

# VRPTW + CF

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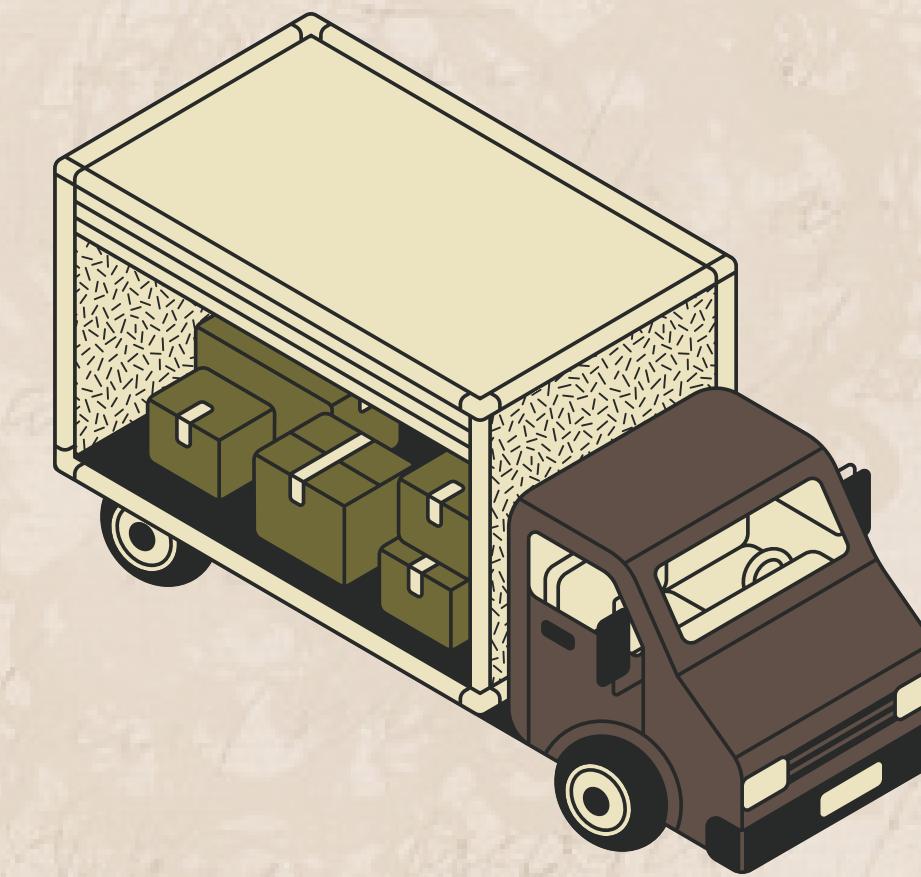
# Plan de la présentation

- 01 Problématique de VPRTW**
- 02 Fondement théroïque de CF**
- 03 Méthodologie de résolution du VRPTW avec CF**
- 04 Analyse**
- 05 Performance sur quelque instance de Solomon's Benchmark**

# 01- Problématique

Problème VRPTW

VEHICLE ROUTING PROBLEM WITH  
TIME WINDOW



# Définition du problème

Le problème de routage de véhicules avec fenêtres temporelles (VRPTW) peut être défini comme le choix des itinéraires pour un nombre de véhicules afin de desservir un groupe de clients dans des plages horaires spécifiques.



# Les Principales contraintes

## Contraintes

la capacité des véhicules

les fenêtres temporelles des clients

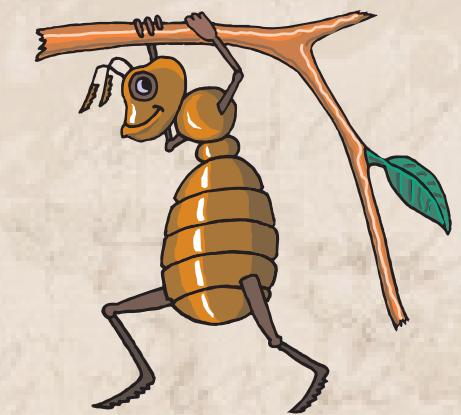
## 02- CF

FONDEMENT THÉORIQUE DE CF  
ANT COLONY  
COLONIE DE FOURMIS



# Ant Colony :

L'Ant Colony Optimization est choisi pour résoudre le VRPTW en raison de sa capacité à trouver des solutions efficaces dans des problèmes d'optimisation combinatoire avec des contraintes temporelles.



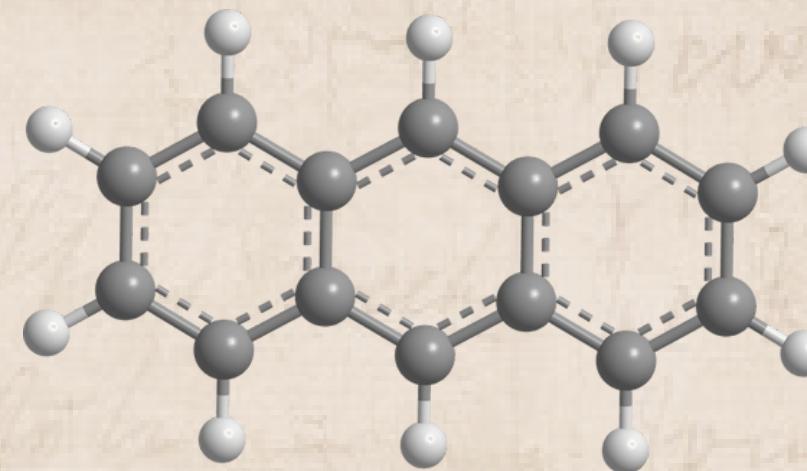
# Ant Colony : Relation avec l'informatique

Les fourmis résolvent des problèmes complexes tels que celui de la recherche du plus court chemin.

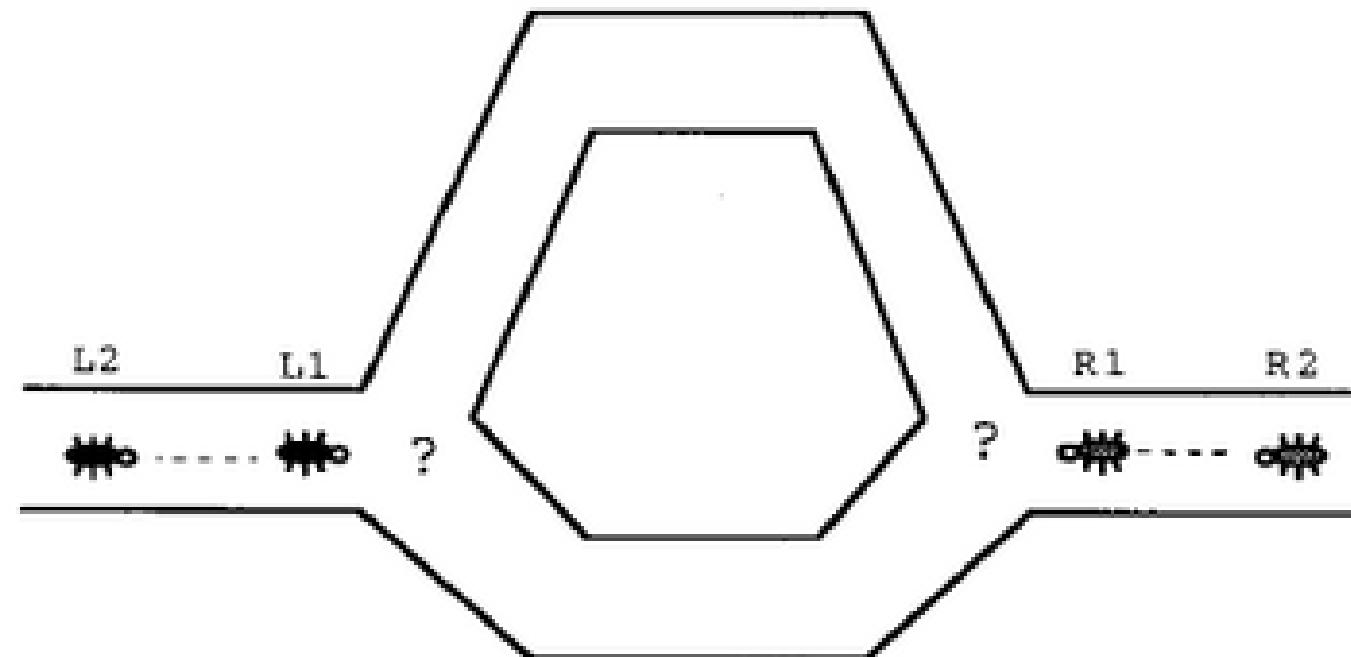


# Ant Colony :Comportement de la fourmi

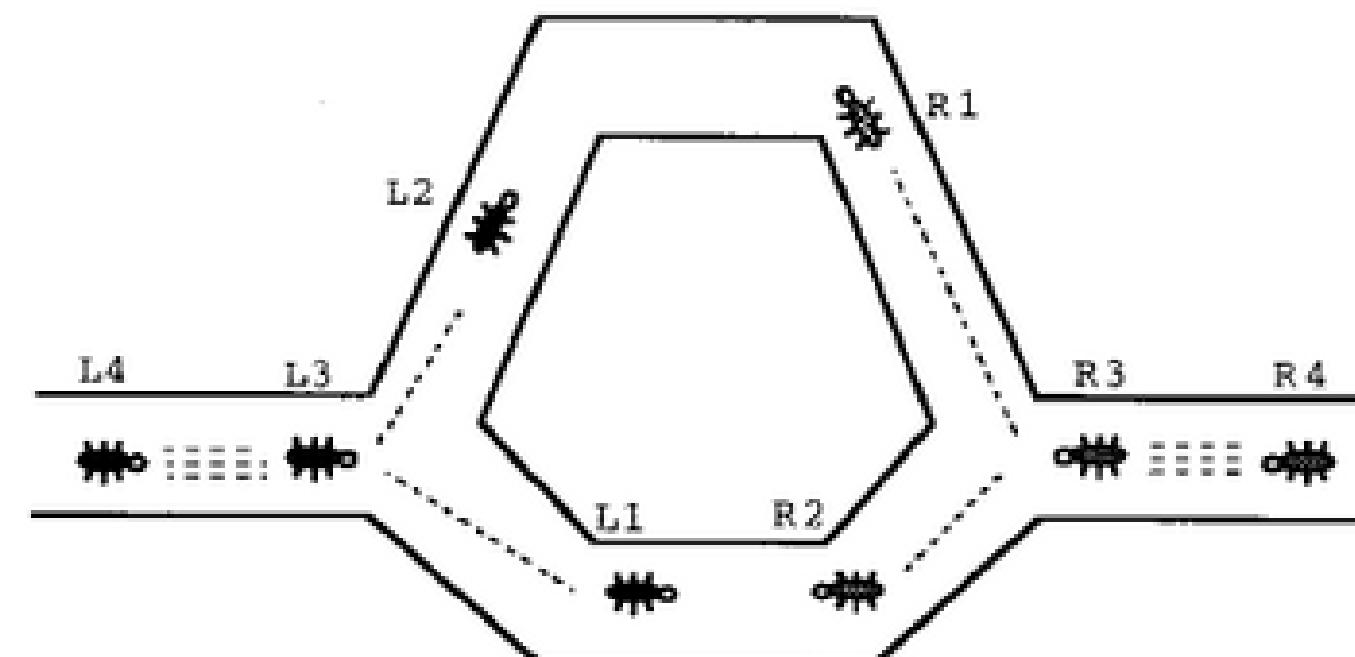
Les fourmis utilisent des **phéromones**



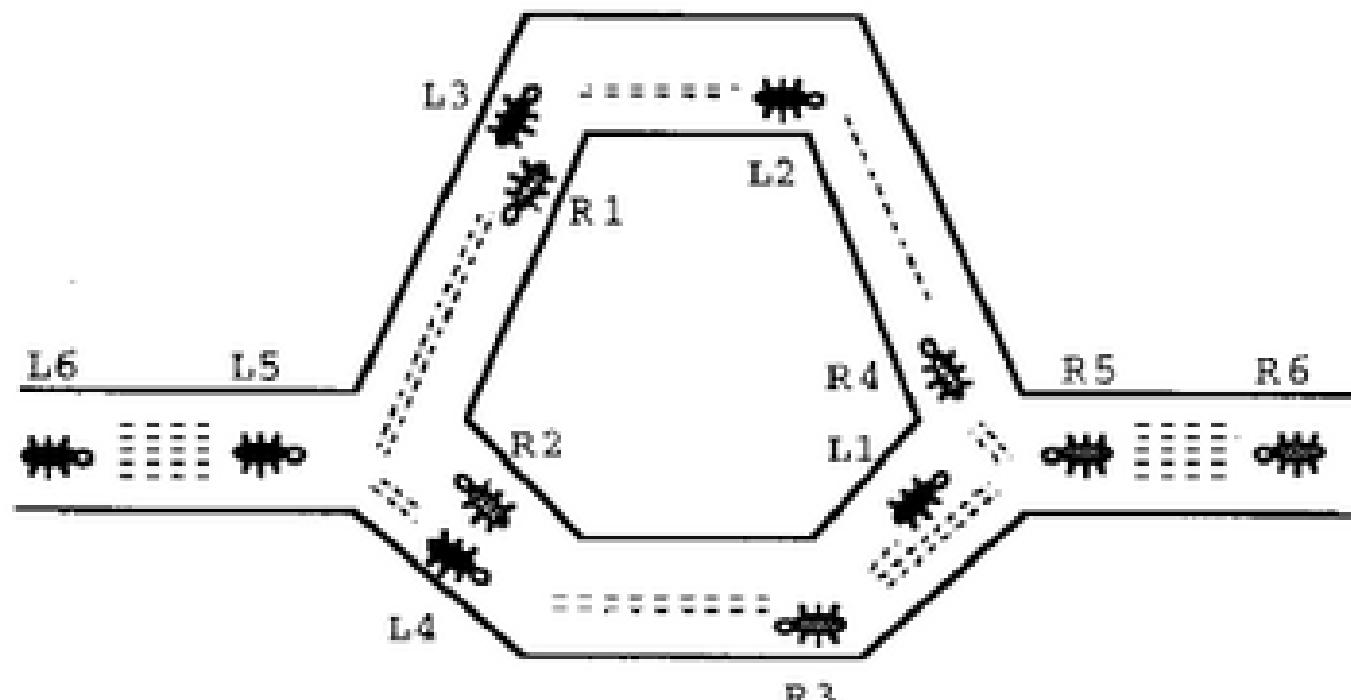
# Ant Colony : Comportement de la fourmi



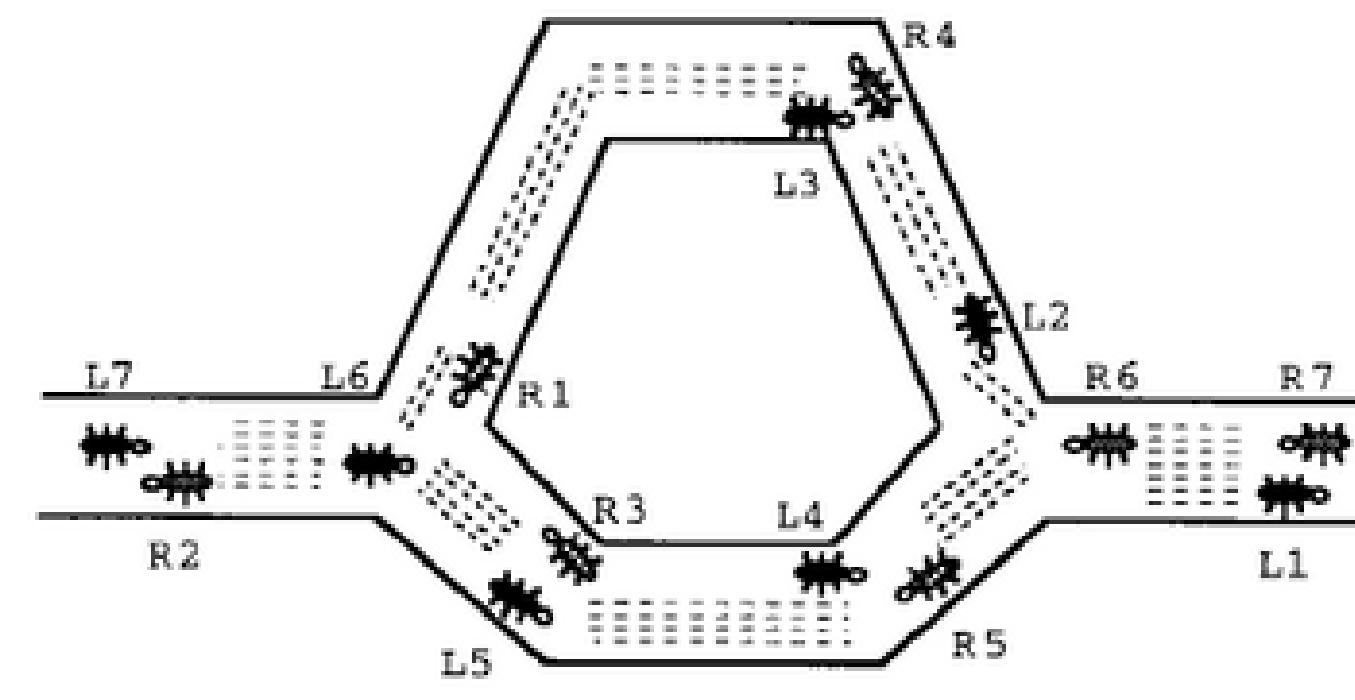
(a)



(b)



(c)



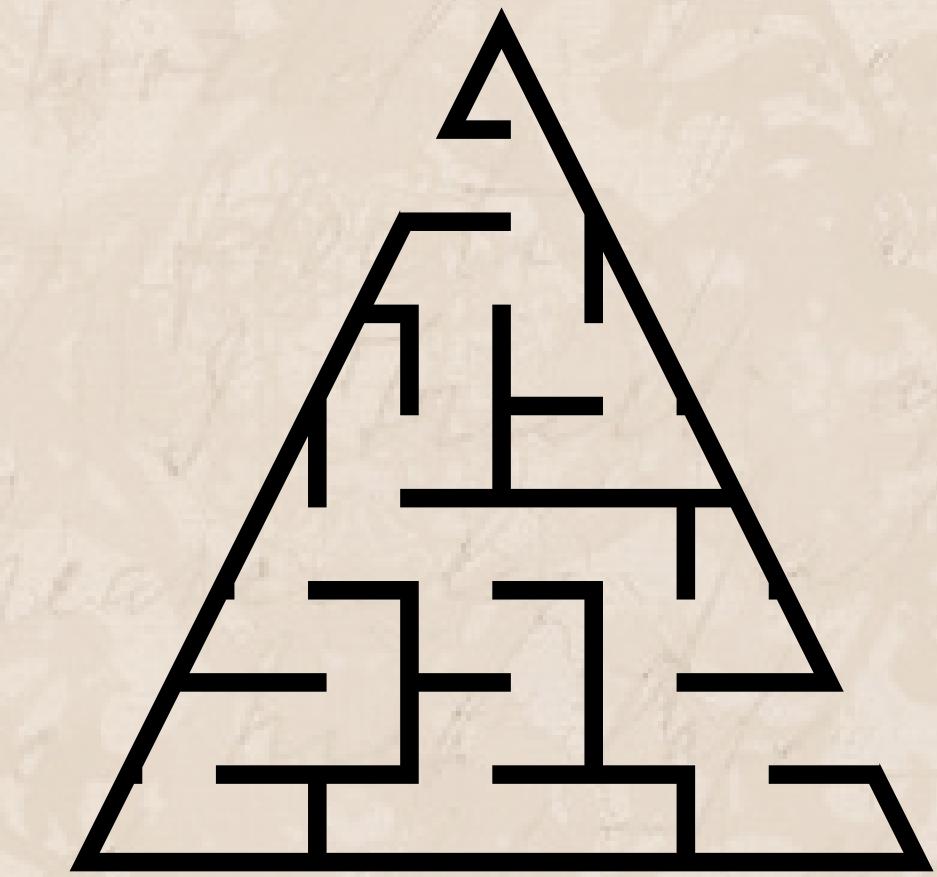
(d)

# Ant Colony :Comportement de la fourmi

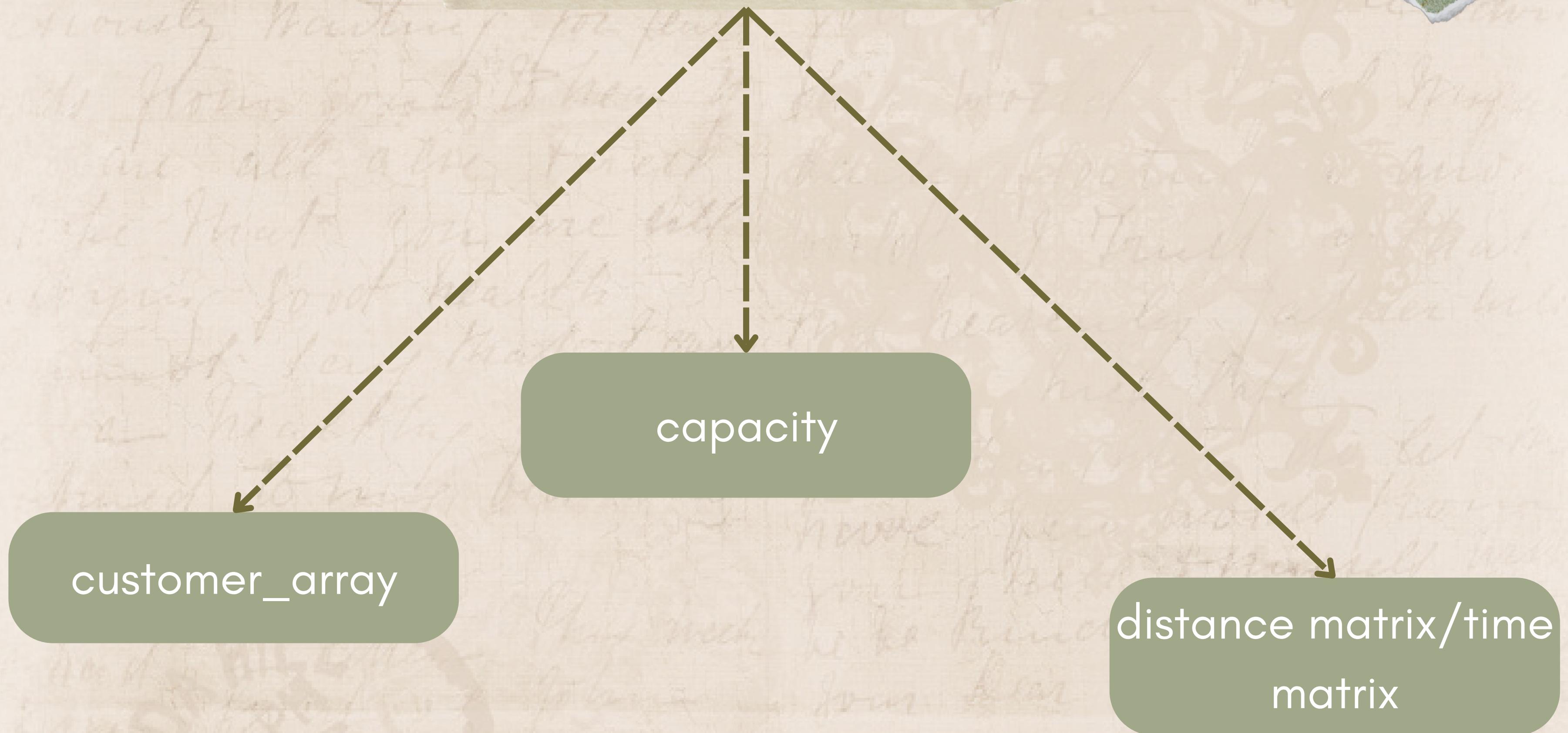
Le comportement des fourmis leur permet de trouver le chemin le plus court vers la nourriture en utilisant les pistes de phéromones, facilitant ainsi le choix optimal parmi plusieurs chemins marqués et constituant la base de méthodes ultérieures de navigation

# 03- Méthodologie

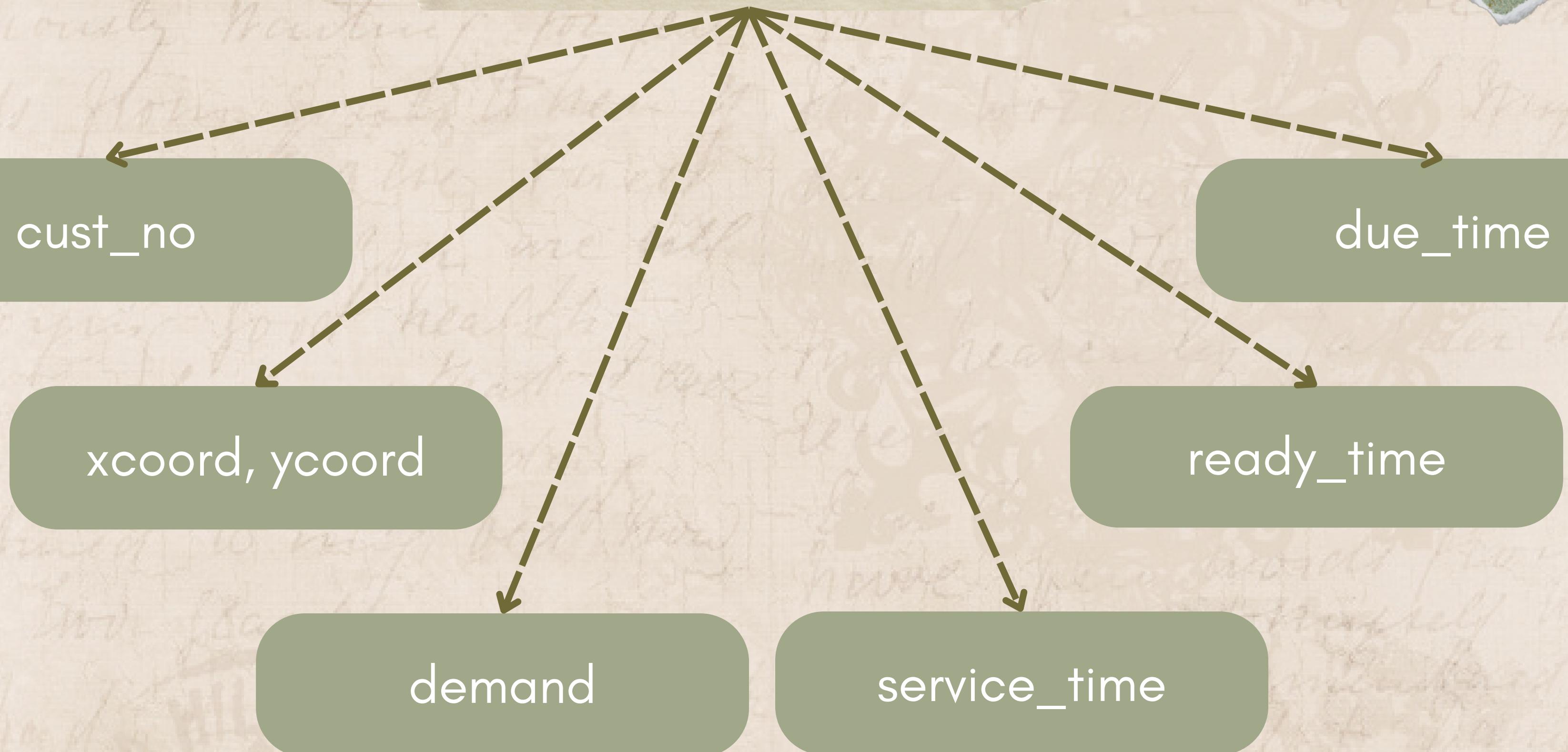
Méthodologie de résolution du  
VRPTW avec CF



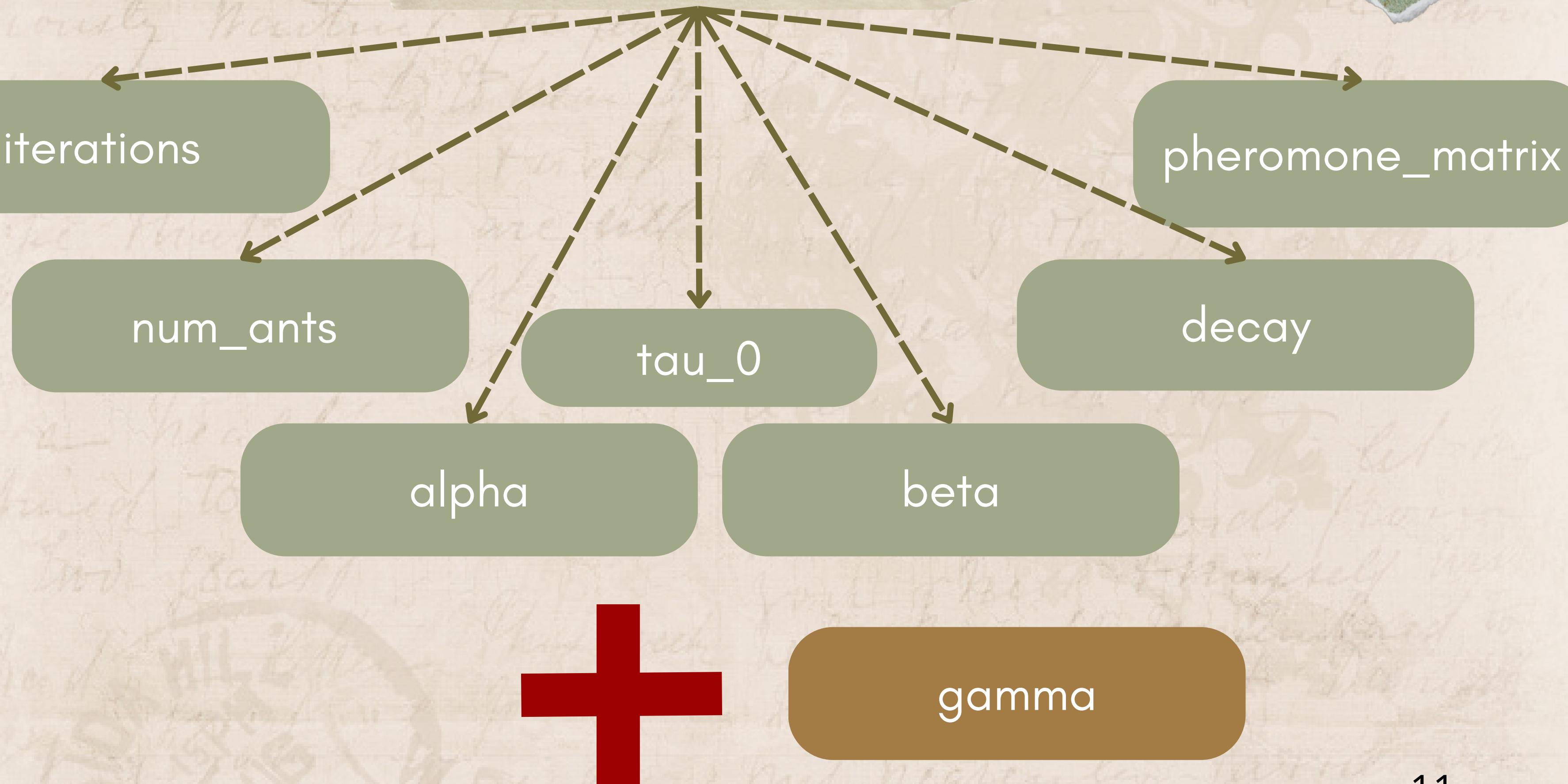
# VPRTW



# Customer



# ACS



# Fonction Main

**Function MAIN**

best\_solution  $\leftarrow$  NULL

**for** I := 1 to iterations:

**for** J := 1 to num\_ants:

path, num\_v, distance  $\leftarrow$  generate\_solution()

current\_solution  $\leftarrow$  (path, num\_v, distance))

local\_update\_pheromones(path)

```
if best_solution == NULL:  
    best_solution ← current_solution  
else:  
    best_solution ← the minimum between  
    current_solution and best_solution depending on number  
    of vehicle in first place then the distance  
end_if;  
end_for;  
global_update_pheromones(best_solution)  
end_for;  
return best_solution
```

# update locale de pheromone

## Function local\_update\_pheromones

pheromone\_matrix  $i,j \leftarrow (1 - \text{decay}) \cdot \text{pheromone\_matrix } i,j + \text{decay} \cdot \text{Tau}_0$

$\text{Tau}_0 = 1 / (\text{number\_of\_customers} * \text{nearest\_neighbor\_solution\_length})$

# update globale de pheromone

**Function** global\_update\_pheromones

pheromone\_matrix ( $i,j$ )  $\leftarrow (1 - \text{decay}) \cdot \text{pheromone\_matrix} (i,j) + \text{decay} \cdot V$

$V = 0$  if  $(i,j)$  is not in best\_solution else  $= 1/\text{distance\_of\_best\_path}$

# generation de solution

## **Function** generate\_solution

```
not_visited ← customer_array  
time ← 0  
num_v ← 1  
capacity ← initial_capacity  
solution ← [[],]  
solution[0].append(depot)  
not_visited.remove(depot)  
possible ← not_visited
```

```
while len(not_visited) > 0 :  
    possible, solution, time, num_v, capacity ← update_possible  
    next_node ← choose_next_node(possible)  
    not_visited.remove(next_node)  
    capacity ← capacity - next_node.demand  
    time = time + time_to_reach_nextNode_from_currentCustomer  
    if time < next_node.ready_time :  
        time ← next_node.ready_time + next_node.service_time  
    else:  
        time ← time + next_node.service_node  
    end_if;  
    solution[-1].append(next_node)  
    end_while;  
solution[-1].append(depot)  
return(solution, num_V, calculate_path_length(solution))
```

# update\_possible

```
Function update_possible
    can_be_served_by_vehicule ← False
    possible ← not_visited
    for customer in not_visited:
        # Check if due date is late
        if time_customer_is_reached > customer.due_date:
            possible.remove(customer)
            continue
        else:
            can_be_served_by_vehicule ← True
        end_if;
        # Check if capacity exceeds the limit
        if capacity < customer.demand :
            possible.remove(customer)
        end_if;
    end_for;
```

```
if len(possible) == 0 :  
    solution[-1].append(depot)  
    capacity ← initial_capacity  
    time ← time + time_to_reach_depot  
    if can_be_served_by_vehicule == False :  
        time← 0  
        num_V ← num_V + 1  
        solution.append([depot,])  
    end_if;  
    solution, possible, time, num_v, capacity = update_possible()  
end_if;  
return solution, possible, time, num_V, capacity
```

## Function choose\_next\_node

```
pheromone_values ← pheromone_matrix(current.cust_no)
```

```
visibility_values ← 0 if customerToReach_not_in_Possible else  
                      (1/ distance_between_Current_and_customerToReach)
```

```
timeP_values ← 0 if customerToReach_not_in_Possible else 1 if  
                  (customerToReach.ready_time <  
                   time_after_reaching_customerToReach) else (1/-  
                  ((time_after_reaching_customerToReach) -  
                   customerToReach.ready_time) )
```

```
for c in customers_array :  
    valueProb(c)  $\leftarrow$  pheromone_values(c)*alpha .  
                           visibility_values(c)*beta . timeP_values*gamma  
end_for;  
  
total_probability  $\leftarrow$  sum(valueProb)  
for c in customers_array :  
    probabilities(c)  $\leftarrow$  valueProb(c)/total_probabilities  
end_for;  
next_node  $\leftarrow$  random_choiceDepending_on_probabilities  
return next_node
```

# Procedure d'optimisation

DR. BY

CH. BY

# improve solution

**Function** improve\_solution(solution):

# Apply tour improvement heuristic to the given solution

solution  $\leftarrow$  optimize\_routes(solution)

# 2-opt heuristic for local search

**for** i in range(len(solution)) :

    route(i)  $\leftarrow$  two\_opt(route(i))

**end\_for;**

**return** solution

# 2-opt

**Function** two\_opt (route):

    route\_dist  $\leftarrow$  route\_total\_distance(route)

    improved\_route  $\leftarrow$  route

**while** improved:

-----

**end\_while;**

**return** improved\_route

```
for i in range(1, len(improved_route) - 3):  
    for j in range(i + 1, len(improved_route) - 2):  
        new_route  $\leftarrow$  Swap edges (i, i+1) with (j, j+1)  
        new_route_dist=route_total_distance(new_route)  
        if (new_route_dist<route_dist) and new_route is valid:  
            improved_route= new_route  
            route_dist=new_route_dist  
            improved= True  
        end_if;  
    end_for;  
end_for;
```

## Function optimize\_routes(routes):

sorted\_routes  $\leftarrow$  Sort routes by their lengths

**while** improved:

**for** i **in** range(len(sorted\_routes)-1):

        new\_route1  $\leftarrow$  sorted\_routes(i)

        # Iterate over each customer in the small route

**for** customer **in** sorted\_routes[i][1:-1]: # Excluding the depot

            # Try to insert the customer into other routes

**for** j **in** range(i+1, len(sorted\_routes)):

                new\_route2  $\leftarrow$  sorted\_routes[j][:]

---

**end\_for;** **end\_for;** **end\_for;**

**end\_while;**

**return** sorted\_routes

```
for k in range(1, len(new_route2)):  
    if customer != depot:  
        new_route2.insert(k, customer)  
        # Check if the resulting route is valid  
        if check_route_valid(new_route2):  
            # Update the route with the inserted customer  
            new_route1.remove(customer)  
            sorted_routes[i] ← new_route1  
            sorted_routes[j] ← new_route2  
            improved ← True  
            customer_inserted ← True  
            break # Move to the next customer
```

```
else:  
    # Revert the insertion if the route is not valid  
    new_route2.remove(customer)  
end_if;  
else :  
    customer_inserted ← True  
    break  
end_if;  
end_for;  
if customer_inserted:  
    break      # Move to the next customer  
end_if;
```

## 04 - Analyse

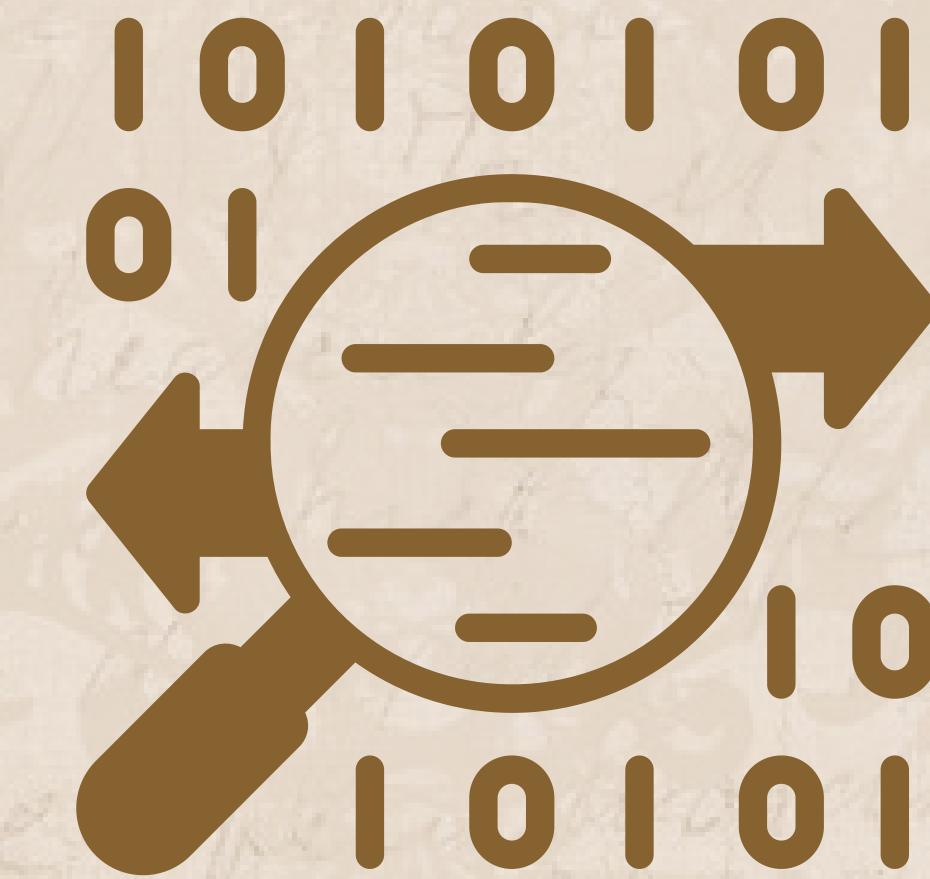
en changeant les paramètres :

alpha

beta

gamma

qu'est ce qu'on obtient ?



# Analyse:

alpha	beta	gamma	Nombre de véhicules	Distances
1.5	2	2	21	2218.994138889913
0.5	2	2	22	2454.760569531627
0.5	0.5	2	26	2999.407832796353
0.5	0.5	0.5	27	2706.029993833985

# Analyse:

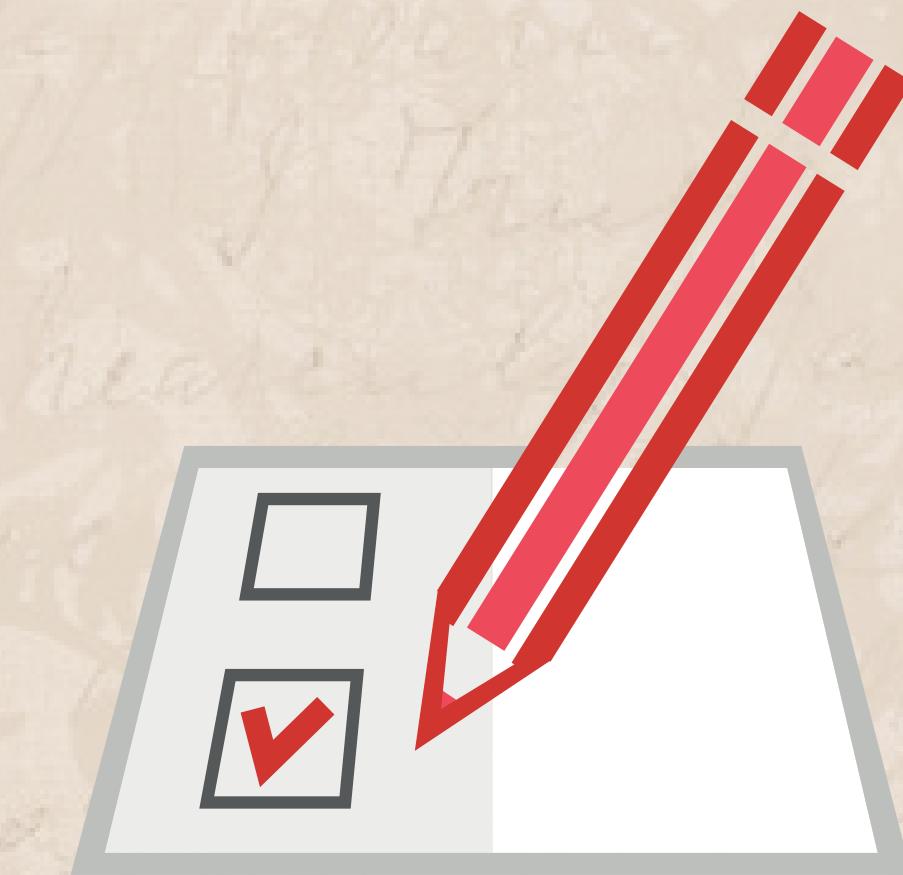
alpha	beta	gamma	Nombre de véhicules	Distances
2	1	0.5	22	2074.173013452977
1.5	3	2	20	2097.1843427051526
3	3	3	23	2499.9212638204167

*Meilleure combinaison : alpha = 1.5 , beta = 3 , gamma = 2*

# Interprétation

Une sélection judicieuse et une adaptation fine de ces paramètres peuvent conduire à des solutions de haute qualité dans divers contextes d'optimisation

# 05- Performance sur quelque instance de Solomon's Benchmark



Instance	Nombre Vehicule	Distance	Nombre Vehicule (B)	Distance (B)
C101	10	828.93686	10	828.93664
C109	10	1180.52763	10	828.93664
C202	4	742.38762	3	591.55626
C205	3	591.55655	3	588.87566
R101	20	2051.47560	19	1650.79864
R108	11	1333.44542	9	960.87528
R207	3	1227.51542	2	890.60784
RC108	12	1467.68051	10	1139.82108
RC203	4	1581.5222	3	1049.62383



Merci Pour Votre Attention