

# MATHEMATICAL FOUNDATION FOR COMPUTER SCIENCE

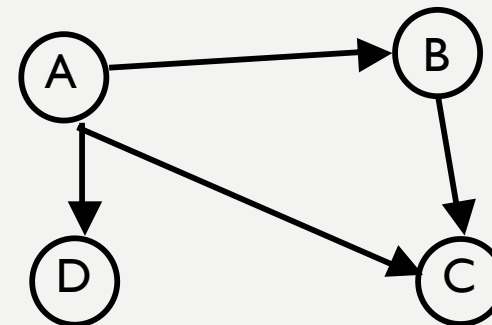
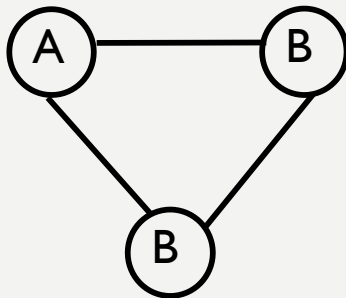
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# GRAPH THEORY

# GRAPHS:

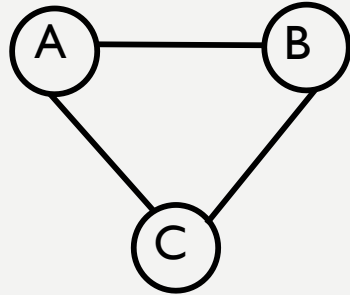
- Graphs are discrete structures consisting of vertices and edges that connect these vertices.
- There are different kinds of graphs, depending on whether edges have directions, whether multiple edges can connect the same pair of vertices, and whether loops are allowed.
- **A graph  $G = (V, E)$  consists of  $V$ , a nonempty set of vertices (or nodes) and  $E$ , a set of edges. Each edge has either one or two vertices associated with it, called its endpoints. An edge is said to connect its endpoints.**
- The set of vertices  $V$  of a graph  $G$  may be infinite. A graph with an infinite vertex set or an infinite number of edges is called an infinite graph, and in comparison, a graph with a finite vertex set and a finite edge set is called a finite graph.



# TYPES OF GRAPHS:

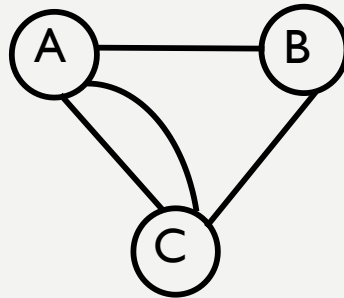
I. **Undirected Graph:** A Graph whose edges are undirected is called undirected Graph

- a. *Simple Graph:* A graph in which each edge connects to the two different vertices and no two edges connect same pair of vertices is called a Simple Graph[no- parallel edges and no loops]

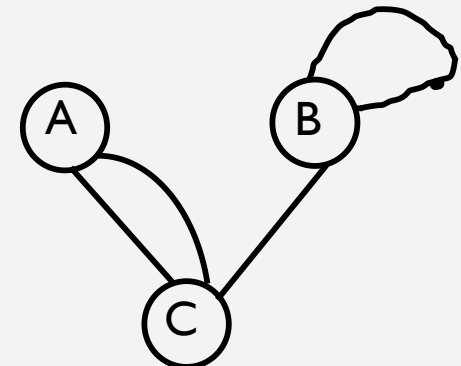


Maximum number of edges possible in a simple graph with  $n$  vertices is  $\frac{n(n-1)}{2}$

- b. *Multi Graph:* If in a graph multiple edges between the same set of vertices are allowed, it is called multigraph.



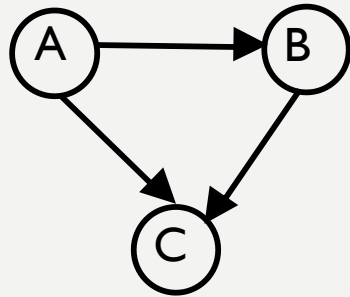
- c. *Pseudo Graph :* It is a multigraph with loops.



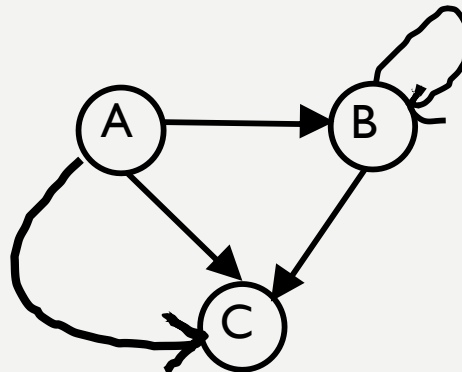
# TYPES OF GRAPHS:

I. **Directed Graph**: A Graph whose edges are directed is called directed Graph

- a. *Simple Directed Graph*: A directed graph in which each edge connects to the two different vertices and no two edges connect same pair of vertices is called a Simple Graph [no- parallel edges and no loops]



- b. *Multiple Directed Graph*: If in a Directed graph multiple edges between the same set of vertices and loops are allowed, it is called multigraph.

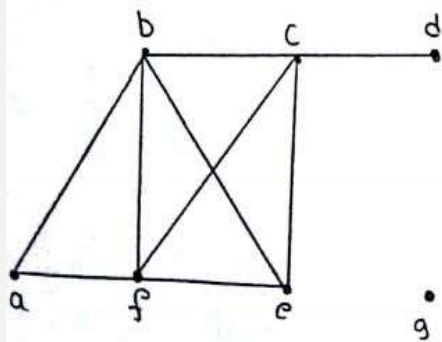


2. **Mixed Graph**: A graph with both directed and undirected edges is called a mixed graph.

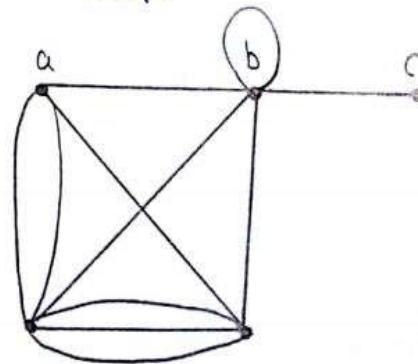
# GRAPH TERMINOLOGIES:

- First, we give some terminology that describes the vertices and edges of undirected graphs
- Two vertices  $u$  and  $v$  in an undirected graph  $G$  are called adjacent (or neighbors) in  $G$  if  $u$  and  $v$  are endpoints of an edge  $e$  of  $G$ . Such an edge  $e$  is called incident with the vertices  $u$  and  $v$  and  $e$  is said to connect  $u$  and  $v$ .
- The **degree** of a vertex in an undirected graph is the number of edges incident with it, except that a loop at a vertex contributes twice to the degree of that vertex. The degree of the vertex  $v$  is denoted by  $\deg(v)$ .

• Example: What are the degree of the vertices in the graphs  $G$  and  $H$ ? ~~graph~~



[G]



[H]

Solution: In  $G$ .

$$\begin{aligned}\deg(a) &= 2 \\ \deg(b) &= 4 \\ \deg(c) &= 4 \\ \deg(d) &= 1 \\ \deg(e) &= 3 \\ \deg(f) &= 4 \\ \deg(g) &= 0\end{aligned}$$

In  $H$

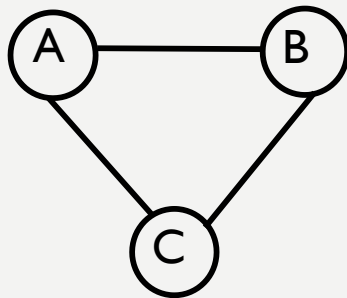
$$\begin{aligned}\deg(a) &= 4 \\ \deg(b) &= 6 \\ \deg(c) &= 6 \\ \deg(d) &= 5 \\ \deg(e) &= 1\end{aligned}$$

# GRAPH TERMINOLOGIES:

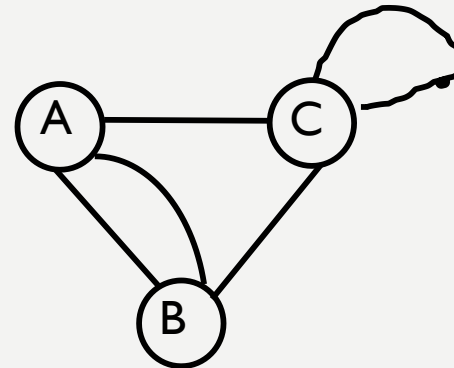
- A vertex of degree zero is called **isolated**. It follows that an isolated vertex is not adjacent to any vertex.
- A vertex is **pendant** if and only if it has degree one. Consequently, a pendant vertex is adjacent to exactly one other vertex.

## Theorem 1: THE HANDSHING THEOREM:

Let  $G = (V, E)$  be an undirected graph with  $m$  edges. Then  $2m = \sum_{v \in V} \deg(v)$



$$2 \cdot 3 = 2 + 2 + 2 \\ 6 = 6$$



$$2 \cdot 5 = 3 + 3 + 4 \\ 10 = 10$$

Q. How many edges are there in a graph with 10 vertices each of degree 6?

Because the sum of the degrees of the vertices is  $6 \cdot 10 = 60$ , it follows that  $2m = 60$  where  $m$  is the number of edges. Therefore,  $m = 30$ .

*Theorem 1 shows that the sum of the degrees of the vertices of an undirected graph is even*

# GRAPH TERMINOLOGIES:

**Theorem 2: An undirected graph has an even number of vertices of odd degree:**

Proof:

Take two sets of vertices,  $V_1$ , a set of vertices with even degree, and  $V_2$ , a set of vertices with odd degree in an undirected graph  $G = (V, E)$  with  $m$  edges. Then,

$$2e = \sum_{v \in V} \deg(v) = \sum_{v_1 \in V} \deg(v) + \sum_{v_2 \in V} \deg(v)$$

Because  $\deg(v)$  is even for  $v \in V_1$ , the first term in the right-hand side of the last equality is even. Furthermore, the sum of the two terms on the right-hand side of the last equality is even, because this sum is  $2m$ . Hence, the second term in the sum is also even. Because all the terms in this sum are odd, there must be an even number of such terms. Thus, there are an even number of vertices of odd degree



# GRAPH TERMINOLOGIES:

- Terminology for graphs with directed edges reflects the fact that edges in directed graphs have directions.
- Let  $(u, v)$  be an edge representing edge of a directed graph  $G$ .  $u$  is called adjacent to  $v$  and  $v$  is called adjacent from  $u$ . The vertex  $u$  is called initial vertex and the vertex  $v$  is called terminal or end vertex. Loop has same initial and terminal vertex.
- In directed graph the in-degree of a vertex  $v$ , denoted by  $\deg^-(v)$ , is the number of edges that have  $v$  as their terminal vertex (incoming edges). The out-degree of a vertex  $v$ , denoted by  $\deg^+(v)$ , is the number of edges that have  $v$  as their initial vertex (outgoing edges). (Note that a loop at a vertex contributes 1 to both the in-degree and the out-degree of this vertex.)

# GRAPH TERMINOLOGIES:

- Example: Find the in-degree and out-degree of each vertex in the graph  $G$  with directed edges shown in the following figure:

Solution:

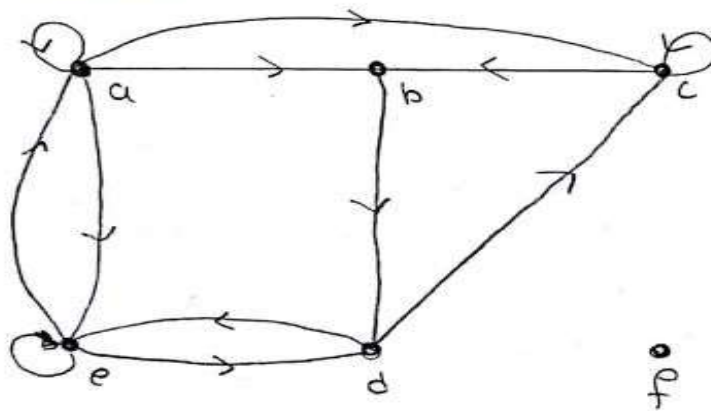


Fig: The directed graph  $G$ .

The in degrees in  $G$  are

$$\deg^-(a) = 2$$

$$\deg^-(b) = 2$$

$$\deg^-(c) = 3$$

$$\deg^-(d) = 2$$

$$\deg^-(e) = 3$$

$$\deg^-(f) = 0$$

The out degrees in  $G$  are

$$\deg^+(a) = 4$$

$$\deg^+(b) = 1$$

$$\deg^+(c) = 2$$

$$\deg^+(d) = 2$$

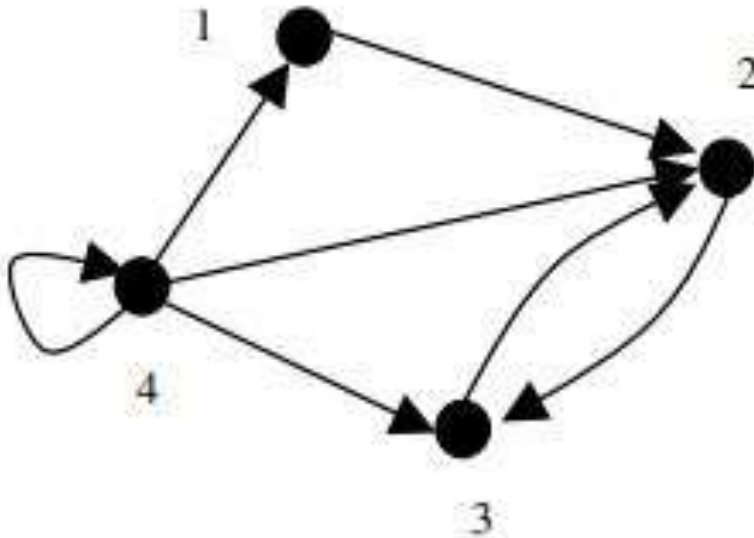
$$\deg^+(e) = 3$$

$$\deg^+(f) = 0$$

# GRAPH TERMINOLOGIES:

**Theorem 3:** Let  $G(V, E)$  be a graph with directed edges. Then,

$$\sum_{v \in V} \deg^-(v) = \sum_{v \in V} \deg^+(v) = |E|.$$

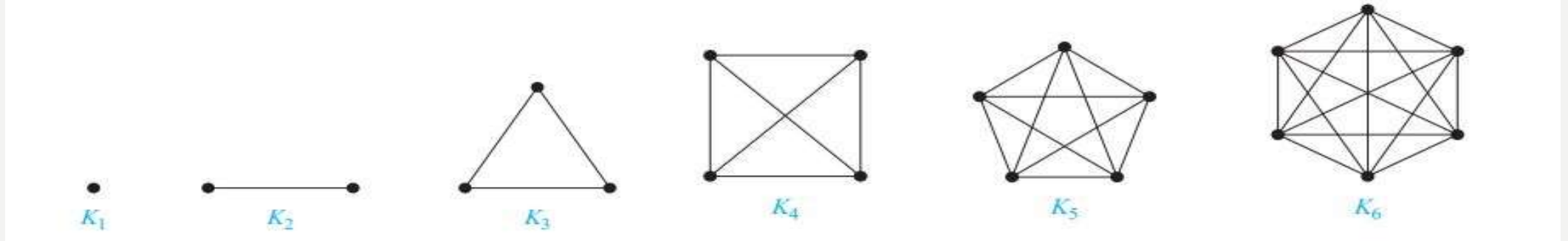


Sum of in-degree =  $1 + 1 + 3 + 2 = 7$   
Sum of out-degree =  $1 + 1 + 1 + 4 = 7$   
Total edges = 7

In-degrees of a graph are  $\deg^-(1) = \deg^-(4) = 1$ ;  $\deg^-(2) = 3$ ;  $\deg^-(3) = 2$  and the out-degrees of a graph are  $\deg^+(1) = \deg^+(2) = \deg^+(3) = 1$ ;  $\deg^+(4) = 4$ .

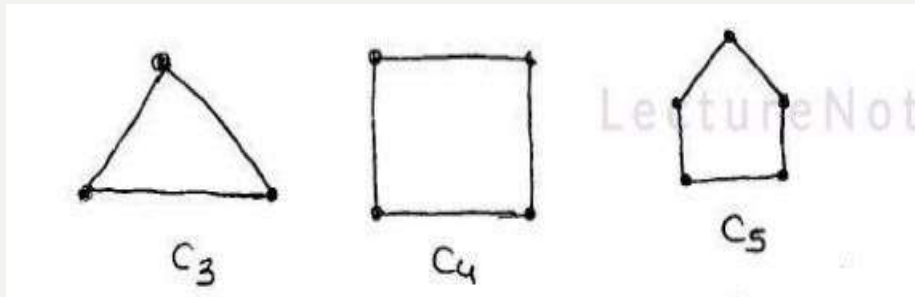
# SOME SPECIAL TYPES OF GRAPHS:

1. **Complete Graph:** A complete graph on  $n$  vertices, denoted by  $K_n$ , is a simple graph that contains exactly one edge between each pair of distinct vertices. The graphs  $K_n$ , for  $n = 1, 2, 3, 4, 5, 6$ , are displayed in below Figures.



- A complete graph is a regular graph but every regular graph is not complete graph
- A complete graph is a simple graph with max number of edges
- Number of edges in  $K_n = \frac{n(n-1)}{2}$
- Degree of each vertex =  $(n-1)$

2. **Cycle Graph:** A graph  $G$  with  $n$  vertices ( $n \geq 3$ ) and  $n$  edges, is called a cycle graph ( $C_n$ ), if all edges form cycle of length  $n$ .



**I. Prove that a complete graph with n vertices contains  $[n(n - 1)]/2$  edges**

Solution:

This is easy to prove by induction.

(a) Base Case:

If  $n = 1$ ,

$$K_1 = [1(1 - 1)]/2 = 0 \text{ (which is true)}$$

(b) Induction Hypothesis:

We assume that it is true for some arbitrary value  $k(k \geq 2)$  i.e.  $[k(k-1)]/2$

(c) Induction Step:

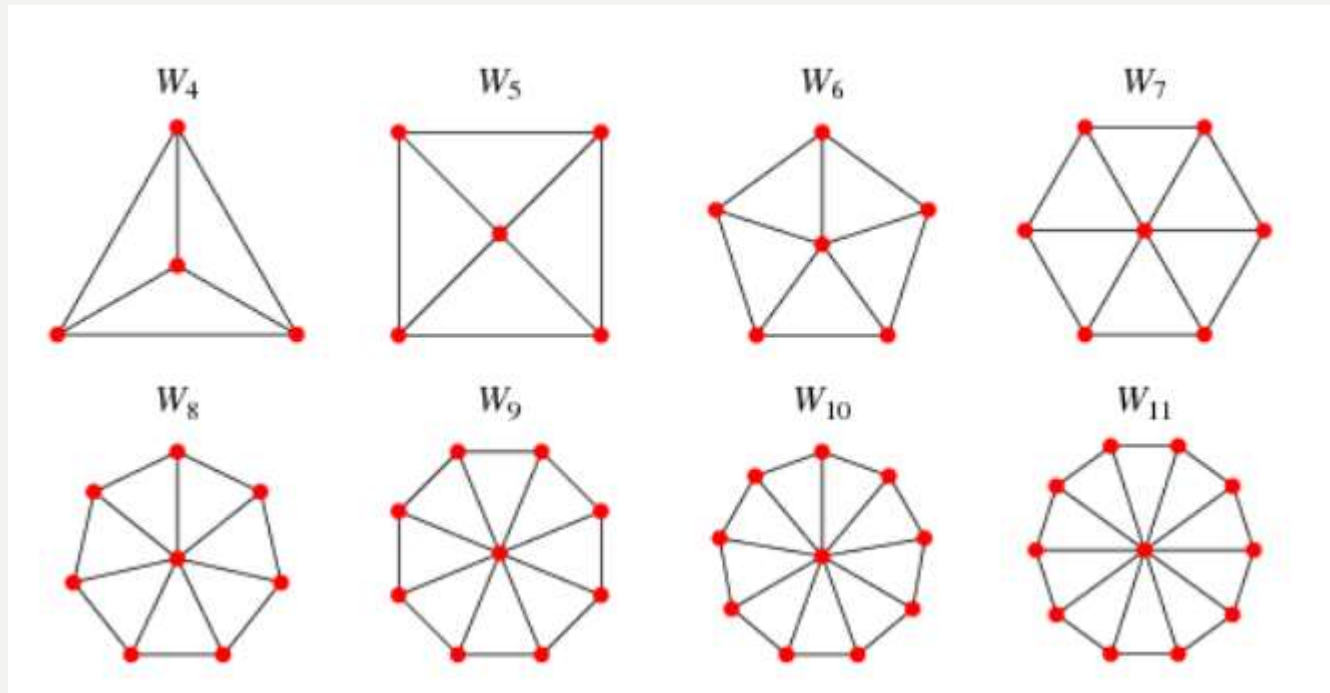
We have to prove it for  $(k+1)$  i.e.  $[(k+1)(k)]/2$

When we add the  $(k + 1)^{\text{st}}$  vertex, we need to connect it to the  $k$  original vertices, requiring  $k$  additional edges.

$$\begin{aligned} \text{We will then have } & \frac{k(k-1)}{2} + k \\ &= \frac{k(k-1)}{2} + k \\ &= \frac{k(k-1) + 2k}{2} \\ &= \frac{k^2 - k + 2k}{2} = \frac{k^2 + k}{2} \\ &= \frac{k(k+1)}{2} \end{aligned}$$

# SOME SPECIAL TYPES OF GRAPHS:

3. **Wheel Graph:** A wheel graph  $W_n$  of  $n$  vertices ( $n \geq 4$ ) can be formed from a cycle graph  $C_{n-1}$  by adding a new vertex (hub) which is adjacent to all vertices of  $C_{n-1}$ . The wheel graphs are displayed below.

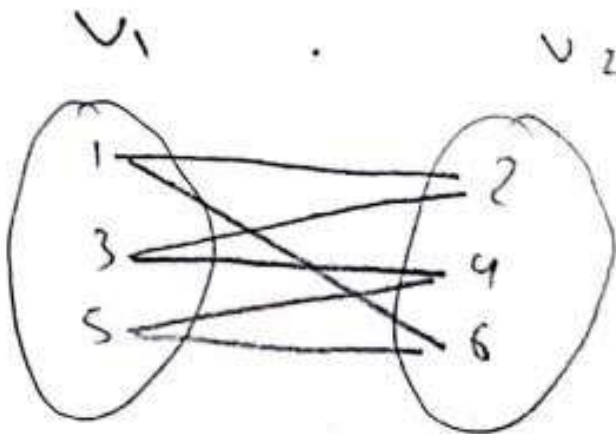
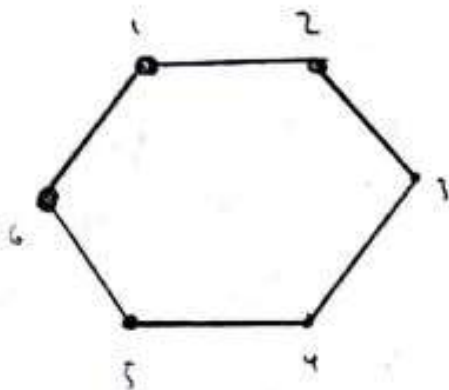


- Number of edges in  $W_n = 2(n-1)$

# SOME SPECIAL TYPES OF GRAPHS:

**4. Bipartite Graphs:** A simple graph  $G$  is called bipartite if its vertex set  $V$  can be partitioned into two disjoint sets  $V_1$  and  $V_2$  such that every edge in the graph connects a vertex in  $V_1$  and a vertex in  $V_2$  (so that no edge in  $G$  connects either two vertices in  $V_1$  or two vertices in  $V_2$ ). When this condition holds, we call the pair  $(V_1, V_2)$  a bipartition of the vertex set  $V$  of  $G$ .

Ex  $C_6$  is bipartite, because its vertex set can be partitioned into the two sets  $V_1 = \{1, 3, 5\}$  and  $V_2 = \{2, 4, 6\}$  and every edge of  $C_6$  connects a vertex in  $V_1$  and a vertex in  $V_2$ .

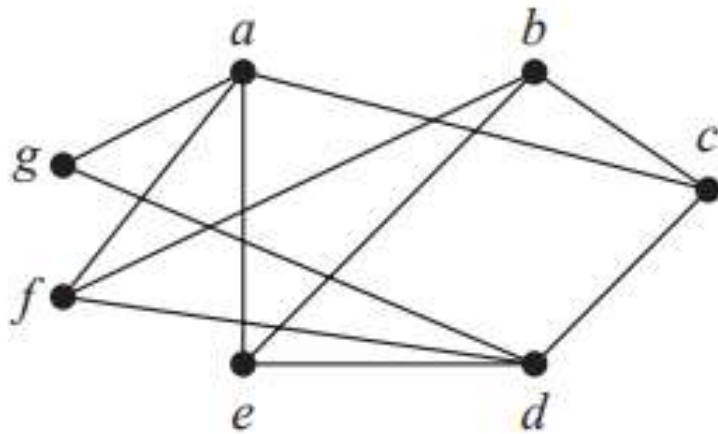


# SOME SPECIAL TYPES OF GRAPHS:

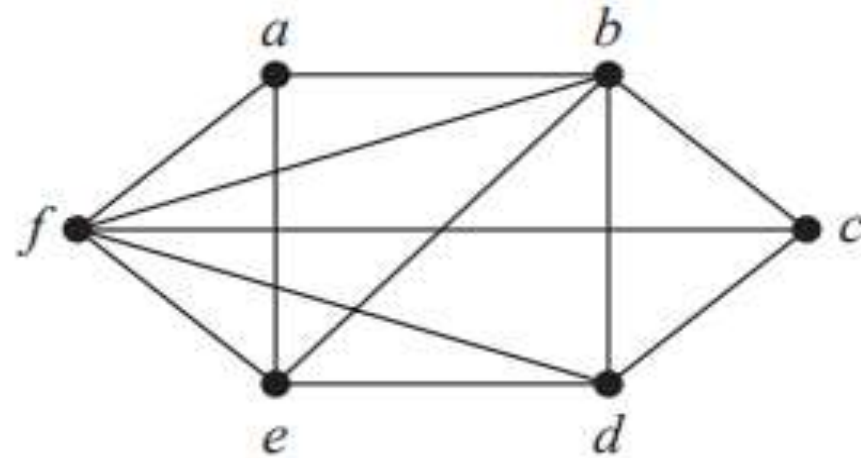
**Theorem 4 :** *A simple graph is bipartite if and only if it is possible to assign one of two different colors to each vertex of the graph so that no two adjacent vertices are assigned the same color.*

The chromatic number of the following bipartite graph is 2.

Is G and H bipartite graph??



G

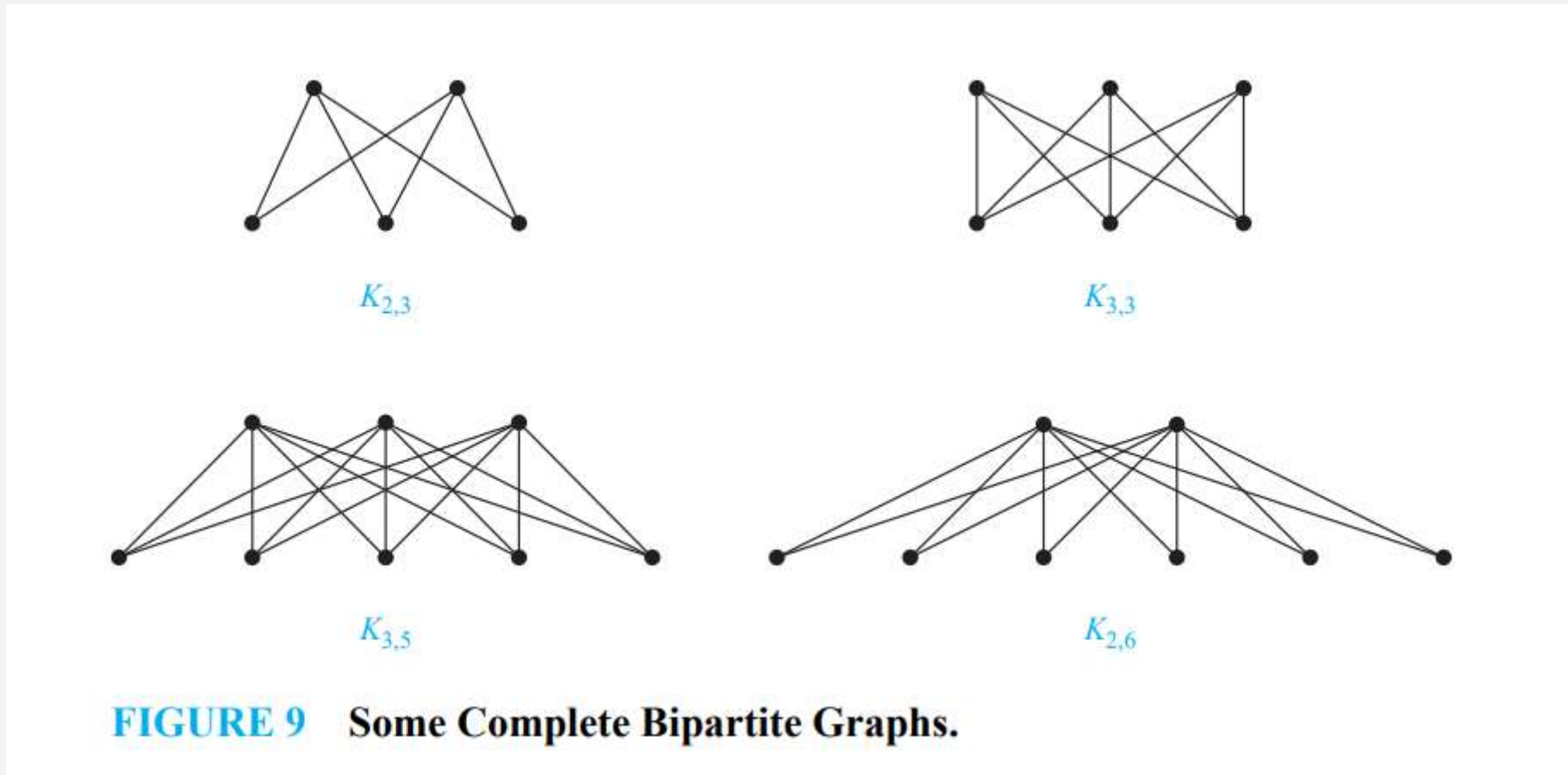


H



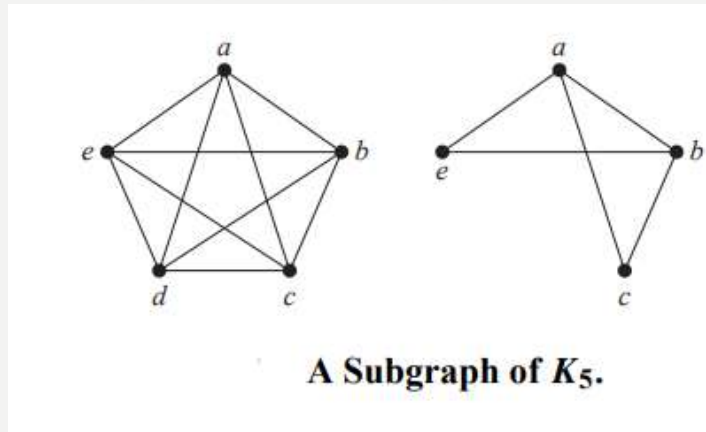
# SOME SPECIAL TYPES OF GRAPHS:

**5. Complete Bipartite Graphs:** A complete bipartite  $K_{m,n}$  graph is a special type of bipartite graph where every vertex of one set is connected to every vertex of another set.

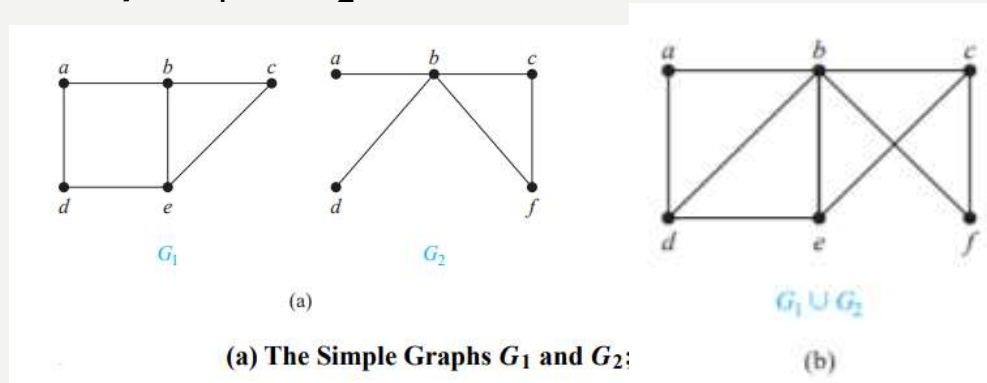


# SOME SPECIAL TYPES OF GRAPHS:

5. **Subgraph Graphs:** A subgraph of a graph  $G = (V, E)$  is a graph  $H = (W, F)$ , where  $W \subseteq V$  and  $F \subseteq E$ . A subgraph  $H$  of  $G$  is a proper subgraph of  $G$  if  $H \neq G$ .



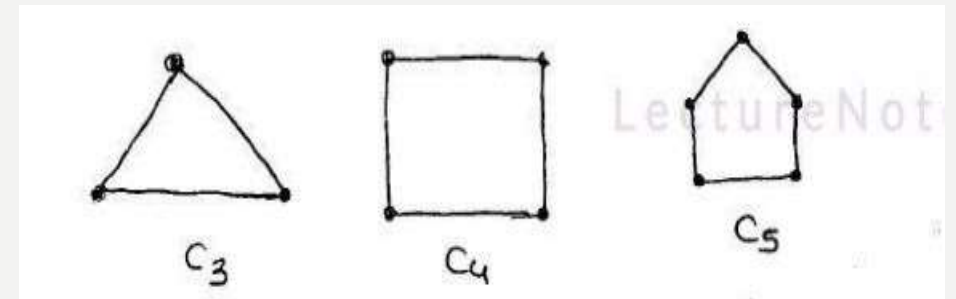
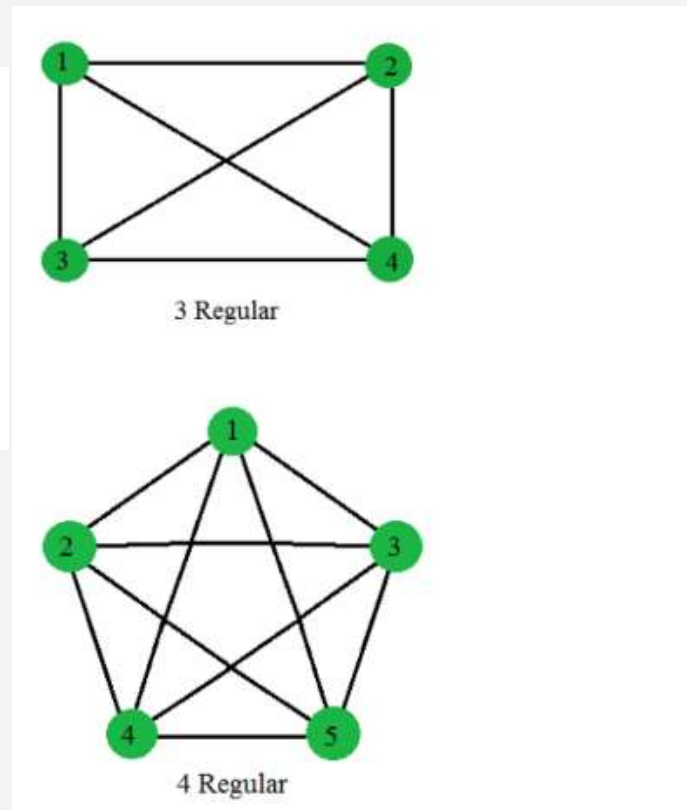
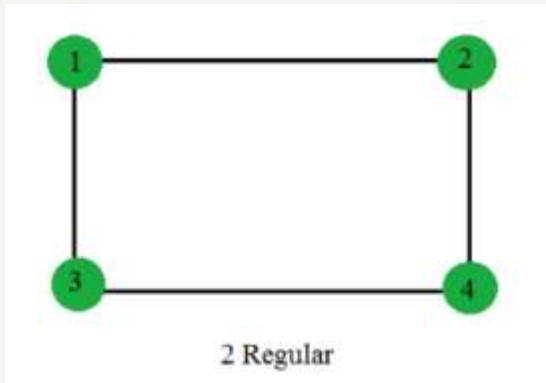
**GRAPH UNIONS:** Two or more graphs can be combined in various ways. The new graph that contains all the vertices and edges of these graphs is called the union of the graphs. The union of two simple graphs  $G_1 = (V_1, E_1)$  and  $G_2 = (V_2, E_2)$  is the simple graph with vertex set  $V_1 \cup V_2$  and edge set  $E_1 \cup E_2$ . The union of  $G_1$  and  $G_2$  is denoted by  $G_1 \cup G_2$ .



# SOME SPECIAL TYPES OF GRAPHS:

**6. Regular Graphs:** A graph is called regular graph if degree of each vertex is equal. A graph is called **K regular** if degree of each vertex in the graph is K.

- A complete graph with N vertices is  $(N-1)$  regular.
- Cycle( $C_n$ ) is always 2 Regular.



2 Regular

I. **Prove that a maximum number of edges possible in a simple graph with n vertices is  $[n(n - 1)]/2$**

Solution:

By Handshaking Theorem, We have

$$2m = \sum_{v \in V} \deg(v) \text{ where } m = \text{number of edges with } n \text{ vertices in the Graph } G$$
$$\deg(v_1) + \deg(v_2) + \deg(v_3) + \dots + \deg(v_n) = 2m \text{-----(i)}$$

We know that the maximum degree of a vertex in a Simple graph can be  $(n-1)$ ,

We can write equation (i) as,

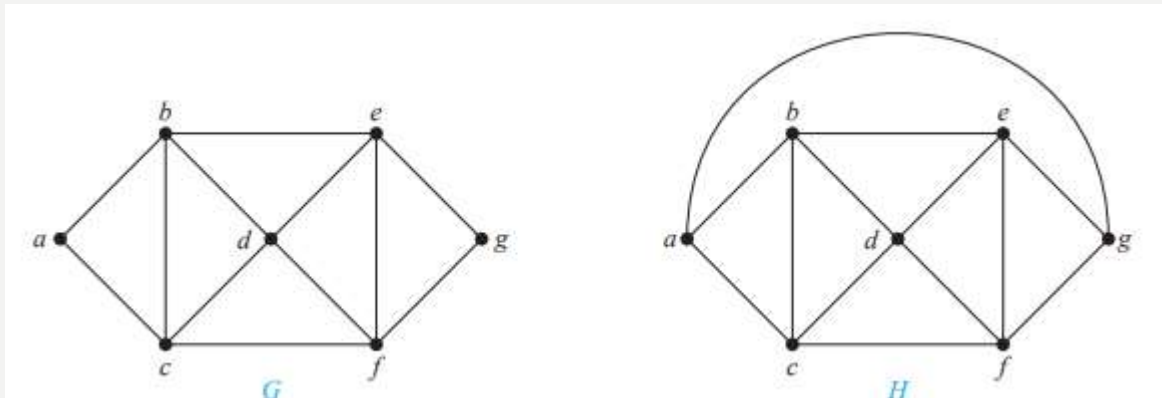
$$(n-1) + (n-1) + (n-1) + \dots \text{up to } n \text{ vertices} = 2m$$

$$n(n-1) = 2m$$

$$m = [n(n - 1)]/2$$

# CHROMATIC NUMBER:

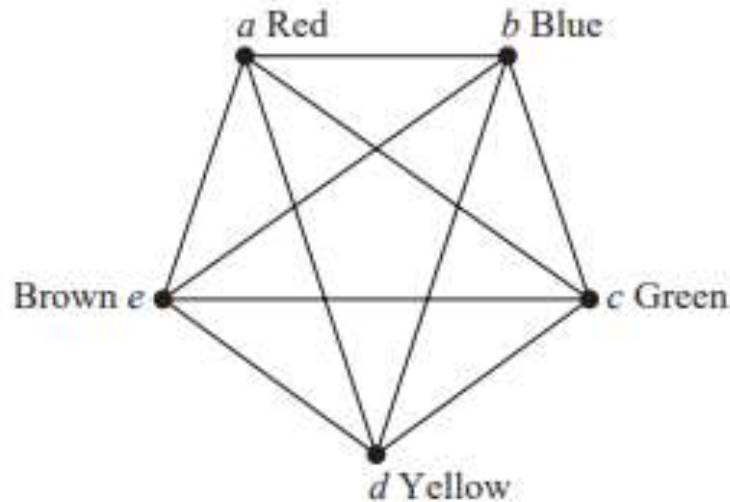
- A graph can be colored by assigning a different color to each of its vertices. However, for most graphs a coloring can be found that uses fewer colors than the number of vertices in the graph. What is the least number of colors necessary?
- The chromatic number of a graph is the least number of colors needed for a coloring of this graph. The chromatic number of a graph  $G$  is denoted by  $\chi(G)$ . (Here  $\chi$  is the Greek letter chi.)
- What are the chromatic numbers of the graphs  $G$  and  $H$  shown in Figure ?



- :The chromatic number of  $G$  is 3, and  $H$  has a chromatic number equal to 4.

# CHROMATIC NUMBER:

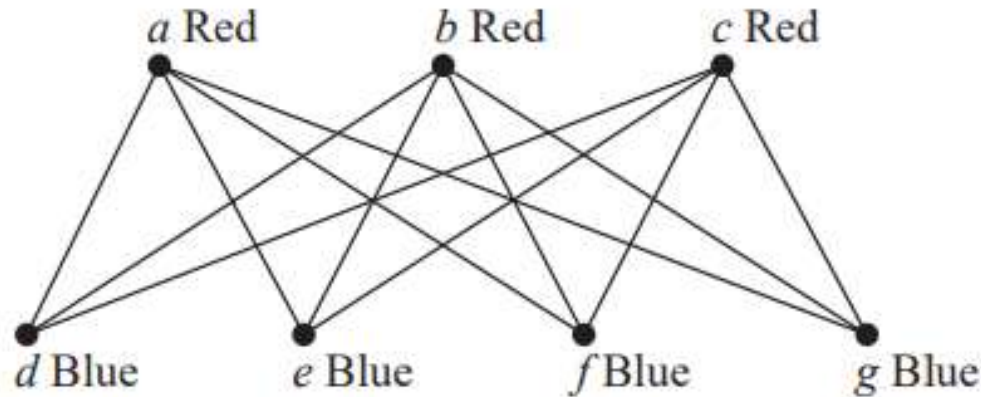
- What is the chromatic number of  $K_n$ ?
  - A coloring of  $K_n$  can be constructed using  $n$  colors by assigning a different color to each vertex.
  - Is there a coloring using fewer colors? The answer is no. No two vertices can be assigned the same color, because every two vertices of this graph are adjacent.
  - **Hence, the chromatic number of  $K_n$  is  $n$ . That is,  $\chi(K_n) = n$ .**
- A coloring of  $K_5$  using five colors is shown in Figure .



**FIGURE 5** A Coloring of  $K_5$ .

# CHROMATIC NUMBER:

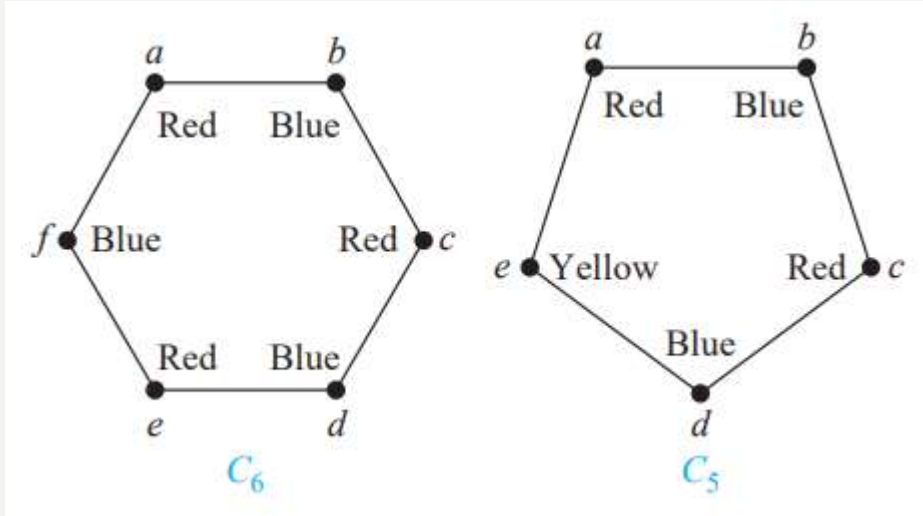
- What is the chromatic number of the complete bipartite graph  $K_{m,n}$ , where  $m$  and  $n$  are positive integers?
  - We can color the set of  $m$  vertices with one color and the set of  $n$  vertices with a second color.
  - Because edges connect only a vertex from the set of  $m$  vertices and a vertex from the set of  $n$  vertices, no two adjacent vertices have the same color.
  - Hence, chromatic number of the complete bipartite graph is 2 i.e.  $\chi(K_{m,n}) = 2$
- A coloring of  $K_{3,4}$  with two colors is displayed in Figure.



**FIGURE 6** A Coloring of  $K_{3,4}$ .

# CHROMATIC NUMBER:

- What is the chromatic number of the graph  $C_n$ , where  $n \geq 3$ ? ( $C_n$  is the cycle with  $n$  vertices.)



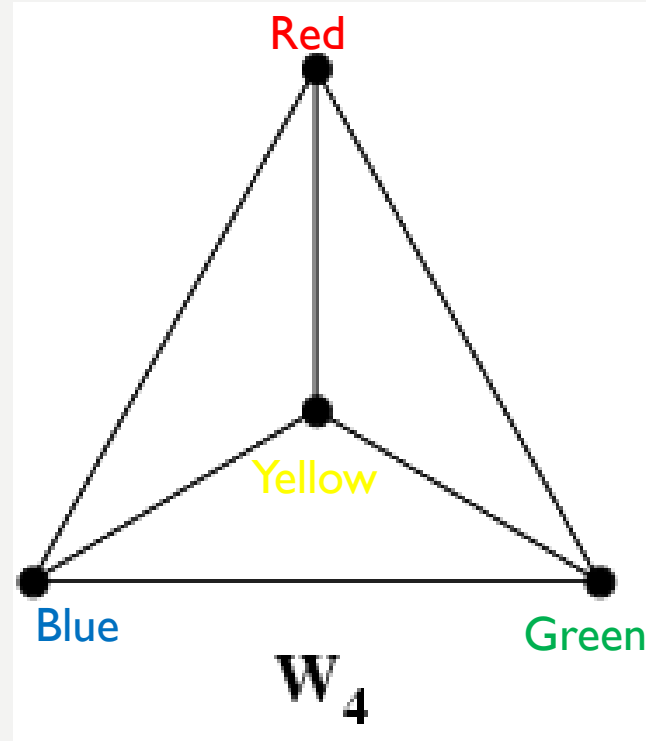
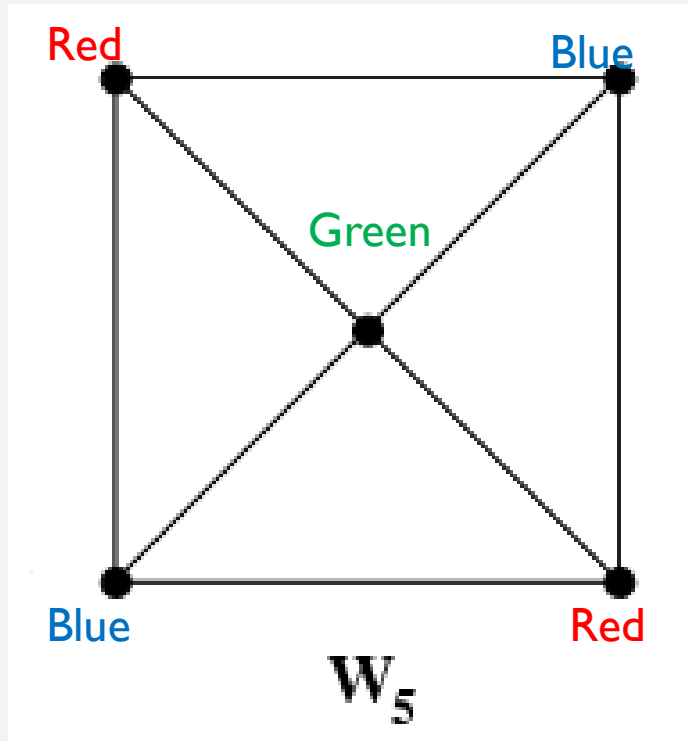
- The chromatic number of  $C_6$  is 2
- The chromatic number of  $C_5$  is 3

- $\chi(C_n) = 2$  if  $n$  is an even positive integer with  $n \geq 4$
- $\chi(C_n) = 3$  if  $n$  is an odd positive integer with  $n \geq 3$ .



# CHROMATIC NUMBER:

- What is the chromatic number of  $W_n$ ?



- The chromatic number of  $W_5$  is 3
- The chromatic number of  $C_4$  is 4

- $\chi(W_n) = 4$  if  $n$  is an even positive integer with  $n \geq 4$
- $\chi(W_n) = 3$  if  $n$  is an odd positive integer with  $n \geq 5$ .