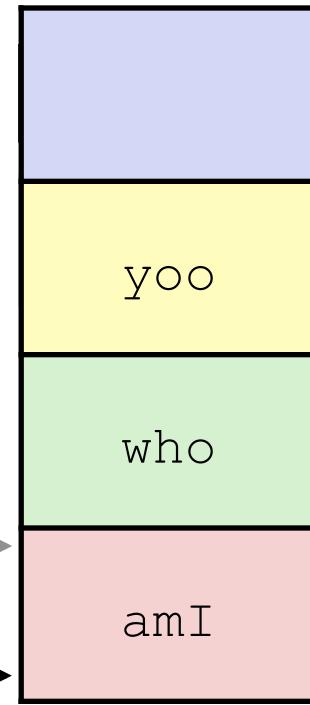
    {
        . . .
        amI();
        . . .
        amI();
        . . .
    }
}
```



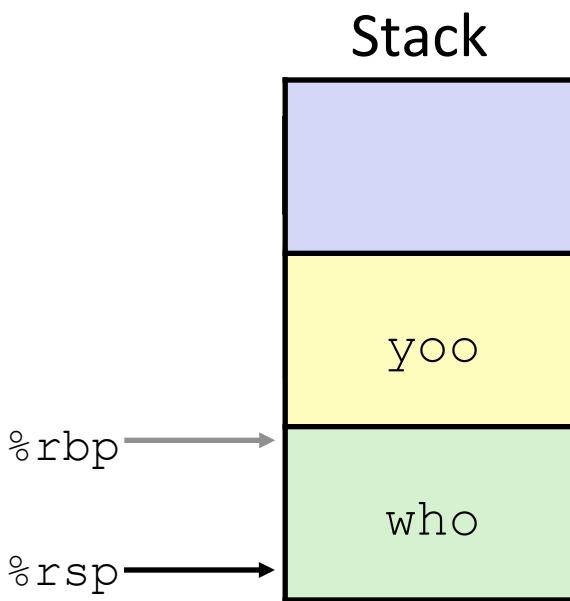
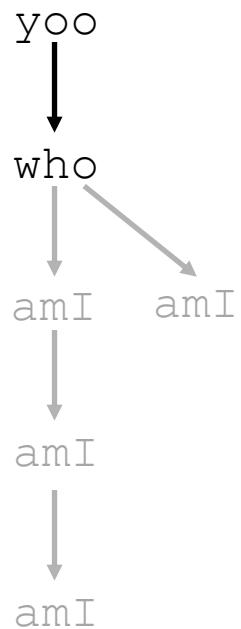
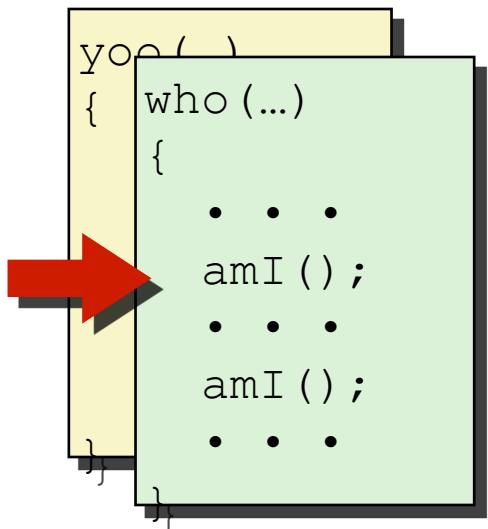
# Example



Stack

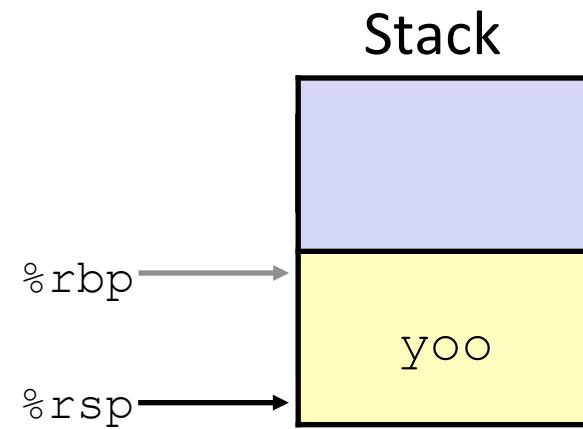
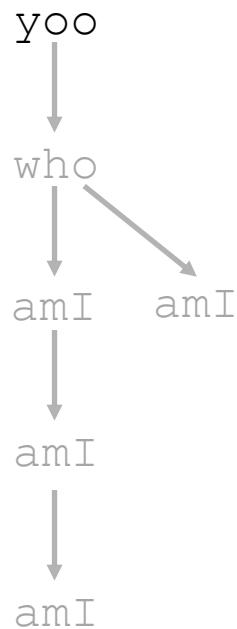


# Example



# Example

```
yoo (...)  
{  
    •  
    •  
    who () ;  
    •  
    •  
}  
}
```

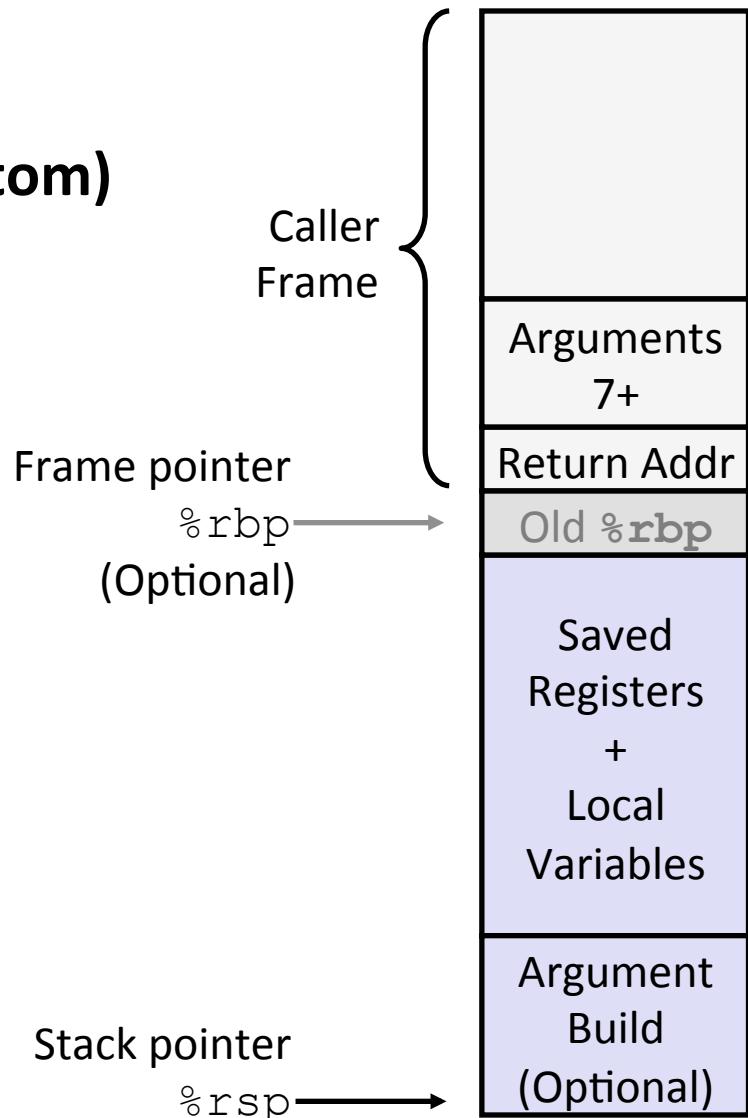


# x86-64/Linux Stack Frame

## Scenario: Caller called a function

### ■ Current Stack Frame (“Top” to Bottom)

- “Argument build:”  
Parameters for function about to call
- Local variables  
If can’t keep in registers
- Saved register context
- Old frame pointer (optional)



### ■ Caller Stack Frame

- Return address
  - Pushed by `call` instruction
- Arguments for this call