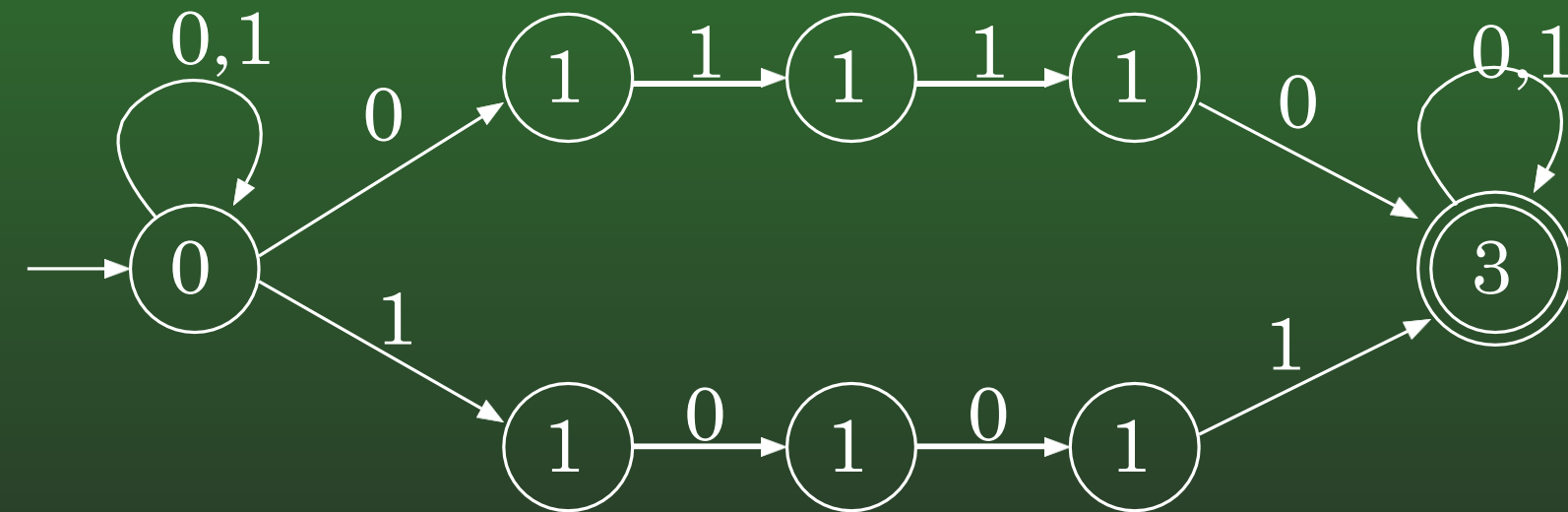


# ***Deterministic Finite Automata vs. Non-Deterministic Finite Automata***

# NFA

Create an NFA for:

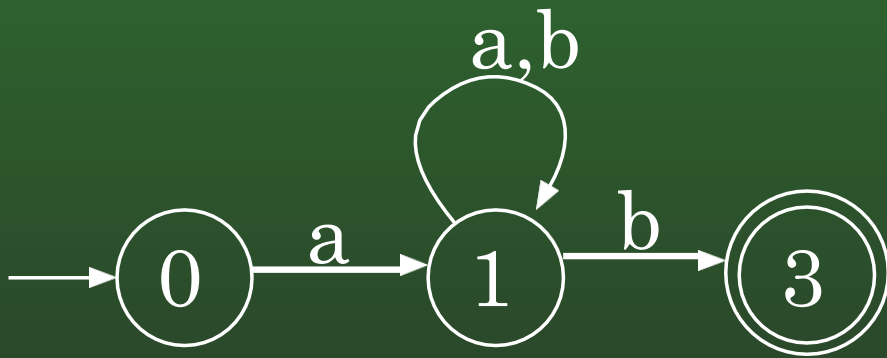
- All strings over  $\{0, 1\}$  that contain 0110 or 1001



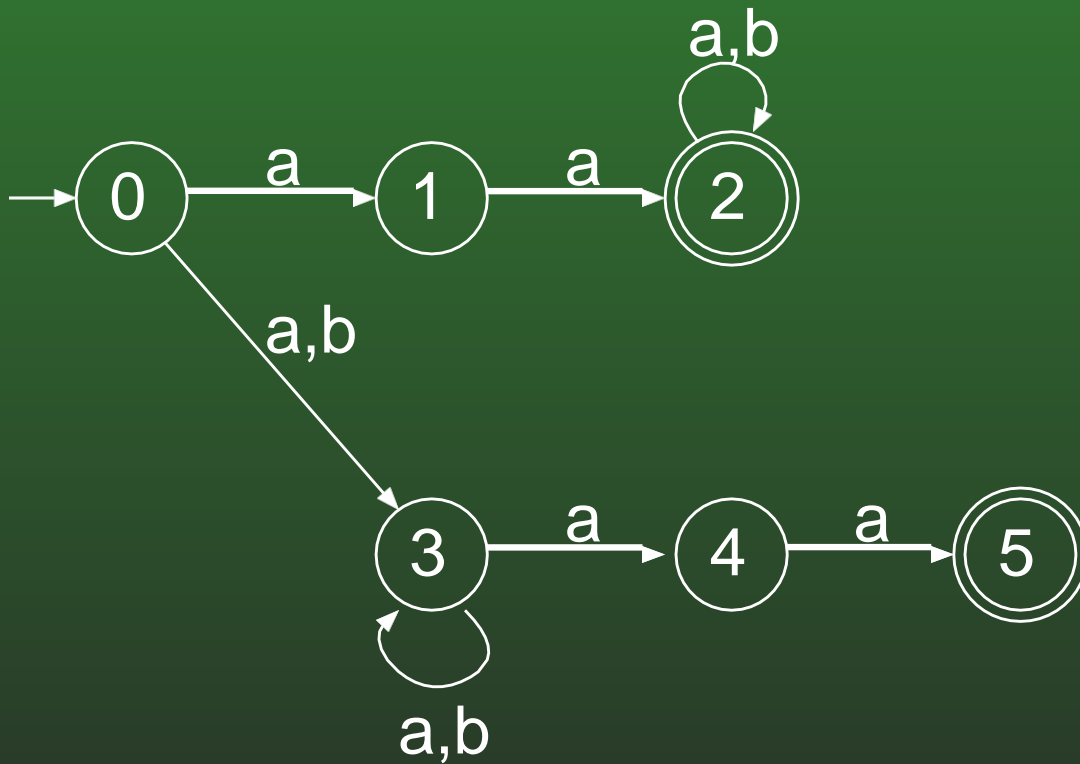
# NFA

Create an NFA for:

- All strings over  $\{a, b\}$  that start with  $a$  and end with  $b$

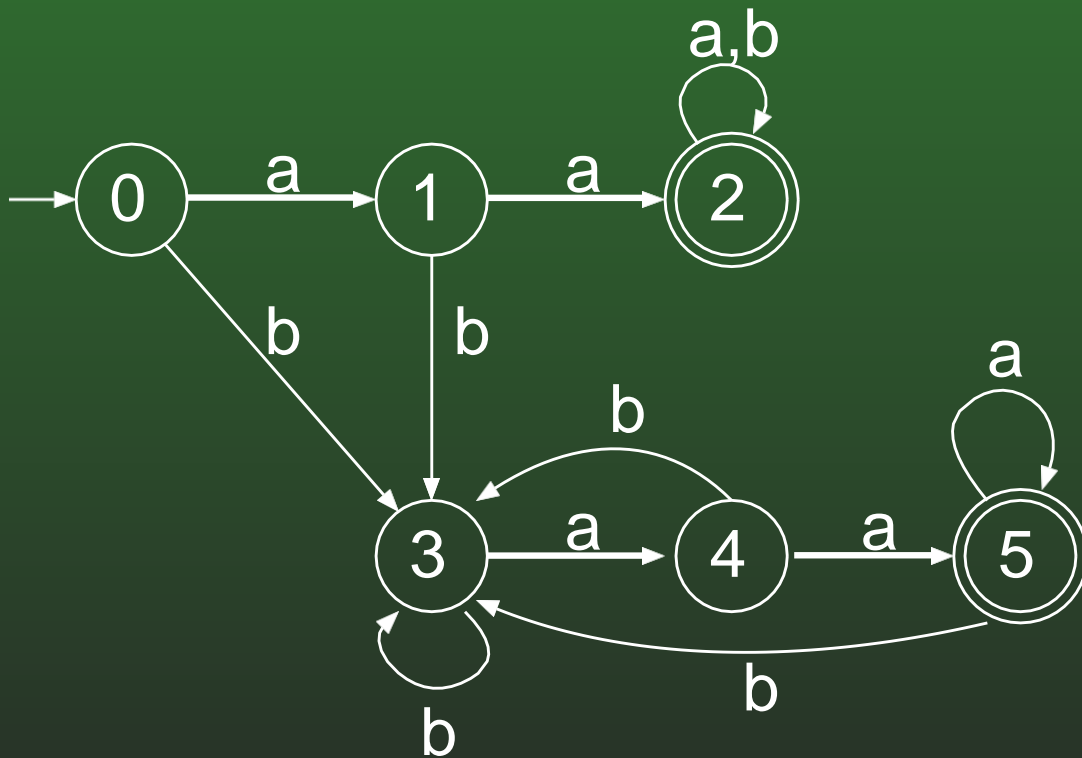


- All strings over  $\{a,b\}$  that begin or end with  $aa$



# NFA $\rightarrow$ DFA

- Can we create a DFA for the same language?
- All strings over  $\{a,b\}$  that begin or end with  $aa$



# $L_{NFA}$ **vs** $L_{DFA}$

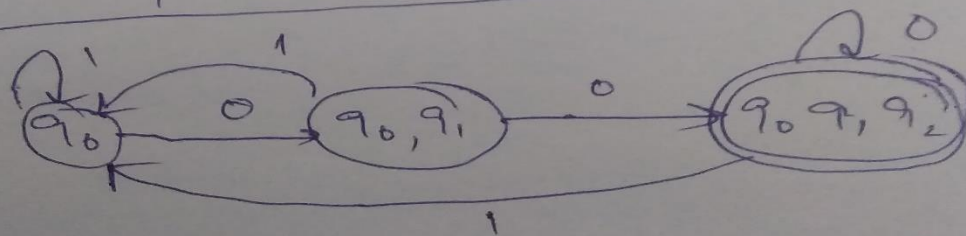
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- What is the relationship between  $L_{NFA}$  and  $L_{DFA}$ ?
  - $L_{DFA} \subseteq L_{NFA} \wedge L_{NFA} \subseteq L_{NFA} (L_{NFA} = L_{DFA})$
- Given any NFA  $M$ , we *can* create a DFA  $M'$  such that  $L[M] = L[M']$

create a NFA over  $\Sigma\{0,1\}$  that accept all strings ends with '00' and convert

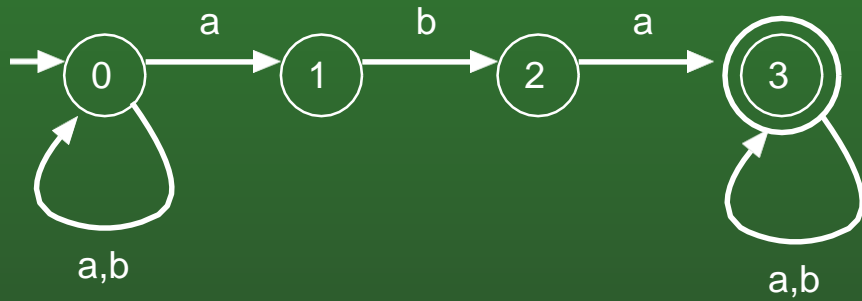


	0	1
q <sub>0</sub>	[q <sub>0</sub> , q <sub>1</sub> ]	[q <sub>0</sub> ]
[q <sub>0</sub> , q <sub>1</sub> ]	[q <sub>0</sub> , q <sub>1</sub> , q <sub>2</sub> ]	[q <sub>0</sub> ]
[q <sub>0</sub> , q <sub>1</sub> , q <sub>2</sub> ]	[q <sub>0</sub> , q <sub>1</sub> , q <sub>2</sub> ]	[q <sub>0</sub> ]



# NFA $\rightarrow$ DFA

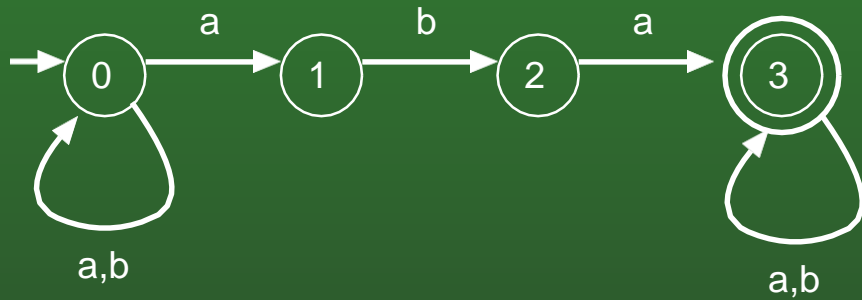
- NFA for all strings over  $\{a,b\}$  containing aba



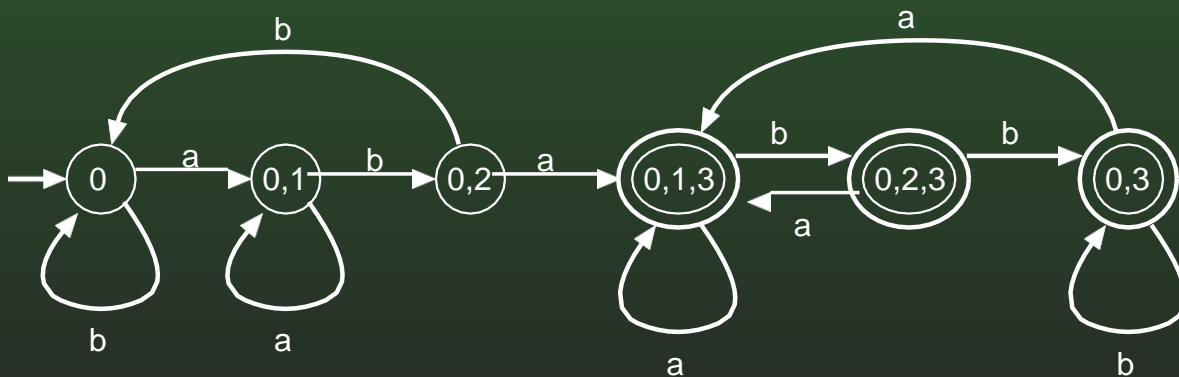


# NFA $\rightarrow$ DFA

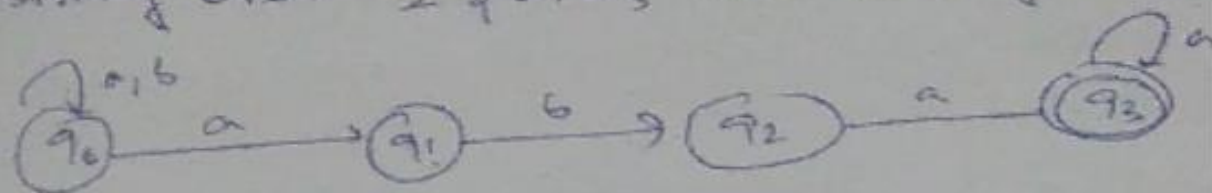
- NFA for all strings over  $\{a,b\}$  containing aba



- Build Equivalent DFA



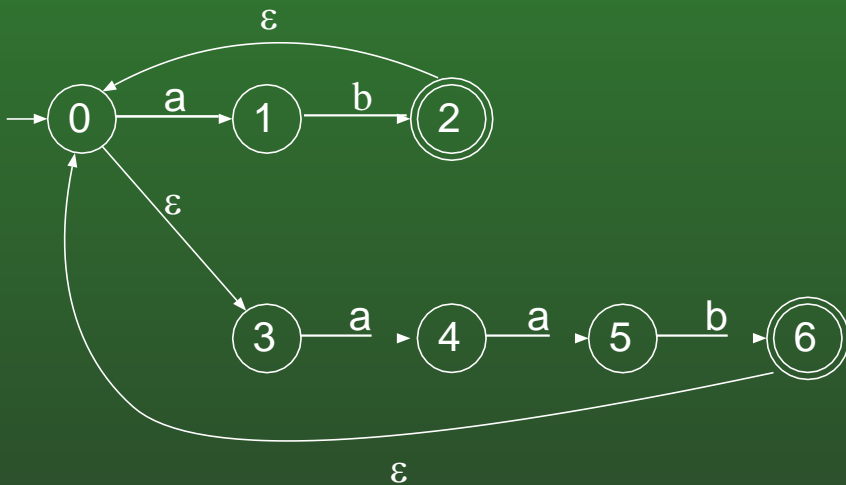
Create DFA ~~which~~ ~~that~~ accepts  
all strings over  $\Sigma \{a, b\}$  with substrings 'aba'



	a	b
q <sub>0</sub>	q <sub>0</sub> q <sub>1</sub>	q <sub>0</sub>
q <sub>0</sub> q <sub>1</sub>	q <sub>0</sub> q <sub>1</sub>	q <sub>0</sub> q <sub>2</sub>
q <sub>0</sub> q <sub>2</sub>	q <sub>0</sub> q <sub>1</sub> q <sub>3</sub>	q <sub>0</sub>
q <sub>0</sub> q <sub>1</sub> q <sub>3</sub>	q <sub>0</sub> q <sub>1</sub> q <sub>3</sub>	q <sub>0</sub> q <sub>2</sub> q <sub>3</sub>
q <sub>0</sub> q <sub>2</sub> q <sub>3</sub>	q <sub>0</sub> q <sub>1</sub> q <sub>3</sub>	q <sub>0</sub> q <sub>3</sub>
q <sub>0</sub> q <sub>3</sub>	q <sub>0</sub> q <sub>1</sub> q <sub>3</sub>	q <sub>0</sub> q <sub>3</sub>

# NFA with $\epsilon$

- What about  $\epsilon$  transitions?



$$\delta: (Q \times \{\Sigma \cup \epsilon\}) \rightarrow P(Q)$$

Any FA with  $\epsilon$  transitions is always NFA  
DFA can never have  $\epsilon$  transitions

# Minimization of DFA

## Example:

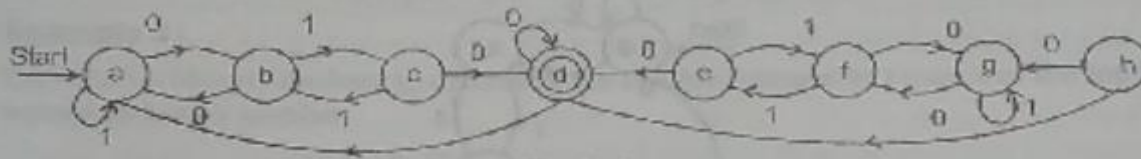


Fig 2.16

## Solution:

### Step 1:

First states are divided into final states group and non-final state group.

### Step 2:

Now the i/p is applied on the states

### Step 3:

Again group the similar states as another group.

### Step 4:

Repeat step 2 and 3 until we get the similar group.

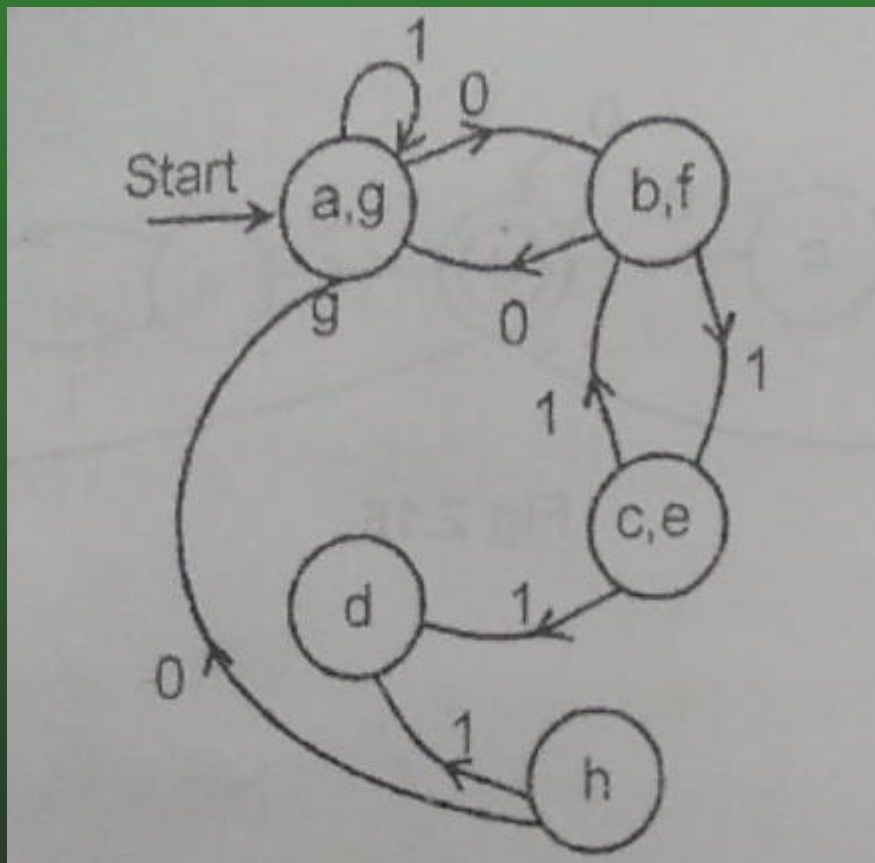
Step 1. G0 (a, b, c, d, e, f, g, h)

Step 2. G01 (a, b, c, e, f, g, h)  
G02 (d)

Step 3. G011 (a, b, g, f)  
G012 (c, e)  
G02 (d) G013 (h)

Step 4. G0111 (a, g)  
G0112 (b, f)  
G012 (c, e)  
G02 (d) G013 (h)

Step 5. G0111 (a, g)  
G0112 (b, f)  
G012 (c, e)  
G02 (d) G013 (h)



Step 1. G0 (a, b, c, d, e, f, g, h)

Step 2. G01 (a, b, c, e, f, g, h)  
G02 ( d)

Step 3. G011 (a, b, g, f )  
G012 (c, e)  
G02 ( d) G013 (h)

Step 4. G0111 (a, g)  
G0112 (b, f)  
G012 (c, e)  
G02 ( d) G013 (h)

Step 5. G0111 (a, g)  
G0112 (b, f)  
G012 (c, e)  
G02 ( d) G013 (h)