

CS & IT ENGINEERING

COMPUTER NETWORKS

Medium Access Control

Lecture No-02



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TOPICS TO BE COVERED

Multiple Access
Protocols Part-2



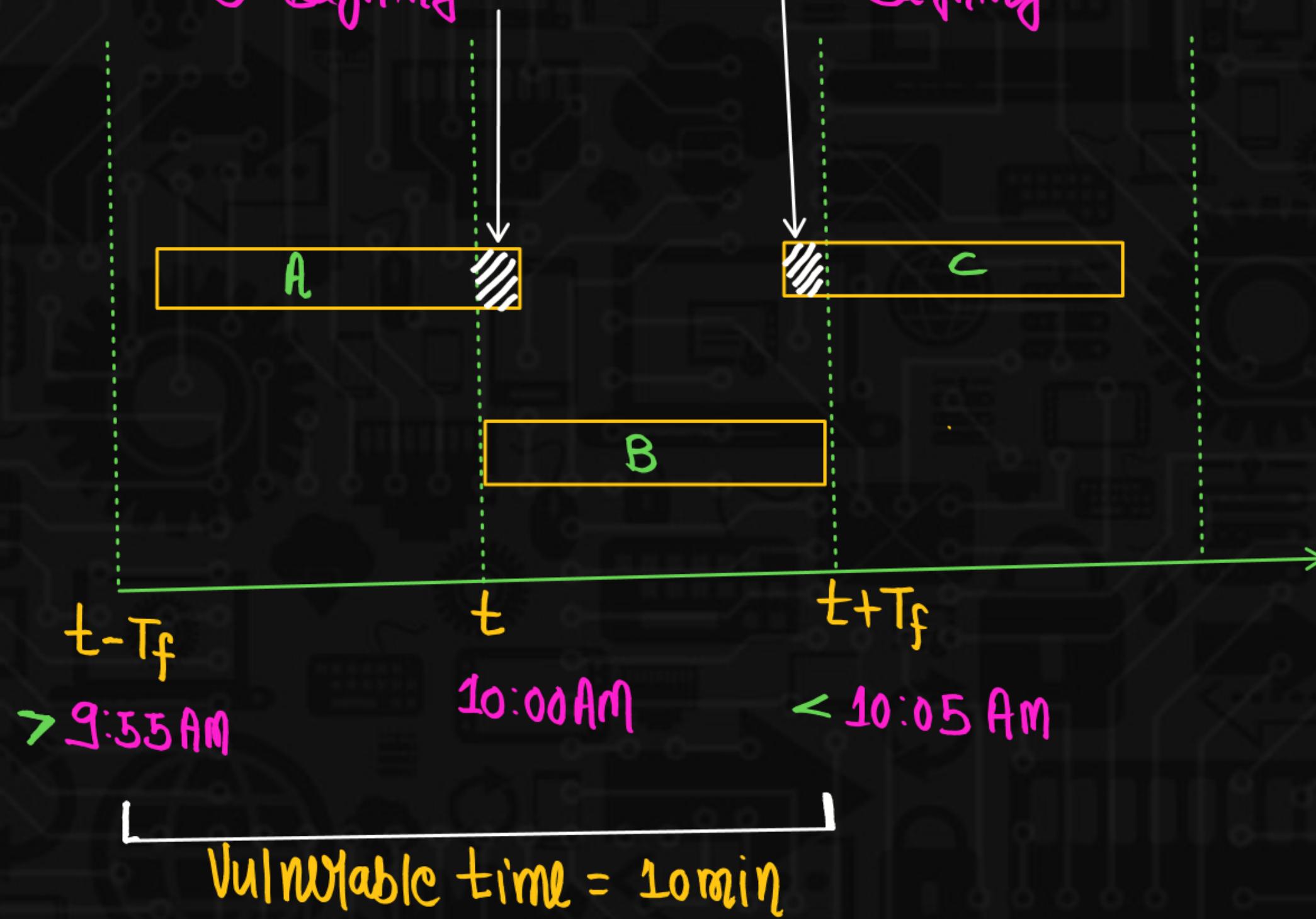
Vulnerable Time
For
Pure ALOHA

Vulnerable time for Pure ALOHA

Vulnerable time is the range of time where collision take place.

A end collide with
B Begning

B end collide with
C Begning



Vulnerable time For Pure Aloha = $2 \times T_f$



$$T_{\text{transmission time For single Frame}}(T_f) = \frac{\text{Frame size}}{\text{Bandwidth}} = \frac{L}{B}$$

assume transmission time of single Frame = 5min

2.



Assume $T_f = 2 \text{ sec}$

A starts at 2.1 sec

A → 2.1 sec to 4.1 sec

C starts at 5.9 sec



2 sec — 6 sec

Vulnerable time = 4 sec

Vulnerable time = $2 \times T_f$

Throughput Of Pure ALOHA

Throughput Of Pure ALOHA

$$C = Q \cdot T_{frame}$$

$$S = G \times e^{-2G}$$

G = Number of frames generated by n/w in one Frame transmission time

For Maximum Throughput

- We put $\frac{ds}{dG} = 0$
- Maximum value of s occurs at $G = \frac{1}{2}$
- Substituting $G = \frac{1}{2}$ in the above equation we get
- Maximum throughput of pure aloha

$$= \frac{1}{2} \times e^{-2 \times \frac{1}{2}} = \frac{1}{2e} = 0.184$$

So maximum throughput of pure aloha = 18.4%

Note

① If 1000 frames are generated by NW in one frame transmission time then maximum 184 frames will be delivered successfully.

② $S_{\max} \rightarrow G_l = \frac{1}{2}$

(i) one half frame should be generated in one frame transmission time to achieve maximum throughput
OR

(ii) one frame should be generated in two transmission time to achieve maximum throughput. Vulnerable time = $Q \times T_f$

Note

- If one frame is generated by the network in two frame transmission time then in this situation we will achieve maximum throughput
- Vulnerable time $2T_F$ is basically representing if one frame is generated by the network in two frame transmission time then there will be no collision. if there is no collision then we will achieve maximum throughput

Example:

A pure aloha network transmits 200 bit frames on a shared channel of 200 kbps. What is the throughput if the system(all station together) produces

Q.1) 1000 Frames generated by n/w in 1sec

Q.2) 500 Frames generated by n/w in 1sec

Q.3) 250 Frames generated by n/w in 1sec

NOTE: Throughput is defined as average number of frames successfully transmitted per second)

$$\text{Frame size} = 200 \text{ bits}$$

$$\text{Bandwidth} = 200 \text{ Kbps}$$

$$T_{\text{Transmission time}} = \frac{\text{Frame size}}{\text{Bandwidth}}$$

$$= \frac{200 \text{ bits}}{200 \times 10^3 \text{ bits/sec}} \\ = 10^{-3} \text{ sec} = 1 \text{ msec}$$

Soln Q.1

1sec → 1000 Frames

$1\text{msec} = 10^{-3}\text{sec}$ → 1000×10^{-3} Frame

$G=1$ [No. of Frames generated by the n/w in one Frame transmission time]

Throughput $S = G \times C^{-2G}$

$$= 1 \times C^{-2 \times 1}$$
$$= \frac{1}{C^2} = 0.135$$

Throughput = 13.5%

Avg Number of Frames
Successfully transmitted

$$\begin{aligned}\text{Per sec} &= 1000 \times 0.135 \\ &= 135\end{aligned}$$



Soln Q.2

$$1 \text{ sec} \longrightarrow 500 \text{ frames}$$

$$\begin{aligned}1 \text{ msec} = 10^{-3} \text{ sec} &\longrightarrow 500 * 10^3 \text{ frames} \\&= 5 * 10^1 \\&= \frac{5}{16} = \frac{1}{4} \text{ frame}\end{aligned}$$

$$G_1 = \frac{1}{4}$$

$$\text{throughput } S = G_1 * e^{-q_1 G_1}$$

$$= \frac{1}{4} * e^{-\frac{1}{4} * \frac{1}{4}}$$

$$= \frac{1}{4} = 0.184$$

$$\text{Throughput} = 18.4 \cdot 10$$

Avg No. of Frames successfully -

$$\begin{aligned}\text{transmitted per sec} &= 500 * 0.184 \\&= 92\end{aligned}$$



Sol Q.3

$$G_1 = \frac{1}{4}$$

$$\text{Throughput } S = 0.152$$

$$\text{Throughput} = 15.2 \cdot 1.$$

Avg. No. of Frames successfully

$$\text{Transmitted} = 250 * 0.152 = 38$$

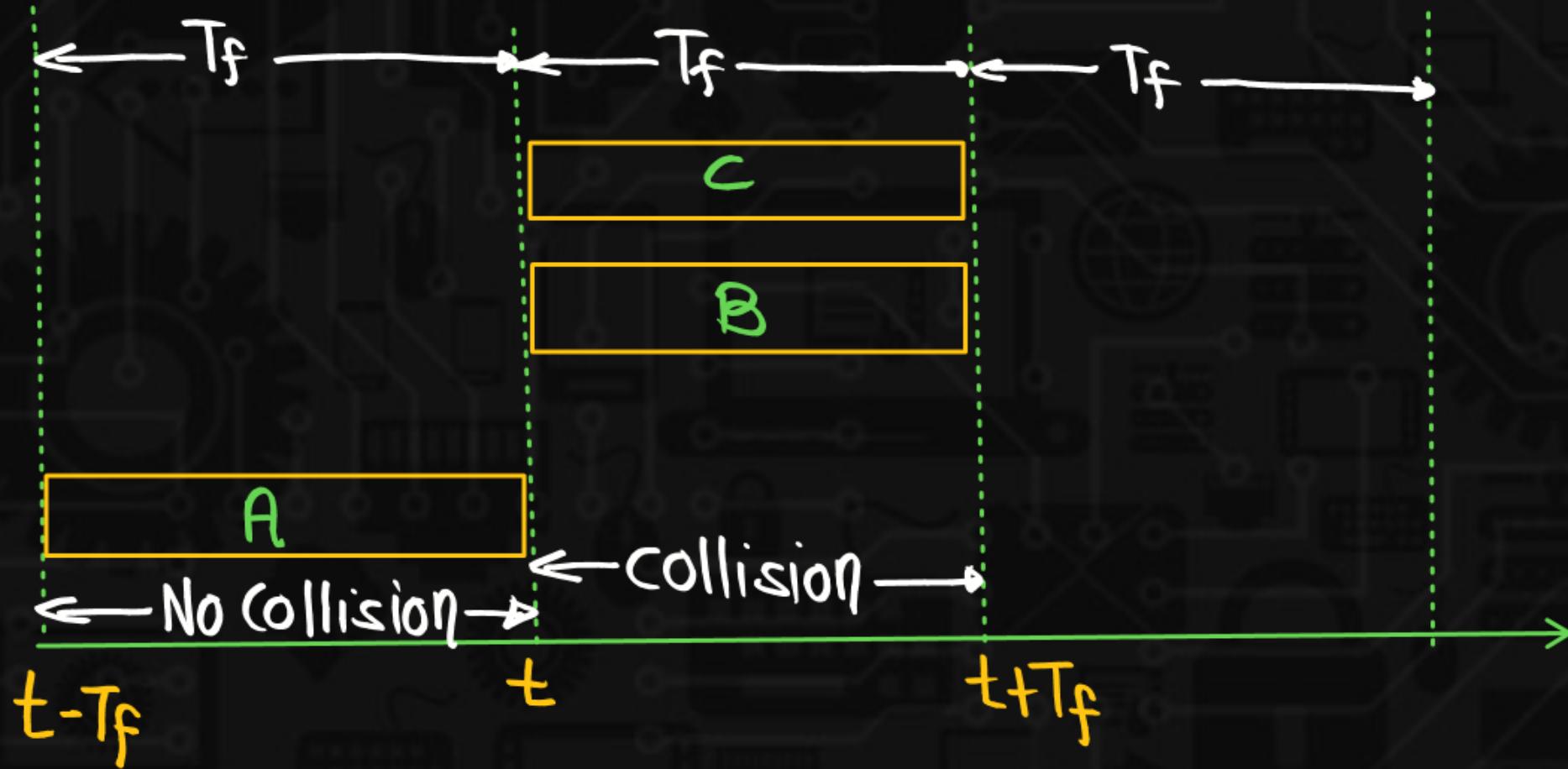
Slotted ALOHA

Slotted ALOHA

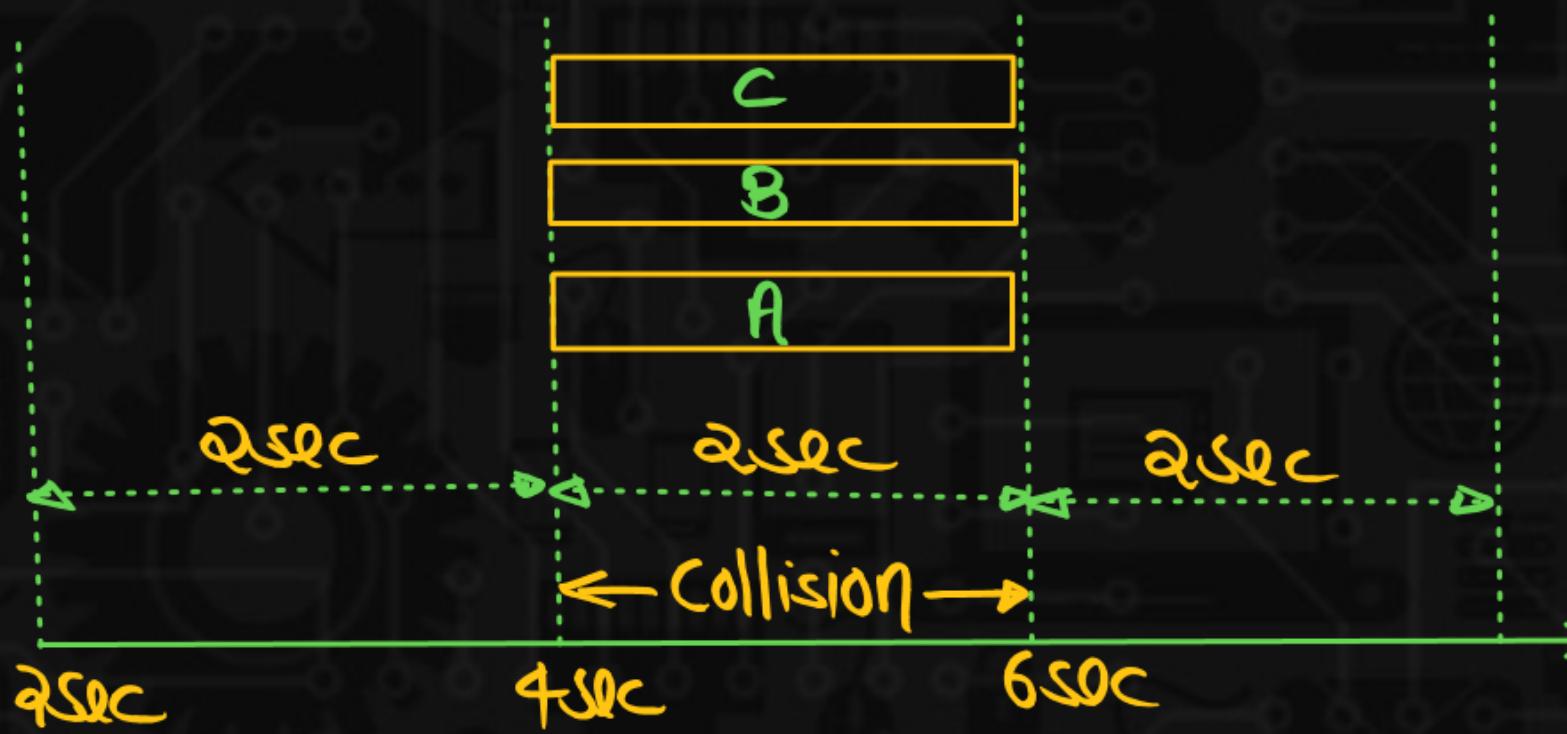
- Slotted aloha divides the time of shared channel into discrete intervals called as time slot (time slot = Transmission time For one frame).
- Any station can transmit its data in any time slot.
- The only condition is that station must start its transmission from the Beginning of time slot.
- If the Beginning of the slot is missed, then station has to wait until the beginning of next time slot.
- A collision may occur if two or more stations try to transmit data at the beginning of the same time slot.

Vulnerable time for Slotted ALOHA

Vulnerable time is the range of time where collision take place.



Vulnerable time for slotted Aloha = T_f

Note

$$A \rightarrow 2.3 \text{ sec}$$

$$B \rightarrow 3.1 \text{ sec}$$

$$C \rightarrow 3.3 \text{ sec}$$

4 sec

Collision Range 4sec → 6sec

Vulnerable time = 2sec

Vulnerable time = T_f Assume $T_f = 2 \text{ sec}$

station A → 2.3 sec → 4 sec

station B → 3.1 sec → 6 sec

Note

- ① For example $T_f = 2\text{ sec}$. and station Has data to send at 2.3 sec
then g_t will send at 4 sec . and station Has data to send at 5.2 sec
then g_t will send at 6 sec
- ② In slotted Aloha collision can place For example if there 3 station
Has data to send at $2.3, 3.1, 3.3\text{ sec}$ then all of them will send
at 4 sec .

Throughput Of Slotted ALOHA

$$S = G \times e^{-G}$$

G = Number of frames generated by n/w in one Frame transmission time

For Maximum Throughput

- We put $\frac{ds}{dG} = 0$
- Maximum value of s occurs at $G = 1$
- Substituting $G = 1$ in the above equation we get
- Maximum throughput of pure alone

$$= 1 \times e^{-1} = \frac{1}{e} = 0.368$$

So maximum throughput of slotted aloha = 36.8%

Note:

① If 1000 frames are generated by the NW in one frame transmission time then maximum 368 frames will be delivered successfully.

Note

$$S_{\max} \rightarrow G = 1$$

- one frame should be generated in one time slot (one frame transmission time) to achieve maximum throughput
- Vulnerable time T_F is basically representing if one frame is generated by the network in one slot time (one frame transmission time) then there will be no collision. if there is no collision then we will achieve maximum throughput

Example:

A slotted aloha network transmits 200 bit frames on a shared channel of 200 kbps. What is the throughput if the system(all station together) produces

Q.1) 1000 Frames generated by n/w in 1sec

Q.2) 500 Frames generated by n/w in 1sec

Ans: 151

Q.3) 250 Frames generated by n/w in 1sec

Ans: 49

NOTE: Throughput is defined as average number of frames successfully transmitted per second)

$$\text{Frame size} = 200 \text{ bits}$$

$$B = 200 \text{ Kbps}$$

$$\text{Transmission time} = \frac{200 \text{ bits}}{200 \times 10^3 \text{ bits/sec}}$$

$$= 10^{-3} \text{ sec} = 1 \text{ msec}$$

Soln: Q. 1

1sec \rightarrow 1000 Frames

$1\text{msec} = 10^{-3}\text{sec}$ \rightarrow ~~1000×10^{-3}~~ Frame
 \rightarrow 1 Frame

$$G_1 = 1$$

$$\begin{aligned}\text{Throughput } S &= G_1 * e^{-G} \\ &= 1 * e^{-1} \\ &= \frac{1}{e} = 0.368\end{aligned}$$

$$\text{Throughput} = 36.8\%$$

Avg. No. of Frames successfully transmitted per sec

$$\begin{aligned}&= 1000 * 0.368 \\ &= 368\end{aligned}$$

Pure Aloha	Slotted Aloha
(1) Any station transmit the data at any time	Any station can transmit the data at the beginning of any time slot
(2) Vulnerable time in which collision may occur = $2 \times T_f$	Vulnerable time in which collision may occur = T_f
(3) Throughput of pure aloha = $G \times e^{-2G}$	Throughput of slotted Aloha = $G \times e^{-G}$
(4) Maximum throughput $s_{max} = 18.4\%$ (When $G = 1/2$)	Maximum throughput $s_{max} = 36.8\%$ (When $G = 1$)
(5) The main Advantage of pure aloha is its simplicity in implementation	The main advantage of slotted aloha is that it reduces the number of collisions to half and double the throughput of pure aloha

