## **Natural Transformations**

$$C \xrightarrow{F} T$$

 $\mathcal{C} \xrightarrow[G]{F} \mathcal{D}$   $\eta$  is a natural transformation  $\forall C \in obj(\mathcal{C})$  we have a morphism

$$F(C) \xrightarrow{\eta_C} G(C)$$

such that given  $f:A\to B$  in  $\mathcal C$   $F(A)\stackrel{\eta_A}{\longrightarrow} G(A)$ 

$$F(A) \xrightarrow{\eta_A} G(A)$$

$$\downarrow^{F(f)} \qquad \downarrow^{G(f)}$$

$$F(B) \xrightarrow{\eta_B} G(B)$$

We also have the identity:  $C \xrightarrow{F} D$ Composition and communitivity work here too!

Exponential of Categories of Categories:  $\mathcal{D}^{\mathcal{C}}$  -wow-