

CSC 503 Homework Assignment 6

Out: September 23, 2015

Due: September 30, 2015

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1. Open formula: $P(x, y) \vee (Q(x, y) \wedge \neg R(y, z))$
Using Distributive law, above formula can be rewritten as
 $(P(x, y) \vee Q(x, y)) \wedge (P(x, y) \vee \neg R(y, z))$
 - a. Thus the clausal form using set notation is : $\{\{P(x, y), Q(x, y)\}, \{P(x, y), \neg R(y, z)\}\}$
 - b. The above formula can be written so that variable can be standardized apart :
 $\{\{P(x_1, y_1), Q(x_1, y_1)\}, \{P(x_2, y_2), \neg R(y_2, z_2)\}\}$
2. Using the rules for negation normal form $\exists x \neg (P(x, y) \leftrightarrow \forall y Q(x, y))$ can be converted to NNF as per following steps:
Step 1: $\exists x \neg ((P(x, y) \rightarrow \forall y Q(x, y)) \wedge (\forall y Q(x, y) \rightarrow P(x, y)))$
Since $v \leftrightarrow w \Rightarrow (v \rightarrow w) \wedge (w \rightarrow v)$
Step 2: $\exists x \neg ((\neg P(x, y) \vee \forall y Q(x, y)) \wedge (\neg \forall y Q(x, y) \vee P(x, y)))$
Since $u \rightarrow w \Rightarrow \neg u \vee w$
Step 3: $\exists x \neg ((\neg P(x, y) \vee \forall y Q(x, y)) \wedge (\exists y \neg Q(x, y) \vee P(x, y)))$
Since $Qx \neg \forall y \phi \Rightarrow Qx \exists y \neg \phi$ NOTE: Qx stands for quantifier x .
Step 4: $\exists x (\neg (\neg P(x, y) \vee \forall y Q(x, y)) \vee \neg (\exists y \neg Q(x, y) \vee P(x, y)))$
Since $\neg (u \wedge w) \Rightarrow \neg u \vee \neg w$ using De-Morgan's law
Step 5: $\exists x ((P(x, y) \wedge \neg \forall y Q(x, y)) \vee (\neg \exists y \neg Q(x, y) \vee \neg P(x, y)))$
Since using De Morgan's Law: $\neg (u \vee w)$ is equivalent to $\neg u \wedge \neg w$
Step 6: $\exists x ((P(x, y) \wedge \exists y \neg Q(x, y)) \vee (\forall y Q(x, y) \wedge \neg P(x, y)))$
3. The formula $\forall x ((\exists y P(x, y)) \rightarrow Q(x, z)) \wedge \exists x ((\forall y R(x, y)) \vee Q(x, y))$ can be converted to prenex normal form as per following steps:
Step 1: $\forall x (\neg (\exists y P(x, y)) \vee Q(x, z)) \wedge \exists x ((\forall y R(x, y)) \vee Q(x, y))$
Since $a \rightarrow b \Rightarrow \neg a \vee b$
Step 2: $\forall x ((\forall y \neg P(x, y)) \vee Q(x, z)) \wedge \exists x ((\forall y R(x, y)) \vee Q(x, y))$
Since $Qx \neg \exists y \phi \Rightarrow Qx \forall y \neg \phi$
Step 3: $\forall x \forall t ((\neg P(x, t/y)) \vee Q(x, z)) \wedge \exists x \forall s (R(x, s/y) \vee Q(x, y))$
Since $Qx (\forall y \phi \vee \psi) \Rightarrow Qx \forall z (\phi(z/y) \vee \psi)$
Step 4: $\forall x \forall t \forall s \exists w ((\neg P(x, t/y)) \vee Q(x, z)) \wedge (R(w/x, s/y) \vee Q(w/x, y))$

4. The formula $\forall x \exists y P(x, y) \rightarrow \forall x \neg \forall y Q(x, y) \wedge \neg \exists x \exists y P(x, y)$ can be converted in Skolem Normal form as per following steps:

Step 1: $\neg (\forall x \exists y P(x, y)) \vee ((\forall x \neg \forall y Q(x, y)) \wedge (\neg \exists x \exists y P(x, y)))$

Since $u \rightarrow w \Rightarrow \neg u \vee w$

Step 2: $(\exists x \neg \exists y P(x, y)) \vee ((\forall x \exists y \neg Q(x, y)) \wedge (\forall x \neg \exists y P(x, y)))$

Since $Qx \neg \forall y \phi \Rightarrow Qx \exists y \neg \phi$ and $Qx \neg \exists y \phi \Rightarrow Qx \forall y \neg \phi$

Step 3: $(\exists x \forall y \neg P(x, y)) \vee ((\forall x \exists y \neg Q(x, y)) \wedge (\forall x \forall y \neg P(x, y)))$

Since $Qx \neg \exists y \phi \Rightarrow Qx \forall y \neg \phi$

Step 4: $(\forall y \neg P(f(y), y)) \vee ((\forall x \neg Q(x, f(x))) \wedge (\forall x \forall y \neg P(x, y)))$

5. $\theta = \{f(y)/x, g(z)/y, v/w\}$
 $\sigma = \{a/x, b/y, f(y)/z, w/v, c/u\}$
 $\theta \sigma = \{f(b)/x, g(f(y))/y, f(y)/z, w/v, c/u\}$