



Eötvös Loránd University
Faculty of Informatics

PROGRAMMING

2019/20 1st semester

Lecture 4



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Content

- › More Patterns of Algorithms (PoA)
 - Copy – calculation with function
 - Multiple item selection
 - Partitioning
 - Intersection
 - Union
- › PoAs – summary



Pattern of Algoritm

What is PoA? It is the general solution of a typical programming task.

- sequence → single value
- sequence → sequence
- sequence → sequences
- sequences → sequence

7. Copy – function evaluation

Example tasks:

- › Let's define the absolute values of **all elements in an array!**
- › Let's convert **all letters** of a text to lowercase!
- › Let's **calculate** the sum of two vectors!
- › Let's calculate $\sin(x)$ values for a given x **series!**
- › We know the ordinal number of N months, let's give their names!



7. Copy – function evaluation

What is common?

We have N „something”, and we have to assign N other things to these. The type of assigned values can differ from the type of original values, but the count (N) remains the same, as well as the order.





7. Copy – function evaluation

Specification:

Input: $N \in \mathbb{N}$, $X_{1..N} \in \mathbb{S}_1^N$

$f : \mathbb{S}_1 \rightarrow \mathbb{S}_2$

Output: $Y_{1..N} \in \mathbb{S}_2^N$

Precondition: –

Postcondition: $\forall i (1 \leq i \leq N) : Y_i = f(X_i)$

another way: $Y_{1..N} = f(X_{1..N})$



7. Copy – function calculation

Algorithm:

Note: It's not a must to use the same index for both arrays. Eg.

Postcondition: $\forall i \ (1 \leq i \leq N) : Y_{p(i)} = f(X_i)$

$p(i)$ could be $2*i$ or $N-i+1$
(defined with the needed array size and index-interval)

Copy

$i=1..N$

$Y[i]:=f(X[i])$

Copy2

$i=1..N$

$Y[p(i)]:=f(X[i])$

7. Copy – function evaluation

Specification:

Input: $N \in \mathbb{N}$, $X_{1..N} \in \mathbb{S}^N$
 $f: \mathbb{S}_1 \rightarrow \mathbb{S}_2$

Output: $Y_{1..N} \in \mathbb{S}_2^N$

Precondition: –

Postcondition: $\forall i (1 \leq i \leq N) : Y_i = f(X_i)$

Specification (a frequent special case)₁:

Input: $N \in \mathbb{N}$, $X_{1..N} \in \mathbb{S}^N$

$g: \mathbb{S} \rightarrow \mathbb{S}$, $A: \mathbb{S} \rightarrow \mathbb{L}$

f is often a function with
a conditional statement

Output: $Y_{1..N} \in \mathbb{S}^N$

Precondition: –

Postcondition: $\forall i (1 \leq i \leq N) : Y_i = f(X_i)$

definition: $f(x) = \begin{cases} g(x), & \text{if } A(x) \\ x, & \text{otherwise} \end{cases}$

7. Copy – function evaluation

Specification:

Input: $N \in \mathbb{N}$, $X_{1..N} \in \mathbb{S}^N$
 $f: \mathbb{S}_1 \rightarrow \mathbb{S}_2$

Output: $Y_{1..N} \in \mathbb{S}_2^N$

Precondition: –

Postcondition: $\forall i (1 \leq i \leq N) : Y_i = f(X_i)$

Specification (a frequent special case)₁:

Input: $N \in \mathbb{N}$, $X_{1..N} \in \mathbb{S}^N$

$g: \mathbb{S} \rightarrow \mathbb{S}$, $A: \mathbb{S} \rightarrow \mathbb{L}$

Output: $Y_{1..N} \in \mathbb{S}^N$

Precondition: –

Postcondition:

$\forall i (1 \leq i \leq N) : (A(X_i) \rightarrow Y_i = g(X_i)) \text{ and}$

not $A(X_i) \rightarrow Y_i = X_i$

the algorithm can be
generated automatically

7. Copy – function calculation

Algorithm₁:

Specification (a frequent special case)₁:

Input: $N \in \mathbb{N}$, $X_{1..N} \in \mathbb{S}^N$

$g: \mathbb{S} \rightarrow \mathbb{S}$, $A: \mathbb{S} \rightarrow \mathbb{L}$

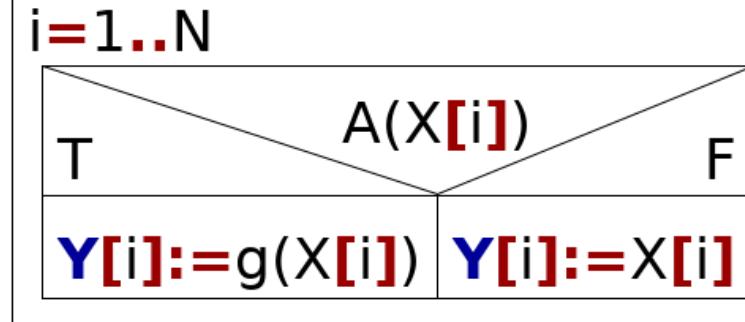
Output: $Y_{1..N} \in \mathbb{S}^N$

Precondition: –

Postcondition:

$\forall i (1 \leq i \leq N) : (A(X_i) \rightarrow Y_i = g(X_i) \text{ and}$
 $\text{not } A(X_i) \rightarrow Y_i = X_i)$

CopySpecial



7. Copy – function evaluation

Specification:

Input: $N \in \mathbb{N}$, $X_{1..N} \in \mathbb{S}_1^N$
 $f: \mathbb{S}_1 \rightarrow \mathbb{S}_2$

Output: $Y_{1..N} \in \mathbb{S}_2^N$

Precondition: –

Postcondition: $\forall i (1 \leq i \leq N) : Y_i = f(X_i)$

Specification (another special case)₂:

Input: $N \in \mathbb{N}$, $X_{1..N} \in \mathbb{S}^N$

Output: $Y_{1..N} \in \mathbb{S}^N$

Precondition: –

Postcondition: $\forall i (1 \leq i \leq N) : Y_i = X_i$

Note: there is no f function, or more precisely it is identical ($f(x) := x$)

7. Copy – function calculation

Algorithm:

Note: It can be replaced by the $Y := X$ value assignment, if the two arrays have the same size.

Except when the indexes are different.

CopySpecial2

1..N

Y[i]:=X[i]

CopySpecial2

1..N

Y[p(i)]:=X[i]

Let's calculate the sum of two vectors!



7. Copy – function evaluation

Specification:

Input: $N \in \mathbb{N}$, $P_{1..N}, Q_{1..N} \in \mathbb{R}^N$

$f : \mathbb{R} \times \mathbb{R} \rightarrow \mathbb{R}$, $f((p_i, q_i)) := p_i + q_i$

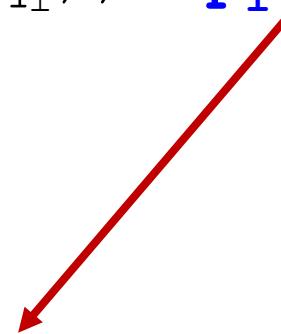
$(P, Q) \in (\mathbb{R} \times \mathbb{R})^N$

Output: $R_{1..N} \in \mathbb{R}^N$

Precondition: –

Postcondition:

$$\forall i (1 \leq i \leq N) : R_i = p_i + q_i$$



Algorithm:

```
Copy
i=1..N
Y[i]:=f(X[i])
```

8. Multiple item selection

Example tasks:

- › Let's **define** all excellent students in a class!
- › Let's **calculate** the divisors of an integer!
- › Let's **list** all the vowels from an English word!
- › Let's collect **all** the people who are taller than 180cm, from a set of people!
- › Let's **list** those days of the year when it didn't freeze at noon!

8. Multiple item selection

What is common?

We have to list **all** elements from N "something", which have a common attribute A.



8. Multiple item selection

Specification:

Input: $N \in \mathbb{N}$, $X_{1\dots N} \in \mathbb{S}^N$, $A: \mathbb{S} \rightarrow \mathbb{L}$

Output: Cnt $\in \mathbb{N}$, $Y_{1\dots N} \in \mathbb{N}^N$

Precondition: –

Postcondition: $Cnt = \sum_{i=1}^N A(x_i)$

Using only the first Cnt elements

See PoA „Counting”

and $\forall i (1 \leq i \leq \text{Cnt}) : A(X_{Y_i})$ and $Y \subseteq \{1, 2, \dots, N\}$

another way: $(\text{Cnt}, Y) = \text{multiselect i}^N$



8. Multiple item selection

Algorithm:

See PoA „Counting”

Comment: Collecting indexes is more general than collecting elements. If we needed elements then we would write: $Y[Cnt] := X[i]$ (and accordingly modified specification, as well – see later)

MultipleItemSelection

Cnt:=0

i=1..N

T

A(X[i])

F

Cnt:=Cnt+1

Y[Cnt]:=i

-

8. Multiple item selection

Specification₂ (selecting values):

Input: $N \in \mathbb{N}$, $X_{1..N} \in \mathbb{S}^N$, $A: \mathbb{S} \rightarrow \mathbb{L}$

Output: $Cnt \in \mathbb{N}$, $Y_{1..N} \in \mathbb{S}^N$

Precondition: –

Postcondition: $Cnt = \sum_{\substack{i=1 \\ A(X_i)}}^N 1$

and $\forall i (1 \leq i \leq Cnt) : A(Y_i)$ and $Y \subseteq X$

another way: $(Cnt, Y) = \text{mutiselect}_{i=1}^N X_i$

Let's list those days
of the year when it
didn't freeze at noon!

8. Multiple item selection

Specification:

Input: $N \in \mathbb{N}$, $T_{1..N} \in \mathbb{R}^N$

$\text{Pos} : \mathbb{R} \rightarrow \mathbb{L}$, $\text{Pos}(x) = (x > 0)$

Output: $\text{Cnt} \in \mathbb{N}$, $F_{1..N} \in \mathbb{R}^N$

Precondition: –

Postcondition: $\text{Cnt} = \sum_{i=1}^N \mathbf{1}_{T_i > 0}$

Input: $N \in \mathbb{N}$, $X_{1..N} \in \mathbb{S}^N$, $A : \mathbb{S} \rightarrow \mathbb{L}$

Output: $\text{Cnt} \in \mathbb{N}$, $Y_{1..N} \in \mathbb{N}^N$

Precondition: –

Postcondition: $\text{Cnt} = \sum_{i=1}^N \mathbf{1}_{A(X_i)}$

and $\forall i (1 \leq i \leq \text{Cnt}) : A(X_{y_i})$ and $Y \subseteq \{1, 2, \dots, N\}$

and $\forall i (1 \leq i \leq \text{Cnt}) : T_{NF_i} > 0$ and $NF \subseteq \{1, 2, \dots, N\}$



Let's list those days of the year when it didn't freeze at noon!

8. Multiple item selection

Algorithm:

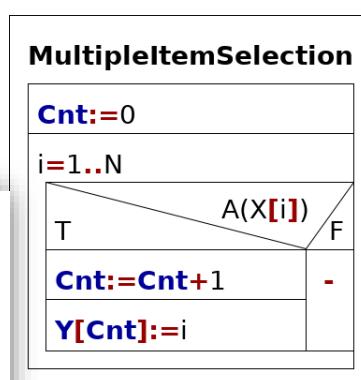
Input: $N \in \mathbb{N}$, $X_{1..N} \in \mathbb{S}^N$, $A: \mathbb{S} \rightarrow \mathbb{L}$

Output: $Cnt \in \mathbb{N}$, $Y_{1..N} \in \mathbb{N}^N$

Precondition: –

Postcondition: $Cnt = \sum_{i=1}^N 1_{A(X_i)}$

and $\forall i (1 \leq i \leq Cnt) : A(X_{y_i})$ and $Y \subseteq (1, 2, \dots, N)$



Specification:

Input: $N \in \mathbb{N}$, $T_{1..N} \in \mathbb{R}^N$

Pos: $\mathbb{R} \rightarrow \mathbb{L}$, $Pos(x) = (x > 0)$

Output: $Cnt \in \mathbb{N}$, $F_{1..N} \in \mathbb{R}^N$

Precondition: –

Postcondition: $Cnt = \sum_{i=1}^N 1_{T_i > 0}$

and $\forall i (1 \leq i \leq Cnt) : T_{NF_i} > 0$ and $NF \subseteq (1, 2, \dots, N)$

NotFreezingDays

Cnt := 0

i = 1 .. N



Cnt := Cnt + 1

NF[Cnt] := i



8. Local selection

Specification:

Input: $N \in \mathbb{N}$, $x_{1..N} \in \mathbb{S}^N$, $A: \mathbb{S} \rightarrow \mathbb{L}$

Output: $Cnt \in \mathbb{N}$, $x'_{1..N} \in \mathbb{S}^N$

Precondition: –

Postcondition: $Cnt = \sum_{i=1}^N 1_{A(x_i)}$

and $\forall i (1 \leq i \leq Cnt) : A(x'_{\textcolor{red}{i}})$ and $x'_{1..Cnt} \subseteq x$

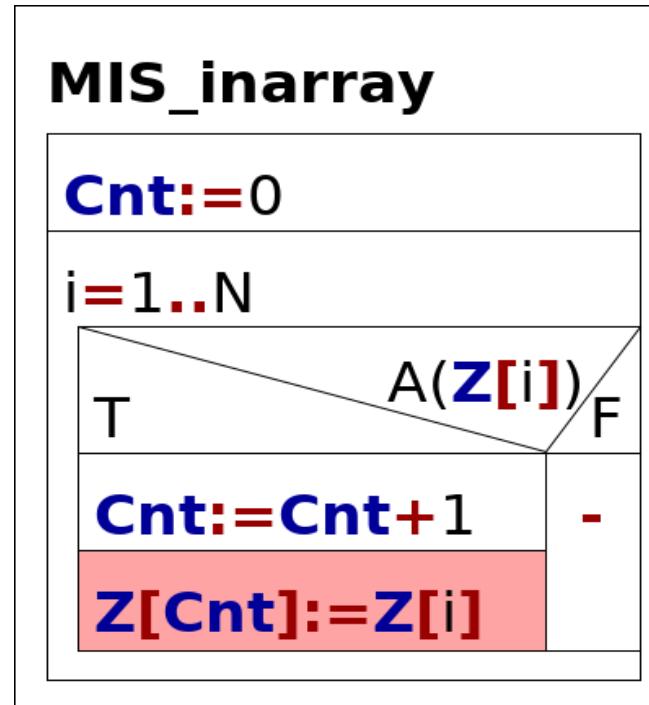
We can use the same variable for x and x' in our programm code.

$x_{1..Cnt}^{\text{output}} \subseteq x_{1..N}^{\text{input}}$ and $\forall i (1 \leq i \leq Cnt) : A(x_{\textcolor{red}{i}})$

8. Local selection

Main idea: we put the selected item to a place which we don't use later.

Algorithm:



9. Partitioning

Example tasks:

- › Let's list all **even** and all **odd** numbers from a series!
- › Let's list those days of the year when it was **freezing** and when it **wasn't** at noon!
- › Let's list all the **vowels** and **consonants** of an English word!
- › Let's collect all the people who are **shorter** than 140cm, who are **between** 140-180 cm, and those who are **taller** than 180cm, from a set of people!
- › From a group of people, **list** the people who were born in **summer**, in **autumn**, in **winter** and in **spring**.

9. Partitioning

What is common?

We have to list all elements from N "something", which have a common attribute A, and then also list those ones not having attribute A. So we “assign” all the elements of the input to one of the output sequences.

9. Partitioning

Specification:

Input: $N \in \mathbb{N}$, $X_{1..N} \in \mathbb{S}^N$, $A: \mathbb{S} \rightarrow \mathbb{L}$

Output: $Cnt \in \mathbb{N}$, $Y_{1..N} \in \mathbb{N}^N$, $Z_{1..N} \in \mathbb{N}^N$

Precondition: –

Postcondition: $Cnt = \sum_{\substack{i=1 \\ A(X_i)}}^N 1$

and $\forall i (1 \leq i \leq Cnt) : A(X_{Y_i})$

and $\forall i (1 \leq i \leq N - Cnt) : \text{not } A(X_{Z_i})$

and $Y \subseteq \{1, 2, \dots, N\}$ and $Z \subseteq \{1, 2, \dots, N\}$

$Y_{1..N} \in \mathbb{N}^{Cnt}$,
 $Z_{1..N} \in \mathbb{N}^{N-Cnt}$

9. Partitioning

Specification₂:

...

Postcondition₂: $(\text{Cnt}, \text{Y}, \text{Z}) = \text{partition}_{\sum_{i=1}^N A(x_i)}$

... **for values:** $(\text{Cnt}, \text{Y}, \text{Z}) = \text{partition}_{X_i}$
 $\sum_{i=1}^N A(x_i)$

9. Partitioning

Algorithm:

Input: $N \in \mathbb{N}$, $X_{1..N} \in \mathbb{S}^N$, $A: \mathbb{S} \rightarrow \mathbb{L}$
Output: $Cnt \in \mathbb{N}$, $Y_{1..N} \in \mathbb{N}^N$, $Z_{1..N} \in \mathbb{N}^N$
Precondition: –
Postcondition: $Cnt = \sum_{i=1}^N 1_{A(X_i)}$
 and $\forall i (1 \leq i \leq Cnt) : A(X_{Y_i})$
 and $\forall i (1 \leq i \leq N - Cnt) : \text{not } A(X_{Z_i})$
 and $Y \subseteq \{1, 2, \dots, N\}$ and $Z \subseteq \{1, 2, \dots, N\}$

Partitioning

Cnt:=0

i=1..N

T

A(X[i])

F

Cnt:=Cnt+1 **Z[i-Cnt]:=i**

Y[Cnt]:=i

Comment:

If we needed the **values**, than we would write **:=x[i]** instead of **:=i** (and accordingly specification should be modified).

9. Partitioning

Modification: There are exactly N number of elements in Y and Z together, so it's sufficient to have only one output array. We will collect the A attributed elements at the beginning of the output array Y.

Specification:

Input: $N \in \mathbb{N}$, $X_{1..N} \in \mathbb{S}^N$, $A: \mathbb{S} \rightarrow \mathbb{L}$

Output: $Cnt \in \mathbb{N}$, $Y_{1..N} \in \mathbb{N}^N$

Precondition: –

Postcondition: $Cnt = \sum_{\substack{i=1 \\ A(X_i)}}^N 1$

and $\forall i (1 \leq i \leq Cnt) : A(X_{Y_i})$

and $\forall i (Cnt+1 \leq i \leq N) : \text{not } A(X_{Y_i})$

and $Y \in \text{Permutation}(1, 2, \dots, N)$

Permutation(1,2,...,N):=set of all permutations of 1..N (set of sets)

9. Partitioning

Specification₂:

...

Postcondition₂: $(\text{Cnt}, Y) = \text{partition}_2 \begin{matrix} N \\ i=1 \\ A(x_i) \end{matrix}$

... for values: $(\text{Cnt}, Y) = \text{partition}_2 \begin{matrix} X_i \\ i=1 \\ A(x_i) \end{matrix}$

9. Partitioning

Algorithm:

Input: $N \in \mathbb{N}$, $X_{1..N} \in \mathbb{S}^N$, $A: \mathbb{S} \rightarrow \mathbb{L}$

Output: $Cnt \in \mathbb{N}$, $Y_{1..N} \in \mathbb{N}^N$

Precondition: –

Postcondition: $Cnt = \sum_{\substack{i=1 \\ A(X_i)}}^N 1$

and $\forall i (1 \leq i \leq Cnt) : A(X_{y_i})$

and $\forall i (Cnt + 1 \leq i \leq N) : \text{not } A(X_{y_i})$

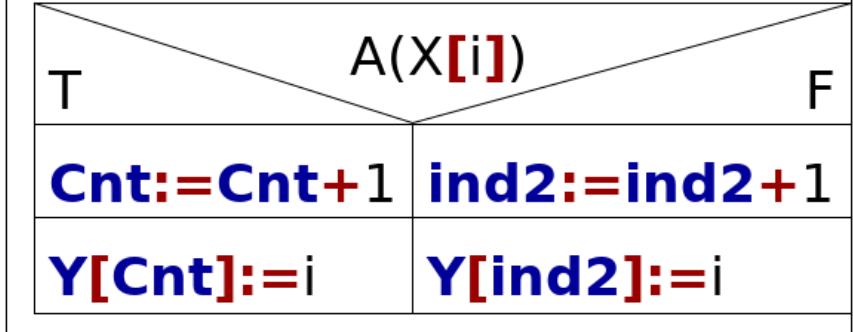
and $Y \in \text{Permutation}(1, 2, \dots, N)$

Partitioning2

Cnt:=0

ind2:=N+1

i=1..N



Comment:

We can use an extra variable (ind2) to know where we are from back in Y.

9. Local partitioning

Specification:

Input: $N \in \mathbb{N}$, $x_{1..N} \in \mathbb{S}^N$, $A: \mathbb{S} \rightarrow \mathbb{L}$

Output: $Cnt \in \mathbb{N}$, $x'_{1..N} \in \mathbb{S}^N$

Precondition: –

Postcondition: $Cnt = \sum_{i=1}^N \mathbf{1}_{A(x_i)}$

and $\forall i (1 \leq i \leq Cnt) : A(x'_i)$

and $\forall i (Cnt+1 \leq i \leq N) : \text{not } A(x'_i)$

and $x' \in \text{Permutation}(X)$

We can use the same variable for x and x' in our programm code.

9. Local partitioning

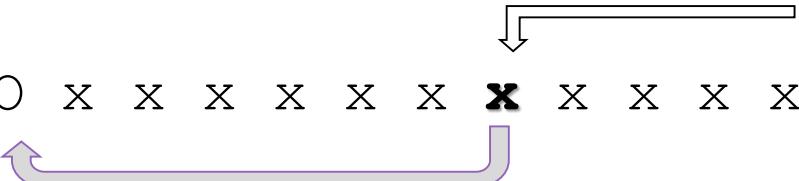
Main idea for the algorithm:

1. Pick out the **first** item of the sequence

O x x x x x x x x x x

2. Let's search an item from **back** which needs to be in the **front-sequence** (attribute A is true for it)

O x x x x x x **x** x x x x



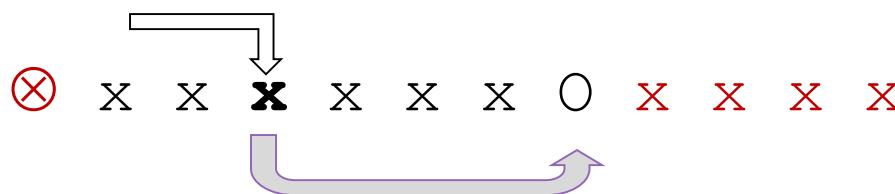
3. Move the found item into the empty (first) place.

⊗ x x x x x x O x x x x

The first item and the items after the new empty place are at the right place.

9. Local partitioning

4. Now we have an empty place at the end-sequence. From the **beginning** (starting after the first item) let's search an item which needs to be in the end-sequence (attribute A is not true for it)



5. Move the founded item into the empty place. And we can search from back away.

$\text{⊗ } \text{x } \text{x } \text{o } \text{x } \text{x } \text{⊗ } \text{x } \text{x } \text{x } \text{x }$

The items before the new empty place and after the new place of the moved item are at right place.

9. Local partitioning

6. ... and so on ...

7. We can finish searching if we reached the empty place from one direction.

x x x x x o x x x x x x

8. We put our first item back into this empty place.

9. Local partitioning

Algorithm:

What we know about X?

- At the beginning: it is the input sequence
- At the end: it is the permutation of the original (input) sequence
- While „running”: to „first” are items with attribute A, from „last” items which are not of attribute A.

Partitioning_onearray

first:=1 [the **first** of the items to be partitioned]

last:=N [the **last** of the items to be partitioned]

y:=X[first]

first < last

SearchFromBack(**first, last, Exists**)

T

Exists

F

X[first]:=X[last]

first:=first+1

SearchFromFront(**first, last, Exists**)

T

Exists

F

X[last]:=X[first]

last:=last-1

X[first]:=y

T

A(y)

F

Cnt:=first

Cnt:=first-1

9. Local partitioning

SearchFromFront(first,last:Integer,Exists:Boolean)

first<last and A(X[first])

first:=first+1

Exists:=first<last

SearchFromBack(first,last:Integer,Exists:Boolean)

first<last and not A(X[last])

last:=last-1

Exists:=first<last

10. Intersection

Example tasks:

- › Let's list all **common** divisors of **two** integers!
- › We investigate birds in winter **and** summer. Let's **collect** the birds which are not migratory!
- › Based on the free hours from the calendars of **two people**, let's **determine** when they can meet.
- › Let's lists the animals which could be seen **both** in the zoos of Budapest and Veszprém.

10. Intersection

What is common?

We have two sets as input (with elements of the same type), and we have to list all elements that are part of both sets.

10. Intersection

Definition of IsSet function:

$\text{IsSet}(x_{1..n}) := \text{not } \exists i \ (1 \leq i \leq n) : x_i \in x_{1..i-1}$
 $\text{IsSet}(x_{1..n}) := i \neq j \rightarrow x_i \neq x_j$

Specification:

Input: $N, M \in \mathbb{N}, X_{1..N} \in \mathbb{S}^N, Y_{1..M} \in \mathbb{S}^M$

Output: $Cnt \in \mathbb{N}, Z_{1..\min(N,M)} \in \mathbb{S}^{\min(N,M)}$

Precondition: $\text{IsSet}(X)$ and $\text{IsSet}(Y)$

Postcondition: $Cnt = \sum_{\substack{i=1 \\ x_i \in Y}}^N 1$

and $\forall i \ (1 \leq i \leq Cnt) : Z_i \in X$ and $Z_i \in Y$

and $\text{IsSet}(Z)$

10. Intersection

Specification₂:

Postcondition₂: $(\text{Cnt}, Z) = \text{Intersection}(N, X, M, Y)$

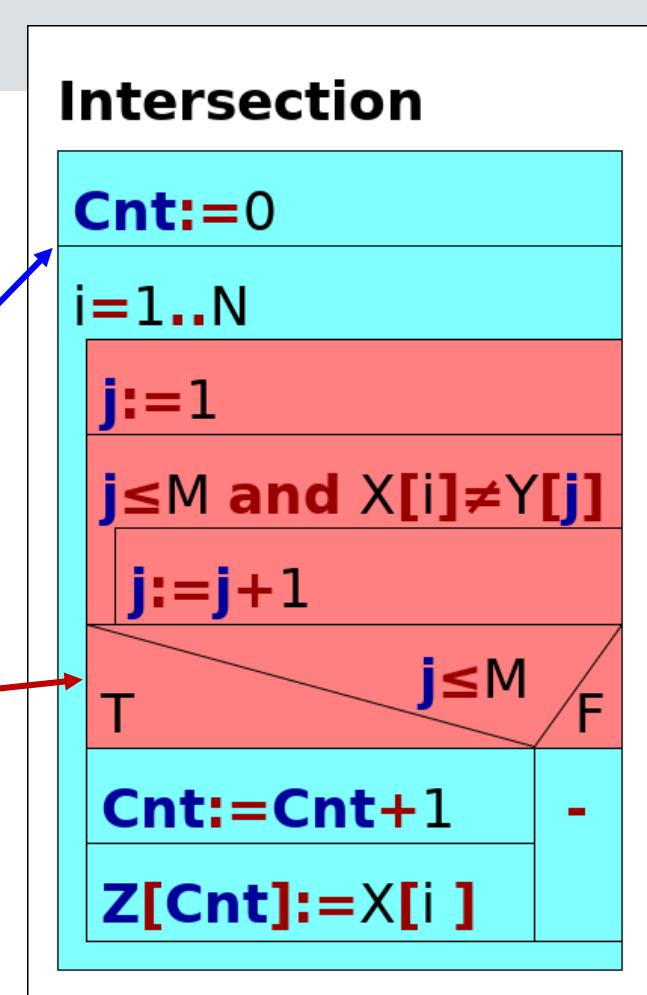
another way: $(\text{Cnt}, Z) = \text{multiselect } X_i$
 $i=1$
 $X_i \in Y$

10. Intersection

Algorithm:

Comment:

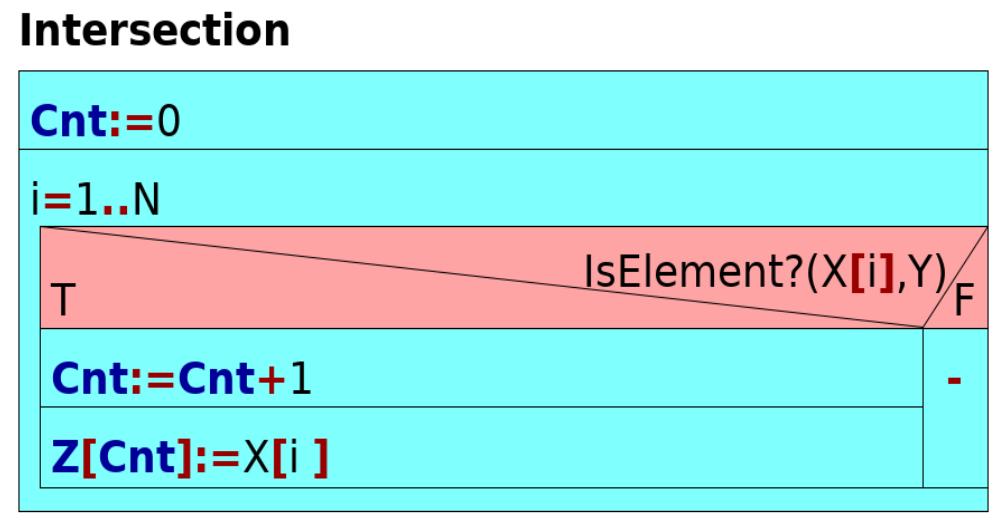
The solution is a **multiple item selection** and a **decision.**



10. Intersection

As the decision PoA gives back a logical value (boolean), it could appear as a logical expression in the conditional statement

Algorithm:



We need to write a function (IsElement)!

10. Intersection

Variations for intersection:

- › We have two sets, we have to determine the **count of common elements**.
- › We have two sets, we have to determine if they have **at least one common element**.
- › We have two sets, we have to give **one common element**.

11. Union

Example tasks:

- › We have two courses, let's list the students visiting **both** courses!
- › Let's collect all the birds we can investigate in winter **as well as** in summer!
- › Based on the free hours from the calendars of two persons, let's **determine** when we can reach **at least one of them**!
- › Let's **lists** the animals which could be seen **either** in the zoos of Budapest or Veszprém.

11. Union

What is common?

We have two sets as input (with elements of the same type), and we have to list all elements that are included at least in one of the sets.

11. Union

Specification:

Input: $N, M \in \mathbb{N}$, $X_{1..N} \in \mathbb{S}^N$, $Y_{1..M} \in \mathbb{S}^M$

Output: $Cnt \in \mathbb{N}$, $Z_{1..N+M} \in \mathbb{S}^{N+M}$

Precondition: $\text{IsSet}(X)$ and $\text{IsSet}(Y)$

Postcondition: $Cnt = N + \sum_{\substack{j=1 \\ Y_j \notin X}}^M 1$

and $\forall i (1 \leq i \leq Cnt) : Z_i \in X \text{ or } Z_i \in Y$

and $\text{IsSet}(Z)$

11. Union

Specification₂:

Postcondition₂: $(\text{Cnt}, \text{Z}) = \text{Union}(\text{N}, \text{X}, \text{M}, \text{Y})$

another way: $(\text{Cnt}, \text{Z}) = \text{x} + \underset{\substack{i=1 \\ \text{Y}_j \notin \text{x}}}{\text{multiselect}} \text{ Y}_j$

11. Union Algorithm

Specification:

Input: $N, M \in \mathbb{N}$, $X_{1..N} \in \mathbb{S}^N$, $Y_{1..M} \in \mathbb{S}^M$

Output: $Cnt \in \mathbb{N}$, $Z_{1..N+M} \in \mathbb{S}^{N+M}$

Precondition: $\text{IsSet}(X)$ and $\text{IsSet}(Y)$

Postcondition: $Cnt = N + \sum_{\substack{j=1 \\ Y_j \notin X}}^M 1$

and $\forall i (1 \leq i \leq Cnt) : Z_i \in X \text{ or } Z_i \in Y$
and $\text{IsSet}(Z)$

multiple item selection

Union

Z:=X; Cnt:=N

j=1..M

i:=1

i≤N and X[i]≠Y[j]

i:=i+1

i>N

Cnt:=Cnt+1

Z[Cnt]:=Y[j]

copy

decision

11. Union

Variations for union:

- › We have two sets, we have to determine the count of all the elements altogether.
- › We have two sets, we have to determine their difference ($X \setminus Y$).
- › We have two sets, we have to give those elements which are only in one of the sets ($(X \setminus Y) \cup (Y \setminus X)$).

Patterns of Algorithms II.

Sequence → Sequence

- Copy – calculation with function
- Multiple item selection
- Sorting (later)

Sequence → Sequences

- Partitioning

Sequences → Sequence

- Intersection
- Union