# Sorting Algorithms

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December 3, 2011

- 1 Project Proposal
- 1.1 Goals
- 1.2 Plan
- 2 Sorts
- 2.1 Merge Sort

#### 2.1.1 Algorithm

Merge sort is a comparison-based sorting algorithm, based on the divide-and-conquer design. Its average and worst cases are both nlog(n), and its best case is  $\Omega(n)$ . Invented in 1945 by John von Neumann, it exploits the fact that combining two lists of sorted data is a linear-time process. [1]

#### 2.1.2 Efficiency

The actual code for the algorithm is available in **Listing 1**. This particular implementation uses an optimization that switches to insertion sort on small arrays. This speeds up the algorithm because it allows the small data set to fit entirely in cache along with the small amount of code associated with the insertion sort algorithm. See **Listing 2** for a parallel implementation of merge sort.

#### 2.1.3 Applications

Merge sort is useful in applications where the data set will not fit entirely into memory. This allows for the data to be read in from disk and sorted as it is read, thus requiring a very small memory footprint. In addition, when time complexity needs to be guaranteed, merge sort is preferable over quicksort, as quicksort's worst case is  $O(n^2)$ .

- 2.2 Quick Sort
- 2.2.1 Algorithm
- 2.2.2 Efficiency
- 2.2.3 Applications
- 2.3 Shell Sort
- 2.3.1 Algorithm
- 2.3.2 Efficiency
- 2.3.3 Applications
- 2.4 Comparing and Contrasting
- 3 Conclusions

## A Code Examples

Listing 1: Go implementation of MergeSort

```
package main
2
3
    var optimize int = 10
 4
    func Merge(arr []int) {
5
       if len(arr) < optimize && len(arr) != 1 {</pre>
 6
 7
         InsertionSort(arr)
 8
       else\ if\ len(arr) > 1
9
         // Start at the midpoint
10
         size := len(arr) / 2
11
12
         \mathtt{left} \ := \ \mathtt{arr} \, [\, : \mathtt{size} \, ]
13
         right := arr[size:]
14
15
         // Perform sub-merges
         Merge(left)
16
17
         Merge(right)
18
19
         // For holding results
20
         data := make([]int, len(arr))
21
22
23
         for len(left) > 0 \&\& len(right) > 0  {
            // Decide which list to pull from
24
25
            // Put the first value into the results and shrink the array
26
27
              \mathtt{data}\,[\,\mathtt{i}\,]\,\,,\,\,\,\mathtt{right}\,=\,\mathtt{right}\,[\,0\,]\,\,,\,\,\,\mathtt{right}\,[\,1\,:\,]
28
29
              // Same thing for the left
30
              data[i], left = left[0], left[1:]
31
            i++ // Next element
```

```
33
34
35
        // Figure out which list isn't empty
36
        var remainder [] int
37
        if len(left) > 0  {
38
39
          remainder = left
40
        } else {
41
          remainder = right
42
43
44
        // Empty remaining values
        for len(remainder) > 0 {
45
          data[i], remainder = remainder[0], remainder[1:]
46
47
          i++
48
49
50
        // Copy result to our answer
        for i := 0; i < len(arr); i++ {</pre>
51
52
          arr[i] = data[i]
53
54
     }
55
56
57
   func InsertionSort(arr []int) {
58
     for i, j := 1, 1; i < len(arr); i, j = i+1, i+1 {
59
        60
          arr[j-1], arr[j] = arr[j], arr[j-1]
61
62
63
     }
64
```

Listing 2: Parallel implementation of MergeSort in Go

```
1
     package main
 2
 3
     import "sync"
 5
     var optimize int = 25
 6
 7
     func Merge(arr []int, group *sync.WaitGroup) {
 8
        if len(arr) < optimize && len(arr) != 1 {</pre>
 9
           InsertionSort(arr)
10
        } else if len(arr) > 1 {
          // Start at the midpoint size := len(arr) / 2
11
12
13
14
           left := arr[:size]
15
           right := arr[size:]
16
           // Perform sub-merges
17
18
           \mathtt{group} \; := \; \underset{}{\mathbf{new}}(\,\mathtt{sync}\,.\,\mathtt{WaitGroup}\,)
19
           \mathtt{group}. \mathtt{Add}(2)
             go Merge(left, group)
go Merge(right, group)
20
21
22
           group.Wait()
23
```

```
24
            // For holding results
            \mathtt{data} \; := \; \frac{\mathsf{make}}{\mathsf{make}}([] \; \mathbf{int} \; , \; \; \mathbf{len}(\mathtt{arr} \, ))
25
26
27
            \mathtt{i} \; := \; 0
            for len(left) > 0 \&\& len(right) > 0  {
28
29
               // Decide which list to pull from
30
               if left [0] > right [0] {
31
                  // Put the first value into the results and shrink the array
                  data[i], right = right[0], right[1:]
32
33
               } else {
34
                  // Same thing for the left
35
                  \mathtt{data}\left[\,\mathtt{i}\,\right]\,,\ \ \mathtt{left}\,=\,\mathtt{left}\left[\,0\,\right]\,,\ \ \mathtt{left}\left[\,1\,:\,\right]
36
37
               i++ // Next element
38
39
40
            // Figure out which list isn't empty
41
            var remainder [] int
42
43
            if len(left) > 0 {
44
               {\tt remainder} \, = \, {\tt left}
45
            } else {
46
               remainder = right
47
48
49
            // Empty remaining values
50
            for len(remainder) > 0 {
               \mathtt{data}\,[\,\mathtt{i}\,]\,,\ \mathtt{remainder}\,=\,\mathtt{remainder}\,[\,0\,]\,,\ \mathtt{remainder}\,[\,1\,:\,]
51
52
53
54
55
            // Copy result to our answer
56
            for i := 0; i < len(arr); i++ {
               arr[i] = data[i]
57
58
59
         {\tt group.Done}\,(\,)
60
61
62
63
      func InsertionSort(arr []int) {
64
         \label{eq:formula} \mbox{for } \mbox{i} \,, \mbox{ } \mbox{j} \; := \; 1 \,, \; \; 1; \; \; \mbox{i} \; < \; \mbox{len}(\mbox{arr}); \; \; \mbox{i} \,, \; \; \mbox{j} \; = \; \mbox{i} + 1, \; \mbox{i} + 1 \; \left\{ \right.
65
            arr[j-1], arr[j] = arr[j], arr[j-1]
66
67
68
69
         }
70
```

### References

[1] Thomas H. Cormen, Charles E. Leiserson, Ronald L. Rivest, and Clifford Stein. *Introduction to algorithms*. MIT Press, Cambridge, Mass, 2nd edition, 2001.