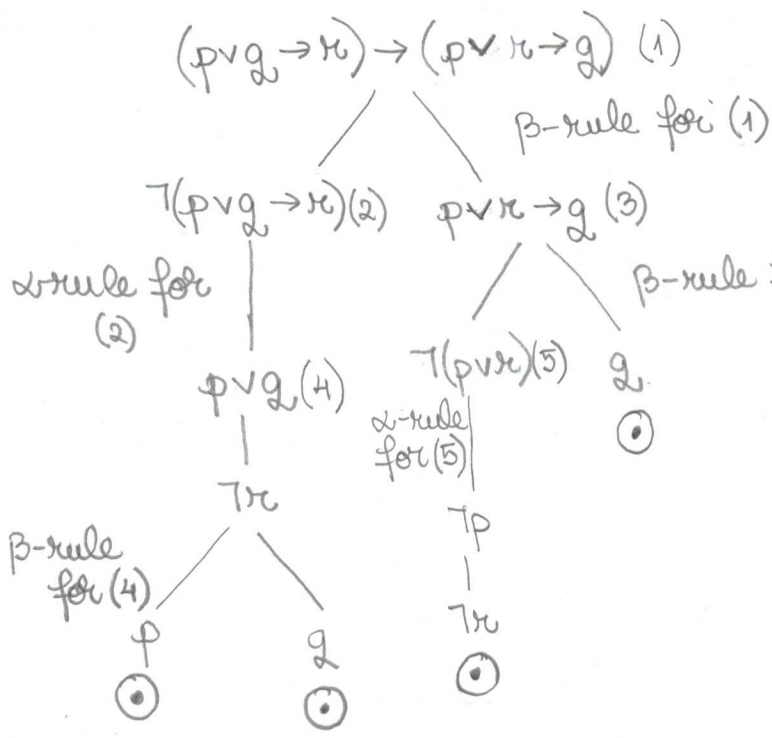


EX1-912

1.2 $U_2 = (p \vee q \rightarrow r) \rightarrow (p \vee r \rightarrow q)$

Use ST method to decide what kind of formula is U_2 (consistent, inconsistent, valid). If consistent, find all its models.



THEORETICAL RESULTS

- 1) consistent formula: has at least 1 model
 $\exists i \in \{p_1, \dots, p_m\} \rightarrow \{T, F\} \text{ s.t. } i(u) = T$
- 2) inconsistent:
- 3) valid:
- 1*) consistent formula has associated a complete and open ST
- 2*) ... 3*) ...

$$\text{DNF}(U) = (\neg r \wedge p) \vee (\neg r \wedge q) \vee (\neg p \wedge \neg r) \vee q$$

ABSORPTION

$$\equiv (\neg r \wedge p) \vee (\neg p \wedge \neg r) \vee q$$

ABSORPTION

$$U \vee (U \wedge V) \equiv U$$

$$U \wedge (U \vee V) \equiv U$$

DNF

ST is disjunction of all its branches; graphical representation of Disjunctive Normal Form (DNF)

We obtain the models of the formula from the open branches.

Cube: $\neg r \wedge p = T$

\downarrow	\downarrow
T	T

$i_{1,2}: \{p, q, r\} \rightarrow \{T, F\}$

$i_1(p) = T$	$i_1(r) = T$
$i_2(q) = T$	$i_2(r) = F$
$i_1(r) = F$	$i_2(r) = F$

Cube: $\neg p \wedge \neg r = T$

\downarrow	\downarrow
T	T

$i_{3,4}: \{p, q, r\} \rightarrow \{T, F\}$

$i_3(p) = F$	$i_4(p) = F$
$i_3(q) = T$	$i_4(q) = F$
$i_4(r) = F$	$i_4(r) = F$

Cube: q

$i_{5,6}: \{p, q, r\} \rightarrow \{T, F\}$

$i_5(p) = T$	$i_6(p) = F$
$i_5(q) = T$	$i_6(q) = T$
$i_5(r) = T$	$i_6(r) = T$

The formula U_2 has 6 distinct models: $i_1, \dots, i_6 \Rightarrow U_2$ consistent, but not a tautology (a tautology would have 8 models) $\Rightarrow U_2$ contingent.

THEORETICAL RESULTS

$$T = (U_1 \wedge U_2) \vee (U_1 \wedge \neg U_2) \vee (\neg U_1 \wedge U_2) \vee (\neg U_1 \wedge \neg U_2)$$

DEFINITION

$$U \vee (V \wedge W) \equiv U \vee (V \wedge W)$$

$$U \wedge (V \vee W) \equiv U \wedge (V \vee W)$$

THEOREM

$$T = (U_1 \wedge U_2) \vee (U_1 \wedge \neg U_2) \vee (\neg U_1 \wedge U_2) \vee (\neg U_1 \wedge \neg U_2)$$

$$T = (U_1 \wedge U_2) \vee (U_1 \wedge \neg U_2) \vee (\neg U_1 \wedge U_2) \vee (\neg U_1 \wedge \neg U_2)$$

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$$T = (U_1 \wedge U_2) \vee (U_1 \wedge \neg U_2) \vee (\neg U_1 \wedge U_2) \vee (\neg U_1 \wedge \neg U_2)$$

4.3-913

$$U_3 = p \rightarrow (q \wedge r) \vee q \wedge \neg p$$

$$\neg U_3 = \neg(p \rightarrow (q \wedge r) \vee q \wedge \neg p)$$

THEORETICAL RESULTS

α, β rules

We obtain the anti-models of U_3 by assessing the open branches of the semantic table associated to $\neg U_3$. (T1)

⊙ = open branch

IDEMPOTENCE:

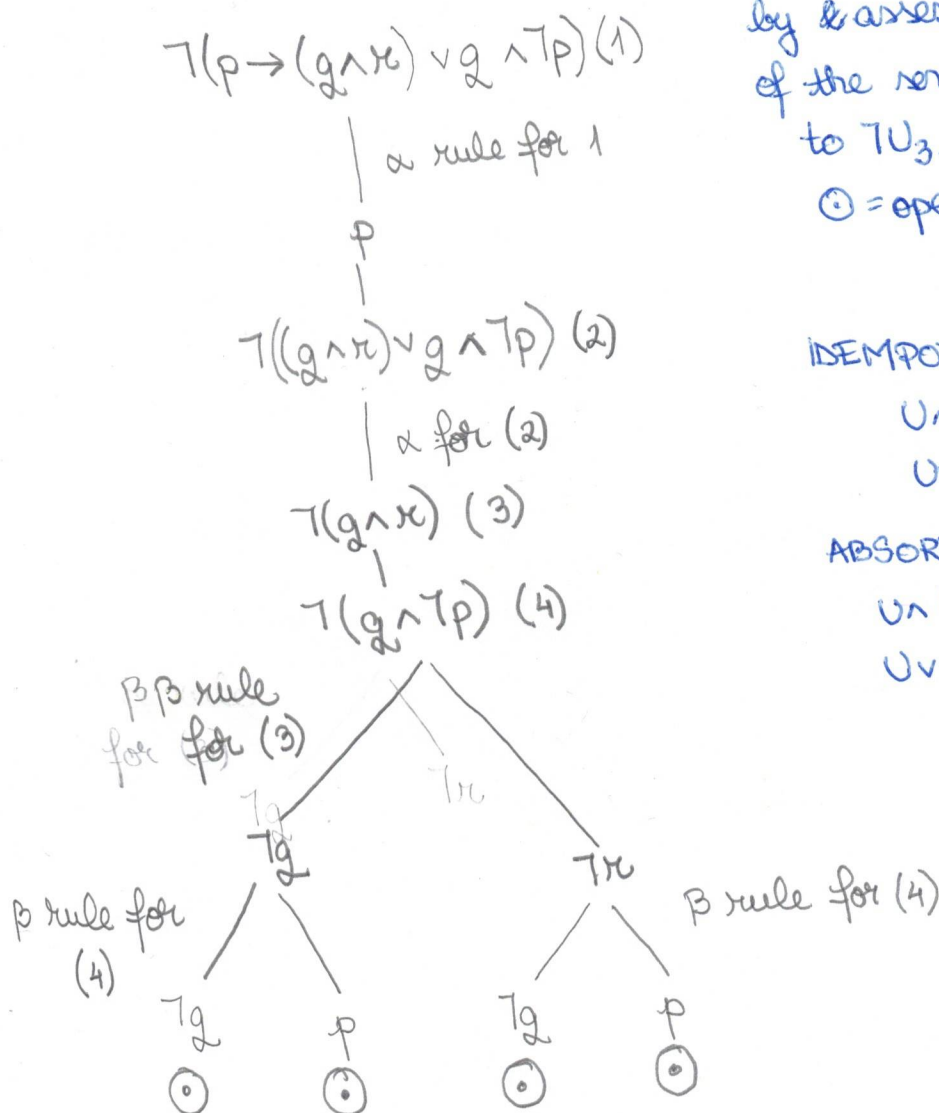
$$U \wedge U \equiv U$$

$$U \vee U \equiv U$$

ABSORPTION

$$U \wedge (U \vee V) \equiv U$$

$$U \vee (U \wedge V) \equiv U$$



$$\begin{aligned} \text{DNF}(\neg U_3) &= (p \wedge \neg q \wedge \neg r) \vee (p \wedge \neg q \wedge p) \vee (p \wedge \neg r \wedge \neg q) \vee (p \wedge \neg r \wedge p) \\ &\equiv (p \wedge \neg q) \vee (p \wedge \neg q) \vee (p \wedge \neg q \wedge \neg r) \vee (p \wedge \neg r) \\ &\equiv (p \wedge \neg q) \vee (p \wedge \neg q \wedge \neg r) \vee (p \wedge \neg r) \\ &\equiv (p \wedge \neg q) \vee (p \wedge \neg r) \end{aligned}$$

Rule 1: $p \wedge \neg q = T$

1	1
T	T

$i_1, i_2: \{p, q, r\} \rightarrow \{T, F\}$

$i_1(p) = T$	$i_2(p) = T$
$i_1(q) = F$	$i_2(q) = F$
$i_1(r) = T$	$i_2(r) = F$

Cube 2: $p \wedge \neg x = T$

\downarrow	\downarrow
T	T

$$i_3, i_4 : \{p, q, x\} \rightarrow \{T, F\}$$

$$i_3(p) = T \quad i_4(p) = T$$

$$i_3(q) = T \quad i_4(q) = F$$

$$i_3(x) = F \quad i_4(x) = F$$

$$i_4 = i_2$$

The models of U_3 are $i_1, i_2, i_3 \stackrel{T_1}{\Rightarrow}$ the i_1, i_2, i_3 anti-models for U_3 .

ST5-

Check whether conclusion C is a logical consequence of the set of hypotheses using the ST. method.

H_1 : All hummingbirds are richly colored.

H_2 : No large birds live on honey.

H_3 : Birds that do not live on honey are dull in color.

C: All hummingbirds are small.

$H_1, H_2, H_3 \vdash C$?

$H_1: (\forall x)(hb(x) \rightarrow rc(x))$

$H_2: \neg(\exists x)(\neg sb(x) \wedge lh(x))$
 $\equiv (\forall x)(\neg sb(x) \rightarrow \neg lh(x))$

$H_3: (\forall x)(\neg lh(x) \rightarrow \neg rc(x))$

C: $(\forall x)(hb(x) \rightarrow sb(x))$

THEORETICAL RESULTS

Let U_1, \dots, U_m, V be predicate formulas. $U_1, U_2, \dots, U_m \models V$ iff there is a closed semantic table associated with formula $U_1 \wedge U_2 \wedge \dots \wedge U_m \wedge \neg V$. (T1)

$(\forall x)(hb(x) \rightarrow rc(x)) \wedge (\forall x)(\neg sb(x) \rightarrow \neg lh(x)) \wedge (\forall x)(\neg lh(x) \rightarrow \neg rc(x))$

$\wedge \neg(\forall x)(hb(x) \rightarrow sb(x))$ (1)

| α for (1)

$(\forall x)(hb(x) \rightarrow rc(x))$ (2)

|

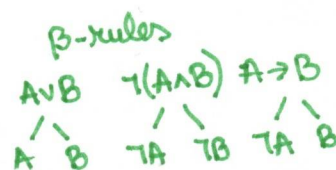
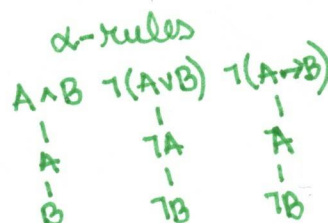
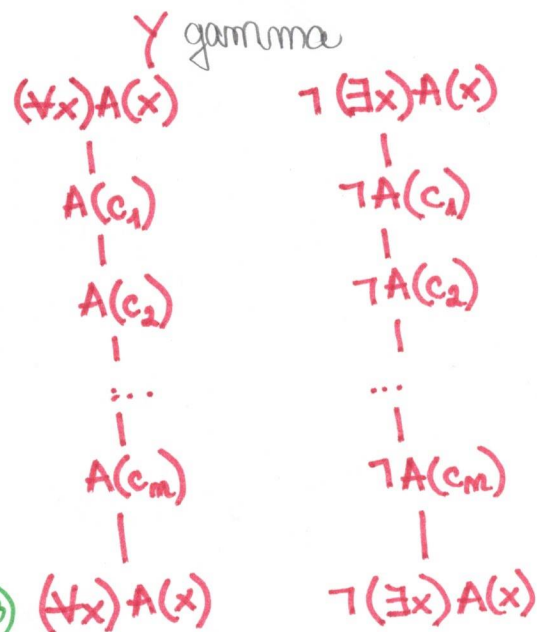
$(\forall x)(\neg sb(x) \rightarrow \neg lh(x))$ (3)

|

$(\forall x)(\neg lh(x) \rightarrow \neg rc(x))$ (4)

|

$\neg((\forall x)(hb(x) \rightarrow sb(x)))$ (5)



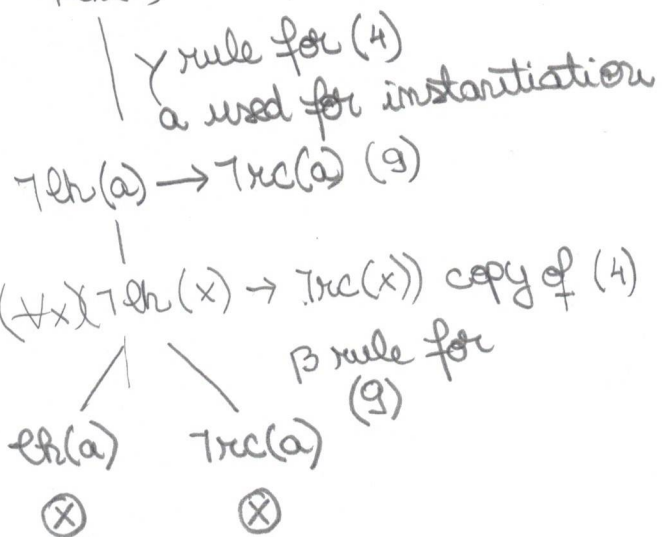
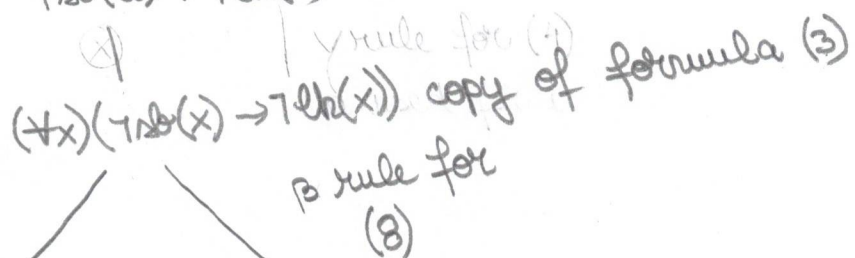
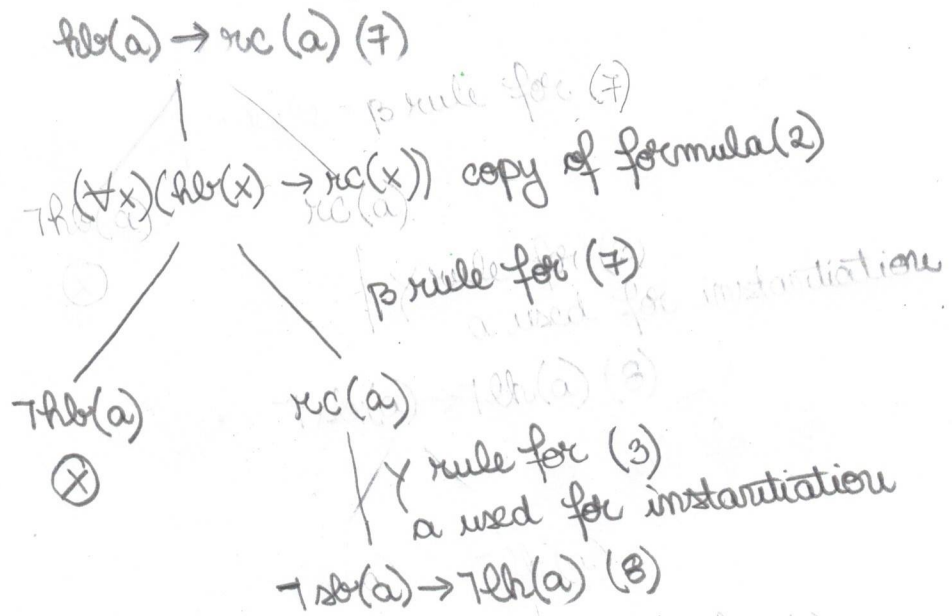
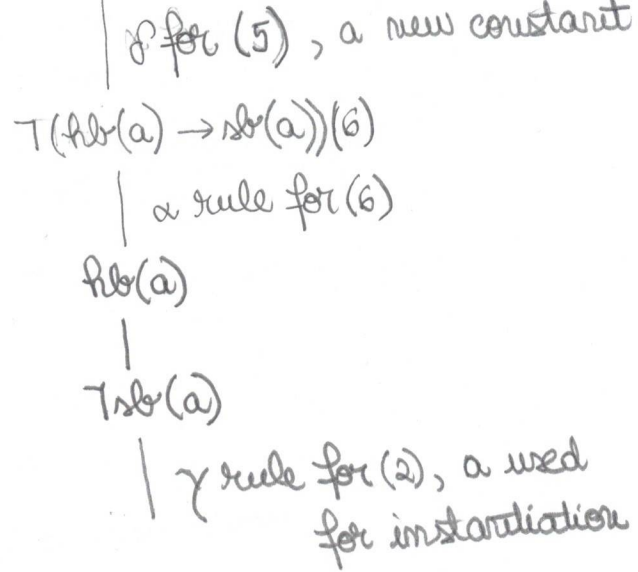
D = universe of birds

$hb: D \rightarrow \{T, F\}$, $hb(x)$ = "x hummingbird"

$rc: D \rightarrow \{T, F\}$, $rc(x)$ = "x richly colored"

$sb: D \rightarrow \{T, F\}$, $sb(x)$ = "x small bird"

$lh: D \rightarrow \{T, F\}$, $lh(x)$ = "x lives on honey"



All 4 branches of the semantic table are closed (they contain pairs $(\text{fb}(a), \neg \text{fb}(a))$, $(\neg \text{sb}(a), \text{sb}(a))$, $(\neg \text{lh}(a), \text{lh}(a))$ and $(\text{rc}(a), \neg \text{rc}(a))$) $\Rightarrow H_1$ semantic table associated with formula $H_1 \wedge H_2 \wedge H_3 \wedge \neg C$ is closed $\Rightarrow H_1, H_2, H_3 \vdash C$.

ST 7.2 Using ST method, prove the following properties

$$\models (\forall x)A(x) \vee (\forall x)B(x) \rightarrow (\forall x)(A(x) \vee B(x)) \quad U_1$$

$$\not\models (\forall x)(A(x) \vee B(x)) \rightarrow ((\forall x)A(x) \rightarrow (\forall x)B(x)) \quad U_2$$

" \forall is semi-distributive over \vee "

$$TU_1 = \neg((\forall x)A(x) \vee (\forall x)B(x) \rightarrow (\forall x)(A(x) \vee B(x))) \quad (1)$$

α rule for (1)

$$(\forall x)A(x) \vee (\forall x)B(x) \quad (2)$$

$$\neg(\forall x)(A(x) \vee B(x)) \quad (3)$$

δ for (3), a new const.

$$\neg(A(a) \vee B(a)) \quad (4)$$

α rule for (4)

$$\neg A(a)$$

$$\neg B(a)$$

β rule for (2)

$$(\forall x)A(x) \quad (5)$$

$$(\forall x)B(x) \quad (6)$$

γ rule for (5)
a used for inst

$$A(a)$$

γ rule for (6)
a used for inst

$$B(a)$$

$$(\forall x)A(x) \text{ copy of } (5) \quad \textcircled{\times}$$

$$(\forall x)B(x) \text{ copy of } (6) \quad \textcircled{\times}$$

THEORETICAL RESULTS

(SOUNDNESS AND COMPLETENESS OF SEMANTIC TABLEAUX METHOD)

A propositional/predicate formula U is a tautology iff there is a closed semantic tableaux associated to TU . (T_1)

α -rules

$$\begin{array}{c} A \wedge B \\ | \\ A \\ | \\ B \end{array}$$

$$\begin{array}{c} \neg(A \vee B) \\ | \\ \neg A \\ | \\ \neg B \end{array}$$

$$\begin{array}{c} \neg(A \rightarrow B) \\ | \\ A \\ | \\ \neg B \end{array}$$

β -rules

$$\begin{array}{cc} A \vee B & \\ / & \backslash \\ A & B \end{array}$$

$$\begin{array}{cc} \neg(A \wedge B) & \\ / & \backslash \\ \neg A & \neg B \end{array}$$

$$\begin{array}{cc} A \rightarrow B & \\ / & \backslash \\ \neg A & B \end{array}$$

δ -rules

$$\begin{array}{c} (\exists x)A(x) \\ | \\ A(c) \end{array}$$

$$\begin{array}{c} \neg(\forall x)A(x) \\ | \\ \neg A(c) \end{array}$$

γ -rules

$$\begin{array}{c} (\forall x)A(x) \\ | \\ A(c_1) \\ \vdots \\ A(c_m) \end{array}$$

$$\begin{array}{c} \neg(\exists x)A(x) \\ | \\ \neg A(c_1) \\ \vdots \\ \neg A(c_m) \end{array}$$

$$(\forall x)A(x) \quad \neg(\exists x)A(x)$$

$$c_1, \dots, c_m$$

$$\text{instantiated by } c_1, \dots, c_m$$

$$\text{inconsistent branch}$$

$\textcircled{\times} \rightarrow$ closed branch (contains formula and its negation)

\rightarrow the tableaux has only closed branches \Rightarrow (2 closed branches containing pairs $(\neg A(a), A(a))$ and $(\neg B(a), B(a)) \Rightarrow TU_1$ has no models $\Rightarrow TU_1$ inconsistent

$$T_1 \Rightarrow \models U_1$$

$$\neg U_2 = \neg((\forall x)(A(x) \vee B(x)) \rightarrow ((\forall x)A(x) \vee (\forall x)B(x))) \quad (1)$$

| α rule for (1)

$$(\forall x)(A(x) \vee B(x)) \quad (2)$$

|

$$\neg((\forall x)A(x) \vee (\forall x)B(x)) \quad (3)$$

| α rule for (3)

$$\neg(\forall x)A(x) \quad (4)$$

|

$$\neg(\forall x)B(x) \quad (5)$$

| δ for (4), a n.c.

$$\neg A(a)$$

| δ for (5), b n.c.

$$\neg B(b)$$

| γ rule for (2), a, b used for instantiation

$$A(a) \vee B(b) \quad (6)$$

|

$$A(a) \vee B(b) \quad (7)$$

|

$$(\forall x)(A(x) \vee B(x)) \quad \text{copy of (2)}$$

β rule for (6)

$$A(a)$$

$$B(a)$$

closed $(\neg A(a), A(a)) \leftarrow \otimes$

β rule for (7)

$$A(b)$$

$$B(b)$$

$$\odot$$

$\otimes \rightarrow$ closed $(\neg B(b), B(b))$

\odot = open branch

n.c. = new constant

The open branch provides models for $\neg U_2$, which are anti-models of U_2 .

$$\neg A(a) \wedge \neg B(b) \wedge B(a) \wedge A(b)$$

$$I_1 = \langle D, m_1 \rangle, D = \{a, b\}$$

$$m_1(A): D \rightarrow \{T, F\}, m_1(A)(a) = F \quad m_1(A)(b) = T$$

$$m_1(B): D \rightarrow \{T, F\}, m_1(B)(a) = T \quad m_1(B)(b) = F$$

↑
generic model

→ anti-model of
 $U_2 \Rightarrow U_2$ is
not a tautology