

## Simplification of a Context-free Grammar (certain definitions)

Let  $G = (N, \Sigma, R, S)$  be a CFG.

A symbol  $X$  in  $N \cup \Sigma$  is called generating if  $X \xrightarrow{*} w$  for some  $w$  in  $\Sigma^*$ .

A symbol  $X$  in  $N \cup \Sigma$  is called reachable if  $S \xrightarrow{*} \alpha X \beta$  for some  $\alpha, \beta$  in  $(N \cup \Sigma)^*$ .

Generating and reachable symbols are called useful and other symbols are called useless.

A production of the form  $A \rightarrow \epsilon$  is called an epsilon-production.

A production of the form  $A \rightarrow B$  is called a unit production.

A CFG is called simple if it has no epsilon-production, no unit production and no useless symbols.

Given a CFG  $G$  we can simplify it ie we can obtain a simple grammar  $G'$  such that  $L(G') = L(G) - \{\epsilon\}$ .

For simplification the following three steps have to be carried out in that order.

- 1) Elimination of epsilon productions.
- 2) Elimination of unit productions.
- 3) Elimination of useless symbols.