

# Category Theory

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## 0 Introduction

I didn't go to the first 3 lectures, so no intro – sorry. I have no idea on what this course is about, let's see

# 1 Definitions and examples

## Definition. (1.1)

A category  $\mathcal{C}$  consists of:

- (a) a collection  $\text{ob } \mathcal{C}$  of *objects*  $A, B, C$ ;
- (b) a collection  $\text{mor } \mathcal{C}$  of *morphisms*  $f, g, h$ ;
- (c) two operations domain, codomain assigning to each  $f \in \text{mor } \mathcal{C}$  a pair of objects, its *domain* and *codomain*; we write  $A \xrightarrow{f} B$  to mean  $f$  is a morphism and  $\text{dom } f = A, \text{cod } f = B$ ;
- (d) an operation assigning to each  $A \in \text{ob } \mathcal{C}$  a morphism  $A \xrightarrow{1_A} A$ ;
- (e) a partial binary operation  $(f, g) \rightarrow fg$  on morphisms, such that  $fg$  is defined iff  $\text{dom } f = \text{cod } g$ , and  $\text{dom}(fg) = \text{dom } g$ ,  $\text{cod}(fg) = \text{cod}(f)$  if  $fg$  is defined, satisfying:
- (f)  $f1_A = f = 1_B f$  for any  $A \xrightarrow{f} B$ ;
- (g)  $(fg)h = f(gh)$  whenever  $fg$  and  $gh$  are defined.

## Remark. (1.2)

- (a) This definition is independent of any model of set theory. If we're given a particular model of set theory, we call  $\mathcal{C}$  *small* if  $\text{ob } \mathcal{C}$  and  $\text{mor } \mathcal{C}$  are sets.
- (b) Some texts say  $fg$  means  $f$  followed by  $g$ , i.e.  $fg$  is defined iff  $\text{cod } f = \text{dom } g$ .
- (c) Note that a morphism  $f$  is an identity iff  $fg = g$  and  $hf = h$  whenever the composites are defined. So we could formulate the definition entirely in terms of morphisms.

## Example. (1.3)

- (a) The category **Set** has all sets as objects, and all functions between sets as morphisms.

Strictly speaking, morphisms  $A \rightarrow B$  are pairs  $(f, B)$  where  $f$  is a set-theoretic function. (See part II logic and sets)

- (b) The category **Gp** has all groups as objects, group homomorphisms as morphisms.

Similarly, **Ring** is the category of rings, **Mod<sub>R</sub>** is the category of  $R$ -modules.

- (c) The category **Top** has all topological spaces as objects, and continuous functions as morphisms.

Similarly, **Unif** has all uniform spaces and uniformly continuous functions as morphisms, **Mf** has all manifolds and smooth maps correspondingly.

- (d) The category **Htpy** has the same objects as **Top**, but morphisms are homotopy classes of continuous functions. More generally, given  $\mathcal{C}$ , we call an equivalence relation  $\simeq$  on  $\text{mor } \mathcal{C}$  a *congruence* if  $f \simeq g \implies \text{dom } f = \text{dom } g$  and  $\text{cod } f = \text{cod } g$ , and  $f \simeq g \implies fh \simeq gh$  and  $kf \simeq kg$  whenever the composites are defined. Then we have a category  $\mathcal{C}/\simeq$  with the same objects as  $\mathcal{C}$ , but congruence classes as morphisms instead.

- (e) Given  $\mathcal{C}$ , the *opposite category*  $\mathcal{C}^{op}$  has the same objects and morphisms as  $\mathcal{C}$ , but  $\text{dom}$  and  $\text{cod}$  are interchanged, and  $fg$  in  $\mathcal{C}^{op}$  is  $gf$  in  $\mathcal{C}$ .

This leads to the *duality principle*: if  $P$  is a true statement about categories, so is the statement  $P^*$  obtained from  $P$  by reversing all arrows.

- (f) A small category with one object is a *monoid*, i.e. a semigroup with 1. In particular, a group is a small cat ( $\boxtimes$ ) with one object in which every morphism is an isomorphism (i.e. for all  $f, \exists g$  s.t.  $fg$  and  $gf$  are identities).

(g) A *groupoid* is a category in which every morphism is an isomorphism. For example, for a topological space  $X$ , the *fundamental groupoid*  $\pi(x)$  has all points of  $X$  as objects, and morphisms  $x \rightarrow y$  are homotopy classes  $rel\{0, 1\}$  of paths  $u : [0, 1] \rightarrow X$  with  $u(0) = x$ ,  $u(1) = y$  (if you know how to prove that the fundamental group is a group, you can prove that  $\pi(x)$  is a groupoid).

(h) A *discrete* cat is one whose only morphism are identities.

A *preorder* is a cat  $\mathcal{C}$  in which, for any pair  $(A, B)$ ,  $\exists$  at most 1 morphism  $A \rightarrow B$ .

A small preorder is a set equipped with a binary relation which is reflexive and transitive.

In particular, a partially ordered set is a small preorder in which the only isomorphisms are identities.

(i) The category **Rel** has the same objects as *set*, but morphisms  $A \rightarrow B$  are arbitrary relations  $R \subseteq A \times B$ . Given  $R$  and  $S \subseteq B \times C$ , we define  $S \cdot R = \{(a, c) \in A \times C \mid (\exists b \in B)((a, b) \in R, (b, c) \in S)\}$ .

The identity  $1_A : A \rightarrow A$  is  $\{(a, a) \mid a \in A\}$ .

Similarly, the category **Part** are for sets and partial functions (i.e. relations s.t.  $(a, b) \in R$  and  $(a, b') \in R \implies b = b'$ ).

(j) Let  $K$  be a field. The category **Mat<sub>K</sub>** has natural numbers as objects, and morphism  $n \rightarrow p$  are  $(p \times n)$  matrices with entries from  $K$ . Composition is matrix multiplication.

(k) We write **Cat** for the category whose objects are all small categories, and whose morphisms are functors between them. (see below for definition of functors)

#### Definition. (1.4)

Let  $\mathcal{C}$  and  $\mathcal{D}$  be categories. A *functor*  $F : \mathcal{C} \rightarrow \mathcal{D}$  consists of:

(a) a mapping  $A \rightarrow FA$  from  $\text{ob } \mathcal{C}$  to  $\text{ob } \mathcal{D}$ ;

(b) a mapping  $f \rightarrow Ff$  from  $\text{mor } \mathcal{C}$  to  $\text{mor } \mathcal{D}$ ,

such that  $\text{dom}(Ff) = F(\text{dom } f)$ ,  $\text{cod}(Ff) = F(\text{cod } f)$ ,  $1_{FA} = F(1_A)$ , and  $(Ff)(Fg) = F(fg)$  whenever  $fg$  is defined.

#### Example. (1.5)

(a) We have *forgetful functors*  $U : \mathbf{Gp} \rightarrow \mathbf{Set}, \mathbf{Ring} \rightarrow \mathbf{Set}, \mathbf{Top} \rightarrow \mathbf{Set}, \mathbf{Ring} \rightarrow \mathbf{AbGp}$  (forget  $\times$ ),  $\mathbf{Ring} \rightarrow \mathbf{Mon}$  (Category of all monoids) (forget  $+$ ).

(b) Given a set  $A$ , the free group  $FA$  has the property:

Given any group  $G$  and any function  $A \xrightarrow{f} UG$  (?), there's a unique homomorphism  $FA \xrightarrow{\tilde{f}} G$  extending  $f$ . Here  $F$  is a functor  $\mathbf{Set} \rightarrow \mathbf{Gp}$ : given  $A \xrightarrow{f} B$ , we define  $Ff$  to be the unique homomorphism extending  $A \xrightarrow{f} B \hookrightarrow UFB$ .

**Functoriality** follows from uniqueness given  $B \xrightarrow{f} C$ .  $F(gf)$  and  $(Fg)(Ff)$  are both homomorphisms extending  $A \xrightarrow{f} B \xrightarrow{g} C \rightarrow UFC$ .

(c) Given a set  $A$ , we write  $PA$  for the set of all subsets of  $A$ .

We can make  $P$  into a functor  $\mathbf{Set} \rightarrow \mathbf{Set}$ , given  $A \xrightarrow{f} B$ , we defined  $Pf(A') = \{f(a) \mid a \in A'\}$  for  $A' \subseteq A$ .

But we also have a functor  $P^* : \mathbf{Set} \rightarrow \mathbf{Set}^{op}$  defined on objects by  $P$ , but  $P^*f(B') = \{a \in A \mid f(a) \in B'\}$  for  $B' \subseteq B$ .

By a *contravariant* functor  $\mathcal{C} \rightarrow \mathcal{D}$ , we mean a functor  $\mathcal{C} \rightarrow \mathcal{D}^{op}$  (or  $\mathcal{C}^{op} \rightarrow \mathcal{D}$ ).

A *covariant* functor is one that doesn't reverse arrows (in *op* I guess?).

- (d) Let  $K$  be a field. We have a functor  $*$  :  $\mathbf{Mod}_K \rightarrow \mathbf{Mod}_K^{op}$  defined by  $V^* = \{ \text{linear maps } V \rightarrow K \}$ , and if  $V \xrightarrow{f} W$ ,  $f^*(\theta : W \rightarrow K) = \theta f$ .
- (e) We have a functor  $op$  :  $\mathbf{Cat} \rightarrow \mathbf{Cat}$ , which is the identity on morphisms (note that this is a covariant).
- (f) A functor between monoids is a monoid homomorphism.
- (g) A functor between posets is an order-preserving map.
- (h) Let  $G$  be a group. A functor  $F \circ G \rightarrow \mathbf{Set}$  consists of a set  $A = F*$  together with an action of  $G$  on  $A$ , i.e. a *permutation representation* of  $G$ . Similarly, a functor  $G \rightarrow \mathbf{Mod}_K$  is a  $K$ -linear representation of  $G$ .
- (i) The construction of the fundamental group  $\pi(X, X)$  of a space  $X$  with basepoint  $X$  is a functor  $\mathbf{Top}_* \rightarrow \mathbf{Gp}$  where  $\mathbf{Top}_*$  is the category of spaces with a chosen basepoint. Similarly, the fundamental groupoid is a functor  $\mathbf{Top} \rightarrow \mathbf{Gpd}$ , where  $\mathbf{Gpd}$  is the category of groupoids and functors between them.

**Definition.** (1.6)

Let  $\mathcal{C}$  and  $\mathcal{D}$  be categories and  $F, G : \mathcal{C} \rightrightarrows \mathcal{D}$  (why two arrows?) two functors. A *natural transformation*  $\alpha : F \rightarrow G$  consists of an assignment  $A \rightarrow \alpha_A$  from  $\text{ob } \mathcal{C}$  to  $\text{mor } \mathcal{D}$  (think about this), such that  $\text{dom}_{\alpha_A} = FA$  and  $\text{cod}_{\alpha_A} = GA$  for all  $A$ , and for all  $A \xrightarrow{f} B$  in  $\mathcal{C}$ , the square

$$\begin{array}{ccc} FA & \xrightarrow{Ff} & FB \\ \downarrow \alpha_A & & \downarrow \alpha_B \\ GA & \xrightarrow{Gf} & GB \end{array}$$

commutes (i.e.  $\alpha_B(Ff) = (Gf)_{\alpha_A}$ ).

(1.3) (l) Given categories  $\mathcal{C}$  and  $\mathcal{D}$ , we write  $[\mathcal{C}, \mathcal{D}]$  for the category whose objects are functors  $\mathcal{C} \rightarrow \mathcal{D}$  and whose morphisms are natural transformations.

**Example.** (1.7)

- (a) Let  $K$  be a field,  $V$  a vector space over  $K$ . There is a linear map  $\alpha_V : V \rightarrow V^{**}$  given by  $\alpha_V(v)\theta = \theta(v)$  for  $\theta \in V^*$ . This is the  $V$ -component of a natural transformation  $1_{\mathbf{Mod}_K} \rightarrow ** : \mathbf{Mod}_K \rightarrow \mathbf{Mod}_K$ .
- (b) For any set  $A$ , we have a mapping  $\sigma_A : A \rightarrow PA$  sending  $a$  to  $\{a\}$ . If  $f : A \rightarrow B$ , then  $Pf\{a\} = \{f(a)\}$ . So  $\sigma$  is a natural transformation  $1_{\mathbf{Set}} \rightarrow P$ .
- (c) Let  $F : \mathbf{Set} \rightarrow \mathbf{Gp}$  be the free group functor (1.5(b)), and  $U : \mathbf{Gp} \rightarrow \mathbf{Set}$  the forgetful functor. The inclusions  $A \rightarrow UFA$  form a natural transformation  $1_{\mathbf{Set}} \rightarrow UF$ .
- (d) Let  $G, H$  be groups and  $f, g : G \rightrightarrows H$  be two homomorphisms. A natural transformation  $\alpha : f \rightarrow g$  corresponds to an element  $h = \alpha_*$  of  $H$ , s.t.  $hf(x) \rightarrow g(x)h$  for all  $x \in G$  or equivalently  $f(x) = h^{-1}g(x)h$ , i.e.  $f$  and  $g$  are conjugate group homomorphisms.
- (e) Let  $A$  and  $B$  be two  $G$ -sets, regarded as functors:  $G \rightrightarrows \mathbf{Set}$ . A natural transformation  $A \rightarrow B$  is a function  $f$  satisfying  $f(g \cdot a) = g \cdot f(a)$  for all  $a \in A$ , i.e. a  $G$ -equivariant map.

**Lemma.** (1.8)

Let  $F, G : \mathcal{C} \rightrightarrows \mathcal{D}$  be two functors, and  $\alpha : F \rightarrow G$  a natural transformation. Then  $\alpha$  is an isomorphism in  $[\mathcal{C}, \mathcal{D}]$  iff each  $\alpha_A$  is an isomorphism in  $\mathcal{D}$ .

*Proof.* Forward is trivial. For backward, suppose each  $\alpha_A$  has an inverse  $\beta_A$ . Given  $f : A \rightarrow B$  in  $\mathcal{C}$ , we need to show that

$$\begin{array}{ccc} GA & \xrightarrow{Gf} & GB \\ \downarrow \beta_A & & \downarrow \beta_B \\ FA & \xrightarrow{Ff} & FB \end{array}$$

□

commutes. But as  $\alpha$  is natural,

$$(Ff)\beta_A = \beta_B\alpha_B(Ff)\beta_A = \beta_B(Gf)\alpha_A\beta_A = \beta_B(Gf)$$

So  $\beta$  is a natural transformation as well.

**Definition.** (1.9)

Let  $\mathcal{C}$  and  $\mathcal{D}$  be categories. By an *equivalence* between  $\mathcal{C}$  and  $\mathcal{D}$ , we mean a pair of functors  $F : \mathcal{C} \rightarrow \mathcal{D}$ ,  $G : \mathcal{D} \rightarrow \mathcal{C}$  together with natural isomorphisms  $\alpha : 1_{\mathcal{C}} \rightarrow GF$  and  $\beta : FG \rightarrow 1_{\mathcal{D}}$ .

We write  $\mathcal{C} \cong \mathcal{D}$  if  $\mathcal{C}$  and  $\mathcal{D}$  are equivalent.

We say a property  $P$  of categories is a *categorical property* if whenever  $\mathcal{C}$  has  $P$  and  $\mathcal{C} \cong \mathcal{D}$ , then  $\mathcal{D}$  has  $P$ .

For example, being a groupoid or a preorder are categorical properties, but being a group or a partial order are not.

**Example.** (1.10)

(a) The category **Part** is equivalent to the category **Set**<sub>\*</sub> of pointed sets (and basepoint preserving functions (as morphisms)):

- We define  $F : \mathbf{Set}_* \rightarrow \mathbf{Part}$  by  $F(A, a) = A \setminus \{a\}$ , and if  $f : (A, a) \rightarrow (B, b)$ , then  $Ff(x) = f(x)$  if  $f(x) \neq b$ , and undefined otherwise;
- and  $G : \mathbf{Part} \rightarrow \mathbf{Set}_*$  by  $G(A) = A^+ = (A \cup \{A\}, A)$ , and if  $f : A \rightarrow B$  is a partial function, we define  $Gf : A^+ \rightarrow B^+$  by  $Gf(x) = f(x)$  if  $x \in A$  and  $f(x)$  defined, and equals  $B$  otherwise.

The composite  $FG$  is the identity on **Part**, but  $GF$  is not the identity. However, there is an isomorphism  $(A, a) \rightarrow ((A \setminus \{a\})^+, A \setminus \{a\})$  sending  $a$  to  $A \setminus \{a\}$  and everything else to itself and this is natural.

Note that there can be no isomorphism from **Set**<sub>\*</sub> to **Part**, since **Part** has a 1-element isomorphism class  $\{\phi\}$  but **Set**<sub>\*</sub> doesn't.

(So we see that equivalent categories can be non-isomorphic. According to a [post](#) on SO, this usually happens when there are multiple copies of the *same* thing in one but not the other. However, we can't generally *discard obsolete copies* in one as that generally requires AC and is not a very useful thing to do anyway – In short, *identifying isomorphic objects is often an extremely bad idea.*)

(b) The category **fdMod**<sub>K</sub> of finite-dimensional vector spaces over  $K$  is equivalent to **fdMod**<sub>K</sub><sup>op</sup>, the functors in both directions are  $*$  (the dual operator) and both isomorphisms are the natural transformations of 1.7(a) (double dual).

(c) **fdMod**<sub>K</sub> is also equivalent to **Mat**<sub>K</sub> (1.3(j)):

We define  $F : \mathbf{Mat}_K \rightarrow \mathbf{fdMod}_K$  by  $F(n) = K^n$ , and  $F(A)$  is the linear map represented by  $A$  w.r.t. the standard bases of  $K^n$  and  $K^p$ .

To define  $G : \mathbf{fdMod}_K \rightarrow \mathbf{Mat}_K$ , choose a basis for each finite dimensional vector

space, and define  $G(V) = \dim V$ ,  $G(V \xrightarrow{f} W)$  to be the matrix representing  $f$  w.r.t. chosen bases.  $GF$  is the identity, provided we choose the standard bases for the spaces  $K^n$ ;  $FG \neq 1$ , but the chosen bases give isomorphisms  $FG(V) = K^{\dim V} \rightarrow V$  for each  $V$ , which form a natural isomorphism.

—Lecture 4—

**Definition.** (1.11)

Let  $\mathcal{C} \xrightarrow{F} \mathcal{D}$  be a functor.

(a) We say  $F$  is *faithful* if, given  $f, f' \in \text{mor } \mathcal{C}$  with  $\text{dom } f = \text{dom } f'$ ,  $\text{cod } f = \text{cod } f'$ , and  $Ff = Ff'$ , then  $f = f'$  (injectivity on morphisms. The name comes more from representation theory);

(b) We say  $F$  is *full* if, given  $FA \xrightarrow{g} FB$  in  $\mathcal{D}$ , there exists  $A \xrightarrow{f} B$  in  $\mathcal{C}$  with  $Ff = g$ . (this is something like surjectivity on morphisms, but see below);

(c) We say  $F$  is *essentially surjective* if, for every  $B \in \text{ob } \mathcal{D}$ , there exists  $A \in \text{ob } \mathcal{C}$  and isomorphism  $FA \rightarrow B$  in  $\mathcal{D}$ .

We say a subcategory  $\mathcal{C}' \subseteq \mathcal{C}$  is *full* if the inclusion  $\mathcal{C}' \rightarrow \mathcal{C}$  is a full functor (basically, if the objects are kept, any morphism between them must be kept). For example, **Gp** is a full subcategory of **Mon** (the category of all monoids), but **Mon** is not a full subcategory of the category **SGp** of semigroups (consider e.g. the homomorphism that sends everything in  $(\mathbb{Z}, \cdot)$  to  $(0, \cdot)$  (which is also a semigroup); but this doesn't preserve 1 so is not a morphism in **Mon**).

**Lemma.** (1.12)

Assuming the axiom of choice, a functor  $F : \mathcal{C} \rightarrow \mathcal{D}$  is part of an equivalence  $\mathcal{C} \simeq \mathcal{D}$  if it's full, faithful, and essentially surjective.

*Proof.*  $\Rightarrow$ : Suppose given  $G, \alpha, \beta$  as in (1.9). Then for each  $B \in \text{ob } \mathcal{D}$ ,  $\beta_B$  is an isomorphism  $FGB \rightarrow B$ , so  $F$  is essentially surjective.

Given  $A \xrightarrow{f} B$  in  $\mathcal{C}$ , we can recover  $f$  from  $Ff$  as composite  $A \xrightarrow{\alpha_A} GFA \xrightarrow{GFf} GFB \xrightarrow{\alpha_B^{-1}} B$ . Hence if  $A \xrightarrow{f'} B$  satisfies  $Ff = Ff'$ , then  $f = f'$ . So  $F$  is faithful;

Lastly, for fullness, given  $FA \xrightarrow{g} FB$ , define  $f$  to be the composite  $A \xrightarrow{\alpha_A} GFA \xrightarrow{Gg} GFB \xrightarrow{\alpha_B^{-1}} B$ . Then  $GFf = \alpha_B f \alpha_A^{-1}$ , which by construction is just  $Gg$ . But  $G$  is faithful for the same reason as  $f$ , so  $Ff = g$ .

$\Leftarrow$ : (need to find suitable  $G, \alpha, \beta$  for  $F$ .) For each  $B \in \text{ob } \mathcal{D}$ , choose  $GB \in \text{ob } \mathcal{C}$  and an isomorphism  $\beta_B : FGB \rightarrow B$  in  $\mathcal{D}$ . Given  $B \xrightarrow{g} B'$ , define  $Gg : GB \rightarrow GB'$  to be the unique morphism whose image under  $F$  is  $FGB \xrightarrow{\beta_B} B \xrightarrow{g} B' \xrightarrow{\beta_{B'}^{-1}} FGB'$ .

Uniqueness implies functoriality: given  $B' \xrightarrow{g'} B''$ ,  $(Gg')(Gg)$  and  $G(g'g)$  have the same image under  $F$ , so they are equal.

By construction,  $\beta$  is a natural transformation  $FG \rightarrow 1_{\mathcal{D}}$ .

Given  $A \in \text{ob } \mathcal{C}$ , define  $\alpha_A : A \rightarrow GFA$  to be the unique morphism whose image under  $F$  is  $FA \xrightarrow{\beta_{FA}^{-1}} FGFA$ .  $\alpha_A$  is an isomorphism, since  $\beta_{FA}$  also has a unique pre-image under  $F$ . And  $\alpha$  is a natural transformation, since any naturality



square for  $\alpha$  (the commutative square when we defined natural transformation) is mapped by  $F$  to a commutative square, and  $F$  is faithful.  $\square$

**Definition.** (1.13)

By a *skeleton* of a category, we mean a full subcategory  $\mathcal{C}_0$  containing one object from each isomorphism class. We say  $\mathcal{C}$  is *skeletal* if it's a skeleton of itself.

For example,  $\mathbf{Mat}_{\mathbf{K}}$  is a skeletal, and the image of  $F : \mathbf{Mat}_{\mathbf{K}} \rightarrow \mathbf{fdMod}_{\mathbf{K}}$  of 1.10(c) is a skeleton of  $\mathbf{fdMod}_{\mathbf{K}}$ .

(there are some examples on wikipedia)

Warning: almost any assertion about skeletons is equivalent to axiom of choice (see q2 on example sheet 1).

**Definition.** (1.14)

Let  $A \xrightarrow{f} B$  be a morphism in  $\mathcal{C}$ .

(a) We say  $f$  is a *monomorphism* (or  $f$  is *monic*) if, given any pair  $C \rightrightarrows_h^g A$ ,  $fg = fh$  implies  $g = h$ .

(b) We say  $f$  is an *epimorphism* (or *epic*) if it's a monomorphism in  $\mathcal{C}^{op}$ , i.e. if  $gf = hf$  implies  $g = h$ .

We denote monomorphisms by  $A \xrightarrow{f} B$ , and epimorphisms by  $A \xrightarrow{f} B$ .

Any isomorphism is monic and epic: more generally, if  $f$  has a left inverse (i.e.  $\exists g$  s.t.  $gf$  is an identity), then it's monic. We call such monomorphisms *split*.

We say  $\mathcal{C}$  is a *balanced* category if any morphism which is both monic and epic is an isomorphism.

**Example.** (1.15)

(a) As usual we consider **Set** first. In **Set**, monomorphisms correspond to injections ( $\Leftarrow$  is easy (ok); for  $\Rightarrow$ , take  $C \rightrightarrows 1 = \{*\}$ ), and epimorphisms correspond to surjections ( $\Leftarrow$  is easy; for  $\Rightarrow$ , use morphisms  $B \rightrightarrows 2 = \{0, 1\}$ ). So **Set** is balanced.

(b) In **Gp**, monomorphisms again correspond to injections (for  $\Rightarrow$  use homomorphisms  $\mathbb{Z} \rightarrow A$ ); epimorphisms again correspond to surjections ( $\Rightarrow$  use **free products with amalgamation** – this is a non-trivial fact about groups, read more if free). So **Gp** is also balanced.

(c) In **Rng** (obvious notation), monomorphisms correspond to injections (proof is much like for **Gp**). However, not all epimorphisms are surjective. For example

the inclusion  $\mathbb{Z} \rightarrow \mathbb{Q}$  is an epimorphism, since if  $\mathbb{Q} \rightrightarrows_g^f R$  (any ring) agree on all integers, they agree everywhere. So **Rng** is not balanced.

(d) One final example is **Top**. Again, monomorphisms are injections and epimorphisms are surjections (and vice versa): proof is similar to **Set** (check). However, **Top** is not balanced since a continuous bijection need not have continuous inverse.

## 2 The Yoneda Lemma

—Lecture 5—

**Definition.** (2.1)

We say a category  $\mathcal{C}$  is *locally small* if, for any two objects  $A, B$ , the morphisms  $A \rightarrow B$  in  $\mathcal{C}$  form a set  $\mathcal{C}(A, B)$ .

If we fix  $A$  and let  $B$  vary, the assignment  $B \rightarrow \mathcal{C}(A, B)$  becomes a functor  $\mathcal{C}(A, -) : \mathcal{C} \rightarrow \mathbf{Set}$ : given  $B \xrightarrow{f} C$ ,  $\mathcal{C}(A, f)$  is the mapping  $g \rightarrow fg$  for all  $g \in \mathcal{C}(B, C)$ . Similarly,  $A \rightarrow \mathcal{C}(A, B)$  defines a functor  $\mathcal{C}(-, B) : \mathcal{C}^{op} \rightarrow \mathbf{Set}$  (for  $A \xrightarrow{f} C \in \text{mor } \mathcal{C}^{op}$ , maps  $g \rightarrow gf$ ).

**Lemma.** (2.2)

- (i) Let  $\mathcal{C}$  be a locally small category,  $A \in \text{ob } \mathcal{C}$  and  $F : \mathcal{C} \rightarrow \mathbf{Set}$  a functor. Then natural transformations  $\mathcal{C}(A, -) \rightarrow F$  are in bijection with elements of  $FA$ ;
- (ii) Moreover, this bijection is natural in  $A$  and  $F$ .

*Proof.* (i) Given  $\alpha : \mathcal{C}(A, -) \rightarrow F$ , we define  $\Phi(\alpha) = \alpha_A(1_A) \in FA$ .<sup>1</sup>

Conversely, given  $x \in FA$ , we define  $\Psi(x) : \mathcal{C}(A, -) \rightarrow F$  by  $\Psi(x)_B(A \xrightarrow{f} B) = (Ff)(x) \in FB$ .<sup>2</sup>

$\Psi(x)$  is natural: given  $g : B \rightarrow C$ , we have

$$\begin{aligned}\Psi(x)_C \mathcal{C}(A, g)(f) &= \Psi(x)_C(gf) = F(gf)(x), \\ (Fg)\Psi(x)_B(f) &= (Fg)(Ff)(x) = F(gf)(x)\end{aligned}$$

Now given  $x \in FA$ ,  $\Phi\Psi(x) = \Psi(x)_A(1_A) = F(1_A)(x) = x$ ; given  $\alpha$ ,

$$\begin{aligned}\Psi\Phi(\alpha)_B(f)\Psi(\alpha_A(1_A))_B(f) &= Ff(\alpha_A(1_A)) \\ &= \alpha_B \mathcal{C}(A, f)(1_A) = \alpha_B(f)\end{aligned}$$

So  $\Psi\Phi(\alpha) = \alpha$ . So  $\Psi\Phi$  and  $\Phi\Psi$  are both identities on their respective domain (so we have a bijection).  $\square$

**Corollary.** (2.3)

The assignment  $A \rightarrow \mathcal{C}(A, -)$  defines a full and faithful functor  $\mathcal{C}^{op} \rightarrow [\mathcal{C}, \mathbf{Set}]$ .

*Proof.* Put  $F = \mathcal{C}(B, -)$  in 2.2(i): we get a bijection between  $\mathcal{C}(B, A)$  and morphisms  $\mathcal{C}(A, -) \rightarrow \mathcal{C}(B, -)$  in  $[\mathcal{C}, \mathbf{Set}]$ <sup>3</sup>. We need to verify this is functorial: but it sends  $f : B \rightarrow A$  to the natural transformation  $g \rightarrow gf$ . So functoriality follows from associativity.  $\square$

<sup>1</sup>Note  $1_A \in \mathcal{C}(A, A)$ , and  $\alpha_A \in \text{mor } \mathbf{Set}$  but  $\text{mor } \mathbf{Set}$  are just functions between sets, so this makes sense.

<sup>2</sup>It seems a bit confusing why this is a natural transformation, but looking carefully it basically defines a function between sets, i.e. is in  $\text{mor } \mathbf{Set}$ .

<sup>3</sup>Think very carefully about this... Given a morphism in  $\mathcal{C}(A, -) \rightarrow \mathcal{C}(B, -)$ , the above gives us a way to identify it uniquely with an element in  $\mathcal{C}(B, A)$  which is in  $\text{mor } \mathcal{C}^{op}$ . But that alone is not enough; we also need the above functor to take that morphism *directly* to the original morphism. Luckily this is the case by the proof of 2.2(i), which is also explained in the later half of the sentence above.

We call this functor (or the functor  $\mathcal{C} \rightarrow [\mathcal{C}^{op}, \mathbf{Set}]$  sending  $A$  to  $\mathcal{C}(-, A)$ ) the *Yoneda embedding* of  $\mathcal{C}$ , and denote it by  $Y$ .

Now let's go back to prove 2.2(ii):

*Proof.* (ii) Suppose for the moment that  $\mathcal{C}$  is small, so that  $[\mathcal{C}, \mathbf{Set}]$  is locally small.<sup>4</sup> Then we have two functors  $\mathcal{C} \times [\mathcal{C}, \mathbf{Set}] \rightarrow \mathbf{Set}$ : one sends  $(A, F)$  to  $FA$ , and the other is the composite:  $\mathcal{C} \times [\mathcal{C}, \mathbf{Set}] \xrightarrow{Y \times 1} [\mathcal{C}, \mathbf{Set}]^{op} \times [\mathcal{C}, \mathbf{Set}] \xrightarrow{[\mathcal{C}, \mathbf{Set}](-; -)} \mathbf{Set}$ .<sup>5</sup>

2.2(ii) says that these are naturally isomorphic. We can translate this into an elementary statement, making sense even when  $\mathcal{C}$  isn't small. Given  $A \xrightarrow{f} B$  and  $F \xrightarrow{\alpha} G$ , the two ways of producing an element of  $GB$  from a natural transformation  $\beta : \mathcal{C}(A, -) \rightarrow F$  give the same result, namely

$$\alpha_B(Ff)\beta_A(1_A) = (Gf)\alpha_A\beta_A(1_A)$$

which is equal to  $\alpha_B\beta_B(f)$ . □

**Definition.** (2.4)

We say a functor  $F : \mathcal{C} \rightarrow \mathbf{Set}$  is *representable* if it's isomorphic to  $\mathcal{C}(A, -)$  for some  $A$ . By a representation of  $F$ , we mean a pair  $(A, x)$  where  $x \in FA$  is such that  $\Psi(x)$  is an isomorphism.

We also call  $x$  a *universal element* of  $F$ .

**Corollary.** (2.5)

If  $(A, x)$  and  $(B, y)$  are both representations of  $F$ , then there's a unique isomorphism  $f : A \rightarrow B$  such that  $(Ff)(x) = y$ .

*Proof.* Consider the composite  $\mathcal{C}(B, -) \xrightarrow{\Psi(y)^{-1}} F \xrightarrow{\Psi(x)} \mathcal{C}(A, -)$ . By (2.3) this is of the form  $Y(f)$  for a unique isomorphism  $f : A \rightarrow B$ , and the diagram

$$\begin{array}{ccc} \mathcal{C}(B, -) & \xrightarrow{Y(f)} & \mathcal{C}(A, -) \\ & \searrow \Psi(y) & \swarrow \Psi(x) \\ & F & \end{array}$$

commutes iff  $(Ff)(x) = y$ . □

**Example.** (2.6)

(a) The forgetful functor  $\mathbf{Gp} \rightarrow \mathbf{Set}$  is representable by  $(\mathbb{Z}, 1)$ ,  $\mathbf{Rng} \rightarrow \mathbf{Set}$  by  $(\mathbb{Z}[X], X)$ , and  $\mathbf{Top} \rightarrow \mathbf{Set}$  by  $(\{\ast\}, \ast)$ .

(b) The functor  $P^* : \mathbf{Set}^{op} \rightarrow \mathbf{Set}$  is representable by  $(\{0, 1\}, \{1\})$ : this is the bijection between subsets and characteristic functions.

(c) Let  $G$  be a group. The unique (up to isomorphism) representable functor  $G(\ast, -) : G \rightarrow \mathbf{Set}$  is the *Cayley representation* of  $G$ , i.e. the set  $UG$  with  $G$  acting by left multiplication.

<sup>4</sup>Elements in  $\text{mor}[\mathcal{C}, \mathbf{Set}]$  correspond to those in  $\text{mor}\mathcal{C}^{op}$  by Yoneda.

<sup>5</sup>The second operator maps two functors to the set of natural transformations between them?

(d) Let  $A$  and  $B$  be two objects of a small category  $\mathcal{C}$ . We have a functor  $\mathcal{C}^{op} \rightarrow \mathbf{Set}$  sending  $C$  to  $\mathcal{C}(C, A) \times \mathcal{C}(C, B)$ . A representation of this, if it exists, is called a (categorical) *product* of  $A$  and  $B$ , and denoted  $(A \times B, (A \times B \xrightarrow{\pi_1} A, A \times B \xrightarrow{\pi_2} B))$ .

This pair has the property that, for any pair  $(C \xrightarrow{f} A, C \xrightarrow{g} B)$ , there's a unique  $C \xrightarrow{h} A \times B$  with  $\pi_1 h = f$  and  $\pi_2 h = g$ .

Products exist in many categories of interest: in **Set**, **Gp**, **Rng**, **Top**, ..., they are *just* cartesian products, in posets they are binary meets (see sheet 1 Q1).

Dually, we have the notion of *coproduct*  $(A + B, A \xrightarrow{\mu_1} A + B, B \xrightarrow{\mu_2} A + B)$ . These also exist in many categories of interest.

—Lecture 6—

(f) (Lecturer didn't like (e) so jumped to (f) directly) Let  $A \xrightleftharpoons[g]{f} B$  be morphisms in locally small category  $\mathcal{C}$ . We have a functor  $F : \mathcal{C}^{op} \rightarrow \mathbf{Set}$  defined by

$$F(C) = \{h \in \mathcal{C}(C, A) \mid fh = gh\}$$

A representation (see (2.4)) of  $F$ , if it exists, is called an *equalizer* of  $(f, g)$ : It consists of an object  $E$  and a morphism  $E \xrightarrow{e} A$  s.t.  $fe = ge$ , and every  $h$  with  $fh = gh$  factors uniquely (see proof of 2.9(i) which gives an insight of what this means) through  $e$ .

In **Set**, we take  $E = \{x \in A \mid f(x) = g(x)\}$  and  $e = \text{inclusion}$ . Similar constructions work in **Gp**, **Rng**, **Top**, ...

Dually, we have the notion of *coequalizer*.

**Remark.** (2.7)

If  $e$  occurs as an equalizer, then it is a monomorphism, since any  $h$  factors through it in at most one way. We say a monomorphism is *regular* if it occurs as an equalizer.

Split monomorphisms are regular (cf sheet1 Q6(i)).

Note that regular epic monomorphisms are isomorphisms: if the equalizer  $e$  of  $(f, g)$  is epic, then  $f = g$ , so  $e \cong 1_{\text{cod } e}$ .

**Definition.** (2.8)

Let  $\mathcal{C}$  be a category,  $\mathcal{G}$  a class of objects of  $\mathcal{C}$ .

(a) We say  $\mathcal{G}$  is a *separating family* for  $\mathcal{C}$ , if given  $A \xrightleftharpoons[g]{f} B$  such that  $fh = gh$  for

all  $G \xrightarrow{h} A$  with  $G \in \mathcal{G}$ , then  $f = g$ .

(i.e. the functors  $\mathcal{C}(G, -)$ ,  $G \in \mathcal{G}$ , are collectively faithful.)

(b) We say  $\mathcal{G}$  is a *detecting family* if, given  $A \xrightarrow{f} B$  such that every  $G \xrightarrow{h} B$  with  $G \in \mathcal{G}$  factors uniquely through  $f$ , then  $f$  is an isomorphism.

If  $\mathcal{G} = \{G\}$ , we call  $G$  a *separator/detector*.

**Lemma.** (2.9)

(i) If  $\mathcal{C}$  is a balanced category, then any separating family is detecting.

(ii) If  $\mathcal{C}$  has equalizers, then any detecting family is separating.

*Proof.* (i) Suppose  $\mathcal{G}$  is separating and  $A \xrightarrow{f} B$  satisfies the condition of 2.8(b). If  $B \xrightleftharpoons[h]{g} C$  satisfy  $gf = hf$ , then  $gx = hx$  for every  $G \xrightarrow{x} B$ , so  $g = h$ , i.e.  $f$  is

epic.

Similarly if  $D \begin{smallmatrix} k \\ \rightrightarrows \\ l \end{smallmatrix} A$  satisfy  $fk = fl$ , then  $ky = ly$  for any  $G \xrightarrow{y} D$ , since both are factorizations of  $fky$  through  $f$ . So  $k = l$ , i.e.  $f$  is monic.

But  $\mathcal{C}$  is balanced. So  $f$  is an isomorphism.

(ii) Suppose  $\mathcal{G}$  is detecting and  $A \begin{smallmatrix} f \\ \rightrightarrows \\ g \end{smallmatrix} B$  satisfies the condition of 2.8(a). Then the equalizer  $E \xrightarrow{e} A$  of  $(f, g)$  is isomorphism, so  $f = g$ .  $\square$

**Example.** (2.10)

(a) In  $[\mathcal{C}, \mathbf{Set}]$ , the family  $\{\mathcal{C}(A, -) | A \in \text{ob } \mathcal{C}\}$  is both separating and detecting (just a restatement of Yoneda Lemma).

(b) In  $\mathbf{Set}$ ,  $1 = \{*\}$  (any one element set) is both a separator and a detector, since it represents the identity functor  $\mathbf{Set} \rightarrow \mathbf{Set}$ .

Similarly,  $\mathbb{Z}$  is both in  $\mathbf{Gp}$ , since it represents the forgetful functor  $\mathbf{Gp} \rightarrow \mathbf{Set}$ . Also,  $2 = \{0, 1\}$  is a coseparator and a codetector in  $\mathbf{Set}$ , since it represents  $P^* : \mathbf{Set}^{op} \rightarrow \mathbf{Set}$ .

(c) In  $\mathbf{Top}$ ,  $1 = \{*\}$  is a separator since it represents the forgetful functor  $\mathbf{Top} \rightarrow \mathbf{Set}$ , but not a detector.

In fact,  $\mathbf{Top}$  has no detecting *set* of objects (note that this doesn't mean it has no detecting family).

For any infinite cardinal  $\kappa$ , let  $X$  be a discrete space of cardinality  $\kappa$ , and  $Y$  the same set with  $co-< \kappa$  topology, i.e.  $F \subseteq Y$  is closed iff  $F = Y$  or  $\text{Card } F < \kappa$  (think about, e.g. cocountable topology, then this name makes sense).

The identity  $X \rightarrow Y$  is continuous, but not a homeomorphism (topologically). So if  $\{G_i | i \in I\}$  is any set of spaces, taking  $\kappa > \text{Card } G_i$  for all  $i$  yields an example to show that the set is not detecting.

(d) (some Algebraic Topology stuff) Let  $\mathcal{C}$  be the category of pointed connected  $CW$ -complexes and homotopy classes of (basepoint-preserving) continuous mappings.

JHC Whitehead proved that  $X \xrightarrow{f} Y$  in this category induces isomorphisms  $\pi_n(X) \rightarrow \pi_n(Y)$  for all  $n$ , then it's an isomorphism in  $\mathcal{C}$ .

This says that  $\{S^n | n \geq 1\}$  is a detecting set of  $\mathcal{C}$ .

But PJ Freyd showed there is no faithful functor  $\mathcal{C} \rightarrow \mathbf{Set}$ , so no separating *set*: if  $\{G_i | i \in I\}$  were separating, then  $x \rightarrow \coprod \mathcal{C}(G_i, x)$  (disjoint unions?) would be faithful.

Note that any functor of the form  $\mathcal{C}(A, -)$  preserves monomorphisms, but they don't normally preserve epimorphisms.

**Definition.** (2.11)

We say an object  $P$  is *Projective* if, given

$$\begin{array}{c} P \\ \downarrow f \\ A \xrightarrow{e} B \end{array}$$

(recall the two head right arrow means epimorphisms) there exists  $P \xrightarrow{g} A$  with  $eg = f$ .

(If  $\mathcal{C}$  is locally small, this says  $\mathcal{C}(P, -)$  preserves epimorphisms).

Dually, an *injective* object of  $\mathcal{C}$  is a projective object of  $\mathcal{C}^{op}$ .

Given a class  $\mathcal{E}$  of epimorphisms, we say  $P$  is  $\mathcal{E}$ -projective if it satisfies the condition for all  $e \in \mathcal{E}$ .

**Lemma.** (2.12)

Representable functors are (pointwise)(?) projective in  $[\mathcal{C}, \mathbf{Set}]$ .

*Proof.* Suppose given

$$\begin{array}{c} \mathcal{C}(A, -) \\ \downarrow \beta \\ F \xrightarrow{\alpha} G \end{array}$$

where  $\alpha$  is pointwise surjective. By Yoneda,  $\beta$  corresponds to some  $y \in GA$ , and we can find  $x \in FA$  with  $\alpha_A(x) = y$ . Now if  $\gamma : \mathcal{C}(A, -) \rightarrow F$  corresponds to  $x$ , then naturality of the Yoneda bijection yields  $\alpha\gamma = \beta$ .  $\square$

—Lecture 7—

First example class: Friday 26th October, 2pm MR3.

Lecture is happy to mark any question we hand in!

### 3 Adjunctions

#### Definition. (3.1)

Let  $\mathcal{C}$  and  $\mathcal{D}$  be two categories and  $\mathcal{C} \xrightarrow{F} \mathcal{D}$ ,  $\mathcal{D} \xrightarrow{G} \mathcal{C}$  two functors.

By an *adjunction* between  $F$  and  $G$  we mean a bijection between morphisms

$FA \xrightarrow{\hat{f}} B$  in  $\mathcal{D}$  and morphisms  $A \xrightarrow{f} GB$  in  $\mathcal{C}$ , which is natural in  $A$  and  $B$ , i.e. given  $A' \xrightarrow{g} A$  and  $B \xrightarrow{h} B'$ , we have  $h\hat{f}(Fg) = \widehat{(Gh)}fg : FA' \rightarrow B'$ .

We say  $F$  is *left adjoint* to  $G$ , and write  $(F \dashv G)$ .

#### Example. (3.2)

(a) The free functor  $\mathbf{Set} \xrightarrow{F} \mathbf{Gp}$  is left adjoint to the forgetful functor  $\mathbf{Gp} \xrightarrow{U} \mathbf{Set}$ , since any function  $f : A \rightarrow UB$  extends uniquely to a homomorphism  $\hat{f} : FA \rightarrow B$ .

Naturality in  $B$  is *easy* (lecturer says so), naturality in  $A$  follows from the definition of  $F$  as a functor.

(b) The forgetful functor  $\mathbf{Top} \xrightarrow{U} \mathbf{Set}$  has a left adjoint  $D$  which equips any set with the discrete topology, *and* also a right adjoint  $I$  which equips a set  $A$  with the discrete (lecturer had *indiscrete* here? Ok, it seems that he actually did mean indiscrete – check) topology  $\{\phi, A\}$ .

(c) The functor  $\text{ob} : \mathbf{Cat} \rightarrow \mathbf{Set}$  (recall  $\mathbf{Cat}$  is the category of small categories) has a left adjoint  $D$  sending  $A$  to the *discrete* category with  $\text{ob}(DA) = A$  and only identity morphisms, and a right adjoint  $I$  sending  $A$  to the category with  $\text{ob}(IA) = A$  and one morphism  $x \rightarrow y$  for each  $(x, y) \in A \times A$ . In this case  $D$  in turn has a left adjoint  $\pi_0$  sending a small category  $\mathcal{C}$  to its set of *connected components*, i.e. the quotient of  $\text{ob}\mathcal{C}$  by the smallest equivalence relation identifying  $\text{dom } f$  with  $\text{cod } f$  for all  $f \in \text{mor } \mathcal{C}$ .

(d) Let  $M$  be the monoid  $\{1, e\}$  with  $e^2 = e$ . An object of  $[M, \mathbf{Set}]$  is a pair  $(A, e)$  (the images of the functor?), where  $e : A \rightarrow A$  satisfies  $e^2 = e$ .

We have a functor  $G : [M, \mathbf{Set}] \rightarrow \mathbf{Set}$  sending  $(A, e)$  to  $\{x \in A | e(x) = x\} = \{e(x) | x \in A\}$  and a functor  $F : \mathbf{Set} \rightarrow [M, \mathbf{Set}]$  sending  $A$  to  $(A, 1_A)$ .

I claim  $(F \dashv G \dashv F)$ : given  $f : (A, 1_A) \rightarrow (B, e)$ , it must take values in  $G(B, e)$ , and any  $g : (B, e) \rightarrow (A, 1_A)$  is determined by its values on the image of  $e$ .

(e) Let  $\mathbf{1}$  be the discrete category with one object  $*$ . For any  $\mathcal{C}$ , there's a unique functor  $\mathcal{C} \rightarrow \mathbf{1}$ : a left adjoint for this picks out an *initial* object of  $\mathcal{C}$ , i.e. an object  $I$  s.t. there exists a unique  $I \rightarrow A$  for each  $A \in \text{ob } \mathcal{C}$ .

Dually, a right adjoint for  $\mathcal{C} \rightarrow \mathbf{1}$  corresponds to a *terminal* object of  $\mathcal{C}$  (think about what this means).

(f) Let  $A \xrightarrow{f} B$  be a morphism in  $\mathbf{Set}$ . We can regard  $PA$  and  $PB$  as posets, and we have functors  $PA \xrightleftharpoons[P^*f]{Pf} PB$ .

I claim  $(PF \dashv P^*f)$ : we have  $Pf(A') \subseteq B' \iff f(x) \in B'$  for all  $x \in A' \iff A' \subseteq P^*f(B')$ .

(g) (*Galois Connection*) Suppose given sets  $A, B$  and a relation  $R \subseteq A \times B$ . We define mappings  $(-)^l, (-)^r$  between  $PA$  and  $PB$  by

$$S^r = \{y \in B | (\forall x \in S)((x, y) \in R)\} \text{ for } S \subseteq A$$

$$T^l = \{x \in A | (\forall y \in T)((x, y) \in R)\} \text{ for } T \subseteq B$$

The mappings are order-reserving (i.e. contravariant functors), and  $T \subseteq S^r \iff$

$$S \times T \subseteq R \iff S \subseteq T^l.$$

We say  $()^r$  and  $()^l$  are *adjoint on the right*.

(h) Let's now consider, as a functor,  $P^* : \mathbf{Set}^{op} \rightarrow \mathbf{Set}$  is self-adjoint on the right, since functions  $A \rightarrow PB$  correspond bijectively to subsets of  $A \times B$ , and hence to functions  $B \rightarrow PA$ .

**Theorem.** (3.3)

Let  $G : \mathcal{D} \rightarrow \mathcal{C}$  be a functor. Then specifying a left adjoint for  $G$  is equivalent to specifying an initial object of  $(A \downarrow G)$  for each  $A \in \text{ob } \mathcal{C}$ , where  $(A \downarrow G)$  has objects  $(B, f)$  with  $A \xrightarrow{f} GB$ , and morphisms  $(B, f) \rightarrow (B', f')$  are morphisms  $B \xrightarrow{g} B'$  such that

$$\begin{array}{ccc} A & \xrightarrow{f} & GB \\ & \searrow f' & \swarrow Gg \\ & GB' & \end{array}$$

commutes.

*Proof.* Suppose given  $(F \dashv G)$ . Consider the morphism  $\eta_A : A \rightarrow GFA$  correspond to  $FA \xrightarrow{\eta} FA$ . Then  $(FA, \eta_A)$  is an object of  $(A \downarrow G)$ . Moreover, given  $g : FA \rightarrow B$  and  $f : A \rightarrow GB$ , the diagram

$$\begin{array}{ccc} A & \xrightarrow{\eta_A} & GFA \\ & \searrow f & \swarrow Gg \\ & GB & \end{array}$$

commutes iff

$$\begin{array}{ccc} FA & \xrightarrow{1_A} & FA \\ & \searrow \hat{f} & \swarrow g \\ & B & \end{array}$$

commutes, i.e.  $g = \hat{f}$ .

So  $(FA, \eta_A)$  is initial in  $(A \downarrow G)$ .

Conversely, suppose given an initial object  $(FA, \eta_A)$  for each  $(A \downarrow G)$ . Given  $A \xrightarrow{f} A'$ , we define  $Ff : FA \rightarrow FA'$  to be the unique morphism making

$$\begin{array}{ccc} A & \xrightarrow{\eta_A} & GFA \\ \downarrow f & & \downarrow GFf \\ A' & \xrightarrow{\eta_{A'}} & GFA' \end{array}$$

commute.

Functoriality follows from uniqueness: given  $f' : A' \rightarrow A''$ ,  $F(f'f)$  and  $(Ff')(Ff)$  are both morphisms  $(FA, \eta_A) \rightarrow (FA'', \eta_{A''} F'f)$  in  $(A \downarrow G)$ .

Note that we haven't finished: we still have to verify natural adjunctions. We'll finish off this next monday.



—Lecture 8—

It's next monday now! Let's finish the proof:

To show  $F \dashv G$ : given  $A \xrightarrow{f} GB$ , we define  $\hat{f} : FA \rightarrow B$  to be the unique morphism  $(FA, \eta_A) \rightarrow (B, f)$  in  $(A \downarrow G)$ . This is a bijection with inverse  $(FA \xrightarrow{g} B) \rightarrow (A \xrightarrow{\eta_a} GFA \xrightarrow{Gg} GB)$ . The latter mapping is natural in  $B$ , as  $G$  is a functor; and also in  $A$ , since by construction,  $\eta$  is a natural transformation  $1_C \rightarrow GF$ .  $\square$

Given an adjunction  $(F \dashv G)$ , the natural transformation  $\eta : 1_C \rightarrow GF$  emergin in the above proof (3.3) is called the *unit* of the adjunction.

Dually, we have a natural transformation traditionally denoted  $\varepsilon : FG \rightarrow 1_D$  s.t.  $\varepsilon_B : FGB \rightarrow B$  corresponds to  $GB \xrightarrow{1_{GB}} GB$ , is called the *counit*.

**Corollary.** (3.4)

If  $F$  and  $F'$  are both left adjoint to  $G : \mathcal{D} \rightarrow \mathcal{C}$ , then they are naturally isomorphic.

*Proof.* For any  $A$ ,  $(FA, \eta_A)$  and  $(F'A, \eta'_A)$  are both initial in  $(A \downarrow G)$ , so there's a unique isomorphism  $\alpha_A : (FA, \eta_A) \rightarrow (F'A, \eta'_A)$ .

In any naturality square for  $\alpha$ , the two ways round are both morphisms in  $(A \downarrow G)$  whose domain is initial, so they are equal. So  $\alpha$  is not only just an isomorphism (but also natural).  $\square$

**Lemma.** (3.5)

Given  $\mathcal{C} \xrightleftharpoons[F]{F} \mathcal{D} \xrightleftharpoons[K]{H} \mathcal{E}$ , with  $(F \dashv G)$  and  $(H \dashv K)$ , we have  $(HF \dashv GK)$ .

*Proof.* We have bijections between morphisms  $A \rightarrow GKC$ , morphisms  $FA \rightarrow KC$  and morphisms  $HFA \rightarrow C$ , which are both natural in  $A$  and  $C$ .  $\square$

**Corollary.** (3.6)

Given a commutative square

$$\begin{array}{ccc} \mathcal{C} & \rightarrow & \mathcal{D} \\ \downarrow & & \downarrow \\ \mathcal{E} & \rightarrow & \mathcal{F} \end{array}$$

of categories and functors, if the functors all have left adjoints, then the diagram of left adjoints commutes up to natural isomorphisms (?).

*Proof.* By (3.5), both ways round the diagram of left adjoints are left adjoint to the composite  $\mathcal{C} \rightarrow \mathcal{F}$ , so by (3.4) they are isomorphic.  $\square$

**Theorem.** (3.7)

Given functors  $\mathcal{C} \xrightleftharpoons[G]{F} \mathcal{D}$ , specifying an adjunction  $(F \dashv G)$  is equivalent to specifying natural transformations  $\eta : 1_C \rightarrow GF$ ,  $\varepsilon : FG \rightarrow 1_D$  satisfying the commutative diagrams,

$$\begin{array}{ccc}
F & \xrightarrow{F_\eta} & FGF \\
& \searrow 1_F & \downarrow \varepsilon_F \\
& & F
\end{array}
\quad \text{and} \quad
\begin{array}{ccc}
G & \xrightarrow{\eta_G} & GFG \\
& \searrow 1_G & \downarrow G_\varepsilon \\
& & G
\end{array}$$

which are sometimes called the *triangular identities* (for obvious reason).

*Proof.* First suppose we are given  $(F \dashv G)$ . Define  $\eta$  and  $\varepsilon$  as in (3.3) and its dual; now consider the composite

$$FA \xrightarrow{F_{\eta_A}} FGF A \xrightarrow{\varepsilon_{FA}} FA$$

under the adjunction, this corresponds to

$$A \xrightarrow{\eta_A} GFA \xrightarrow{1_{GFA}} GFA$$

But this also corresponds to  $1_{FA}$ , so  $\varepsilon_{FA} \cdot F_{\eta_A} = 1_{FA}$ .

The other identity is dual to this one (check).

Conversely, suppose we are given  $\eta$  and  $\varepsilon$  satisfying the triangular identities. Given  $A \xrightarrow{f} GB$ , let  $\Phi(f)$  be the composite  $FA \xrightarrow{Ff} FGB \xrightarrow{\varepsilon_B} B$ ; and given  $FA \xrightarrow{g} B$ , let  $\Psi(g)$  be  $A \xrightarrow{\eta_A} GFA \xrightarrow{Gg} GB$ . Then  $\Phi$  and  $\Psi$  are both natural; we now need to show they are inverse to each other. Let's do  $\Psi\Phi$ , say: now

$$\begin{aligned}
\Psi\Phi(A \xrightarrow{f} GB) &= A \xrightarrow{\eta_A} GFA \xrightarrow{GFf} GFGB \xrightarrow{G\varepsilon_B} GB \\
&= A \xrightarrow{f} FB \xrightarrow{\eta_{GB}} GFGB \xrightarrow{G\varepsilon_B} GB \\
&= f
\end{aligned}$$

and dually,  $\Phi\Psi(g) = g$ . □

**Lemma.** (3.8)

Suppose given  $\mathcal{C} \xrightleftharpoons[G]{F} \mathcal{D}$  and natural isomorphisms  $\alpha : 1_{\mathcal{C}} \rightarrow GF$ ,  $\beta : FG \rightarrow 1_{\mathcal{D}}$ .

Then there are isomorphisms  $\alpha' : 1_{\mathcal{C}} \rightarrow GF$ ,  $\beta' : FG \rightarrow 1_{\mathcal{D}}$  which satisfy the triangular identities. So  $(F \dashv G)$  (and  $(G \dashv F)$ ).

*Proof.* We define  $\alpha' = \alpha$  and, in attempt to fix  $\beta'$ , define  $\beta'$  to be the composite

$$FG \xrightarrow{(FGB)^{-1}} FGF G \xrightarrow{(F\alpha_G)^{-1}} FG \xrightarrow{\beta} 1_{\mathcal{D}}$$

Note that  $FGB = \beta_{FG}$ , since

$$\begin{array}{ccc}
FGFG & \xrightarrow{FGB} & FG \\
\downarrow \beta_{FG} & & \downarrow B \\
FG & \xrightarrow{\beta} & 1_{\mathcal{D}}
\end{array}$$

commutes by naturality of  $\beta$ , and  $\beta$  is monic. So it doesn't matter which way we choose above.

Now  $(\beta'_F)(F\alpha')$  is the composite

$$\begin{aligned}
F &\xrightarrow{F\alpha} FGF \xrightarrow{(\beta_{FGF})^{-1}} FGF GF \xrightarrow{(F\alpha_{GF})^{-1}} FGF \xrightarrow{\beta_F} F \\
&= F \xrightarrow{(\beta_F)^{-1}} FGF \xrightarrow{FGF\alpha} FGF GF \xrightarrow{(F\alpha_{GF})^{-1}} FGF \xrightarrow{\beta_F} F \\
&= F \xrightarrow{(\beta_F)^{-1}} FGF \xrightarrow{\beta_F} F \\
&= 1_F
\end{aligned}$$

Since  $GF\alpha = \alpha_{GF}$  (similar reasoning as previous).

Now similarly  $(G\beta')(G'\alpha')$  is

$$\begin{aligned}
G &\xrightarrow{\alpha_G} GFG \xrightarrow{(GF\alpha_G)^{-1}} GFG GF \xrightarrow{(GF\alpha_G)^{-1}} GFG \xrightarrow{G\beta} G \\
&= G \xrightarrow{(G\beta)^{-1}} GFG \xrightarrow{\alpha_{GFG}} GFG GF \xrightarrow{(GF\alpha_G)^{-1}} GFG \xrightarrow{G\beta} G \\
&= G \xrightarrow{(G\beta)^{-1}} GFG \xrightarrow{G\beta} G \\
&= 1_G
\end{aligned}$$

□

**Lemma.** (3.9)

Suppose  $G : \mathcal{D} \rightarrow \mathcal{C}$  has a left adjoint  $F$  with counit  $\varepsilon : FG \rightarrow 1_{\mathcal{D}}$ , then:

- (i)  $G$  is faithful iff  $\varepsilon$  is pointwise epic;
  - (ii)  $G$  is full and faithful iff  $\varepsilon$  is an isomorphism.
- (and of course the dual results for unit – change epic to monic).

*Proof.* (i) Given  $B \xrightarrow{g} B'$ ,  $Gg$  corresponds, under the adjunction, to the composite  $FGB \xrightarrow{\varepsilon_B} B \xrightarrow{g} B'$ . Hence the mapping  $g \rightarrow Gg$  is injective on morphisms with domain  $B$  (and specified codomain) iff  $g \rightarrow g\varepsilon_B$  is injective, i.e. iff  $\varepsilon_B$  is an epimorphism.

(ii) The proof of this is actually very similar:  $G$  is full and faithful iff  $g \rightarrow g\varepsilon_B$  is bijective, but that forces  $\varepsilon$  to be an isomorphism: if  $\alpha : B \rightarrow FGB$  is such that  $\alpha\varepsilon_B = 1_{FGB}$ , then this must be a two sided inverse as  $\varepsilon_B\alpha\varepsilon_B = \varepsilon_B$ , whence  $\varepsilon_B\alpha = 1_B$ . So  $\varepsilon_B$  is an isomorphism, for all  $B$ . □

—Lecture 9—

**Definition.** (3.10) By a *reflection*, we mean an adjunction in which the right adjoint is full and faithful (equivalently, the counit is an isomorphism).

We say a full subcategory  $\mathcal{C}' \subseteq \mathcal{C}$  is *reflective* if the inclusion  $\mathcal{C}' \rightarrow \mathcal{C}$  has a left adjoint.

**Example.** (3.11)

(a) The category **AbGp** of abelian groups is reflective in **Gp**, the left adjoint sends a group  $G$  to its *abelianization*  $G/G'$ , where  $G'$  is the subgroup generated by all commutators  $[x, y] = xyx^{-1}y^{-1}$ ,  $x, y \in G$ , which is always a normal subgroup of  $G$  (see part II Galois Theory).

The unit of the adjunction is the quotient map  $G \rightarrow G/G'$ .

(b) Given an abelian group  $A$ , let  $A_t$  denote the torsion subgroup, i.e. the subgroup of elements of finite order. The assignment  $A \rightarrow A/A_t$  gives a left adjoint to the inclusion  $\mathbf{tfAbGp} \rightarrow \mathbf{AbGp}$  where  $\mathbf{tfAbGp}$  is the full subcategory of torsion-free abelian groups.  $A \rightarrow A_t$  is right adjoint to the inclusion  $\mathbf{tAbGp} \rightarrow \mathbf{AbGp}$ , so this subcategory is coreflective.

(c) Let  $\mathbf{KHaus} \subseteq \mathbf{Top}$  be the full subcategory of compact Hausdorff spaces (see part IB Metric and Topological Spaces). The inclusion  $\mathbf{KHaus} \rightarrow \mathbf{Top}$  has a left adjoint  $\beta$ , the *Stone-Ćech compactification*.

(d) Let  $X$  be a topological space. We say  $A \subseteq X$  is *sequentially closed* if  $x_n \rightarrow x_\infty$  and  $x_n \in A$  for all  $n$  implies  $x_\infty \in A$ .

We say  $X$  is *sequential* if all sequentially closed sets are closed. Given a non-sequential space  $X$ , let  $X_s$  be the same set with topology given by the sequentially open sets in  $X$ ; the identity  $X_s \rightarrow X$  is continuous, and defines the counit of an adjunction between the inclusion  $\mathbf{Seq} \rightarrow \mathbf{Top}$  and the functor  $X \rightarrow X_s$ .

(e) If  $X$  is a topological space, the poset  $CX$  of closed subsets of  $X$  is reflective in the full power set  $\mathcal{P}X$ , with reflector given by closure, and the poset  $OX$  of open subsets is coreflective, with reflector given by interior.

## 4 Limits

**Definition.** (4.1)

(a) Let  $\mathcal{J}$  be a category (almost always small, and often finite). By a *diagram of shape  $\mathcal{J}$*  in  $\mathcal{C}$ , we mean a functor  $D : \mathcal{J} \rightarrow \mathcal{C}$ . The objects  $D(j), j \in \text{ob } \mathcal{J}$ , are called *vertices* of the diagram, and the morphism  $D(\alpha), \alpha \in \text{mor } \mathcal{J}$  are called *edges* of  $D$ .

For example, if  $\mathcal{J}$  is the category



with 4 objects and 5 non-identity morphisms, a diagram of shape  $\mathcal{J}$  is a commutative square

$$\begin{array}{ccc} A & \xrightarrow{f} & B \\ \downarrow g & & \downarrow h \\ C & \xrightarrow{k} & D \end{array}$$

If  $\mathcal{J}$  is  $\begin{array}{ccc} \cdot & \longrightarrow & \cdot \\ \downarrow & \searrow & \downarrow \\ \cdot & \longrightarrow & \cdot \end{array}$ , a diagram of shape  $\mathcal{J}$  is a not-necessarily-commutative square.

(b) Given  $D : \mathcal{J} \rightarrow \mathcal{C}$ , a *cone* over  $D$  consists of an object  $A$  of  $\mathcal{C}$  (the *apex* of the cone) together with morphisms  $A \xrightarrow{\lambda_j} D(j)$  for each  $j \in \text{ob } \mathcal{J}$ , such that

$$\begin{array}{ccc} & A & \\ \lambda_j \swarrow & & \searrow \lambda_{j'} \\ D(j) & \xrightarrow{D(\alpha)} & D(j') \end{array} \quad \text{commutes for all } j \xrightarrow{\alpha} j' \text{ in } \text{mor } \mathcal{J}.$$

(The  $\lambda_j$  are called the *legs* of the cone).



Given cones  $(A, (\lambda_j)_{j \in \text{ob } \mathcal{J}})$  and  $(B, (\mu_j)_{j \in \text{ob } \mathcal{J}})$ , a *morphism* of cones between

them is a morphism  $A \xrightarrow{f} B$  s.t.  $\begin{array}{ccc} A & \xrightarrow{f} & B \\ \lambda_j \searrow & & \swarrow \mu_j \\ & D(j) & \end{array}$  commutes for all  $j$ .

We write  $\mathbf{Cone}(D)$  for the category of cones over  $D$  (I guess with morphisms being all the possible ones from above?).

(c) A *limit* for  $D$  is a terminal object of  $\mathbf{Cone}(D)$ , if this exists.

Dually, we have the notion of cone under a diagram, and of colimit (= initial cone under  $D$ ).

Alternatively, if  $\mathcal{C}$  is locally small, and  $\mathcal{J}$  is small, we have a functor  $\mathcal{C}^{op} \rightarrow \mathbf{Set}$  sending  $A$  to the set of cones with apex  $A$ . A limit for  $D$  is a representation of this functor.

If  $\triangle A$  denotes the constant diagram of shape  $\mathcal{J}$  with all vertices  $A$  and all edges  $1_A$ , then a cone over  $D$  with apex  $A$  is the same thing as a natural transformation  $\triangle A \rightarrow D$ .

$\triangle$  is a functor  $\mathcal{C} \rightarrow [\mathcal{J}, \mathcal{C}]$  and  $\mathbf{Cone}(D)$  is the category  $(\triangle \downarrow D)$  in the notation of (3.3<sup>op</sup>) (the dual case of (3.3)... very creative notation by the lecturer). So to say that every diagram of shape  $J$  in  $\mathcal{C}$  has a limit is equivalent to saying that  $\triangle$  has a right adjoint. (We say  $\mathcal{C}$  has *limits* of shape  $\mathcal{J}$ ).

Dually,  $\mathcal{C}$  has colimits of shape  $J$  iff  $\triangle : \mathcal{C} \rightarrow [\mathcal{J}, \mathcal{C}]$  has a left adjoint.

**Example.** (4.2)

(a) (Lecturer says he'll give a very simple example) Suppose  $\mathcal{J} = \phi$  (a diagram of that here. It's easy to draw, but a bit hard to see). There's a unique diagram of shape  $\mathcal{J}$  in  $\mathcal{C}$ , a cone over it is just an object (with no legs), and a morphism of cones is a morphism of  $\mathcal{C}$  (any one). So a limit for the empty diagram is a terminal object of  $\mathcal{C}$ .

Dually, a colimit for it is an initial object.

(Indeed a very simple example)

—Lecture 10—

(b) Let  $\mathcal{J}$  be the category with two objects and no non-identity morphisms. A diagram of shape  $\mathcal{J}$  is a pair of objects  $A, B$ ; a cone over it is a span



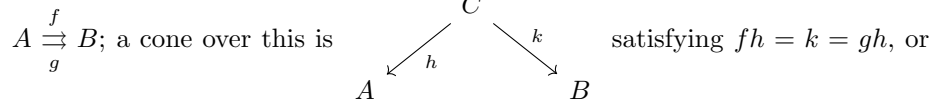
; and a limit for it is a product



as defined in 2.6(e). Dually, a colimit for it is a coproduct

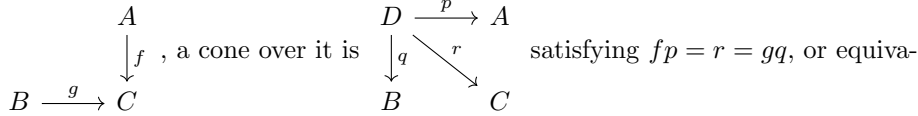
(c) More generally, if  $\mathcal{J}$  is a small discrete category, a diagram of shape  $\mathcal{J}$  is a  $\mathcal{J}$ -indexed family  $(A_j | j \in \mathcal{J})$ , and a limit for it is a product  $(\prod_{j \in \mathcal{J}} A_j \xrightarrow{\pi_j} A_j | j \in \mathcal{J})$  (Dually,  $(A_j \xrightarrow{\nu_j} \sum_{j \in \mathcal{J}} A_j | j \in \mathcal{J})$ , or  $\coprod_{j \in \mathcal{J}} A_j$ , but we usually use the first notation).

(d) Let  $\mathcal{J}$  be the category  $\cdot \xrightarrow[f]{g} B$ . A diagram of shape  $\mathcal{J}$  is a parallel pair



equivalently a morphism  $C \xrightarrow{h} A$  satisfying  $fh = gh$ . A (co)limit for the diagram is a (co)equalizer as defined in 2.6(f).

(e) Let  $\mathcal{J}$  be the category  $\begin{array}{c} \cdot \\ \downarrow \\ \cdot \end{array} \xrightarrow{\quad} \begin{array}{c} \cdot \\ \downarrow \\ \cdot \end{array}$ . A diagram of shape  $\mathcal{J}$  is a cospan



A limit for the diagram is called a *pullback* of  $(f, g)$ . In **Set**, the apex of the pullback is the *fibre product*

$$A \times_C B = \{(x, y) \in A \times B | f(x) = g(y)\}$$

Dually, colimits of shape  $\mathcal{J}^{op}$  are called *pushouts*. Given  $\begin{array}{ccc} A & \xrightarrow{f} & B \\ \downarrow g & & \\ & C & \end{array}$ , we push

$g$  along  $f$  to get the RH side of the colimit square.

(f) (not very important for this course, but might explain why the term *limit* is used) Let  $J$  be the poset of natural numbers. A diagram of shape  $J$  is a *direct system*  $A_0 \xrightarrow{f_0} A_1 \xrightarrow{f_1} A_2 \xrightarrow{f_2} \dots$

A colimit for this is called a *direct limit*: it consists of  $A_\infty$  equipped with morphisms  $A_n \xrightarrow{g_n} A_\infty$  satisfying  $g_n = g_{n+1}$  for all  $n$ , and universal among such. Dually, we have *inverse system* and *inverse limit*.

**Theorem.** (4.3)

- (i) Suppose  $\mathcal{C}$  has equalizers and all finite (respectively, small) products. Then  $\mathcal{C}$  has all finite (respectively, small) limits.
- (ii) Suppose  $\mathcal{C}$  has pullbacks and a terminal object, then  $\mathcal{C}$  has all finite limits.

*Proof.* (i) Suppose given  $D : \mathcal{J} \rightarrow \mathcal{C}$ . Form the products  $P = \prod_{j \in \text{ob } \mathcal{J}} D(j)$  and  $Q = \prod_{\alpha \in \text{mor } \mathcal{J}} D(\text{cod } \alpha)$ .

We have morphisms  $P \xrightarrow[f]{g} Q$  defined by  $\pi_\alpha f = \pi_{\text{cod}(\alpha)}$ ,  $\pi_\alpha g = D(\alpha)\pi_{\text{dom } \alpha}$  for all  $\alpha$ .

Let  $E \xrightarrow{e} P$  be an equalizer of  $(f, g)$ . The composites  $\lambda_j = \pi_j e : E \rightarrow D(j)$  form a cone over  $D$ : given  $\alpha : j \rightarrow j'$  in  $\mathcal{J}$ ,

$$D(\alpha)\lambda_j = D(\alpha)\pi_j e = \pi_{\alpha} g e = \pi_{\alpha} f e = \pi_{j'} e = \lambda_{j'}$$

Given any cone  $(A, (\mu_j | j \in \text{ob } \mathcal{J}))$  over  $D$ , there's a unique  $\mu : A \rightarrow P$  with  $\pi_j \mu = \mu_j$  for each  $j$ , and  $\pi_\alpha f \mu = \mu_{\text{cod } \alpha} = D(\alpha)\mu_{\text{dom } \alpha} = \pi_\alpha g \mu$  for all  $\alpha$ , and hence  $f \mu = g \mu$ . So there is a unique  $\nu : A \rightarrow E$  with  $e \nu = \mu$ . So  $(E, (\lambda_j | j \in \text{ob } \mathcal{J}))$  is a limit cone.

- (ii) It's enough to construct finite products and equalizers. But if 1 is the terminal

object, then a pullback for 
$$\begin{array}{ccc} & A & \\ & \downarrow & \\ B & \longrightarrow & 1 \end{array}$$
 has the universal property of a product  $A \times B$ , and we can form  $\prod_{i=1}^n A_i$  inductively as  $A_1 \times (A_2 \times (A_3 \times \dots (A_{n-1} \times A_n)))$ .

Now, to form the equalizer of  $A \xrightarrow[f]{g} B$ , consider the cospan 
$$\begin{array}{ccc} & A & \\ & \downarrow (1_A, f) & \\ A & \xrightarrow{(1_A, g)} & A \times B \end{array}$$

. A cone over this consists of 
$$\begin{array}{ccc} P & \xrightarrow{h} & A \\ & \downarrow k & \\ & A & \end{array}$$
 satisfying  $(1_A, f)h = (1_A, g)k$ , or

equivalently  $1_A h = 1_A k$ , and  $fh = gk$ , or equivalently, a morphism  $P \xrightarrow{h} A$  satisfying  $fh = gh$  (think). So a pullback for  $(1_A, f)$  and  $(1_A, g)$  is an equalizer of  $(f, g)$ .

We say a category  $\mathcal{C}$  is *complete* if it has all small limits. Dually, *cocomplete* means it has all small colimits.

**Set** is both complete and cocomplete: products are cartesian products, coproducts are disjoint unions.

Similarly, **Gp**, **AbGp**, **Rng**, **Mod<sub>R</sub>**, ... are all complete and cocomplete (nice to know that). **Top** is also complete and cocomplete, ...  $\square$

**Definition.** (4.4)

Let  $F : \mathcal{C} \rightarrow \mathcal{D}$  be a functor.



- (a) We say  $F$  *preserves limits* of shape  $\mathcal{J}$  if, given  $D : \mathcal{J} \rightarrow \mathcal{C}$  and a limit cone  $(L, (\lambda_j | j \in \text{ob } \mathcal{J}))$  in  $\mathcal{C}$ ,  $(FL, (F\lambda_j | j \in \text{ob } \mathcal{J}))$  is a limit for  $FD$ .
- (b) We say  $F$  *reflects limits* of shape  $\mathcal{J}$  if, given  $D : \mathcal{J} \rightarrow \mathcal{C}$  and a cone  $(L, (\lambda_j)_j)$  s.t.  $(FL, (F\lambda_j)_j)$  is a limit for  $FD$ , then  $(L, (\lambda_j)_j)$  is a limit for  $D$ .
- (c) (this is stronger than both of above and implies them. Note that a lot of textbook gets this wrong; the definitions given by them are usually not categorical) We say  $F$  *creates limits* of shape  $\mathcal{J}$  if, given  $D : \mathcal{J} \rightarrow \mathcal{C}$  and (lecturer had *and and* here; if this is any other course I won't bother put this, but this is Category Theory) a limit  $(M, (\mu_j)_j)$  for  $FD$ , there exists a cone  $(L, (\lambda_j)_j)$  over  $D$  whose image under  $F$  is isomorphic to the limit cone, and any such cone is a limit in  $\mathcal{C}$ .

## 5 Example Class 1

Many people wrote too much for questions. Do have the confidence to use the duality principle when it's usable!

### 5.1 Question 1

We do need to verify that  $(AB)_{ik} = \vee_j (a_{ij} \wedge b_{jk})$ , and

$$\begin{aligned} ((AB)C)_{il} &= \vee_k ((\vee_j (a_{ij} \wedge b_{jk})) \wedge c_{kl}) \\ &= \vee_k \vee_j (a_{ij} \wedge b_{jk} \wedge c_{kl}) \\ &= (A(BC))_{il} \text{ by symmetry} \end{aligned}$$

And of course identity matrices are identities.

Define a functor  $F : \mathbf{Mat}_L \rightarrow \mathbf{Rel}_f$  by  $F(n) = \{1, 2, \dots, n\}$ ; if  $A : n \rightarrow p$  is a  $p \times n$  matrix in  $\mathbf{Mat}_L$ , then  $FA = \{(i, j) | a_{ji} = 1\}$ .

Important: we have to verify this is functorial, which many people didn't bother to do. Why is it functorial? Well again we just have to verify explicitly that

$$(AB)_{ik} = 1 \iff (\exists j)(a_{ij} = b_{jk} = 1)$$

so  $F(AB) = FA \circ FB$ . This does require verifications, because you are multiplying the matrices over lattices; for example say if you are doing it for the finite field with 2 elements then it won't work.

Now note that to prove they are equivalent we don't really need to find both the two functors and natural transformations; instead we can use a theorem in chapter 1, that  $F$  is part of an equivalence if it is full, faithful and essentially surjective. So we just have to verify that  $F$  is f.f. and es. Indeed it is, since any finite set is isomorphic to  $F(n)$  for some  $n$ . So by (1.12),  $\mathbf{Mat}_L \simeq \mathbf{Rel}_f$ .

### 5.2 Question 2

Part (i) was easy, but many people had problems on part (ii).

(i) Given  $(A_i | i \in I)$ , define  $\mathcal{C}$  by  $\text{ob } \mathcal{C} = \{(i, a) | i \in I, a \in A_i\}$ , and  $\mathcal{C}((i, a)(j, b)) = \phi$  if  $i \neq j$ , and is  $\{*\}$  if  $i = j$ .

$\mathcal{C}$  is a groupoid; its isomorphic classes of objects are of the form  $\{i\} \times A_i$ ,  $i \in I$ . If we've got a skeleton then we can pick out one from each of these, which is equivalent to  $\text{AC}$ .

(ii) Now take  $\text{ob } \mathcal{C} = I \times \{0, 1\}$ , and morphisms  $(i, m) \rightarrow (i, n)$  are formal finite sums as given in the hint, and of course composition is just addition. Again this is a groupoid because every morphism can be inverted by just reverting the sign of every coefficient. So  $\mathcal{C}$  has isomorphic classes  $\{i\} \times \{0, 1\}$ ,  $i \in I$ . So this has a skeleton, say we take  $\mathcal{C}_0$  to be the full subcategory on objects  $I \times \{0\}$ .

But then by assumption we have an equivalence  $\mathcal{C}_0 \xrightleftharpoons[F]{G} \mathcal{C}$ , then  $FG(i, 0) = FG(i, 1)$  for all  $i$ . So if we have a natural transformation  $\beta : FG \rightarrow 1_{\mathcal{C}}$  which is also an

isomorphism, then either  $\beta_{(i,0)}$  or  $\beta_{(i,1)}$  is a non-zero formal finite sum. So we just put  $A_i = \{x \in A_i | x \text{ occurs in either } \beta_{(i,0)} \text{ or } \beta_{(i,1)} \text{ with non-zero coefficient}\}$ .

### 5.3 Question 3

(i) Quite a lot of people forgot to verify that it is actually a subgroup! Suppose given automorphisms  $F, G, H$  with isomorphisms  $\alpha : F \rightarrow 1_{\mathcal{C}}, \beta : G \rightarrow 1_{\mathcal{C}}$ . Then  $(F^{-1}\alpha)^{-1} : F^{-1} \rightarrow F^{-1}F = 1_{\mathcal{C}}$  is an isomorphism, so  $F^{-1}$  is inner. Now  $FG \xrightarrow{F\beta} F \xrightarrow{\alpha} 1_{\mathcal{C}}$  is isomorphism, so  $FG$  is inner.

Now we have to verify normality:  $HFH^{-1} \xrightarrow{H\alpha_H^{-1}} HH^{-1} = 1_{\mathcal{C}}$  is iso, so  $HFH^{-1}$  is inner.

You don't have to spend a lot of time to verify that all these are nat transforms (ok).

(ii) Note that an isomorphism is in particular an equivalence. Now if  $F$  is an automorphism, it is full and faithful, so for any  $A$ , morphism  $A \rightarrow F1$  are in bijection with morphisms  $F^{-1}A \rightarrow 1$ , so there's just one of them.

Hence if  $F$  is an isomorphism of **Set**,  $F(1) = 1$ , and hence there's a unique  $\alpha : \mathbf{Set}(1, -) \rightarrow F$  (this is just Yoneda).

$$\begin{array}{ccc} & 1 & \xrightarrow{\alpha_1} F1 \\ & \downarrow x & \downarrow Fx \\ \text{We need to show } \alpha \text{ is iso: but for any } A, \text{ and any } 1 & \xrightarrow{x} A, & \\ & A & \xrightarrow{\alpha_A} FA \end{array}$$

commutes, but  $F$  is full and faithful, so  $\alpha_A$  is bijective. Hence by (1.8),  $\alpha$  is an isomorphism.

(iii) Note that if  $X$  has  $\geq 3$  points, then it has  $\geq 4$  continuous endo maps (constants and identity). If  $X$  has  $\leq 1$  point, then its only endo is  $1_X$ . So the only possibility is  $X$  having 2 points. Say  $X = \{0, 1\}$  with discrete or indiscrete topology, all 4 maps  $X \rightarrow X$  are continuous. So the only possible topology is the Sierpinski space given in the question (or the other way), which there are 3 continuous maps  $X \rightarrow X$ .

Hence if  $F$  is an autom of  $\mathcal{C} \subseteq \mathbf{Top}$ , we must have  $FS \cong S$ .

We also have  $F1 \cong 1$ , and  $U : \mathcal{C} \rightarrow \mathbf{Set}$  is iso to  $\mathcal{C}(1, -)$ . So there's a unique nat  $\alpha : U \rightarrow UF$ , and  $\alpha_X$  is bijective for all  $X$ , as before. Now we don't know if  $\alpha_X$  is continuous or not for a given  $X$ . So we consider naturality squares

$$\begin{array}{ccc} UX & \xrightarrow{\alpha_X} & UFX \\ \downarrow Uf & & \downarrow UFF \\ US & \xrightarrow{\alpha_S} & UFS \end{array} \quad . \text{ If } \alpha_S \text{ is discontinuous, then for any } X, \alpha_X \text{ maps open}$$

subsets of  $X$  bijectively to closed subsets of  $FX$ , but this is impossible if not every intersection of open sets in  $X$  is open.

So then  $\alpha_S$  is a homeomorphism, and  $\alpha_X$  is a homeomorphism for all  $X$ .

(iv) This basically says that if we restrict to finite topological spaces, then we do have the other case in the above happening. Let  $F : \mathbf{Top}_f \rightarrow \mathbf{Top}_f$  sending  $X$  to the same set with closed sets as new opens.

Then  $FF = 1_{\mathbf{Top}_f}$ , so  $F$  is an automorphism, and it is not isomorphic to the identity as there exists finite spaces  $X$  with  $X \not\cong FX$  (we could find one with 3 points – lecturer didn't give the explicit example). So  $F$  is not inner.

Now if  $G$  is any non-inner autom, then  $GF$  is inner(????); so  $|\text{Aut}(\mathbf{Top}_f)| :$

$$|Inn(Top_f)| = 2.$$

### 5.4 Question 4

$$(i) \text{ Suppose we are given } \begin{array}{ccc} C & \xrightarrow{h} & A \\ \downarrow g & & \downarrow f \\ D & \xrightarrow{k} & B \\ & \downarrow p & \downarrow q \\ & E & \end{array} \text{ where } f \text{ is an equalizer of } (p, q).$$

Then  $pfh = pkg = qkg = qfh$ , and  $g$  is epic, so  $pk = qk$ , so exists unique  $t$  with  $ft = k$ . Then  $ftg = kg = fh$ , and  $f$  is monic, so  $tg = h$ .

$$(ii) \begin{array}{ccccc} A & \xrightarrow{f} & B & \xleftarrow{l} & D \\ & \searrow k & \downarrow g & \downarrow h & \swarrow m \\ & & C & & \end{array} \quad \text{Note that } f \text{ isn't regular monic. Why not? Because}$$

it is not iso and not an equalizer of  $(g, h)$ , since  $l$  doesn't factor through it. It is trivially monic and strong monic: the only squares with  $f$  on RHS is to put the identity before  $f$ , but it's trivial to verify those cases.

In some sense we can see that this is a minimal counter-example to the statement (think a bit, there's not anything better you can do).

(iii) This actually has nothing to do with the previous two parts. We have

$$\begin{array}{ccc} A & \xrightarrow{f} & B \\ & \searrow k & \downarrow g \\ & & C \end{array} \quad \text{We want a commutative square } \begin{array}{ccc} \cdot & \longrightarrow & \cdot \\ \downarrow & & \downarrow \\ \cdot & \longrightarrow & \cdot \end{array} \text{ with one vertical edge}$$

$f$  and the other one not. There are still a lot of possibilities, but almost any of

$$\text{it work. Say we pick } \begin{array}{ccc} \cdot & \xrightarrow{f} & \cdot \\ \downarrow f & & \downarrow g \\ \cdot & \xrightarrow{h} & \cdot \end{array} \quad \text{Try } f, f, g, h: \text{ the only possible composites are}$$

$$\text{with } \begin{array}{ccc} \cdot & \xrightarrow{h} & \cdot \\ \downarrow 1_B & & \downarrow 1_C \\ \cdot & \xrightarrow{h} & \cdot \end{array} \quad \text{or} \quad \begin{array}{ccc} \cdot & \xrightarrow{h} & \cdot \\ \downarrow h & & \downarrow 1_C \\ \cdot & \xrightarrow{1_C} & \cdot \end{array} \quad \text{So } (f, g) \text{ is vacuously epic.}$$

### 5.5 Question 5

This question has a lots of boring parts so lecturer is not going to write out all of it. The only a little bit tricky part is the strong part of (ii). We'll do

$$\text{that: suppose } gf \text{ is strong monic, and suppose we are given } \begin{array}{ccc} \cdot & \xrightarrow{l} & \cdot \\ \downarrow h & & \downarrow f \\ \cdot & \xrightarrow{k} & \cdot \\ & \downarrow g & \\ & \cdot & \end{array} \quad \text{Then}$$

$$\begin{array}{ccc} \cdot & \xrightarrow{l} & \cdot \\ \downarrow k & \nearrow t & \downarrow gf \\ \cdot & \xrightarrow{gk} & \cdot \end{array}$$

commutes. A lot of people used  $g$  is monic, but we weren't given it here(oops)! Here  $\exists t$  with  $th = l$  (and  $gft = gk$ ). Then  $fth = fl = kh$  and  $h$  is epic, so  $ft = k$ .

Suppose  $gf$  is regular monic, say it's the equalizer of  $\cdot \xrightarrow[k]{l} \cdot$ . To show that  $f$  is an equalizer of  $(kg, fg)$ , suppose we are given  $\cdot \xrightarrow{m} \cdot$  with  $kgm = lgm$ . Then  $\exists n$  with  $gm = gfn$ . But now  $g$  is monic. So  $m = fn$ .

(iv) This part is also fairly problematic. The first thing we need to work out is what equalizers look like in this category. Given  $A \xrightarrow[g]{f} B$  in  $\mathcal{C}$ , their equalizer in  $\mathbf{AbGp}$  is the subgroup of  $\{a \in A \mid f(a) = g(a)\}$ , and this belongs to  $\mathcal{C}$ . Since  $\mathcal{C}$  is full, it's also an equalizer in  $\mathcal{C}$ . Now  $\mathbf{Z} \xrightarrow{\times 2} \mathbf{Z}$  is an equalizer of  $\mathbf{Z} \xrightarrow[q]{f} \mathbf{Z}/2\mathbf{Z}$ ,

so it's regular monic in  $\mathcal{C}$ . But if  $\mathbf{Z} \xrightarrow{\times 4} \mathbf{Z}$  were an equalizer of  $\mathbf{Z} \xrightarrow[g]{f} A$ , then  $f(1) - g(1)$  must have order 4 in  $A$ , so  $A \notin \text{ob } \mathcal{C}$ .

The last part of this is to find a counter-example to a previous part. We consider  $\mathbf{Z} \xrightarrow{f} \mathbf{Z} \oplus \mathbf{Z}/2\mathbf{Z} \xrightarrow{g} \mathbf{Z}$ , where  $f(n) = (2n, [n])$ ,  $g(p, q) = p$ .

Note that  $gf$  is regular monic, but  $f$  isn't, as  $(1, [1])$  has order 4 modulo  $\text{Im}(f)$ .

## 5.6 Question 6

(i) Suppose  $e = fg$ ,  $gf$  an identity. We claim that  $e$  is an equalizer of  $e$  and  $1_{\text{dom } e}$ : we have  $ef = fgf = f$ , so  $f$  has equal composites with  $e$  and  $1_{\text{dom } e}$ ; now if  $h$  satisfies  $eh = h$ , then  $h = fgh$ , so it factorizes through  $h$ ; moreover this factorization is unique, as  $f$  is a split monomorphism.

Conversely, if  $f$  is an equalizer of  $(e, 1_{\text{dom } e})$ , then of course  $e$  must factor through it since  $ee = e$ . Say  $e = fg$ . Now  $fgf = ef = f$ , and  $f$  is monic, so  $gf = 1_{\text{dom } f}$ .

(ii)  $\mathcal{E} \subseteq \text{Idem } \mathcal{C}$ , morphisms  $e \rightarrow d$  are morphisms  $\text{dom } e \xrightarrow{f} \text{dom } d$  in  $\mathcal{C}$  with  $dfe = f$ . Note that this is equivalent to the two separate equations: one way is clear, now  $df = d(dfe) = dfe = f$  (remember  $d$  is idempotent!!!). Similarly  $fe = f$ .

Composition in  $\mathcal{C}[\check{\mathcal{E}}]$  is composition in  $\mathcal{C}$ : if  $e \xrightarrow{f} d \xrightarrow{g} c$ , then  $cgf = gf = gfe$ , so  $gf : e \rightarrow c$ . The identity on  $e$  is  $e \xrightarrow{e} e$ .

(iii) We define  $I : \mathcal{C} \rightarrow \mathcal{C}[\check{\mathcal{E}}]$  by  $IA = 1_A$ ,  $If = f$  (check this works).  $I$  is f and f since all morphisms  $A \rightarrow B$  in  $\mathcal{C}$  are morphisms  $1_A \rightarrow 1_B$  in  $\mathcal{C}[\check{\mathcal{E}}]$ .

For every  $A \xrightarrow{e} A$  in  $\mathcal{E}$ ,  $Ie$  splits as  $1_A \xrightarrow{e} e \xrightarrow{e} 1_A$ .

So if  $T = \hat{T}I$ ,  $T$  must send idempotents in  $\mathcal{E}$  to split idempotents.

Now we have to show the converse, where we do have to use choice here. Suppose  $Te$  is split for every  $A \xrightarrow{e} A$  in  $\mathcal{E}$ . Choose a splitting  $TA \xrightarrow{g_e} \hat{T}e \xrightarrow{f_e} TA$  of it, and define  $\hat{T} : \mathcal{C}[\check{\mathcal{E}}] \rightarrow \mathcal{D}$  on morphisms by  $\hat{T}(e \xrightarrow{h} d) = \hat{T}e \xrightarrow{f_e} TA \xrightarrow{Th} TB \xrightarrow{g_d} \hat{T}d$  (verify that this is functorial): provided we split  $T(1_A)$  as  $TA \xrightarrow{1_{TA}} TA \xrightarrow{1_{TA}} TA$ , we have  $\hat{T}I = T$ .

(iv) Now suppose  $\mathcal{E} = \{\text{all idempotents of } \mathcal{C}\}$ . If  $e \xrightarrow{d} e$  is idempotent in  $\mathcal{C}[\check{\mathcal{E}}]$ , then  $dd = d$  in  $\mathcal{C}$ , so  $d \in \mathcal{E}$ , and  $e \xrightarrow{d} e$  splits as  $e \xrightarrow{d} d \xrightarrow{d} e$ .

(v) If  $\mathcal{D}$  is Cauchy-complete, consider the functor  $[\hat{\mathcal{C}}, \mathcal{D}] \xrightarrow{\Phi} [\mathcal{C}, \mathcal{D}]$  sending  $\hat{T}$  to  $\hat{T}I$  and  $\alpha \text{ to } \alpha_I$ .  $\Phi$  is surjective on objects by (iii); so we need to show, given  $S, T : \hat{\mathcal{C}} \rightrightarrows \mathcal{D}$ , any nat trans  $\alpha : SI \rightarrow TI$  extends uniquely to a nat trans  $S \rightarrow T$ .

Given  $A \xrightarrow{e} A$  in  $\mathcal{E}$ , we have a morphism  $Se \xrightarrow{S(e \xrightarrow{e} 1_A)} SA \xrightarrow{\alpha_A} TA \xrightarrow{T(1_A \xrightarrow{e} e)} Te$  which we take to be  $\alpha_e$ .

This is the only possibility that makes the naturality squares for both  $e \xrightarrow{e} 1_A$

$$\text{and } 1_A \xrightarrow{e} e \text{ commute: } \begin{array}{ccc} Se & \xrightarrow{S(e \xrightarrow{e} 1)} & SA \\ \downarrow \alpha_e & & \downarrow \alpha_A \\ Te & \xrightarrow{T(e \xrightarrow{e} 1)} & TA \end{array} \text{ commutes since } \alpha_A \text{ is natural w.r.t.}$$

$1_A \xrightarrow{e} 1_A$ . So this is the only possible way to extend it to a nat trans, and we of course have to verify naturality w.r.t. any  $e \xrightarrow{f} d$  in  $\mathcal{C}[\check{\mathcal{E}}]$ .

So  $\Phi$  is part of an equivalence by (1.12).

## 5.7 Question 7

For any  $F : \mathcal{C} \rightarrow \mathbf{Set}$ ,  $\coprod_{(A,x), A \in \text{ob } \mathcal{C}, x \in FA} \mathcal{C}(A, -) \rightarrow F$  is pointwise surjective, so  $F$  irreducible implies that there exists  $\mathcal{C}(A, -) \twoheadrightarrow F$ .

Conversely, given  $\mathcal{C}(A, -) \xrightarrow{\alpha} F$  and an epi  $\coprod_{i \in I} G_i \xrightarrow{f} F$ , we have

$$\begin{array}{ccc} & & \mathcal{C}(A, -) \\ & \swarrow \gamma & \downarrow \alpha \\ \coprod_{i \in I} G_i & \xrightarrow{\beta} & F \end{array}$$

$\gamma$  corresponds to an element of  $\coprod_{i \in I} G_i A$ , which lives in  $G_i A$  for some  $i$ . Then

$$\begin{array}{ccc} & \mathcal{C}(A, -) & \\ \swarrow \gamma & \downarrow \alpha & \\ G_i & \xrightarrow{\beta_i} & F \end{array} \text{ forces } \beta_i \text{ to be epic.}$$

(ii) If  $F$  is irreducible and projective, then we get

$$\begin{array}{ccc} & F & \\ \swarrow f & \downarrow 1 & \\ \mathcal{C}(A, -) & \xrightarrow{\alpha} & F \end{array}, \text{ so } \alpha \text{ is}$$

split epic.

Conversely, if  $\mathcal{C}(A, -) \xleftarrow[\alpha]{\beta} F$  is split epic, then given

$$\begin{array}{ccc} & F & \\ & \downarrow f & \\ & \mathcal{C}(A, -) & \\ & \downarrow \alpha & \\ G & \longrightarrow & H \end{array} \text{ so}$$

$F$  is projective.

The composite  $\mathcal{C}(A, -) \xrightarrow{\alpha} F \xrightarrow{\beta} \mathcal{C}(A, -)$  is idempotent;  $Y : \mathcal{C}^{op} \rightarrow [\mathcal{C}, \mathbf{Set}]$  is full and faithful, so it's of the form  $Y(e)$  for a unique idempotent  $A \xrightarrow{e} A$  in  $\mathcal{C}$ .

But now if this splits as  $A \xrightarrow{g} B \xrightarrow{f} A$ , then  $\mathcal{C}(A, -) \xrightarrow{\mathcal{C}(f, -)} \mathcal{C}(B, -) \xrightarrow{\mathcal{C}(g, -)} \mathcal{C}(A, -)$  is a splitting of  $\beta\alpha$ . But then by Q6(i), so we must have  $F \cong \mathcal{C}(B, -)$ .

(iii) We know  $[\mathcal{C}, \mathbf{Set}] \simeq [\hat{\mathcal{C}}, \mathbf{Set}]$  since  $\mathbf{Set}$  is Cauchy-complete. If  $\hat{\mathcal{C}} \simeq \hat{\mathcal{D}}$  by functors  $F$  and  $G$ , then  $T \rightarrow TF$  and  $T \rightarrow TG$  give an equivalence  $[\hat{\mathcal{C}}, \mathbf{Set}] \simeq [\hat{\mathcal{D}}, \mathbf{Set}]$ . But any equivalence  $[\hat{\mathcal{C}}, \mathbf{Set}] \simeq [\hat{\mathcal{D}}, \mathbf{Set}]$  restricts to an equivalence between the full subcategories of irreducible projectives, which are equivalent to  $\hat{\mathcal{C}}^{op}$  and  $\hat{\mathcal{D}}^{op}$ .

## 5.8 Question 8

This question is actually quite quick.

$\mathcal{C}(A, -)$  is a monofunctor just means that for any  $f : B \rightarrow C$ , the map  $g \rightarrow fg$  is an injection  $\mathcal{C}(A, B) \rightarrow \mathcal{C}(A, C)$ . This holds for all  $A$  iff all  $f \in \text{mor } \mathcal{C}$  are monic – that's the equivalence between (i) and (ii).

We now prove (ii)  $\implies$  (iii): Since we have an epi  $\coprod \mathcal{C}(A, -) \twoheadrightarrow F$  and disjoint unions of monofunctors are monofunctors (?).

For (iii)  $\implies$  (ii), if we have  $F \xrightarrow{\alpha} \mathcal{C}(A, -)$  with  $F$  a monofunctor, we have a splitting  $\mathcal{C}(A, -) \xrightarrow{\beta} F$ , and any subfunctor of a monofunctor is a monofunctor.

Given  $f : A \rightarrow B$ , consider the push out

$$\begin{array}{ccc} \mathcal{C}(B, -) & \xrightarrow{\mathcal{C}(f, -)} & \mathcal{C}(A, -) \\ \downarrow \mathcal{C}(f, -) & & \downarrow \\ \mathcal{C}(A, -) & \longrightarrow & F \end{array} \quad . \text{ Explicitly,}$$

$F(C) \cong \mathcal{C}(A, C) \times \{0, 1\} / \sim$  where  $(g, 0) \sim (g, 1) \iff f$  factors through  $g$ .

If this is a monofunctor, we must have  $(1_A, 0) \simeq (1_A, 1)$  (check the relation in the middle), since  $Ff$  sends them to the same thing.

If all morphisms of  $\mathcal{C}$  are split monic, then  $\mathcal{C}$  is a groupoid. And of course the converse holds (check).