

# Machine-Level Programming III: Procedures

**adapted for CS367@GMU**

# Mechanisms in Procedures

## ■ Passing control

- To beginning of procedure code
- Back to return point

## ■ Passing data

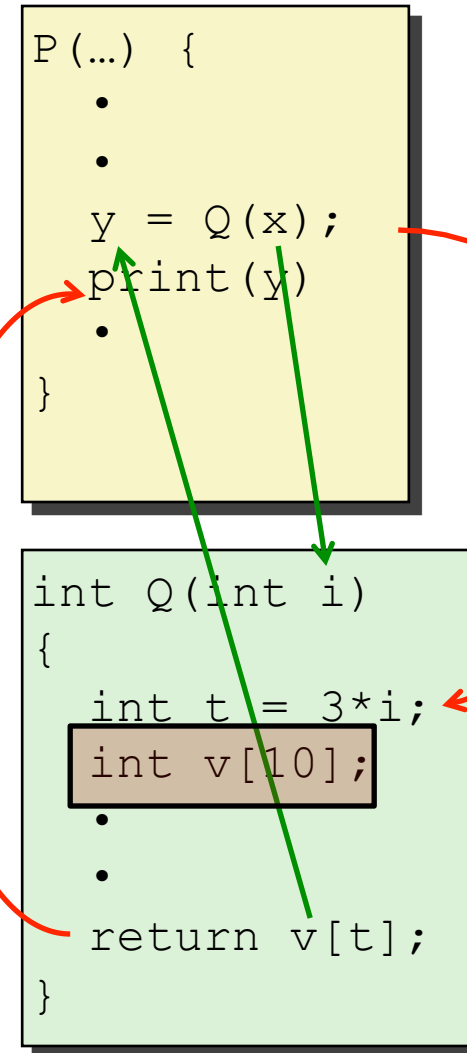
- Procedure arguments
- Return value

## ■ Memory management

- Allocate during procedure execution
- Deallocate upon return

## ■ Mechanisms all implemented with machine instructions

## ■ x86-64 implementation of a procedure uses only those mechanisms required



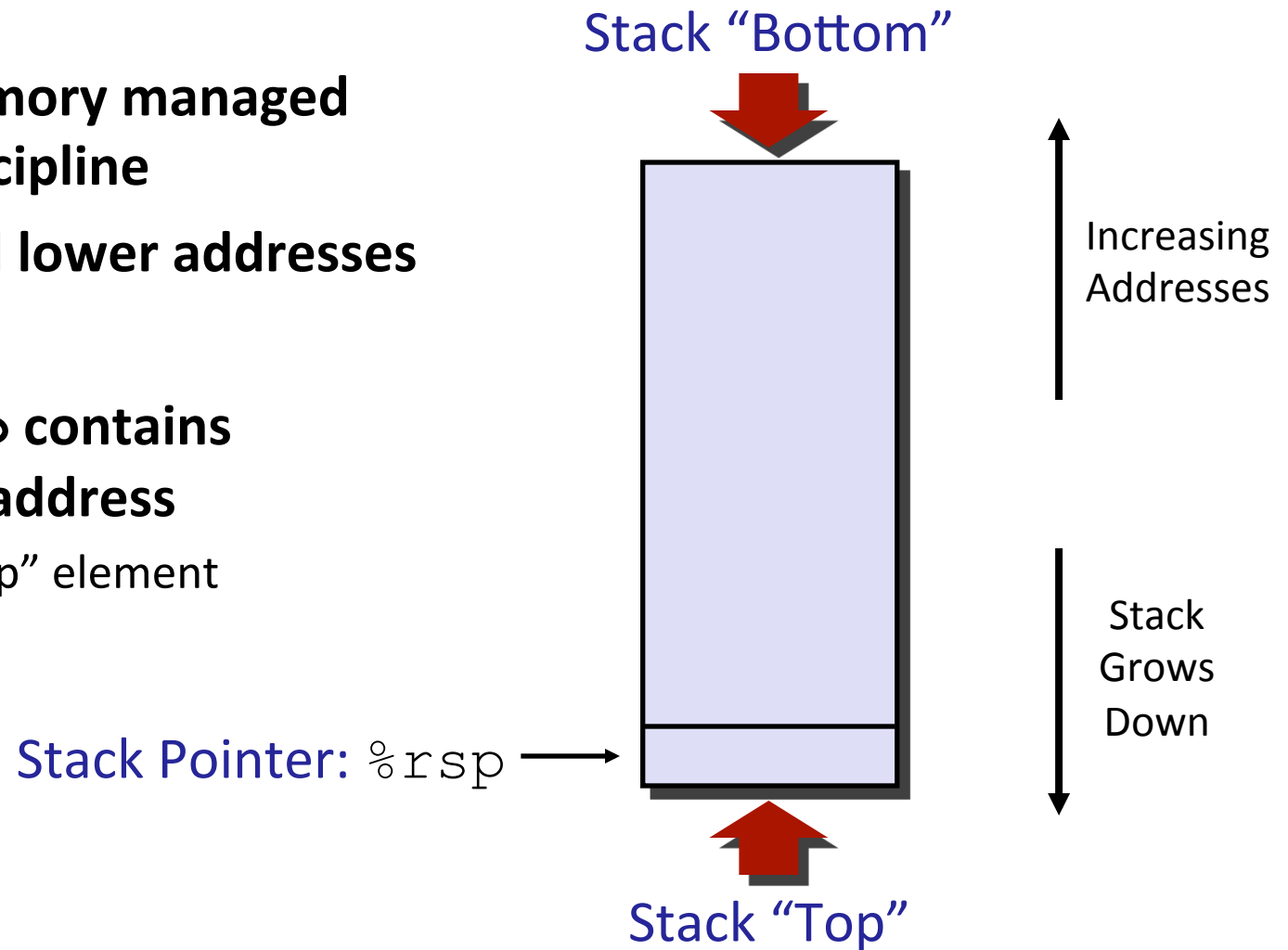
# Today

## ■ Procedures

- **Stack Structure**
- **Calling Conventions**
  - Passing control
  - Passing data
  - Managing local data
- **Illustration of Recursion**

# x86-64 Stack

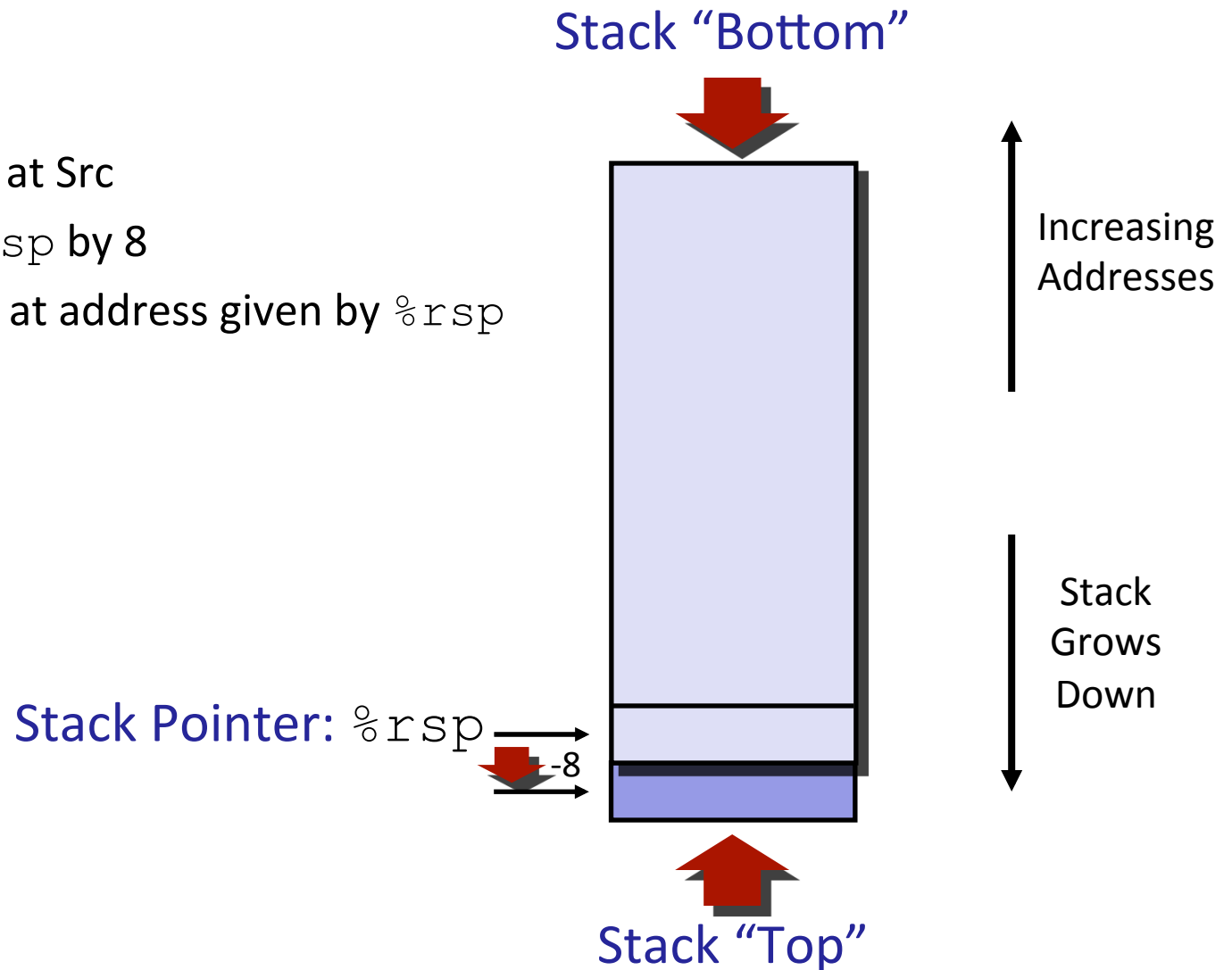
- Region of memory managed with stack discipline
- Grows toward lower addresses
- Register `%rsp` contains lowest stack address
  - address of “top” element



# x86-64 Stack: Push

## ■ `pushq Src`

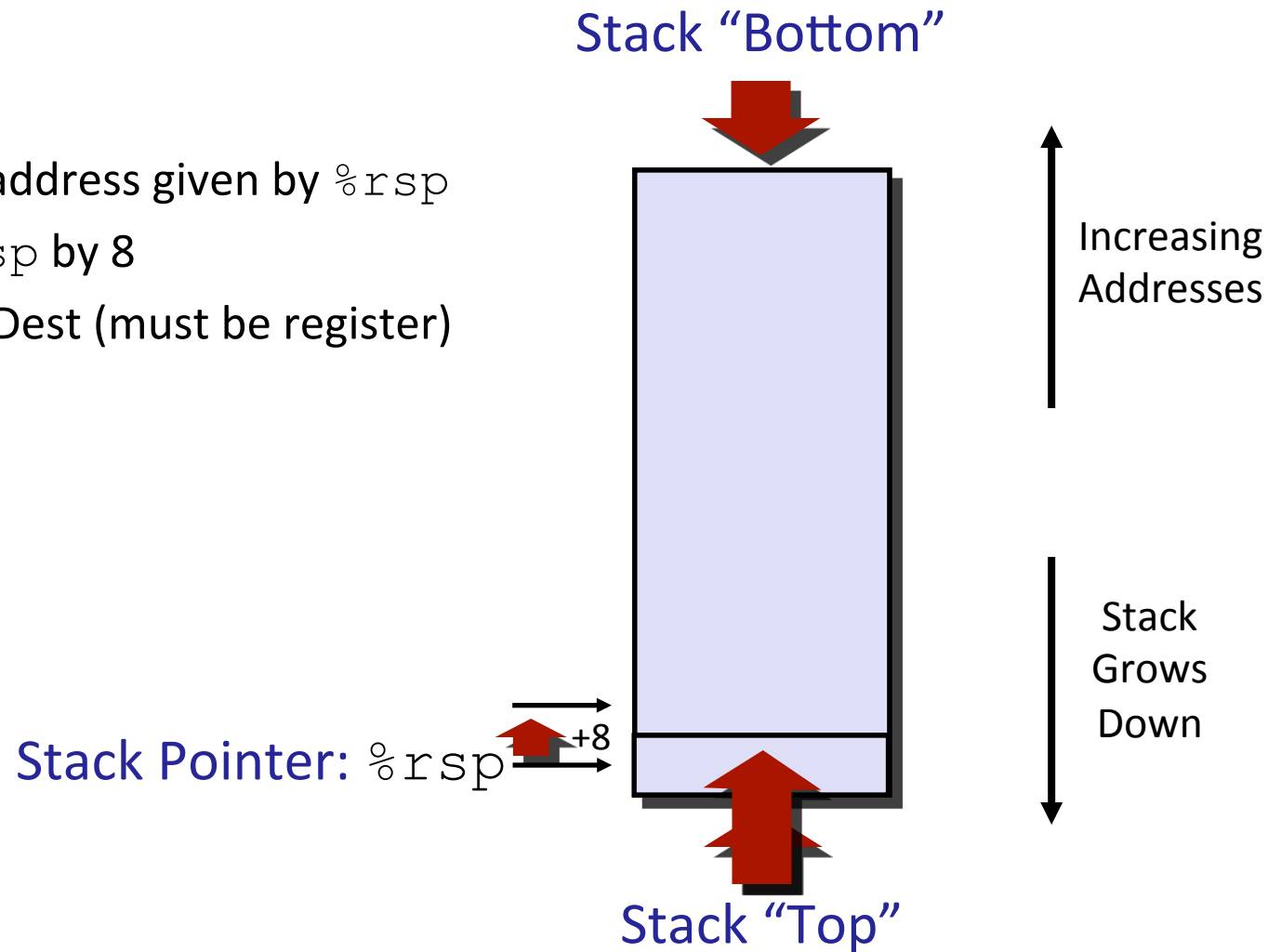
- Fetch operand at `Src`
- Decrement `%rsp` by 8
- Write operand at address given by `%rsp`



# x86-64 Stack: Pop

## ■ `popq Dest`

- Read value at address given by `%rsp`
- Increment `%rsp` by 8
- Store value at `Dest` (must be register)



# Today

## ■ Procedures

- Stack Structure
- Calling Conventions
  - **Passing control**
  - Passing data
  - Managing local data
- Illustration of Recursion

# Code Examples

```
void multstore
(long x, long y, long *dest)
{
    long t = mult2(x, y);
    *dest = t;
}
```

```
00000000000400540 <multstore>:
400540: push    %rbx                # Save %rbx
400541: mov     %rdx,%rbx           # Save dest
400544: callq   400550 <mult2>      # mult2(x,y)
400549: mov     %rax, (%rbx)         # Save at dest
40054c: pop     %rbx                # Restore %rbx
40054d: retq                      # Return
```

```
long mult2
(long a, long b)
{
    long s = a * b;
    return s;
}
```

```
00000000000400550 <mult2>:
400550: mov     %rdi,%rax           # a
400553: imul    %rsi,%rax           # a * b
400557: retq                      # Return
```



# Procedure Control Flow

- Use stack to support procedure call and return
- **Procedure call:** `call label`
  - Push return address on stack
  - Jump to label
- **Return address:**
  - Address of the next instruction right after call
  - Example from disassembly
- **Procedure return:** `ret`
  - Pop **code** address from stack
  - Jump to address, resume execution there

# Control Flow Example #1

```
00000000000400540 <multstore>:  
.  
.  
400544: callq 400550 <mult2>  
400549: mov  %rax, (%rbx)  
.  
.
```

```
00000000000400550 <mult2>:  
400550: mov  %rdi,%rax  
.  
.  
400557: retq
```

0x130

0x128

0x120

%rsp

%rip

0x120

0x400544

# Control Flow Example #2

```
00000000000400540 <multstore>:
```

•  
•  
•

```
400544: callq 400550 <mult2>
```

```
400549: mov    %rax, (%rbx) ←
```

•  
•

```
00000000000400550 <mult2>:
```

```
400550: mov    %rdi,%rax ←
```

•  
•

```
400557: retq
```

0x130

0x128

0x120

0x118

%rsp

%rip

0x400549

0x118

0x400550

# Control Flow Example #3

```
00000000000400540 <multstore>:
```

•  
•  
•  
•  
•

```
400544: callq 400550 <mult2>
```

```
400549: mov    %rax, (%rbx) ←
```

```
00000000000400550 <mult2>:
```

```
400550: mov    %rdi, %rax
```

•  
•

```
400557: retq ←
```

0x130

0x128

0x120

0x118

%rsp

%rip

0x400549

0x118

0x400557

# Control Flow Example #4

```
00000000000400540 <multstore>:  
.  
.  
400544: callq 400550 <mult2>  
400549: mov  %rax, (%rbx)  
.  
.
```

```
00000000000400550 <mult2>:  
400550: mov  %rdi,%rax  
.  
.  
400557: retq
```

0x130

0x128

0x120

%rsp

%rip

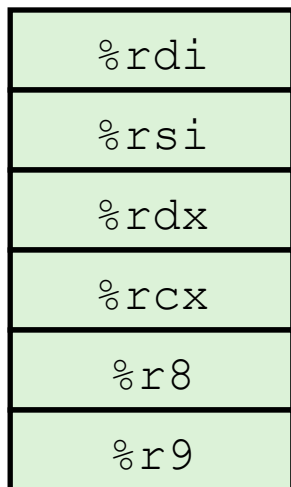
0x120

0x400549

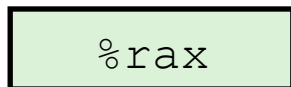
# Procedure Data Flow

## Registers

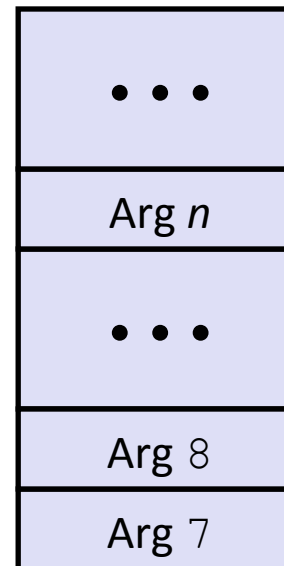
### ■ First 6 arguments



### ■ Return value



## Stack



### ■ Only allocate stack space when needed

# Data Flow Examples

```
void multstore
(long x, long y, long *dest)
{
    long t = mult2(x, y);
    *dest = t;
}
```

```
0000000000400540 <multstore>:
    # x in %rdi, y in %rsi, dest in %rdx
    . . .
400541: mov     %rdx,%rbx        # Save dest
400544: callq   400550 <mult2>    # mult2(x,y)
    # t in %rax
400549: mov     %rax, (%rbx)      # Save at dest
    . . .
```

```
long mult2
(long a, long b)
{
    long s = a * b;
    return s;
}
```

```
0000000000400550 <mult2>:
    # a in %rdi, b in %rsi
400550: mov     %rdi,%rax        # a
400553: imul    %rsi,%rax        # a * b
    # s in %rax
400557: retq                      # Return
```

# Stack-Based Languages

## ■ Languages that support recursion

- e.g., C, Pascal, Java
- Code must be “Reentrant”
  - Multiple simultaneous instantiations of single procedure
- Need some place to store state of each instantiation
  - Arguments
  - Local variables
  - Return pointer

## ■ Stack discipline

- State for given procedure needed for limited time
  - From when it's called to when it returns
- Callee returns before caller does (strictly nested calls)

## ■ Stack allocated in **Frames**

- state for single procedure instantiation



# Call Chain Example

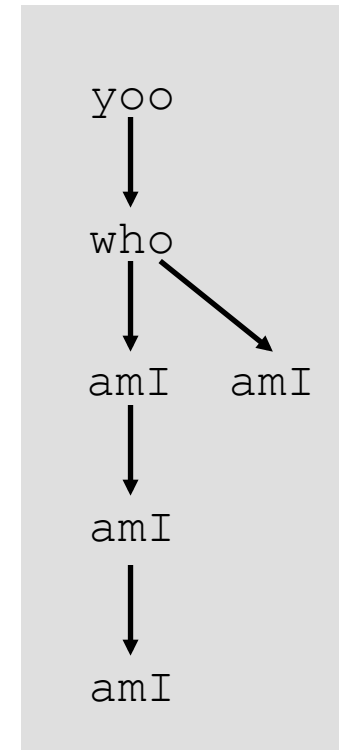
```
yoo (...)  
{  
  .  
  .  
  who ();  
  .  
  .  
}
```

```
who (...)  
{  
  . . .  
  amI ();  
  . . .  
  amI ();  
  . . .  
}
```

```
amI (...)  
{  
  .  
  .  
  amI ();  
  .  
  .  
}
```

Procedure `amI ()` is recursive

Example  
Call Chain



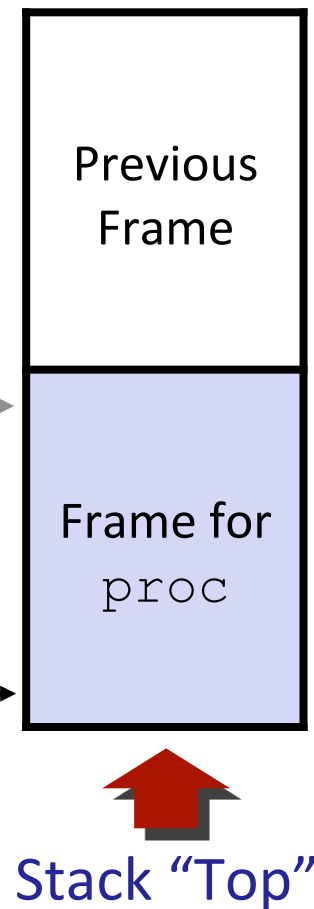
# Stack Frames

## ■ Contents

- Return information
- Local storage (if needed)
- Temporary space (if needed)

Frame Pointer: `%rbp`  
(Optional)

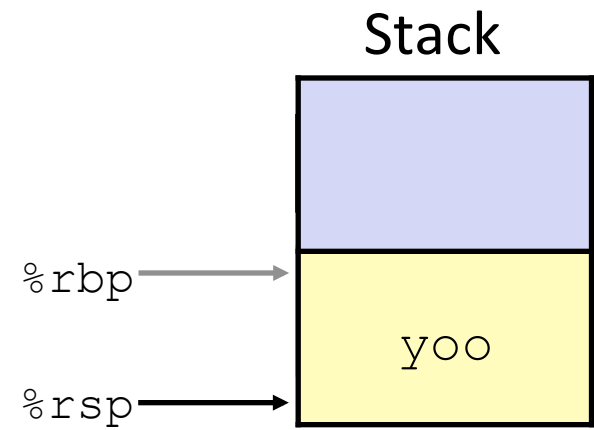
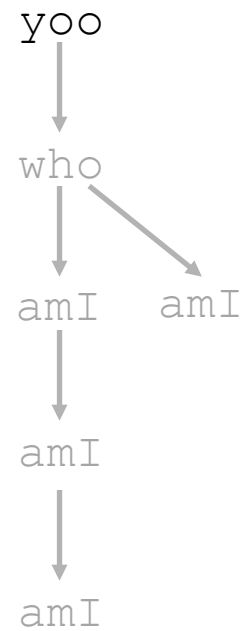
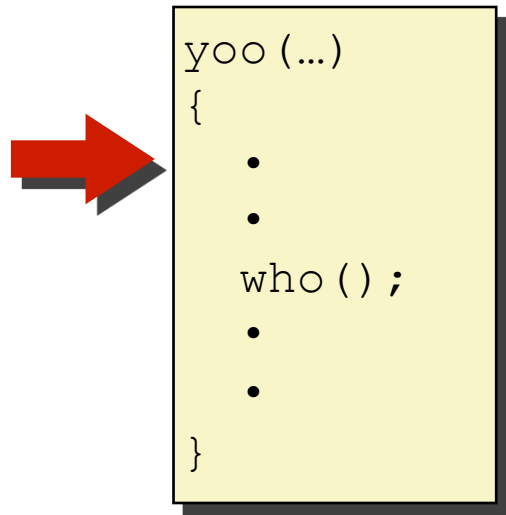
Stack Pointer: `%rsp`



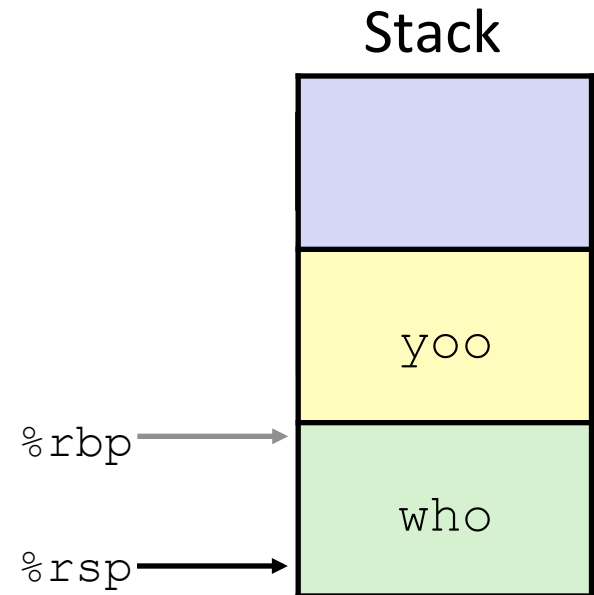
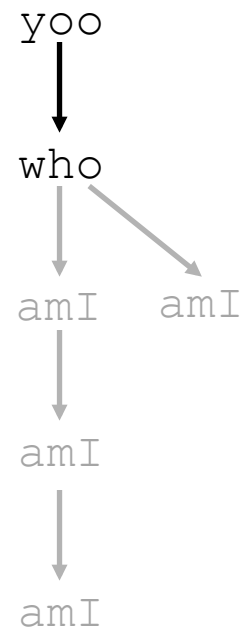
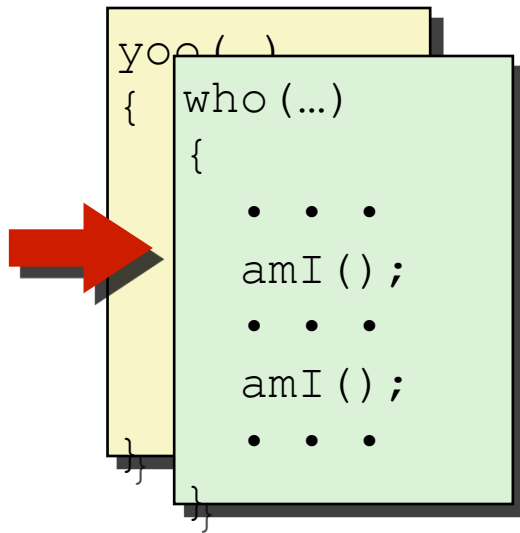
## ■ Management

- Space allocated when enter procedure
  - “Set-up” code
  - Includes push by **call** instruction
- Deallocated when return
  - “Finish” code
  - Includes pop by **ret** instruction

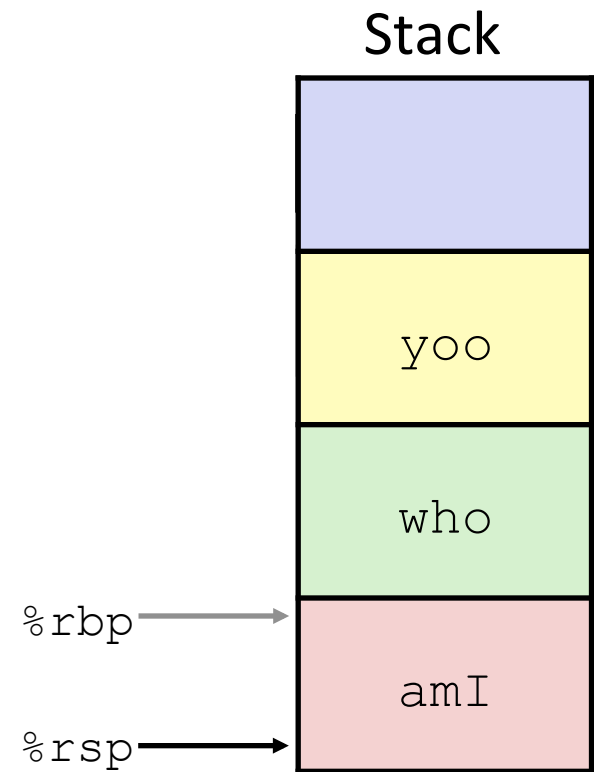
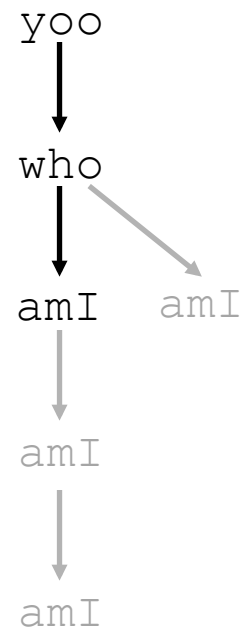
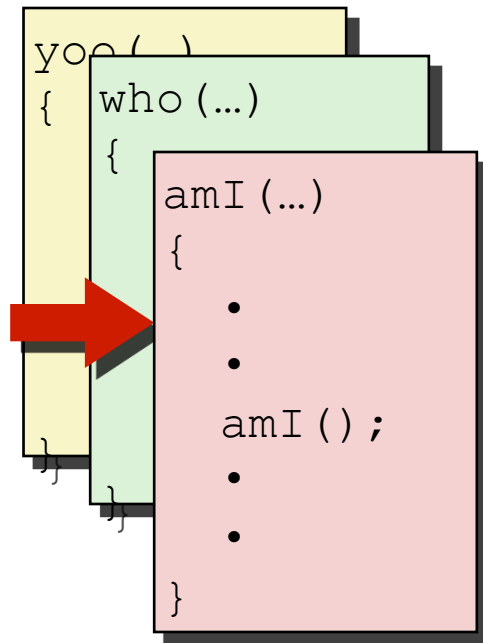
# Example



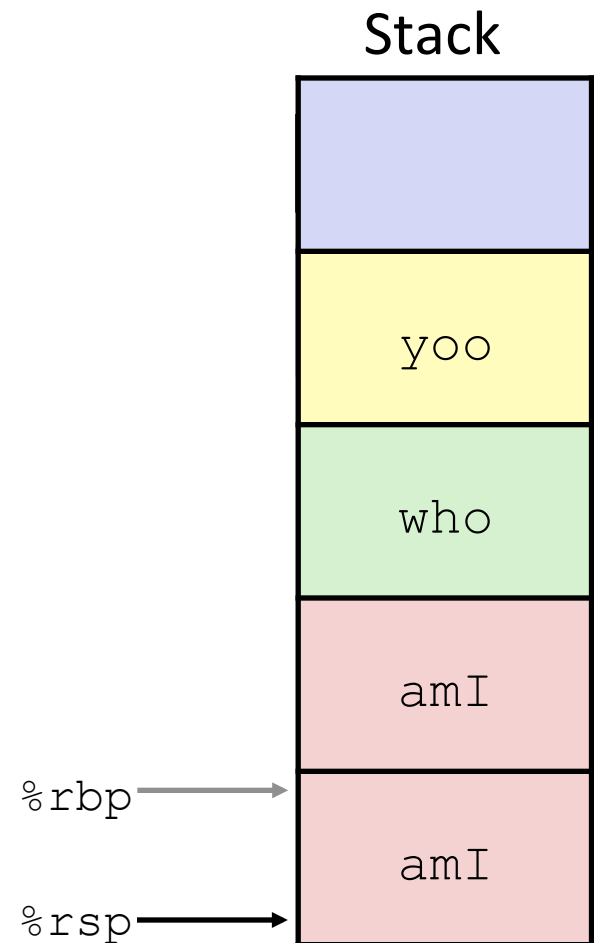
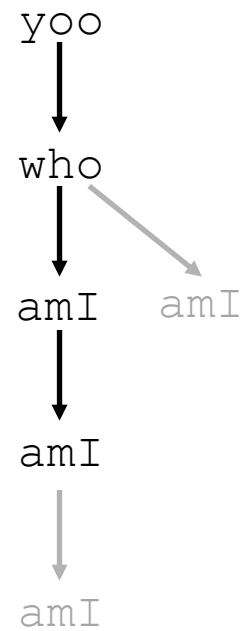
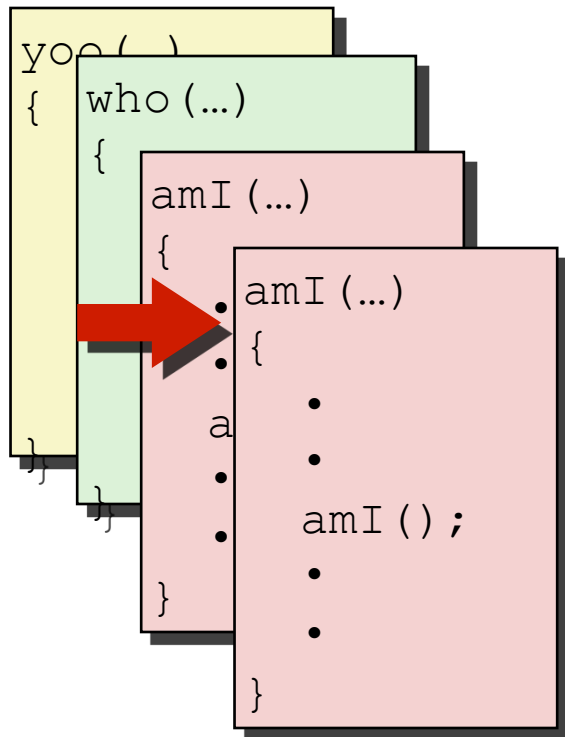
# Example



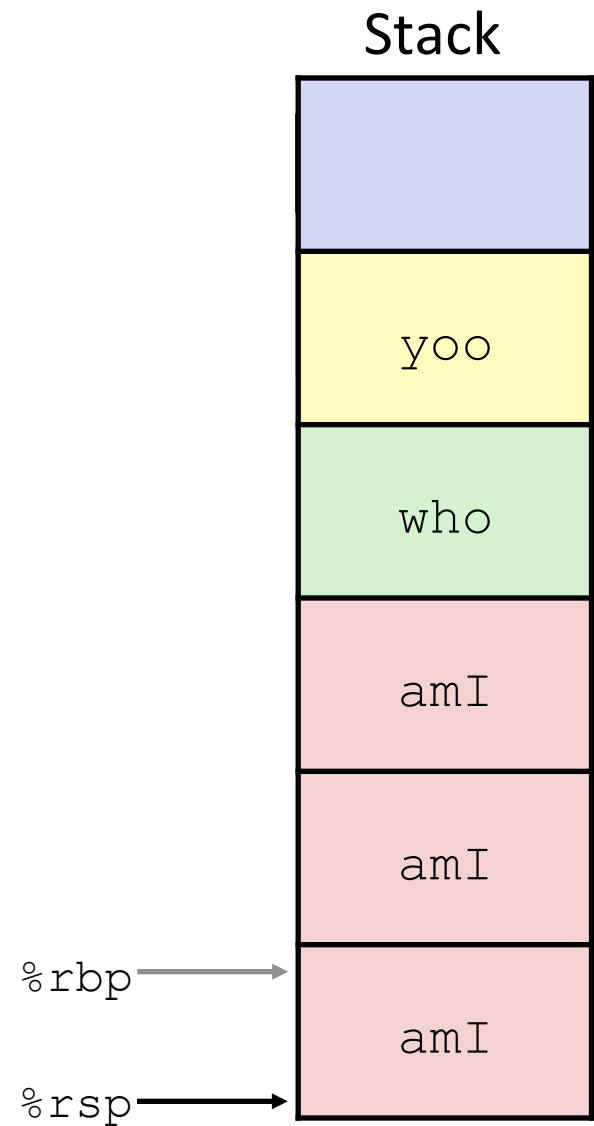
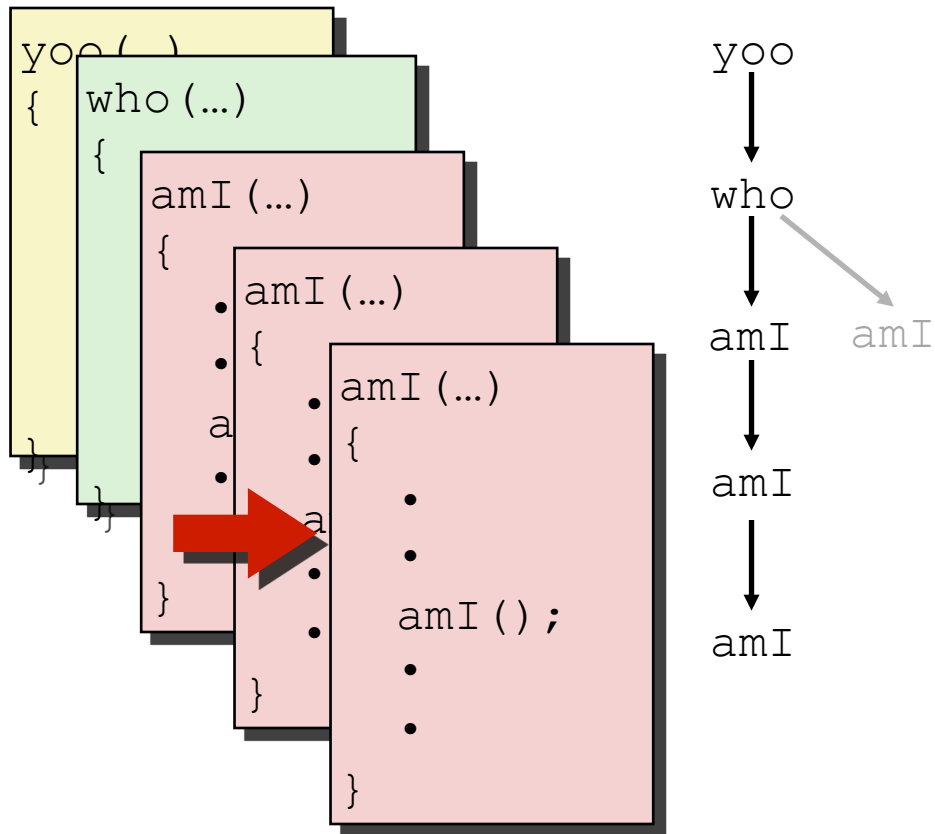
# Example



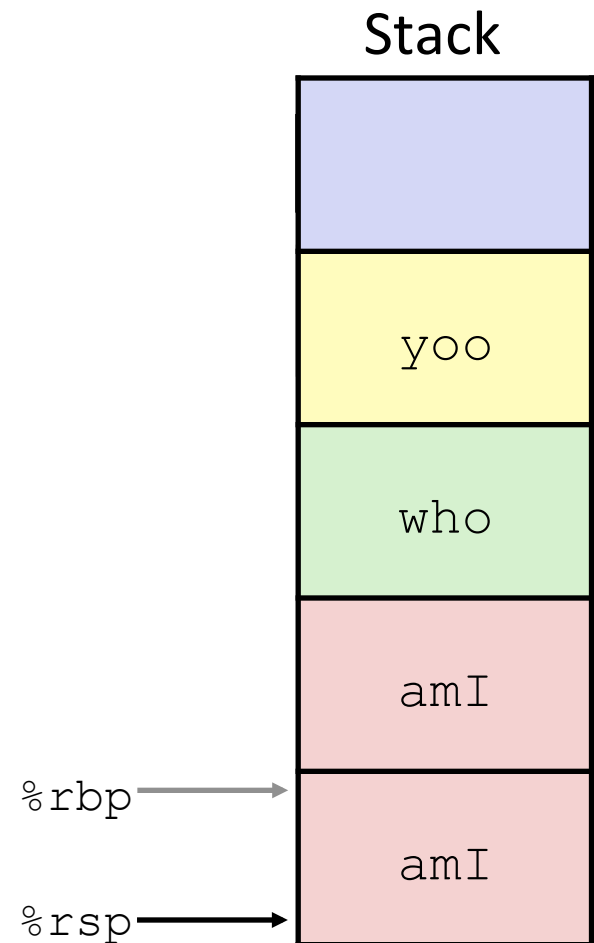
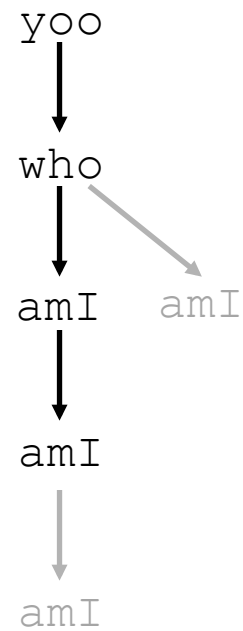
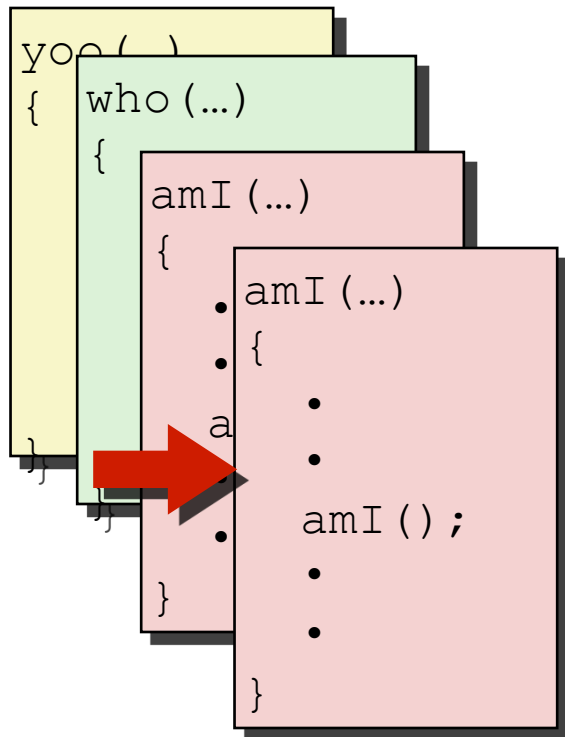
# Example



# Example

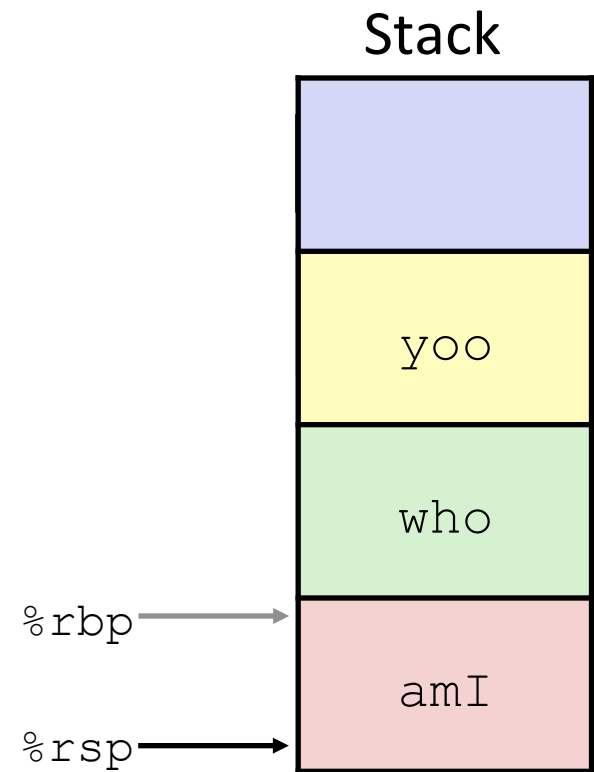
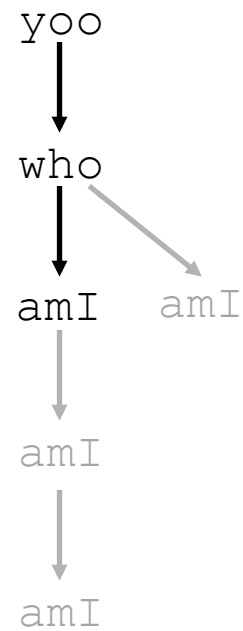
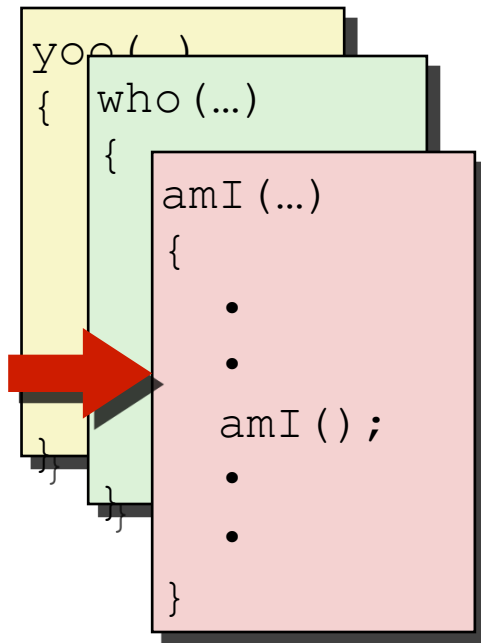


# Example

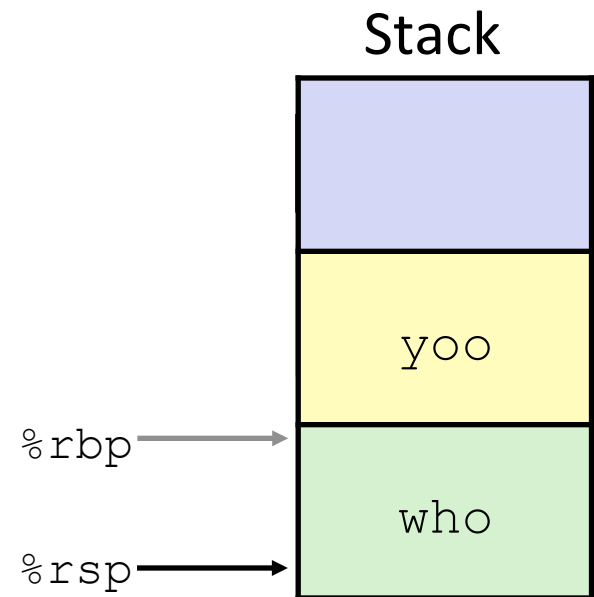
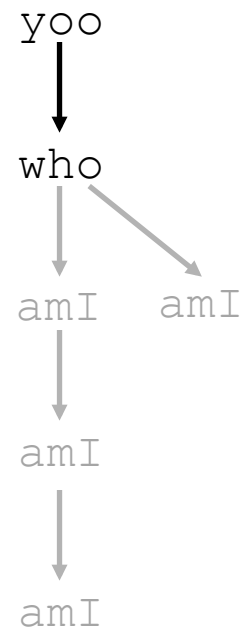
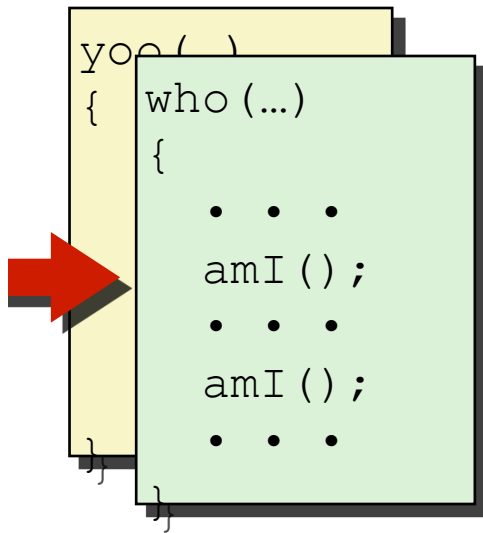




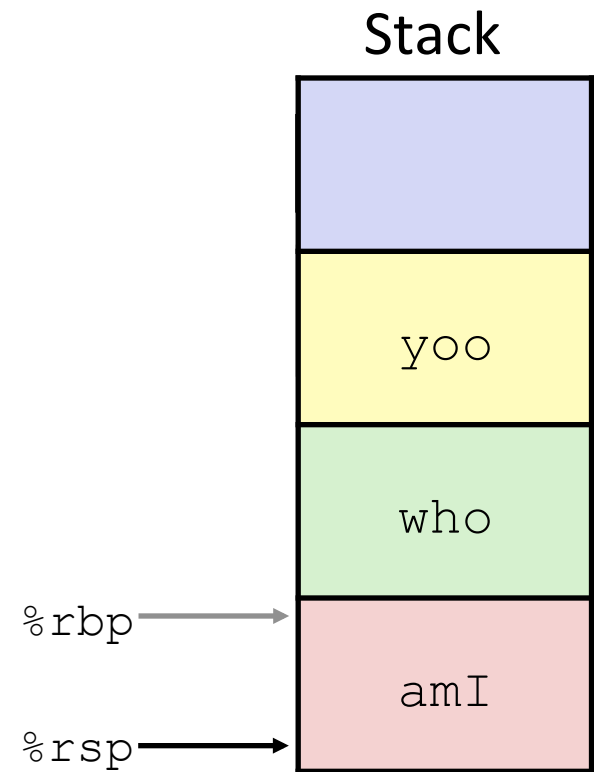
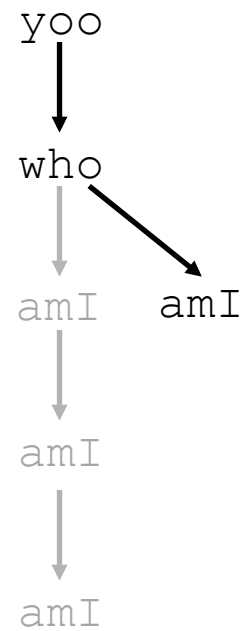
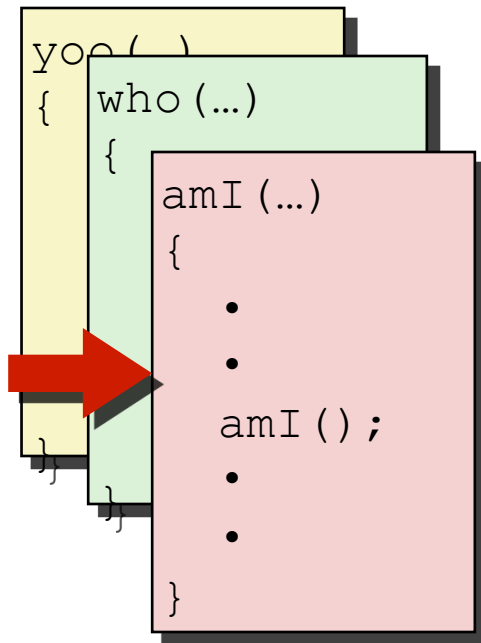
# Example



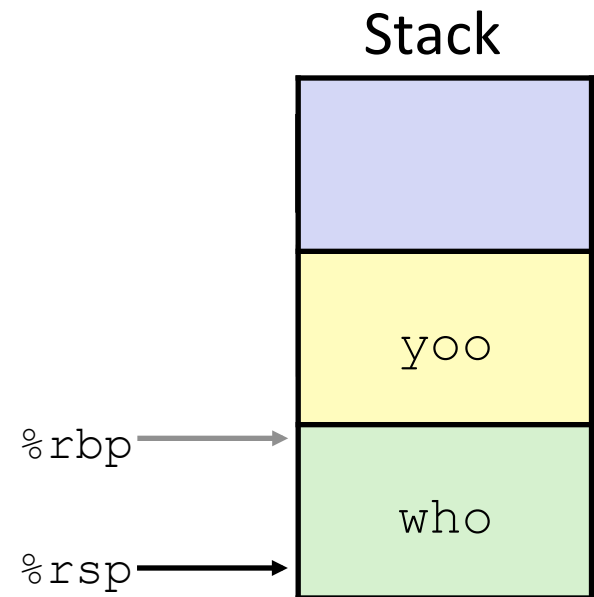
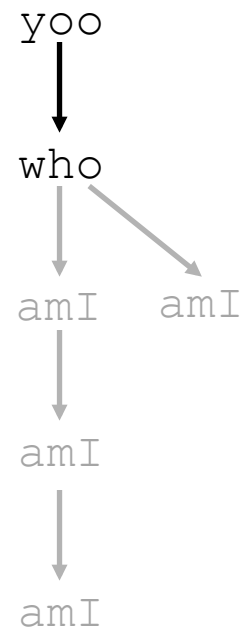
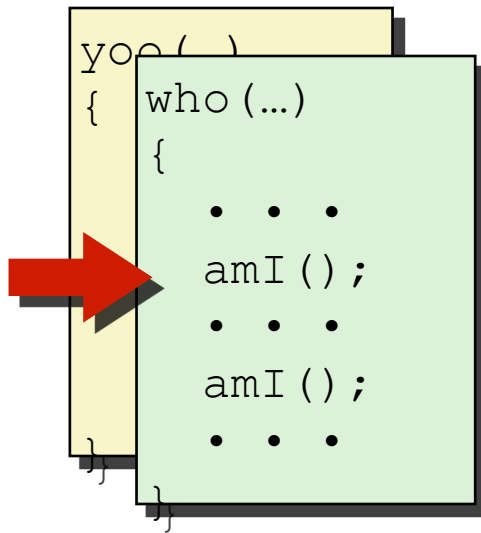
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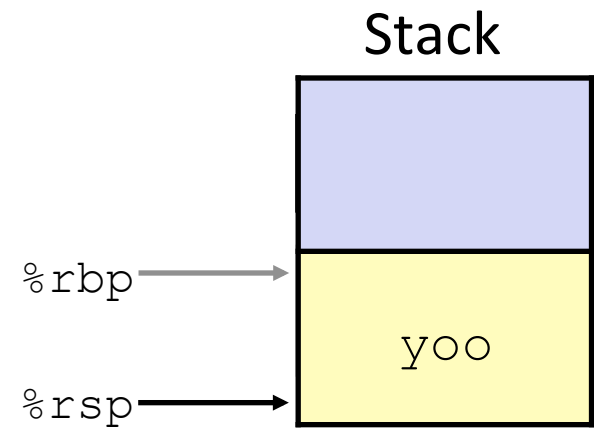
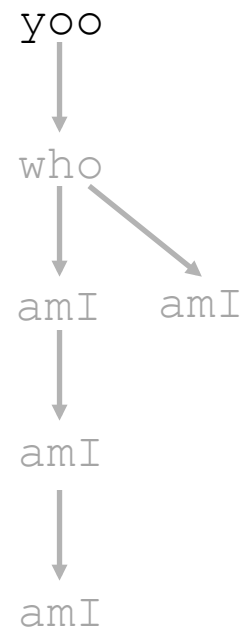
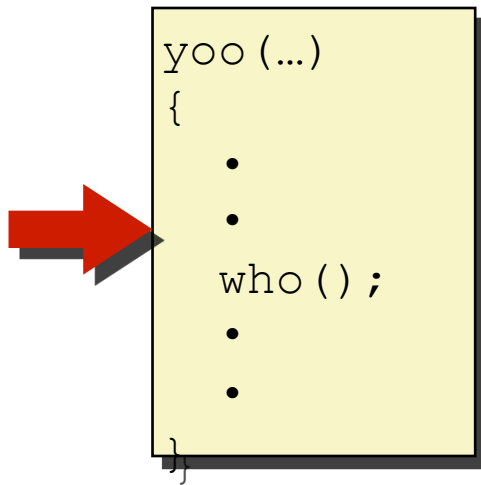
# Example



# Example



# Example



# x86-64/Linux Stack Frame

## Scenario: Caller called a function

### ■ Current Stack Frame (“Top” to Bottom)

- “Argument build:”  
Parameters for function about to call
- Local variables  
If can’t keep in registers
- Saved register context
- Old frame pointer (optional)

### ■ Caller Stack Frame

- Return address
  - Pushed by `call` instruction
- Arguments for this call

