

# CS & IT ENGINEERING

Theory of Computation  
Finite Automata



Lecture No. 14



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## TOPICS TO BE COVERED

01 Assignment Questions

02 GATE Questions

03 Doubts

Q.1

Number of prefixes of "n" length string is \_\_\_\_\_  
(assume all the symbols in given string are different)

- A. n
- B. n+1**
- C. n+2
- D. n-1

$$w = \boxed{abc} = n = 3$$

Diagram illustrating the prefixes of the string  $w = abc$ :

- The string  $w$  is shown in a box.
- The empty prefix is labeled  $\epsilon$ .
- The first prefix is  $a$ .
- The second prefix is  $ab$ .
- The third prefix is  $abc$ .
- A curly brace groups all four prefixes ( $\epsilon, a, ab, abc$ ) and is labeled  $\gamma = n + 1$ .

Q.2

P  
W

$((ab)^* \setminus \emptyset)$  is equivalent to

- A.  $(ab)^*$
- ~~B.  $\emptyset$~~
- C.  $(a+b)^*$
- D. None

$$(ab)^* + \phi = (ab)^*$$

$$(ab)^* \cdot \phi = \phi$$

Q.3

$$(\emptyset^* \cup (\underline{b}b^*)) =$$

- A. Epsilon
- B.  $\emptyset$
- C.  $b^+$
- D.  $b^*$

$$\mathcal{E} + b^+ = b^*$$

P  
W

$$\phi^* = \epsilon$$

$$\phi^+ = \phi$$

$$\phi^0 = \epsilon$$

Q.4

P  
W

Which of the following is TRUE?

$a, aba, ababa, \dots$

$a, aba, ababa, \dots$

A.  $(ab)^*a = a(ba)^*$

B.  $(aa)^*b = a(ab)^*$

C.  $(ba)^*a = b(aa)^*$

D. All of the above

a

b

**Q.5**

OR operator in regular expression satisfies

- A. Associative
- B. Commutative
- C. both A and B
- D. Neither A nor B

**Q.6**

P  
W

Concatenation operator in regular expression satisfies

- A. Associative
- B. Commutative
- C. both A and B
- D. Neither A nor B

Q.7

Which of the following distribution is valid in regular expressions?

- A. OR over CONCATENATION
- B. CONCATENATION over OR
- C. Both A and B
- D. Neither A nor B

$$\begin{aligned}\cancel{a + (b+c)} &= \\ a \cdot (b+c) &= ab+ac\end{aligned}$$

**Q.8**

P  
W

Match the following groups over  $\Sigma = \{a, b\}$ .

**Group-1:**

1. OR identity  $\rightarrow b$
2. OR dominator  $\rightarrow c$
3. Concatenation identity  $\rightarrow a$
4. Concatenation dominator  $\rightarrow b$

**Group-2:**

- a. Epsilon
- b. Empty Expression
- c.  $(a+b)^*$
- d.  $(aa)^*$

- A. 1-a, 2-b, 3-c, 4-d
- ~~B. 1-b, 2-c, 3-a, 4-b~~
- C. 1-d, 2-a, 3-b, 4-c
- D. None

Q.9

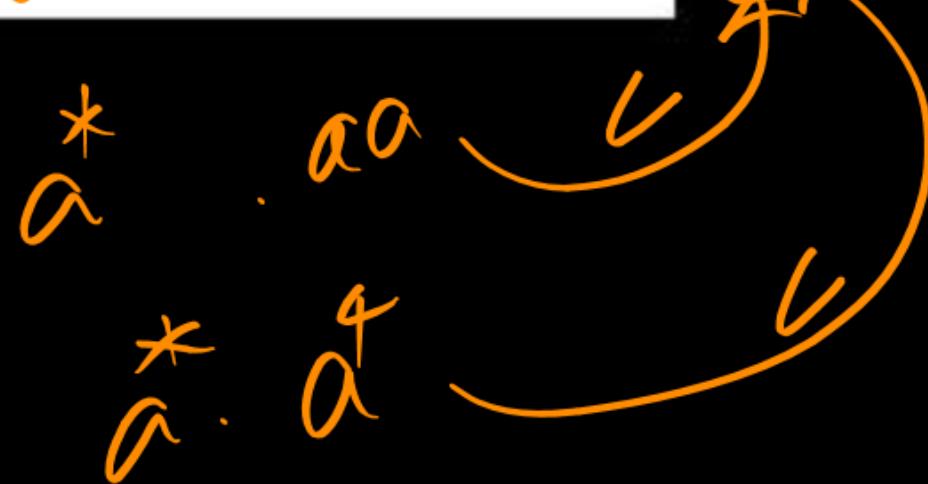
If  $R1 = a^*$ , and  $R2 = (aa)^*$  then  $R1.R2 = \underline{\quad}$

P  
W

- A.  $R1$
- B.  $R2.R1$
- C.  $R1+R2$
- D. All of these

$$a^* \cdot (aa)^* = a^*$$

Diagram illustrating the NFA for  $a^* \cdot (aa)^*$ . The start state is labeled  $\epsilon$ . There are two states: a black state and a white state. Transitions: from  $\epsilon$  to the black state labeled  $a^*$ ; from the black state to the white state labeled  $aa$ ; from the white state back to the black state labeled  $a^*$ .



Q.10

If  $R_1 = a(aa)^*$ , and  $R_2 = (aa)^*$  then  $R_1 + R_2 = \underline{\quad}$

- A.  $R_1$
- B.  $R_2 \cdot R_1$
- C.  $R_1^*$
- D. None

$$\begin{aligned}
 R_1 + R_2 &= a(aa)^* \cup (aa)^* \\
 &= \overset{*}{a}
 \end{aligned}$$

$$R_2 \cdot R_1 = (aa)^* \cdot a(aa)^* \rightarrow a, a^3, a^5, \dots$$

$$\begin{aligned}
 R_1^* &= [a (aa)^*]^* = \{\varepsilon, a, a^2, \dots\} \rightarrow R_1 \\
 &= \overset{*}{a}
 \end{aligned}$$

# Home Work:

shortest length string

Find minimum string in the following Expressions.

$$1) \underbrace{(ab)^+}_{ab} \underbrace{aaa}_{\overset{\circ}{aaa}} \underbrace{(b+ab)}_b \longrightarrow \text{min} = abaaab$$

$$2) \underbrace{(ab)^*}_{\epsilon} \underbrace{aaa}_{\overset{\circ}{aaa}} \underbrace{(b+\epsilon)}_{\epsilon} \longrightarrow aac$$

$$3) \underbrace{(a+ab)^+}_a \underbrace{(bb+aaa)^+}_b \underbrace{(ab)^*}_{\epsilon} \longrightarrow abb$$

$$4) \underbrace{(ab)^+ + \epsilon}_{\epsilon} \longrightarrow \epsilon$$

5)  $\underbrace{(ab)(a+b)^*aaa}_{ba} + \underbrace{ba}_{ba} \rightarrow ba$

6)  $\left[ \underbrace{(ab^* + a^+ + ba)ab}_{aab} \right]^+ \cdot \underbrace{(a+b)^*}_{\epsilon} \rightarrow aab$

7)  $\left[ \underbrace{((ab)^+ a)^* aaa}_{\epsilon} \right]^+ \rightarrow \alpha \alpha \alpha$

8)  $\left[ \underbrace{((aba)^+ aba)^+}_{\epsilon} ab \right]^* \rightarrow \epsilon$

$$9) (\dot{a})^+ \xrightarrow{\epsilon} \epsilon$$

$$10) \frac{((ab)^* \dot{a}^+)^+}{\epsilon \quad a} \xrightarrow{bb^+ aa^*} abba \quad \left. \begin{array}{l} \\ \\ \end{array} \right\} 14)$$

$$11) \left[ \frac{(ab)^+}{ab} \cdot \frac{aba}{aba} \right]^+ \cdot \frac{aaa}{aaa}$$

$$12) \left( \frac{aaa^* aba^*}{\epsilon \quad \epsilon} \right)^+ \xrightarrow{\quad} aaaab$$

$$13) \left( \underline{a} + \underline{aa} + \underline{aaa} \right)^+ \xrightarrow{a} \left. \begin{array}{l} \\ \\ \end{array} \right\} 15)$$

14)

$$\frac{\dot{a}^+ \dot{b}^+}{a \quad b} \longrightarrow ab$$

$$\frac{\dot{a}^* \dot{b}^*}{\epsilon \quad \epsilon} \longrightarrow \epsilon$$

16)

$$\frac{\dot{a}^* \dot{b}^* c^+ d^*}{\epsilon \quad \epsilon \quad \epsilon \quad \epsilon} \longrightarrow c$$

17)

$$\left( a + \frac{b}{\epsilon} + c^+ \right)^+ \longrightarrow \epsilon$$

18)

$$\left( a + b + \epsilon^- \right)^+ \longrightarrow \epsilon$$

$$(a+\varepsilon)^* \quad \Sigma = \{a\}$$

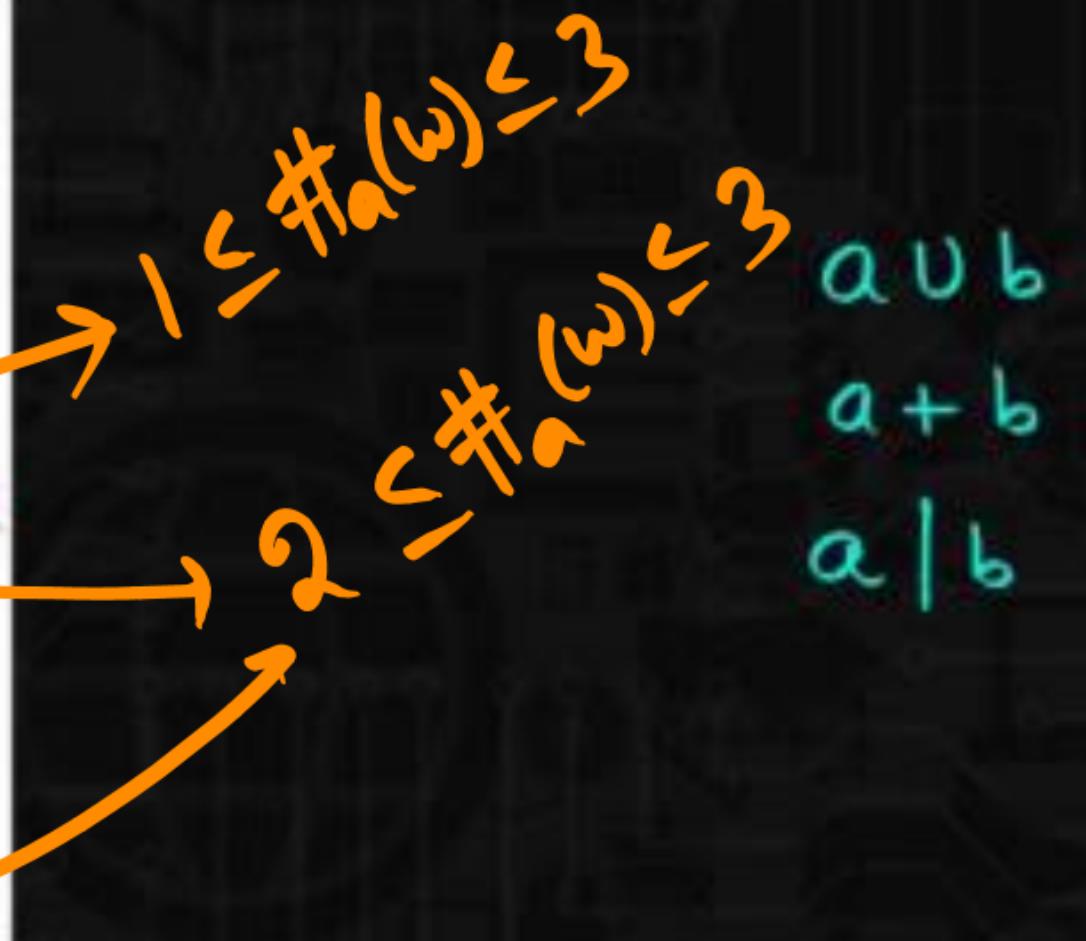
P  
W

1)

$$L = \{w \in \{a, b\}^*: \#_a(w) \leq 3\}.$$

*at most 3*

- A.  $b^* (a \cup \varepsilon) b^* (a \cup \varepsilon) b^* (a \cup \varepsilon) b^*$
- B.  $b^* \underline{(a)} b^* \underline{(a \cup \varepsilon)} b^* \underline{(a \cup \varepsilon)} b^*$
- C.  $b^* \underline{(a)} b^* \underline{(a)} b^* \underline{(a \cup \varepsilon)} b^*$
- D.  $b^* (a \cup \varepsilon) b^* \underline{(a)} b^* \underline{(a)} b^*$



$$\overset{*}{\epsilon} (a+\varepsilon) \overset{*}{b} \quad (a+\varepsilon) \overset{*}{\epsilon} \quad (a+\varepsilon) \overset{*}{b}$$

$$\Sigma = \{a, b\}$$

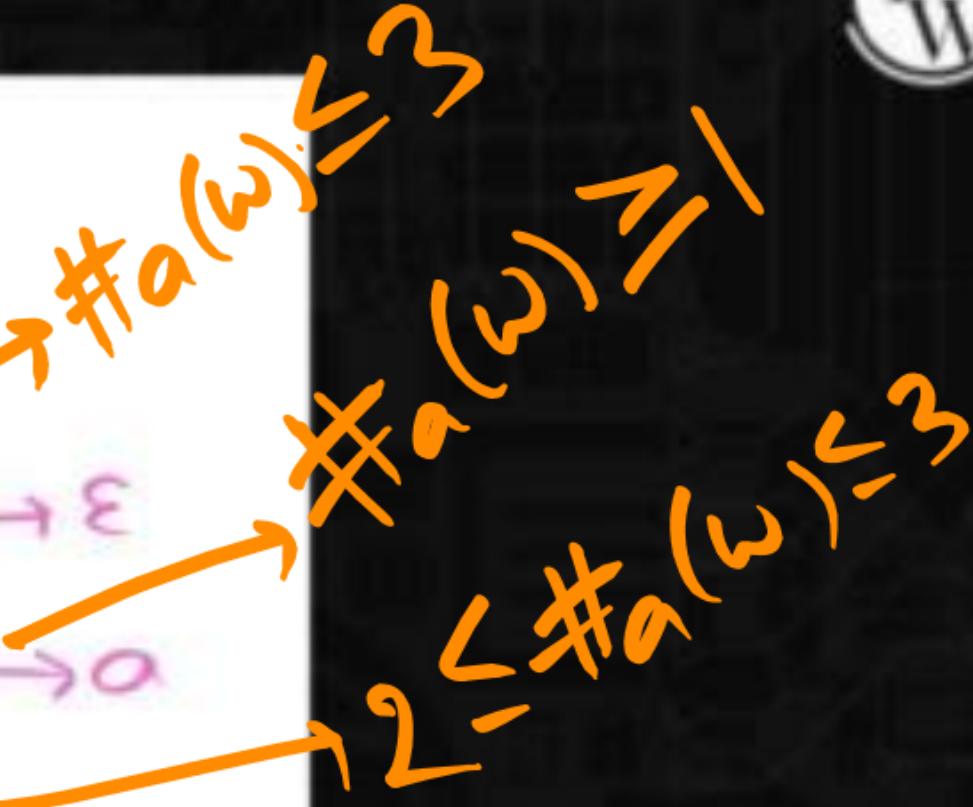
2)

$$L = \{w \in \{a, b\}^*: \#_a(w) \geq 3\}.$$

At least 3

- A.  $b^* (a \cup \epsilon) b^* (a \cup \epsilon) b^* (a \cup \epsilon) b^* \rightarrow \epsilon$
- B.  $\underline{(a \cup b)}^* \underline{(a)} b^* \underline{(a \cup \epsilon)} b^* \underline{(a \cup \epsilon)} b^* \rightarrow a$
- C.  $b^* \underline{(a)} b^* \underline{(a)} b^* \underline{(a \cup \epsilon)} b^* \rightarrow aa$
- D.  $(a \cup b)^* \underline{a} (a \cup b)^* \underline{a} (a \cup b)^* \underline{a} (a \cup b)^*$

X      a      X      a      X      a      X



3)

$$L_1 = a^* b^*$$

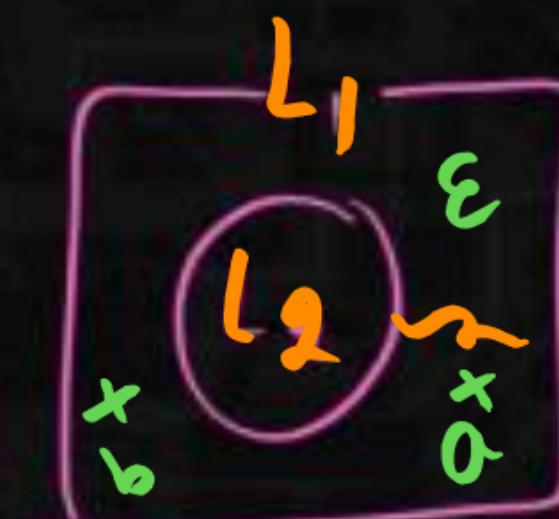
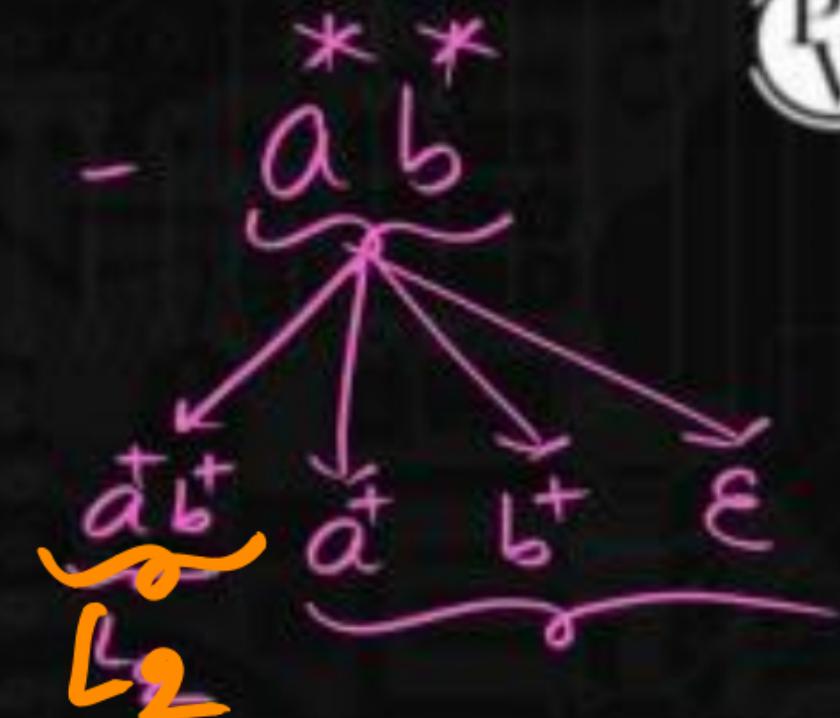
$$L_2 = \underbrace{a^+ b^+}_{\text{in}}$$

Find  $L_2 - L_1$ .

- A.  $a^*$   
 B.  $b^*$   
 C.  $a^* + b^*$   
 D. None

$$L_2 - L_1 = \emptyset$$

$$L_2 - L_1 = a^+ b^+ - = \emptyset$$

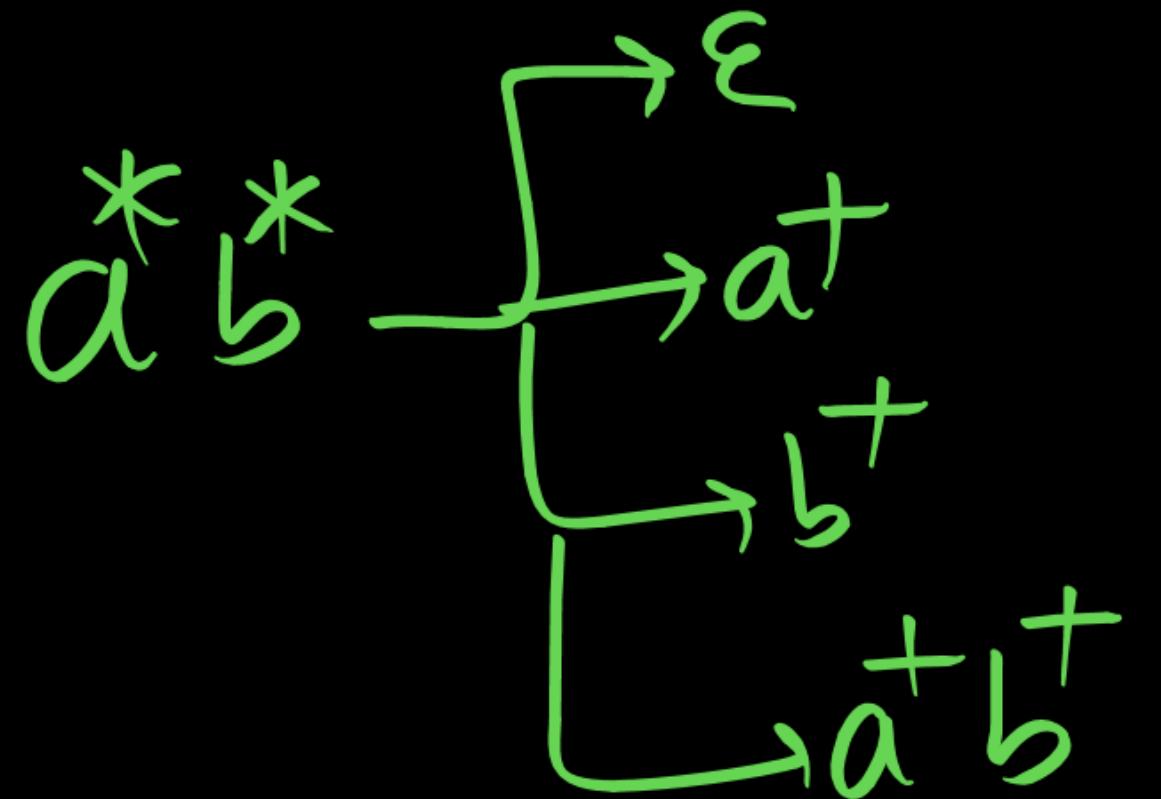


$$L_2 \subset L_1$$

Note:  $L_1 - L_2 = a^* b^* - a^+ b^+ = a^+ + b^+ + \epsilon$

```

graph LR
    l1((l1)) -- "l2" --> l2((l2))
    l1 -- "*" --> e1(( ))
    l2 -- "a" --> a1((a))
    l2 -- "b" --> b1((b))
    a1 -- "a^+" --> a2(( ))
    b1 -- "b^+" --> b2(( ))
    a2 -- "*" --> e2(( ))
  
```



$\overset{+}{a} \overset{+}{b}$

If  $X \subset Y$  then

$$\text{i)} X \cup Y = Y$$

$$\text{ii)} X \cap Y = X$$

4)

$L_1 = a^* + b^*$  and  $L_2 = \overbrace{a^* b^*}$ .

Which of the following is TRUE?

~~A.~~  $L_1 = L_2$

$$L_1 < L_2$$

~~B.~~  $L_1 \cup L_2 = (a+b)^*$

$$L_1 \cup L_2 = L_2$$

~~C.~~  $L_1^* = L_2^* = (a+b)^*$

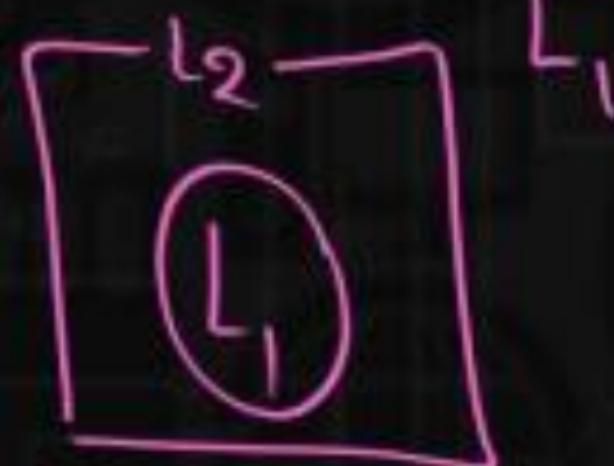
D. None

$$L_1^* = (a^* + b^*)^* = (a+b)^*$$

$$L_2^* = (a^* b^*)^* = (a+b)^*$$

$$L_1 = a^* + b^*$$

$$L_2 = a^* b^*$$



5)

$$L_1 = a^* + b^* \text{ and } L_2 = a^*b^*.$$

Which of the following is TRUE?

- A.  $L_1$  is subset of  $L_2$
- B.  $L_2$  is subset of  $L_1$
- C.  $L_1 \cup L_2 = L_1$
- D. None

$$\begin{matrix} L_1 & \subset & L_2 \\ L_2 & \supset & L_1 \end{matrix}$$

$$L_1 \subset L_2$$

$$L_1 \cup L_2 = L_2$$

$$L_1 \cap L_2 = L_1$$

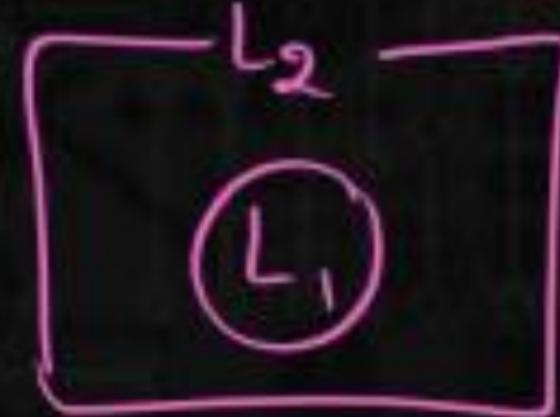
$$L_1 \cup L_2 = L_2$$

$$(a+b)^* = (\overset{+}{a} \overset{+}{b})^*$$

6)

$L_1 = a^+b^+$  and  $L_2 = a^*b^*$ .

Which of the following is FALSE?



- A.  $L_1$  is subset of  $L_2$   $\top$
- B.  $L_1^* = L_2^*$   $\text{False}$
- C.  $L_1 \cup L_2 = L_2$   $\top$
- D. None

$$\begin{array}{c} (\overset{+}{a} \overset{+}{b})^* \neq (\overset{*}{a} \overset{*}{b})^* \\ = (a+b)^* \end{array}$$

$$\begin{array}{c} L_1 \subset L_2 \\ a^+b^+ \in L_1 \\ a^*b^+ \in L_2 \end{array}$$

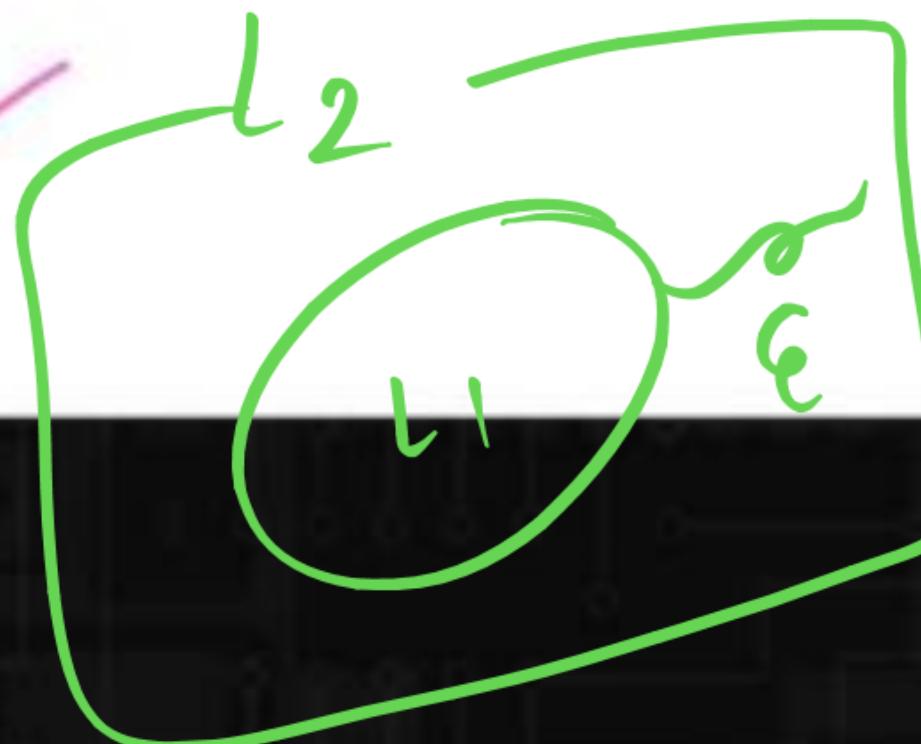
$$\begin{array}{l} L_1 \subset L_2 \\ x a \leftarrow L_1^* = (\overset{+}{a} \overset{+}{b})^* \\ \checkmark a \leftarrow L_2^* = (\overset{*}{a} \overset{*}{b})^* \end{array}$$

$$(a+b)^+ = (a+b)^*$$

?)  $L_1 = a^+ + b^+$  and  $L_2 = a^* + b^*$ .

Which of the following is TRUE?

- A.  $L_1 = L_2$   $\times$   $(a+b)^*$
- B.  $L_1^+ = L_2^+$   $\times$
- C.  $L_1 \cup L_2 = L_2$  ✓
- D. None



$$\begin{aligned}L_1^+ &= (a+b)^+ \\L_2^+ &= (a+b)^*\end{aligned}$$

$$L_1 + \epsilon = L_2$$

$$L_1 \subset L_2$$

$$L_2 = L_1 \cup \{\epsilon\}$$

8)

$$L1 = a^+ \text{ and } L2 = a^*$$

Which of the following is TRUE?

- A.  $L1^+ = L2^*$
- B.  $L1^+ = L2^+$
- C.  $L1^* = L2^+$
- D. None

Handwritten derivation:

$$(a^+)^+ = (a^*)^*$$

$$(a^*)^+ \neq (a^*)^+$$

$$(a^*)^+ = (a^*)^+$$

Annotations:

- $(a^+)^+$  is circled in green.
- $(a^*)^+$  is circled in green.
- $(a^*)^+$  is circled in green.
- $a^*$  is circled in green.
- $a^*$  is circled in green.
- $a^*$  is circled in green.

- i)  $a^+ a^+ = aa^+ = aaaa^*$
- ii)  $a^* a^* = a^*$
- iii)  $a^+ a^* = a^*$
- iv)  $(a^+)^+ = a^+$
- v)  $(a^*)^+ = a^*$
- vi)  $(a^*)^+ = a^*$
- vii)  $(a^*)^+ = a^*$

$$\begin{aligned}
 (\overset{+}{a})^+ &= (\overset{+}{a})^1 \cup (\overset{+}{a})^2 \cup (\overset{+}{a})^3 \cup \dots \\
 &= \overset{+}{a} \cup \overset{+}{a}\overset{+}{a} \cup \overset{+}{a}\overset{+}{a}\overset{+}{a} \cup \dots \\
 &= \overset{+}{a}
 \end{aligned}$$

9)

$$A = a^* \text{ and } B = b^*$$

$$AB = ?$$

- A.  $\{ a^n b^n \mid n \geq 0 \}$
- B.  $\{ a^m b^n \mid m, n \geq 0 \}$
- C.  $(a+b)^*$
- D. None

$$A = a^* = \{ a^n \mid n \geq 0 \}$$

$$B = b^* = \{ b^n \mid n \geq 0 \}$$

$$A \cdot B = a^* b^*$$

$$= \{ a^{n_1} b^{n_2} \mid n_1, n_2 \geq 0 \}$$

$$\neq \{ a^b \mid n \geq 0 \}$$

$$a^* = \{ a^n \mid n \geq 0 \}$$

$$b^* = \{ b^n \mid n \geq 0 \}$$

$$a^* \xrightarrow{*} \varepsilon \\ a \cdot b \xrightarrow{*} a$$

$$\{a^n\} \quad \{b^n\}$$

$$a^n b^n \neq a^* b^*$$

$$\{a^{n_1} b^{n_2}\}$$

$$\{a^n\} = \{a^n \mid n > 0\}$$

$$a^n b^n \Rightarrow \{ \epsilon, ab, a^2 b^2, a^3 b^3, \dots \}$$

$$a^* b^* \Rightarrow \{ \epsilon, a, aa, aaa, \dots, b, bb, bbb, \dots \\ aab, abl, \dots \}$$

$$a^* b^* = a^{n_1} b^{n_2} = a^m b^n$$

$$(ab)^* \neq \underbrace{ab}_\Sigma^n$$

$\downarrow$

$\downarrow$

$a b$

$a^i b^j = aabb$

10)

$$A = aa^* \text{ and } B = bb^*$$

$$(A \cup B)^* = ?$$

- A.  $\{ a^n b^n \mid n \geq 0 \}$
- B.  $\{ a^m b^n \mid m, n \geq 0 \}$
- C.  $(a+b)^*$
- D. None

$$\begin{aligned} A &= a^+ \\ B &= b^+ \\ (A+B)^* & \end{aligned}$$

$$\begin{aligned} &= (a^+ + b^+)^* \\ &= (a+b)^* \end{aligned}$$

$$\begin{aligned} a^* a &= a a^* = a^+ \\ b b^* &= b^* b = b^+ \end{aligned}$$

$$\begin{aligned} (A \cup B)^* &= (a^+ + b^+)^* \\ &= (a+b)^* \end{aligned}$$

10

Given the language  $L = \{ab, aa, baa\}$ , which of the following strings are not in  $L^*$ ?

- 1) abaabaaabaa  $\in L^5$   
 2) aaaabaaaaa  $\in L^4$   
 3) baaaaabaaaaab  $\notin L^*$

- A. 1 only
  - B. 2 only
  - C. 3 only
  - D. None

$$L = ab + aa + baa$$

$$L^* = (ab + aa + baa)^*$$

$$L = \{ab, aa, baa\}$$

$$L^* = (ab + aa + baa)^*$$

aaaab   
ab 

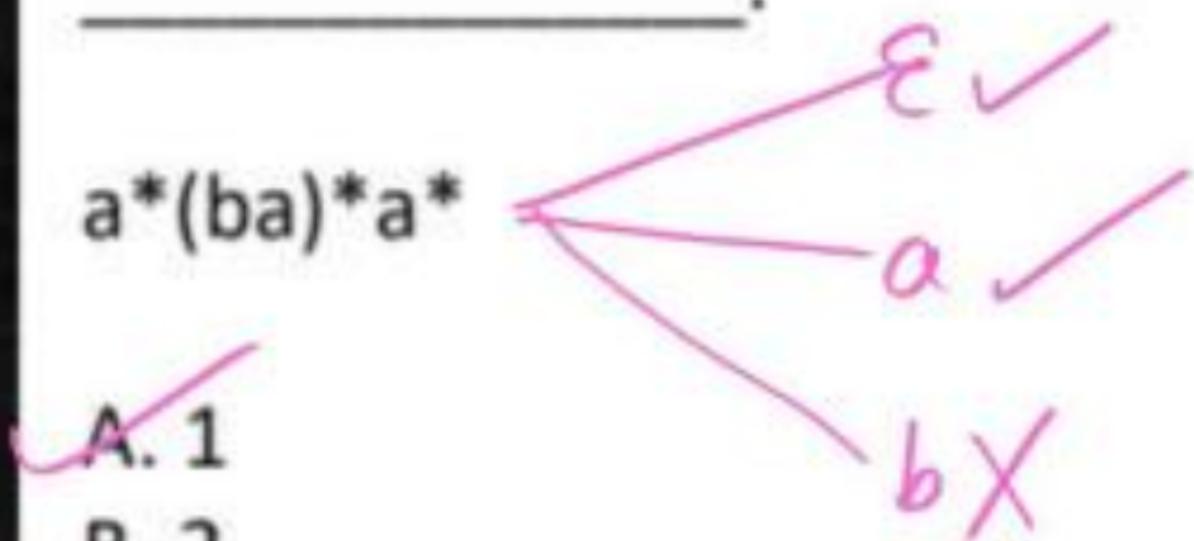
2) aaaabaaa

3) baaaaaabaaaad)<sup>E</sup><sub>L\*</sub>

12)

The length of the shortest string NOT in the language (over  $\Sigma = \{a, b\}$ ) of the following regular expression is \_\_\_\_\_.

$a^*(ba)^*a^*$



- A. 1
- B. 2
- C. 3
- D. 4

$\cancel{a^*} (ba)^* \cancel{a^*}$

$\epsilon$

$a$

b

aa

ab

ba

bb

$a^*(ba)^*a^*$

1 length

Starts with a

(13)

$L = a(a+b)^*$  is equivalent to \_\_\_\_\_

- A.  $(ab^*)^+$
- B.  $(a^+b^*)^+$
- C. a $^*(ab^*)^+$
- D. All of the above

$$\left. \begin{array}{l} A. (ab^*)^+ \\ B. (a^+b^*)^+ \\ C. \underline{a}^* (ab^*)^+ \end{array} \right\} = a(a+b)^*$$

$$a(a+b)^*$$

$$(ab^*)^+$$

$$(a^+b^*)^+$$

14)

$L = (a+b)^*b$  is equivalent to \_\_\_\_\_

- A.  $(ab^*)^+$
- B.  $(a^+b^*)^+$
- C.  $b^*(ab^*)^*$
- D. None

Note:

$$\left. \begin{array}{l}
 \text{i)} b^* (a b^*)^* \\
 \text{ii)} a^* (b a^*)^* \\
 \text{iii)} (b^* a)^* b^* \\
 \text{iv)} (a^* b)^* a^*
 \end{array} \right\} = (a+b)^*$$

$$(a+b)^*b = (\cancel{a}b)^+$$

$$= (\cancel{a}b^t)^+ = (\cancel{a}b)^+ b^*$$

$$= (a+b)^*$$

$$= (b^*a)^* b^*$$

$$= (a^*b)^* a^*$$

$$= b^* (a^*b)^*$$

$$= a^* (b^*a)^*$$

Imp forms



15)

$$(b + ba)(b + a)^*(ab + b) = b(\epsilon + a) \times (a + \epsilon)b$$

- A.  $(a+b)^*$
- B.  $a(a+b)^*a$
- C.  $b(a+b)^*b$
- D. None

$$\boxed{b \times b}$$

$$(a+b)^* (\epsilon + \text{Any}) = (a+b)^*$$

- I)  $a(a+b)^*a \rightarrow$  starts and ends with 'a'  
min 2 length
- II)  $b(a+b)^*b$
- III)  $a + a(a+b)^*a \Rightarrow$  starts ends with 'a'
- IV)  $b + b(a+b)^*b$

$$(a+b)^* \cdot (\epsilon + \underline{\text{Any}}) = (a+b)^*$$

$$(a+b)^* \cdot (ab)^* = (a+b)^*$$

$\circlearrowleft$

$\epsilon + ab + abab$

16)

$$\{w \in \{a, b\}^*: \#_a(w) \equiv_3 0\}.$$

A.  $(b^*ab^*ab^*a)^*b^*$

B.  $(b^*ab^*ab^*a)^*$   $\xrightarrow{by}$

C.  $(ab^*ab^*a)^*$   $\xrightarrow{bX}$   
 $\xrightarrow{baab+}$

D.  $(ab^*ab^*a)^*b^*$   $\xrightarrow{baab+}$



(17)

$$(a \cup b)^* \underbrace{(a \cup \epsilon)}_{\epsilon} \overline{\epsilon} b^* =$$

A.  $(a+b)(a+b)^* = (a+b)^+$

B.  $(a+b)^*$

C.  $(aa+b)^*$

D. None

$$(a+b)^* \underbrace{(a+\epsilon)b^*}_{\epsilon} = (a+b)^*$$

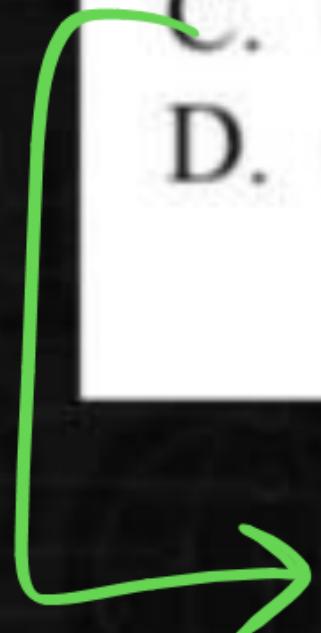
$$(a+b)^* \cdot \epsilon \cdot \epsilon$$

18)

 $L = \{w \in \{a, b\}^* \mid w \text{ has } bba \text{ as a substring}\}$ 

Which of the following describes L ?

- A.  $(a \cup b)^* bba (a \cup b)^*$
- B.  $(a \cup b)^+ bba (a \cup b)^*$
- C.  $(a \cup b)^+ bba (a \cup b)^+$
- D.  $(a \cup b)^* bba (a \cup b)^+$



bba  
valid  
but

A = ~~X bba X~~

B C A  
C C A  
D C A

19)

$$L = \{w \in \{a, b\}^*\} = (a+b)^*$$

- Same*
- a* ←
1.  $(a + b)^*$
  2.  $(a + b + \text{epsilon})^+$
  3. Epsilon +  $(a + b)^+$
  4.  $(a^*b^*)^*$
  5.  $(b^*a^*)^*$
  6.  $(\underline{a}^+\underline{b}^+)^* \neq (a+b)^*$

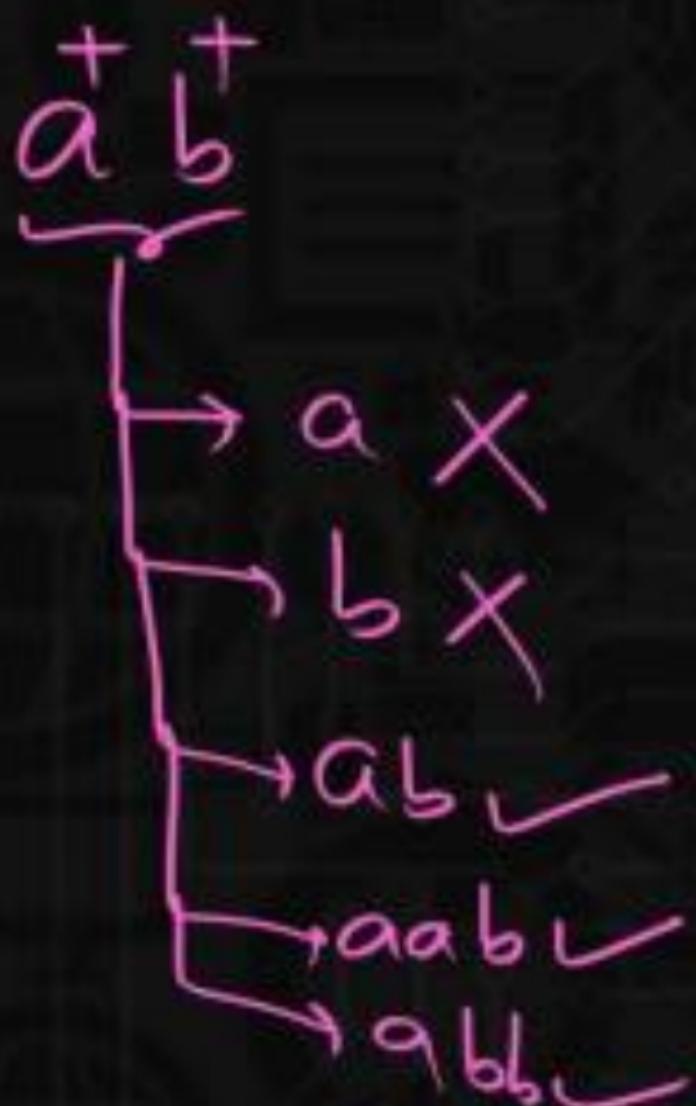
How many of above are equivalent to given L ?

A. 4

B. 5

C. 6

D. 3



0

Which Two of the following four regular expressions are equivalent?

$$(i) (00)^*(\epsilon + 0) = \overset{*}{0}$$

$\epsilon$  

$$(ii) (00)^*$$



$$(iii) \overset{*}{0}$$

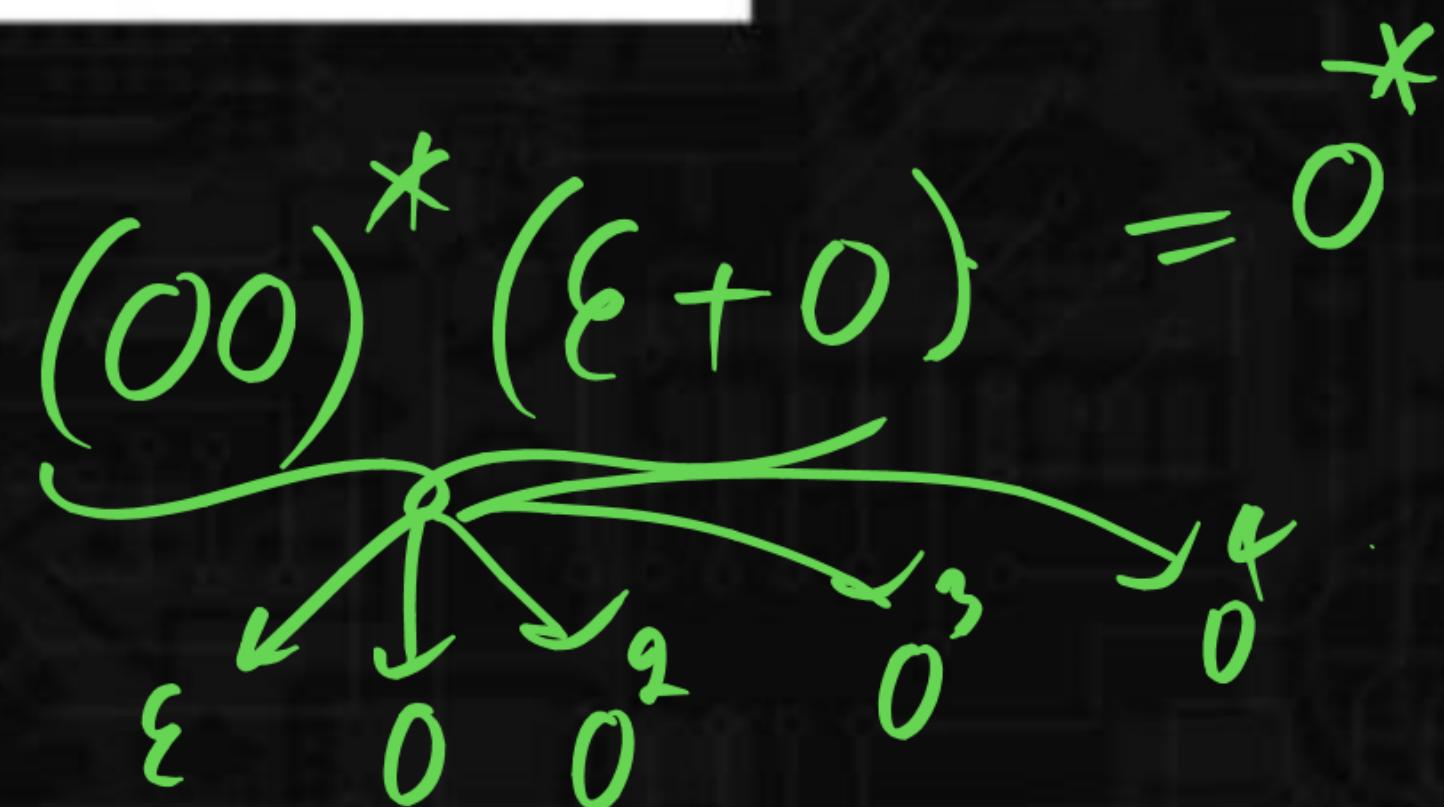


$$(iv) 0 (00)^*$$



- (a) (i) and (ii)
- (b) (ii) and (iii)
- (c) (i) and (iii)
- (d) (iii) and (iv)

**(GATE - 96)**

$$(00)^* (\epsilon + 0) = \overset{*}{0}$$


$$\text{I) } \left( (00)^* \right)^* = (00)^*$$

$$\text{II) } \left( 0 (00)^* \right)^* = 0^*$$

②

If the regular set A is represented by  $A = (01+1)^*$  and the regular set 'B' is represented by  $B = ((01)^*1^*)^*$ , which of the following is true? **(GATE - 98)**

- (a)  $A \subset B$
- (b)  $B \subset A$
- (c) A and B are incomparable
- (d)  $A = B$

$$01 = x$$

$$1 = y$$

$$A = (x+y)^*$$

$$B = (x^*y^*)^*$$

$$A = B$$

$$(01+1)^* = (x+y)^*$$

$$((01)^*1^*)^* = (x^*y^*)^*$$

3

The string 1101 does not belong to the set represented by (GATE - 98)

(a)  $110^* (0+1)^*$   $\leftarrow 110(1)$

~~(c)  $(10)^* (01)^* (00+11)^*$~~

$1101 \notin J$

(b)  $1(0+1)^* 101 \leftarrow 1( )^* 101$

~~(d)  $(00+(11)^* 0)^*$~~

$1101 \notin J$

$110110 \neq 1101$   
 SubStrng       $\hookrightarrow$   
 Strng

1101 is string

$\in$  

$$\begin{aligned} 1101 &\in (a) \\ 1101 &\in (b) \end{aligned}$$

4

Let S and T be languages over  $\Sigma = \{a, b\}$  represented by the regular expressions  $(a + b^*)^*$  and  $(a + b)^*$ , respectively.

P  
W

Which of the following is true? (GATE - 2000)

(a)  $S \subset T$

(b)  $T \subset S$

(c)  $S = T$

(d)  $S \cap T = \emptyset$

$$S \subseteq T \quad \checkmark \quad (a+b^*)^* = (a+b)^*$$

$$S \supseteq T \quad \checkmark \quad (a^*+b)^* = (a+b)^*$$

$$\begin{array}{ll} \times S \cap T = \emptyset \\ \times S \cup T = \emptyset \\ \checkmark S \cap T = T \end{array}$$

$$\begin{array}{ll} \times S \subset T \\ \checkmark S \subseteq T \\ \times S \supset T \\ \checkmark S \supset T \end{array}$$

$$\begin{array}{ll} \times S \cap T = \emptyset \\ \times S \cup T = \emptyset \\ \checkmark S \cap T = S \\ \checkmark S \cup T = T \\ \checkmark S \cap T = S \end{array}$$

5  
H.W.

Consider the set  $\Sigma^*$  of all strings over the alphabet  $\Sigma = \{0, 1\}$ .  
 $\Sigma^*$  with the concatenation operator for strings (GATE - 03)

P  
W

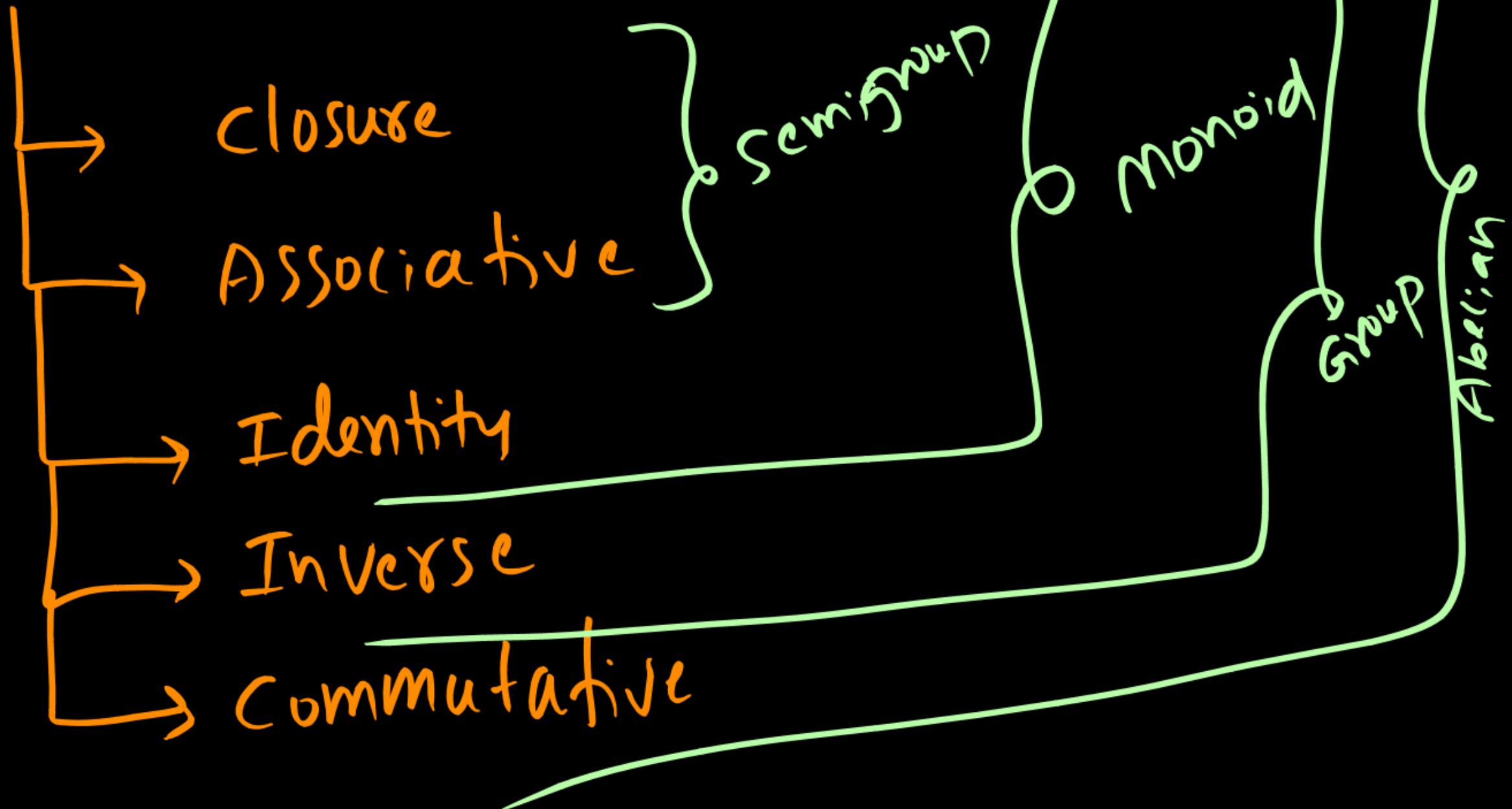
- (a) Does not form a group
- (b) Forms a non-commutative group
- (c) Does not have a right identity element
- (d) Forms a group if the empty string is removed from  $\Sigma^*$

H.W. : Group 9

(Indirect hints)  
Group Theory  
closure ?  
Associative ?  
Identity ?  
Inverse ?  
Commutative ?

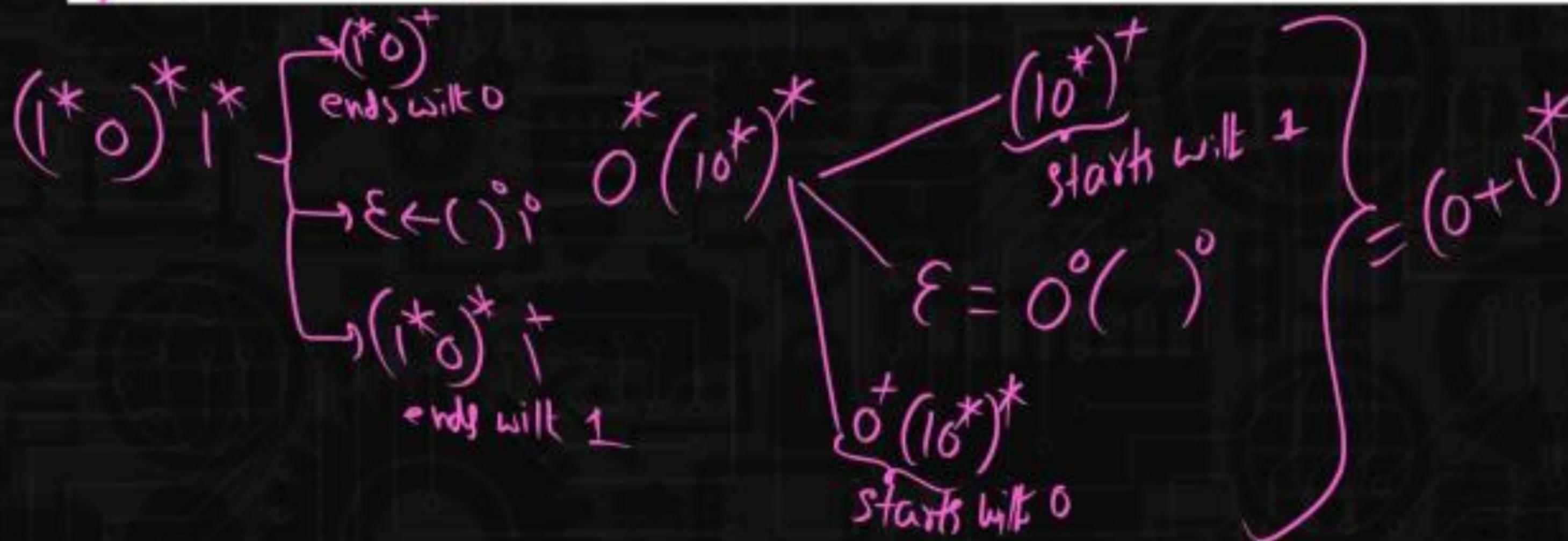
P  
W

# Group Theory



- 6) The regular expression  $0^*(10^*)^*$  denotes the same set as (GATE - 03)

- (a)  ~~$(1^*0)^*1^* = (0+1)^*$~~       (b)  ~~$0^+(0+10)^*$~~
- (c)  ~~$(0+1)^*10(0+1)^*$~~       (d) None of the above



$$(01^*)^+ \rightarrow 0X$$

$$(10^*)^+ \rightarrow 1X$$

$$(1^*0)^+ \rightarrow X_0$$

$$(0^*1)^+ \rightarrow X_1$$

$$(01^*)^* = 0X + \epsilon$$

Not starting with 1

$$(10^*)^* = 1X + \epsilon$$

not starting with 0

$$(1^*0)^* = X_0 + \epsilon$$

not ending with 1

$$(0^*1)^* = X_1 + \epsilon$$

not ending with 0

$$L = 0(0+1)^*$$

$$\bar{L} = \epsilon + 1(0+1)^* = \underbrace{(10^*)}_\text{Not starts with 0}^*$$

$$L \cup \bar{L} = \Sigma^*$$

$$L \cap \bar{L} = \emptyset$$

$$i^* \underbrace{(0i^*)^*}_{\text{Not starting with } 1} = (0+1)^*$$

$i^0 \rightarrow \text{Not starting with } 1$

$i^+ \rightarrow \text{Starting with } 1$

$$\begin{aligned} & \overset{*}{\circ} (\overset{*}{\circ})^* \\ & \overset{*}{\circ} (\overset{*}{\circ})^* \\ & (\overset{*}{\circ} \overset{*}{\circ})^* \quad \overset{*}{-} \\ & (\overset{*}{\circ} \overset{*}{\circ})^* \quad \overset{*}{\circ} \end{aligned}$$

A diagram showing four expressions in yellow. The first two are grouped by a curly brace at the top right. The third and fourth are grouped by a curly brace at the bottom right. The fifth expression is isolated on the right.

P  
W

$$0 \times + 1 \times + \epsilon = \times$$

$$\times_0 + \times_1 + \epsilon = \times$$

This is \*

7)

Which one of the following languages over the alphabet  $\{0, 1\}$  is described by the regular expression  $(0+1)^*0(0+1)^*0(0+1)^*$  (GATE - 09)

- X o X o X
- (a) The set of all strings containing the substring 00  $\Rightarrow$  XooX
  - (b) The set of all strings containing at most two 0's  $\Rightarrow$   $\cancel{(0+1)^*} (0+1)^* (0+1)^*$
  - (c) The set of all strings containing at least two 0's
  - (d) The set of all strings that begin and end with either 0 or 1  $\Rightarrow (0+1)^*$

at least two 0's

Explain

8) Consider the languages  $L_1 = \phi$  and  $L_2 = \{a\}$ . Which one of the following represents  $L_1 L_2^* \cup L_1^*$ ? (GATE - 13)

(a)  $\{\epsilon\}$

$$\underbrace{\phi \cdot L_2^*}_\phi \cup \phi^*$$

(c)  $a^*$

$$\phi \cup \{\epsilon\}$$

(b)  $\phi$

(d)  $\{\epsilon, a\}$

$$L_1 L_2^* \cup L_1^*$$

$$\underbrace{\phi \cdot \phi^*}_\phi \cup \phi^*$$

$$\phi \cup \{\epsilon\} \Rightarrow \{\epsilon\}$$

$$\boxed{\phi^* = \epsilon}$$

$$\boxed{\phi^+ = \phi}$$

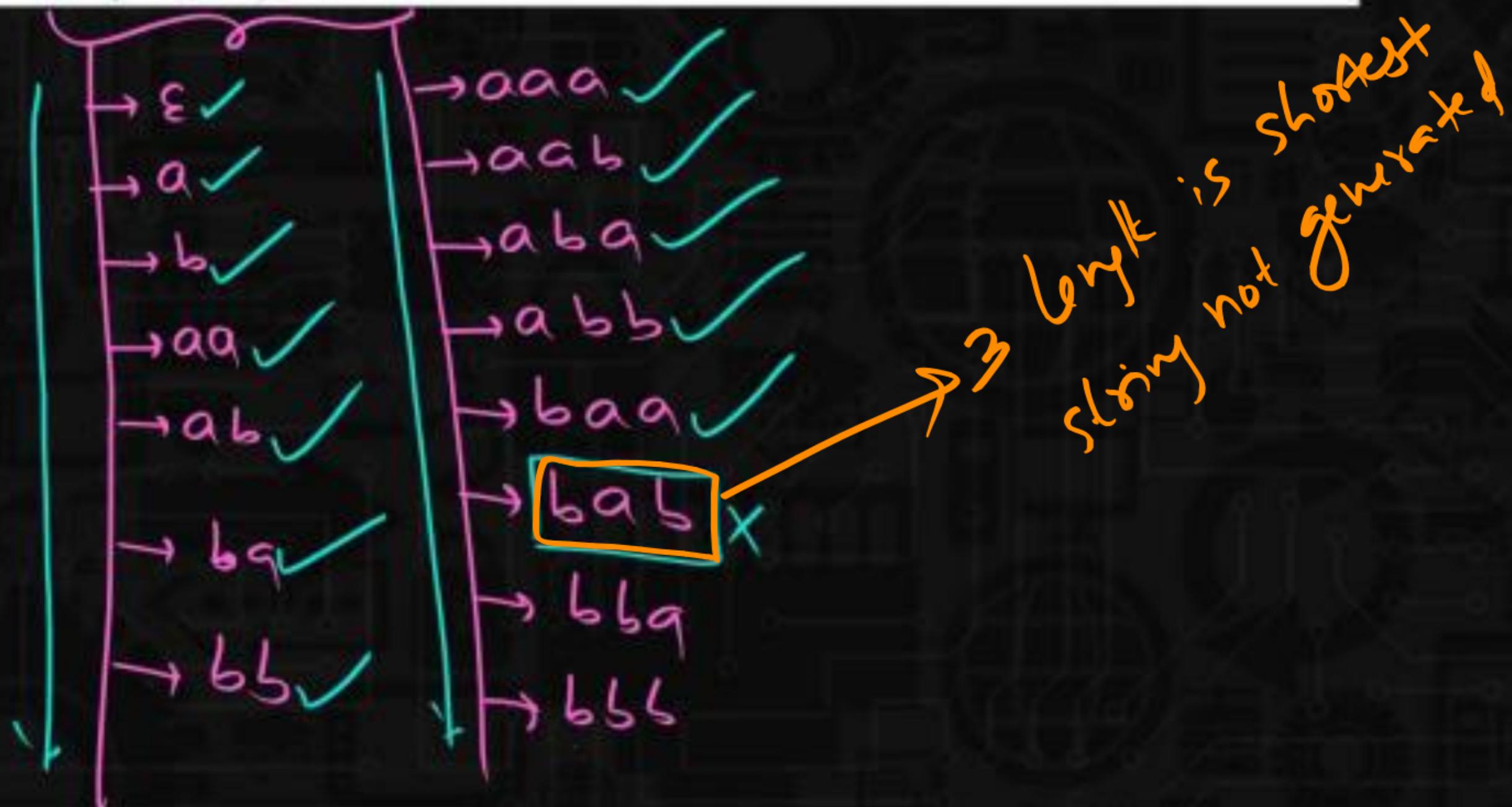
9

The length of the shortest string NOT in the language (over  $\Sigma = \{a, b\}$ ) of the following regular expression is 3.

P  
W

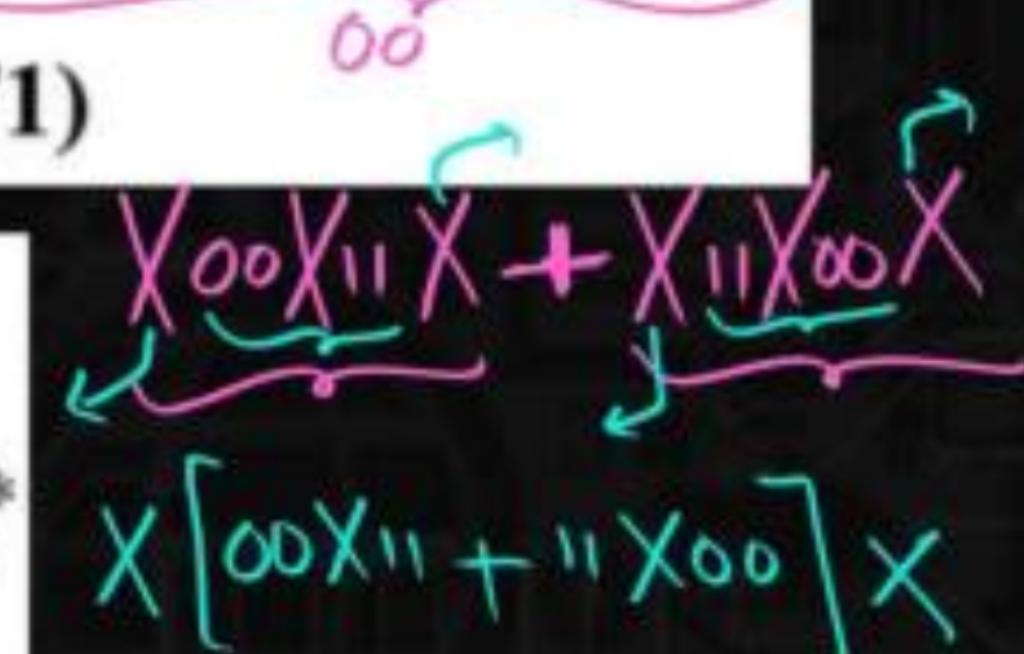
(GATE - 14-SET3)

$$a^*b^*(ba)^*a^*$$



Which one of the following regular expressions represents the language: the set of all binary strings having two consecutive 0's and two consecutive 1's? (GATE - 16 - SET1)

- (a)  ~~$(0+1)^* \underline{00} \underline{11} (0+1)^* + (0+1)^* \underline{11} \underline{00} (0+1)^*$~~
- (b)  ~~$(0+1)^* (\underline{00}(0+1)^* \underline{11} + \underline{11}(0+1)^* \underline{00})(0+1)^*$~~
- (c)  ~~$(0+1)^* \underline{00}(0+1)^* + (0+1)^* \underline{11}(0+1)^*$~~
- (d)  ~~$\underline{00}(0+1)^* \underline{11} + \underline{11}(0+1)^* \underline{00}$~~



Contain 00  
and  
Contain 11

4)

Let  $r = 1(1+0)^*$ ,  $s = 11^*0$  and  $t = 1^*0$  be three regular expressions. Which one of the following is true? (GATE - 91)

- (a)  $L(s) \subseteq L(r)$  and  $L(s) \subseteq L(t)$
- (b)  $L(r) \not\subseteq L(s)$  and  $L(s) \subseteq L(t)$
- (c)  $L(s) \subseteq L(t)$  and  $L(s) \subseteq L(r)$
- (d)  $L(t) \not\subseteq L(s)$  and  $L(s) \subseteq L(r)$ .

$$\left. \begin{array}{l} r = 1(1+0)^* \\ s = 11^*0 \\ t = 1^*0 = 0 + 1^*0 \end{array} \right\} s \subseteq t$$

sct

↓

sct

(12)

Which of the following regular expression identities are true?

(GATE - 92)

(a)  $r^*(*) = r^*$

(c)  $(r+s)^* = r^* + s^*$

(b)  $(r^*s^*)^* = (r+s)^*$

(d)  $r^*s^* = r^* + s^*$

$$\delta^*$$

$$(\delta)^*$$

$$\gamma^{(*)}$$

not deg exp  
not wff

$$X = \gamma + S$$

$$Y = \gamma^* S^*$$

$$Z = S^* \gamma^*$$

$$X^* = Y^* = Z^*$$

$$\gamma \neq Y \neq Z$$

13) Which one of the following regular expressions represents the set of all binary strings with an odd number of 1's?

- A.  $((0 + 1)^* 1 (0 + 1)^* 1)^* 1 0^*$
- B.  $(0^* 1 0^* 1 0^*)^* 0^* 1$
- C.  $1 0^* (0^* 1 0^* 1 0^*)^*$
- D.  $(0^* 1 0^* 1 0^*)^* 1 0^*$

$P(x^{n-1})$

GATE 2020

$$\begin{aligned} L &= \{ 1, 01, 10, 100, 010, 001, 111, \dots \} \\ \therefore R &= 0^* (0^* | 0^* 1 | 0^* 1 0^*)^* 0^* | 0^* \end{aligned}$$

14)

Which one of the following regular expressions over  $\{0, 1\}$  denotes the set of all strings **not** containing 100 as a substring?

(GATE - 97)

- (a)  $0^*(1+0)^*$
- (c)  $0^*1^*01^*$  *C Actual*

- (b)  $0^*\ 1010^*$
- (d)  $0^* (10+1)^*$  *= Actual*

$\bar{L}^+$

$$L = \{ \epsilon, 0, 1, 00, 01, 10, 11, 000, 001, 010, 011, 101, 110, 111, 0000, 0001, \dots \}$$

Model-6:

H.V.

⑤

$$L = b^* \bar{a}^*$$

⑥

$$L = \bar{b}^* \bar{a}^*$$

⑦

$$L = b^* \bar{a}^+$$

⑧

$$L = \bar{b}^* \bar{a}$$

$$L = b^* \bar{a}^+ \bar{b}^*$$

⑨

$$L = \bar{a}^+ \bar{b}^+ c^+$$

⑩

$$L = \bar{b}^+ \bar{a}^+ c^+$$

⑪

$$L = \bar{c}^+ \bar{a}^+ \bar{b}^+$$

Final

Model-7:

$$\textcircled{19} \quad \left\{ a^m b^n \mid m = \text{even}, n = \text{odd} \right\}$$

Model-8:

$$\textcircled{3} \quad \{\omega \mid " , n_a(\omega) \leq 1, n_b(\omega) = 2\}$$

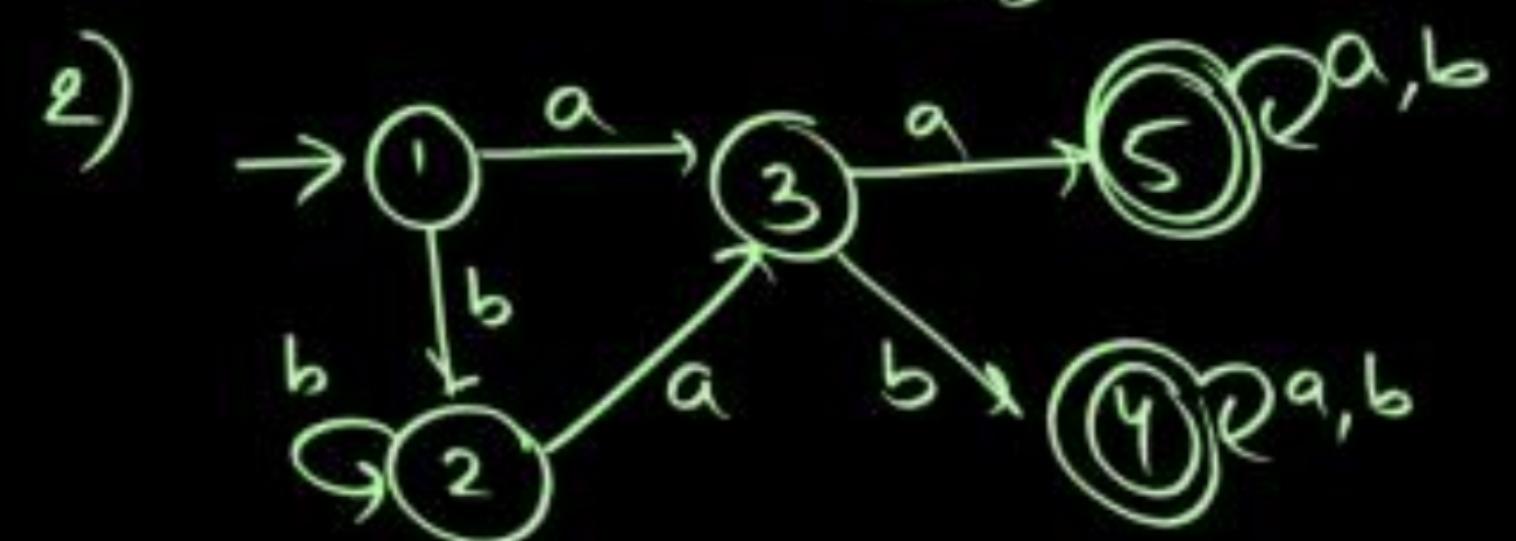
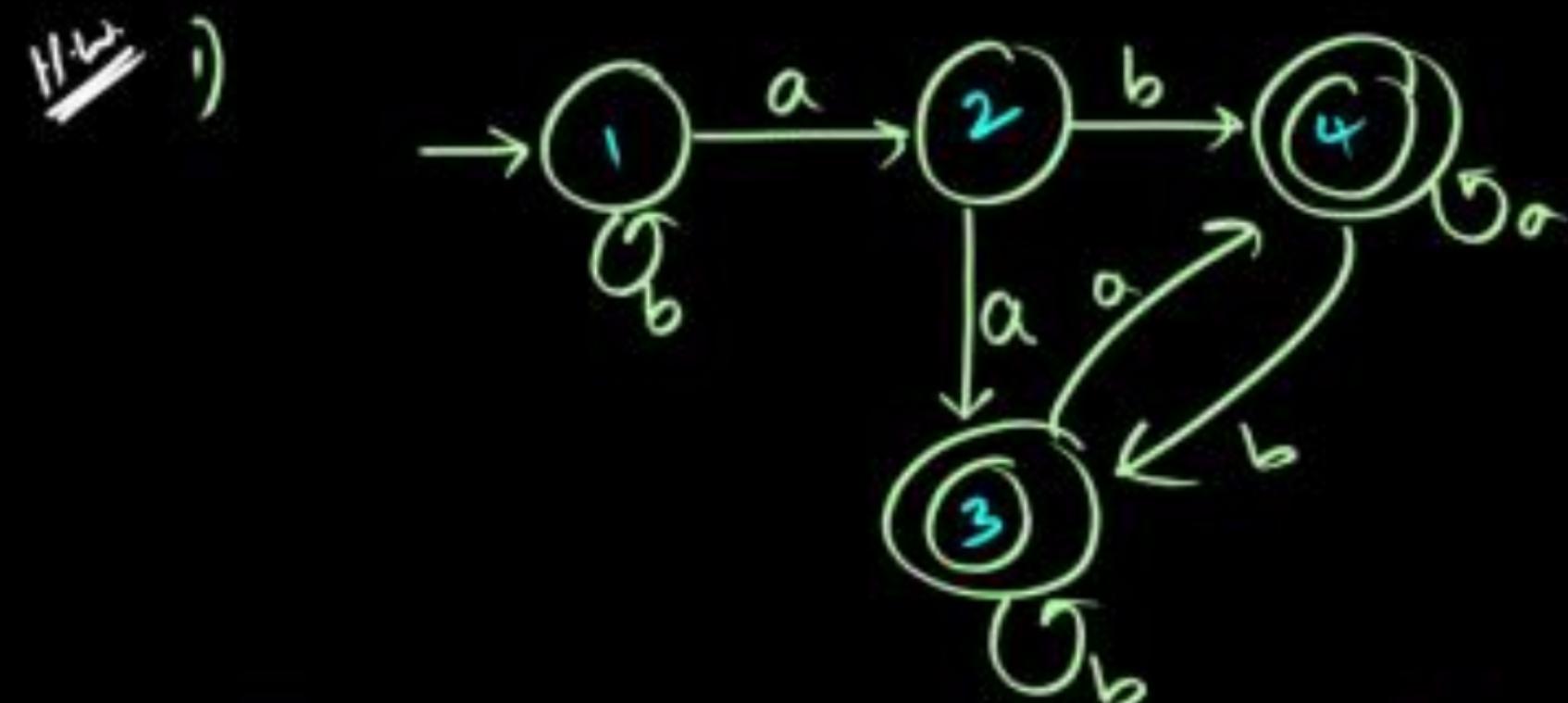
$$\textcircled{4} \quad \{\omega \mid " , n_a(\omega) \leq 1, n_b(\omega) \leq 2\}$$

$$\textcircled{5} \quad \{\omega \mid " , n_a(\omega) \geq 1, n_b(\omega) \leq 2\}$$

Model-12:

- ④  $L = \{w \mid \text{" , 2^{nd} symbol is 'a' AND 4^{th} symbol is 'b'}\}$
- ⑤  $L = \{w \mid \text{" , 2^{nd} symbol is 'a' OR 4^{th} symbol is 'b'}\}$

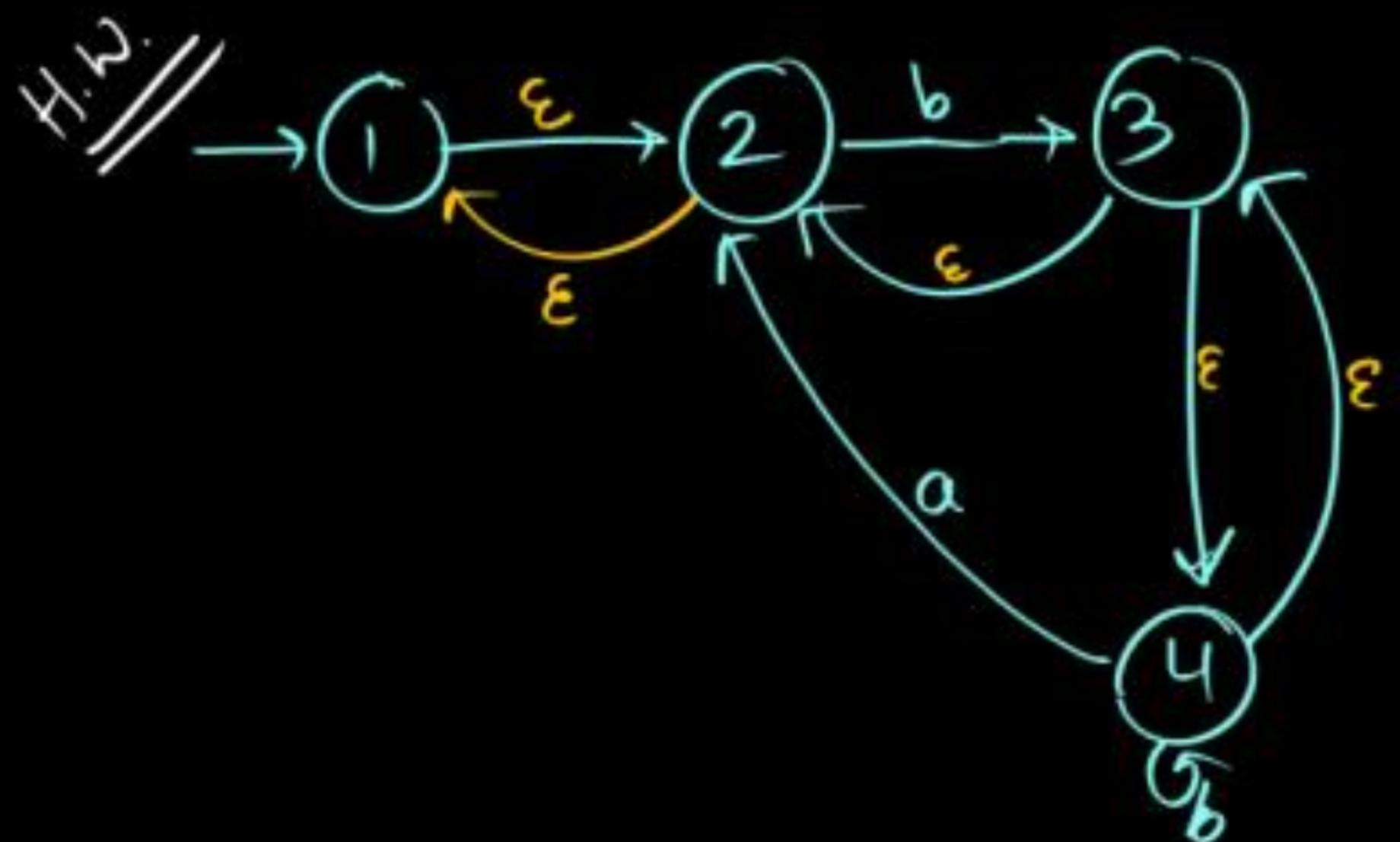
## DFA Minimization



Model-13:

part ③

$$\left\{ \omega \mid \omega \in \{0,1\}^*, \text{ every prefix } P \text{ of } \omega \text{ satisfies } \right. \\ \left. \left\{ \epsilon, 0, 1, \cancel{00}, \alpha, 10, \cancel{X}, \dots \mid |n_0(P) - n_1(P)| \leq 1 \right\} \right\}$$



$$\delta(1, aab) =$$

$$\delta(2, bb) =$$

$$\delta(3, bab) =$$

$$\delta(4, ab) =$$

Let  $L$  be the language represented by the regular expression  $\Sigma^*0011\Sigma^*$  where  $\Sigma = \{0,1\}$ . What is the minimum number of states in a DFA that recognizes  $\bar{L}$  (complement of  $L$ )?

(GATE – 15 – SET3)

- (a) 4
- (b) 5
- (c) 6
- (d) 8

Consider the language  $L$  given by the regular expression  $(a+b)^*b(a+b)$  over the alphabet  $\{a, b\}$ . The smallest number of states needed in a deterministic finite-state automaton (DFA) accepting  $L$  is \_\_\_\_\_. **(GATE – 17 – SET1)**

The minimum possible number of states of a deterministic finite automaton that accepts the regular language **(GATE – 17 – SET2)**

$L = \{w_1aw_2 \mid w_1, w_2 \in \{a, b\}^*, |w_1|=2, |w_2| \geq 3\}$  is \_\_\_\_\_.

The lexical analysis for a modern computer language such as java needs the power of which one of the following machine models in a necessary and sufficient sense? **(GATE - 11)**

- (a) Finite state automata
- (b) Deterministic pushdown automata
- (c) Non-deterministic pushdown automata
- (d) Turing machine

The number of substrings (of all lengths inclusive) that can be formed from a character string of length n is (GATE - 89 & 94)

- (a) n
- (b)  $n^2$
- (c)  $\frac{n(n-1)}{2}$
- (d)  $\frac{n(n+1)}{2} + 1$

Consider the following language.

$$L = \{x \in \{a, b\}^* \mid \text{number of } a\text{'s in } x \text{ divisible by 2 but not divisible by 3}\}$$

The minimum number of states in DFA that accepts  $L$  is \_\_\_\_\_

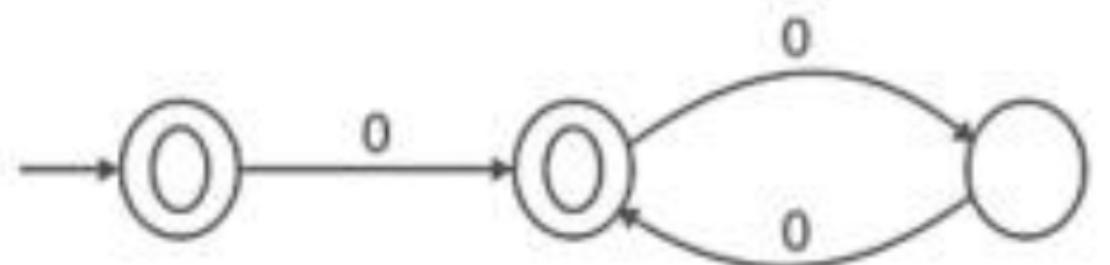
Given a language  $L$ , define  $L^i$  as follows:

$$L^0 = \epsilon$$

$$L^i = L^{i-1} \cdot L \text{ For all } i > 0$$

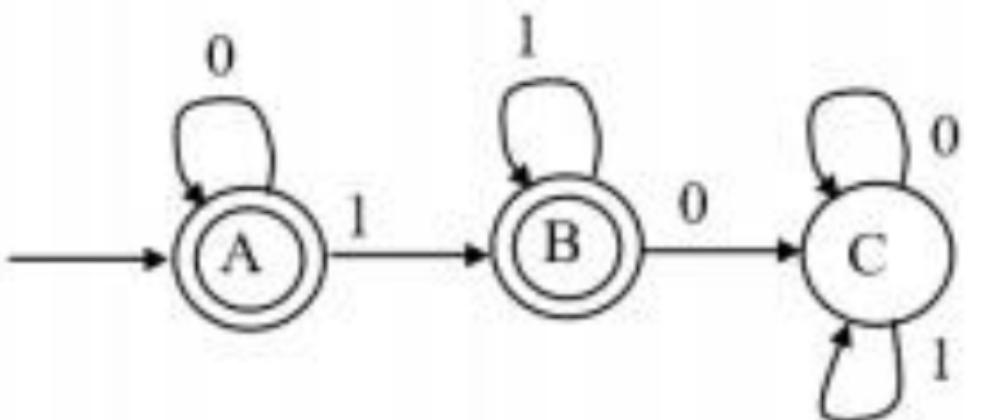
The order of a language is defined as the smallest such that  $L^k =$

$L^{k+1}$ . Consider the language  $L_1$  (over alphabet 0) accepted by the following automaton:



The order of  $L_1$  is \_\_\_\_\_

The regular expression for the language recognized by the finite state automation of the below figure is \_\_\_\_\_ (GATE - 94)



Let  $L$  be the set of all binary strings whose last two symbols are the same .The number of states in the minimum state deterministic finite-state automaton accepting  $L$  is

(GATE - 98)

What can be said about a regular language L over {a} whose minimal finite state automaton has two states? **(GATE - 2000)**

- (a) L must be  $\{a^n \mid n \text{ is odd}\}$
- (b) L must be  $\{a^n \mid n \text{ is even}\}$
- (c) L must be  $\{a^n \mid n \geq 0\}$
- (d) Either L must be  $\{a^n \mid n \text{ is odd}\}$  or L must be  $\{a^n \mid n \text{ is even}\}$

Consider a DFA over  $\Sigma = \{a, b\}$  accepting all strings which have number of a's divisible by 6 and number of b's divisible by 8. What is the minimum number of states that the DFA will have?

**(GATE - 01)**

- (a) 8
- (b) 14
- (c) 15
- (d) 48

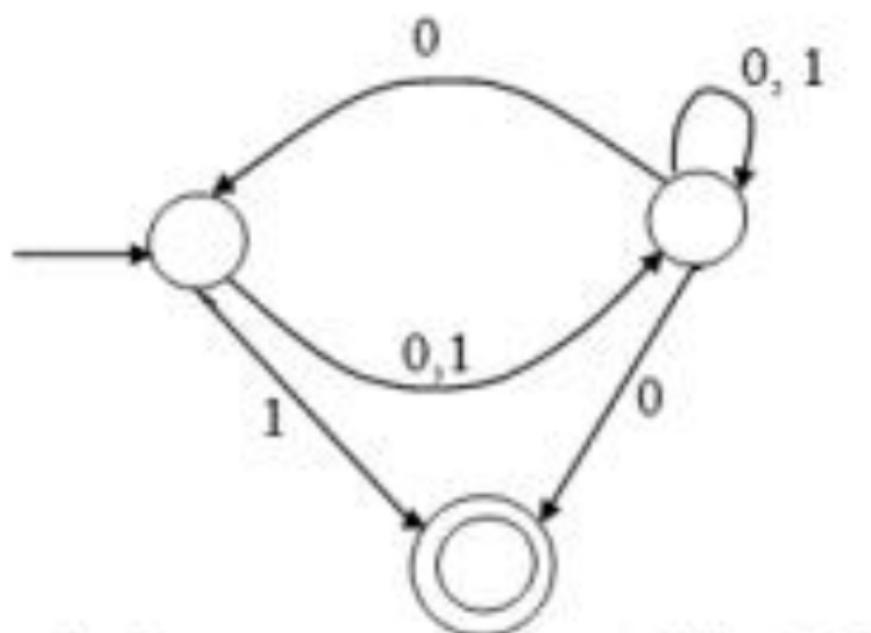
The smallest finite automation which accepts the language

$L = \{x \mid \text{length of } x \text{ is divisible by 3}\}$  has

(GATE - 02)

- (a) 2 states
- (b) 3 states
- (c) 4 states
- (d) 5 states

Consider the NFA M shown below.



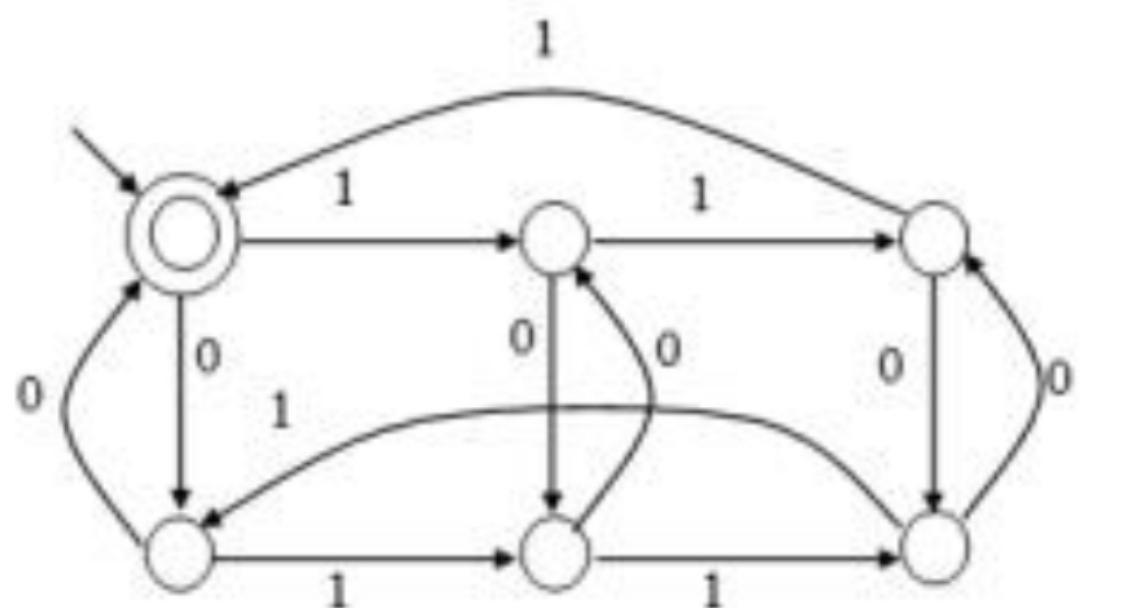
Let the language accepted by M be L. Let  $L_1$  be the language accepted by the NFA,  $M_1$  obtained by changing the accepting state of M to a non-accepting state and by changing the non-accepting state of M to accepting states. Which of the following statements is true?

- (a)  $L_1 = \{0, 1\}^* - L$   
(c)  $L_1 \subseteq L$

- (b)  $L_1 = \{0, 1\}^*$   
(d)  $L_1 = L$

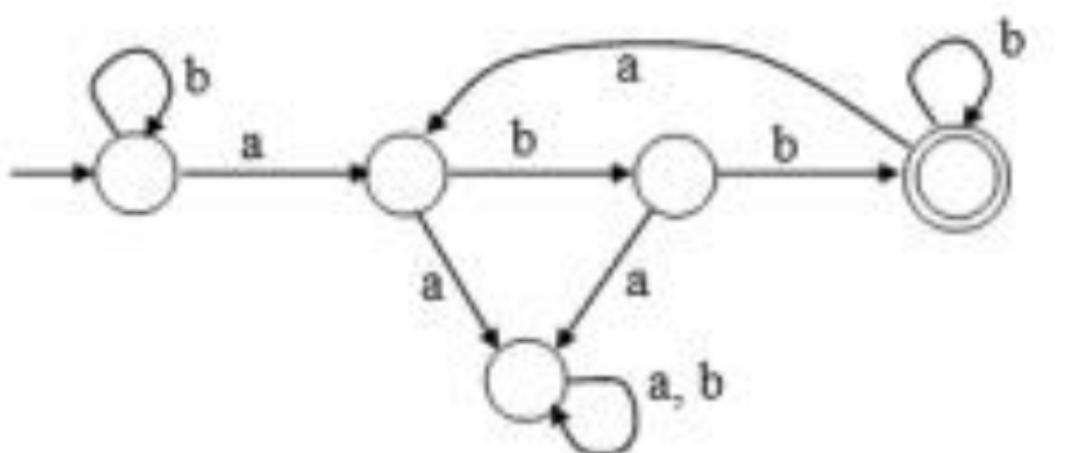
The following finite state machine accepts all those binary strings in which the number of 1's and 0's are respectively

(GATE - 04)



- (a) Divisible by 3 and 2
- (b) Odd and even
- (c) Even and odd
- (d) Divisible by 2 and 3

Consider the machine M:



The language recognized by M is:

(GATE - 05)

- (a)  $\{w \in \{a,b\}^* \mid \text{every } a \text{ in } w \text{ is followed by exactly two } b's\}$
- (b)  $\{w \in \{a,b\}^* \mid \text{every } a \text{ in } w \text{ is followed by at least two } b's\}$
- (c)  $\{w \in \{a,b\}^* \mid w \text{ contains the substring 'abb'}\}$
- (d)  $\{w \in \{a,b\}^* \mid w \text{ does not contain 'aa' as a substring}\}$

If  $s$  is a string over  $(0+1)^*$  then let  $n_0(s)$  denote the number of 0's in  $s$  and  $n_1(s)$  the number of 1's in  $s$ . Which one of the following languages is not regular?

**(GATE - 06)**

- (a)  $L = \{s \in (0+1)^* \mid n_0(s) \text{ is a 3-digit prime}\}$
- (b)  $L = \{s \in (0+1)^* \mid \text{for every prefix } s' \text{ of } s, |n_0(s') - n_1(s')| \leq 2\}$
- (c)  $L = \{s \in (0+1)^* \mid |n_0(s) - n_1(s)| \leq 4\}$
- (d)  $L = \{s \in (0+1)^* \mid n_0(s) \bmod 7 = n_1(s) \bmod 5 = 0\}$

Consider the regular language

$L = (111+11111)^*$ . The minimum number of states in any DFA accepting this language is

**(GATE - 06)**

- (a) 3      (b) 5      (c) 8      (d) 9

A minimum state deterministic finite automaton accepting the language

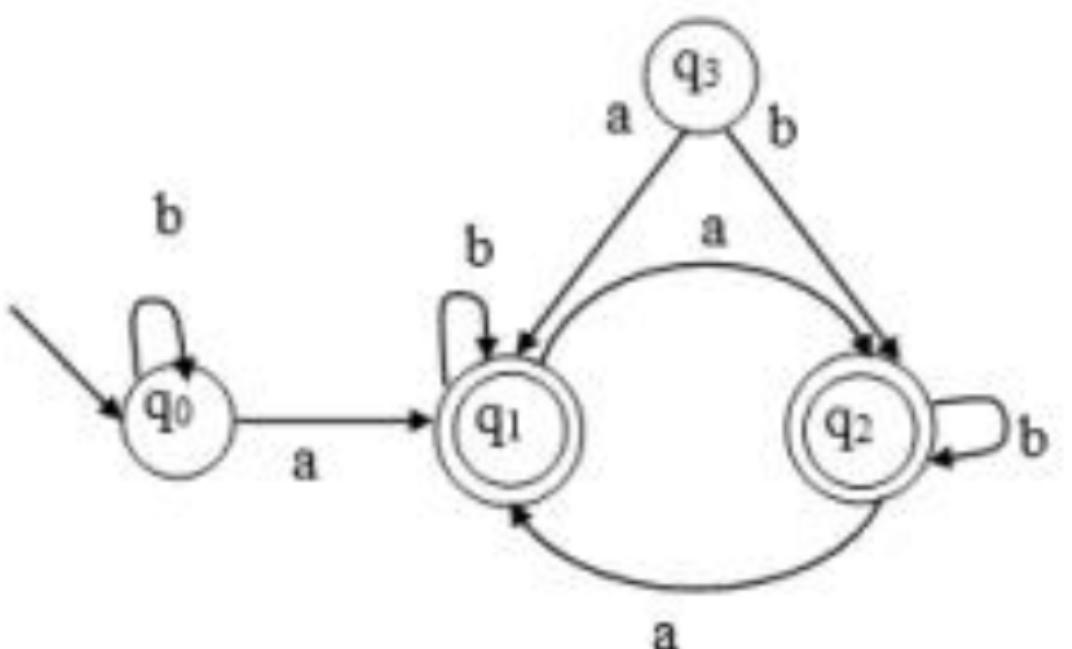
$L = \{w \mid w \in \{0, 1\}^*, \text{number of } 0\text{'s and } 1\text{'s in } w \text{ are divisible by } 3 \text{ and } 5, \text{ respectively}\}$  has

(GATE - 07)

- (a) 15 states
- (b) 11 states
- (c) 10 states
- (d) 9 states

Consider the following finite state automaton

(GATE - 07)

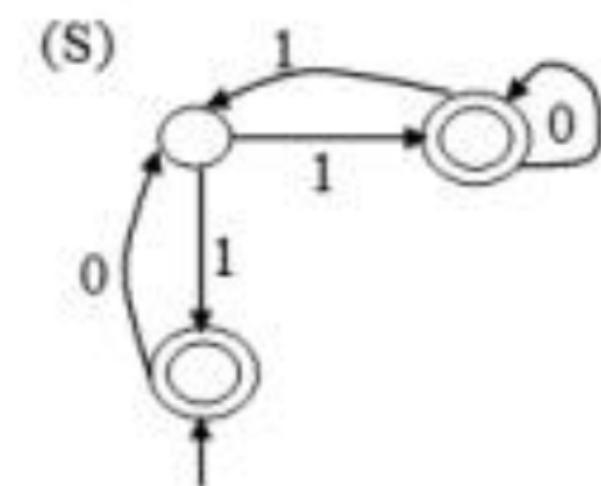
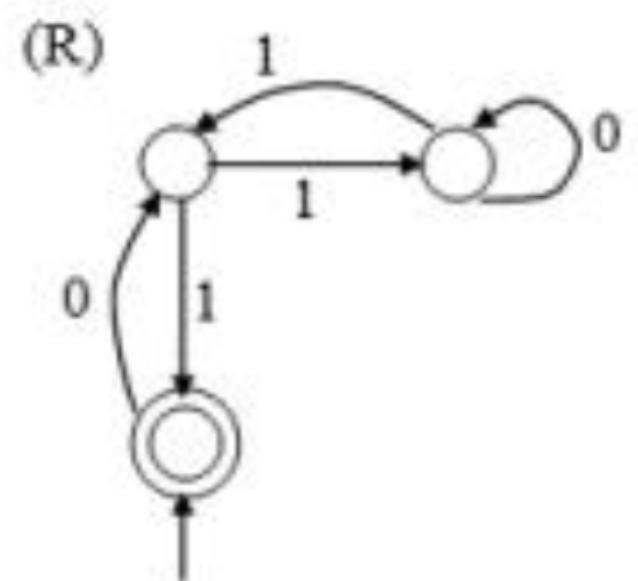
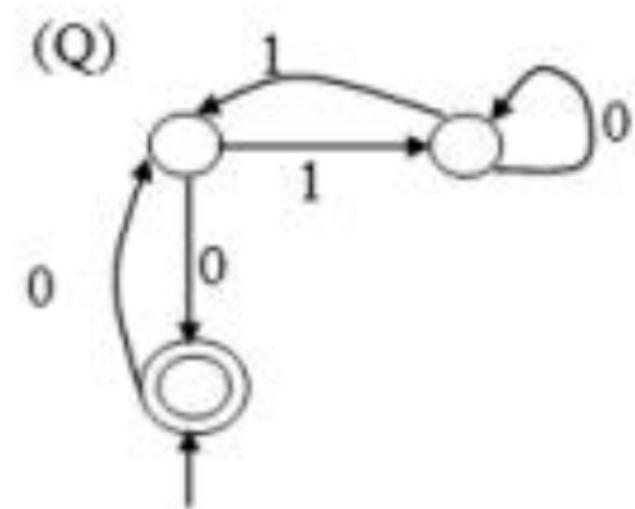
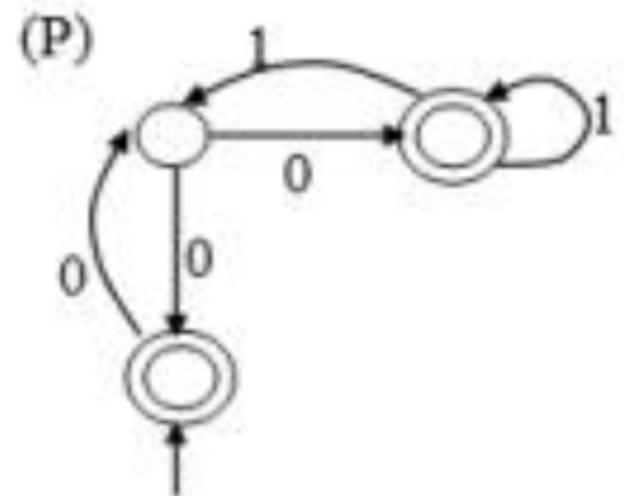


The language accepted by this automaton is given by the regular expression

- (a)  $b^*ab^*ab^*ab^*$
- (b)  $(a+b)^*$
- (c)  $b^*a(a+b)^*$
- (d)  $b^*ab^* ab^*$

Match the following NFAs with the regular expressions they correspond to

P  
W



**(GATE - 08)** 1.  $\epsilon + 0(01^* 1 + 00)^* 01^*$

2.  $\epsilon + 0(10^* 1 + 00)^* 0$

3.  $\epsilon + 0(10^* 1 + 10)^* 1$

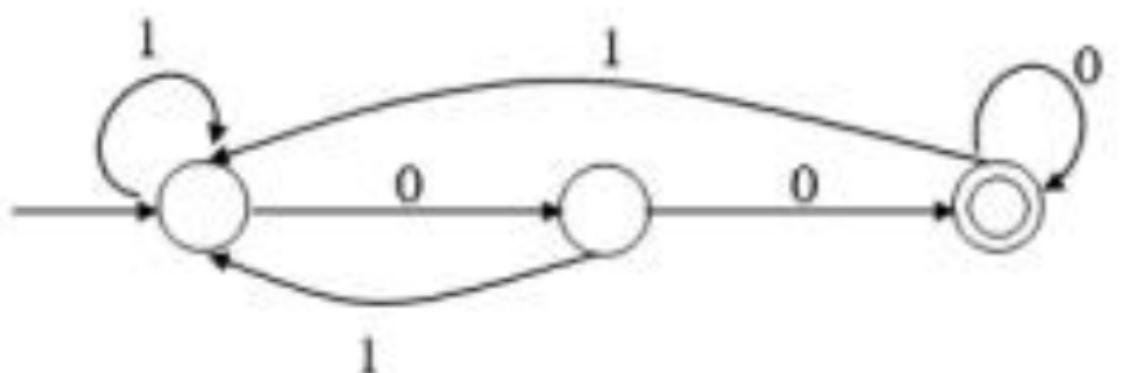
4.  $\epsilon + 0(10^* 1 + 10)^* 10^*$

(a) P-2, Q-1, R-3, S-4

(b) P-1, Q-3, R-2, S-4

(c) P-1, Q-2, R-3, S-4

(d) P-3, Q-2, R-1, S-4



The above DFA accepts the set of all strings over  $\{0, 1\}$  that

**(GATE - 09)**

- (a) Begins either with 0 or 1
- (b) End with 0
- (c) End with 00
- (d) Contains the substring 00.

Let  $L = \{w \in (0+1)^* | w \text{ has even number of } 1's\}$ , i.e  $L$  is the set of all bit strings with even number of 1's. Which one of the regular expressions below represents  $L$ ?

(GATE - 10)

- (a)  $(0^* 1 0^*)^*$
- (b)  $0^* (1 0^* 1 0^*)^*$
- (c)  $0^* (1 0^* 1)^* 0^*$
- (d)  $0^* 1 (1 0^* 1)^* 1 0^*$

Let  $w$  be any string of length  $n$  in  $\{0, 1\}^*$ . Let  $L$  be the set of all substrings of  $w$ . What is the minimum number of states in a non-deterministic finite automation that accepts  $L$ ?

(GATE-10)

- (a)  $n - 1$
- (b)  $n$
- (c)  $n+1$
- (d)  $2^{n-1}$

Definition of the language L with alphabet {a} is given as following.

$$L = \{a^{nk} \mid k > 0, \text{ and } n \text{ is a positive integer constant}\}$$

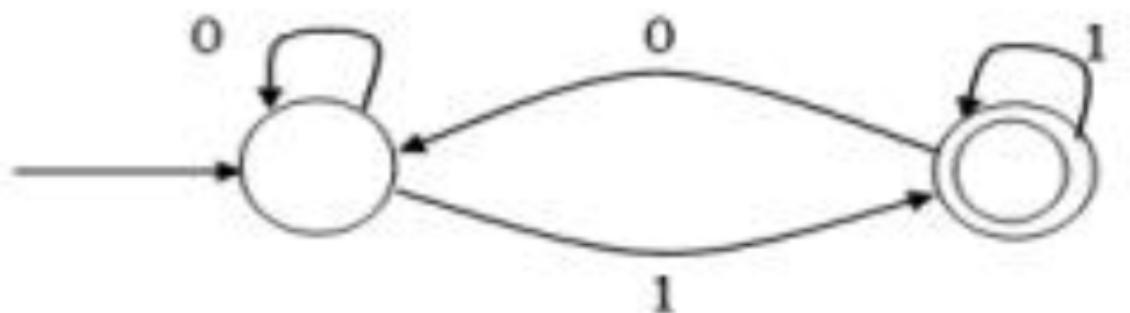
What is the minimum number of states needed in a DFA to recognize L?

(GATE - 11)

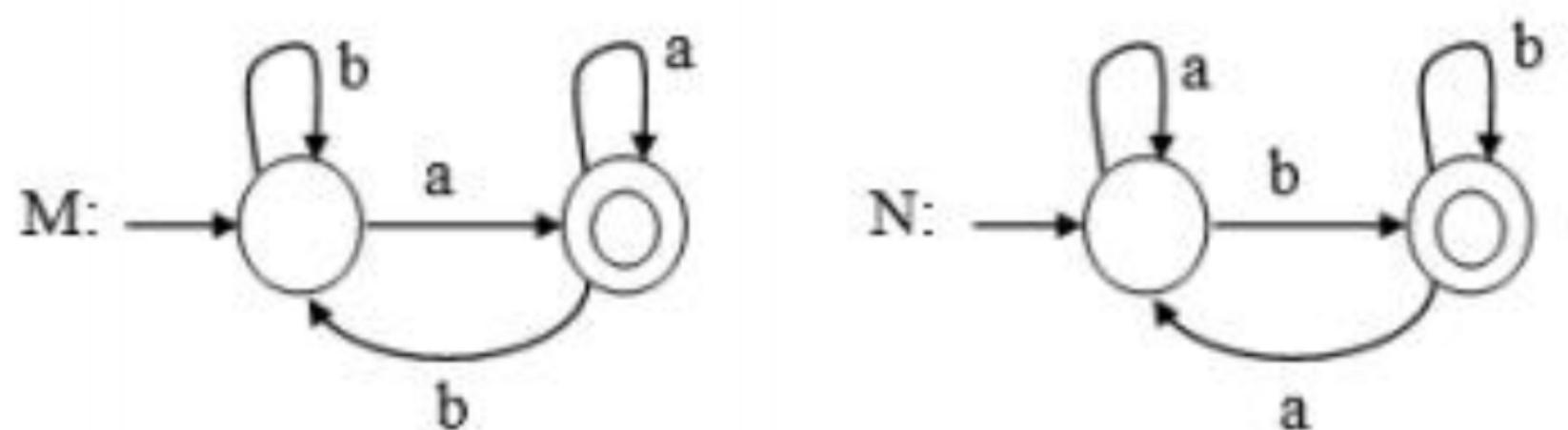
- (a)  $k + 1$
- (b)  $n + 1$
- (c)  $2^{n+1}$
- (d)  $2^{k+1}$

Which of the regular expression given below represent the following DFA?

(GATE – 14-SET1)



- I.  $0^*1(1+00^*1)^*$
  - II.  $0^*1^*1+11^*0^*1$
  - III.  $(0+1)^*1$
- 
- (a) I and II only
  - (b) I and III only
  - (c) II and III only
  - (d) I, II, and III



Consider the DFAs M and N given above. The number of states in a minimal DFA that accepts the languages  $L(M) \cap L(N)$  is \_\_\_\_\_.

**(GATE – 15 – SET1)**

Consider the alphabet  $\Sigma = \{0, 1\}$ , the null/empty string  $\lambda$  and the set of strings  $X_0$ ,  $X_1$ , and  $X_2$  generated by the corresponding non-terminals of a regular grammar  $X_0$ ,  $X_1$ , and  $X_2$  are related as follows.

$$X_0 = 1 X_1$$

$$X_1 = 0 X_1 + 1 X_2$$

$$X_2 = 0 X_1 + \{\lambda\}$$

Which one of the following choices precisely represents the strings in  $X_0$ ?

(GATE - 15- SET2)

- (a)  $10(0^*+(10)^*)1$
- (b)  $10(0^*+(10^*))^*1$
- (c)  $1(0+10)^*1$
- (d)  $10(0+10)^*1 + 110(0+10)^*1$

Let  $\delta$  denote the transition function and  $\hat{\delta}$  denote the extended transition function of the  $\epsilon$ -NFA whose transition table is given below:

$\delta$	$\epsilon$	a	b
$\rightarrow q_0$	$\{q_2\}$	$\{q_1\}$	$\{q_0\}$
$q_1$	$\{q_2\}$	$\{q_2\}$	$\{q_3\}$
$q_2$	$\{q_0\}$	$\phi$	$\phi$
$q_3$	$\phi$	$\phi$	$\{q_2\}$

Then  $\hat{\delta}(q_2, aba)$  is **(GATE - 17 - SET2)**

- (a)  $\phi$
- (b)  $\{q_0, q_1, q_3\}$
- (c)  $\{q_0, q_1, q_2\}$
- (d)  $\{q_0, q_2, q_3\}$

Let  $L = (a+b)(a+b)(a+b)^*$  over alphabet  $\Sigma = \{a, b\}$ . [NAT]

If  $x$  is total number of final states and  $y$  is total number of non-final states, then  $|x-y| = \underline{\hspace{2cm}}$ .

If  $x$  and  $y$  are number of equivalence classes for  $L_1$  and  $L_2$  respectively then  $L_1 + L_2 = \underline{\hspace{2cm}} 0$ .

$$L_1 = aa(a+b)^*$$

$$L_2 = (a+b)^* aa$$

**GATE 1998**

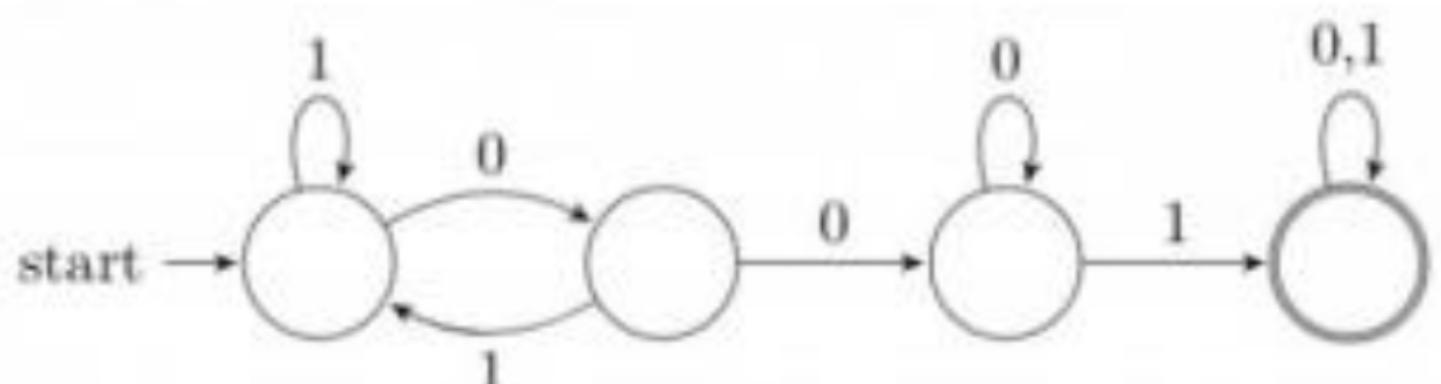
Which of the following set can be recognized by a Deterministic Finite state Automaton?

- A. The numbers  $1, 2, 4, 8, \dots, 2^n, \dots$  written in binary
- B. The numbers  $1, 2, 4, 8, \dots, 2^n, \dots$  written in unary
- C. The set of binary string in which the number of zeros is the same as the number of ones.
- D. The set  $\{1, 101, 11011, 1110111, \dots\}$

Consider a DFA over  $\Sigma=\{a,b\}$  accepting all strings which have number of a's divisible by 6 and number of b's divisible by 8. What is the minimum number of states that the DFA will have?

- A. 8
- B. 14
- C. 15
- D. 48

Consider the following deterministic finite state automaton M.



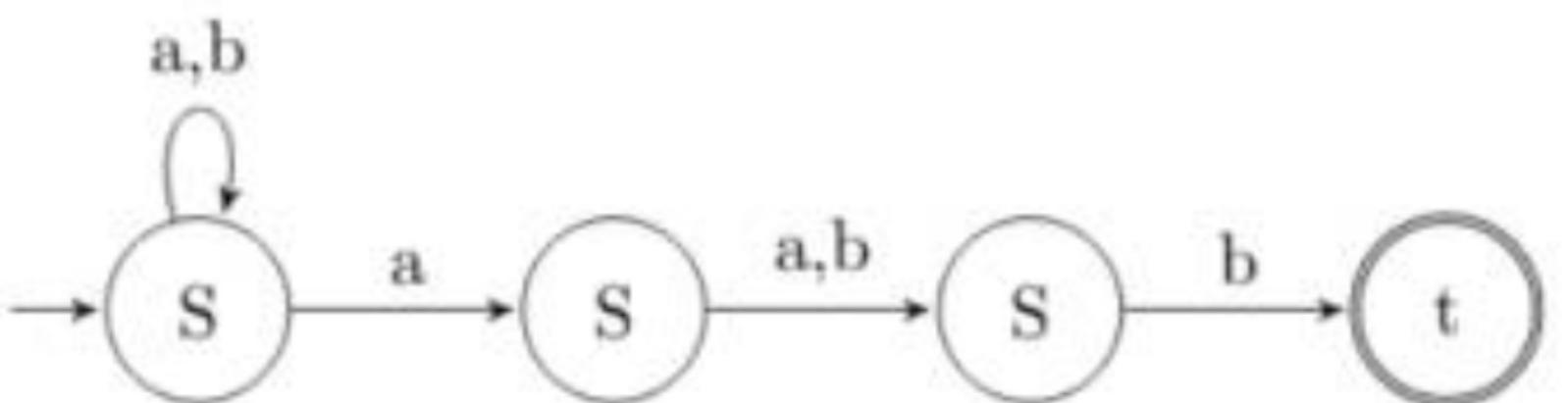
Let  $S$  denote the set of seven bit binary strings in which the first, the fourth, and the last bits are 1. The number of strings in  $S$  that are accepted by  $M$  is

- A. 1
- B. 5
- C. 7
- D. 8

Which one of the following regular expressions is NOT equivalent to the regular expression  $(a+b+c)^*$ ?

- A.  $(a^*+b^*+c^*)^*$
- B.  $(a^*b^*c^*)^*$
- C.  $((ab)^*+c^*)^*$
- D.  $(a^*b^*+c^*)^*$

Which regular expression best describes the language accepted by the non-deterministic automaton below?



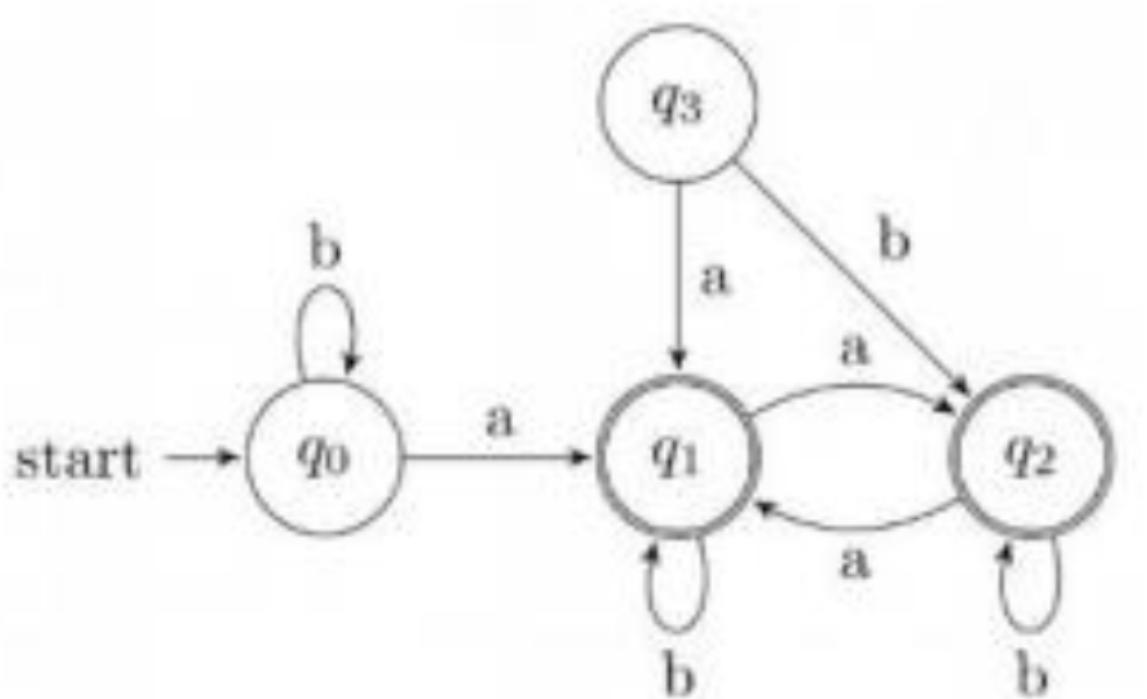
- A.  $(a+b)^* a(a+b)b$
- B.  $(abb)^*$
- C.  $(a+b)^* a(a+b)^* b(a+b)^*$
- D.  $(a+b)^*$

A minimum state deterministic finite automaton accepting the language

$L = \{w | w \in \{0,1\}^*, \text{ number of } 0\text{s and } 1\text{s in } w \text{ are divisible by } 3 \text{ and } 5, \text{ respectively } \}$  has

- A. 15 states
- B. 11 states
- C. 10 states
- D. 9 states

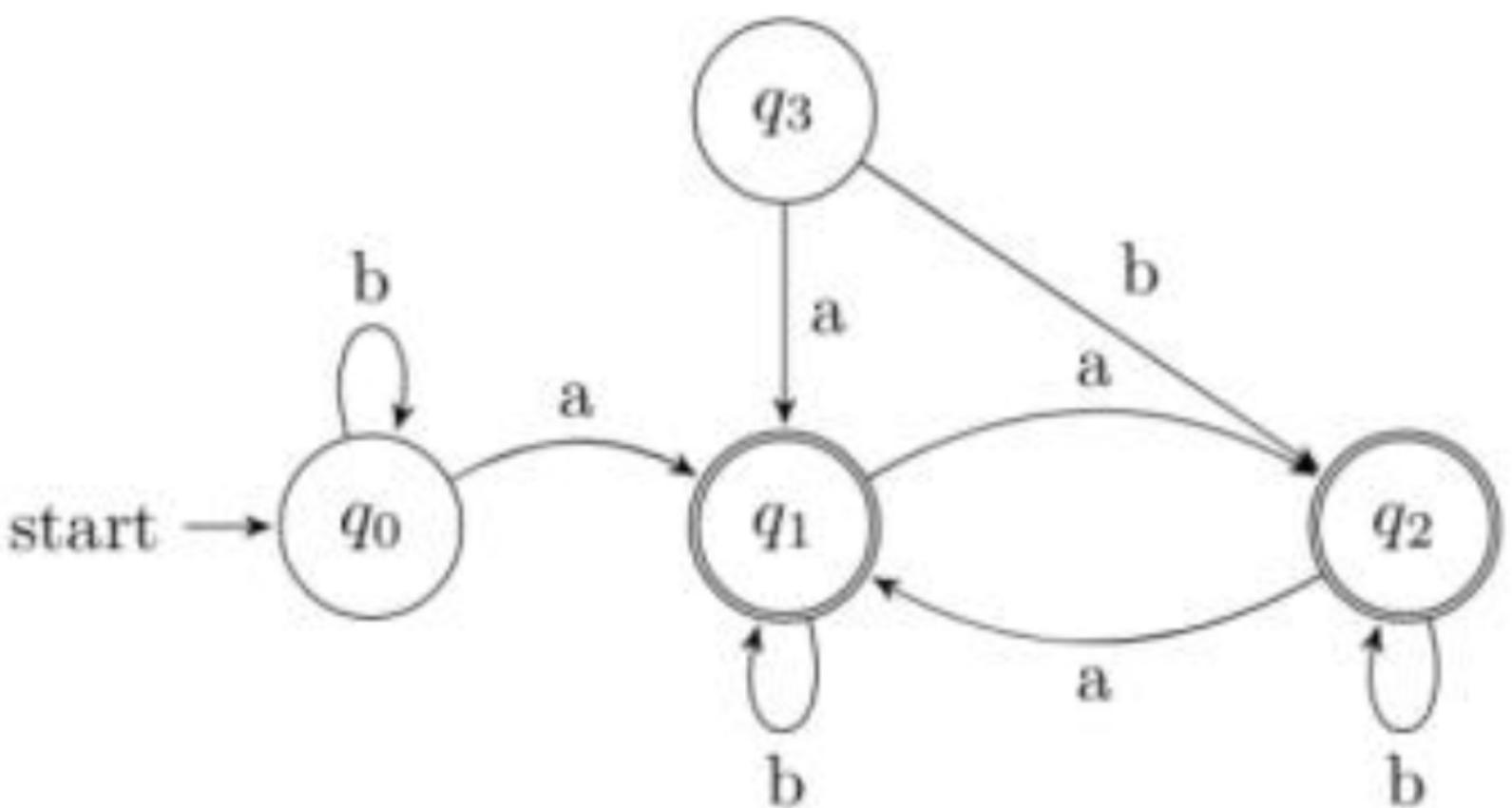
Consider the following Finite State Automaton:



The language accepted by this automaton is given by the regular expression

- A.  $b^*ab^*ab^*ab^*$
- B.  $(a+b)^*$
- C.  $b^*a(a+b)^*$
- D.  $b^*ab^*ab^*$

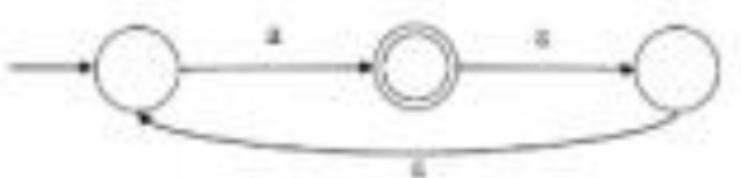
Consider the following Finite State Automaton:



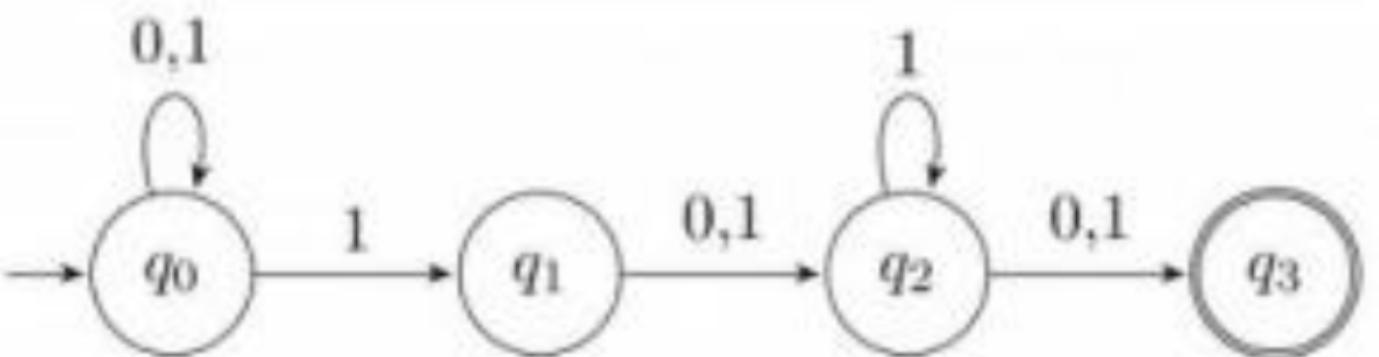
The minimum state automaton equivalent to the above FSA has the following number of states:

- A. 1
- B. 2
- C. 3
- D. 4

What is the complement of the language accepted by the NFA shown below?  
Assume  $\Sigma = \{a\}$  and  $\epsilon$  is the empty string.



- A.  $\emptyset$
- B.  $\{\epsilon\}$
- C.  $a^*$
- D.  $\{a, \epsilon\}$



What is the set of reachable states for the input string 0011?

- A. {q0,q1,q2}
- B. {q0,q1}
- C. {q0,q1,q2,q3}
- D. {q3}

The number of states in the minimal deterministic finite automaton corresponding to the regular expression  $(0+1)^*(10)$  is \_\_\_\_.

The number of states in the minimum sized DFA that accepts the language defined by the regular expression.

$(0+1)^*(0+1)(0+1)^*$  is \_\_\_\_\_.

Let  $\Sigma$  be the set of all bijections from  $\{1, \dots, 5\}$  to  $\{1, \dots, 5\}$ , where  $\text{id}$  denotes the identity function, i.e.  $\text{id}(j) = j, \forall j$ . Let  $\circ$  denote composition on functions. For a string  $x = x_1 x_2 \dots x_n \in \Sigma^n$ ,  $n \geq 0$ , let  $\pi(x) = x_1 \circ x_2 \circ \dots \circ x_n$ . Consider the language  $L = \{x \in \Sigma^* \mid \pi(x) = \text{id}\}$ . The minimum number of states in any DFA accepting  $L$  is \_\_\_\_\_.

