

## 0.1 Exponentials

**Definition 1.** For a category  $\mathcal{C}$  with binary products, an exponential  $C^B$  is associated with two objects and an evaluation arrow  $\epsilon : A \times B \rightarrow C$  if

$$A \times B \xrightarrow{f} C$$

then there exists a unique

$$A \times B \xrightarrow{\tilde{f}} C^B$$

where  $\tilde{f}$  is the transpose of  $f$  such that

$$\begin{array}{ccc} C^B \times B & \xrightarrow{\epsilon} & C \\ \uparrow \tilde{f} \times 1_B & \nearrow f & \\ A \times B & & \end{array} \quad \text{and} \quad \begin{array}{ccc} C^B & & \\ \uparrow g & & \\ A & & \end{array}$$

Check the transpose of the transpose is the thing...

$$\begin{aligned} \epsilon \circ f \times 1_B &= \tilde{f} \\ \tilde{\tilde{f}} &= \epsilon \circ (\tilde{f} \times 1_B) = f \\ \bar{g} &= \epsilon \circ (g \times 1_B) \therefore \tilde{\bar{g}} = g \end{aligned}$$

In set:

$$\begin{aligned} C^B &= \{f : B \rightarrow C\} \\ \epsilon : C^B \times B &\rightarrow C \\ \epsilon(f, b) &= f(b) \end{aligned}$$

So...

$$\begin{aligned} f &: A \times B \rightarrow C \\ f(a, b) &= c \\ f(\tilde{a}) &= f(a, *) : B \rightarrow C \\ g &: A \rightarrow C^B \\ \bar{g} &: A \times B \rightarrow C \\ \bar{g}(a, b) &= (g(a))(b) \end{aligned}$$

This is just currying!

## 0.2 More Exponentials

Another way to write  $A \times B \rightarrow C$  is as  $A \rightarrow C^B$ , meaning  $Hom(A \times B, C) \cong Hom(A, C^B)$  is natural. This is an example of a left adjoint  $(- \times B)$  and a right adjoint  $(-^B)$ .

TODO: Example 6.6 with graphs pg 124

$-^A : C \rightarrow C$  is a functor. For a functor we need to know what it does with objects and with morphisms.

$$\text{Define } \beta : B^A \rightarrow C^A \text{ and } \epsilon : < TODO >. \text{ then } B \longmapsto B^A \quad B \xrightarrow{\beta} C \quad B^A \xrightarrow{\beta \circ \epsilon} C^A$$