

Theme:

Date: / /
Sat Sun Mon Tue Wed Thu Fri

Lectures,

Q; critical section Problem?

⇒ Critical section problem happen when multiple processes are running at the same time and they all want to access or change the same shared data.

Each process has a part of its code called the critical section where it is uses that shared data. The problem making sure that only one process at a time can be in its critical section at a time so the shared data doesn't get messed up or given wrong result.

Each process must ask permission to enter critical section in entry section. may follow critical section with exit section, then remainder section.

```
do {  
    entry section  
    critical section  
    exit section  
} while (true);
```

Date: / /
Sat Sun Mon Tue Wed Thu Fri

Theme:

Q; Explain three critical section problem solution ?

→ There are three critical section Problem Solution. 1. Mutual Exclusion.
2. Progress . 3. Bounded waiting

Mutual exclusion: This means - that only one process can be in critical section at a time . If one process is already using the shared resource no other process should be allowed to enter the critical section until the 1st one finishes.

Progress: If no process is in the critical section and some processes want to enter, then one of them should be allowed to go in without unnecessary delay . The processes must not block each other forever if no one is inside .

Bounded waiting: once a process has asked to enter the critical section, there should be a limit on how many other processes can go ahead of it. It should not have to wait forever while others keep entering the critical Section again and again

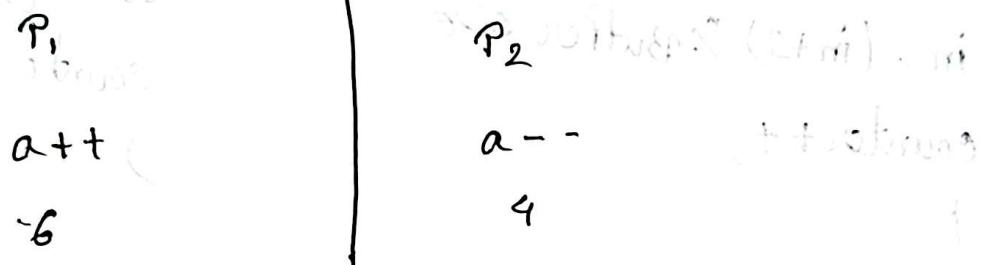
Theme:

Q; What is synchronization? Explain with example?

→ Synchronization in operating systems means making sure that multiple process or threads that run at the same time work together properly without interfering with each other, especially when they share data or resources.

Example :

`int a = 5; [Shared variable]`



Q; Explain Race condition with produce & consumer example?

→ A Race condition happens when two or more processes try to change the same data at the same time, and the final result depends on who finish first.

Example :

Lets The counter = 5;

- Producer wants to do `counter++` (increase to 6)
 - Consumer wants to do `counter--` (decrease to 4)
- without synchronization;

1. Producer read = 5
2. Consumer read = 5
3. Producer add1 = 6
4. Consumer remove 1 = 4
5. Producer store = 6
6. Consumer store = 4

The final result is = 4 (But it should be 5)

Date: / /
 Sat Sun Mon Tue Wed Thu Fri

Theme:

Buffer example (producer and consumer with one buffer)

Producer Consumer
 while (true) while (true)
 {
 while (counter < Buffer size) {
 Buffer[in] = next_producer; in = buffer[in];
 in = (in + 1) % Buffer size out = (out + 1) % Buffer size;
 counter++; counter++;
 } }
 } }

Producer produces & consumer takes out items from buffer

at the same time it just makes consumer wait for producer
 they have to wait until each other comes out of loop
 - left-hand side no break

left margin) is a producer of the above consumer.
 (right margin) -> taking of the above consumer.
 (middle margin) breaking
 S - last remaining -> a - Buffer producer.
 P - producer remains -> a - Buffer consumer.
 P - static remaining -> a - static producer
 (and break if last) -> a - static consumer.

Theme:

Date: / /
Sat Sun Mon Tue Wed Thu Fri

Peterson's Solution:

```
Ri          Rj          Pj  
do {  
    flag[i] = True;  
    turn = j;  
    while (flag[i] & & turn == i);  
    //critical section  
    flag[i] = False;  
    //reminder section  
} while (true)
```

Here using two variable -
 1. flag [which process is ready to enter]
 2. turn [whose turn it is]

Here, 3 critical problem solution are worked : (Mutual exclusion, Progress, Bounded waiting)

Mutual exclusion → 1. Process can access shared resource at a time.
Progress → 2. If one process is waiting for a resource, then another process can access it.

Bounded waiting → 3. If one process is waiting for a resource, then another process can access it after some time.

→ If one process is waiting for a resource, then another process can access it after some time.

Date: / /
 Sat Sun Mon Tue Wed Thu Fri

Theme:

57 Test and set lock solution

boolean TestAndSet (boolean *target)

{

 boolean rv = *target;

 *target = True;

 return rv;

}

P₁

do {

 while (TestAndSet (&lock));

 // critical section

 lock = false;

 // Reminder section

} while (true);

P₂

do {

 while (TestAndSet (&lock));

 // critical section

 lock = false;

 // Reminder section

} while (true);

There are two variable here. ^{Test} [lock]

Here use just one variable → lock;

Test and Set [function name]

lock 0 → CS → Available
 lock 1 → unavailable

Only mutual exclusion & program are fulfill Here. Bounded waiting never meet Here.

Date: / /
 Sat Sun Mon Tue Wed Thu Fri

Theme:

④ compare & swap instruction:

int compare-And-swap (int *locked, 0, 1)

```

{
    int value = *locked;
    if (value == 0)
        *locked = 1;
    return value
}
```

P₁

```

do {
    while (compare And swap (&lock, 0, 1) != 0);
    // CS
    lock = 0;
    // RS
} while (true);
```

④ The variable is lock

lock 0 = available
lock 1 = unavailable

④ Bounded waiting can't be fulfilled! Here

Date: / /
 Sat Sun Mon Tue Wed Thu Fri

Theme:

■ Mutex lock :

P₁

white (true)

{

while (lock != 0);

lock = 1;

// CS

lock = 0;

RS

}

Atomic Instruction

acquire (int *L)

{

while (*L != 0);

*L = 1;

}

release (int *L)

{

*L = 0;

}

P₁

white (true)

{

acquire (&lock);

// CS

release (&lock);

// RS

}

■ Function → acquire and release

variable → lock

■ Bounded waiting can't fulfill here.

Date: / /

 Sat Sun Mon Tue Wed Thu Fri

Theme:

② Semaphore :

wait (semaphore s)

{
while ($s \leq 0$);

$s--;$

}

signal (semaphore s)

{
 $s++;$

}

P₁

while (true)

{
wait (s);

//cs

signal (s);

//RS

}

P₂

while (true)

{
wait (s);

//cs

signal (s);

//RS

}

Function → wait and signal

variable → s

[
 $s = 0$ = not available
 $s = 1$ available]

spinlock problem solve

Date: / /
Sat Sun Mon Tue Wed Thu Fri

Theme:

Lab - 9;

if else statement:

Syntax

if [condition]

then

 statement

else

 statement

fi

#!/bin/bash

if else if statement:

Syntax

if [condition]

then

 statement

else elif [condition]

 statement

then

 statement

else

 statement

AND

0	0 - 0
0	1 - 0
1	0 - 0

Theme:

1 - 1 - 0

0	0	0
0	1	1
1	0	1
1	1	1

$a = 10$

Date: / /
 Sat Sun Mon Tue Wed Thu Fri

Operators:

-eq or ==

$[\$a > 18] \text{ || } [\$a < 30]$

-then

echo "True"

else

echo "False"

fi

Count 1p

if [\$count

$a = 10$

if $[\$a > 18] \& \& [\$val < 30]$

then

echo "False"

else

echo "True"

fi

then

echo "equal"

else

echo "not equal"

fi

Date: / /
Sat Sun Mon Tue Wed Thu Fri

Theme:

Check if a number is even or odd.

#!/bin/bash

if [\$num % 2 = 0]

then

echo "even"

else

echo "odd"

fi

Take two numbers as input and print the greater one

#!/bin/bash

if [\$a > \$b]

Then

Echo "a is greater than b"

else

Echo "b is greater than a"

fi

Check a number is positive, negative or zero

#!/bin/bash

if [\$num -gt 0]

Then

Echo "positive number"

elif [\$num -lt 0]

Echo "negative number"

than

Echo "zero"

fi

Date: / /

Sat Sun Mon Tue Wed Thu Fri

Theme:

Q) Take mark as input:

A+ : 90 - 100

A : 85 - 89

B+ : 80 - 84

B : 75 - 79

C+ : 70 - 74

C : 65 - 69

D+ : 60 - 64

D : 55 - 59

E/F : Below 55

#!/bin/bash

read -p "Enter your marks: " marks

if [\$marks -gt 90] && [\$marks -lt 100];

Then echo "Grade A+"

elif [\$marks -gt 85] && [\$marks -lt 89];

Then echo "Grade A"

elif [\$marks -gt 80] && [\$marks -lt 84];

Then echo "Grade B+"

elif [\$marks -gt 75] && [\$marks -lt 79];

Then echo "Grade B"

else

fi "Grade F"

Lecture 7;

Q; what is multilevel Queue scheduling?

→ Multilevel Queue scheduling is a way to manage processes by putting them into different groups based on their type or priority. Each queue can use a different scheduling method. Ex: high priority task may use Round Robin, lower priority task may use FCFS. The CPU always picks tasks from the highest priority queue first, and it only looks at the next queue if the higher is empty.

Q; What is feedback Queue scheduling?

→ It is a smart way to manage processes where ~~task~~ tasks can move different queue based on how they behave. At first every task starts in the highest priority queue. If a task take too long it moved to a lower-priority queue. If it keep taking long, it moves even lower. This way fast task finish quickly, and slow tasks get a fair chance. It keep fair balance.

Date: / /
Sat Sun Mon Tue Wed Thu Fri

Theme:

Q; what is load balancing?

⇒ In a multiprocessor system, if one CPU is overloaded with many tasks while others are lightly loaded, the overall system performance suffers. Load balancing ensures that all processors are kept busy. This helps the computer run fast and smoother.

There are 2 main types of load balancing;

1. Pull Migration, 2. Push Migration.

↓
An free CPU pull tasks from a overloaded CPU. → A CPU that is too busy pushes some of its tasks to other CPUs.

Q; what is processor affinity?

⇒ When a process uses a CPU, that CPU stores some helpful data in its cache. The cache helps the process run faster next time. Now if the process is suddenly moved to different CPU, it can't use the old cache. It has to start fresh. This slows things down. To avoid that, the operating system tries to keep the process on the same CPU as much as possible. This is called process affinity.

Date: / /
 Sat Sun Mon Tue Wed Thu Fri

Theme:

There are two type of Process affinity;

1. Soft Affinity \Rightarrow The OS tries to keep the process on the same CPU.
But it can move the process if needed. Its flexible.
2. Hard Affinity



The process is locked specific CPU. Its not flexible.

Q; what is latency?

\Rightarrow In OS latency means the time delay between an event occurring and the system's response to that event.

There are two types of latency; 1. Interrupt latency and Dispatch latency.

Q; what is little's law?

$$n = \lambda \times w$$

n = average queue length

w = average waiting time in queue.

λ = average arrival rate into queue.

What is your little's law?

Date: / /

 Sat Sun Mon Tue Wed Thu Fri

Theme:

④ Banker's Algorithm:

S; Here; Total resource A=10; B=5; C=7

Process	Allocation			Max Need			Available		
	A	B	C	A	B	C	A	B	C
P ₀	0	1	0	7	5	3	3	3	2
P ₁	2	0	0	3	2	2			
P ₂	3	0	2	9	0	2			
P ₃	2	1	1	2	2	2			
P ₄	0	0	2	4	3	3			
	= 7	2	5						

Total Resource:

$$A = 7 + 3 = 10$$

$$B = 2 + 3 = 5$$

$$C = 5 + 2 = 7$$

Here;

Need Matrix: (Max Allocation)

	A	B	C
P ₀	7	4	3
P ₁	1	2	2
P ₂	6	0	0
P ₃	0	1	1
P ₄	4	3	1

$$[\text{Max Allocation} = [\text{Max Need}] - [\text{Allocation}]]$$

[Total Available + Processed Allocation = New Available]

Date: / /

 Sat Sun Mon Tue Wed Thu Fri

Theme:

need \leq available

$$P_0 \rightarrow 7 4 3 \leq 3 3 2 [F]$$

$$P_1 \rightarrow 1 2 2 \leq 3 3 2 [T]$$

$$P_2 \rightarrow 6 0 0 \leq 5 3 2 [T]$$

$$P_3 \rightarrow 0 1 1 \leq 5 3 2 [T]$$

$$P_4 \rightarrow 4 3 1 \leq 7 4 3 [T]$$

$$P_0 \rightarrow 7 4 3 \leq 7 4 5 [T]$$

$$P_2 \rightarrow 6 0 0 \leq 7 5 5 [T]$$

Process	Allocation			Available		
	A	B	C	A	B	C
P ₀	0	1	0	3	3	2
P ₁	2	0	0	5	3	2 $\rightarrow P_1$
P ₂	3	0	2	7	4	3 $\rightarrow P_3$
P ₃	2	1	1	7	4	5 $\rightarrow P_4$
P ₄	0	0	2	7	5	5 $\rightarrow P_0$
				10	5	7 $\rightarrow P_2$

$\therefore P_1 \rightarrow P_3 \rightarrow P_4 \rightarrow P_0 \rightarrow P_2$

Q: Can the request 3,4,0 by P₄ be granted immediately?

: a) request \leq need

$$3 4 0 \leq 4 3 1 [T]$$

b) request \leq available

$$3 4 0 \leq 3 3 2 [F]$$

Equation:

$$1) \text{Allocation} = \text{Allocation} + \text{Request}$$

$$2) \text{Available} = \text{Available} - \text{Request}$$

$$3) \text{Need} = \text{Need} - \text{Request}$$

Theme:

Date: / /
 Sat Sun Mon Tue Wed Thu Fri

Q; can the request 2, 1, 0 by P_1 be granted immediately

a) request \leq need

$$2 \ 1 \ 0 \leq 4 \ 3 \ 1 [T]$$

b) request \leq available

$$2 \ 1 \ 0 \leq 3 \ 3 \ 2 [T]$$

Process	Allocation			Max Need			Available		
	A	B	C	A	B	C	A	B	C
P_0	0	1	0	7	5	3	1	2	2
P_1	2	0	0	3	2	2	3	2	2 $\rightarrow P_1$
P_2	3	0	2	9	0	2	5	3	3 $\rightarrow P_3$
P_3	2	1	1	2	2	2	7	4	5 $\rightarrow P_4$
P_4	2	1	2	4	3	3	7	5	5 $\rightarrow P_0$
P_5	2	1	2	10	5	7	5	7	$\rightarrow P_2$
	=	9	3	5					

For P_5 , Allocation = Allocation + Request

$$\begin{aligned} &= 0 \ 0 \ 2 + 2 \ 1 \ 0 \\ &= 2 \ 1 \ 2 \end{aligned}$$

Need Matrix:

$$\begin{matrix} A & B & C \\ P_4 & 2 & 2 & 1 \end{matrix}$$

safe state:
 $P_1 \rightarrow P_3 \rightarrow P_5 \rightarrow P_0 \rightarrow P_2$

Need \leq available

$$P_0 \rightarrow 7 \ 4 \ 3 \leq 2 \ 2 \ 1 [F]$$

$$P_1 \rightarrow 1 \ 2 \ 2 \leq 2 \ 2 \ 1 [T]$$

$$P_2 \rightarrow 6 \ 0 \ 0 \leq 4 \ 2 \ 1 [F]$$

$$P_3 \rightarrow 0 \ 1 \ 1 \leq 4 \ 2 \ 1 [T]$$

$$P_4 \rightarrow 2 \ 2 \ 1 \leq 6 \ 3 \ 2 [T]$$

$$P_5 \rightarrow 7 \ 4 \ 3 \leq 8 \ 4 \ 4 [T]$$

$$P_0 \rightarrow 6 \ 0 \ 0 \leq 8 \ 5 \ 4 [T]$$

Here, The resource

$$A = 10$$

$$B = 5$$

$$C = 7$$

\therefore New available = Available - Request

$$= 3 \ 3 \ 2 - 2 \ 1 \ 0$$

$$= 1 \ 2 \ 2$$

Theme:

Date _____
 Date _____
 Date _____
 Date _____
 Date _____
 Date _____
 Date _____

Q: Can the request $\begin{pmatrix} 3 & 3 & 1 \end{pmatrix}$ by P_4 be granted immediately?

a) Request \leq need

$$\begin{pmatrix} 3 & 3 & 1 \end{pmatrix} \leq \begin{pmatrix} 3 & 3 & 1 \end{pmatrix} [T]$$

b) Request \leq available

$$\begin{pmatrix} 3 & 3 & 1 \end{pmatrix} \leq \begin{pmatrix} 3 & 3 & 2 \end{pmatrix} [T]$$

Process	Allocation	MaxNeed	Available
P_0	$\begin{pmatrix} 0 & 1 & 0 \end{pmatrix}$	$\begin{pmatrix} 7 & 5 & 3 \end{pmatrix}$	$\begin{pmatrix} 0 & 0 & 1 \end{pmatrix}$
P_1	$\begin{pmatrix} 2 & 0 & 0 \end{pmatrix}$	$\begin{pmatrix} 3 & 2 & 2 \end{pmatrix}$	
P_2	$\begin{pmatrix} 3 & 0 & 2 \end{pmatrix}$	$\begin{pmatrix} 9 & 0 & 2 \end{pmatrix}$	
P_3	$\begin{pmatrix} 2 & 1 & 1 \end{pmatrix}$	$\begin{pmatrix} 2 & 2 & 2 \end{pmatrix}$	
P_4	$\begin{pmatrix} 3 & 3 & 3 \end{pmatrix}$	$\begin{pmatrix} 4 & 3 & 3 \end{pmatrix}$	
	$\begin{pmatrix} 10 & 5 & 6 \end{pmatrix}$		
	$\begin{pmatrix} 10 & 5 & 7(001) \end{pmatrix}$		

For, P_4 Allocation = Request + Allocation

$$= \begin{pmatrix} 3 & 3 & 1 \end{pmatrix} + \begin{pmatrix} 0 & 0 & 2 \end{pmatrix}$$

$$= \begin{pmatrix} 3 & 3 & 3 \end{pmatrix}$$

P_4 . Need matrix = Need - Request

$$\begin{matrix} & \begin{pmatrix} 4 & 3 & 1 \end{pmatrix} & - & \begin{pmatrix} 3 & 3 & 1 \end{pmatrix} \\ \therefore & \begin{pmatrix} 1 & 0 & 0 \end{pmatrix} & = & \end{matrix}$$

Need \leq available

$$P_0 \rightarrow \begin{pmatrix} 7 & 4 & 3 \end{pmatrix} \leq \begin{pmatrix} 0 & 0 & 1 \end{pmatrix} [F]$$

$$P_1 \rightarrow \begin{pmatrix} 1 & 2 & 2 \end{pmatrix} \leq \begin{pmatrix} 0 & 0 & 1 \end{pmatrix} [F]$$

$$P_2 \rightarrow \begin{pmatrix} 6 & 0 & 0 \end{pmatrix} \leq \begin{pmatrix} 0 & 0 & 1 \end{pmatrix} [F]$$

$$P_3 \rightarrow \begin{pmatrix} 0 & 1 & 1 \end{pmatrix} \leq \begin{pmatrix} 0 & 0 & 1 \end{pmatrix} [F]$$

$$P_4 \rightarrow \begin{pmatrix} 1 & 0 & 0 \end{pmatrix} \leq \begin{pmatrix} 0 & 0 & 1 \end{pmatrix} [F]$$

P_4 , available = Available - Request

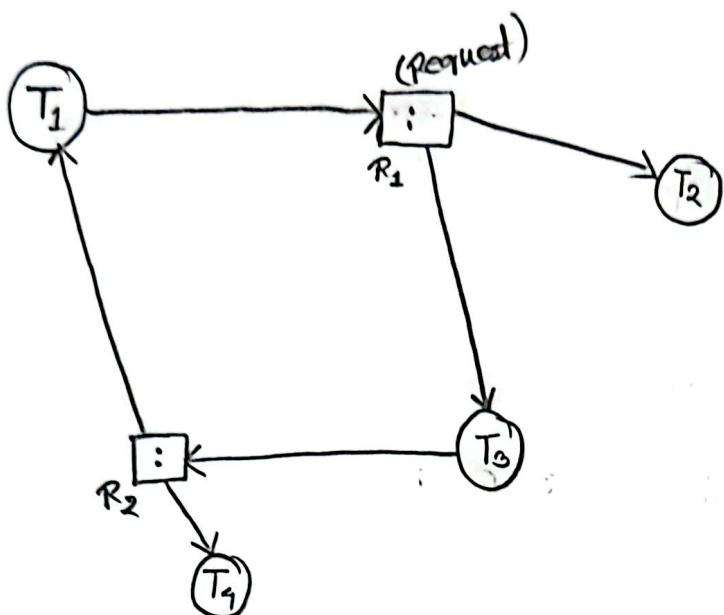
$$= \begin{pmatrix} 3 & 3 & 2 \end{pmatrix} - \begin{pmatrix} 3 & 3 & 1 \end{pmatrix}$$

$$= \begin{pmatrix} 0 & 0 & 1 \end{pmatrix}$$

Date: / /
 Sat Sun Mon Tue Wed Thu Fri

Theme:

Resource Allocation Graph:



	Allocation		Request		Available	
	R ₁	R ₂	R ₁	R ₂	R ₁	R ₂
T ₁	0	1	1	0	0	0
T ₂	1	0	0	0	1	0 → T ₂
T ₃	1	0	0	1	1	1 → T ₄
T ₄	0	1	0	0	1	2 → T ₁
	2	2			2	2 → T ₃

T₄ -> চুক্তি স্থাপিত
Allocation -> মাঠে, আব
T₄ এর স্থাপিত হলে
স্থাপিত Request -> মাঠ

Here; Resource : 2 2

: request ≤ available

$$T_1 \rightarrow 0 \ 0 \leq 0 \ 0 [F]$$

$$T_2 \rightarrow 0 \ 0 \leq 0 \ 0 [T]$$

$$T_3 \rightarrow 0 \ 1 \leq 1 \ 0 [F]$$

$$T_4 \rightarrow 0 \ 0 \leq 1 \ 0 [T]$$

$$T_1 \rightarrow 1 \ 0 \leq 1 \ 1 [T]$$

$$T_3 \rightarrow 0 \ 1 \leq 1 \ 2 [T]$$

∴ safest state: $T_2 \rightarrow T_4 \rightarrow T_1 \rightarrow T_3$

Theme:

Q; What is Address binding?

= Address binding is the process of mapping logical address to physical address in memory.

1) source code. 2) compiler. 3) object file. 4) Linker. 5) Executable file. 6) Loader. 7) CPU.

3 types of time here; 1. compile-time, 2. Load-time, 3. execute time



Memory management
unit.

Q; what is loading?

= Loading is a process of transferring a program from secondary memory to main memory.

There are 2 types of loading: 1) static loading 2. Dynamic loading

1) static loading: Here, whole program together load in the RAM; It is fast but the wastage of memory are huge.

2) Dynamic loading: whenever what need, just load that into RAM. It's also fast but didn't waste any memory.

Theme:

Date: / /
Sat Sun Mon Tue Wed Thu Fri

Q; What is swapping?

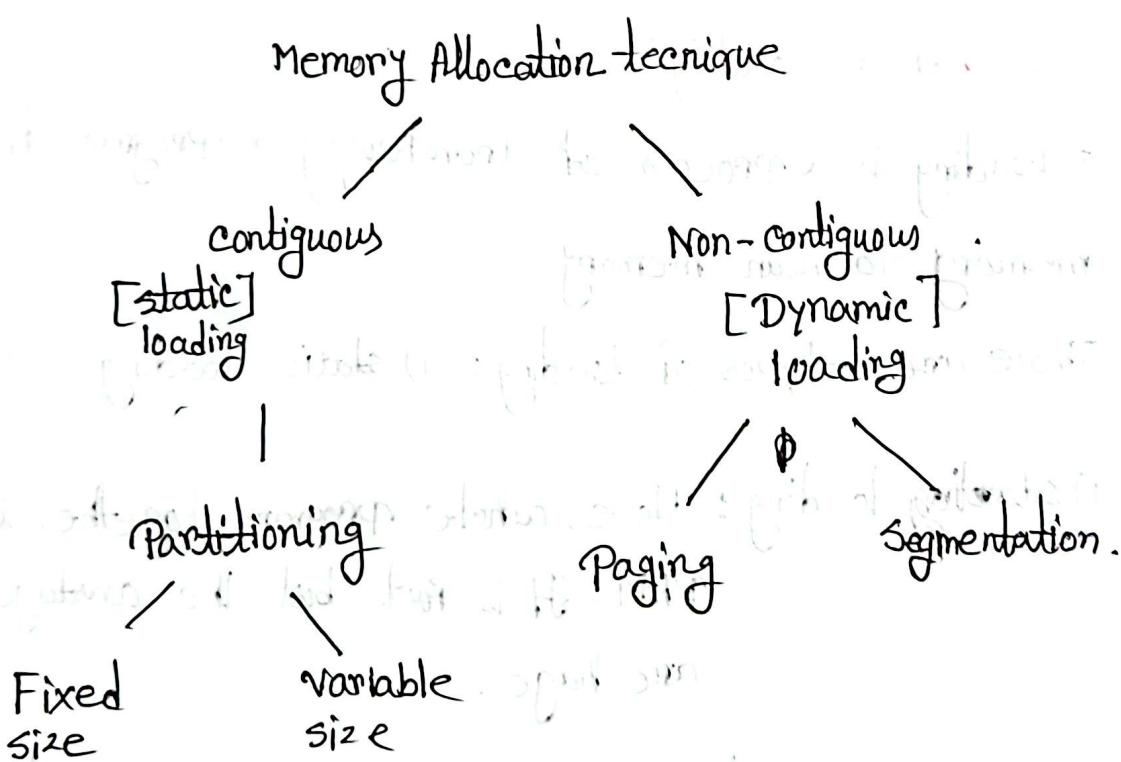
= Swapping is a memory management technique where a process temporarily moves from main memory to a backing store.

where OS removes a program from the RAM to free up the memory.

RAM \rightarrow Backing store = swap in

Backing store \rightarrow RAM = swap out

Q; Memory Allocation technique;



Theme:

Date: / /
Sat Sun Mon Tue Wed Thu Fri

Q; What is fragmentation?

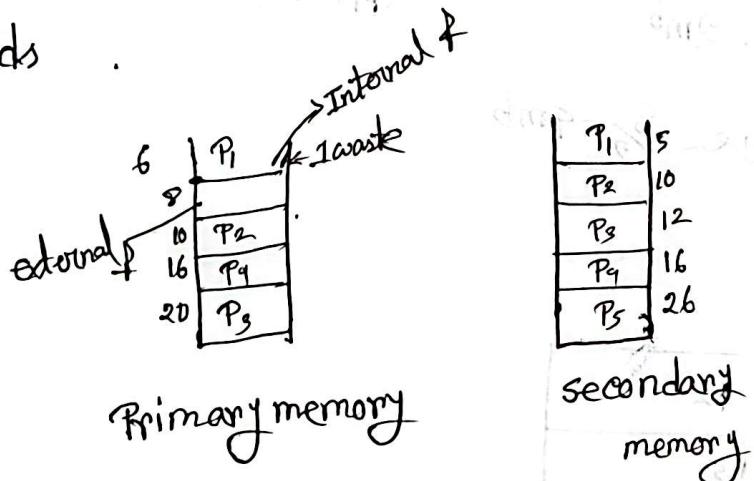
⇒ Fragmentation is a memory management problem.

There are 2 types of fragmentation:

1. External fragmentation.
2. Internal fragmentation.

External fragmentation: There is enough total free memory to run a program. but it is scattered in small, non-contiguous block. so the program can't be allocated memory as one large continuous block.

Internal fragmentation: Extra wasted inside memory block. It happen when a block is bigger than what the program actually needs.

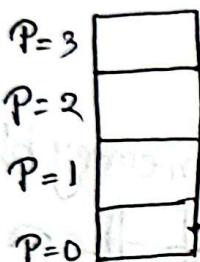


Theme:

S; what is Paging?

= Paging is a memory management technique that allows a program's memory to be divided into fixed size blocks called pages. The physical memory also divided into blocks of the same size called frames.

Page

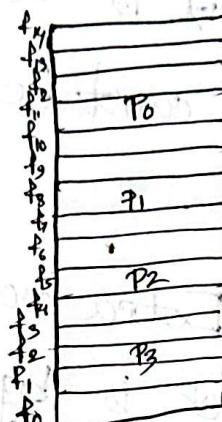


$P_1 = 8 \text{ mb}$
Block size = 2 mb

$$\therefore \text{No of page} = \frac{8}{2} = 4 \text{ page}$$

Page Table

P_3	f_2
P_2	f_5
P_1	f_8
P_0	f_{11}



$$PM = 30 \text{ mb}$$



Conversion

Theme:

Q; what is segmentation?

= segmentation is a memory management technique where a program is divided into variable-size parts called segments. It tracks each segment's location. It makes memory access more logical but it can suffer from external fragmentation.

(With fixed size block will cause internal fragmentation)

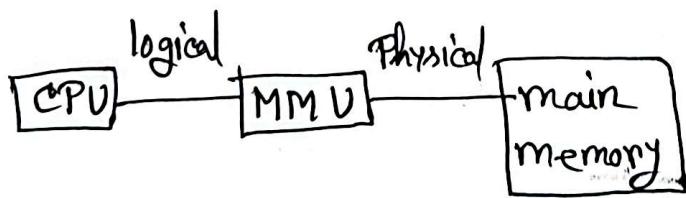
S; Physical address vs logical address

Physical address

1. Actual address in main memory.
2. Data is stored here.
3. This address used by MMU.

logical address

1. Generated by CPU.
2. It is used as a reference to access the memory located by CPU.
3. It's a virtual address.



Date: / /

Sat Sun Mon Tue Wed Thu Fri

Theme:

Lab;

- Always variable name এর আগে \$ রয়ে।
 - echo \$Name
 - read command - takes input - from - the keyboard
- Ex: read variablename
- Taking input in same line (Not in next line)

Ex: read -P "Enter variable"

- sum = \$(expr \$a + \$b)
- sub = \$(expr \$a - \$b)
- multi = \$(expr \$a * \$b)
- div = \$(expr \$a / \$b)
- mod = \$(expr \$a % \$b)

Relational operators

- eq (=)
- ne (!=)
- gt (>)
- lt (<)
- ge (>=)
- le (<=)



Date: / /
Sat Sun Mon Tue Wed Thu Fri

Theme:

④ If else statement:

if [condition]

then

statement

else

statement

fi

→ Ascii value check in

④ switch case:

case expression is

Pattern)

statement

;

pattern)

statement

فڑ

*)

echo Default

;

esac

④ If else if statement:

if [condition]

then

statement

elif [condition]

then

statement

else

statement.

fi

switch case:

var = "Name"

case \$var in

"karim")

echo "karim is good"

;

"Rahim")

echo "Rahim is bad,"

;

*)

echo "unknown"

;

esac

Theme:

OR		
0	0	0
0	1	1
1	0	1
1	1	1

AND		
0	0	0
0	1	0
1	0	0
1	1	1

Date: / /
Sat Sun Mon Tue Wed Thu Fri

OR operator:

val = 10
if [\$val -gt 18] || [\$val -lt 30]

then

True

else

False

fi

AND operator:

val = 10
if [\$val -gt 18] && [\$val -lt 30]

Then

False

else

True

fi

Date: / /
Sat Sun Mon Tue Wed Thu Fri

Theme:

□ while Loop:

while [condition]

do

statement

statement

done

Ex :

n = 1

while [\${n - le 10}]

do

echo "\$n"

n = \${n + 1}

done

□ For Loop:

for [condition]

do

statement

done

Ex :

for ((i=0; i<10; i++))

do

echo \$i

done

□ Pwd :

for c in Pwd uname

do

echo "command name : \$c"

echo "command output : "

\$c

done

Date: / /
Sat Sun Mon Tue Wed Thu Fri

Theme:

Until Loop

Syntax

until [condition]

Do

statement

done

Ex.

a = 0

until [\$a -lt 10]

do

echo \$a

a = `expr \$a + 1'

done

Break and continue

Break

for ((i=1 ; i<10 ; i++))

do

if [\$i -gt 5]

then

break

fi

echo "\$i"

done

continue;

for ((i=1 ; i<10 ; i++))

do

if [\$i -eq 5]

then

continue

fi

echo "\$i"

done

Date: / /
Sat Sun Mon Tue Wed Thu

Theme:

Using IFS (Internal File separator)

filename = user.txt

while IFS = ',' read name age address

do

echo \$name "lives in" \$address

done < \$filename

File existence check

#!/bin/bash

read -P "Enter a file name: " filename

if [-P \$filename]

then

echo "File Exist"

else

echo "File doesn't exist"

fi

arr = (10 20 5 15 35)

min = \${arr[0]}

for num in "\${arr[@]}"

do

if ((num < min))

then

min = \$num

fi

done

echo "Minimum value :" \$min

Date: / /
Sat Sun Mon Tue Wed Thu Fri

Theme:

□ Check if a number even or odd :

echo "Enter a number:"

read num

if [\$num % 2 == 0]

then

echo " Its a even number"

else

echo " Its a odd number"

fi

□ Take two number as input and print the greater one

echo "Enter two numbers:"

read a

read b

if [\$a -gt \$b]

echo "a is greater than b"

else

echo "b is greater than a"

fi

Date: / /

Sat Sun Mon Tue Wed Thu Fri

Theme:

Even or odd in a Loop :
read -P "Enter the number" i

For ((i=0 ; i<=5 ; i++))
do

if [\$i % 2 = 0]

then

echo "even number" \$i

else

echo " odd number " \$i

fi

done

Even number between 1 to 20 ; and sum them

Sum-even = 0

Sum-odd = 0

Total sum = 0

for ((i=0 ; i<=20 ; i++))

do

if ((i % 2 == 0))

then

echo "\$i is an even number"
((sum-even = sum-even + i))

else

echo "\$i is an odd number"

((sum-odd = sum-odd + i))

fi

done

total-sum = \$ ((sum-even + sum-odd))

echo " sum of even numbers: " \$sum-even

echo "sum of odd numbers: " \$sum-odd

echo "sum of total number: " \$total-sum

Date: / /
Sat Sun Mon Tue Wed Thu Fri

Theme:

② Pyramid using *

```
read -P "Enter a number of Rows: " rows
for ((i=1; i<=rows; i++)) [Print the pyramid]
do
    for ((j=i; j<=rows; j++)) [Print space]
    do
        echo -n " "
    done
    for ((k=1; k<=2*i; k++)) [Print stars]
    do
        echo -n "*"
    done
    echo
done
```

③ basic Pyramid

```
read -P "Enter number of rows: " rows
```

```
for ((i=1; i<=rows; i++))
do
    for ((j=1; j<=i ; j++))
    do
        echo -n "*"
    done
    echo
done
```

Theme:

Date: / /
Sat Sun Mon Tue Wed Thu Fri

□ wc - c

count = \$(wc - c < p1.txt)

for var in \$count

do

echo "Number of character: " \$var

done

□ wc - L

count = \$(wc - L < p1.txt)

for var in \$count

do

echo "Number of Line: " \$var

done

□ wc - w

count = \$(wc - w < p1.txt)

for var in count

do

echo "Number of word: " \$var

done

□ file existence

read -P "Enter a file name: " filename

if [-f \$filename]

then

echo "File exist"

else

echo "File doesn't exist"

fi

Date: / /
Sat Sun Mon Tue Wed Thu Fri

Theme:

④ Fibonacci number:

read -P "Enter the 'number of terms': " n

a = 0

b = 1

echo "Fibonacci series up to \$n terms: "

for ((i=0; i<n; i++))

do

echo \$a

\$fn=\$((a+b))

a=\$b

b=\$fn

done

echo

⑤ Prime number:

read -P "Enter a number: " num

if [\$num -le 1]

then

echo "\$num is not a prime number"

fi

for ((i=2; i<=num; i++))

do

if ((num % i == 0)) :

then

echo "\$num is not a prime number"

fi

done

echo "\$num is a prime number"