

CS & IT ENGINEERING

Data Structures

Tree

Lecture No.- 03



By- Pankaj Sharma Sir

Topics to be Covered



Topic

Tree Part-03

*binary
trees*

Recap of Previous Lecture



Topic

Tree Part-02

Tree traversal

- pre-order
- In-order
- Post-order



Topic : Tree

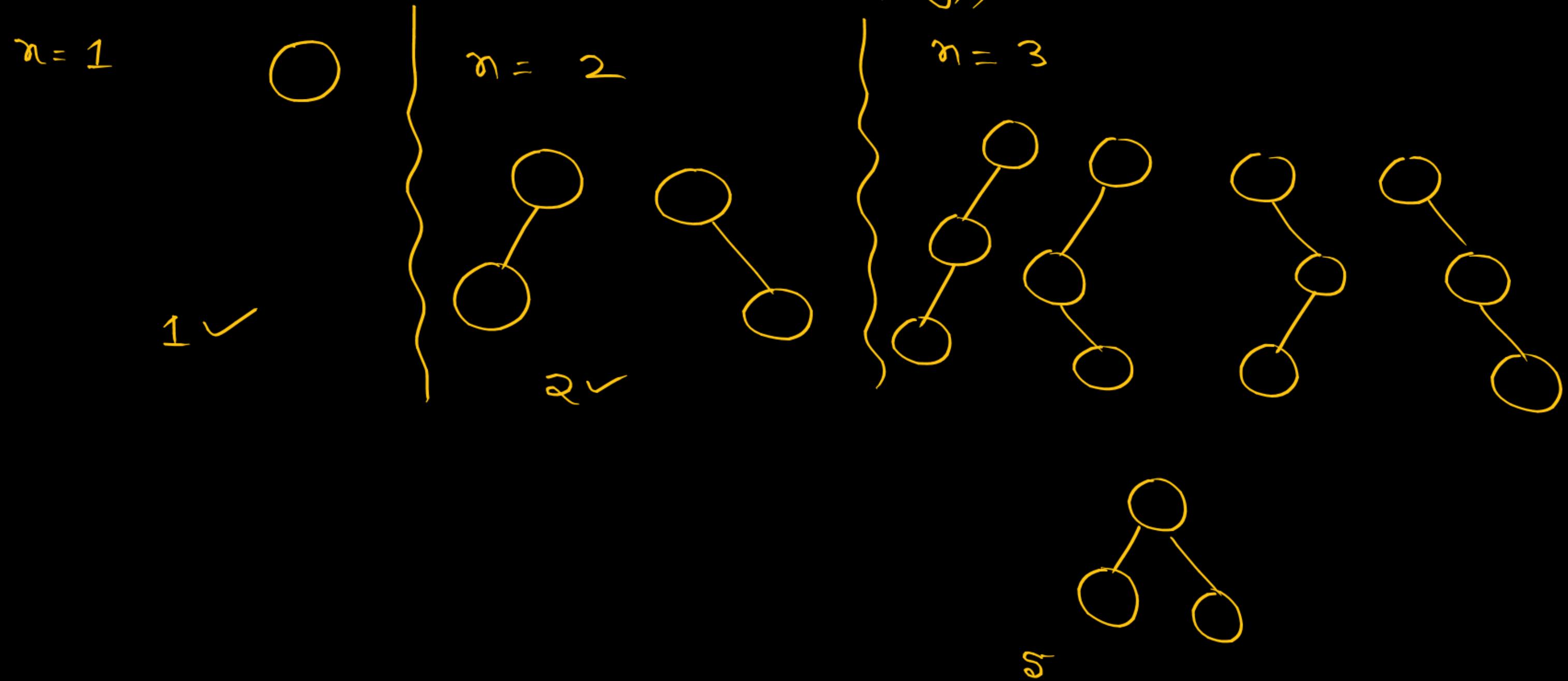
P
W

Pankaj sir PW

- ① Unlabelled
- ② labelled



unlabelled binary trees with n nodes
(shape/structure/topology)



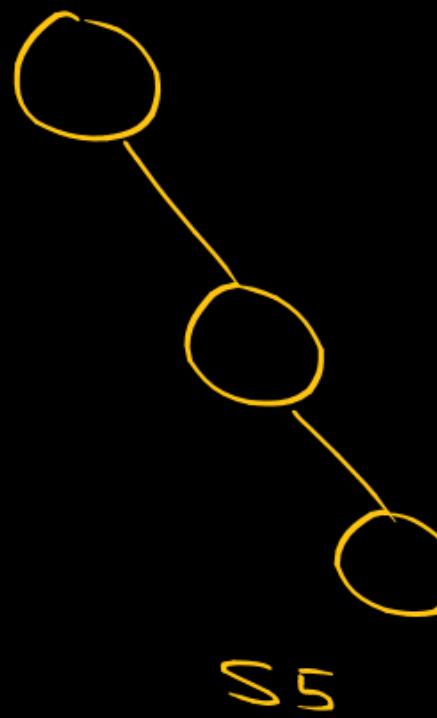
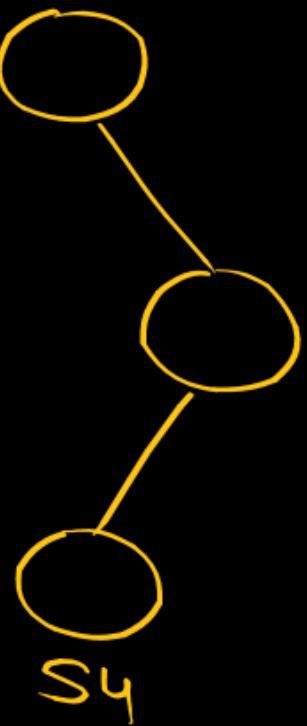
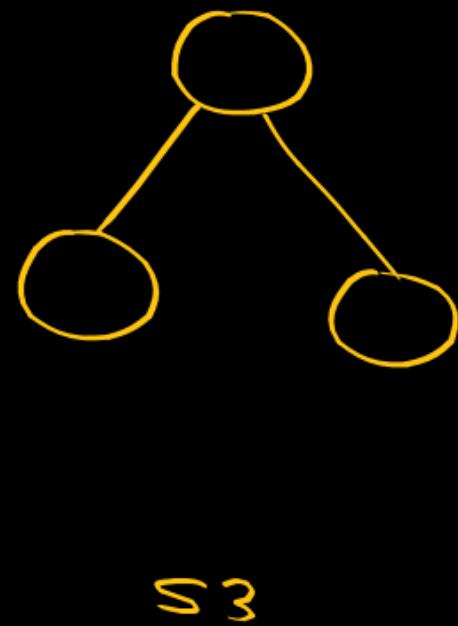
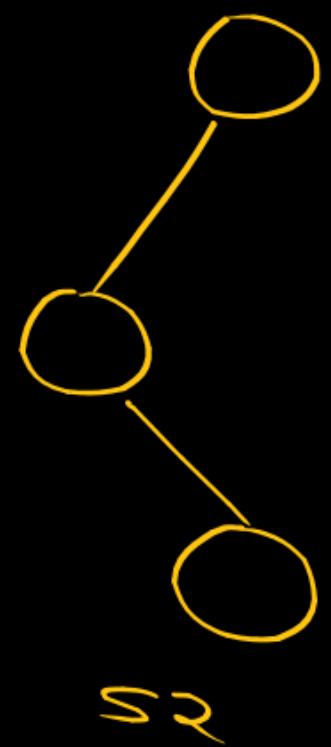
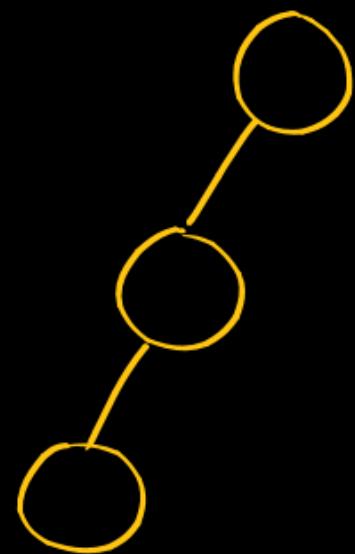
$$\# \text{ unlabelled } \underline{\text{binary}} \text{ tree with } n \text{ nodes} = \frac{{}^n C_n}{n+1}$$

$$n=3$$

$$\frac{{}^6 C_3}{3+1} = \frac{{}^6 C_3}{4} = \frac{1}{4} \times \frac{6!}{3!3!} = \frac{1}{4} \times \frac{6 \times 5 \times \cancel{4} \times \cancel{3} \times \cancel{2} \times \cancel{1}}{\cancel{3} \times \cancel{3} \times \cancel{2} \times \cancel{1}} = 5$$

With 3 nodes \Rightarrow 5 $\underline{\text{binary}}$ tree structure/shape are possible.

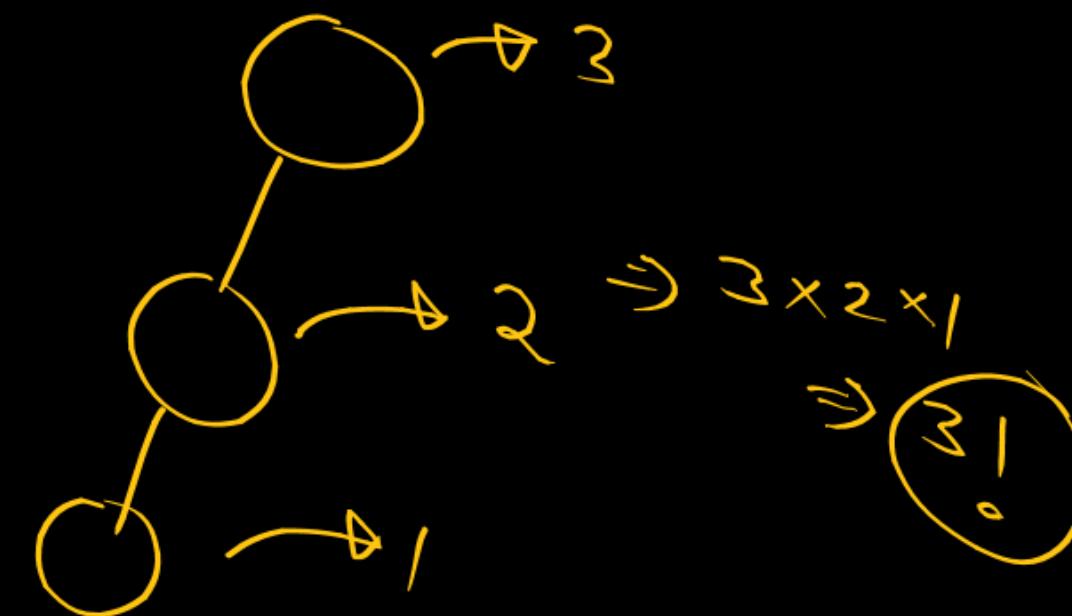
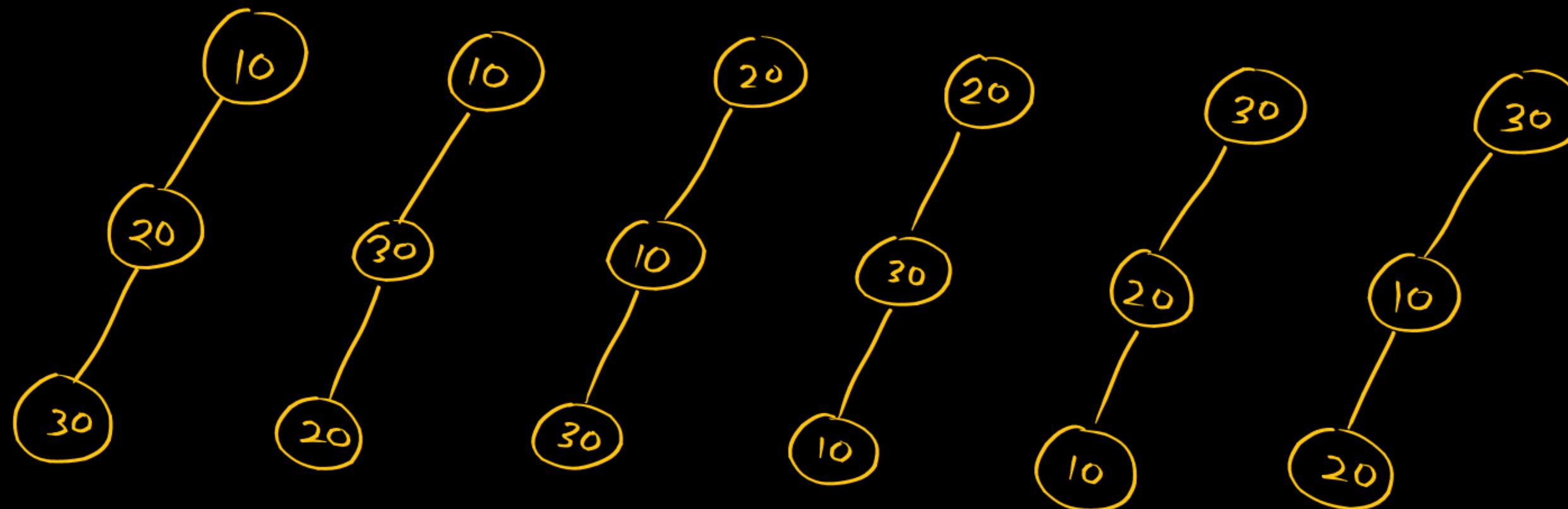
$n=3$ \rightarrow 5 structures are possible



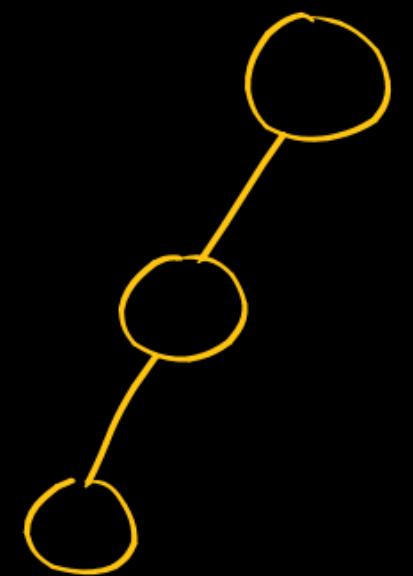
Labelled binary tree with 3 distinct keys.

SI

⑩ 20,30

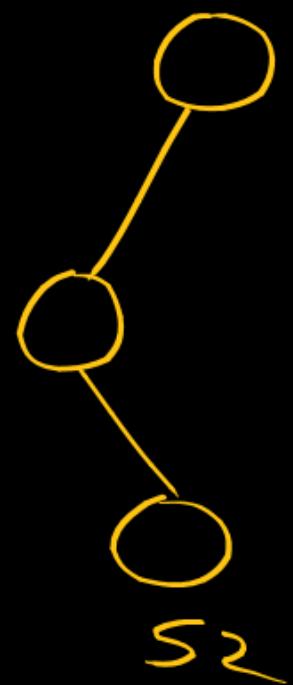


$n=3$



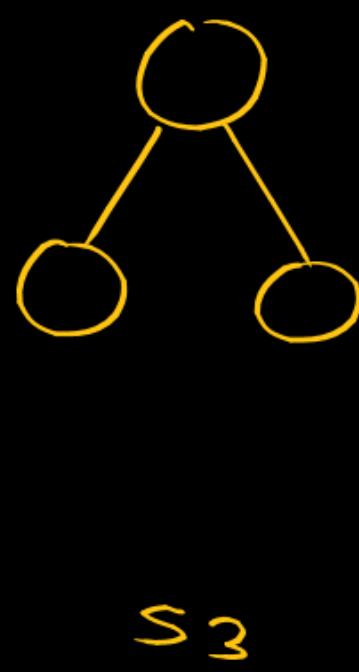
s_1

$$\xrightarrow{3!}$$



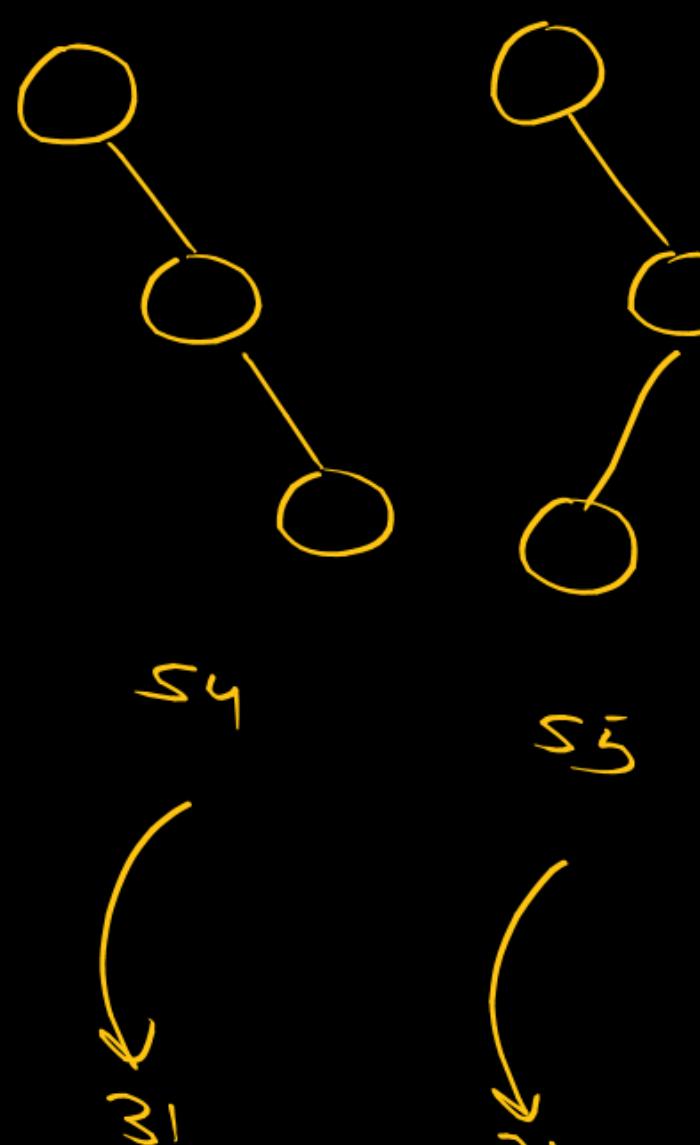
s_2

$$\xrightarrow{3!}$$



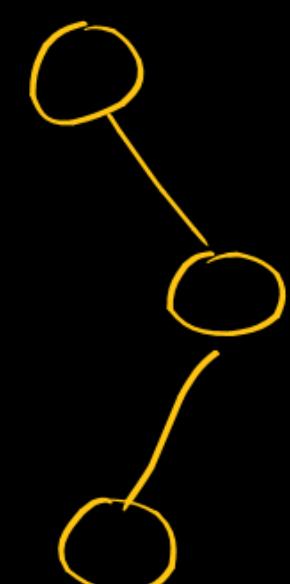
s_3

$$\xrightarrow{3!}$$



s_4

$$\xrightarrow{3!}$$



s_5

$$\xrightarrow{3!}$$

labelled binary trees
with 3 distinct keys

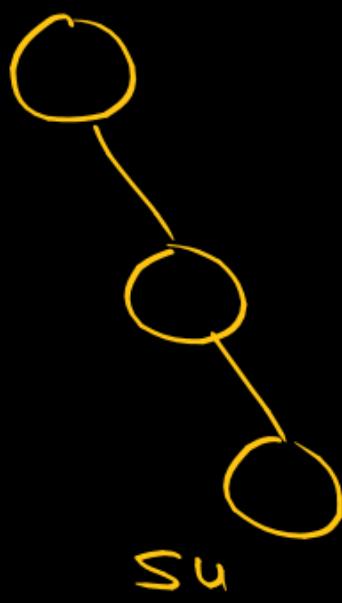
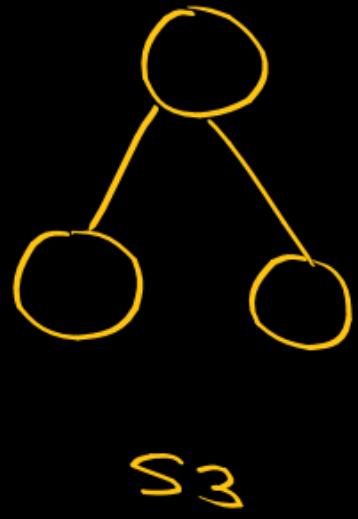
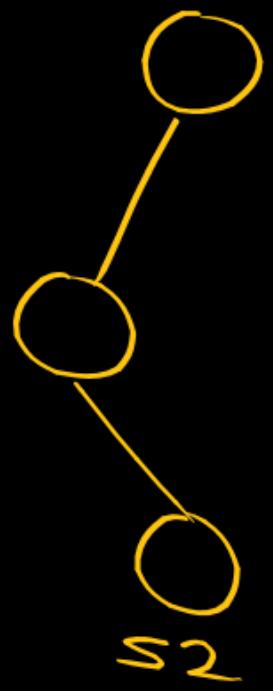
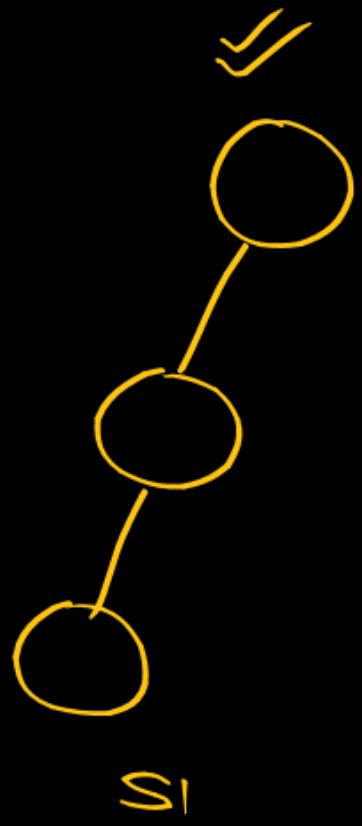
$$= \left(\begin{array}{l} \# \text{ of unlabelled binary} \\ \text{trees with 3 nodes} \end{array} \right)$$

$$\times 3!$$

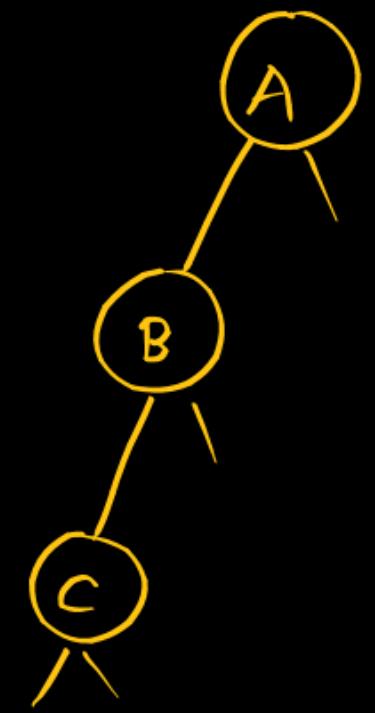
$$= 5 \times 6 \\ = 30$$

$$\# \text{labelled binary trees with } n \text{ distinct keys} = \frac{2^n}{n+1} C_n \times n!$$

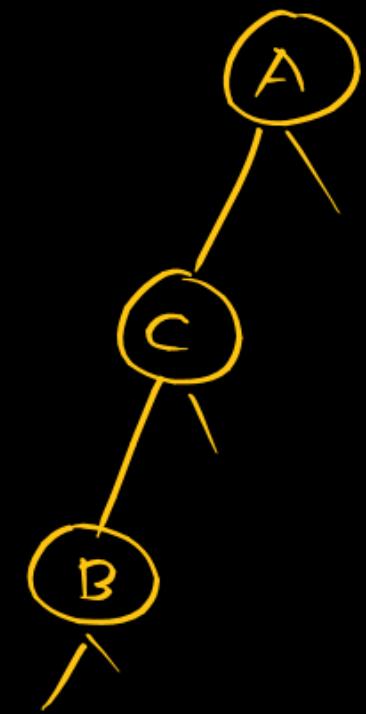
binary trees possible with preorders ABC.



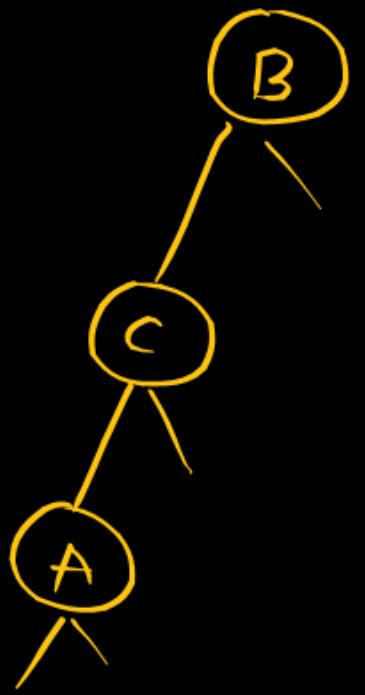
For every structure 1 such binary tree is possible



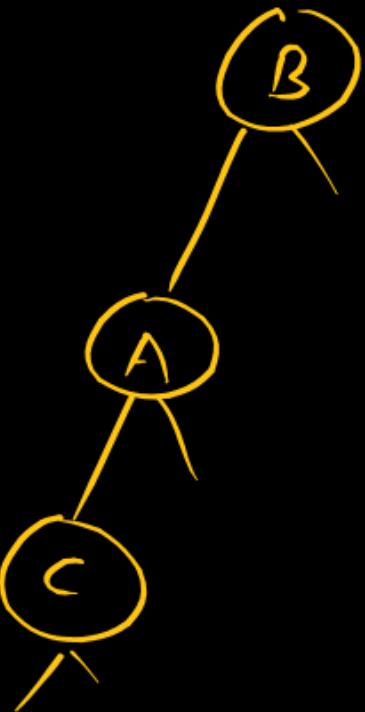
Pre: ABC
✓



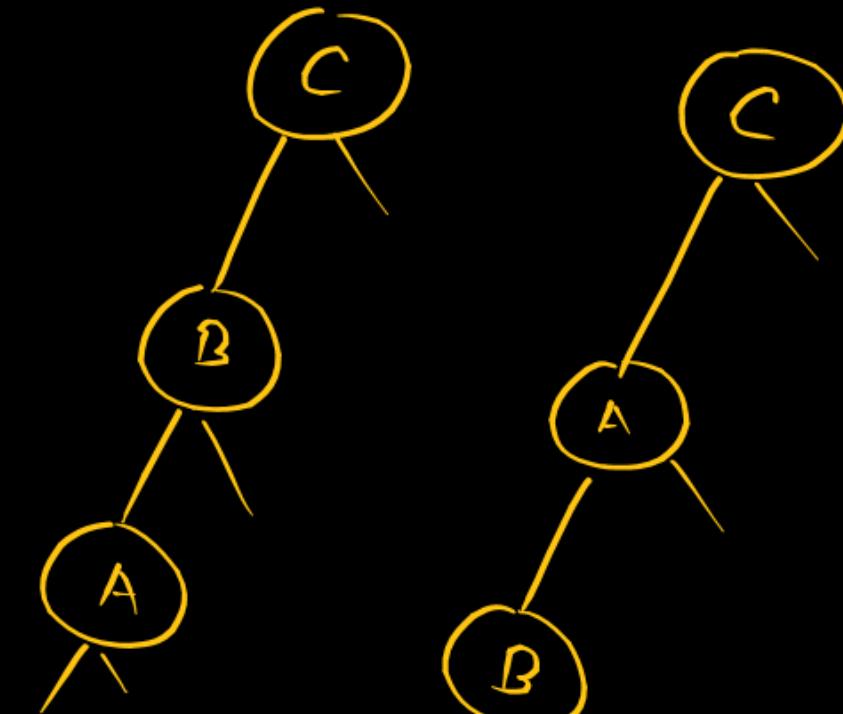
Pre: ACB
✗



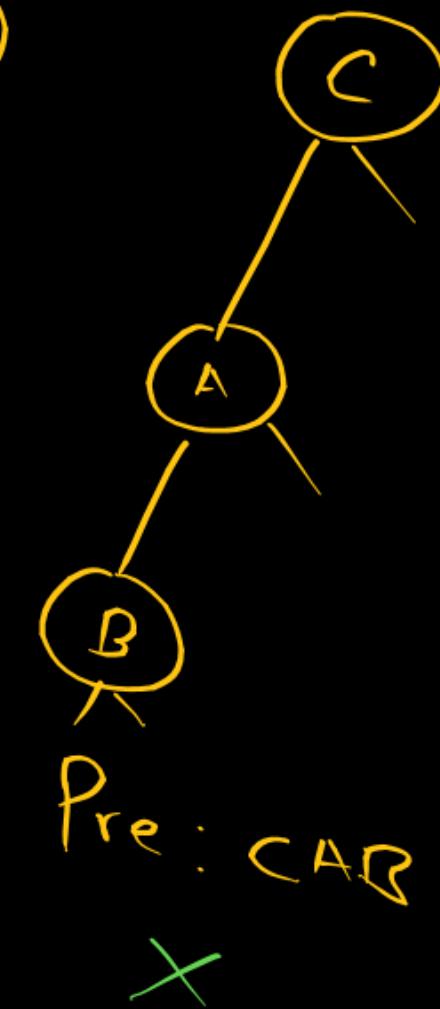
Pre: BCA
✗



Pre: BAC
✗

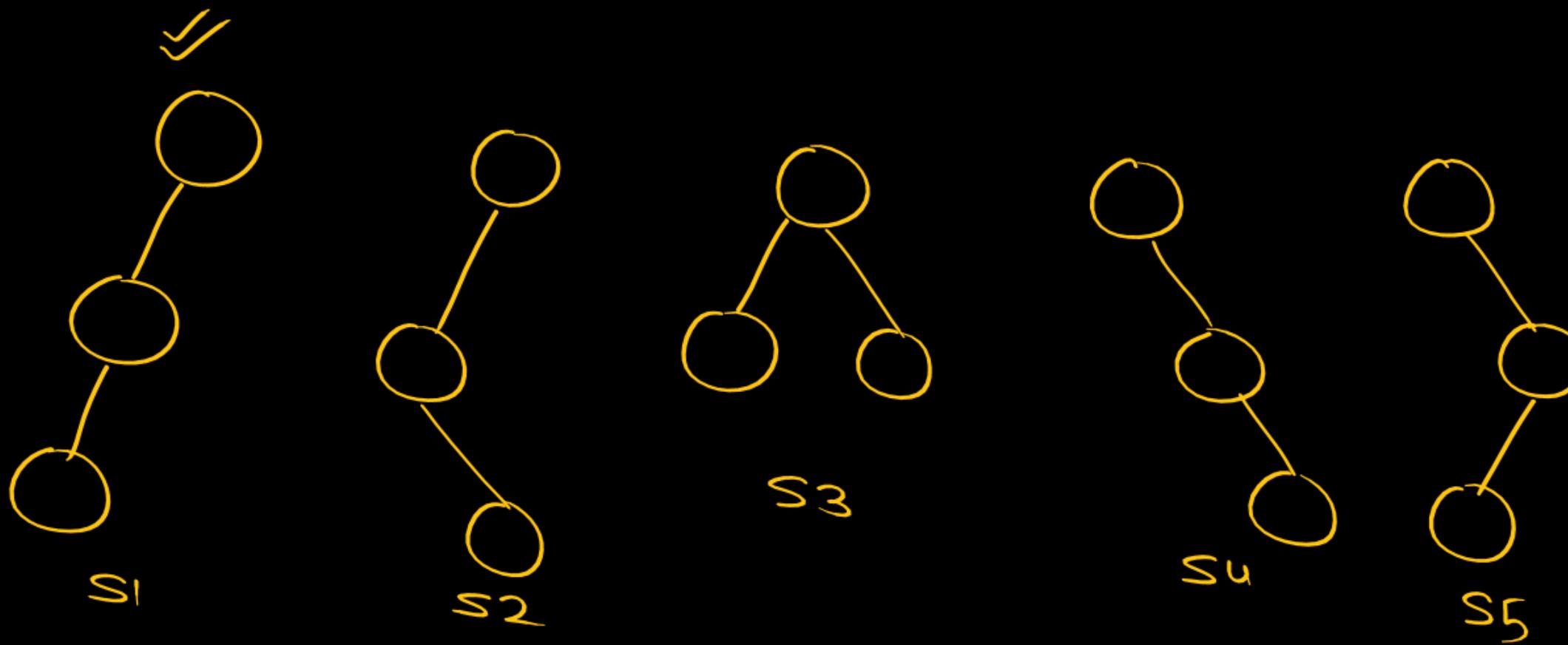


Pre: CBA
✗



Pre: CAB
✗

binary trees possible with preorders ABC.

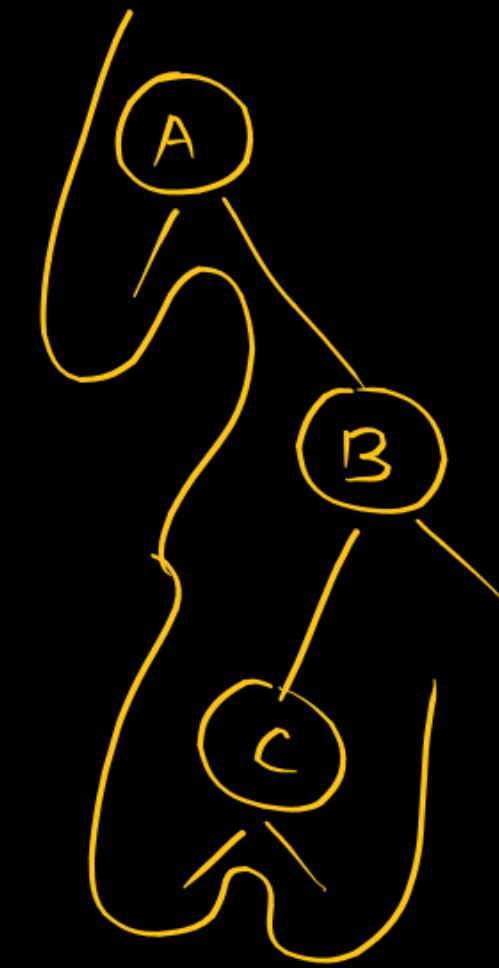
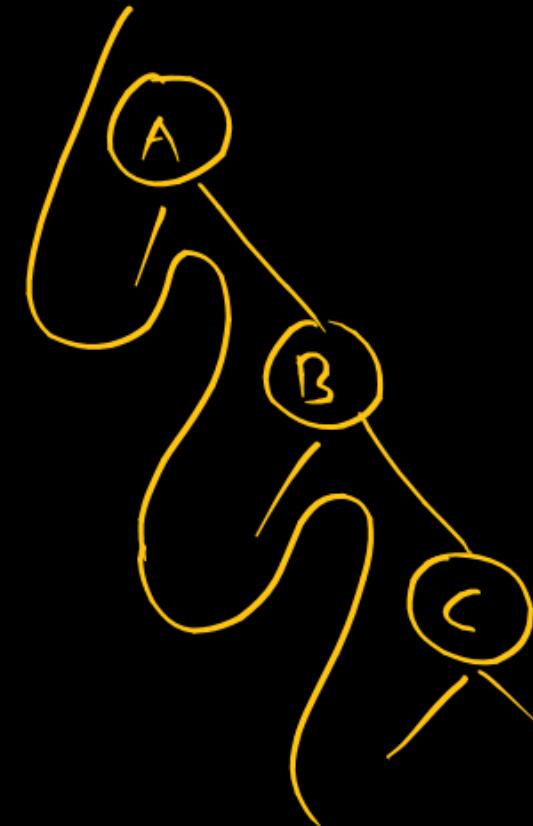
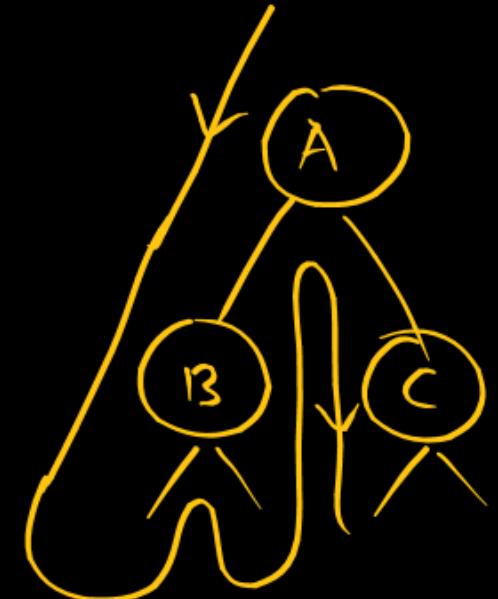
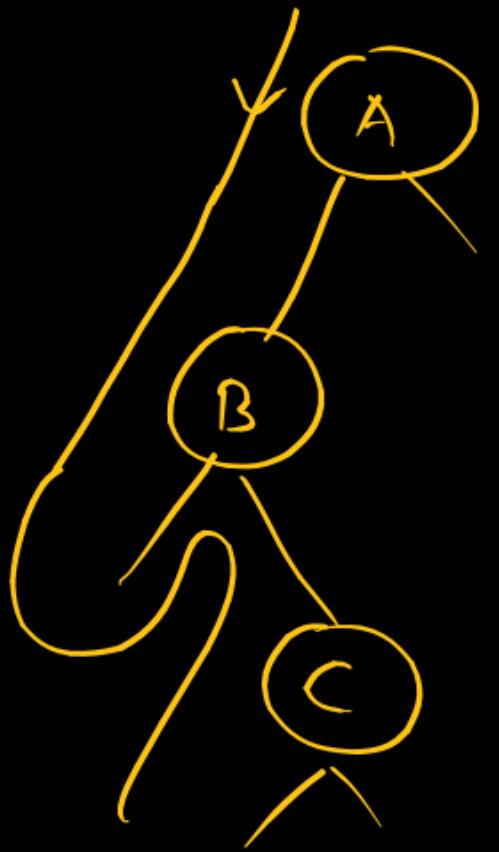
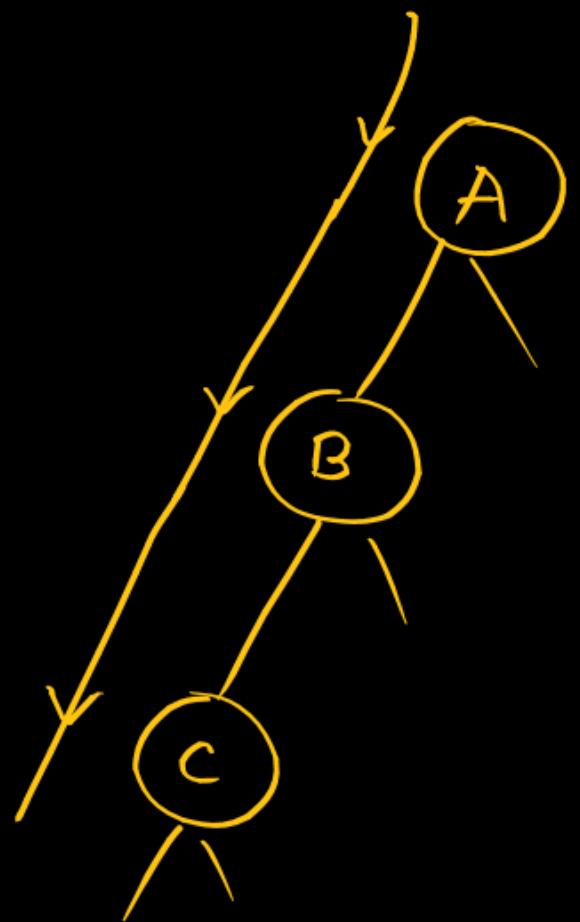


⇒ # binary trees with preorder as ABC = # binary tree structure with 3 nodes

$$\# \text{ binary trees with a given preorder of } n \text{ length} = \frac{2^n C_n}{n+1}$$

Q # binary trees with preorder DABCE = $\frac{10 C_5}{6}$

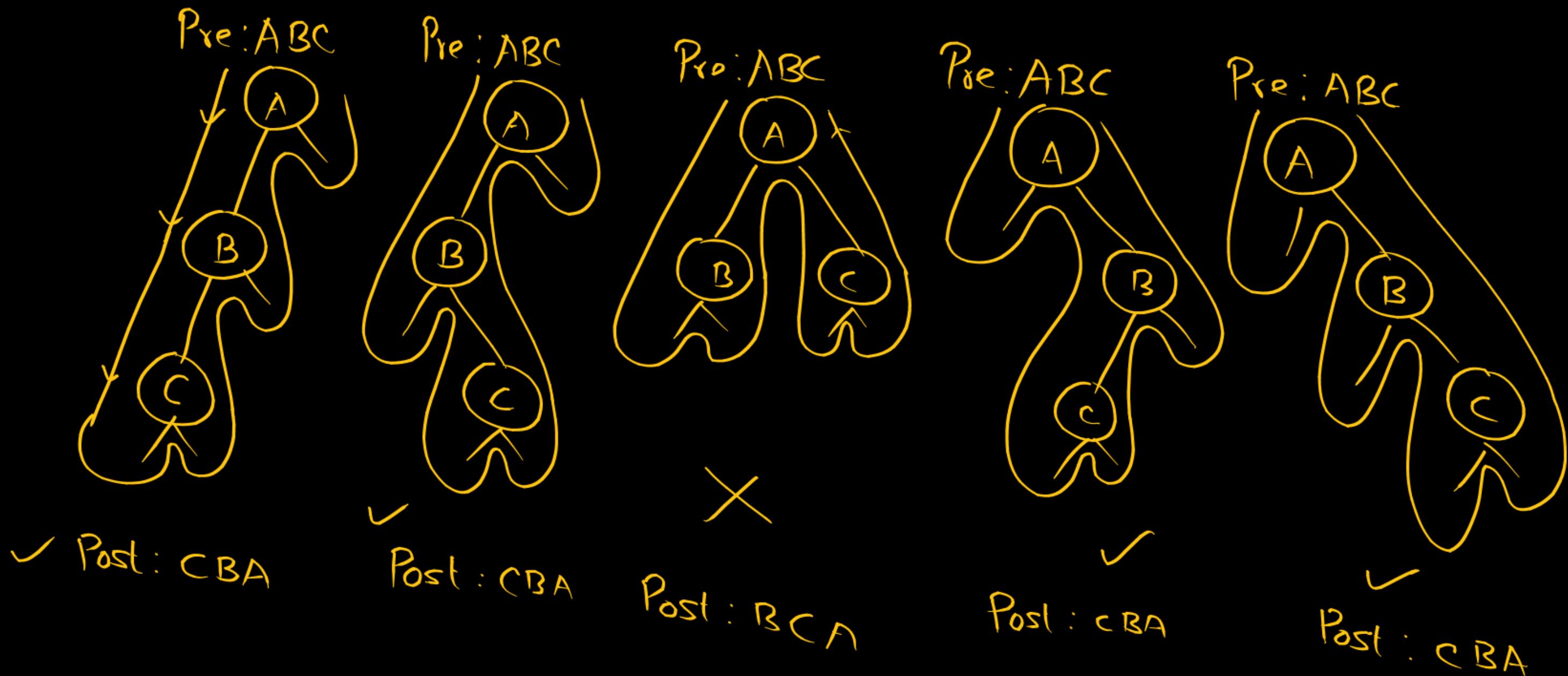
Pre : ABC



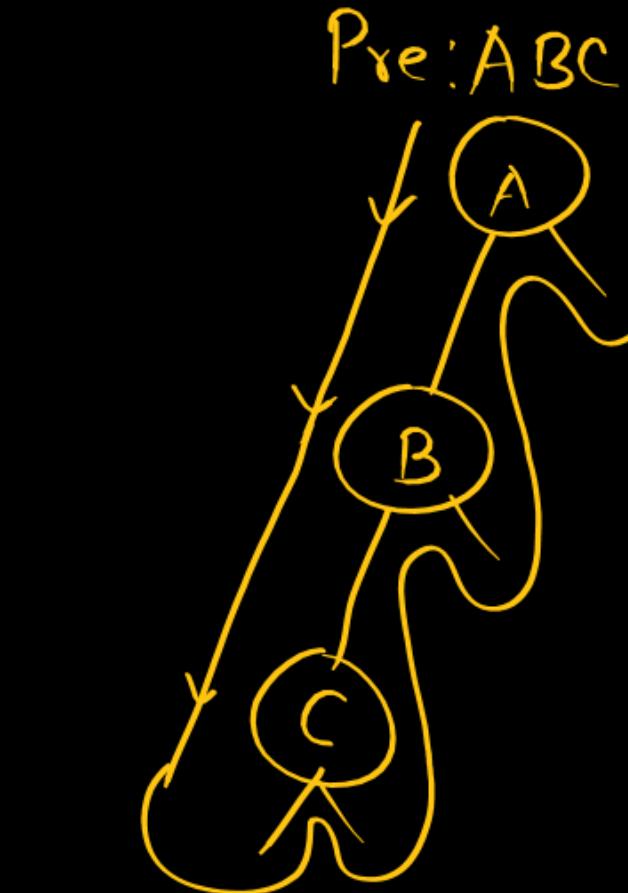
Given Postorder (length n), no. of binary tree = $\frac{2^n C_n}{n+1}$

With any 1 given traversal (Pre/In/Post), the no. of binary tree possible = $\frac{2^n C_n}{n+1}$

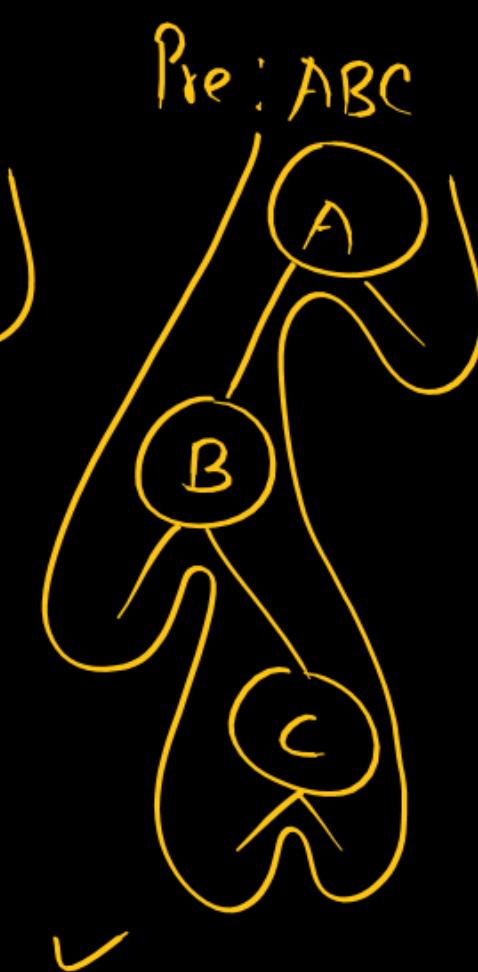
binary trees possible with preorder ABC & postorder CBA



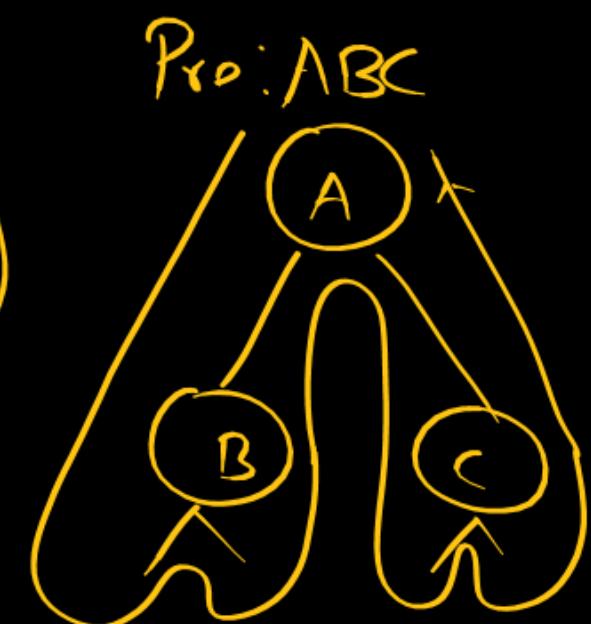
binary trees possible with preorder ABC & postorder CBA



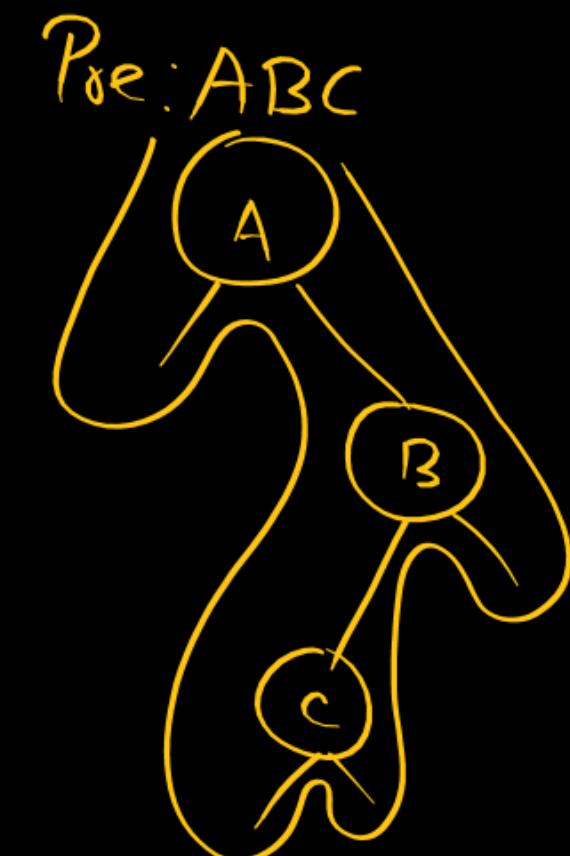
✓ Post:CBA



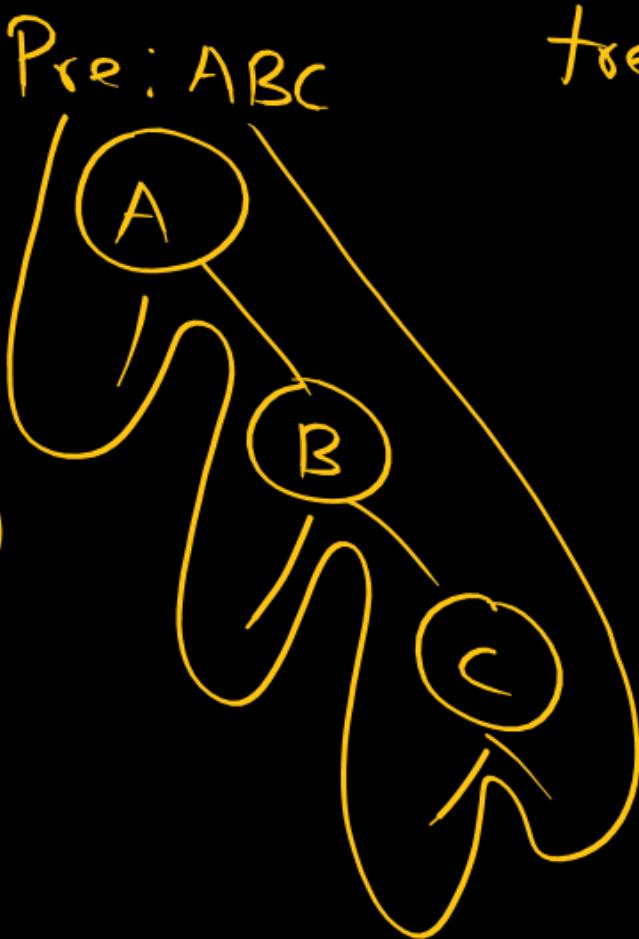
✓ Post:CBA



✗ Post:BCA



✓ Post:CBA



✓ Post:CBA

⇒ Many such
trees

Q

Binary tree

Pre : ABC
Post : BAC

No binary tree

Q)

Randomly ✓

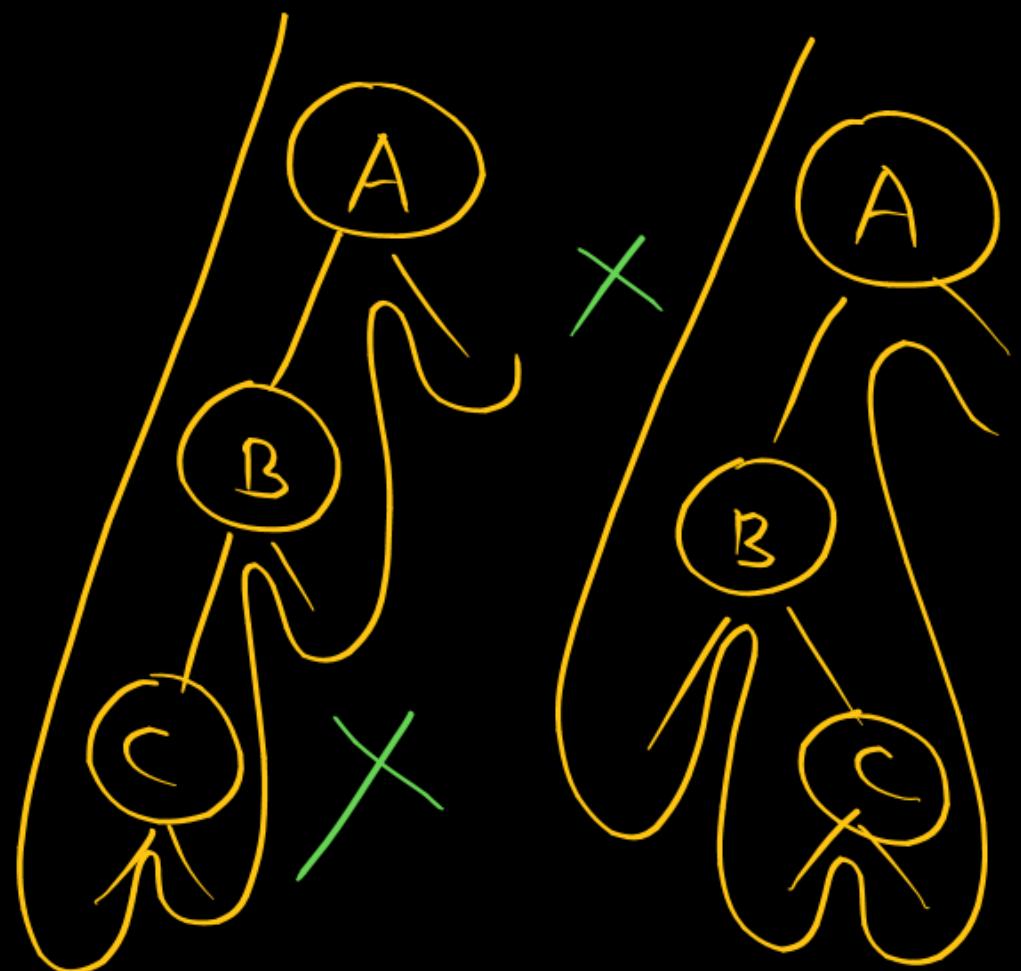
Pre :
Post :

Given } \Rightarrow Many trees

Given } \Rightarrow Zero or more

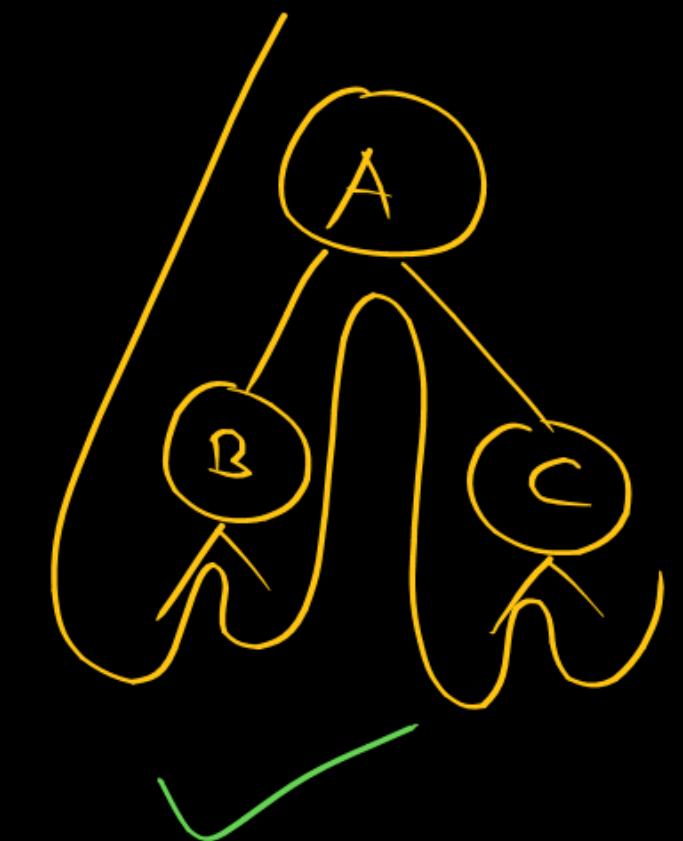
binary trees with Preorder: ABC &
Inorder : BAC

⇒ Unique binary tree

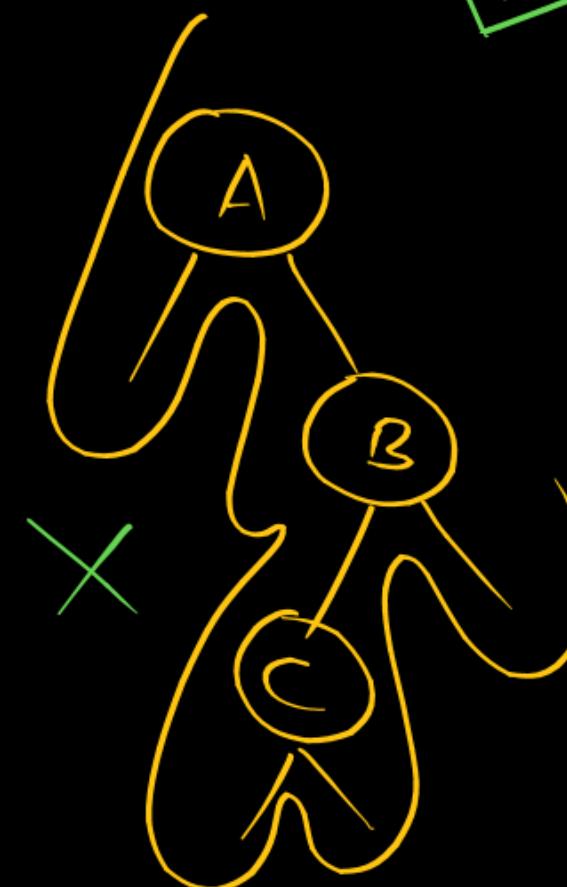


In: CBA

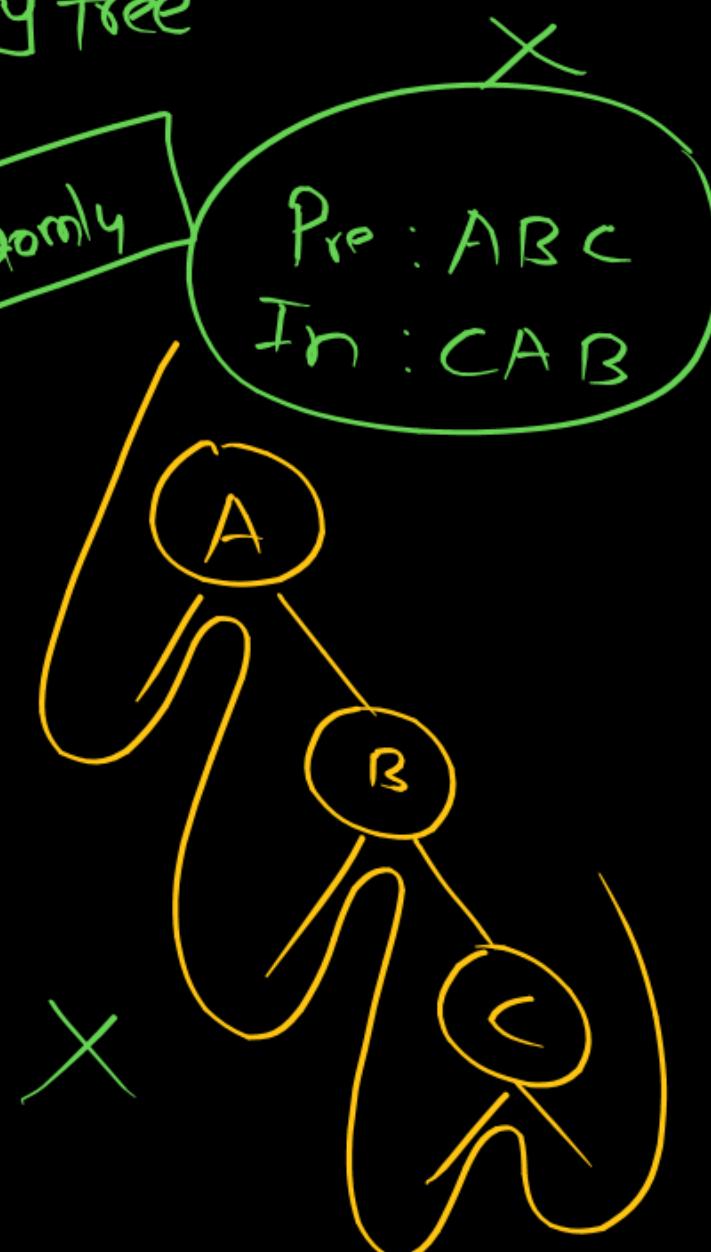
In: BCA



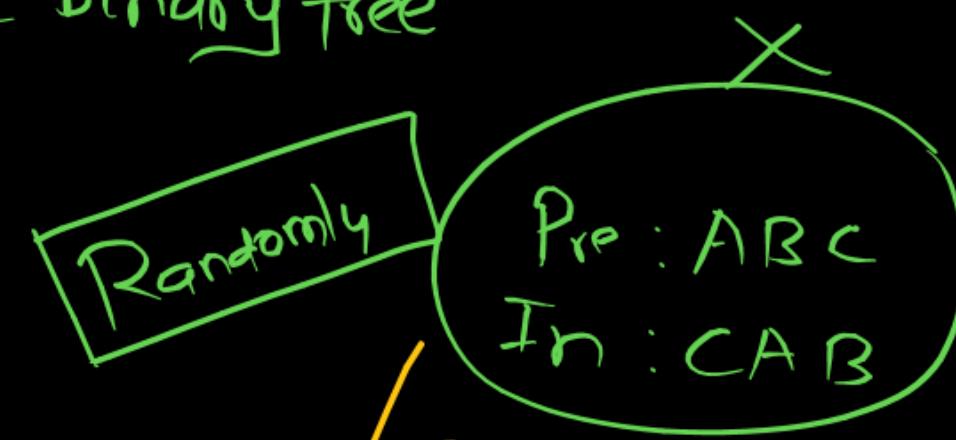
In: BAC



In: ACR



In: ABC

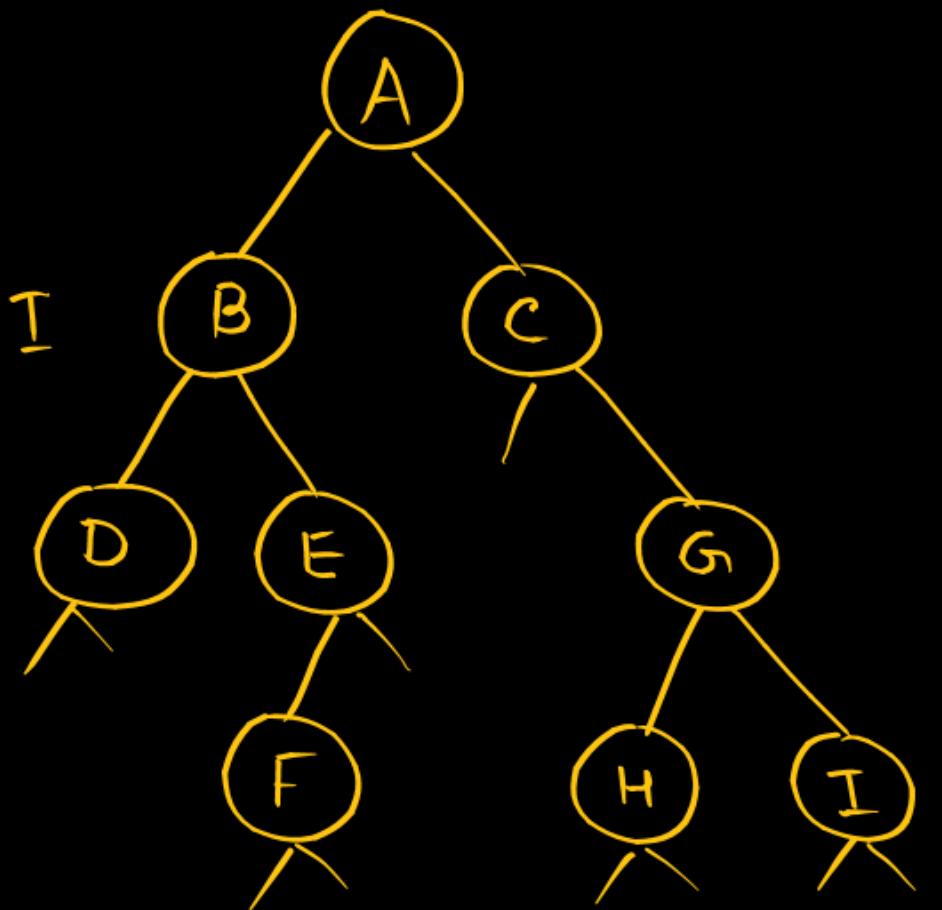


Randomly

Pre : ABC
In : CAB

Pre : A B D E F C G H I

In : D B F E A C H G I



Pre : A B D E F C G H I

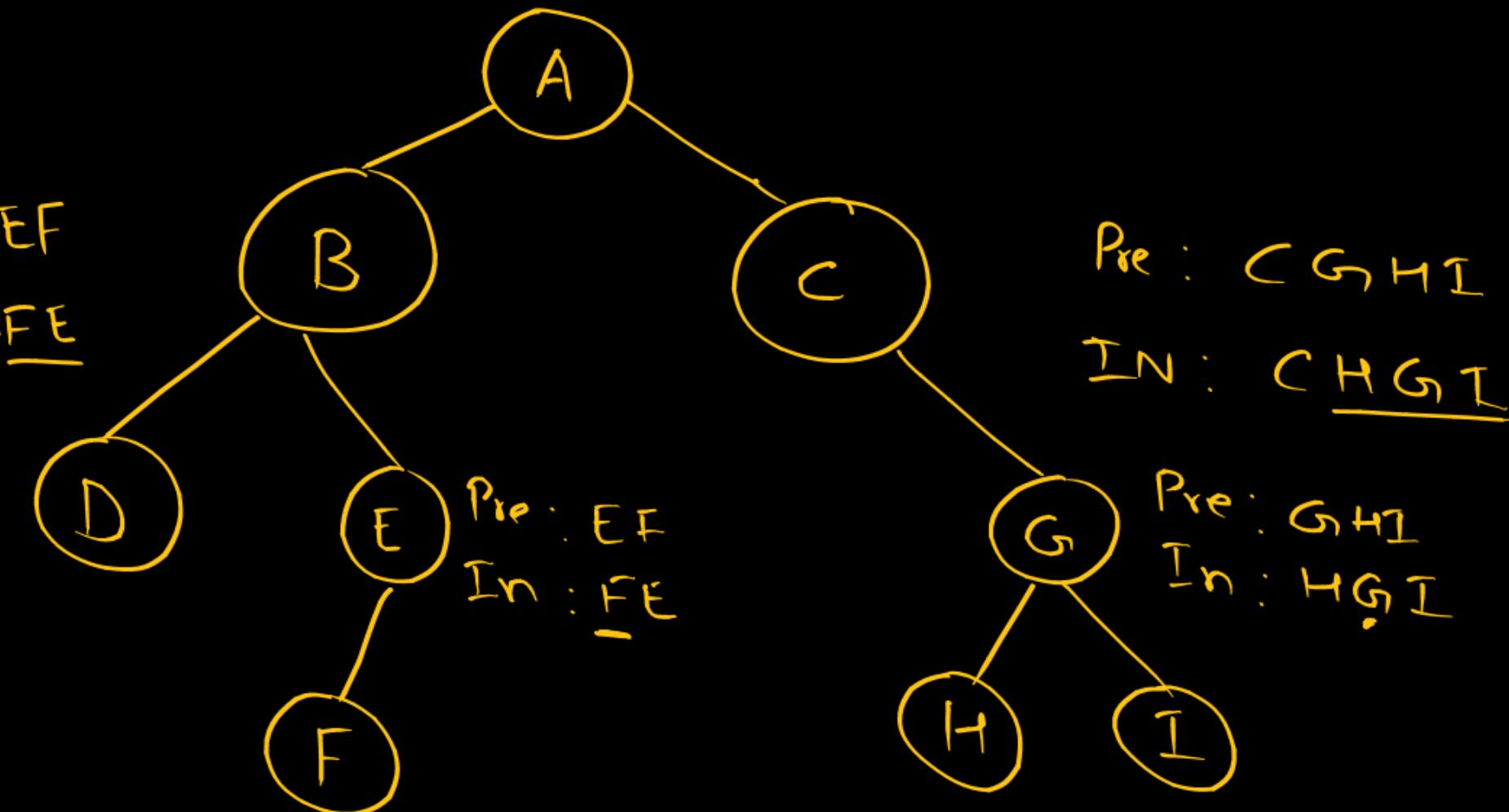
In : DBF E A C H G I
LT RT

Root, LT, RT
LT, Root, RT

Pankaj sir PW

Pre: BDEF
In: DBFE

Pre: D
In: D



Pre: CGHI
IN: CHGI

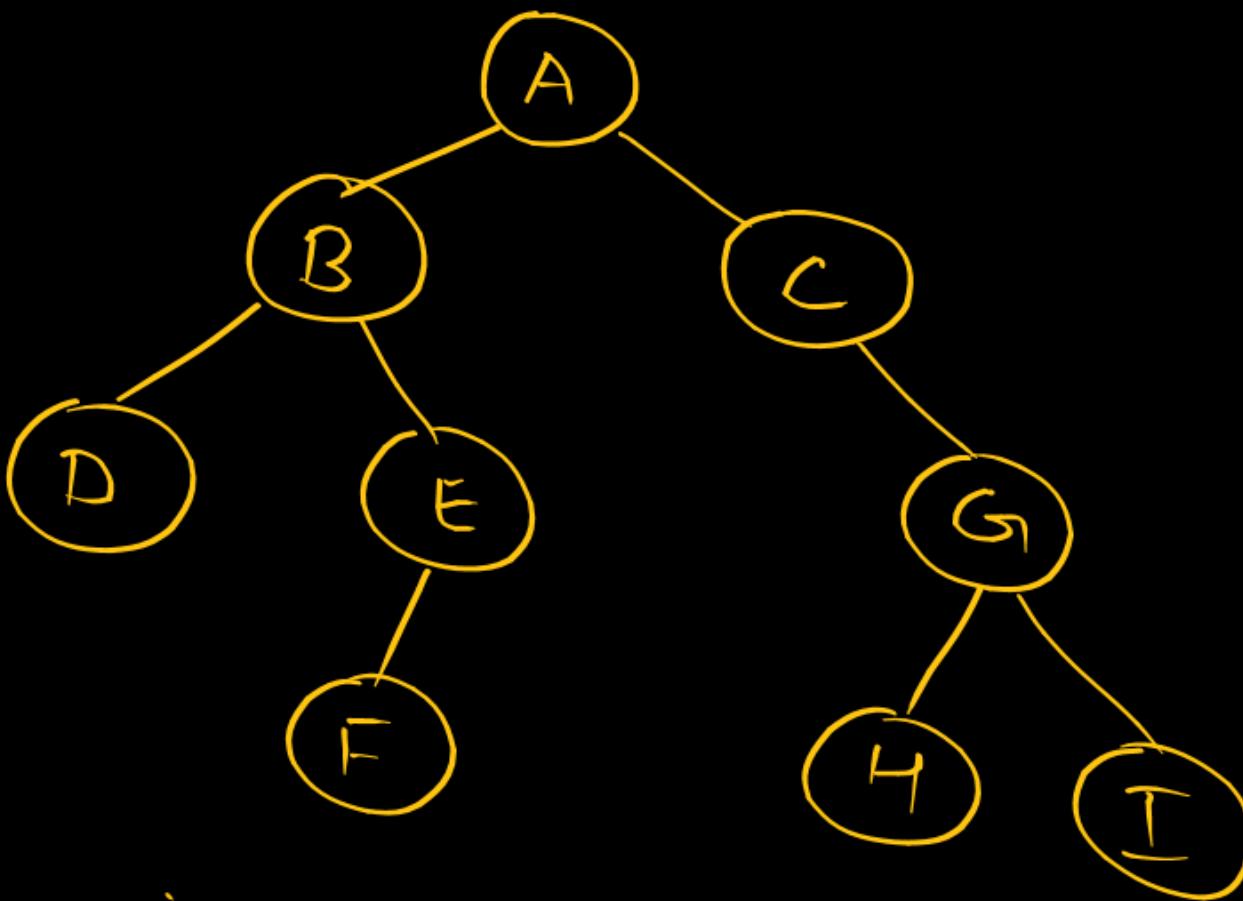
Pre: GHI
In: HGII

Pre : A B D E F C G H I

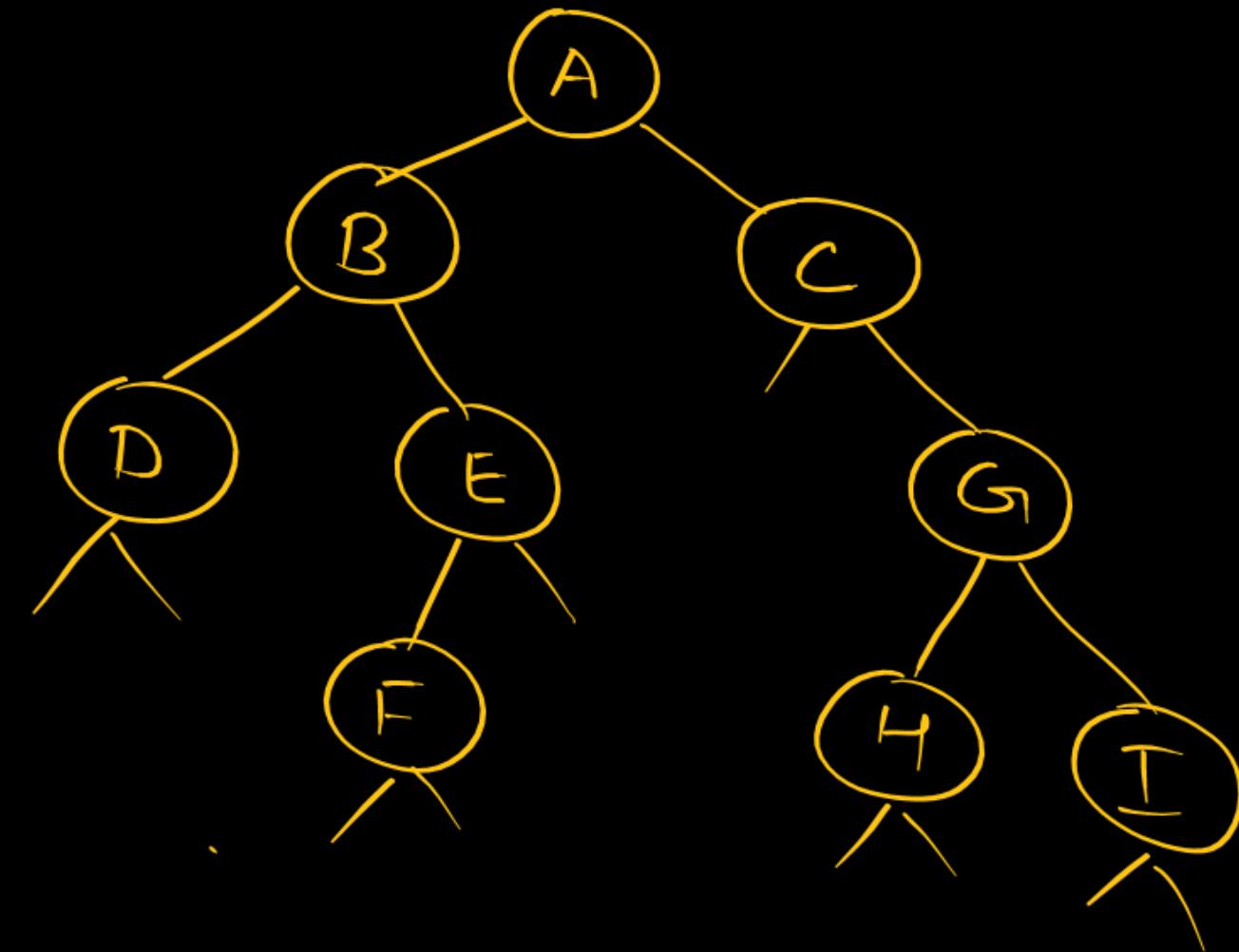
In : DBF E A C H G I

short-trick

time/PWpankajsirP



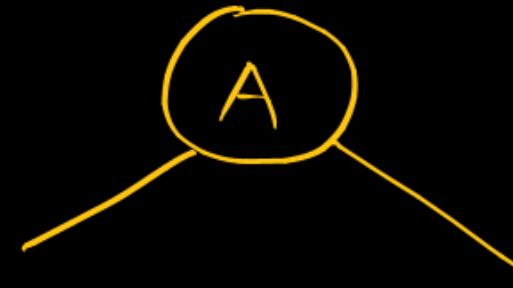
Post : D F E B H I G C A
In : D B F E A C H G I



Post : D F E B H I G C A

In : D B F E A C H G I

LT RT



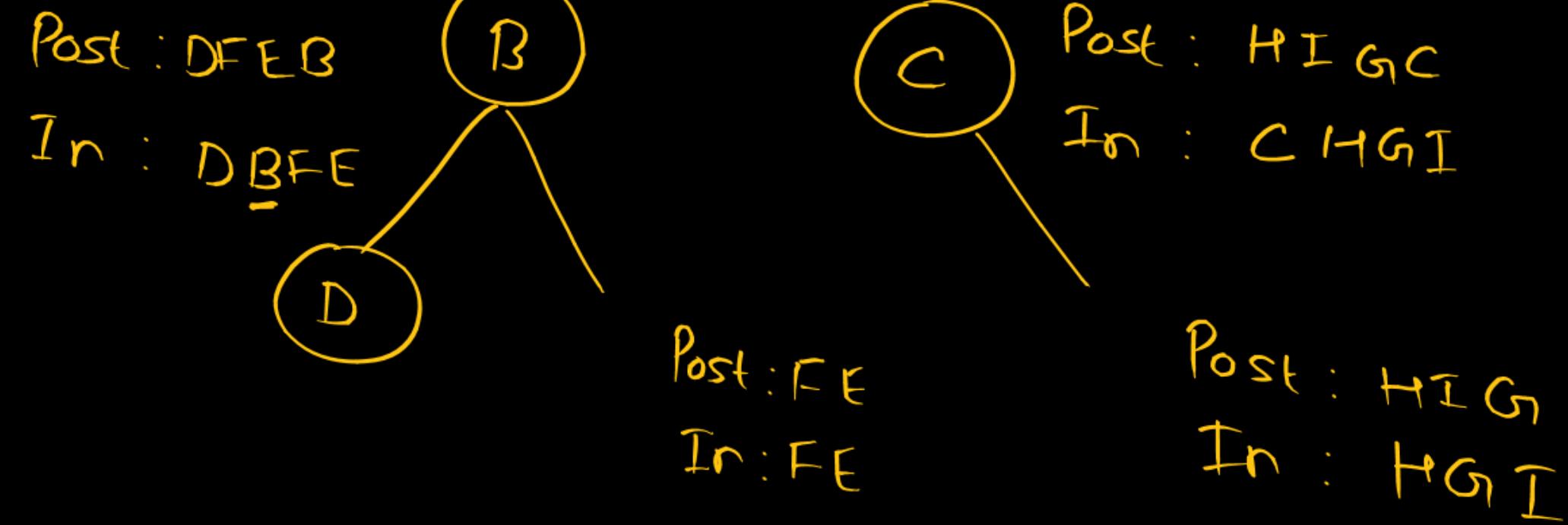
Post : D F E B

In : D B F E

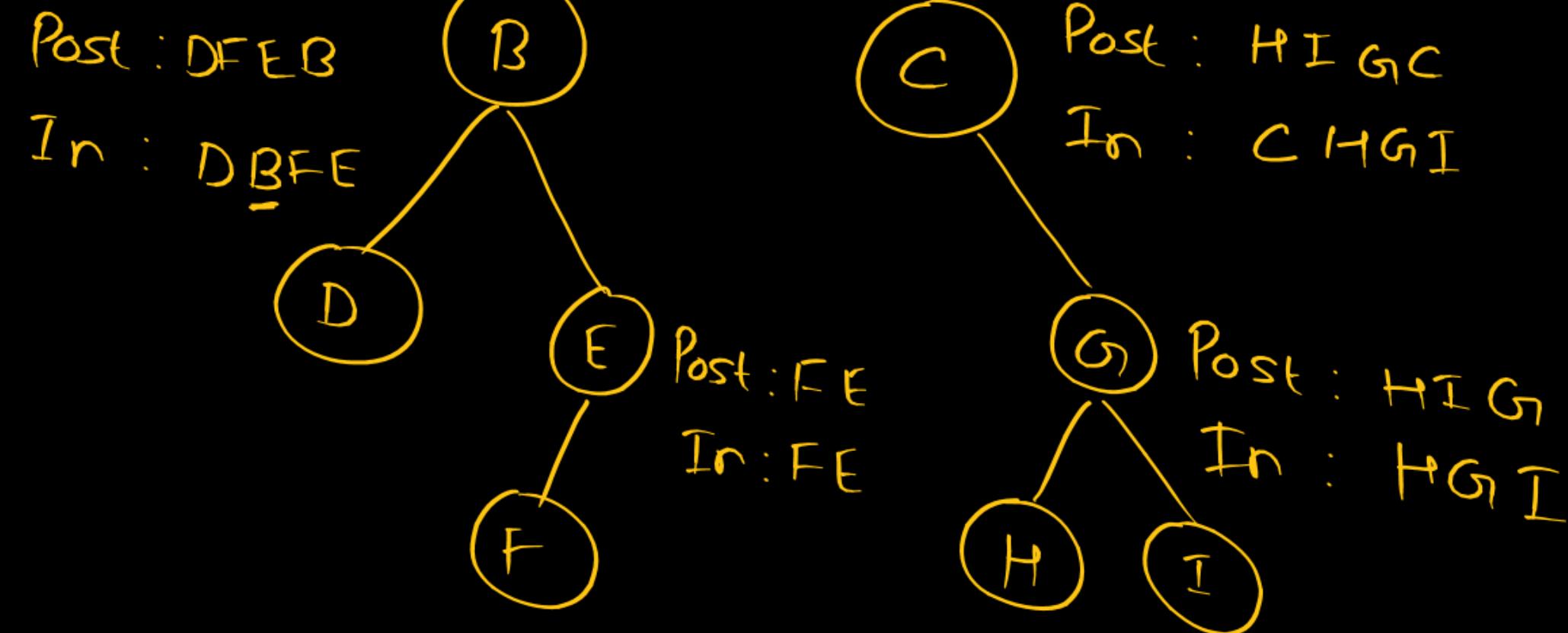
Post : H I G C

In : C H G I

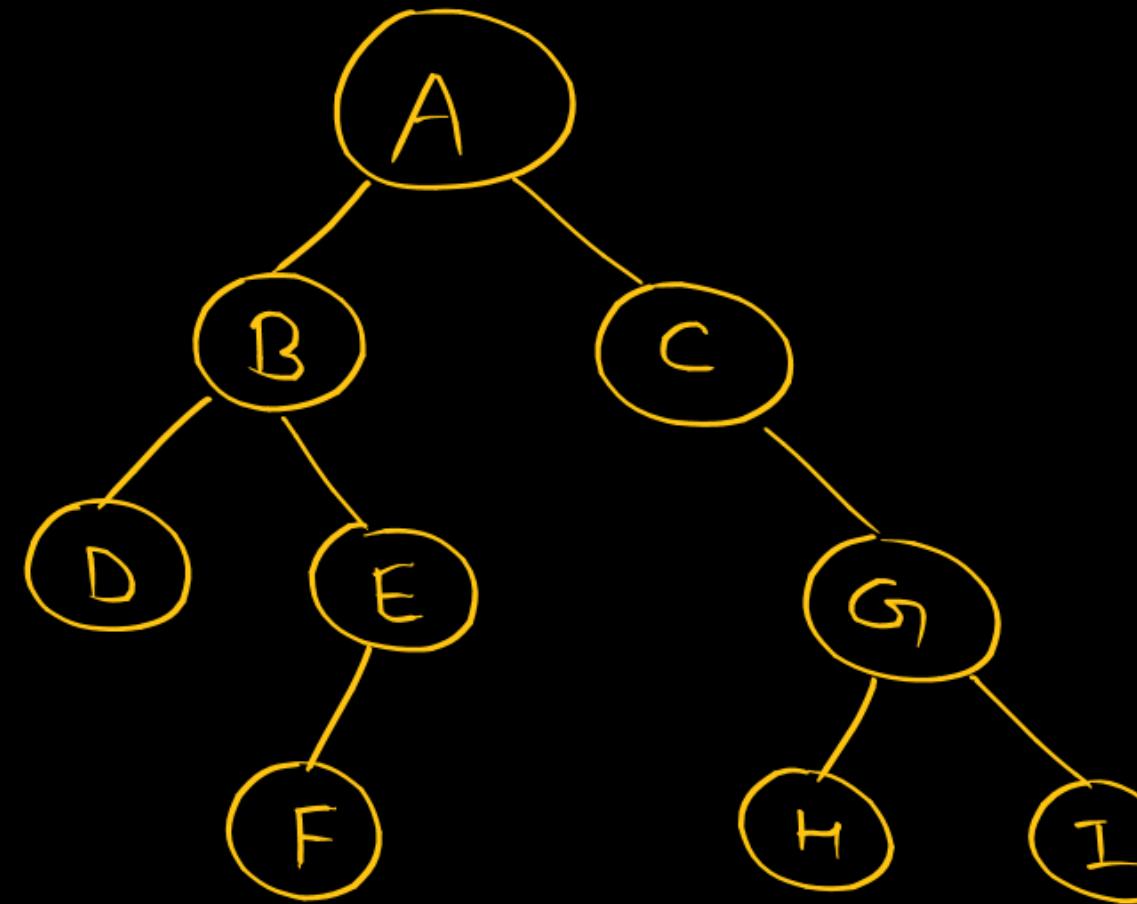
Post : D F E B H I G C A
In : D B F E A C H G I
LT RT



Post : D F E B H I G C A
In : D B F E A C H G I
LT RT



Post : D F E B H I G C A
In : D B F E A C H G I



binary Tree \rightarrow 1 leaf node



Every node has 0 or 1 child

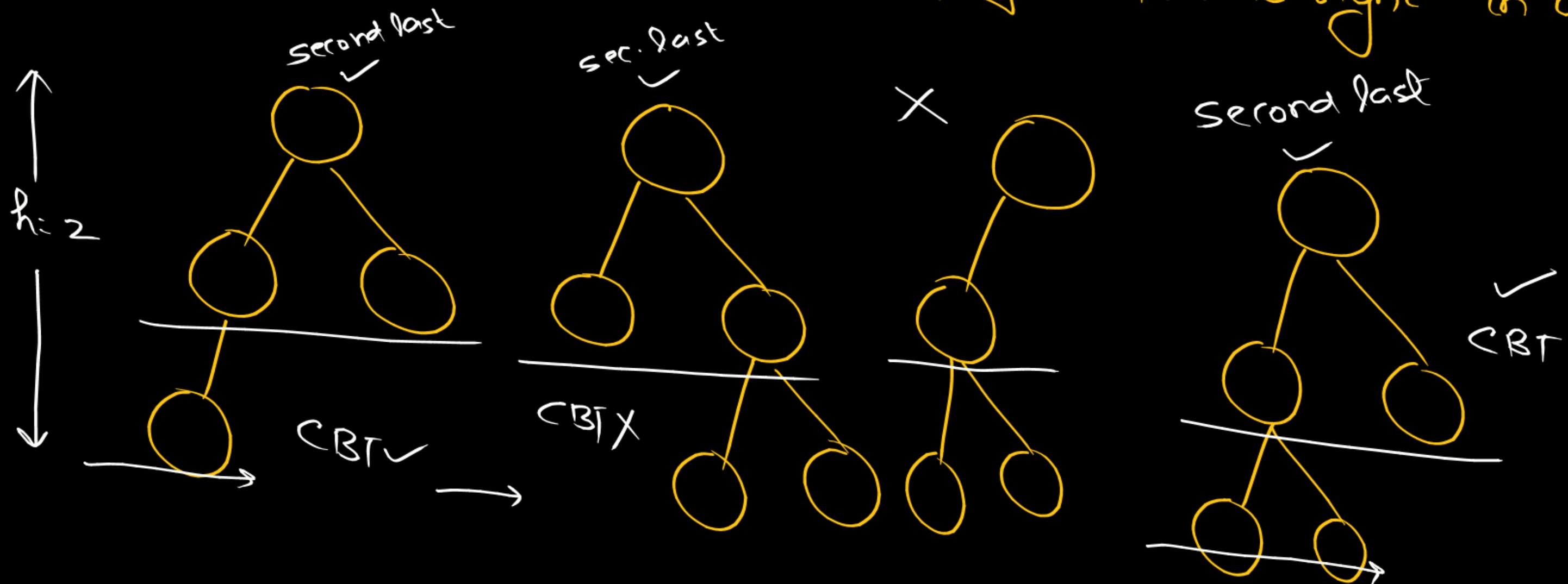
Q # binary tree with $R=3$ & 1 leaf node.

H.W

Q # binary tree with $R=7$ & 1 leaf node

Complete binary Tree

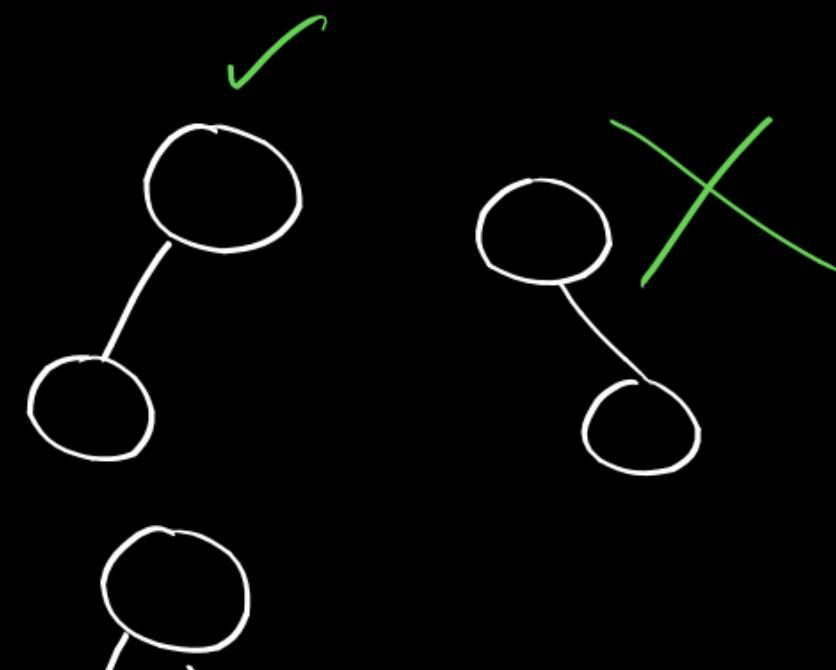
A CBT is a binary tree which is full upto second last level & nodes at last level are filled from left to right in order.



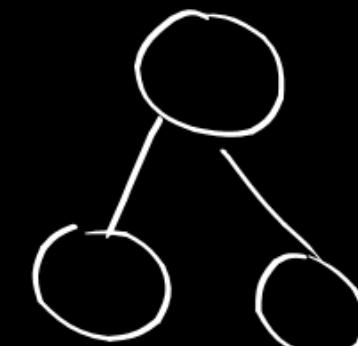
1) Structures of CBT
with 1 node



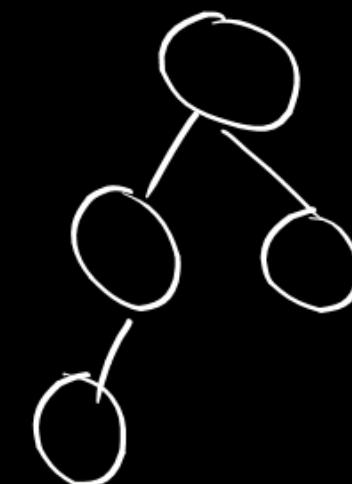
2) Structures of CBT
with 2 nodes



3) Structures of CBT
with 3 nodes



4) Structure of CBT
with 4 nodes



* The structure of a CBT
with k nodes is always
fix

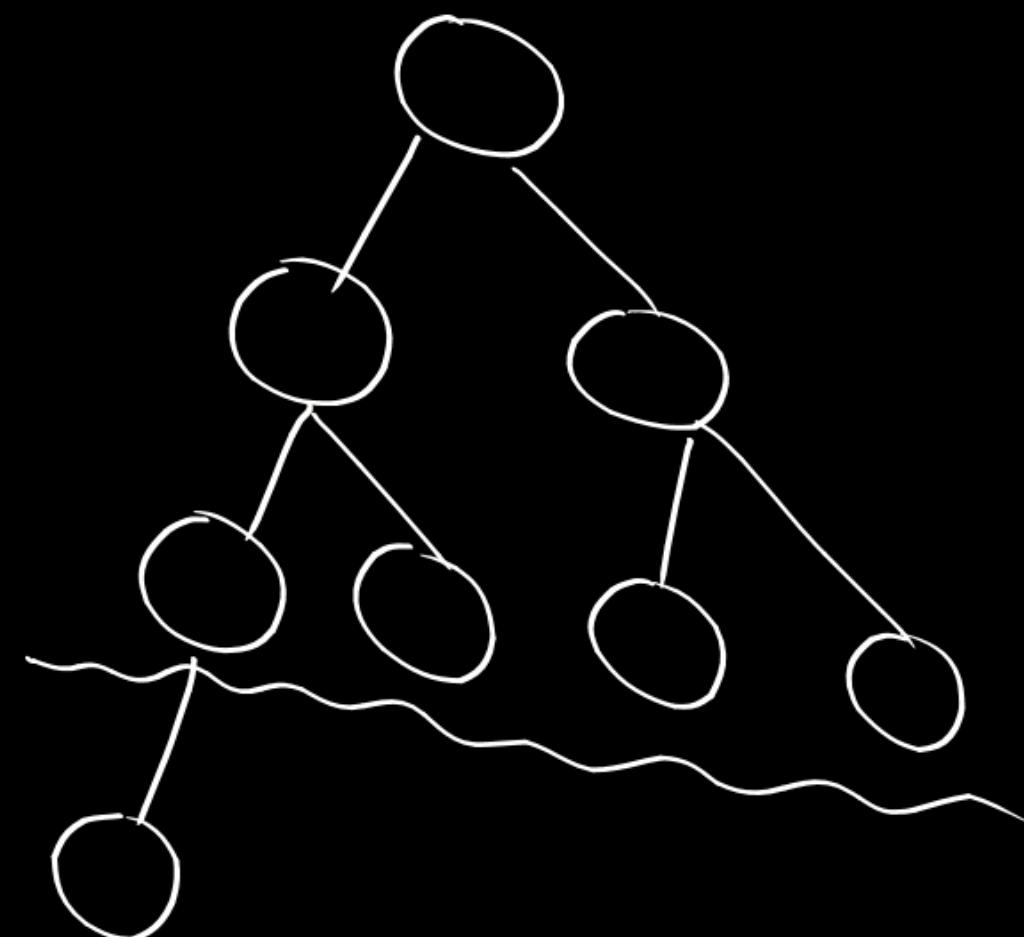
CBT

* Max. no. of nodes in a CBT of h height = $2^{h+1} - 1$

* Min no. of nodes in a CBT of R height

$$\underline{R=3}$$

Why
1 node?



$$2^0$$

$$2^1$$

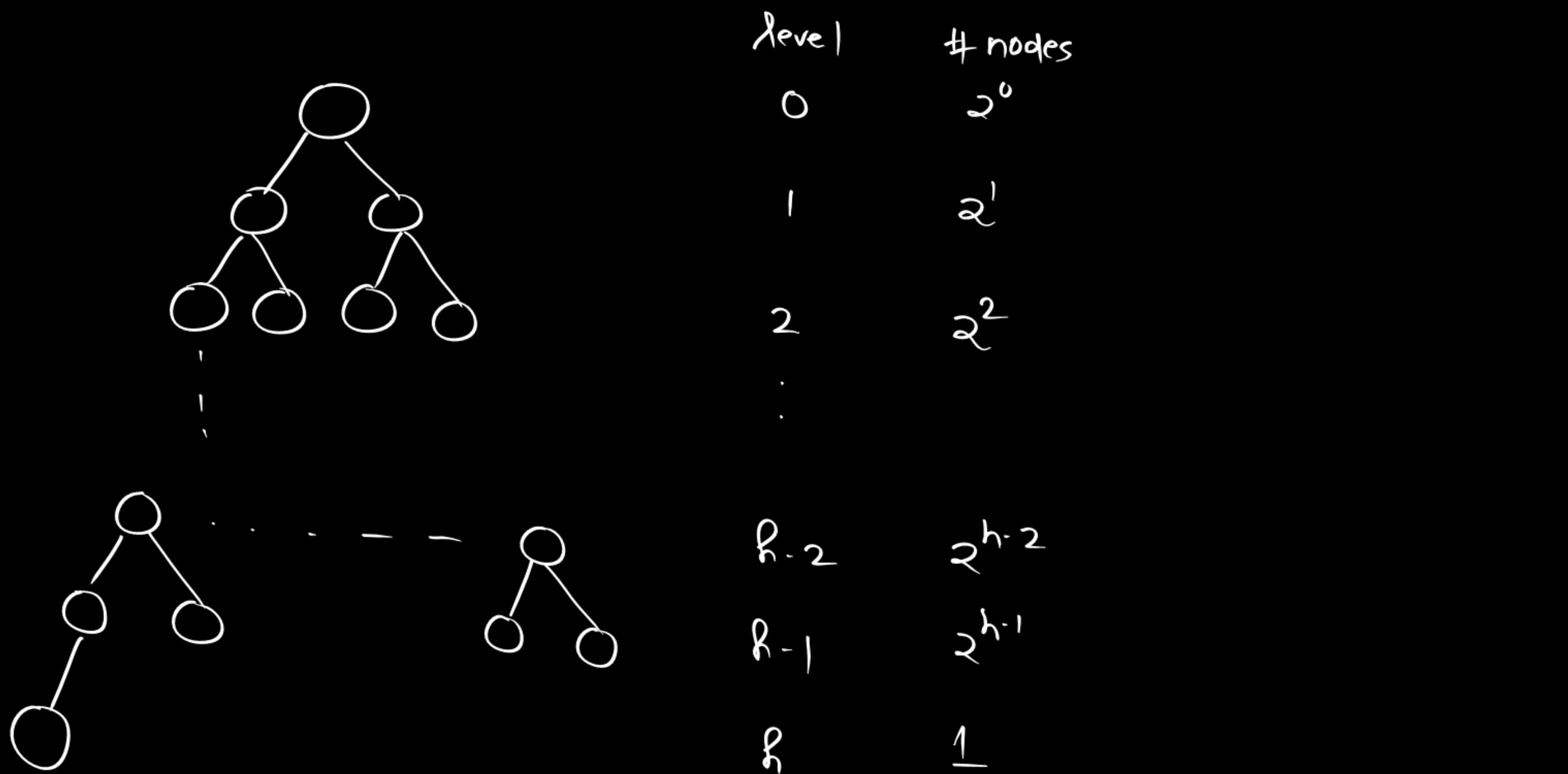
$$2^2$$

level
0

1

2

$$\begin{aligned} N &= (2^0 + 2^1 + 2^2) + 1 \\ &= 2^3 - 1 + 1 \\ &= 2^3 \end{aligned}$$



$$n_{\min} = \left(2^0 + 2^1 + 2^2 + \dots + 2^{h-1} \right) + 1$$

from G.P

$$\left(\frac{2^h - 1}{2 - 1} \right) + 1 = 2^h - 1 + 1 = 2^h$$

$$2^h \leq n \leq 2^{h+1} - 1$$

$n < 10$ & $n \geq 3$

$$2^h \leq n$$

$$\log 2^h \leq \log n$$

$$h \log 2 \leq \log n$$

$$h \leq \frac{\log n}{\log 2}$$

$$h \leq \boxed{\log_2 n}$$

- ①

$$\boxed{\log_2(n+1)-1 \leq h \leq \log_2 n}$$

$$n \leq 2^{h+1} - 1$$

$$(n+1) \leq 2^{h+1}$$

$$\Rightarrow \boxed{3 \leq n \leq 10}$$

$$\log(n+1) \leq (h+1) \log 2$$

$$(h+1) \log 2 \geq \log(n+1)$$

$$(h+1) \geq \frac{\log_2(n+1)}{(h+1)}$$

$$\boxed{h \geq \log_2(n+1)-1} \quad (2)$$

THANK - YOU