## 1 Exponential

**Definition 1.** For a cotegory C with binary products, an exponential  $C^B$  is associated with two objects and an evaluation arrow  $\epsilon: A \times B \to \text{if}$ 

$$A \times B \stackrel{f}{\longrightarrow} C$$

then there exists a unique

$$A\times B\stackrel{\tilde{f}}{-\!\!\!-\!\!\!-\!\!\!-} C^B$$

where  $\tilde{f}$  is the transpose of f such that

$$\begin{array}{c} C^B \times B \xrightarrow{\varepsilon} C & C^B \\ \tilde{f} \times 1_B & \\ A \times B & A \end{array}$$
 and  $g \upharpoonright A \times B$ 

Check the transpose of the transpose is the thing...

$$\varepsilon \circ f \times 1_B = \tilde{f}$$

$$\bar{\tilde{f}} = \varepsilon \circ (\tilde{f} \times 1_B) = f$$

$$\bar{g} = \varepsilon \circ (g \times 1_B) : \tilde{\tilde{g}} = g$$

In set:

$$C^{B} = \{f : B \to C\}$$
$$\varepsilon : C^{B} \times B \to C$$
$$\varepsilon(f, b) = f(b)$$

So...

$$\begin{split} f:A\times B &\to C \\ f(a,b) &= c \\ \tilde{f(a)} &= f(a,*):B \to C \\ g:A \to C^B \\ \bar{g}:A\times B \to C \\ \bar{g}(a,b) &= (g(a))(b) \end{split}$$

This is just currying!

## 1.1 More Exponentials

Another way to write  $A \times B \to C$  is as  $A \to C^B$ , meaning  $Hom(A \times B, C) \cong Hom(A, C^B)$  is natural. This is an example of a left adjoint  $(- \times B)$  and a right adjoint  $(-^B)$ .

TODO: Example 6.6 with graphs pg 124

 $-^A:C\to C$  is a functor. For a functor we need to know what it does with objects and with morphisms.

Define 
$$\beta: B^A \to C^A$$
 and  $\varepsilon: < TODO >$ . then  $B \longmapsto B^A$   $B \stackrel{\beta}{\longrightarrow} C$   $B^A \stackrel{\beta \circ \varepsilon}{\longrightarrow} C^A$