

NOSQL AND CASSANDRA

@rantav

NoSQL



ABOUT MYSELF

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reversim.com

prettyprint.me

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speakerdeck.com/u/rantav/



ABOUT MYSELF

Full stack developer

User of Cassandra

Contributor to Cassandra

Creator of Hector

User of MongoDB





NOSQL



NOSQL

A Multi headed beast?

WHAT'S WRONG WITH SQL?

- **Nothing really wrong**
- But **some problems** - easier to solved with no:sql
- Examples:
 - High scale (reads, write, data size)
 - Data structures (sorted sets, hashes)
 - Graph processing



RELATIONAL DATABASES - SWISS ARMY KNIFE?

- You used to be able to do **everything** with Relational Databases
- Until you reach:
 - **Scaling** limits
 - **\$\$\$** limits
 - **Flexibility** limits
 - **Performance** limits



NOSQL USE CASES

- Large amounts of data (linear horizontal scalability)
- Massive writes
- Speedy key-value access
- Graphs
- High availability
- Easier for ops (well, for some...)
- Easier for programmers

TAXONOMY



KEY-VALUE CACHING

- **memcached**
- repcached
- coherence
- infinispan
- eXtreme scale
- jboss cache
- velocity
- Terracotta



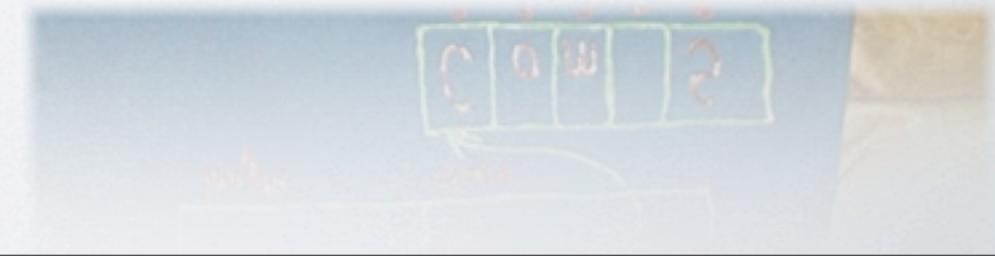
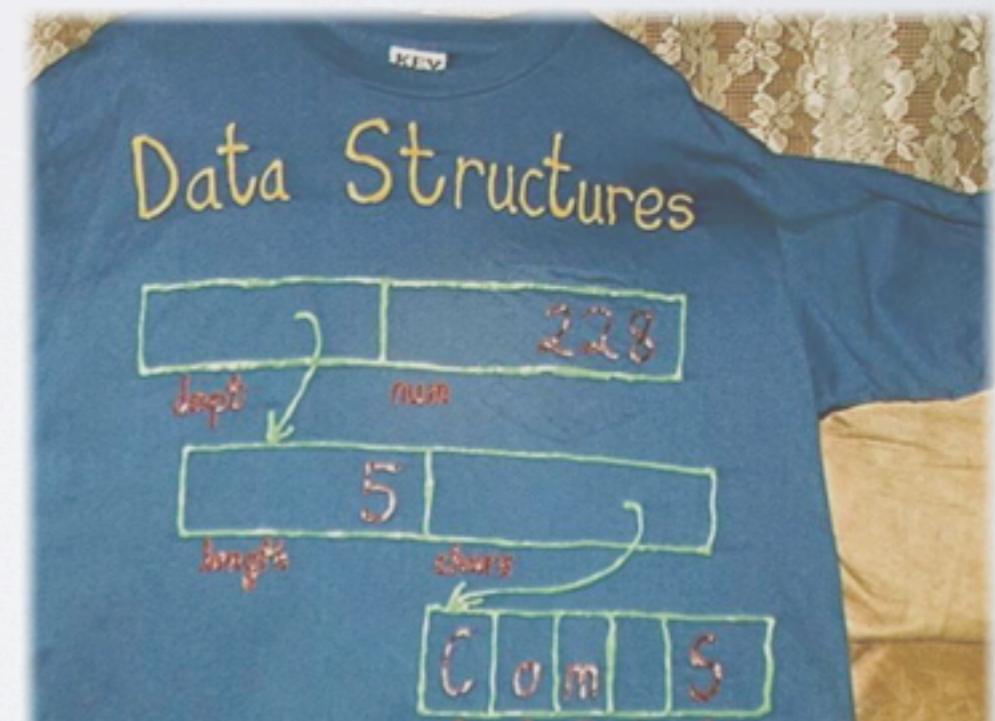
KEY-VALUE STORE

- **DynamoDB**
- voldemort
- Dynomite
- SubRecord
- Mo8onDb
- Dovetaildb
- **Tokyo Tyrant**
- lightcloud
- NMDB
- luxio
- **memcachedb**
- **CouchBase**
- actord
- **Riak**



DATA STRUCTURE STORES

- Redis



DOCUMENT STORE

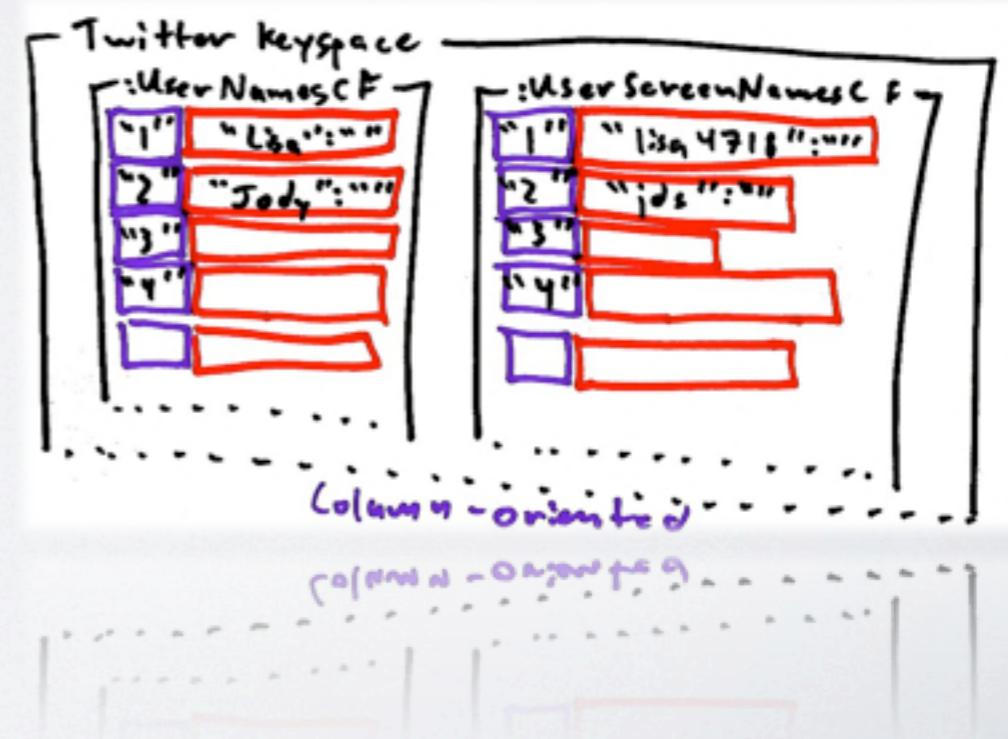
- **CouchDB**
- **MongoDB**
- Jackrabbit
- XML Databases
- ThruDB
- CloudKit
- Perservere
- Scalaris



COLUMNAR STORES

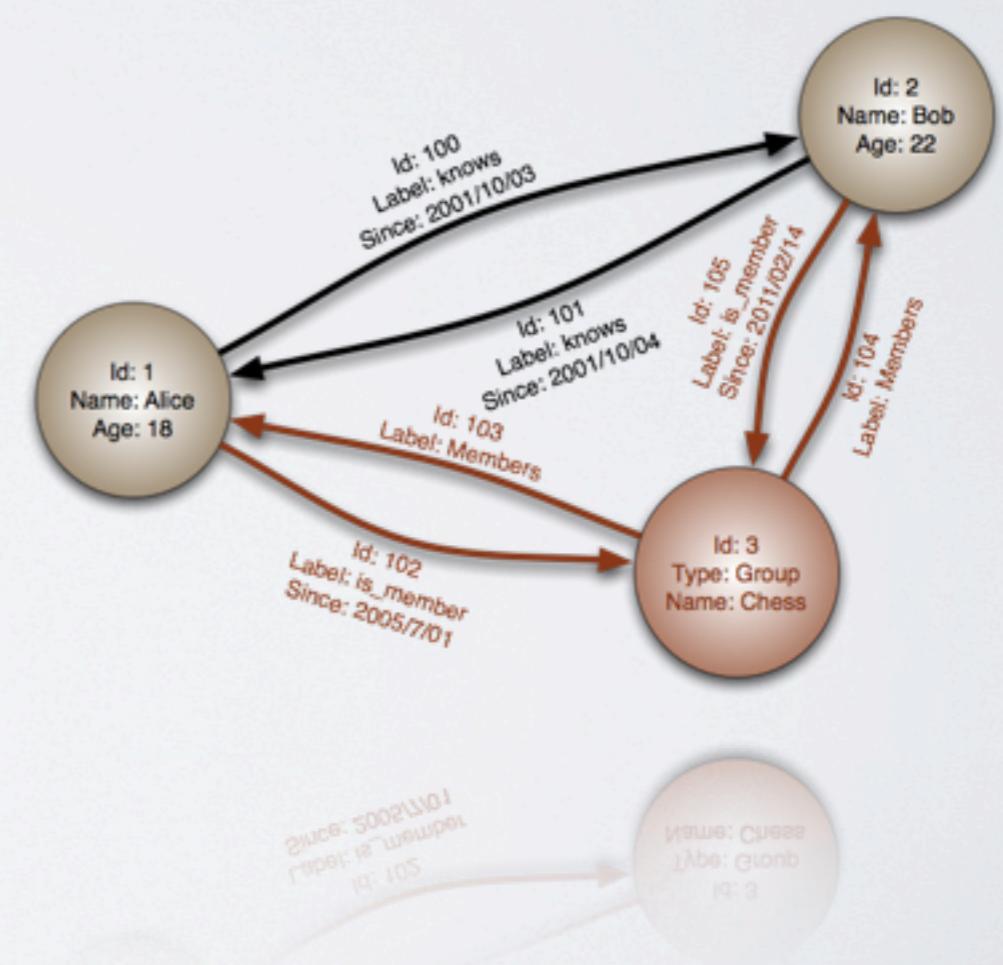
- **BigTable**
- **Hbase**
- **Cassandra**
- Hypertable
- KAI
- OpenNeptune

- Qbase
- KDI



GRAPH DATABASES

- Neo4J
- VertexDB
- FlockDB
- DEX
- OrientDB





- Decentralized (p2p)
- Very high scale
- Column oriented
- Replication and Multi-DC
- Fault tolerant
- Tunable consistency
- Tunable Caching
- MapReduce
- Query language (CQL)
- Thrift



- Document Database
- Ad-hoc queries
- Indexing
- Replication, Sharding
- MapReduce
- File storage
- Capped collections
- Flexible schema



redis

- Key-Value store
- Data structure server
- (Mostly) in-memory
- Very high performance
- Master-slave setup
- Pub/Sub
- Atomic, fast operations:
 - Sets
 - Sorted Sets
 - Strings
 - Hashes

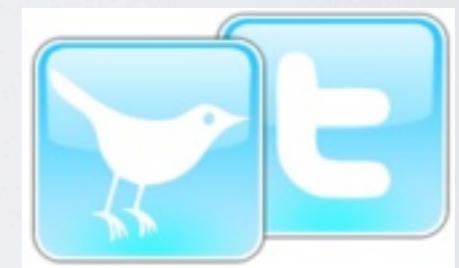
Reminder



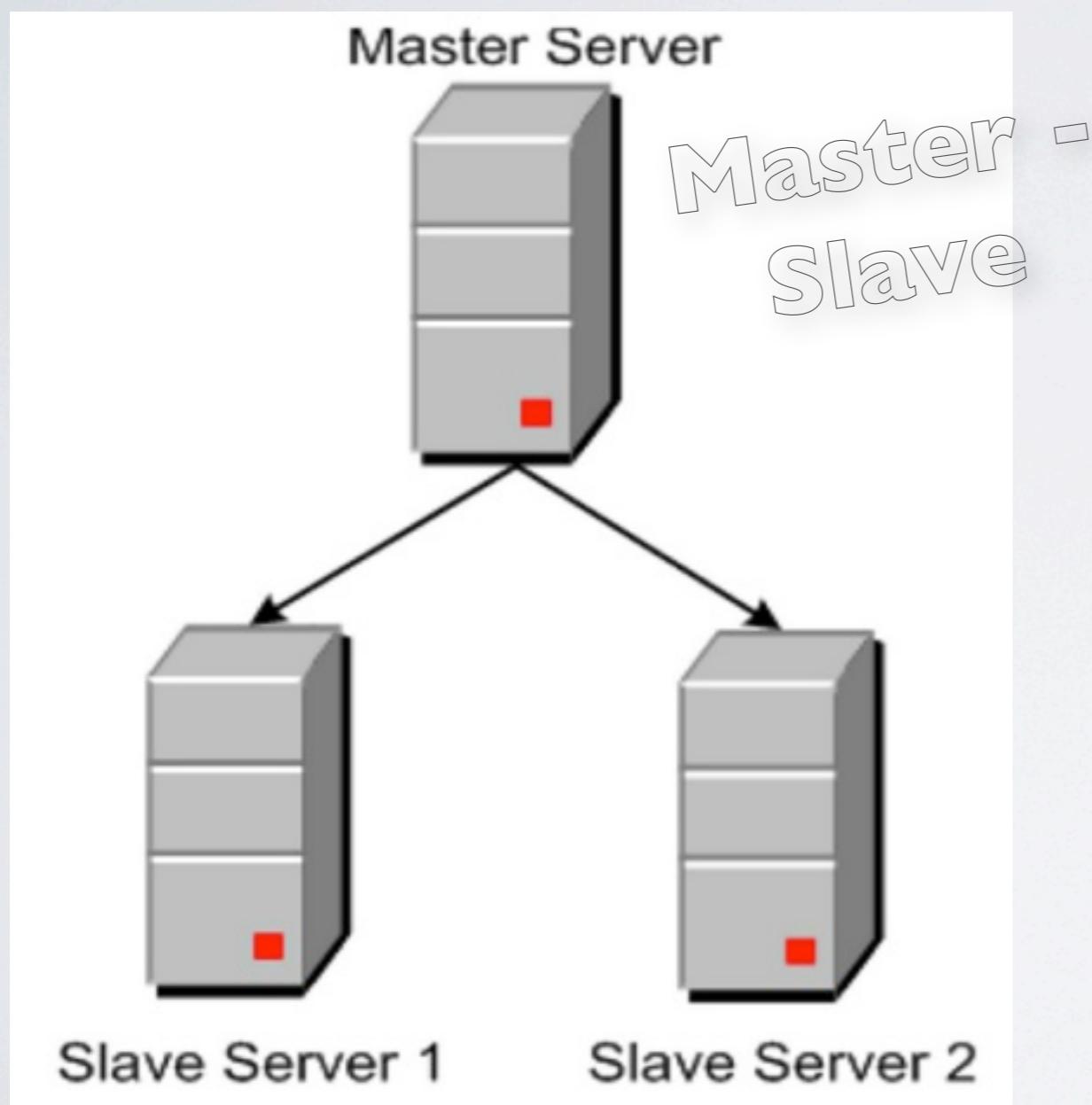
- ACID
- Atomicity
- Consistency
- Isolation
- Durability
- Transactional
- SQL
- Stored Procedures
- (Usually) optimized for reads
- MySQL Cluster
- NoSQL and SQL mix

TODAY'S FOCUS: SCALING

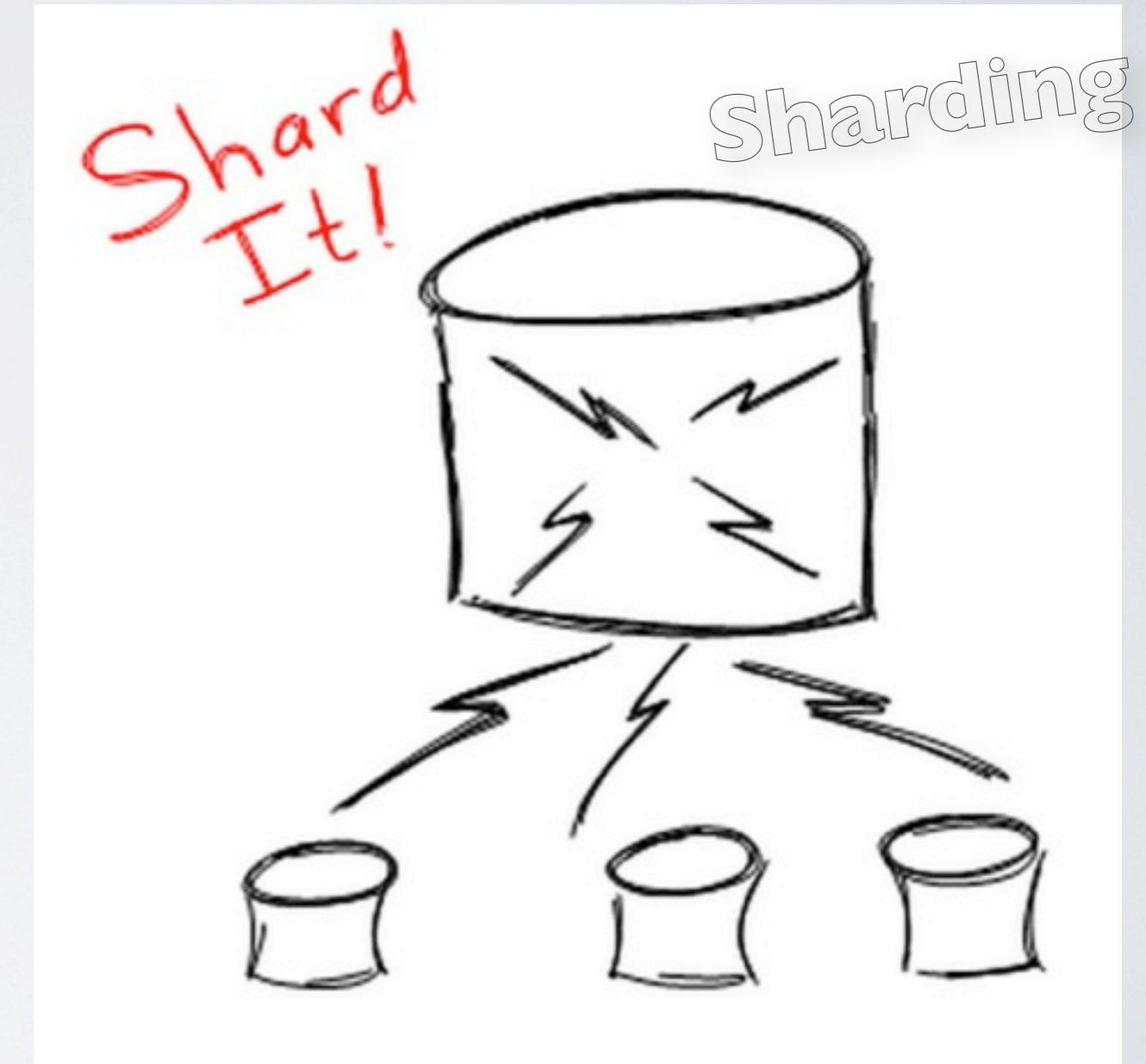
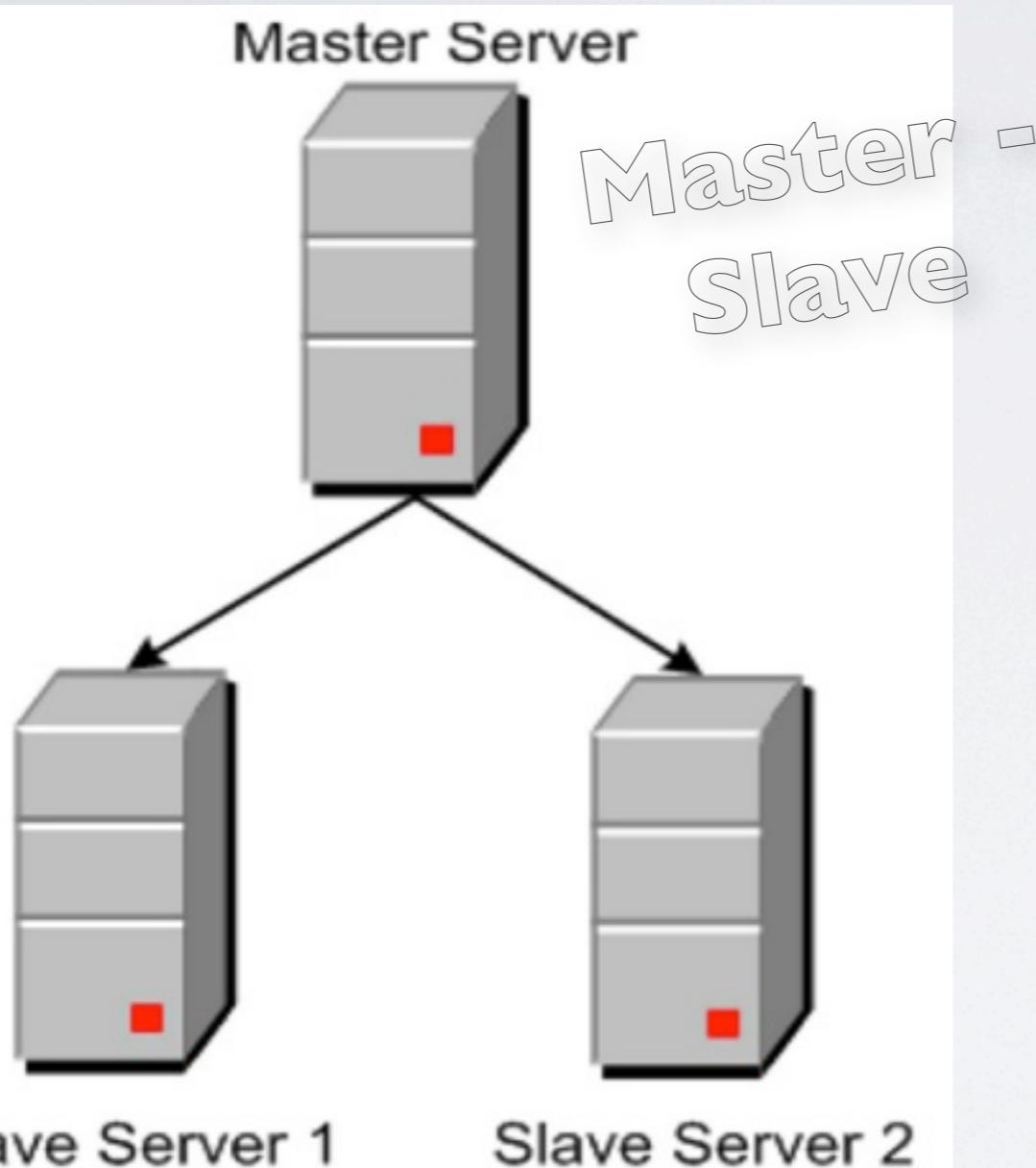
- “Internet scale” data size
- High read/write rates
- Frequent schema changes
- **Social** apps: not **banks**
⇒ can relax ACID



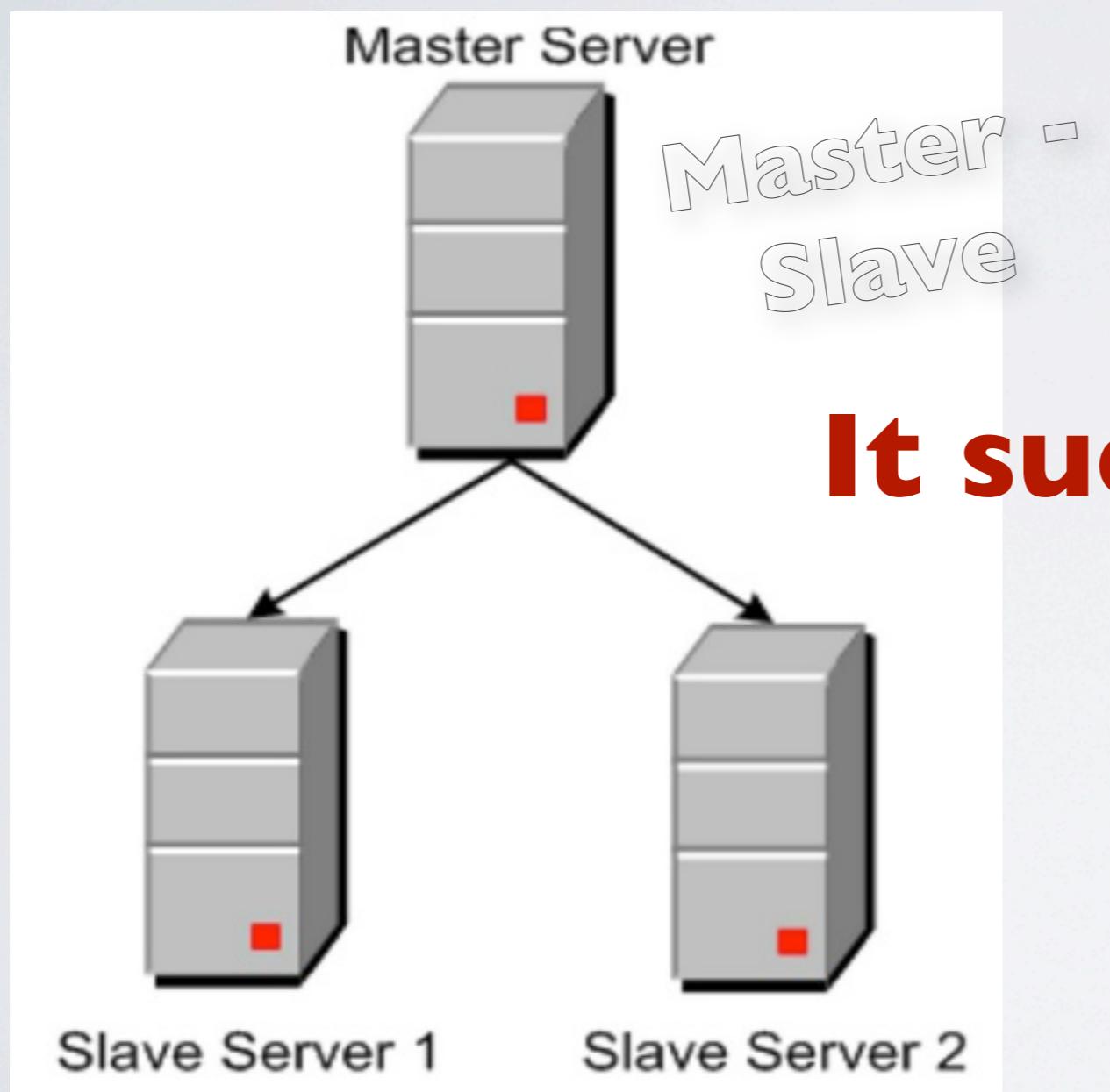
TRADITIONAL SCALING



TRADITIONAL SCALING



TRADITIONAL SCALING



CASSANDRA

- Developed at Facebook
- Now used by many .com
- BigTable data model: **Columnar**
- Dynamo **Eventual Consistency**
- Apache **open source**
- Implemented in **Java**

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Goal:
Easy Scaling

WHAT IS EVENTUAL CONSISTENCY?

- Full Consistency: Once you write, it can be read
- Eventual Consistency: Given enough time, **eventually** data can be read

Observation:

Eventual Consistency \neq ACID

ACADEMIC

DIFFERENT LEVELS OF EVENTUAL CONSISTENCY

- Causal consistency
- Read-your-writes consistency
- Session consistency
- Monotonic read consistency
- Monotonic write consistency

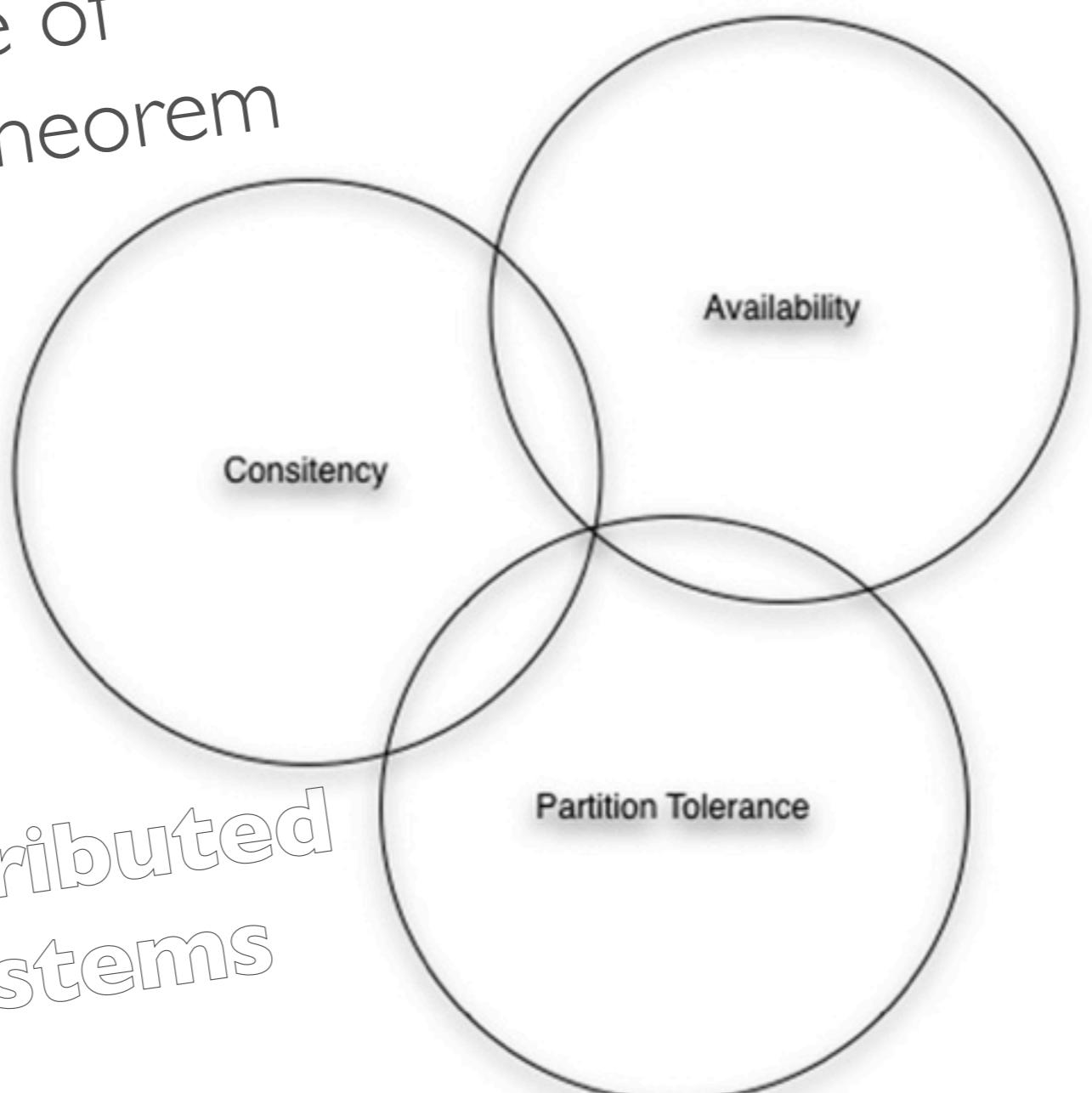
ACADEMIC

WHY EVENTUAL CONSISTENCY AT ALL?

WHY EVENTUAL CONSISTENCY AT ALL?

Because of
the **CAP** theorem

Distributed
Systems



ACADEMIC

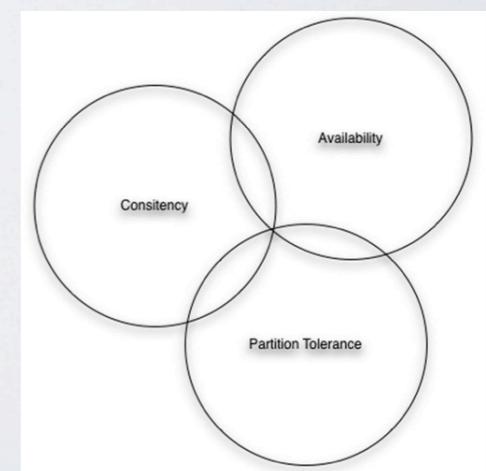
EVENTUAL CONSISTENCY DOWN TO EARTH

How it's done in Cassandra

N - Number of replicas (nodes) for any data item

W - Number or nodes a write operation blocks on

R - Number of nodes a read operation blocks on



NRW - TYPICAL VALUES

W=1 ⇒ Block until **first** node writes successfully

W=N ⇒ Block until **all** nodes writes successfully

W=0 ⇒ **Async** writes

NRW - TYPICAL VALUES

R=1 ⇒ Block until the **first** node returns an answer

R=N ⇒ Block until **all** nodes return an answer

R=0 ⇒ Doesn't make sense

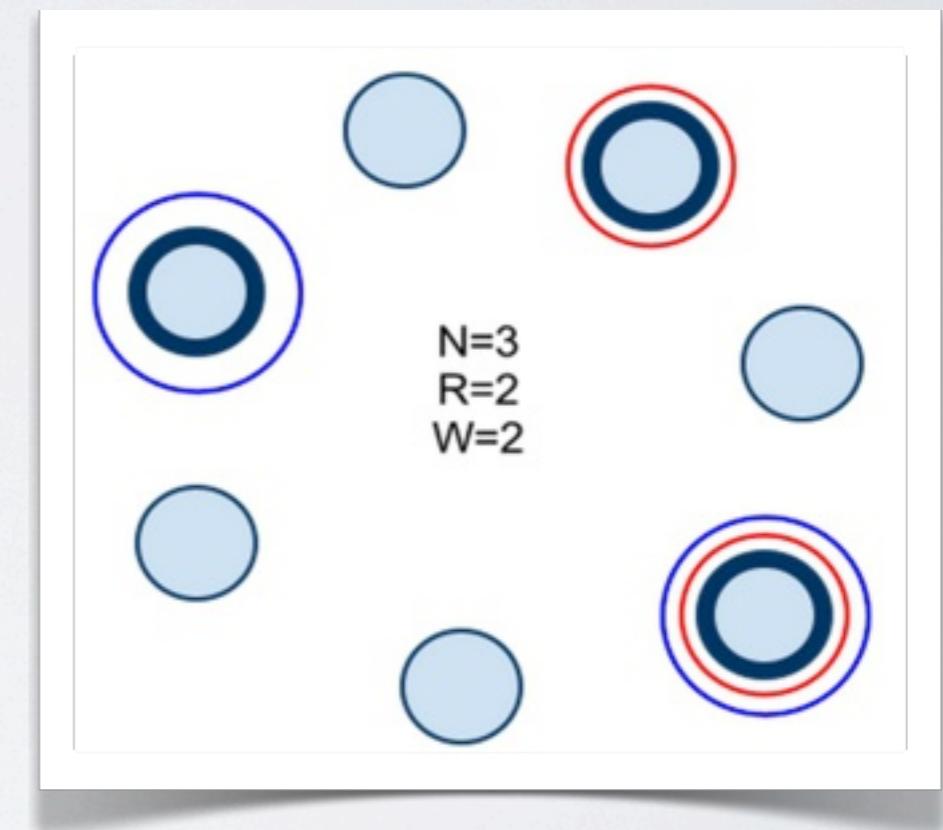
NRW - TYPICAL VALUES

QUORUM

$$R = N/2 + 1$$

$$W = N/2 + 1$$

⇒ Fully consistent



NRW - OBSERVATION

Observation:

When $R + W < N \Rightarrow$ Weak consistency

DATA MODEL IN CASSANDRA

Forget **SQL**



DATA MODEL - VOCABULARY

- **Keyspace** – like namespace for unique keys.
- **Column Family** – very much like a table... but not quite.
- **Key** – a key that represent row (of columns)
- **Column** – representation of value with:
 - Column name
 - Value

DATA MODEL

Users:		CF
ran:		ROW
emailAddress: <u>foo@bar.com</u> ,		COLUMN
webSite: <u>http://bar.com</u>		COLUMN
f.rat:		ROW
emailAddress: <u>f.rat@mouse.com</u>		COLUMN
Stats:		CF
ran:		ROW
visits: 243		COLUMN

DATA MODEL

Users:		CF
ran:		ROW
emailAddress: <u>foo@bar.com</u> ,		COLUMN
webSite: <u>http://bar.com</u>		COLUMN
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Stats:		CF
ran:		ROW
visits: 243		COLUMN

Sparse Columns

DATA MODEL

Songs:

Meir Ariel:

Shir Keev: 6:13 Go wild with column

Tikva: 4:11 names

Erol: 6:17

Suetz: 5:30

Dr Hitchakmut: 3:30

Mashina:

Rakevet Layla: 3:02

Optikai: 5:40

THE API

- Thrift
- CQL
- Client libraries:
 - Hector
 - PyCassa
 - Helenus
 - Aquiles
 - Many more...

READ API

get

get_range_slices

get_slice

get_indexed_slices

get_count

execute_cql_query

multiget

multiget_slice

THE TRUE API

`get(keyspace, key, column_path, consistency)`

`get_slice(ks, key, column_parent, predicate, consistency)`

`multiget(ks, keys, column_path, consistency)`

WRITE API

insert

add

remove

remove_counter

batch_mutate

THE TRUE WRITE API

`insert(..., consistency)`

`add(..., consistency)`

`remove(..., consistency)`

`remove_counter(..., consistency)`

`batch_mutate(..., consistency)`

METADATA API

describe_schema_versions

describe_keyspaces

describe_cluster_name

describe_version

describe_ring

DDL API

system_add_column_family

system_drop_column_family

system_add_keyspace

system_drop_keyspace

system_update_keyspace

system_update_column_family

CQL API

Thrift: execute_cql_query

```
cqlsh> SELECT key, state FROM users;
```

```
cqlsh> INSERT INTO users (key, full_name,  
birth_date, state) VALUES ('bsanderson',  
'Brandon Sanderson', 1975, 'UT');
```

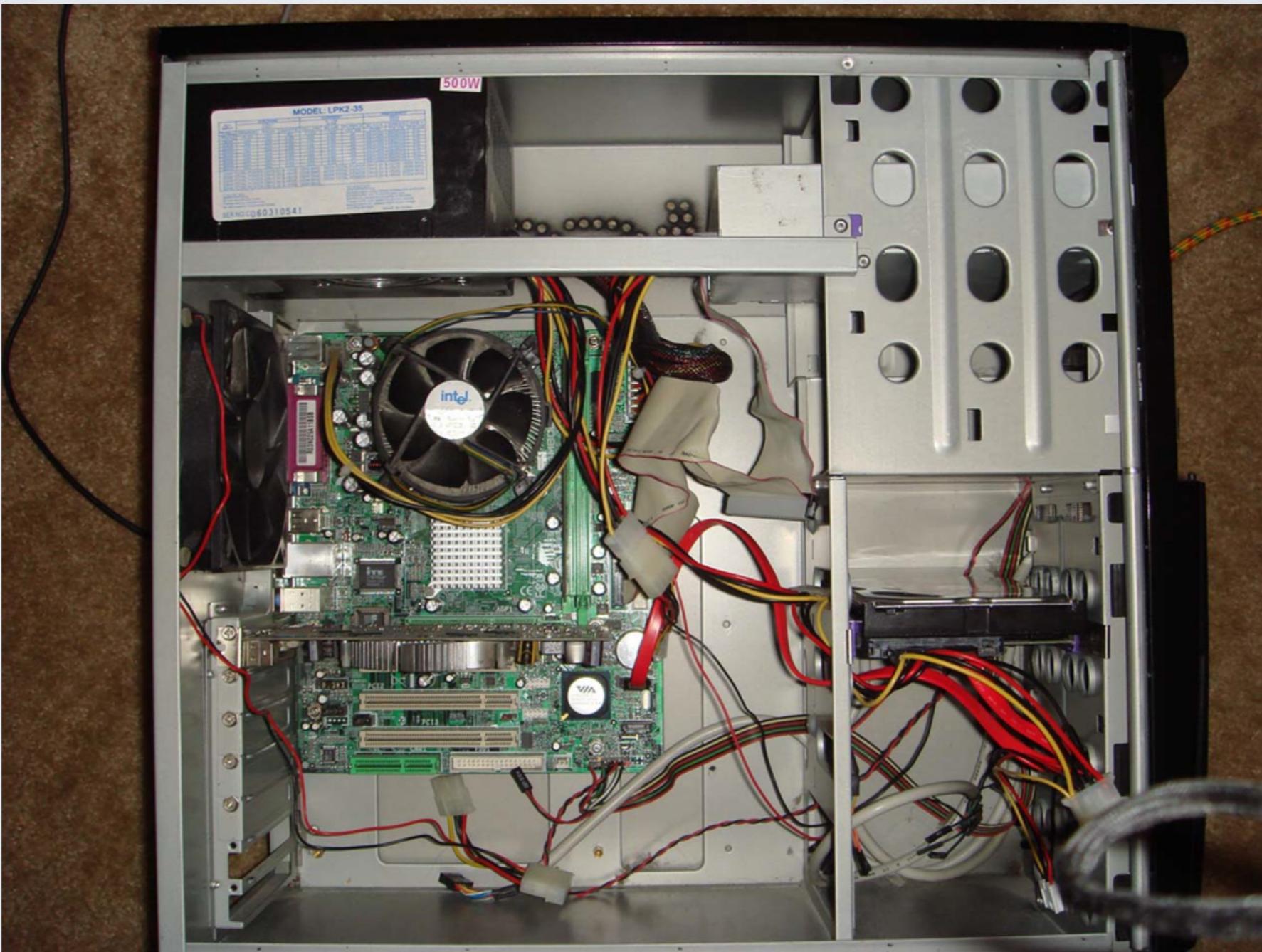
JAVA CODE (HECTOR)

```
/**  
 * Insert a new value keyed by key  
 *  
 * @param key    Key for the value  
 * @param value  the String value to insert  
 */  
public void insert(final String key, final String value) {  
    Mutator m = createMutator(keyspaceOperator);  
    m.insert("key",  
             "columnFamily",  
             createColumn("columnName", value));  
}
```

JAVA CODE (HECTOR)

```
/**  
 * Get a string value.  
 *  
 * @return The string value; null if no value exists for the given key.  
 */  
public String get(final String key) throws HectorException {  
    ColumnQuery<String, String> q = createColumnQuery(keyspaceOperator, serializer, serializer);  
    Result<HColumn<String, String>> r = q.setKey(key).  
        setName("column").  
        setColumnFamily("columnFamily").  
        execute();  
    HColumn<String, String> c = r.get();  
    return c == null ? null : c.getValue();  
}
```

CASSANDRA INTERNALS



AGENDA

- Background and history
- Architectural Layers
- Transport: (Thrift)
- Write Path
- Read Path
- Compactions
- Bloom Filters
- Gossip
- Deletions
- and more...

BACKGROUND

- Required reading:
 - **BigTable** <http://labs.google.com/papers/bigtable.html>
 - **Dynamo** http://www.allthingsdistributed.com/2007/10/amazons_dynamo.html



FROM DYNAMO



- Symmetric **p2p** architecture
- **Gossip** based discovery and error detection
- **Distributed** key-value store
- Pluggable **partitioning**
- Pluggable **topology discovery**
- **Eventual consistent** and Tunable per operation

FROM BIGTABLE



- **Sparse** Column oriented sparse array
- **SSTable** disk storage
- Append-only **commit log**
- **Memtable** (buffering and sorting)
- **Compactions**
- High write performance

ARCHITECTURAL LAYERS

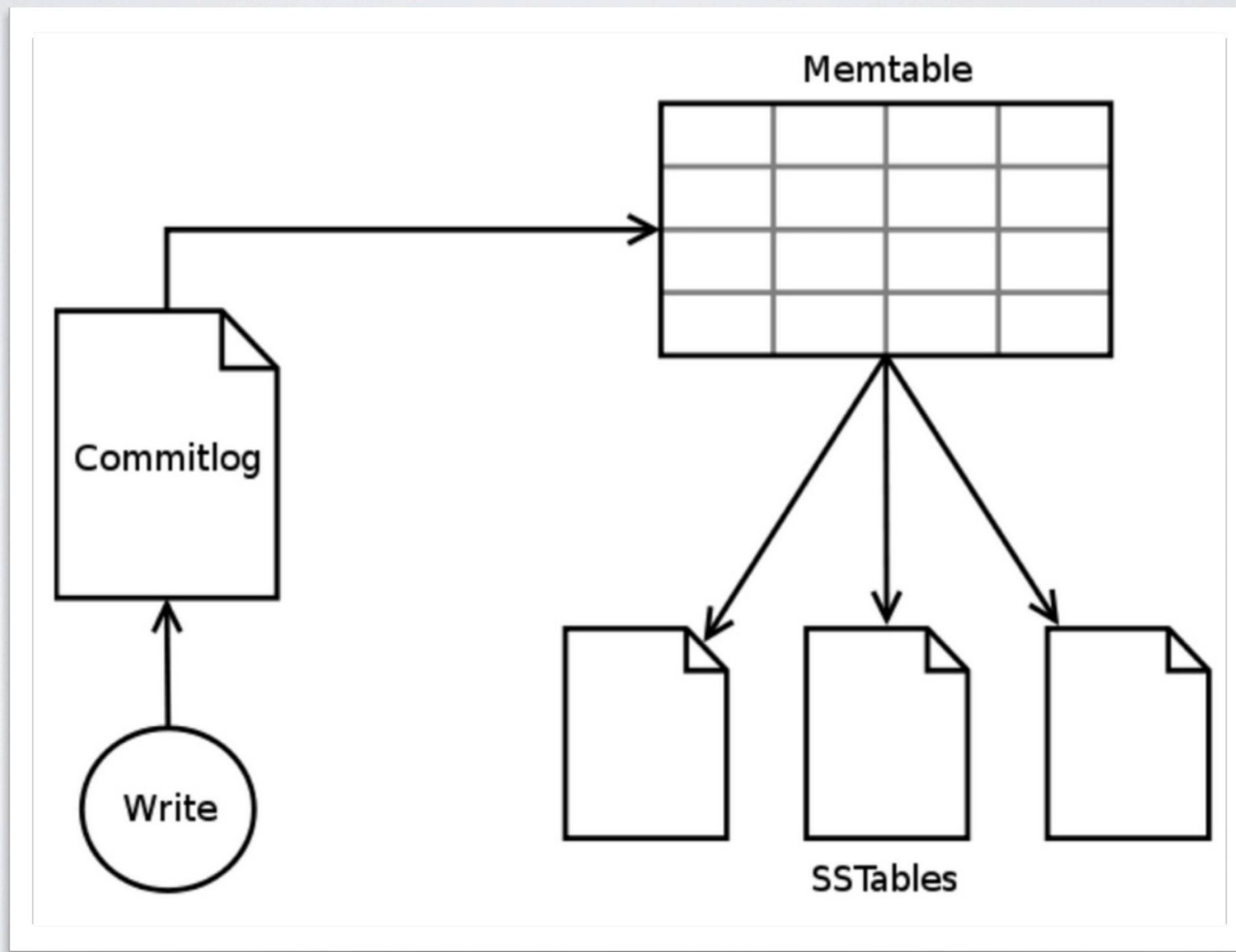
Cluster Management

- Messaging service
- Gossip
- Failure detection
- Cluster state
- Partitioner
- Replication

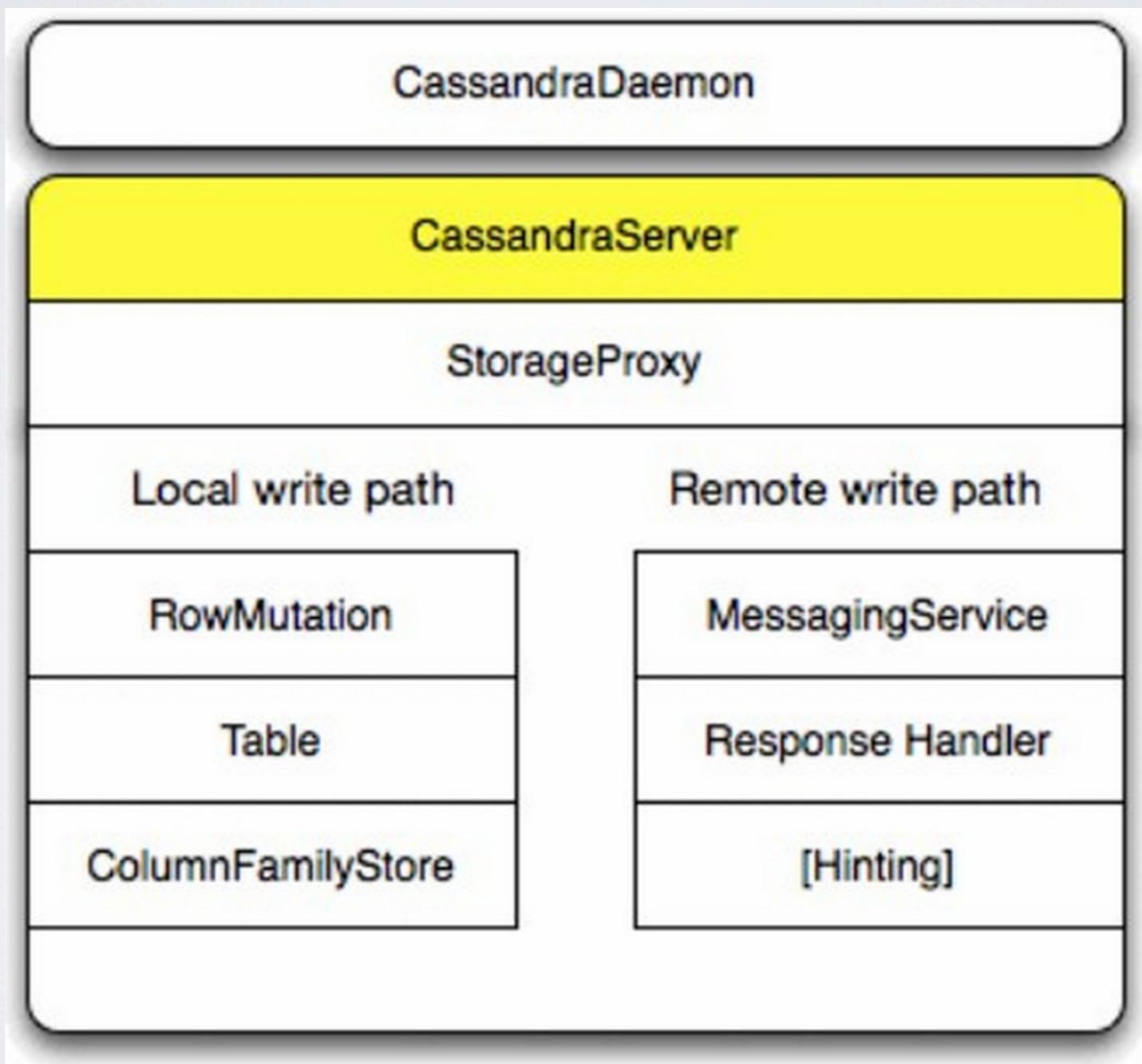
Host Management

- Commit log
- Memtable
- SSTable
- Indexes
- Compaction

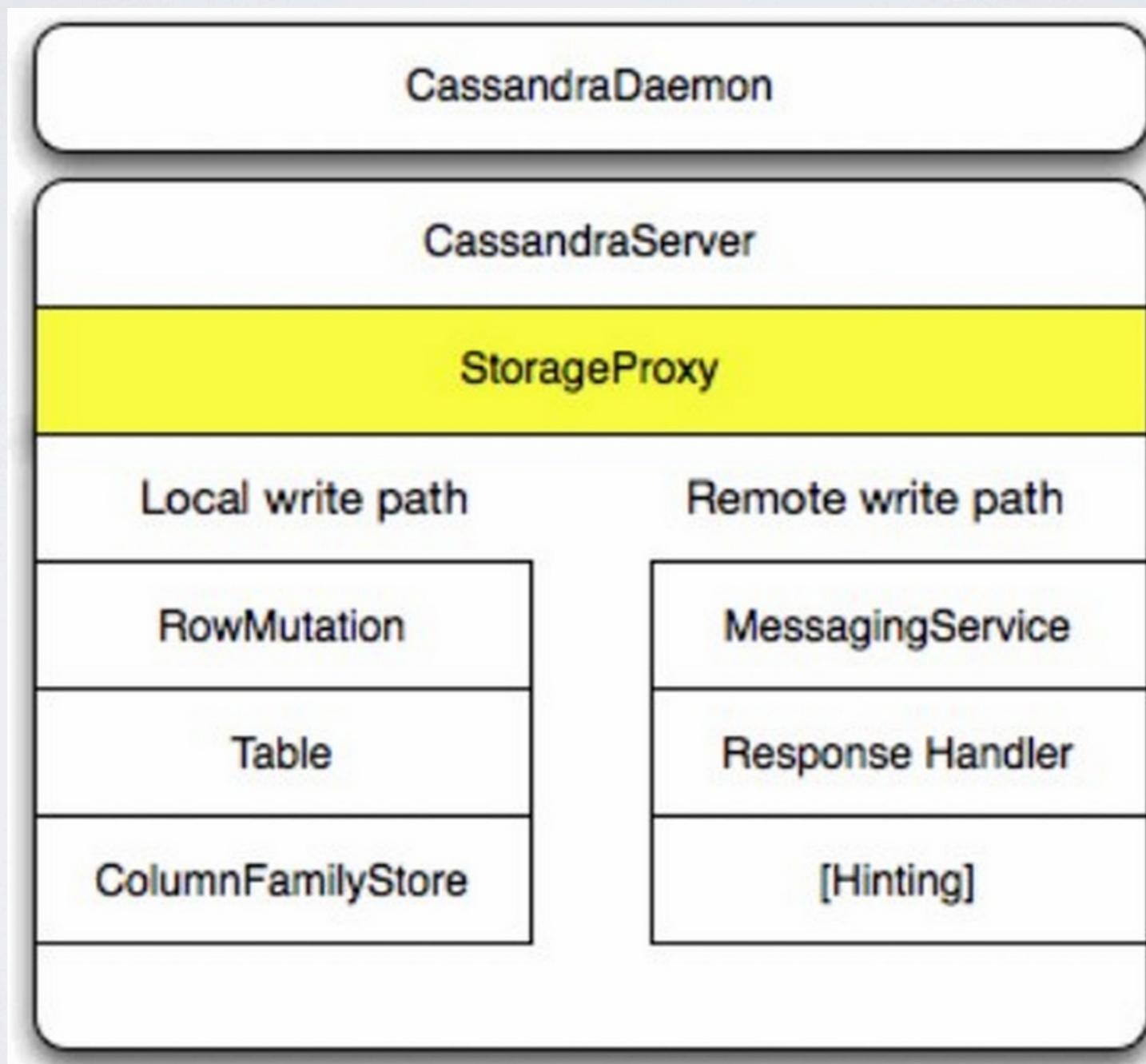
LOCAL WRITE PATH



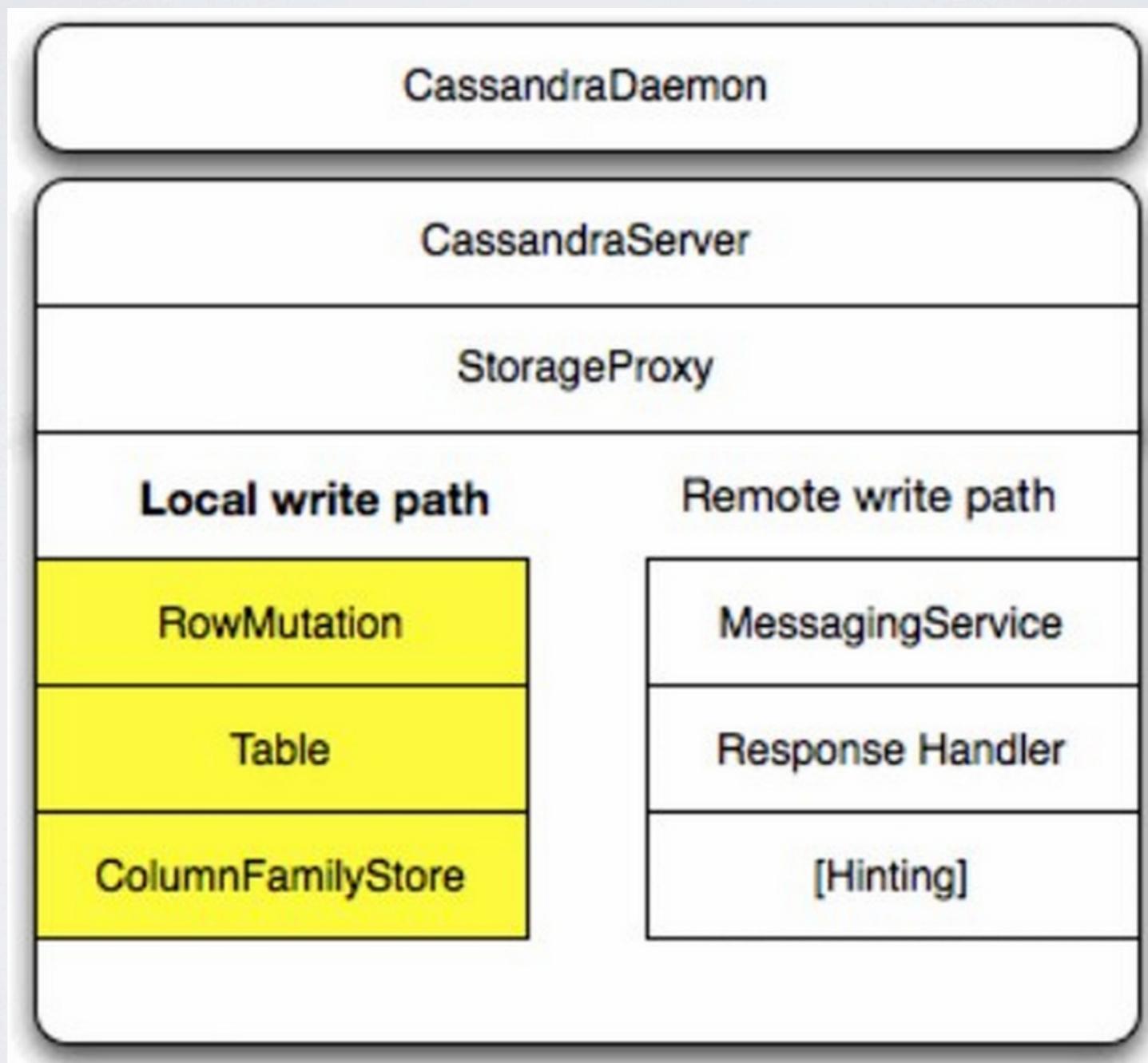
NETWORK WRITE PATH



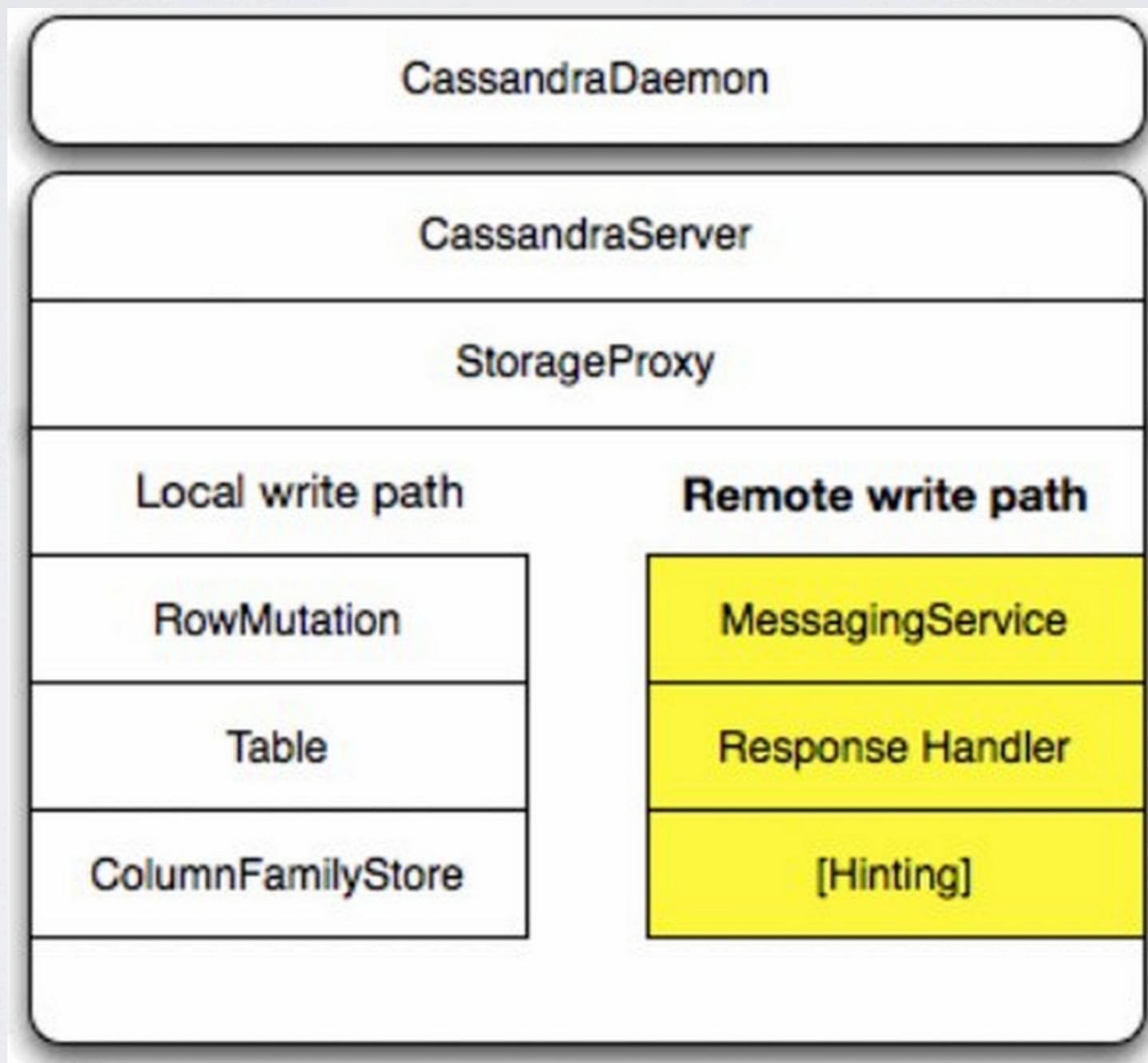
NETWORK WRITE PATH



NETWORK WRITE PATH



NETWORK WRITE PATH



WHAT ARE MEMTABLES?

- In-memory representation of recently written data
- When the table is full, it's sorted and then flushed to disk ⇒ sstable

WHAT ARE SSTABLES?

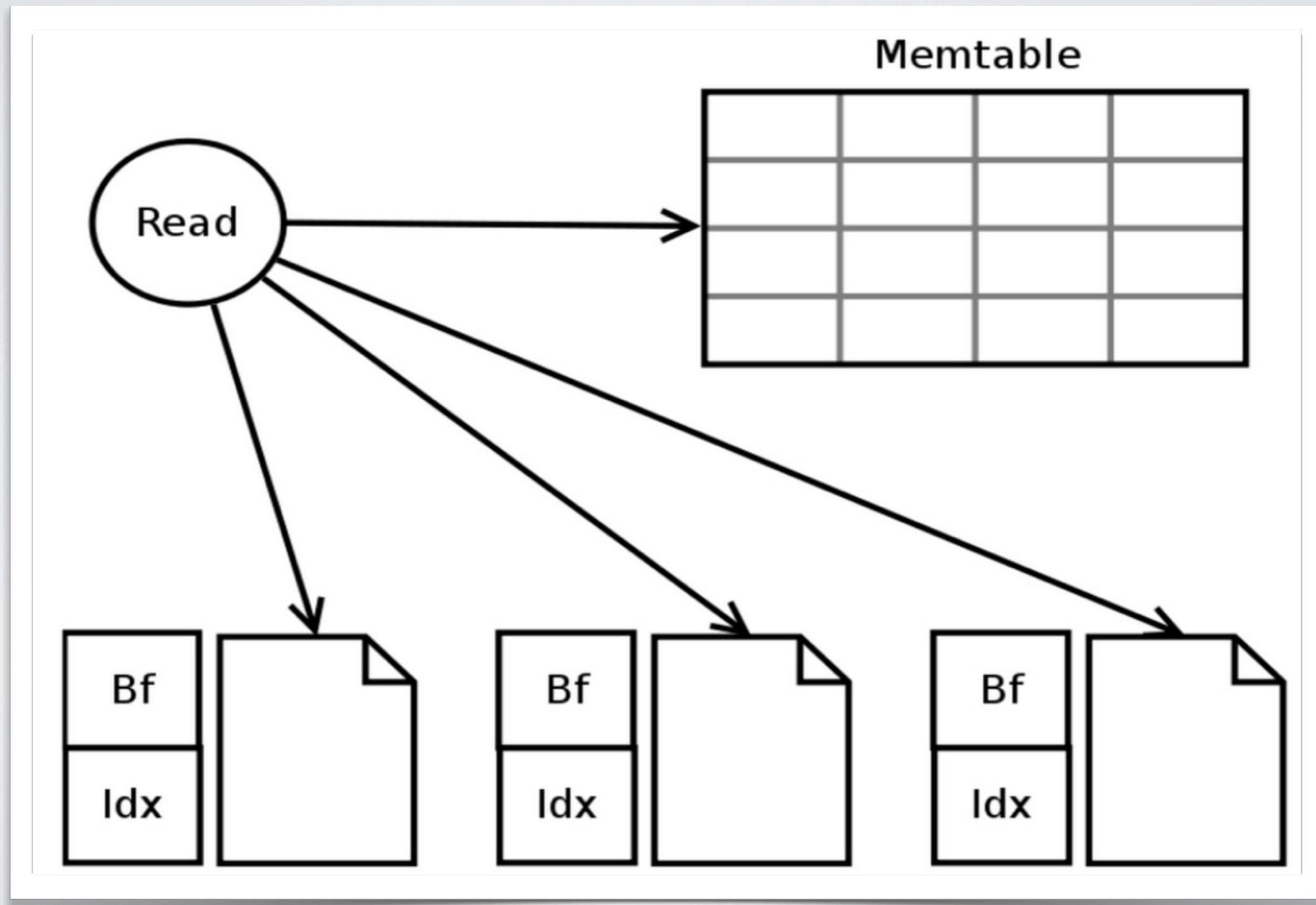
Sorted Strings Tables

- Immutable
- On-disk
- Sorted by a string key
- In-memory index of elements
- Binary search (in memory) to find element location
- Bloom filter to reduce number of unneeded binary searches.

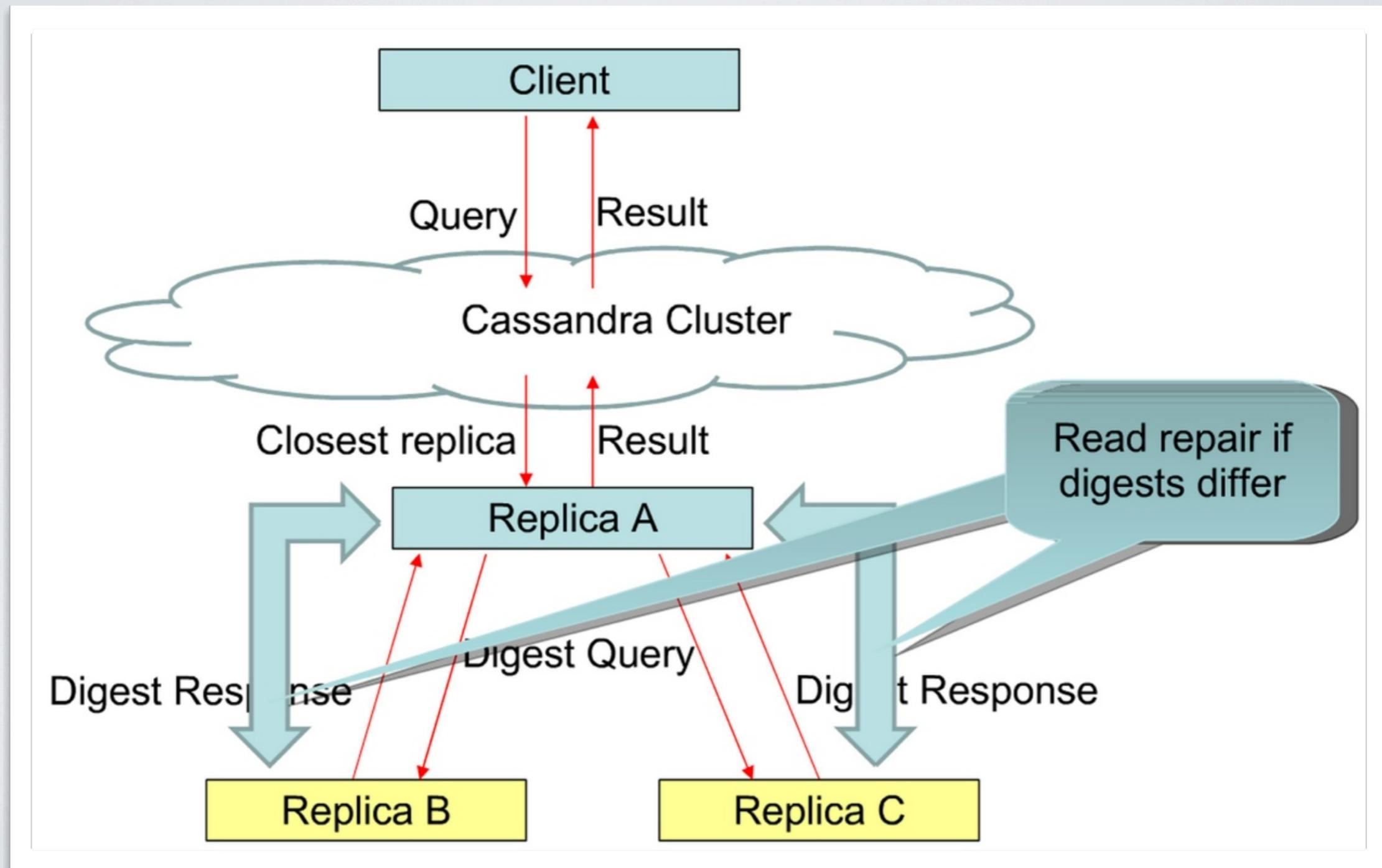
WRITE PROPERTIES

- **No Locks** in the critical path
- **Always available** to writes, even if there are failures.
- **No reads**
- **No seeks**
⇒Fast
- **Atomic** within a Row

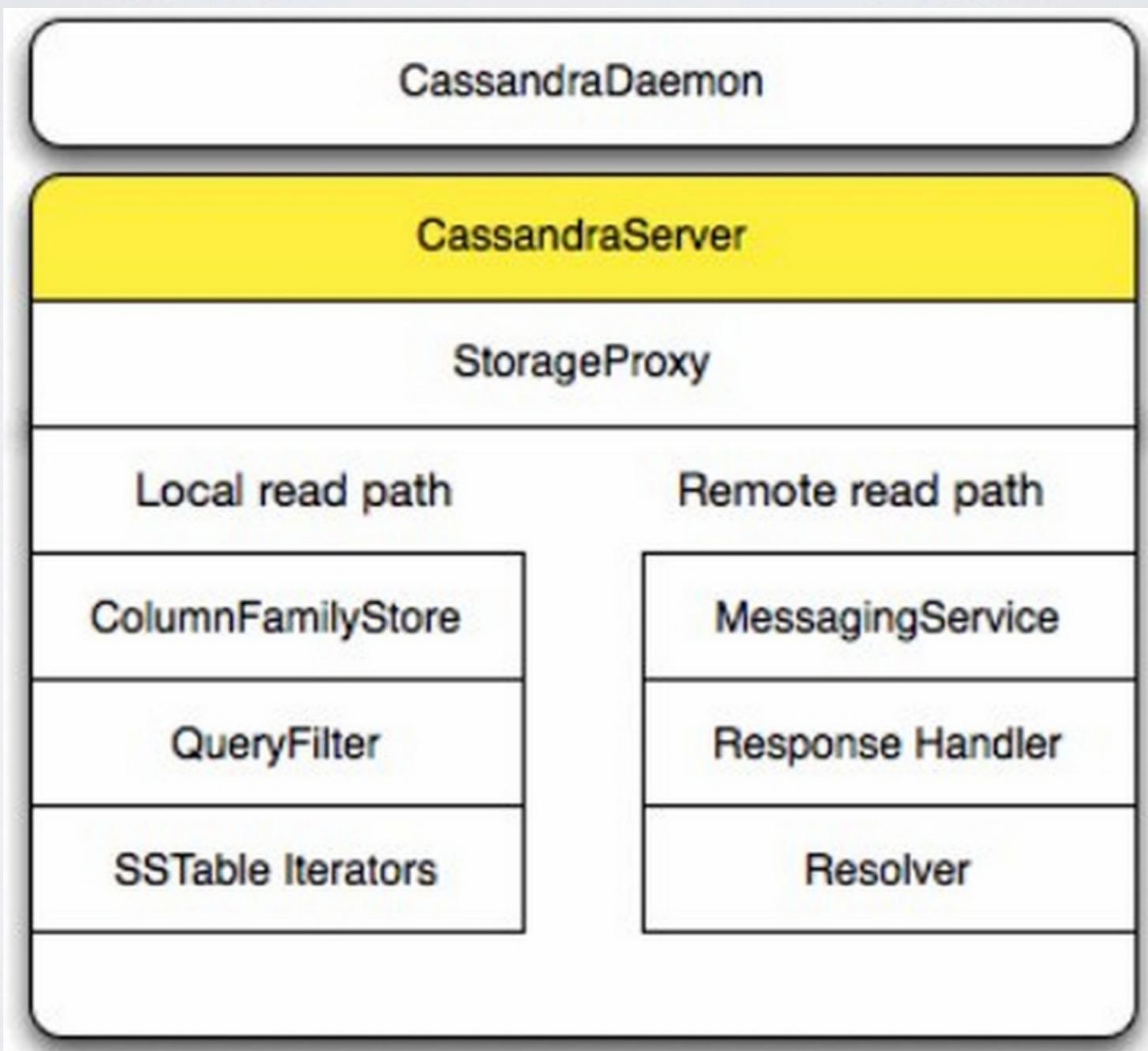
LOCAL READ PATH



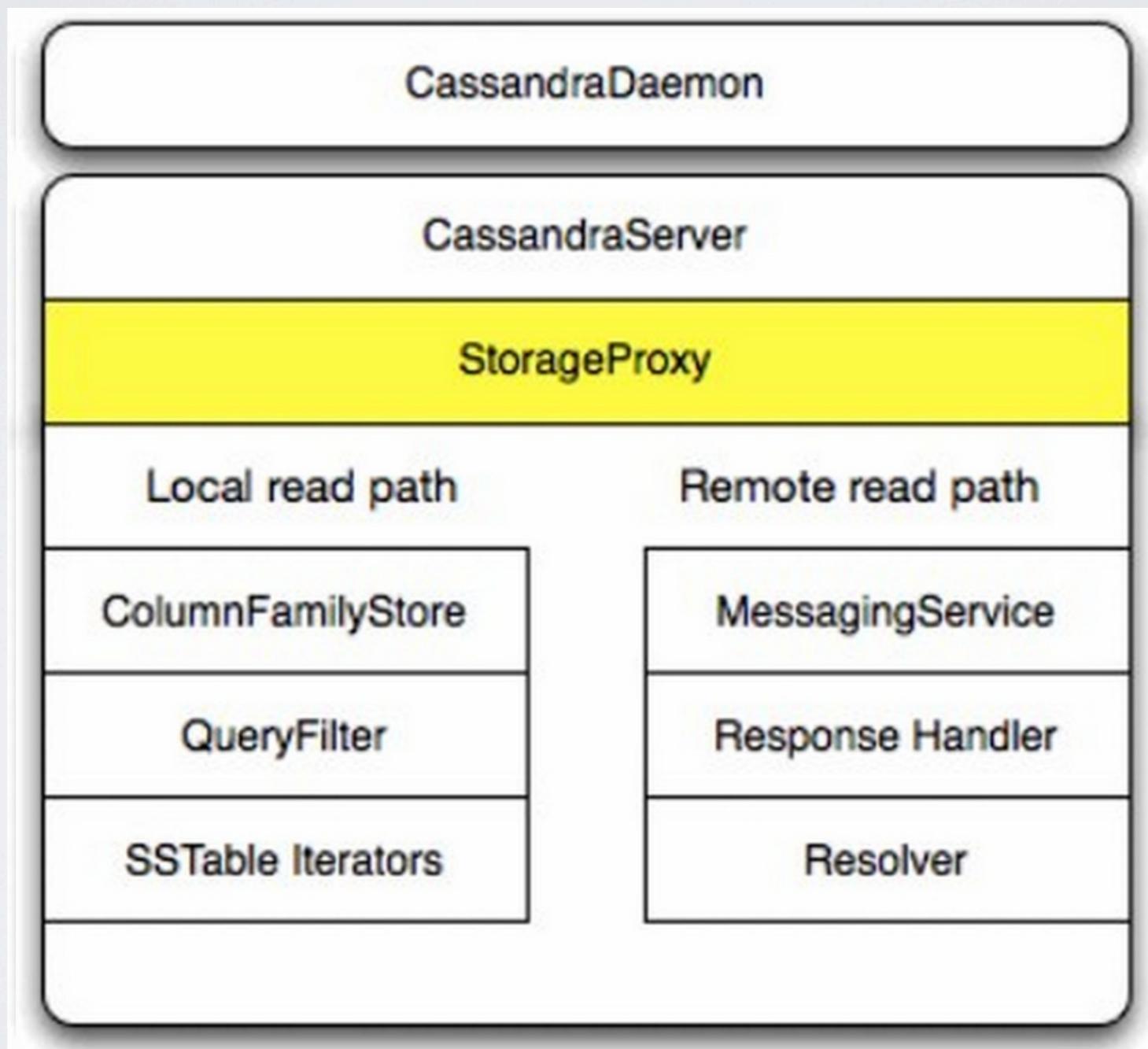
NETWORK READ PATH



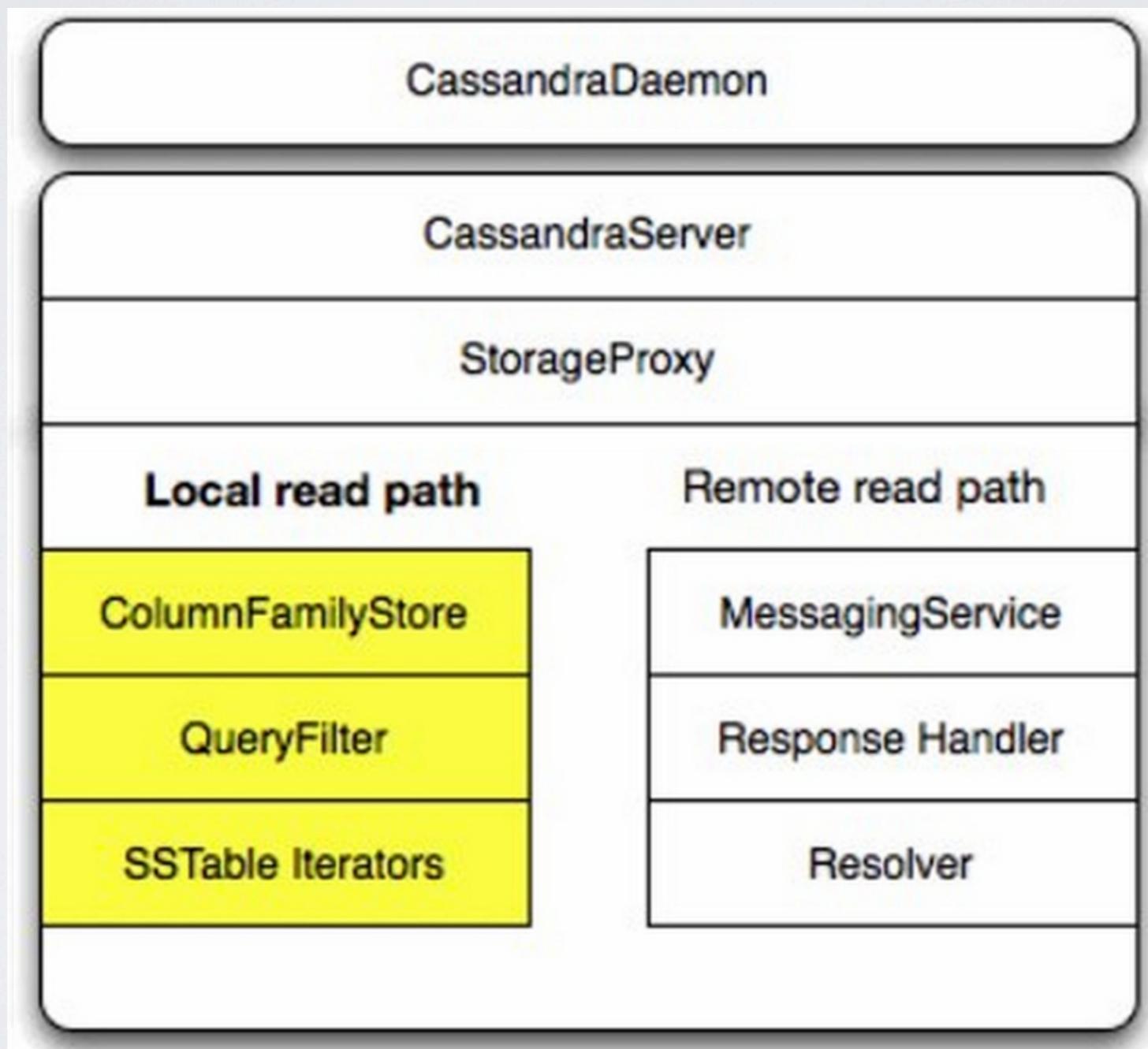
NETWORK READ PATH



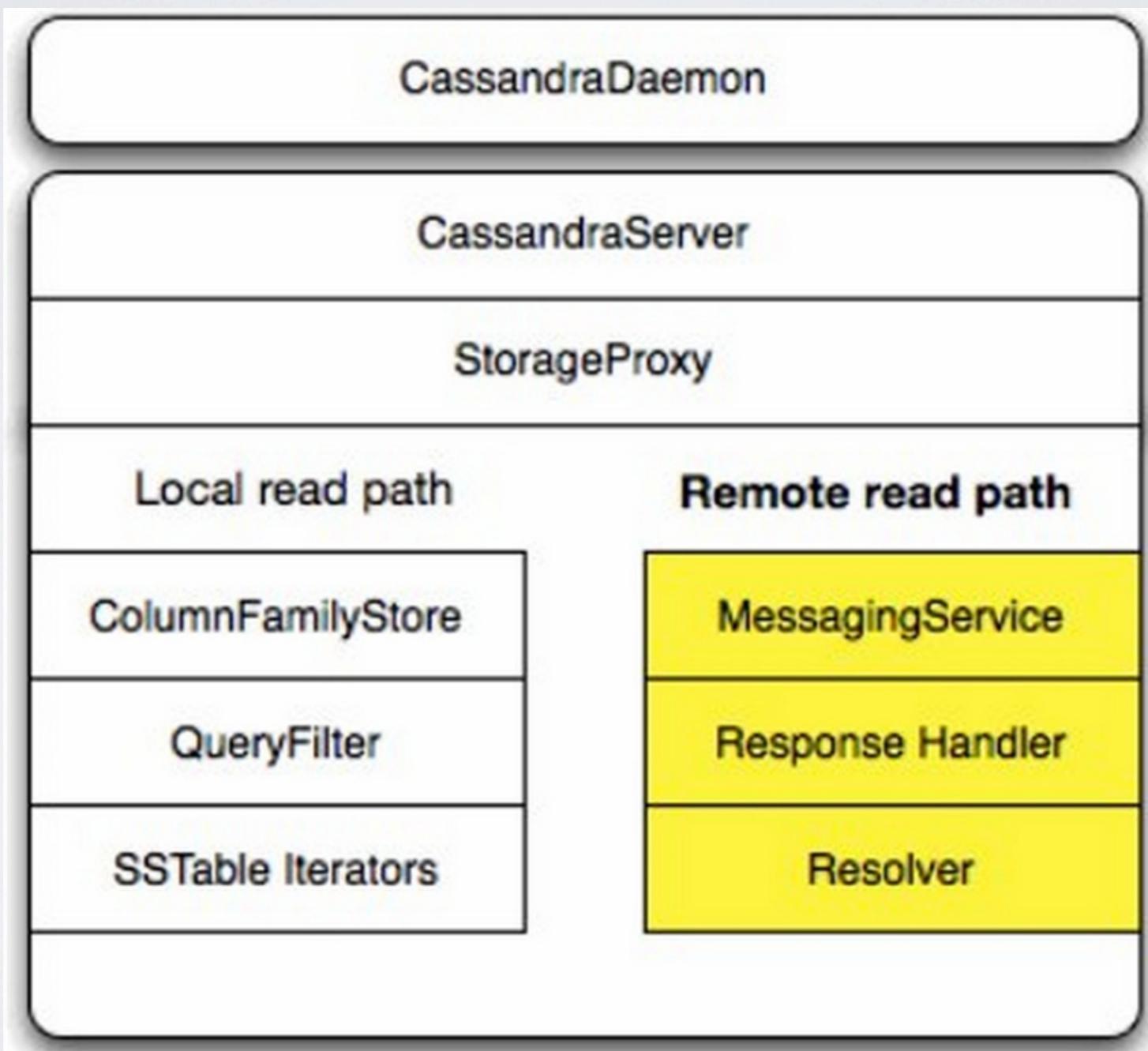
NETWORK READ PATH



NETWORK READ PATH



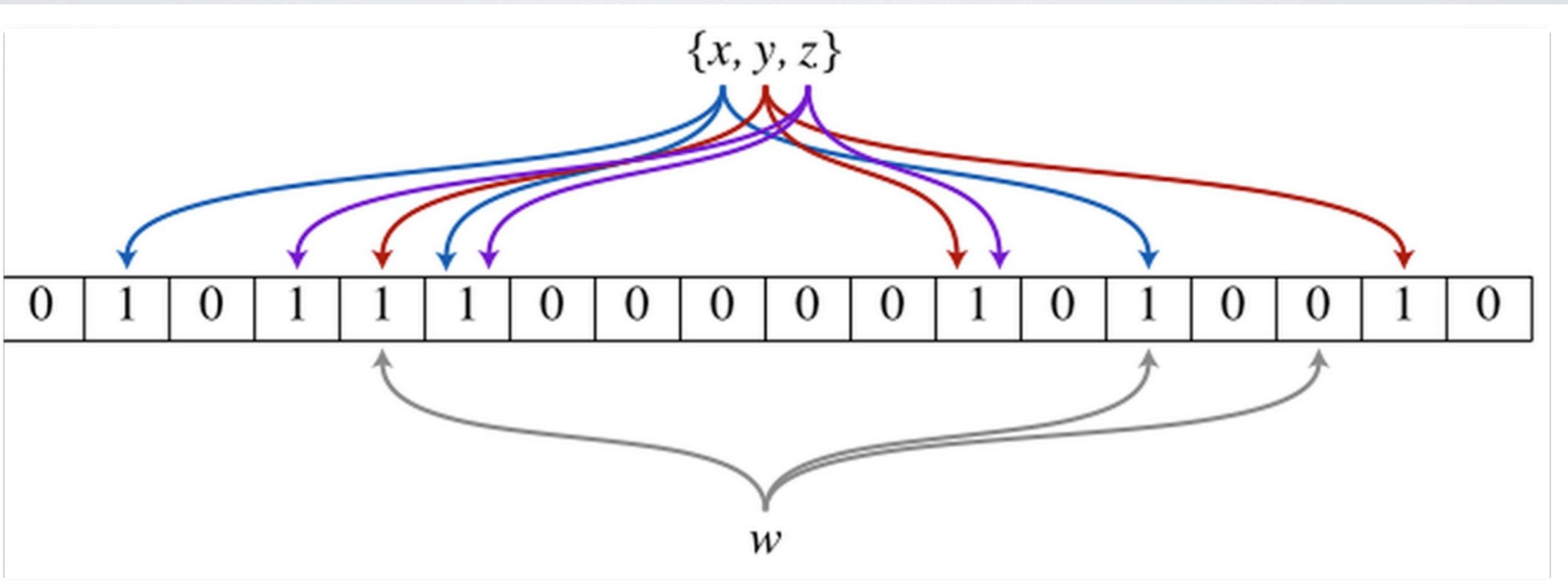
NETWORK READ PATH



READ PROPERTIES

- Read multiple SSTables
- Slower than writes (but still fast)
- Seek can be mitigated with more RAM
- Uses probabilistic bloom filters to reduce lookups.
- Extensive (optional) caching
 - Key Cache
 - Row Cache
- Excellent monitoring

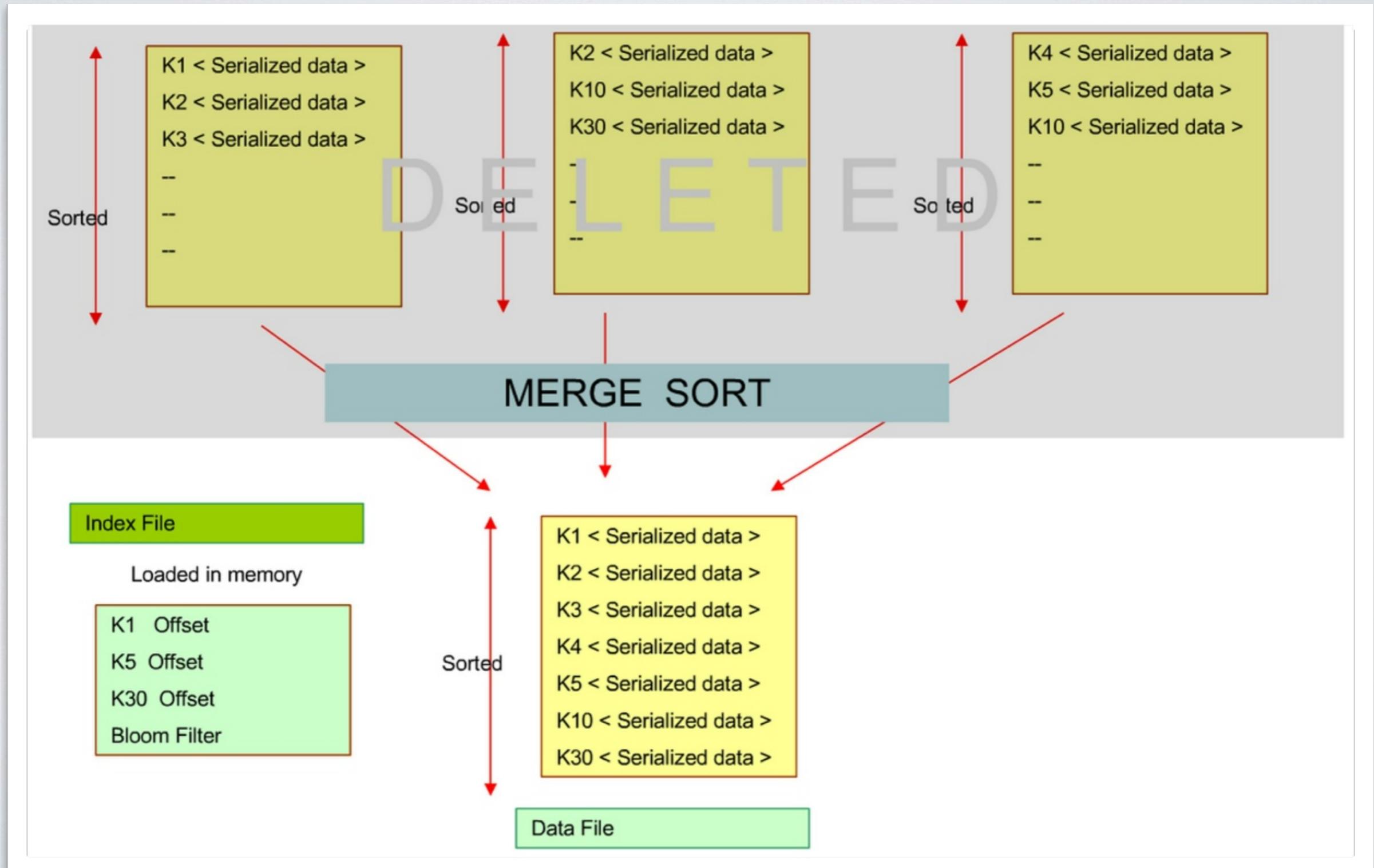
BLOOM FILTERS



BLOOM FILTERS

- Space efficient probabilistic data structure
- Test whether an element is a member of a set
- Allow false positive, but not false negative
- k hash functions
- Union and intersection are implemented as bitwise OR, AND

COMPACTIONS



COMPACTIONS

- Merge keys
- Combine columns
- Discard tombstones
- Merge bloom filters (bitwise OR operation)

GOSSIP



- p2p
- Enables seamless nodes addition.
- Rebalancing of keys
- Fast detection of nodes that goes down.
- Every node knows about all others - no master.

DELETIONS

The problem

1. Node A goes down
 2. Client sends **DELETE x** for data in A
 3. Node B, which is a replica of A accepts the **DELETE x**
 4. Node A goes up again
- ⇒ A thinks B misses x so sends **WRITE x** to B

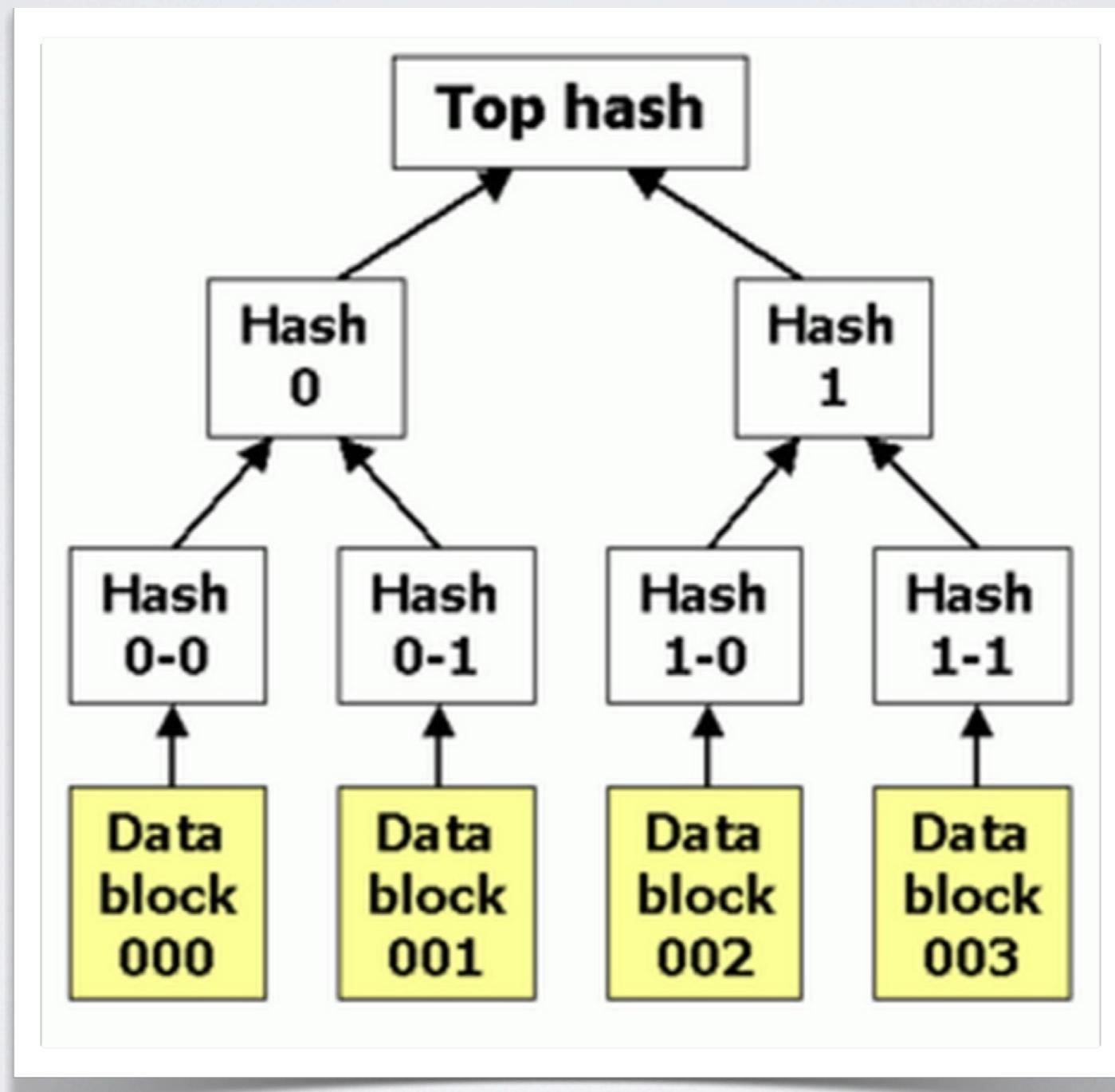
DELETIONS



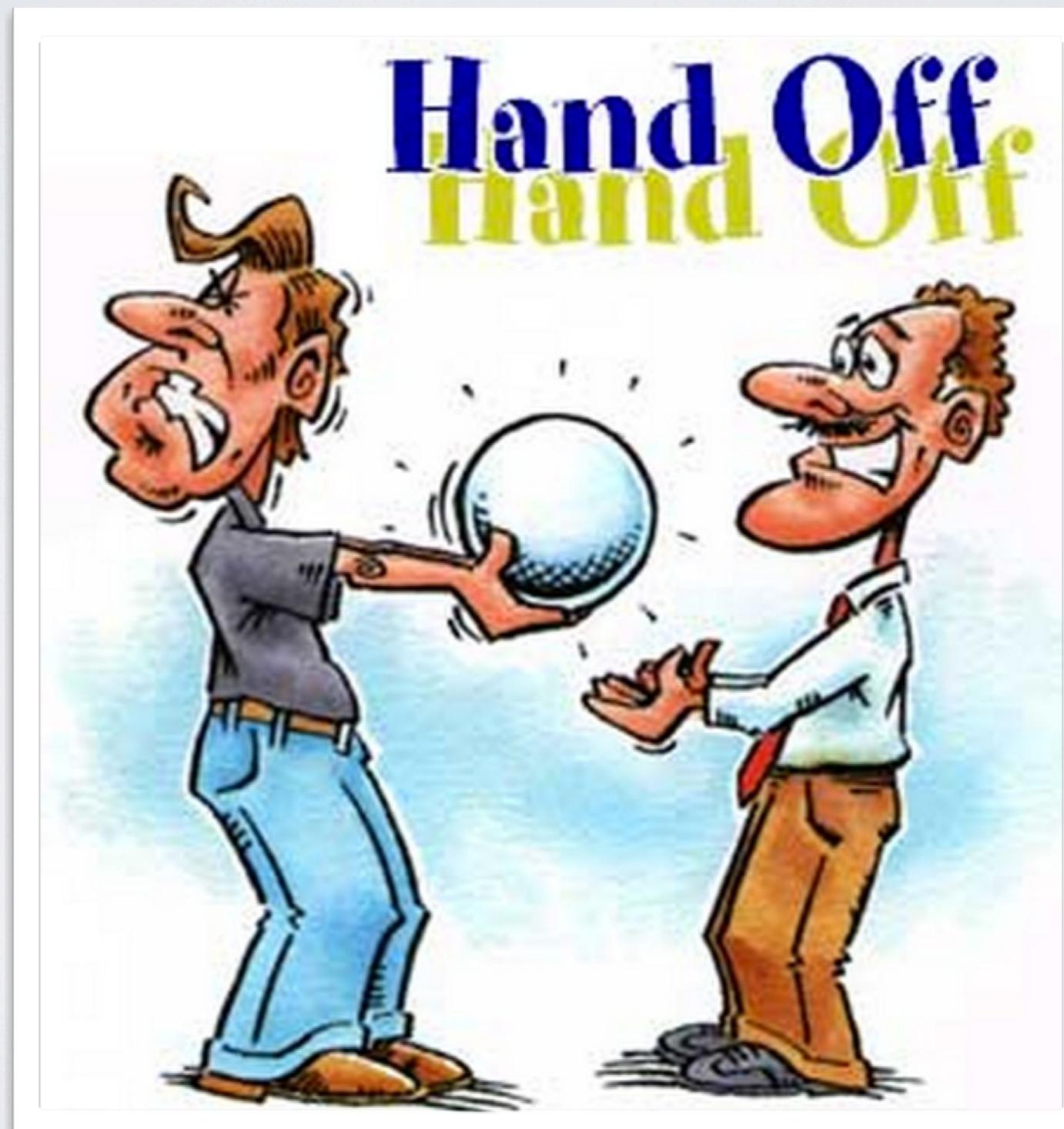
The solution

- Deletion marker (tombstone) necessary to suppress data in older SSTables, until compaction
- Read repair complicates things a little
- Eventually consistent complicates things more
- Solution: configurable delay before tombstone GC, after which tombstones are not repaired

ANTI ENTROPY SERVICE AND MERKLE TREES



HINTED HANDOFF



REFERENCES

- This presentation: <https://speakerdeck.com/u/rantav/p/nosql-and-cassandra-at-ibm>
- <http://highscalability.com/blog/2009/11/5/a-yes-for-a-nosql-taxonomy.html>
- <http://en.wikipedia.org/wiki/NoSQL#Taxonomy>
- http://en.wikipedia.org/wiki/Apache_Cassandra
- <http://www.datastax.com/dev/blog/whats-new-in-cassandra-1-0-performance>
- <http://wiki.apache.org/cassandra/ClientOptions>

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- <http://bret.appspot.com/entry/how-friendfeed-uses-mysql>
- <http://www.julianbrowne.com/article/viewer/brewers-cap-theorem>
- http://www.allthingsdistributed.com/2008/12/eventually_consistent.html
- <http://wiki.apache.org/cassandra/DataModel>

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