



Eötvös Loránd University  
Faculty of Informatics

# PROGRAMMING

2019/20 1<sup>st</sup> semester

## Lecture 6



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# Content

- › Function in condition
- › Stacking compound data structures
  - Matrix – vector of vectors
  - Vector of records
- › Transforming to set
- › The SET type



# Function in condition



# Function in condition

**Task:** Let's give a vowel from an English word.

**Specification** (selection):

**Input:**  $N \in \mathbb{N}$ , Word  $\in \mathbb{C}^h^N$

**Output:**  $v_w \in \mathbb{N}$

**Precondition:**  $N > 0$  and  $\forall i (1 \leq i \leq N) : \text{IsVowel}(\text{Word}_i)$

**Postcondition:**  $1 \leq v_w \leq N$  and  $\text{IsVowel}(\text{Word}_{v_w})$

**Definition:**  $\text{IsVowel} : \mathbb{C}^h \rightarrow \mathbb{L}$ ,

$\text{IsVowel}(L) := \forall i (1 \leq i \leq 5) : L = \text{Vowels}_i$ ,  
 $\text{Vowels} \in \mathbb{C}^h^5 = ("a", "e", "i", "o", "u")$

# Function in condition

## Algorithm

### VowelInWord

**Vw:=1**

**not IsVowel(Word[Vw])**

**Vw:=Vw+1**

**IsVowel(L:Character):Boolean**

**i:=1**

**i≤5 and L≠Vowels[i]**

**i:=i+1**

**IsVowel:=i≤5**



# Matrix



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## Matrix - pictures

**Task:** Let's double the size of a raster image by pixel-multiplication!

This means each point will be enlarged to a  $2 \times 2$  same colored area.



## Matrix - pictures

How should we represent an image?

- › The picture consists of ordered pixels, so it can be represented with some kind of sequence.

However, it would be hard to assign a single sequence number to each pixel.

- › It's easier to define the **row** and the **column** number of a pixel, so why don't we use **double indexing**? The double indexed array is called "**matrix**".

# Matrix - pictures

## Specification

**Input:**  $N, M \in \mathbb{N}, P_{1..N, 1..M} \in \mathbb{N}^{N \times M}$

**Output:**  $DP_{1..2^i, 1..2^j} \in \mathbb{N}^{2^i \times 2^j}$

**Precondition:** –

**Postcondition:**

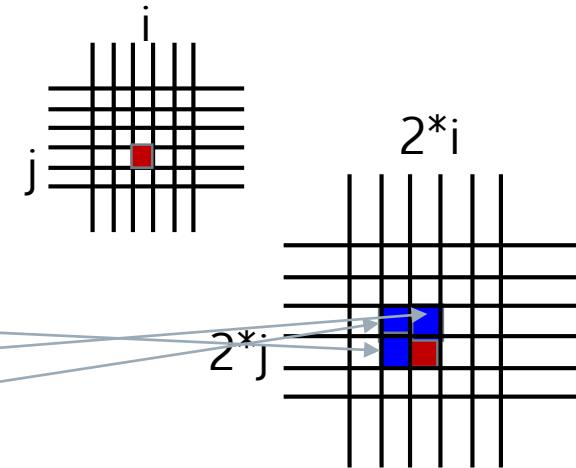
$\forall i (1 \leq i \leq N) : \forall j (1 \leq j \leq M) :$

$DP_{2^i, 2^j} = P_{i, j}$  and

$DP_{2^i-1, 2^j} = P_{i, j}$  and

$DP_{2^i-1, 2^j-1} = P_{i, j}$  and

$DP_{2^i, 2^j-1} = P_{i, j}$



A variation of the **copy** pattern, but from one item, 4 are created.

# Matrix - pictures

## Algorithm

**Comment:** In different programming languages, different notation may be required to access an element in a matrix  
e.g.: C++: `P[i][j]`

### DoublePicture

`i=1..N`

`j=1..M`

`DP[2*i,2*j]:=P[i,j]`

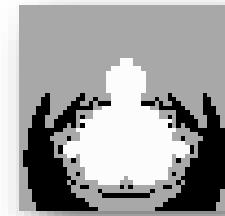
`DP[2*i-1,2*j]:=P[i,j]`

`DP[2*i,2*j-1]:=P[i,j]`

`DP[2*i-1,2*j-1]:=P[i,j]`

**Task:** Let's halve the size of a raster image by pixel-averaging (to size  $N/2 \times M/2$ ). This means each point of the small image will be the „average” of a 2x2 area

„average”: the average of color codes



# Matrix - pictures

## Specification

**Input:**  $N, M \in \mathbb{N}, P_{1..N, 1..M} \in \mathbb{N}^{N \times M}$

**Output:**  $S P_{1..N/2, 1..M/2} \in \mathbb{N}^{N/2 \times M/2}$

**Precondition:**  $\text{isEven}(N), \text{isEven}(M)$

**Postcondition:**

$$\forall i (1 \leq i \leq N/2) : \forall j (1 \leq j \leq M/2) :$$

$$S P_{i,j} = (P_{2*i, 2*j} + P_{2*i-1, 2*j} + P_{2*i, 2*j-1} + P_{2*i-1, 2*j-1}) / 4$$

**Definition:**  $\text{isEven} : \mathbb{N} \rightarrow \mathbb{L}$

$$\text{isEven}(x) := (x \bmod 2) = 0$$

# Matrix - pictures

## Algorithm

### HalvePicture

i=1..N/2

j=1..M/2

$SP[i,j] := (P[2*i, 2*j] + P[2*i-1, 2*j] + P[2*i, 2*j-1] + P[2*i-1, 2*j-1]) / 4$

### Comment:

1. With colored images, there might be a problem with averaging. What color is the average of a red and a blue pixel?
2. If using RGB the color is:  
Record (red, green, blue: Integer); and what will be the average?

## Matrix - pictures

**Task:** Let's apply one kind of a Rank-filter to the picture of Crab Nebula (astronomy). This rank filter substitutes each pixel with the maximum of itself and its adjacent 8 pixels.



# Matrix - pictures

**Specification (copy + maximum selection)**

**Input:**  $N, M \in \mathbb{N}, P_{1..N, 1..M} \in \mathbb{N}^{N \times M}$

**Output:**  $RP_{1..N, 1..M} \in \mathbb{N}^{N \times M}$

**Precondition:** –

**Postcondition:**  $\forall i (1 < i < N) : \forall j (1 < j < M) :$

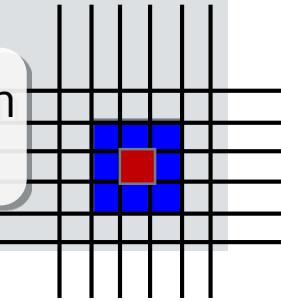
$$RP_{i,j} = \max_{p=i-1}^{i+1} \max_{q=j-1}^{j+1} P_{p,q}$$

and  $\forall j (1 \leq j \leq M) : RP_{1,j} = P_{1,j}$  and  $RP_{N,j} = P_{N,j}$

and  $\forall i (1 \leq i \leq N) : RP_{i,1} = P_{i,1}$  and  $RP_{i,M} = P_{i,M}$

# Matrix - pictures

See: maximum selection  
for value



## RankFilter

$i = 2..N-1$

$j = 2..M-1$

**Max:=0**

$p = i-1..i+1$

$q = j-1..j+1$

T

$P[p,q] > Max$

F

**Max:=P[p,q]**

**RP[i,j]:=Max**

## Algorithm

### RankFilter(cont...)

$j = 1..M$

**RP[1,j]:=P[1,j]**

**RP[N,j]:=P[N,j]**

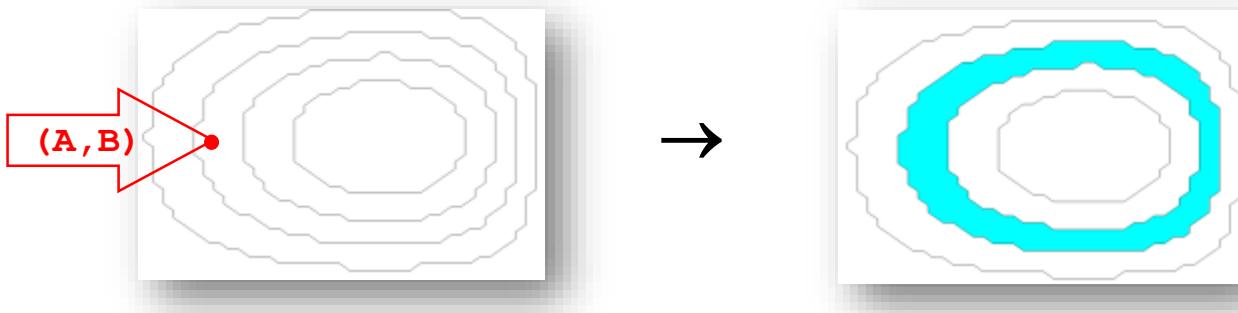
$i = 1..N$

**RP[i,1]:=P[i,1]**

**RP[i,M]:=P[i,M]**

# Matrix - pictures

**Task:** Let's color a range of white pixels of an image to light blue starting from an inner point (A,B).



**To be colored:** the „inner points”, if  $\text{Inner}(i, j) = \text{True}$  where

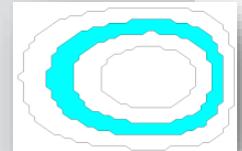
$\text{Inner}: \mathbb{N} \times \mathbb{N} \rightarrow \mathbb{L}$

$\text{Inner}(i, j) := (i = A \text{ and } j = B \text{ or }$

$\text{White}(i, j) \text{ and }$

$\text{Inner}(i-1, j) \text{ or } \text{Inner}(i+1, j) \text{ or } \text{Inner}(i, j-1)$

$\text{or } \text{Inner}(i, j+1)) )$



# Matrix - pictures

## Specification

**Input:**  $N, M \in \mathbb{N}$ ,  $P_{1..N, 1..M} \in \mathbb{N}^{N \times M}$ ,  $A, B \in \mathbb{N}$

**Output:**  $BP_{1..N, 1..M} \in \mathbb{N}^{N \times M}$

**Precondition:**  $A \in (1..N)$ ,  $B \in (1..M)$ ,  $P_{AB} = \text{white}$

**Postcondition:**  $\forall i (1 \leq i \leq N) : \forall j (1 \leq j \leq M) :$

$\text{Inner}(i, j) \rightarrow BP_{i,j} = \text{lightblue}$  and

not  $\text{Inner}(i, j) \rightarrow BP_{i,j} = P_{i,j}$

```
BP:=P  
Paint(A,B)
```

The  $BP$  matrix is global –  
in the function `Paint` is  
aviable

# Matrix - pictures

## Algorithm

**Paint(i,j:integer)**

**BP[i,j]:=lightblue**

T

**BP[i-1,j]=white and i>1**

F

Paint(i-1,j)

-

T

**BP[i+1,j]=white and i<N**

F

Paint(i+1,j)

-

T

**BP[i,j-1]=white and j>1**

F

Paint(i,j-1)

-

T

**BP[i,j+1]=white and j<M**

F

Paint(i,j+1)

-



# Vector (1D array) of records



## Vector of records - time

**Task:** There were N earthquakes on a given day. We know the exact times of each earthquake, in their original order in time. Let's give how much was the duration between earthquakes.

Towards the solution:

- › First, let's define time
- › We can define time by three values (hour, min, sec):

$$\text{Time} = H \times M \times S, \quad H, M, S = \mathbb{N}$$

- › Algorithmic pattern: copy

# Vector of records - time

## Specification

**Input:**  $N \in \mathbb{N}$ ,  $E_{1..N} \in \text{Time}^N$ ,

$\text{Time} = H \times M \times S$ ,  $H, M, S = \mathbb{N}$

**Output:**  $T_{1..N-1} \in \mathbb{N}^{N-1}$

**Precondition:**  $\forall i (1 \leq i \leq N) : 0 \leq E_i.h \leq 23$  and  
 $0 \leq E_i.m \leq 59$  and  $0 \leq E_i.s \leq 59$  and

$\forall i (1 \leq i < N) : E_i < E_{i+1}$

**Postcondition:**  $\forall i (1 \leq i \leq N-1) : T_i = E_{i+1} - E_i$

## Definition:

$- : \text{Time} \times \text{Time} \rightarrow \mathbb{N}$

$i1 - i2 := \dots ? \dots$

$< : \text{Time} \times \text{Time} \rightarrow \mathbb{L}$

$i1 < i2 := \dots ? \dots$

# Vector of records - time

## Difference of times

### 1. idea for solution:

We could regard it as the difference of a **3-digit** number where the **base** of each digit is **different**. Then we convert them to seconds.

### 2. idea for solution:

We express times in seconds, and then we could calculate the difference of two whole numbers.

```
inseconds(i) := i.h * 3600 + i.m * 60 + i.s
```

How big integer do we need? ( $24 * 3600 = 86\ 400$ ) What should be the type? (>2 byte)

# Vector of records - time

## Specification

**Input:**  $N \in \mathbb{N}$ ,  $E_{1..N} \in \text{Time}^N$ ,

$\text{Time} = H \times M \times S$ ,  $H, M, S = \mathbb{N}$

**Output:**  $T_{1..N-1} \in \mathbb{N}^{N-1}$

**Precondition:**  $\forall i (1 \leq i \leq N) : 0 \leq E_i.h \leq 23$  and  
 $0 \leq E_i.m \leq 59$  and  $0 \leq E_i.s \leq 59$  and

$\forall i (1 \leq i < N) : E_i < E_{i+1}$

**Postcondition:**  $\forall i (1 \leq i \leq N-1) : T_i = E_{i+1} - E_i$

## Definition:

$- : \text{Time} \times \text{Time} \rightarrow \mathbb{N}$

$$i_1 - i_2 := (i_1.h * 3600 + i_1.m * 60 + i_1.s) - (i_2.h * 3600 + i_2.m * 60 + i_2.s)$$

# Vector of records - time

## Algorithm

Postcondition:  $\forall i \ (1 \leq i \leq N-1) : T_i = E_{i+1} - E_i$

Definition:

-  $:Time \times Time \rightarrow \mathbb{N}$   
 $i_1 - i_2 := i_1.h * 3600 + i_1.m * 60 + i_1.s - (i_2.h * 3600 + i_2.m * 60 + i_2.s)$

### EarthquakeFreq

$i=1..N$

$S[i]:=E[i].h*3600+E[i].m*60+E[i].s$

$i=1..N-1$

$T[i]:=S[i+1]-S[i]$

- We use an **S** vector as a helper (auxiliary) object.
- The operator „-“ for Times is put indirectly to the algorithm through **S**.

# Vector for records - time

Algoritmh2:

## EarthquakeFreq2

i=1..N

S[i]:=inseconds(E[i])

i=1..N-1

T[i]:=S[i+1]-S[i]

### Comments:

- The `inseconds` function has to be implemented!
- If the difference is needed in (hour, minute, second) format, a back transformation is needed.

# Vector for records - time

## Algorithm3:

### EarthquakeFreq3

```
i=1..N
```

```
T[i]:=inseconds(E[i+1])-inseconds(E[i])
```

#### Comment:

By using the `inseconds` function, you can eliminate the `S` auxiliary vector and the loop for creating the `S` vector. However, we convert almost all `E[i]` to seconds twice.

# Vector for records - time

Algorithm4:

**EarthquakeFreq4**

$i = 1..N-1$

$T[i] := E[i+1] - E[i]$

**Comment:**

The operator „-“ must be defined, in which you could use the inseconds function or its body.



# Sequence to set transformation



## Sequence → set transformation

In some tasks, like the ones using the intersection and union pattern of algorithm, the starting data are in a **set**. If we get a sequence as input, we might need to make a set from it.

**Example:** We know about  $N$  shoppings which products customers bought ( $I_{n[1..N]}$ ). List the products bought by customers ( $P[1..Cnt]$ ).

**Solution:** **multiple item selection** programming pattern – select all the items from the input that have not been put in the result set of the selection yet (**decision**).

# Sequence → set transformation

## Algorithm

### SequenceToSet

**Cnt:=0**

**i=1..N**

T

$In[i] \notin P[1..Cnt]$

F

**Cnt:=Cnt+1**

**P[Cnt]:=In[i]**

-



# The SET Type



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# The SET Type

## Value set

- › The iteration of the base set (that is defined by the element type) – „*which items can be in the set?*”
- › The **element type** is usually a finite, discrete type, sometimes even the count of elements is limited (<256)
- › If it does not exist in the language, then the implementation might allow more elements

# The SET Type

## Operations (mathematical)

- › Intersection:  $A \cap B$
- › Union:  $A \cup B$
- › Difference:  $A \setminus B$
- › Complement:  $A'$  (not always implementable)
- › Element of set:  $a \in A$
- › Subset:  $A \subseteq B$ ,  $A \subset B$

# The SET Type

## Operations (implementation)

- › **ToSet** (puts an element into the set):  $S := S \cup \{ e \}$
- › **FromSet** (discards an element from the set):  
 $S := S \setminus \{ e \}$
- › **Read** (reading in the set)
- › **Write** (writing out the set)
- › **Empty** (creating an empty set) or **Empty'SetType** pre-defined constant
- › **Empty?** (boolean function)



# The SET Type

## Representation

- › Listing the elements
- › A boolean vector - bitmap



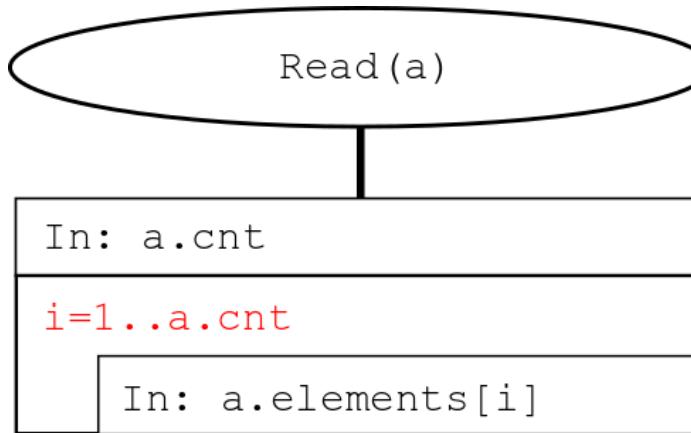
# The set type – by listing elements

Representation: by listing the elements

```
Set (ElementType) =  
  Record (cnt:integer,  
          elements:Array (1..MaxCnt:ElementType))
```

We give the set by listing its elements in an vector whose length is the same as the count of elements of the set (more precisely, in the first Cnt elements).

# The set type – by listing elements - READ

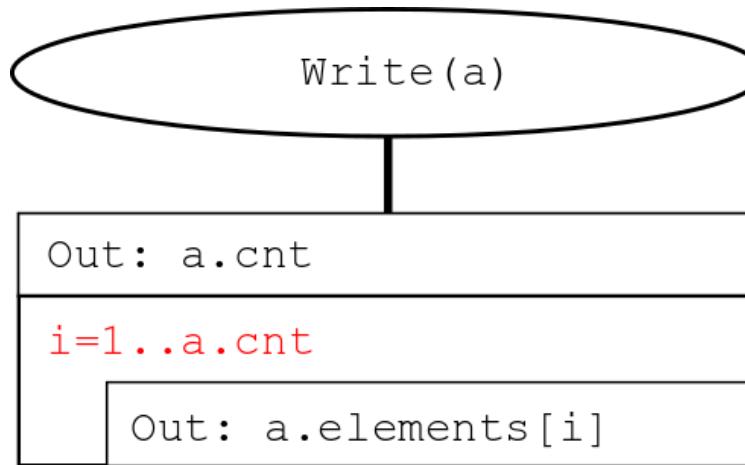


We assume that we read in a **set**.

**Calculation of the operation need:**

The loop will run as many times as many elements there are in the set – thus, the runtime is proportional to the count of elements of the set.

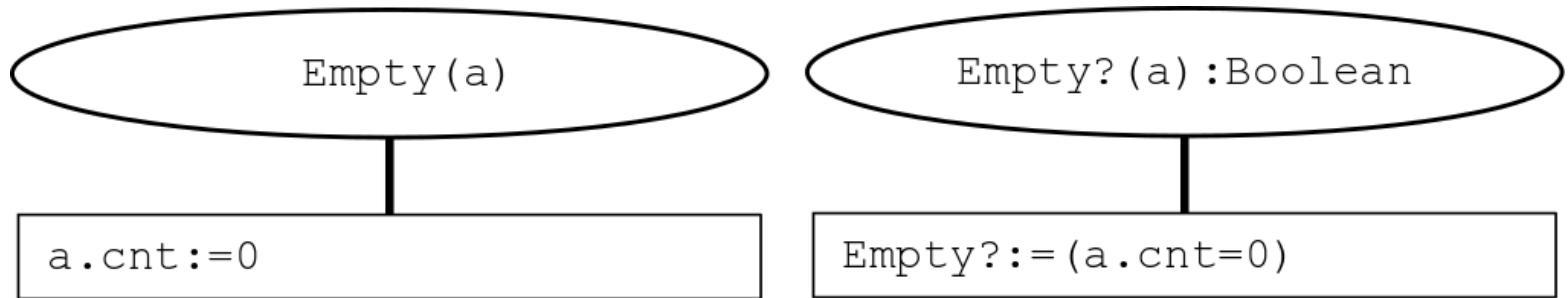
# The set type – by listing elements - WRITE



## Calculation of the operation need:

The loop will run as many times as many elements there are in the set – thus, the runtime is proportional to the count of elements of the set.

# The set type – by listing elements – **EMPTY,** **EMPTY?**



## Calculation of the operation need:

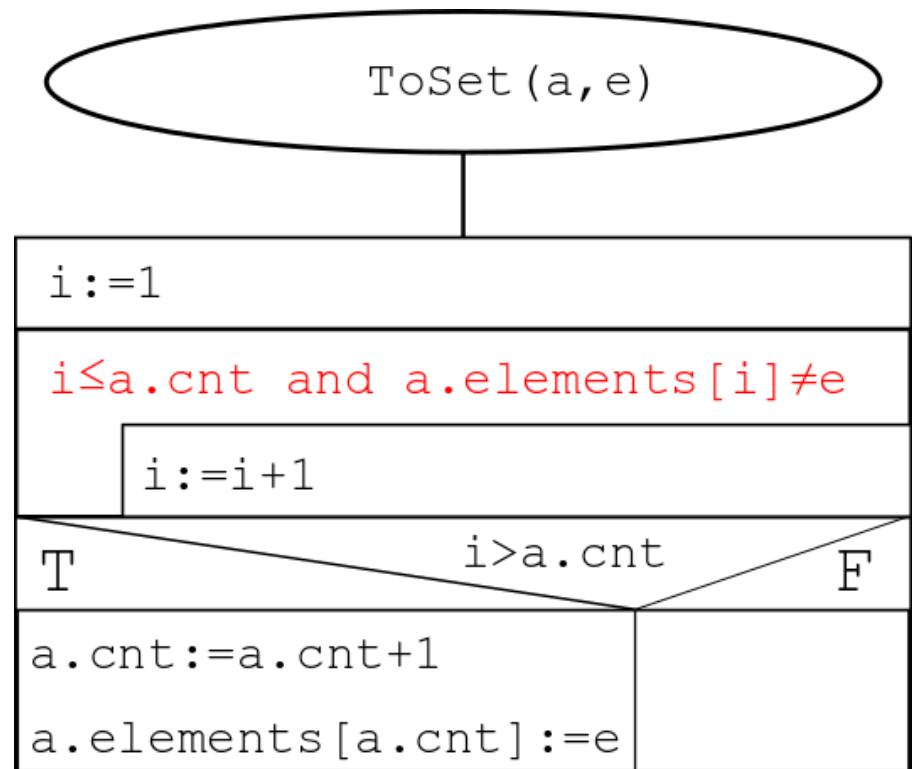
It does not depend on the count of elements of the set.

# The set type – by listing elements – ToSet

Applying the  
**Decision PoA**

**Calculation of the operation need:**

The loop will run as many times as many elements there are in the set – thus, the runtime is proportional to the count of elements of the set.

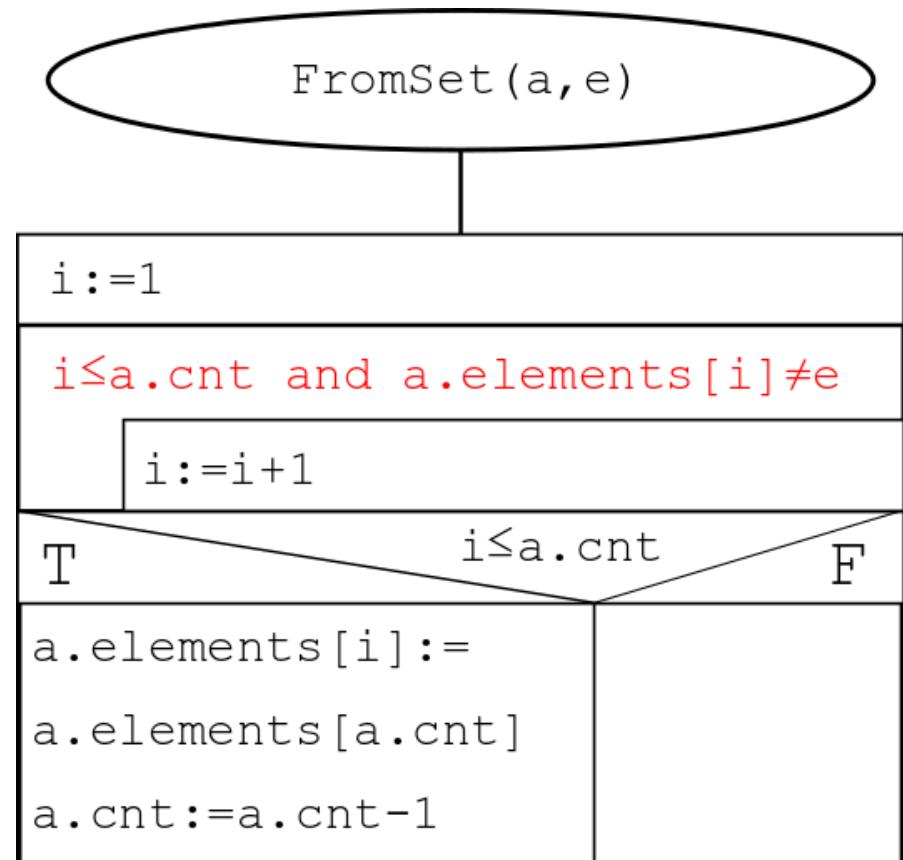


# The set type – by listing elements – FromSet

## Applying the **Search** PoA

### Calculation of the operation need:

The loop will run as many times as many elements there are in the set – thus, the runtime is proportional to the count of elements of the set.

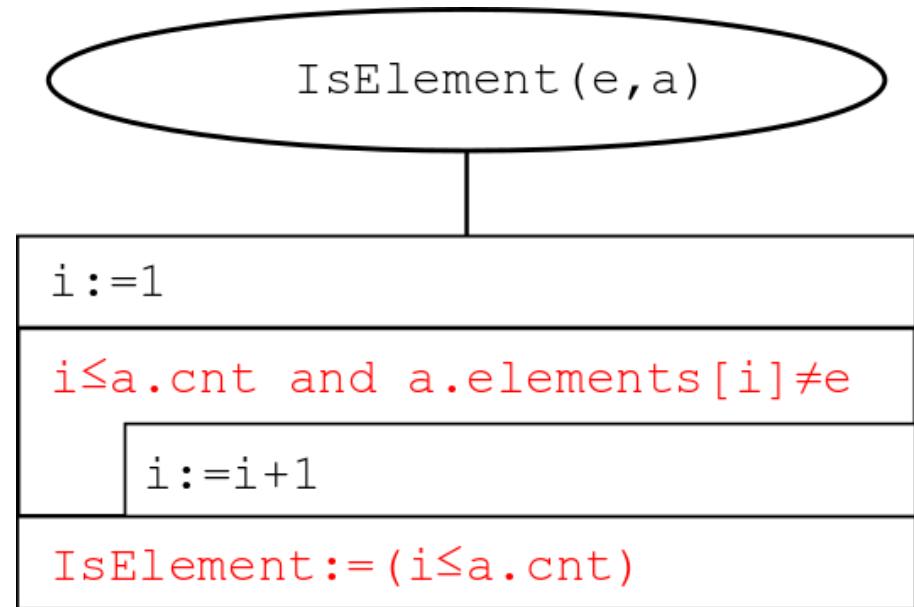


# The set type – by listing elements – IsElement

Applying the  
**Decision PoA**

**Calculation of the operation need:**

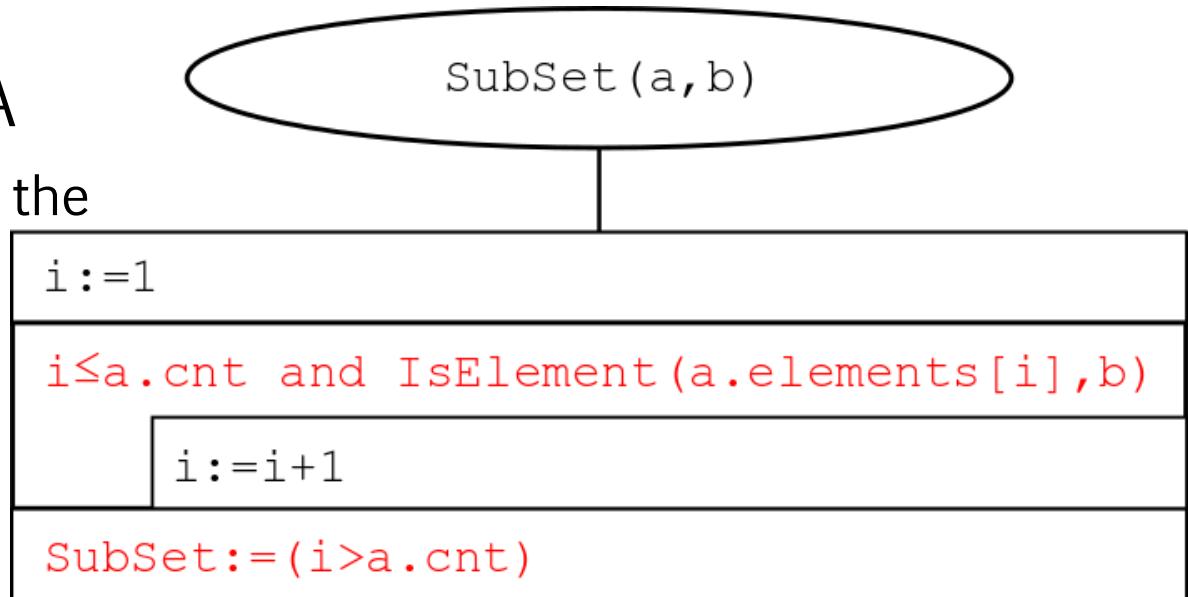
The loop will run as many times as many elements there are in the set – thus, the runtime is proportional to the count of elements of the set.



# The set type – by listing elements – SubSet

Applying the  
**Decision PoA**

with **decision** in the  
conditional  
statement



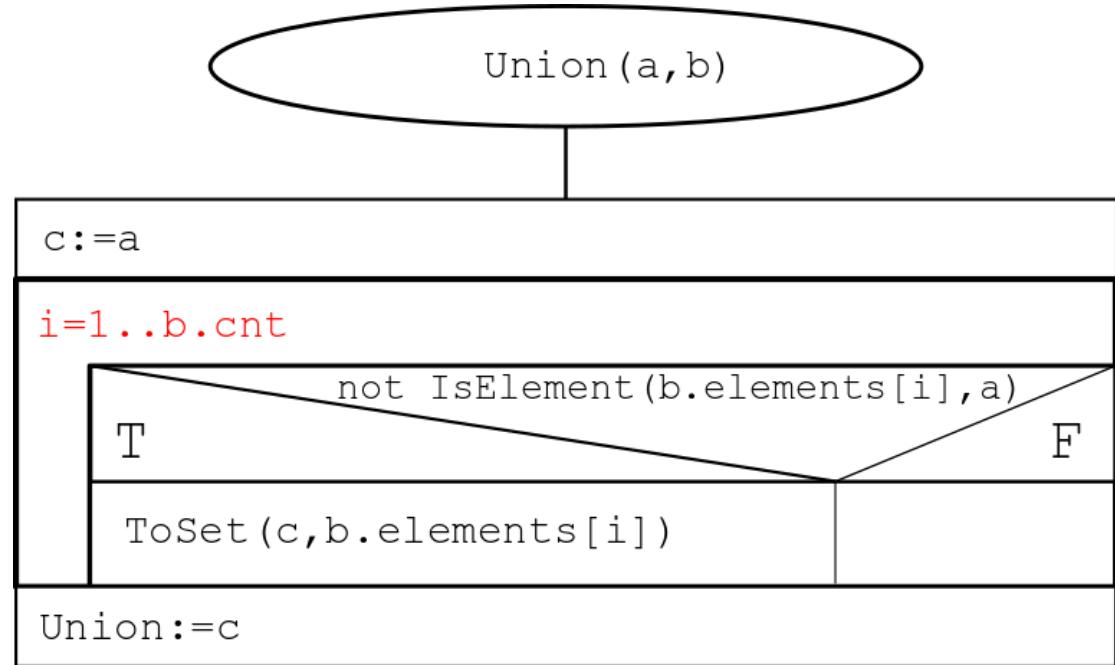
**Calculation of the operation need:**

The loop will run as many times as there are elements in set A, the `IsElement` function as many times as there are elements in set B, thus, the runtime is proportional to the product of the count of elements in the two sets.

# The set type – by listing elements – Union

Applying the  
**Copy, Multiple  
item selection,  
Decision**

PoA's



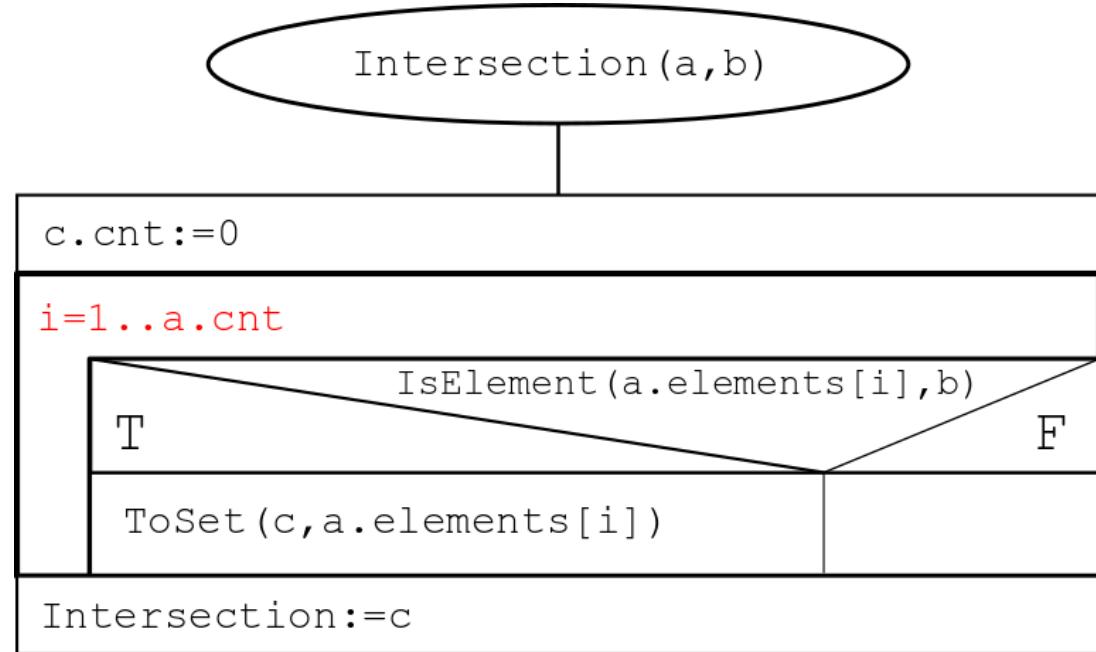
## Calculation of the operation need:

The loop will run as many times as many elements there are in set B, the IsElement function as many times as many elements there are in set A, thus, the runtime is proportional to the product of the count of elements in the two sets.

# The set type – by listing elements – Intersection

Applying the  
**Multiple item selection,**  
**Decision**

PoA's



## Calculation of the operation need:

The loop will run as many times as many elements there are in set A, the IsElement function as many times as many elements there are in set B, thus, the runtime is proportional to the product of the count of elements in the two sets.

# The set type – by listing elements

## Notes:

- › **Problem:** it is not checked if only elements that should be in the set are actually in the set.
- › No limit on the type of elements stored in the set, as we can store **anything** in an vector
- › **No limit on the count of elements** of the base set that the elements of the set derive from. We only limit the count of elements of the specific sets.

## The set type – as boolean vector

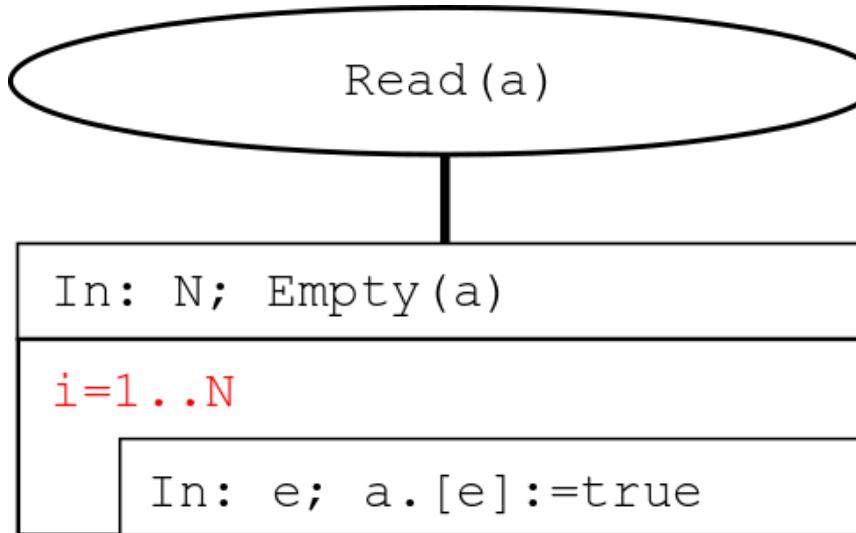
Bit map – boolean vector:

```
Set (ElementType) = Array (Min' ElementType .  
. Max' ElementType : Boolean)
```

We interpret the set as a vector of {true, false} elements, where we use the value of the element as index.

Such a set is always sorted.

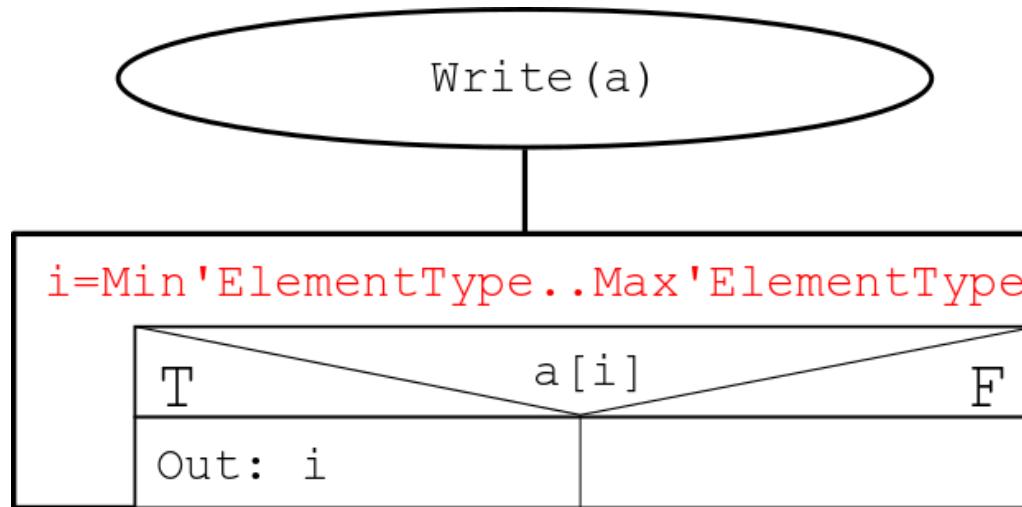
# The set type – as boolean vector - READ



## Calculation of the operation need:

The operation need of `Empty(a)` and the loop. The loop will run as many times as many elements there are in the set – thus, the runtime is proportional to the count of elements of the set

# The set type – as boolean vector - WRITE

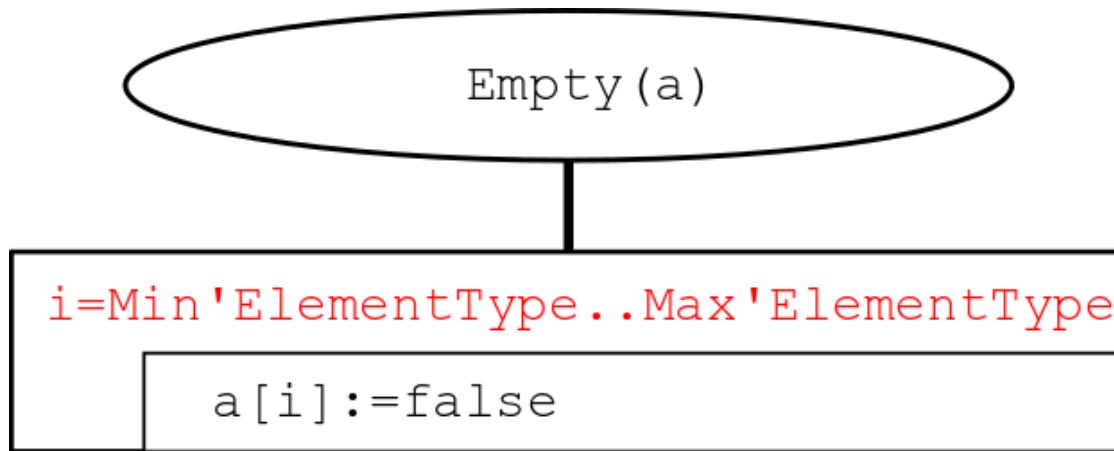


**Calculation of the operation need:**

The loop will run as many times as many elements there might be in the set – thus, the runtime is proportional to the cardinality of element type of the set.

*What if we stored the maximum and minimum element of the set?*

# The set type – as boolean vector - EMPTY

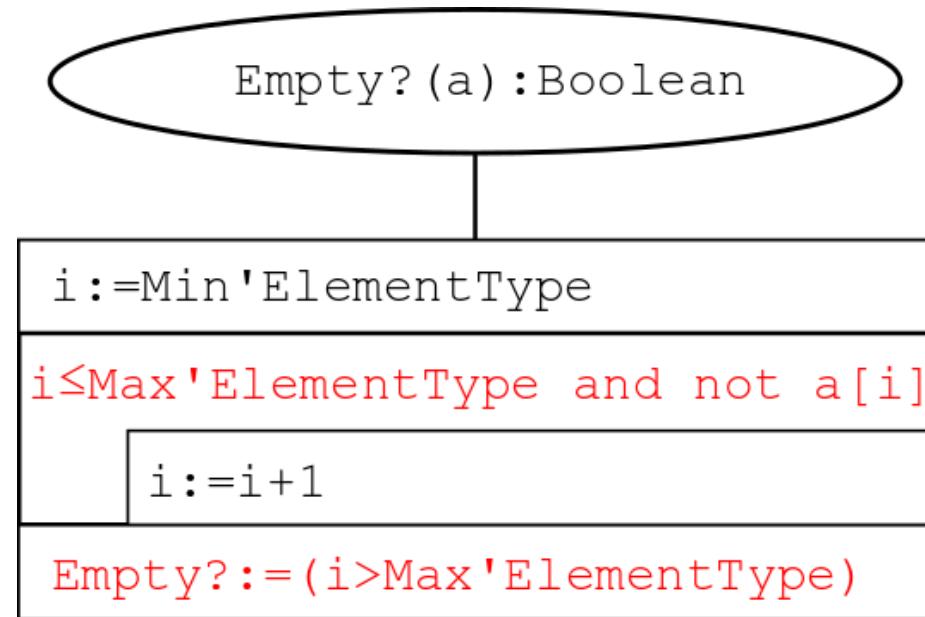


## Calculation of the operation need:

The loop will run as many times as many elements there might be in the set – thus, the runtime is proportional to the cardinality of element type of the set.

# The set type – as boolean vector – EMPTY?

Applying the  
**Decision PoA**



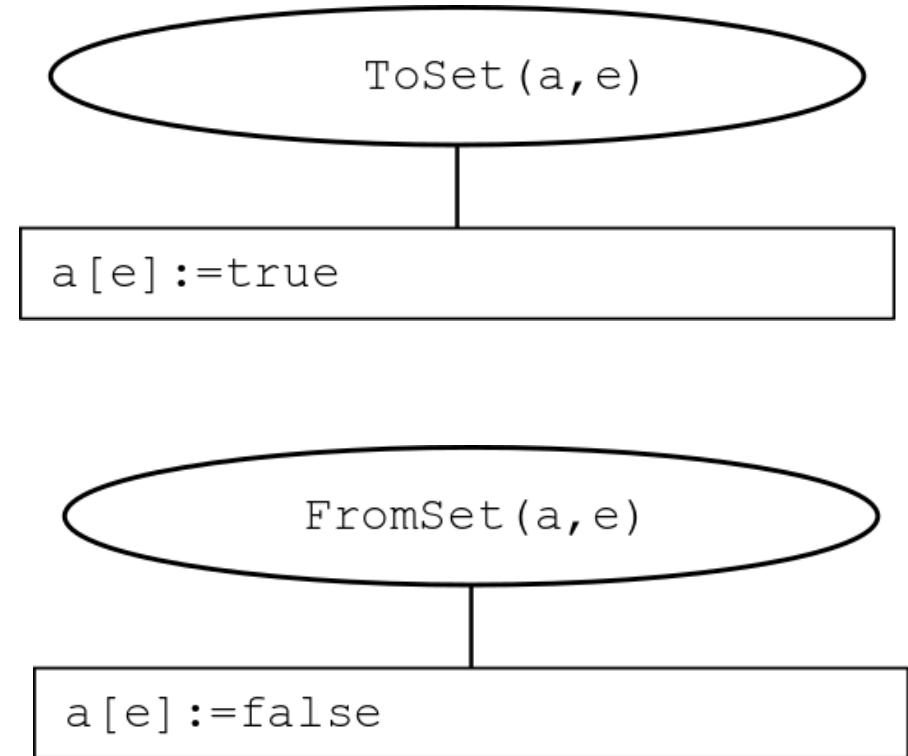
**Calculation of the operation need:**

The loop will run as many times as many elements there might be in the set – thus, the runtime is proportional to the cardinality of element type of the set.

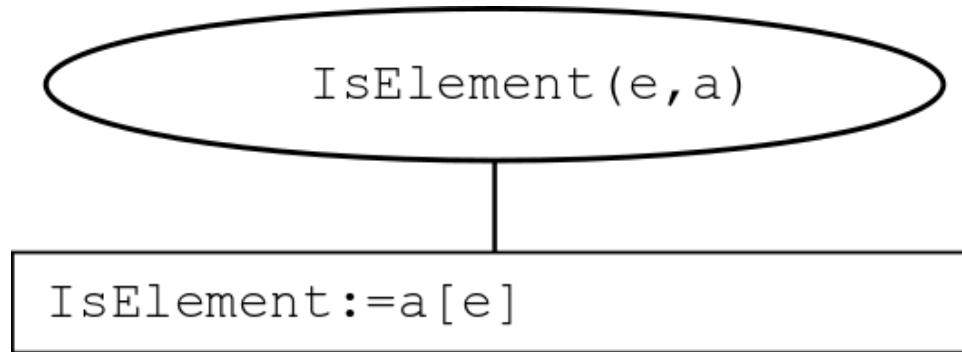
# The set type – as boolean vector – ToSet, FromSet

**Calculation of the operation need:**

It does not depend on  
the count of elements  
of the set.



# The set type – as boolean vector – IsElement

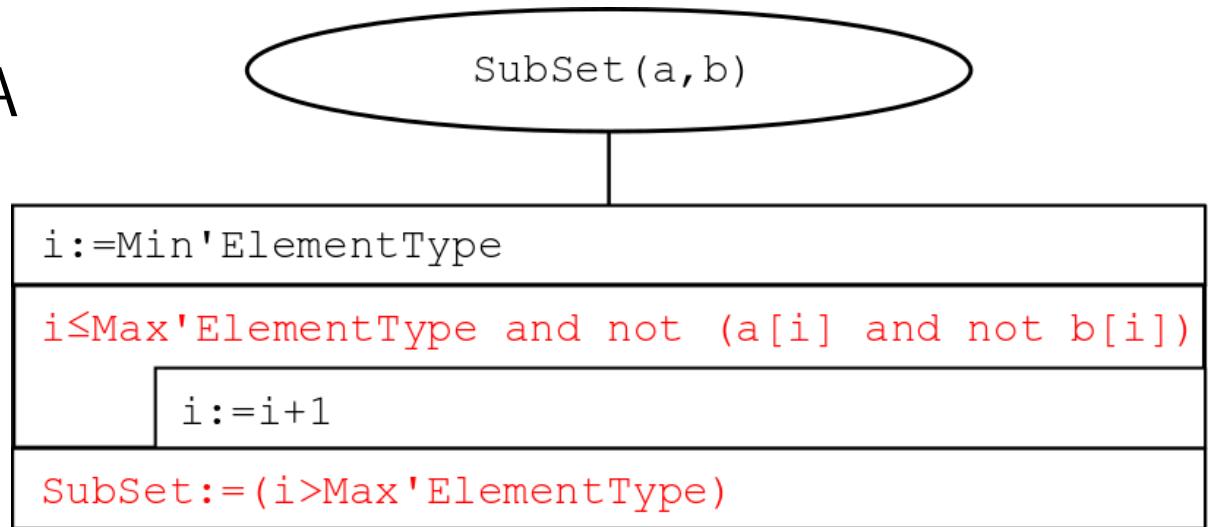


**Calculation of the operation need:**

It does not depend on the count of elements of the set.

# The set type – as boolean vector – SubSet

## Applying the **Decision PoA**

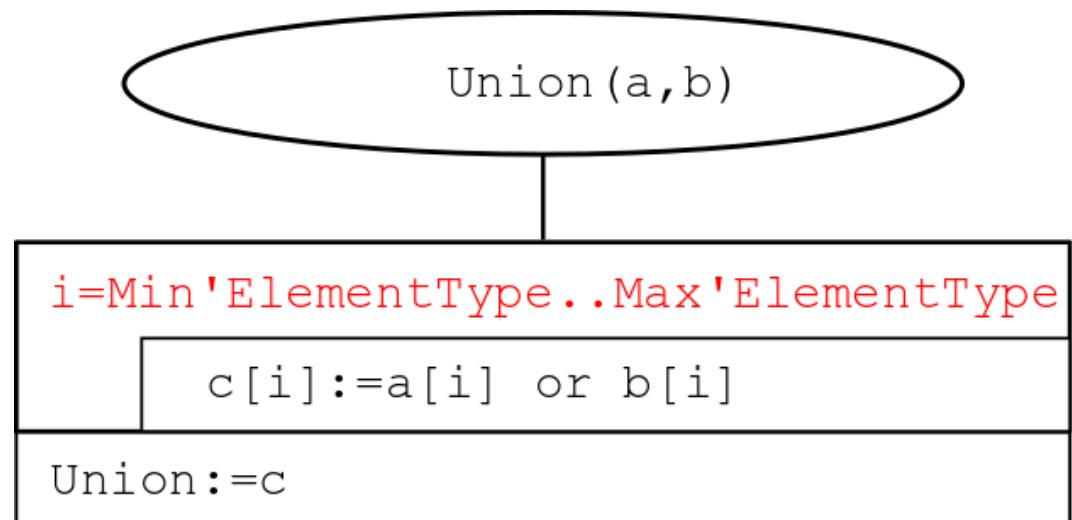


## Calculation of the operation need:

The loop will run as many times as many elements there might be in the set – thus, the runtime is proportional to the cardinality of element type of the set.

# The set type – as boolean vector – Union

Applying the  
**Copy PoA**

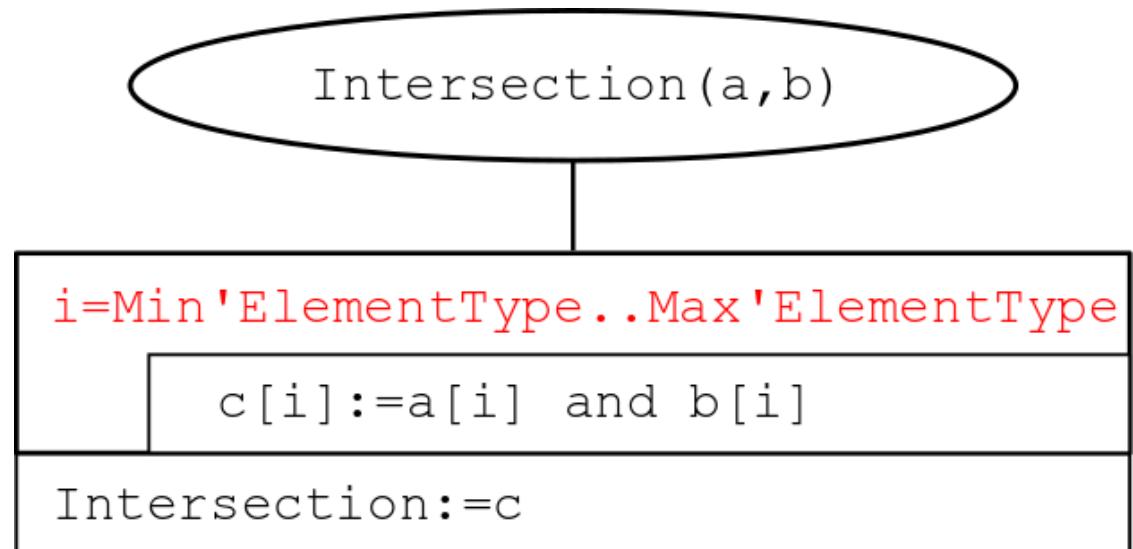


**Calculation of the operation need:**

The loop will run as many times as many elements there might be in the set – thus, the runtime is proportional to the cardinality of element type of the set.

# The set type – as boolean vector – Intersection

## Applying the Copy PoA



## Calculation of the operation need:

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