
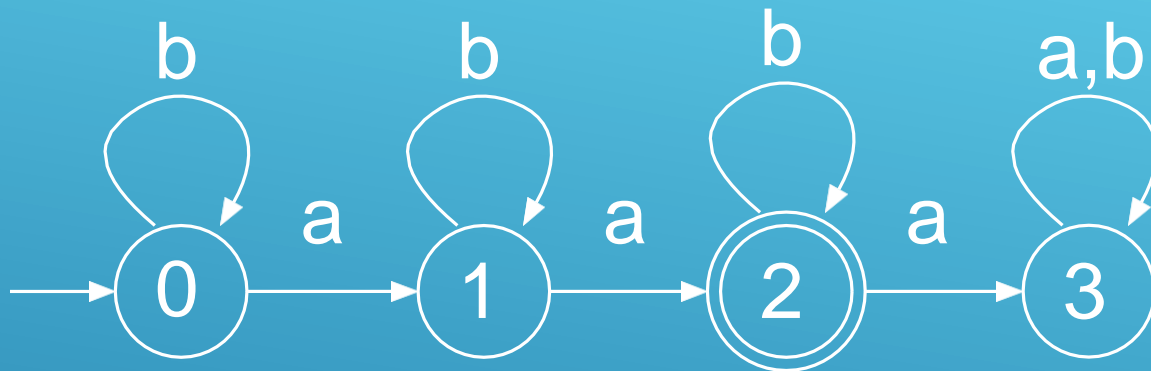


DFA DEFINITION


- A DFA is a 5-tuple $M = (Q, \Sigma, \delta, q, F)$
 - Q Set of states
 - Σ Alphabet
 - $\delta : (Q \times \Sigma) \rightarrow Q$ is a Transition function
 - $q \in Q$ Initial state
 - $F \subseteq Q$ Set of final states
- 
- A series of four parallel white diagonal lines in the bottom right corner of the slide, slanting upwards from left to right.

DFA DEFINITION



- $Q = \{q_0, q_1, q_2, q_3\}$
- $\Sigma = \{a, b\}$
- $\delta =$
 $\{((q_0, a), q_1), ((q_0, b), q_0), ((q_1, a), q_2), ((q_1, b), q_1),$
 $((q_2, a), q_3), ((q_2, b), q_2), ((q_3, a), q_3), ((q_3, b), q_3)\}$
- $q = q_0$
- $F = \{q_2\}$

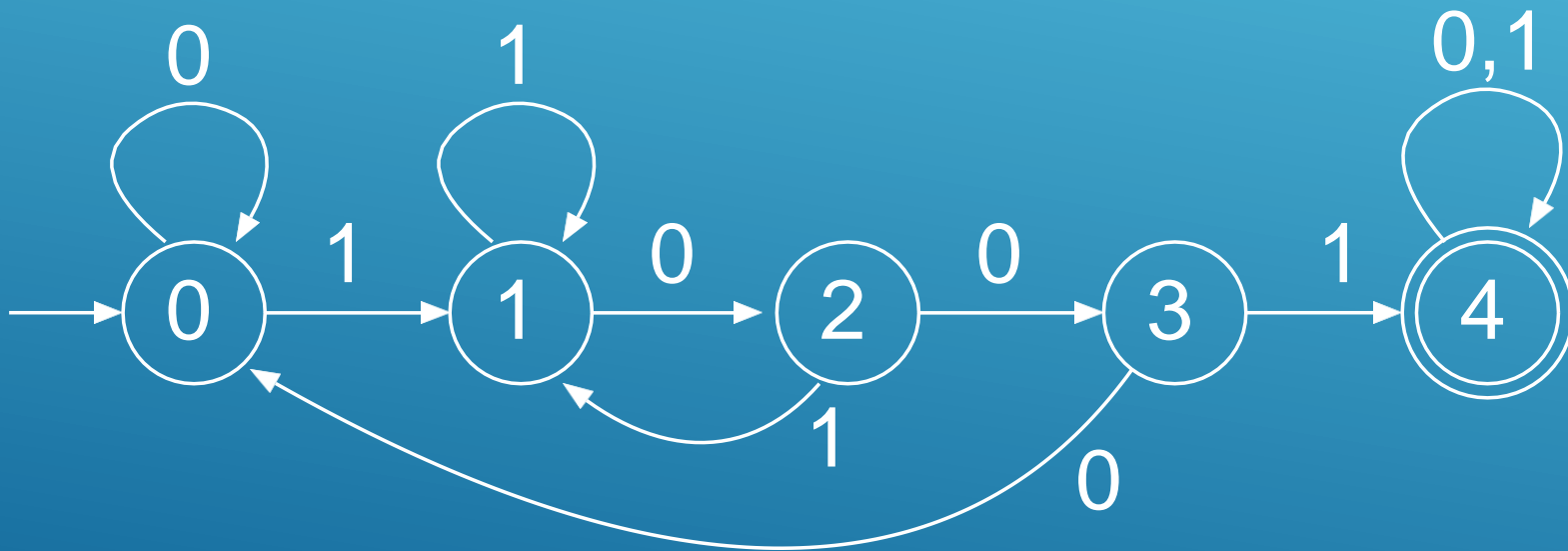
WHY DFA

- Why are these machines called “Deterministic Finite Automata”
 - Deterministic Each transition is completely determined by the current state and next input symbol. That is, for each state / symbol pair, there is exactly *one* state that is transitioned to
 - Every DFA has a finite number of states
- 
- A series of four parallel white diagonal lines in the bottom right corner of the slide, slanting upwards from left to right.

EXAMPLE. DFA

Create a DFA for:

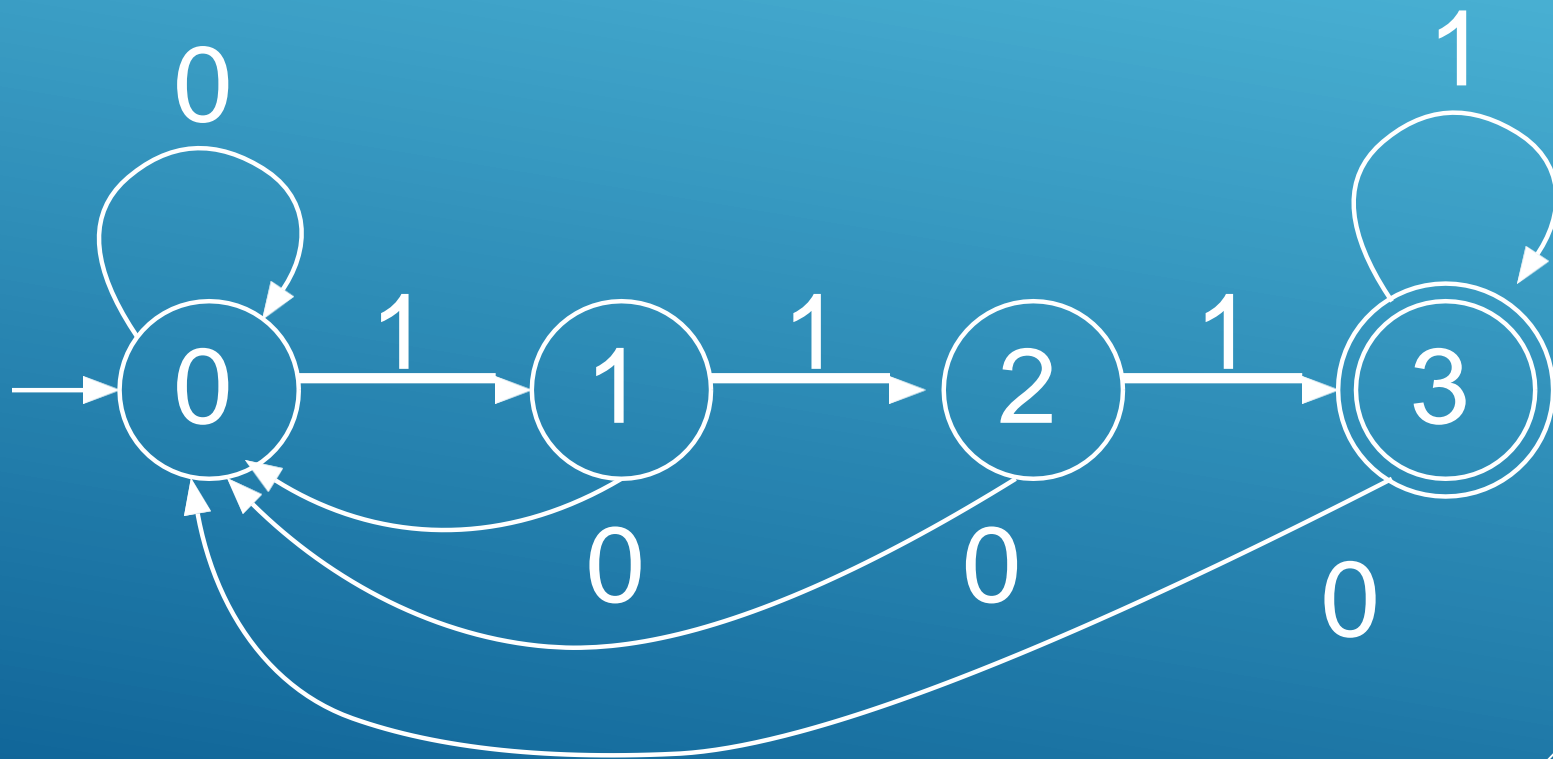
- All strings over $\{0, 1\}$ that contain the substring 1001



EXAMPLE. DFA

Create a DFA for:

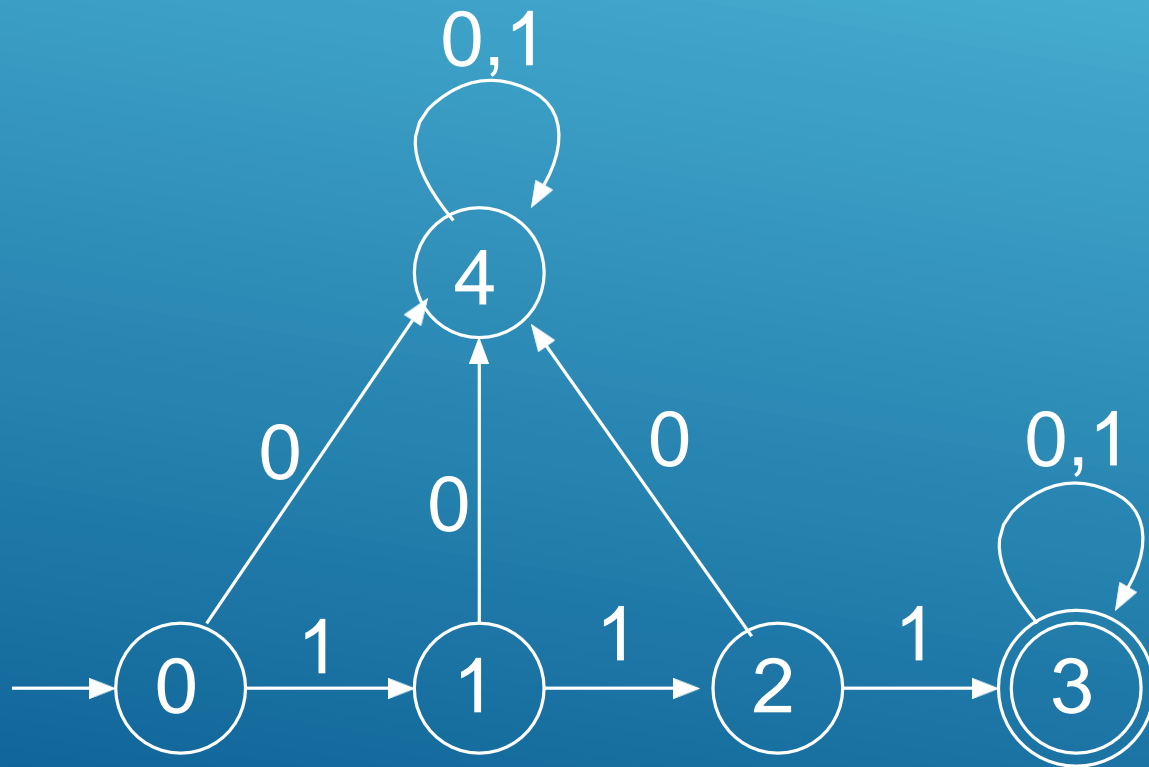
- All strings over $\{0, 1\}$ that end with 111



EXAMPLE. DFA

Create a DFA for:

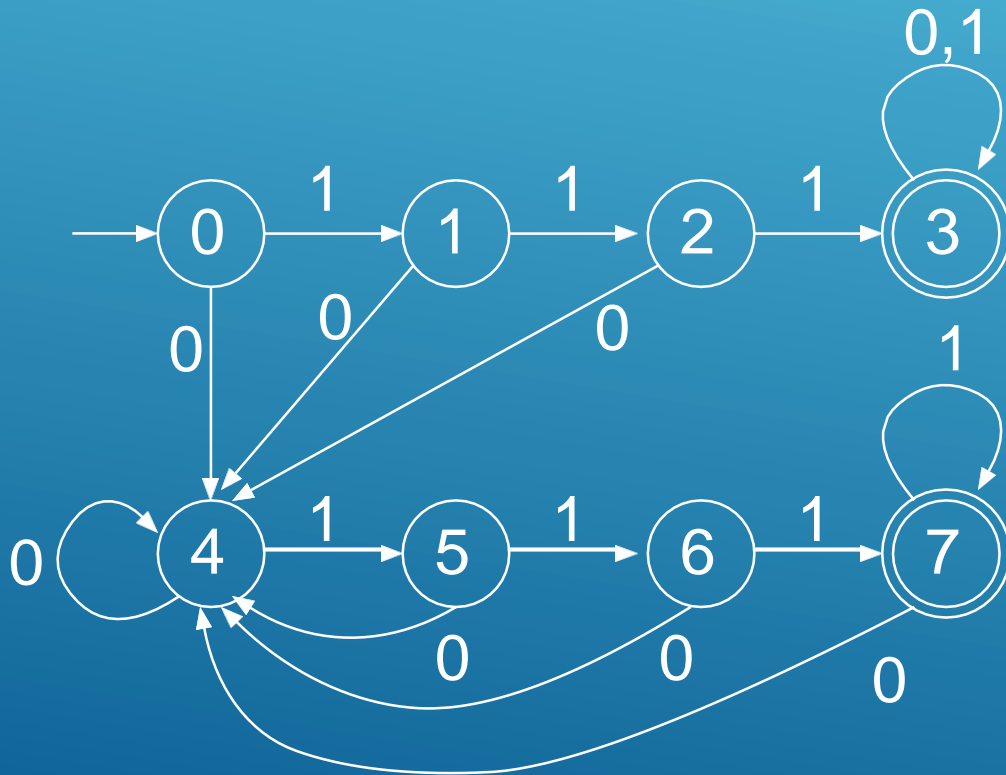
- All strings over $\{0, 1\}$ that begin with 111



EXAMPLE. DFA

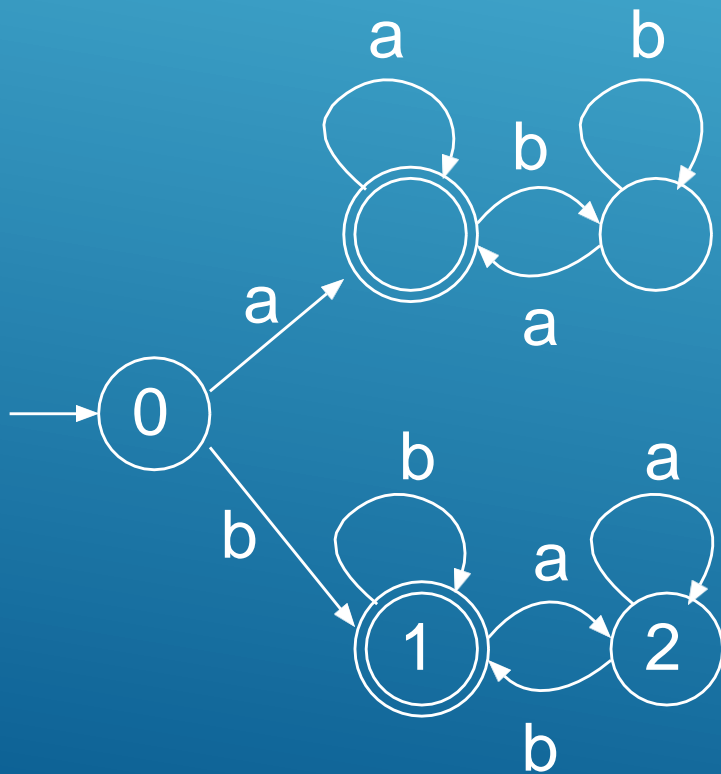
Create a DFA for:

- All strings over $\{0, 1\}$ that begin or end with 111



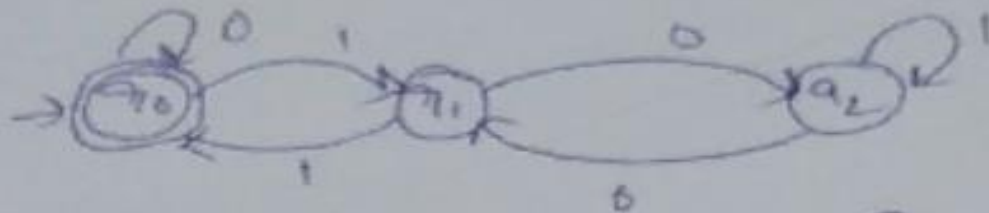
EXAMPLE. DFA

- ▶ Create a DFA for:
- ▶ All strings over $\{a, b\}$ that begin and end with the same letter



Construct a minima DFA which interpreted as binary number is divisible by '3' over $\Sigma \{0, 1\}$.

divisible by 3, therefore 3 states.

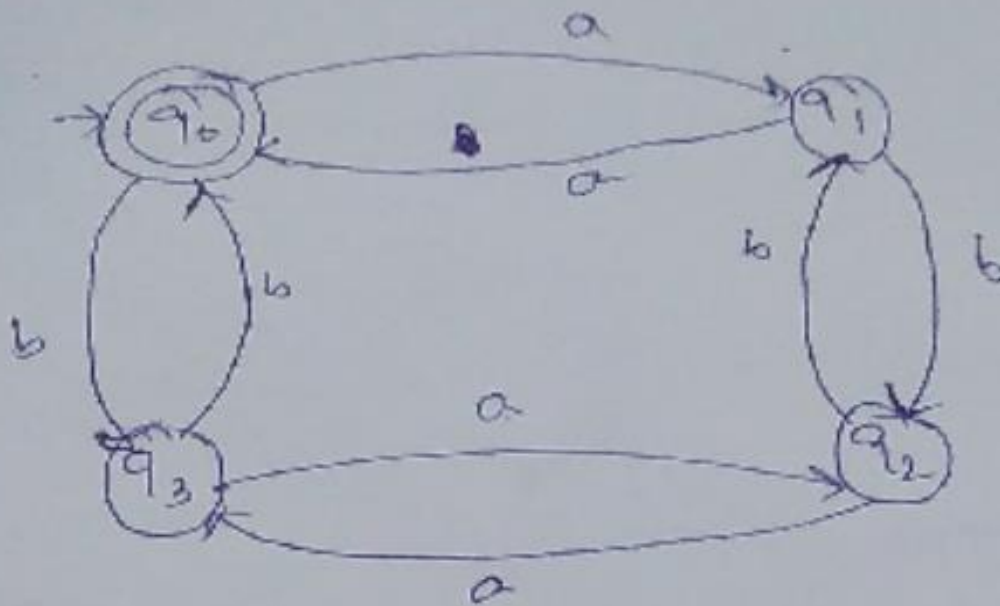


min of δ \Rightarrow

q_0	$\left\{ \begin{array}{l} 0 \times 2 + 0 = 0 \rightarrow q_0 \\ 0 \times 2 + 1 = 1 \rightarrow q_1 \end{array} \right.$
q_1	$\left\{ \begin{array}{l} 1 \times 2 + 0 = 2 \rightarrow q_2 \\ 1 \times 2 + 1 = 0 \rightarrow q_0 \end{array} \right.$
q_2	$\left\{ \begin{array}{l} 2 \times 2 + 0 = 1 \rightarrow q_1 \\ 2 \times 2 + 1 = 2 \rightarrow q_2 \end{array} \right.$

Construct a DFA which accepts strings of even number of a's and even number of b's over $\Sigma \{a, b\}$.

Construct DFA over $\Sigma \{a, b\}$ which DFA accepts strings of even number of a's and even number of b's



odd a's	odd b's	→ q2	}
even a's	odd b's	→ q3	
odd a's	even b's	→ q1	



NONDETERMINISM

- A *nondeterministic finite automaton* has the ability to be in several states at once.
- Transitions from a state on an input symbol can be to any set of states.
- ▶ Start in one start state.
- ▶ Accept if any sequence of choices leads to a final state.

NFA DEFINITION

NFA is a 5-tuple $M = (Q, \Sigma, \delta, q, F)$

Q Set of states

Σ Alphabet

$\delta : (Q \times \Sigma) \rightarrow P(Q)$ is a Transition function

$q \in Q$ Initial state

$F \subseteq Q$ Set of final states

EXAMPLE NFA

- Set of all strings with two consecutive a's or two consecutive b's:

