## Concurrent Systems — Exam June 2014

Name:	
Duration: 120 minutes — No document authorized	
<ul><li>1.</li><li>a) What does it mean for a lock to be <i>reentrant</i>?</li></ul>	
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<b>b)</b> What is a <i>Future</i> ?	-
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	-
c) Explain informally the notions of <i>work</i> and <i>critical path length</i> the parallelism of a multi-threaded program)?	- (as used when studying
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d) How do AtomicStampedReference and AtomicMarkableRe classical AtomicReference in Java? What problem can they help		differ	from
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e) Explain informally the principle of exponential backoff.	-		
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f) Give 3 examples of read-modify-write operations.	-		
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a

A semaphore<sup>1</sup> is an abstract data type used for controlling access, by multiple processes, to a common resource in a parallel programming environment. A useful way to think of a semaphore is as a record of how many units of a particular resource are available, coupled with operations to safely (i.e., without race conditions) adjust that record as units are required or become free, and, if necessary, wait until a unit of the resource becomes available. Semaphores that allow an arbitrary resource count are called counting semaphores.

Consider the following implementation of a counting semaphore:

```
public class CountingSemaphore {
 private int available;
 public CountingSemaphore(int n) {
    available = n;
 public synchronized void acquire() {
    while(available == 0)
      try { wait(); } catch(InterruptedException e) { }
   available--;
 public synchronized void release() {
    available++;
   notify();
  }
}
Propose an equivalent implementation of a counting semaphore that is lock-free.
```

<sup>&</sup>lt;sup>1</sup> <u>https://en.wikipedia.org/wiki/Semaphore\_(programming)</u>

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```
3.
Consider the simple barrier seen in the course:
public class Barrier {
  AtomicInteger count;
  int size;
  public Barrier(int n) {
    size = n;
    count = new AtomicInteger(size);
}
  public void await() {
    if (count.getAndDecrement() == 1)
      count.set(size);
    else
      while (count.get() != 0) { }
  }
}
What is the problem (limitation) with this implementation? Describe a scenario where
this would problematic. Explain informally how one can fix this problem by modifying
the barrier implementation.
```

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1 Q:b	
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2 Q.deq()	
2 Q:a	
) Can you think o	of a history that would be linearizable but not sequentially consistent?

Consider the code below. Explain what it does, which properties it provides, and how a well-behaved client should use it.

```
import java.util.concurrent.atomic.*;
public class C {
 private AtomicInteger c, n;
 public C() {
   c = new AtomicInteger();
   n = new AtomicInteger();
 public void f() {
   int i = n.getAndIncrement();
   while (c.get() != i);
 public void g() {
   c.getAndIncrement();
}
```

A read-write lock allows either a single writer or multiple readers section.	to execute in a critical
a) Provide an implementation of a read-write lock in Java. You methods and the wait/notify mechanism if you wish. The class methods lockRead(), unlockRead(), lockWrite(), and implementation does not need to be FIFO, starvation-free, nor reen HINT: You might want to keep track of the number of readers and	should provide the 4 inlockWrite(). This trant.
<pre>public class ReadWriteLock {</pre>	

not starve (i.e., readers cannot prevent writers from acquiring the lock infinitely). This implementation does not need to be FIFO nor reentrant. public class ReadWriteLock { }

b) Modify the implementation of your read-write lock so that it ensures that writers do