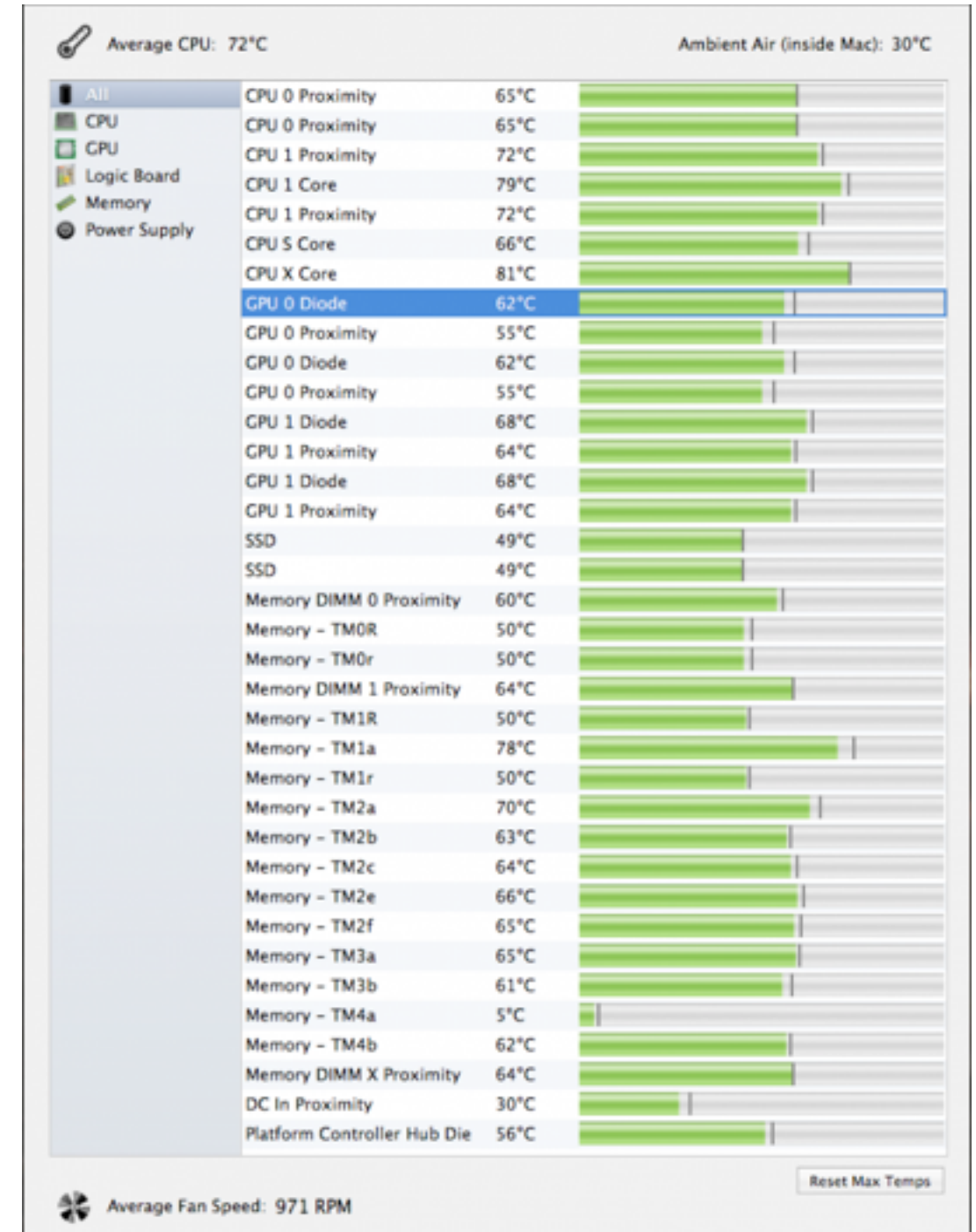
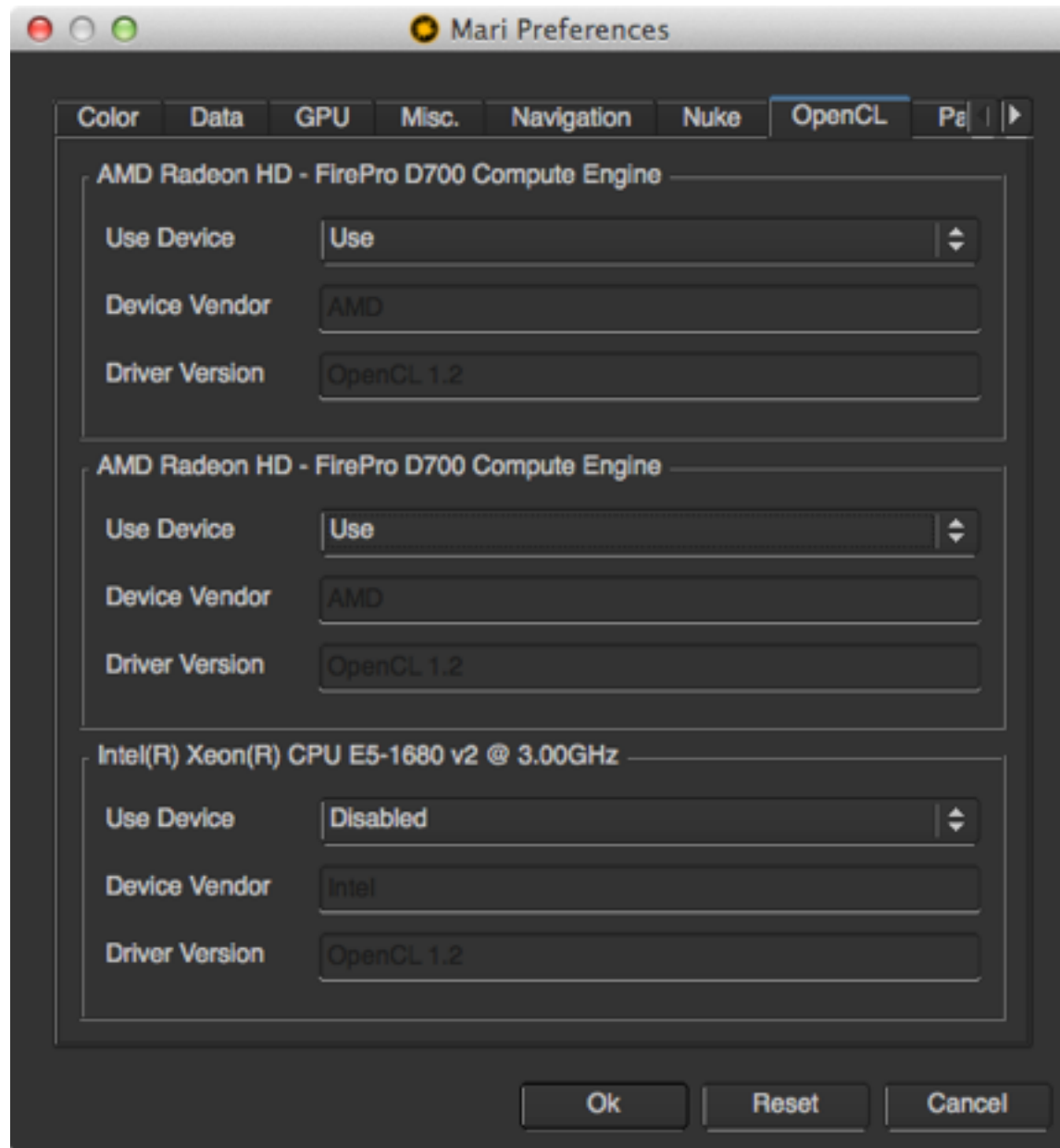


# Which one?



# Heat is a concern!



# Short Discussion of Parallelism

- Types of Parallelism
  - task parallelism
  - data parallelism

# Task Parallelism

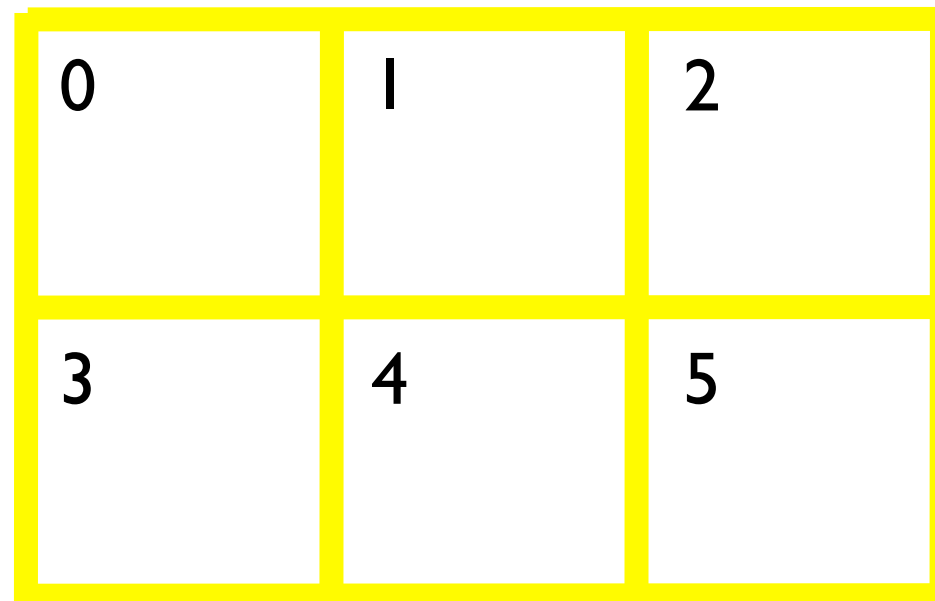
- Simple to understand
  - make a list of things to do
  - you have a set of  $T$  tasks
  - you have a set of  $P$  processors
  - give each task to a separate machine
  - time  $O(T/P)$

# Amdahl's Law

- Scaling  $O(T/P)$  is a lie
- Assume  $P$  is infinite, time can't be zero
  - always serial costs: startup, I/O, interprocess communication (MPI), ...
- WARNING, file system (NFS) is a serial resource, shared by **everyone**
- write to local /tmp or /scratch when possible

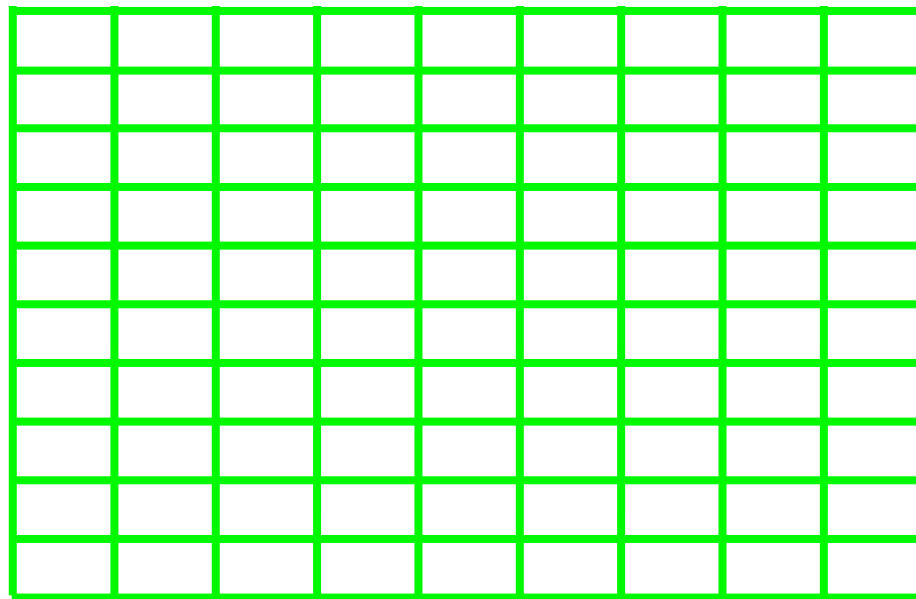
# Processor Topology

- A set of  $P$  processes can be arranged in a virtual grid of  $M_p * N_p = P$  processes
- 2 rows by 3 columns = 6 processes



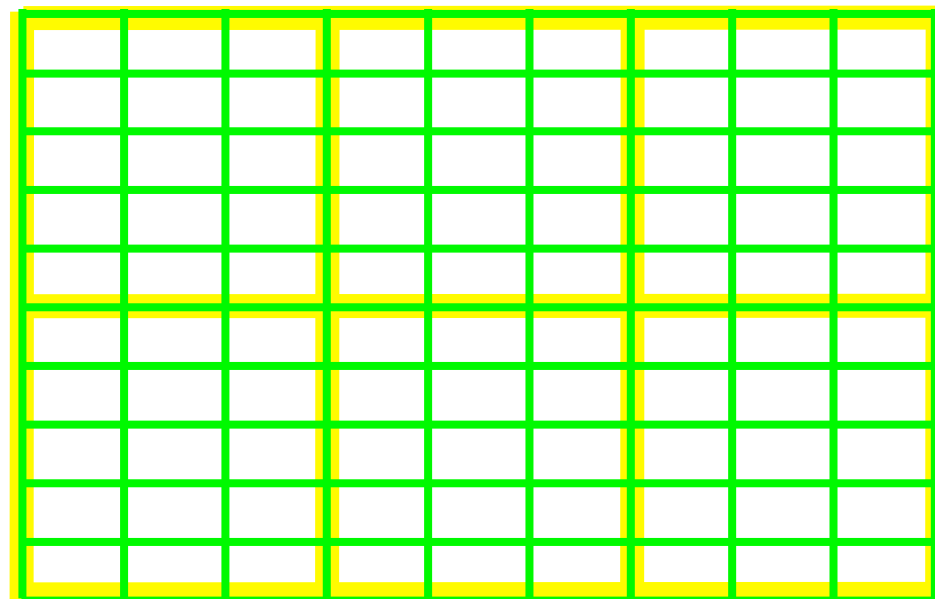
# Data Grid Topology

- Cells of data can be arranged in a grid of  $M \times N$   
= L cells
- 10 rows by 9 columns = 90 cells



# Data Parallelism

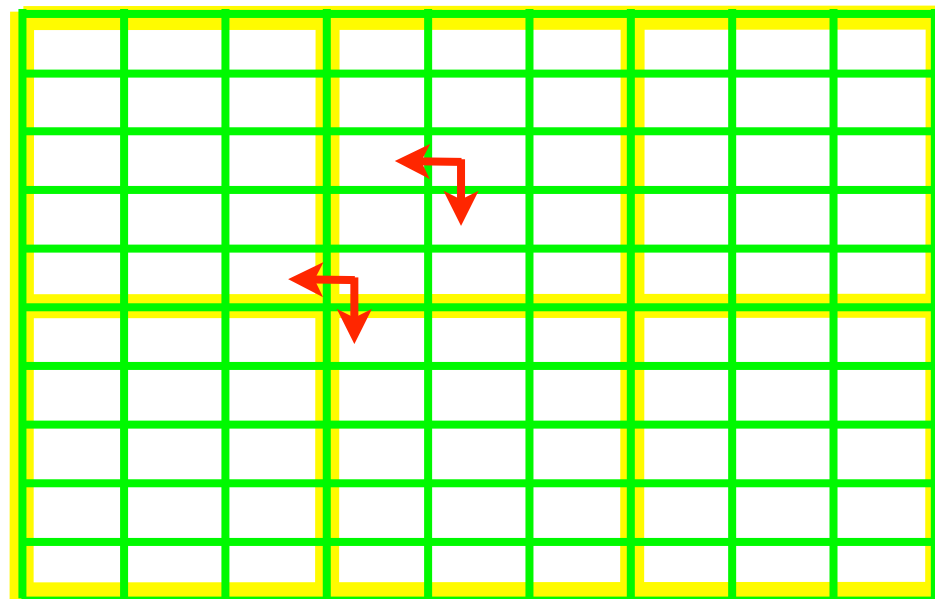
- Distribute a group of cells to a processor
- Scales well (up to 100 K + processors)
  - weak scaling: as cell count increases, use more processors
- 5 rows by 3 columns = 15 cells per processor





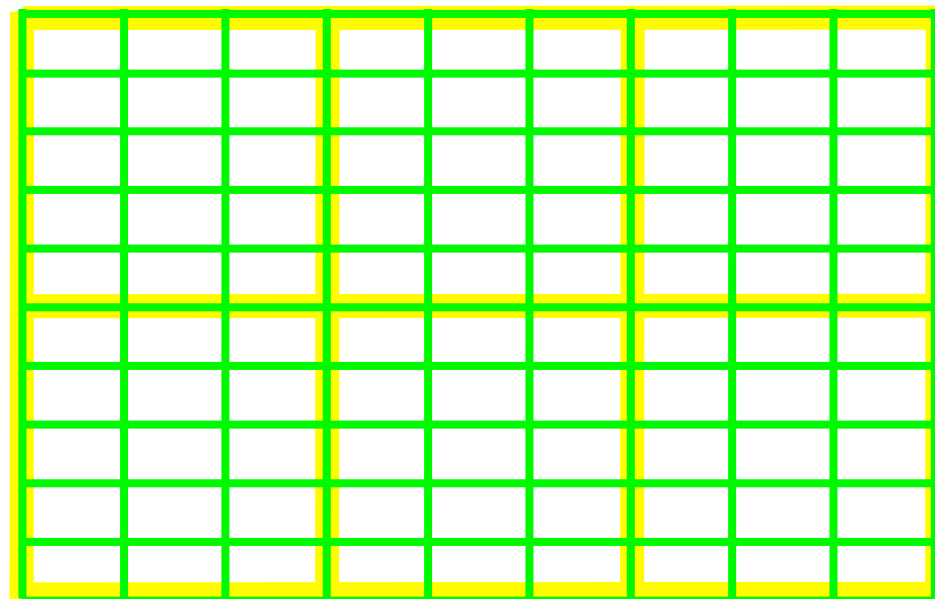
# Data Parallelism: Communication

- Communication is usually between neighboring processors
  - communication costs scale as processor surface
  - computation done on volume (# of cells)
- Global communication often scales  $O(\log(P))$ 
  - calculate an average over all cells (e.g., average temperature)



# Data Parallelism Scaling

- Weak scaling beats Amdahl's law because number of cells per processor is constant
- Unless global communication, communication costs are also fixed



# Access to ACISS

- Sign up for ACISS:
  - <http://prodigal.nic.uoregon.edu:4063/newuser>
- Try simple commands on ACISS
  - <https://blogs.uoregon.edu/casspr/2012/05/25/mpi-workshop-simple-task/>