

## Elimination of useless symbols

Given a CFG  $G$  we can find an equivalent CFG  $G_1$  such that  $G_1$  has no useless symbols.

We have to first find the set of generating symbols  $\mathfrak{G}$ .

For this we carry out the procedure.

1)  $\mathfrak{G} \leftarrow \Sigma$  since  $a \rightarrow^* a$  for every  $a$  in  $\Sigma$ .

2) if  $X \rightarrow X_1 X_2 \dots X_k$  and  $X_i$  are in  $\mathfrak{G}$  then insert  $X$  in  $\mathfrak{G}$ .

3) Repeat 2) until no more insertions can be made.

$\mathfrak{G}$  gives the set of generating symbols. Other symbols can be dropped.

To identify  $R$ , the set of reachable symbols in  $G$  we carry out the following steps.

1)  $R = \{S\}$

2) If  $A$  is in  $R$  and  $A \rightarrow X_1 X_2 \dots X_k$ , insert  $X_1, X_2, \dots, X_k$  in  $R$ .

3) Repeat 2) until no more insertions can be made.

$R$  is the set of all reachable symbols. Other symbols can be dropped.

Example :  $S \rightarrow a A a \mid a B b$

$A \rightarrow a S \mid b D$

$B \rightarrow a B a \mid b$

$C \rightarrow a b b \mid D D$

$D \rightarrow a D a$

Here  $a, b$  are generating. Hence  $B, C$  are generating. Hence  $S$  is generating and therefore  $A$  is generating. So  $D$  is non-generating.

Getting rid of  $D$  we get :  $S \rightarrow a A a \mid a B b$

$A \rightarrow a S$

$B \rightarrow a B a \mid b$

$$C \rightarrow a b b$$

S is reachable. Hence a, A, B and b are reachable. C is not reachable.

So getting rid of C we get an equivalent grammar with no useless symbols.  $S \rightarrow a A a \mid a B b$

$$A \rightarrow a S$$
$$B \rightarrow a B a \mid b$$

Note that we have to eliminate the non-generating symbols first. For example if we start with the grammar

$$S \rightarrow A B \mid a$$
$$B \rightarrow b$$

and get rid of the unreachable symbols first we find that all the symbols are reachable. After that if we get rid of the non-generating symbols, A will get dropped and we will be left with the grammar

$$S \rightarrow a$$
$$B \rightarrow b$$

which still has B unreachable. However if we get rid of the non-generating symbols first, we get the grammar

$$S \rightarrow a$$
$$B \rightarrow b$$

Here B is unreachable. Dropping B we get the grammar

$$S \rightarrow a$$

which is equivalent to the original grammar and has no useless symbols.

Points to note :

We saw that we can determine whether  $\epsilon$  is generated by a grammar G by checking whether S is nullable.

Similarly we can determine whether  $L(G)$  is non-empty by checking whether  $S$  is generating.